# Merkle-tree-based integrity verification protocol for geo-distributed storage systems

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#### Content

- What problem are we addressing?
- How did we solve it?
- What about scalability?

#### **Problem**

How can we verify that none of the files are corrupted?

```
root@node:/root# ls -lh
total 200G
                  root
                        21M Jan 1 00:00 photo1.png
                  root
                        4.4K Jan 10 08:30 photo2.png
                  root
                           2G Jan 7 09:50 photo3.png
                  root
                        7.1M Mar 9 11:00 photo4.png
-rw-r--r-- 1 root
                  root
                        24K Mar 9 11:01 photo5.png
-rw-r-- 1 root
                  root
                       130M Jun 25 18:25 photo6.png
```

#### **Problem**

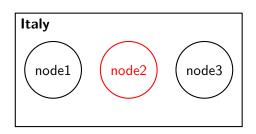
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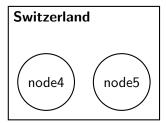
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We can use checksums.

#### **Problem**

But, what if we have a file split in pieces and distributed across different regions? We would need to perform a checksum on each node.





### Cubbit

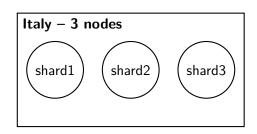


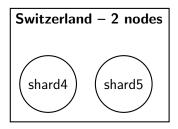
Cubbit, a geo-distributed storage system.

#### Cubbit – How it works

Each file is split in n + k shards (i.e., pieces) using Reed-Solomon.

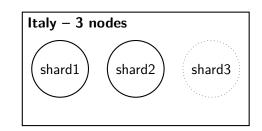
At least *n* shards are required to reconstruct the file, while *k* shards provide redundancy.

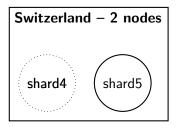




# Cubbit - Problems using checksum

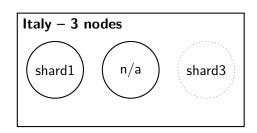
If nodes are offline, it becomes impossible to check all shards for a file.

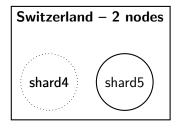




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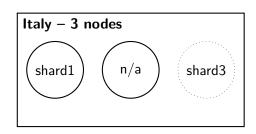
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- ② During an upload some agents can be offline, but they could be online during the check.

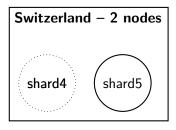




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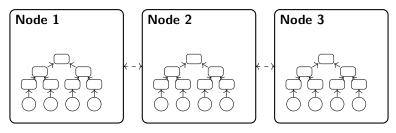
- If nodes are offline, it becomes impossible to check all shards for a file.
- ② During an upload some agents can be offline, but they could be online during the check.
- Check for each reconstructed file or for each shard?





#### Solution

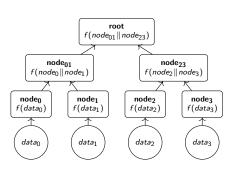
Each node uses a Merkle tree structure to organize shards during integrity verification. All nodes agree on which file is corrupted through Raft consensus. Data are organized using Reed-Solomon codes.



#### Merkle trees

It is a binary tree T of height H with  $2^H$  leaves and  $2^H - 1$ internal nodes. Each leaf stores the cryptographic hash of the underlying data, rather than the raw data itself. The same cryptographic hash function is applied recursively at internal nodes, which store the hash of the concatenation of their two children.

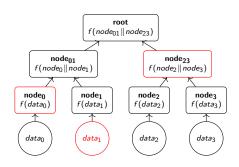
$$n_{parent} = f(n_{left}||n_{right})$$



#### Merkle trees - Proof

Merkle trees has the ability to prove that a given piece of data is part of a larger set, without revealing or recomputing the entire dataset. For  $data_1$  we have  $\pi = \{node_0, node_{23}\}.$ 

We can make a proof verification comparing the result with the presumed root in  $O(\log n)$ .



#### Merkle trees – Lib

A dedicated Rust library was developed to efficiently compute Merkle tree root hashes.

Hash function	5 MB	10 MB	15 MB
SHA-256	89.901 ms	178.42 ms	268.53 ms
Keccak-256	521.49 ms	1.1334 s	1.3438 s
BLAKE3	73.091 ms	154.68 ms	219.79 ms



#### Raft

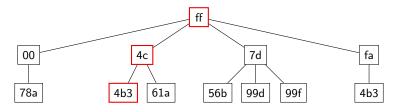
A consensus protocol, where each server on a cluster is a follower, a candidate or a leader. There is a single leader, responsible for sending messages to other servers via the Raft log.



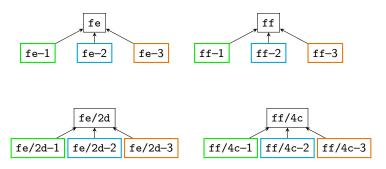
## Upload flow

Let's save our Christmas holidays family picture on Cubbit.

- Select a file. (e.g., xmas2024.png)
- Gateway applies Reed-Solomon codes and sends each shard to a different agent. (e.g., xmas2024.1 to Agent 1)
- Each agent chooses the same pseudo-random string for the filename. (e.g., ff4c4b3)
- Finally, every agent saves the file using a two-level folder structure. (e.g. ff/4c/4b3.1 for Agent 1)



# Aggregated roots



The leader builds Merkle trees for each folder, with the agents' Merkle root hashes as leaves.

## Map of hashes

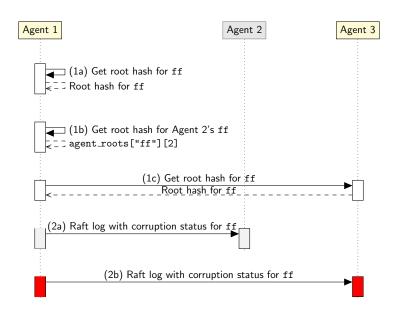
```
roots = {
    "fe": "root hash of fe",
    "ff": "root hash of ff",
   "ff/4c": "root hash of ff/4c",
   # ...
}
agent_roots = {
    "fe": [
        "Agent 1 root hash of fe",
        "Agent 2 root hash of fe",
        "Agent 3 root hash of fe"
   ],
   # ...
    "ff/4c": [
        "Agent 1 root hash of ff/4c",
        "Agent 2 root hash of ff/4c",
        "Agent 3 root hash of ff/4c"
   ],
```

## Map of hashes

#### This map allows us to:

- Use agent\_roots[<some\_folder>][i-th-agent] when the i-th agent is offline.
- If an agent is offline during the upload, the roots[<some\_folder>] value is still computed using the current status of every agent.

## Check corruptions



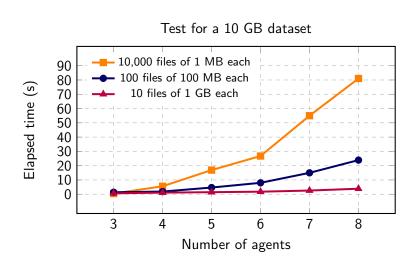
## Tests – Set up

8 Agents with 4 CPU(s), 8-16 GB of RAM, and 45 GB of Disk.

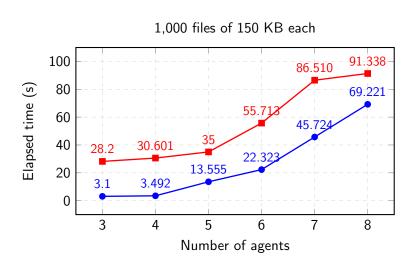


For each test, we consider the number of online/offline agents and the file sizes and counts.

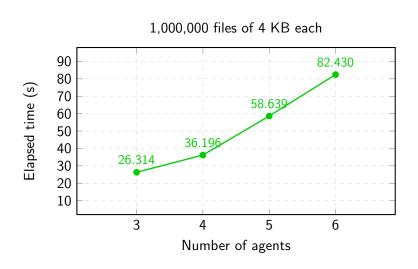
#### Same dataset but different file size and count



## Test with offline agents



# Test with very large number of tiny files



#### Results

- Scenarios with many small files amplify the cost of synchronization and consensus.
- Verification time increases with the number of agents and files.
- Confirmed the correctness and resilience of the approach under different scenarios and network conditions.

#### Conclusion

By combining Merkle trees, Raft consensus, and Reed-Solomon codes, we built a scalable and fault-tolerant protocol for distributed data integrity verification without relying on full file scans or constant node availability.

