CMPE 110 Computer Architecture

Fall 2016, Homework #4

Submission Guidelines

- → The homework must be submitted to ecommons by 11:59pm on **Thursday**, **December 1**st
- → Anything later is a late submission.

Format

- → The homework must be typed and submitted as a single file in PDF format.
- → Homework file should be saved as cmpe110-hw4-yourcruzid.pdf.
- → Please write your name and your UCSC email address.
- → The homework should be "readable" without too much effort.
- → To increase readability, use <u>heading styles</u> and a <u>monospaced</u> font (e.g. courier) for code.

Content

- → Keep your responses coherent and organized or you may lose points.
- → Clearly state all assumptions and provide details on how to reach a solution. An answer without explanation gets no credit.

Points

64 = 20 + 24 + 20

1. Address Translation (20 points)

Below, you are provided with a snapshot of a 4 entry, fully associative TLB (translation-lookaside buffer) and a page table.

→ If needed, you may assume that the value of the page table register is 0

PARAMETER	VALUE
page size	4 KB
associativity	fully associative
replacement policy	LRU - the entry that has not been used for the longest time will be evicted if a new entry is to be added

Initial TLB State

TLB INTERPRETATION

COLUMN	VALUE	MEANING
VALID	0	entry is not valid (TLB miss)
VALID	1	entry is valid (TLB hit)
DIRTY	0	data has not been modified
DIKTT	1	data has been modified
LRU	1	most recently used
LKO	4	least recently used

TLB TABLE

VALID	DIRTY	LRU	TAG (VPN)	PHYSICAL PAGE #
1	1	2	0101	0111
1	0	3	1100	1011
1	0	1	0001	1001
1	0	4	0010	0101

Initial Page Table State

ADDRESS (VPN)	VALID	PHYSICAL PAGE #
0000	0	Disk
0001	1	1001
0010	1	0101
0011	1	0010
0100	1	0101
0101	0	Disk
0110	1	0111
0111	1	0110
1000	1	1001
1001	1	1000
1010	0	Disk
1011	1	1010
1100	1	1011
1101	1	1100
		···

Problem

Show the final state of the TLB after processing the 2 virtual address sequence given below. If a page fault happens, the physical page number to be entered into the page table would be 0000; if another page fault happens, use physical page number 0001.

VIRTUAL ADDRESSES SUPPLIED

```
1. (MSB) 1010 1111 1110 1010 (LSB) (load instruction)
```

^{2. (}MSB) 0011 1100 1001 0001 (LSB) (store instruction)

Solution

- 4KB pages => 12 bits for offset, 4 bits for VPN
 1. (MSB) 1010 1111 1110 1010 (LSB) (load instruction)
- 2. (MSB) 0011 1100 1001 0001 (LSB) (store instruction)

INITIAL STATE

VALID	DIRTY	LRU	TAG	PHYSICAL PAGE #
1	1	2	0101	0111
1	0	3	1100	1011
1	0	1	0001	1001
1	0	4	0010	0101

AFTER INSTRUCTION 1: 1010 1111 1110 1010

VALID	DIRTY	LRU	TAG	PHYSICAL PAGE #
1	0	3	0101	0111
1	0	4	1100	1011
1	0	2	0001	1001
1	0	1	1010	0000

AFTER INSTRUCTION 2: 0011 1100 1001 0001

VALID	DIRTY	LRU	TAG	PHYSICAL PAGE #
1	0	4	0101	0111
1	1	1	0011	0010
1	0	3	0001	1001
1	0	2	1010	0000

2. Cache Coherence (24 points)

Consider a processor with 4 cores. In each of the following parts you will fill out a table for a different cache coherence protocol. For the following memory references, show the state of the cache line containing the variable x in each core's cache. Consider the cacheline only in L1.

2.A VI Protocol (8 points)

Assume for this part that we use the VI protocol for cache coherence. Fill out the table as specified at the beginning of the question.

	CORE	REQUEST TYPE	CO \$ LINE STATE	C1 \$ LINE STATE	C2 \$ LINE STATE	C3 \$ LINE STATE
1	0	read x				
2	1	read x				
3	2	read x				
4	3	write x				
5	1	read x				

2.B MSI Protocol (8 points)

Assume for this part that we use the MSI protocol for cache coherence. Fill out the table as specified at the beginning of the question.

	CORE	REQUEST TYPE	CO \$ LINE STATE	C1 \$ LINE STATE	C2 \$ LINE STATE	C3 \$ LINE STATE
1	0	read x				
2	1	read x				
3	2	read x				
4	3	write x				
5	1	read x				

2.C MESI Protocol (8 points)

Assume for this part that we use the MESI protocol for cache coherence. Fill out the table as specified at the beginning of the question.

	CORE	REQUEST TYPE	CO \$ LINE STATE	C1 \$ LINE STATE	C2 \$ LINE STATE	C3 \$ LINE STATE
1	0	read x				
2	1	read x				
3	2	read x				
4	3	write x				
5	1	read x				

Question 2. Cache Coherence Protocols (15 points) <u>Solution.</u>

Question 2.A VI Protocol (5 points)

	Core	Request Type	C0 Cache Line State	C1 Cache Line State		
1	0	Read x	V	-	-	-
2	1	Read x	Ι	V	-	-
3	2	Read x	Ι	I	V	-
4	3	Write x	Ι	I	I	V
5	1	Read x	Ι	V	I	I

Question 2.B MSI Protocol (5 points)

	Core	Request Type	C0 Cache Line State		C2 Cache Line State	
1	0	Read x	S	-	-	-
2	1	Read x	S	S	-	-
3	2	Read x	S	S	S	-
4	3	Write x	Ι	I	I	M
5	1	Read x	Ι	S	I	S

Question 2.C MESI Protocol (5 points)

	Core	Request Type		C1 Cache Line State	C2 Cache Line State	C3 Cache Line State
1	0	Read x	E	-	-	-
2	1	Read x	S	S	-	-
3	2	Read x	S	S	S	-
4	3	Write x	Ι	Ι	I	M
5	1	Read x	I	S	I	S

3. Virtual Memory (20 points)

An ISA supports a 16-bit, byte-addressable virtual address space. The corresponding physical memory has 256 bytes. Each page contains 32 bytes. A simple, one-level translation scheme is used and the page table resides in physical memory. The initial contents of the frames of physical memory are shown below.

PARAMETER	VALUE
virtual address size	16 bit
addressing type	byte addressable - each byte of memory has its own address
page size	32 B
physical memory size	256 B

Physical Page Contents

LRU	FRAME (PHYSICAL PAGE) #	FRAME CONTENTS
3	0	Empty
?	1	Page 7
?	2	Page 1
?	3	Page 4
2	4	Empty
?	5	Page 3
1	6	Empty
?	7	Page Table

Questions

A three-entry TLB that uses Least Recently-Used (LRU) replacement is added to this system. Initially, this TLB contains the entries for pages 3, 4, and 7.

INITIAL TLB

LRU	TAG (VPN)	PHYSICAL PAGE #
?	3	5
?	4	3
?	7	1

3.A TLB MISSES / PAGE FAULTS

For the following sequence of references, put a circle around those that generate a TLB hit and put a rectangle around those that generate a page fault. What is the hit rate of the TLB for this sequence of references?

Note: LRU policy is used to select pages for replacement in physical memory. Assume initial conditions where physical page number (PPN) 0 was used less recently than PPN 4, which was used less recently than PPN 6:

LRU (PPN 0) > LRU (PPN 4) > LRU (PPN 6)

Virtual Page References

3.B TLB ENTRIES

At the end of this sequence, what three entries are contained in the TLB?

3.C FRAME CONTENTS

What are the contents of the 8 physical frames?

Solution

3.A TLB MISSES / PAGE FAULTS

TLB Hit Rate = 8/16

3.B TLB ENTRIES

9, 12, 7

3.C FRAME CONTENTS

PPN	VPN
0	10
1	7
2	9
3	4
4	11
5	12
6	8
7	Page Table