Monads and C++ Template Metaprogramming

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Motivation

- Eric Niebler's (Joel Falcou's) Proto
- Example: Lambda DSL
 - Expressions turned into types
 - Types manipulated at compile time
 - The result: a runtime function object
- Mixing compile-time TMP with runtime execution



Plan

- Convert Haskell Expression monad to C++
 - Expression tree (compile-time)
 - State (runtime)
 - Action (constructed at compile-time, executed at runtime)
 - Metafunctions Bind/Return (compile-time)
 - "Compile" metafunction (compile-time)
- 2-argument lambda EDSL



Teaser

```
template<class L, class R>
struct Compile<Plus<L, R>> : Prog {
  int operator()(Args args) {
    return Bind<Compile<L>, Bind<Compile<R>, Return>> (
       Compile<L>(),
       [](int left) -> Bind<Compile<R>, Return> {
          return Bind<Compile<R>, Return>(
          Compile<R>(),
            [left](int right) -> Return {
               return Return(left + right);
     )(args);
```

Domain Specific Languages

- Write a C++ expression which, by
- Abusing C++ operator overloading
- Is interpreted as something completely different:
- An expression tree of very specific type
- The type is processed at compile time resulting in an object that can be
- Executed at runtime to produce a desired result
- What does it have to do with Haskell?
 - Hint: the expression monad



Expression Tree (compile time)

```
template<int c> struct Const {};

template<class E1, class E2>
struct Plus {};

template<class E1, class E2>
struct Times {};

struct Arg1 {};
struct Arg2 {};
```



State (runtime)

```
type Args = [Int]
```

```
struct Args
{
    Args(int i, int j) {
        _a[0] = i;
        _a[1] = j;
    }
    int operator[](int n) { return _a[n]; }
    int _a[2];
};
```



Type Constructor

```
Args -> a
```

```
newtype Prog a = PR (Args -> a)
```

- Prog created at compile time using a metafunction
- Action executed at runtime
- PR is really a concept with operator() as associated function

```
struct PR {
      // int operator()(Args args);
};
```



Monadic Metafunction

- Metafunctions may "return"
 - Values
 - Types
 - Other metafunctions
- New type of metafunction "returning" a function
 - (Further generalization: metafunction returning template function)

```
getArg :: Int -> Prog Int
getArg n = PR (λ args -> args !! n)
```



Bind: Construction



Bind: Action

```
template < class P1, class P2>
struct Bind : PR {
    ...
    int operator()(Args args) {
        int v = _prog(args);
        P2 prog2 = _cont(v);
        return prog2(args);
    }
    ...
};
```

Return

```
return :: a -> Prog a return v = PR (λ args -> v)
```

```
struct Return : PR
{
    Return(int v) : _v(v) {}
    int operator()(Args args)
    {
        return _v;
    }
    int _v;
};
```



The Compile Metafunction

```
compile :: Exp -> Prog Int
```

- Metafunction Compile
 - Takes compile-time Exp
 - Returns a Prog
- Every specialization will define its own operator()

```
template<class Exp>
struct Compile {};
```



Simple Specializations

```
compile (Const c) = return c
template<int c>
struct Compile<Const<c>> : Return
   Compile() : Return(c) {}
};
compile Arg1 =
    getArg 0
template<>
struct Compile<Arg1> : GetArg<0> {};
```

Composite Specializations

```
template<class L, class R>
struct Compile<Plus<L, R>> {
  int operator()(Args args)
    return Bind<...> (
      Compile<L>(),
      [](int left) {
        return Bind<...>(
          Compile<R>(),
          [left](int right) {
            return Return(left + right);
                          compile (Plus exL exR) =
                            bind compile exL
    )(args);
                                  λ left ->
                                     bind compile exR
                                          λ right ->
                                              return (left+right)
```

The Plus Node: Types

```
template<class L, class R>
struct Compile<Plus<L, R>> {
  int operator()(Args args)
   return Bind<Compile<L>, Bind<Compile<R>, Return>> (
      Compile<L>(),
      [](int left) -> Bind<Compile<R>, Return> {
        return Bind<Compile<R>, Return>(
          Compile<R>(),
          [left](int right) -> Return {
            return Return(left + right);
                          compile (Plus exL exR) =
                            bind compile exL
    )(args);
                                 λ left ->
                                    bind compile exR
                                         λ right ->
                                             return (left+right)
```

Test: Arg1 * Arg2 + 13

```
testExp =
  let exp = Plus (Times Arg1 Arg2) (Const 13)
  compile exp
```

```
test =
  let args = [3, 4]
    act = testExp
  in
    run act args
```

```
void main () {
    Args args(3, 4);
    Compile<Plus<Times<Arg1, Arg2>, Const<13>>> act;
    int v = act(args);
    std::cout << v << std::endl;
}</pre>
```

2-Arg Lambda EDSL

```
int x = (arg1 + arg2 * arg2)(3, 4);
```

- Trick C++ into converting this expression into a tree
- arg1 and arg2: objects of types for which overloaded operators + and * exist
- Their return types correspond to expression trees
- Expression trees are (2-argument) function objects



Expression Wrapper

- For any expression E
 - Compile it to an action
 - Run the action and return the result

```
template < class E>
struct Lambda {
    int operator()(int x, int y) {
        Args args(x, y);
        Compile < E > prog;
        return prog(args);
    }
};
```



From Expression to Lambda

```
int x = (arg1 + arg2 * arg2)(3, 4);
```

Special Lambda objects for Arg1 and Arg2

```
const Lambda (Arg2) arg1; const Lambda (Arg2) arg2;
```

- Overloaded operators
 - Generate types at compile time
 - Generate function objects at runtime

```
template<class E1, class E2>
Lambda<Plus<E1, E2>> operator+ (Lambda<E1> e1, Lambda<E2> e2)
{
    return Lambda<Plus<E1, E2>>();
}
```



Compile-Time vs. Runtime

```
(arg1 + arg2 * arg2) (3, 4)
// Compile time
Lambda<Arg1> Lambda<Arg2> Lambda<Arg2> // original types
             Lambda<Times<Arg2, Arg2>> // type returned by *
Lambda<Plus<Arg1, Times<Arg2, Arg2>> // type returned by +
// runtime, after template expansion
int Lambda<Plus<Arg1, Times<Arg2, Arg2>>::operator()(int x, int y)
{
    Args args(x, y);
    Compile<Plus<Arg1, Times<Arg2, Arg2>> prog;
    return prog(args);
```



Haskell Lambda EDSL

```
newtype Lambda = L Exp
toFun (L ex) =
 \ x y ->
     run (compile ex) [x, y]
instance Num Lambda where
    (L e1) + (L e2) = L (Plus e1 e2)
    (L e1) * (L e2) = L (Times e1 e2)
    fromInteger n = L (Const n)
test =
    let arg1 = L Arg1
        arg2 = L Arg2
    in
        (toFun (arg1 + 2 * arg2 * arg2)) 2 3
```

Conclusion

- Almost mechanical translation from Haskell state monad to C++ EDSL
- Haskell code easy to understand (after some initial pains) and test
- What seemed to be a bunch of template hacks gains strong theoretical foundations
 - Makes possible reasoning and proofs of correctness
- Reusable abstraction with unexplored potential
 - C++ state monad orthogonal to the construction of EDSL
 - Bind and Return used in defining monadic metafunction "Compile"
 - Monadic metafunction plugged into EDSL



Future Directions

- Explore metafunctions that "return" template functions
 - This is what Proto Transform does
- Better story on Const
 - Const values are part of state (even though they're known at compile time)
- Factor out Transform and Domain of Proto
 - It's really Transform that is a monad?



Bibliography

- Mike Vanier's blog (Haskell monads): http://mvanier.livejournal.com/3917.html
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- My blog (extended treatment of current presentation):
 http://bartoszmilewski.wordpress.com/2011/01/09/mo-nads-for-the-curious-programmer-part-1/
- Brian McNamara, Yannis Smaragdakis, <u>Syntax</u> sugar for FC++: lambda, infix, monads, and more.

