

## EEE 424 Coding Assignment 1

In this assignment the aim was to compute and compare different Discrete Fourier Transform (DFT) methods.

### Q1)

In this part of the assignment 5 different DFT approaches are implemented and applied to 3 different length arrays ( $N = 32, 256, 4096$ ). Whole code of this question is provided in **Appendix A**.

These five different approaches are:

- DFT summation: Direct DFT Summation formula
- DFT matrix: DFT Matrix algorithm
- FFT-DIT: Fast Fourier Transform (FFT) using the decimation-in-time algorithm
- FFT-DIF: FFT using decimation-in-frequency algorithm
- fft: MATLAB's built-in FFT command

Since the steps (a) to (h) are repeated for different lengths of arrays, steps (a)-(h) are written as a function called **steps\_a\_h(x)** which is provided in **Appendix A**.

First, a complex array of length  $N = 32$  is created with the given code snippet:

```
rng(2, "twister")
N = 32;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;
```

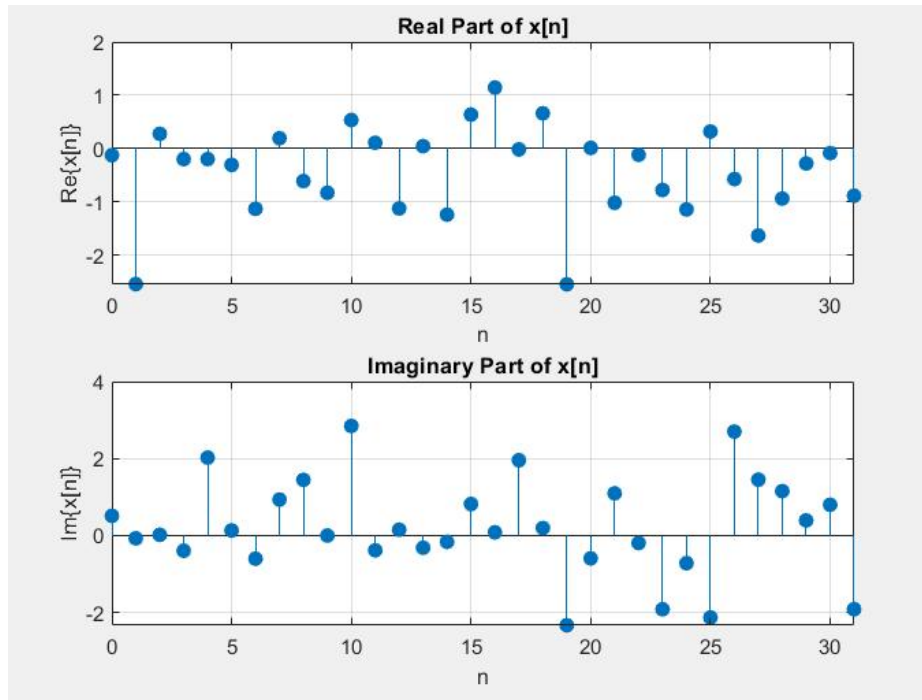
### a)

Real and imaginary parts of  $x[n]$  of length  $N = 32$  are plotted with the following code snippet:

```
n = 0:N-1;

figure;
subplot(2,1,1);
stem(n, real(x), 'filled');
title('Real Part of x[n]');
xlabel('n');
ylabel('Re\{x[n]\}');
xlim([0 N-1]);
grid on;

subplot(2,1,2);
stem(n, imag(x), 'filled');
title('Imaginary Part of x[n]');
xlabel('n');
ylabel('Im\{x[n]\}');
xlim([0 N-1]);
grid on;
```



**Fig.1** Real and Imaginary parts of  $x[n]$ ,  $N=32$

b)

The definition of DFT in summation form:

$$X[k] = \sum_{n=0}^{N-1} x[n] \exp\left(-j2\pi \frac{kn}{N}\right)$$

In order to compute  $N = 32$  point DFT of the array  $x[n]$  using the direct definition in summation form the following function is defined:

```
function X = DFT_summation(x)
N = length(x);
X = zeros(1, N);
for k = 0:N-1
    for n = 0:N-1
        X(k+1) = X(k+1) + x(n+1) * exp(-1i * 2 * pi * k * n / N);
    end
end
end
```

In order to compute DFT of  $x[n]$ , the function is called as follows:

```
t1 = tic;
Xsum = DFT_summation(x);
elapsed_time = toc(t1);
fprintf('Elapsed time for DFT_summation is %.6f seconds.\n', elapsed_time);
plotDFT(Xsum, 'Summation Formula')
```

c)

DFT matrix defined as:

$$\mathbf{W} = \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \omega^1 & \omega^2 & \dots & \omega^{N-1} \\ 1 & \omega^2 & \omega^4 & \dots & \omega^{2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{N-1} & \omega^{2(N-1)} & \dots & \omega^{(N-1)(N-1)} \end{bmatrix}$$

where:

$$\omega = e^{-j\frac{2\pi}{N}}$$

In order to create DFT matrix and compute the DFT of the array  $x[n]$  using the DFT matrix the following function is defined:

```
function X = DFT_matrix(x)
N = length(x);
WN = exp(-1i * 2 * pi / N);
DFT_mat = zeros(N, N);
for k = 0:N-1
    for n = 0:N-1
        DFT_mat(k+1, n+1) = WN^(k * n);
    end
end
X = (DFT_mat * x.').';
end
```

d)

DFT of  $x[n]$  calculated with the DFT matrix as follows:

$$X = \mathbf{W} \cdot \mathbf{x}$$

In order to compute DFT of  $x[n]$ , the DFT matrix function is called as follows:

```
t2 = tic;
X_mat = DFT_matrix(x);
elapsed_time = toc(t2);
fprintf('Elapsed time for DFT_matrix is %.6f seconds.\n', elapsed_time);
plotDFT(X_mat, 'DFT matrix')
```

e)

Decimation-in-time Fast Fourier Transform is defined as recursively dividing the even and odd samples:

$$X_k = \sum_{n=0}^{N-1} x_n W_N^{kn} = \sum_{n=0}^{N/2-1} x_{2n} W_N^{k(2n)} + \sum_{n=0}^{N/2-1} x_{2n+1} W_N^{k(2n+1)}$$

where:

$$W_N = e^{-i\frac{2\pi}{N}}$$

Key Steps in the FFT DIT Algorithm:

1. Divide: The sequence  $x$  is split into two smaller subsequences: one containing the even-indexed elements and the other containing the odd-indexed elements.
2. Conquer: The FFT of the even and odd subsequences is computed recursively.
3. Combine: The results of the even and odd subsequences are combined using a "twiddle factor"  $W$ , which is a complex exponential factor, to form the final result.

In order to compute the DFT of the array  $x[n]$  using the FFT decimation-in-time algorithm the following function is defined:

```
function X = FFT_DIT(x)
N = length(x);
if N == 1
    X = x;
else
    % Divide
    x_even = x(1:2:end);
    x_odd = x(2:2:end);
    % Conquer
    X_even = FFT_DIT(x_even);
    X_odd = FFT_DIT(x_odd);
    % Combine
    WN = exp(-1i * 2 * pi / N);
    W = 1;
    X = zeros(1, N);
    for k = 1:N/2
        X(k) = X_even(k) + W * X_odd(k);
        X(k + N/2) = X_even(k) - W * X_odd(k);
        W = W * WN;
    end
end
end
```

In order to compute DFT of  $x[n]$ , the FFT DIT function is called as follows:

```
t3 = tic;
X_DIT = FFT_DIT(x);
elapsed_time = toc(t3);
fprintf('Elapsed time for FFT_DIT is %.6f seconds.\n', elapsed_time);
plotDFT(X_DIT, 'FFT-DIT')
```

f)

N point DFT defined as:

$$X[k] = \sum_{n=0}^{N-1} x_n W_N^{kn}$$

where:

$$W_N = e^{-i \frac{2\pi}{N}}$$

Decimation-in-frequency Fast Fourier Transform can be calculated by putting  $k = 2r$  in the summation formula then we get:

$$X[2r] = \sum_{n=0}^{N-1} x[n] W_N^{n(2r)}$$

Split into two sums:

$$X[2r] = \sum_{n=0}^{N/2-1} x[n] W_N^{2rn} + \sum_{n=0}^{N/2-1} x[n + N/2] W_N^{2r(n+N/2)}$$

Rearranging the terms:

$$X[2r] = \sum_{n=0}^{N/2-1} (x[n] + x[n + N/2]) W_{N/2}^{rn}$$

This  $N/2$  DFT of first and second half summed.  $X[2r+1]$  can be found similarly.

This can be repeated recursively.

In order to compute the DFT of the array  $x[n]$  using the FFT decimation-in-frequency algorithm the following function is defined:

```
function X = FFT_DIF(x)
N = length(x);
if N == 1
    X = x;
else
    X = x;
    % Butterfly stage
    for k = 1:N/2
        temp = X(k);
        X(k) = temp + X(k + N/2);
        X(k + N/2) = (temp - X(k + N/2)) * exp(-1i * 2 * pi * (k - 1) / N);
    end
    % Recursive stage
    X(1:N/2) = FFT_DIF(X(1:N/2));
    X(N/2+1:N) = FFT_DIF(X(N/2+1:N));
end
end
```

In order to compute DFT of  $x[n]$ , the FFT DIF function is called as follows:

```
t4 = tic;
X_DIF = bitrevorder(FFT_DIF(x));
elapsed_time = toc(t4);
fprintf('Elapsed time for FFT_DIF is %.6f seconds.\n', elapsed_time);
plotDFT(X_DIF, 'FFT-DIF')
```

**g)**

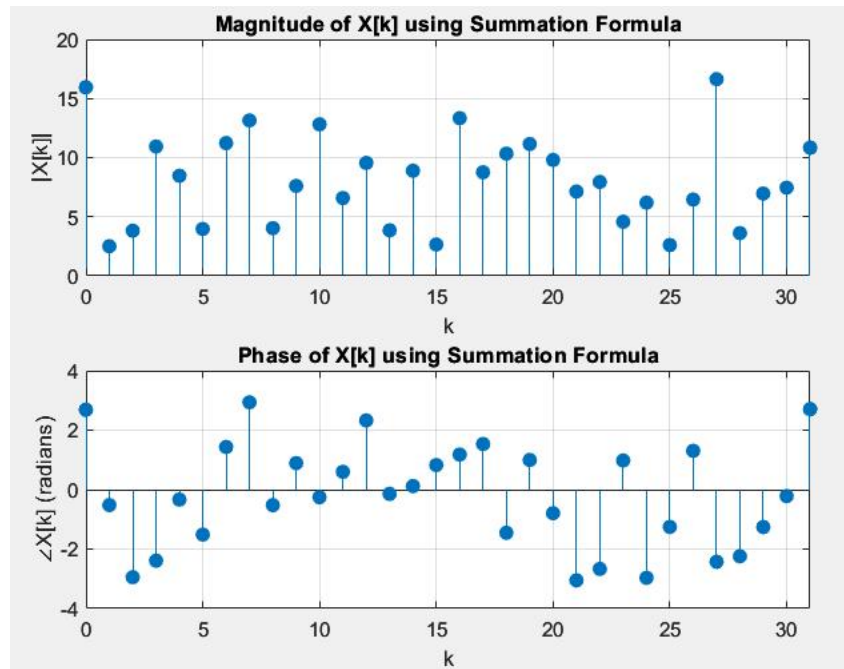
In order to compute DFT of  $x[n]$  using the MATLAB's built-in FFT command, the `fft` command is called as follows:

```
t5 = tic;
X_fft = fft(x);
elapsed_time = toc(t5);
fprintf('Elapsed time for fft is %.6f seconds.\n', elapsed_time);
plotDFT(X_fft, 'FFT')
```

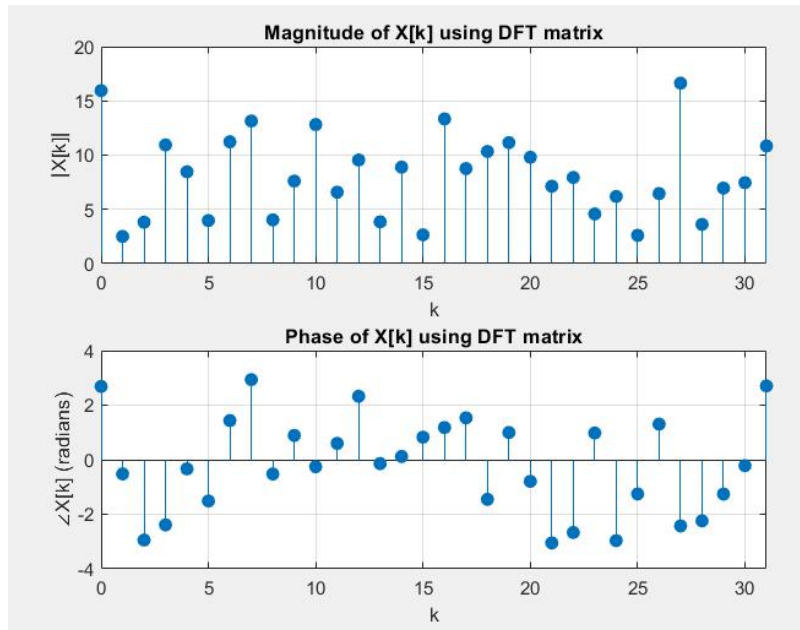
**h)**

For plotting purposes, **plotDFT(X, str)** function is defined. That takes the name of the algorithm (str) and DFT of  $x[n]$  ( $X[k]$ ) and plots the magnitude and phase graphs of  $X[k]$ . The function is visible in **Appendix A**.

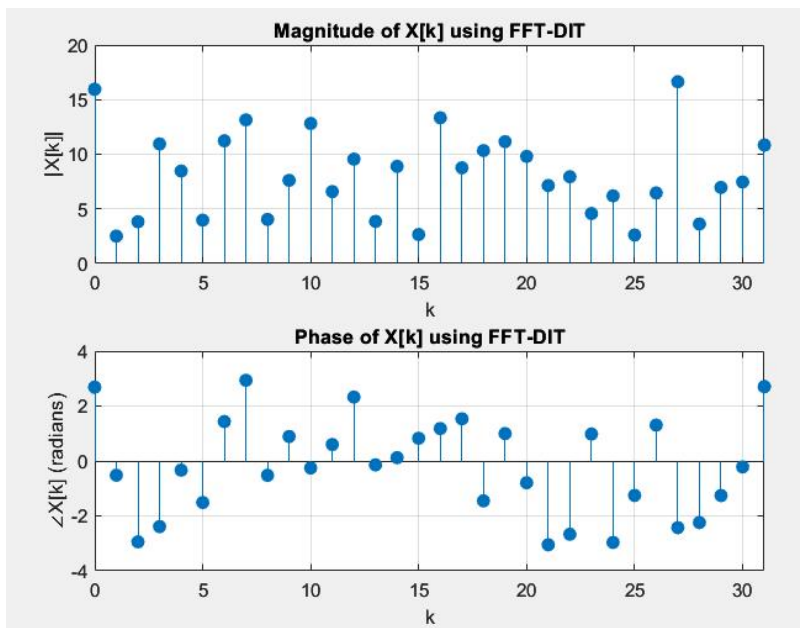
The magnitude and phase graphs of  $X[k]$  for each algorithm when  $N = 32$  are below in Figures 2-6:



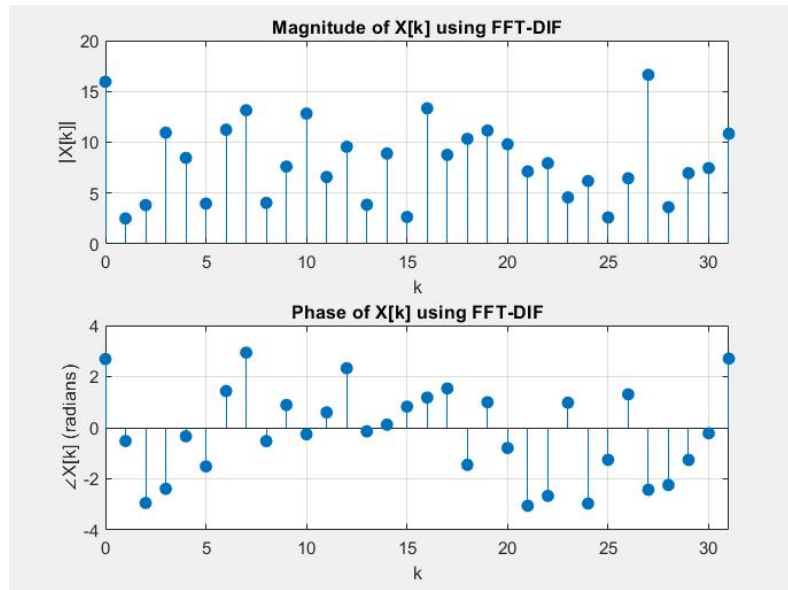
**Fig.2** DFT using summation formula,  $N = 32$



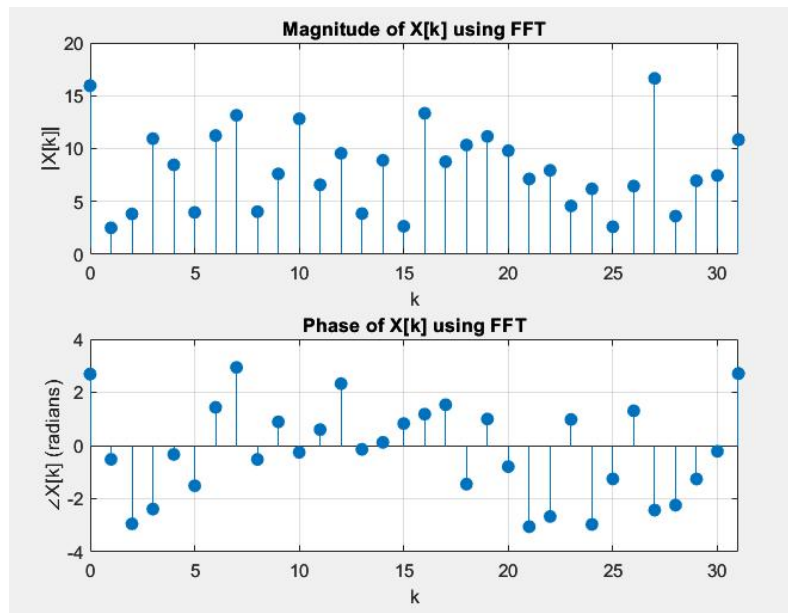
**Fig.3** DFT using DFT matrix,  $N = 32$



**Fig.4** DFT using FFT-DIT,  $N = 32$



**Fig.5** DFT using FFT-DIF,  $N = 32$



**Fig.6** DFT using FFT,  $N = 32$

Comparing Figures 2-6 it is visible that each algorithm produces equivalent results.

The algorithms are also compared using the difference of the norms between arrays. Each custom algorithm is compared with the built-in FFT command to check whether the norm is infinitesimally small or not (close to machine precision). The following code snippet computes the differences:

```
disp(['Summation Formula: ', num2str(norm(X_fft - Xsum))]);
disp(['DFT matrix: ', num2str(norm(X_fft - X_mat))]);
disp(['FFT-DIT: ', num2str(norm(X_fft - X_DIT))]);
disp(['FFT-DIF: ', num2str(norm(X_fft - X_DIF))]);
disp(['fft: ', num2str(norm(X_fft - X_fft))]);
```

And the output is as following:



```

Summation Formula: 3.1489e-13
DFT matrix:       4.8211e-13
FFT-DIT:          1.4408e-14
FFT-DIF:          1.2255e-14
fft:              0

```

**Fig.7** Difference of the norms between the algorithms,  $N = 32$

From Fig.7 it is evident that the difference of the norms are very close to zero, therefore, all the algorithms generates the same result.

i)

Now, steps (a) - (h) repeated for a array of length with  $N = 256$ . The vector is crated using the provided code snippet as follows:

```

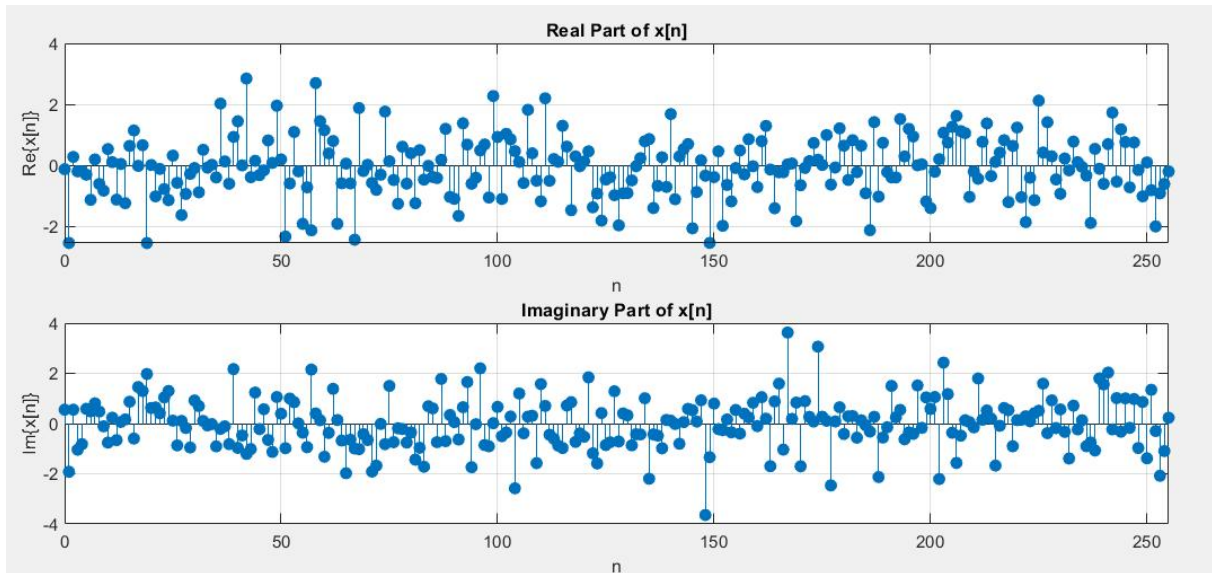
rng(2,"twister")
N = 256;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;

```

And the steps (a)-(h) are repeated using the above mentioned custom function **steps\_a\_h(x)** (see Appendix A):

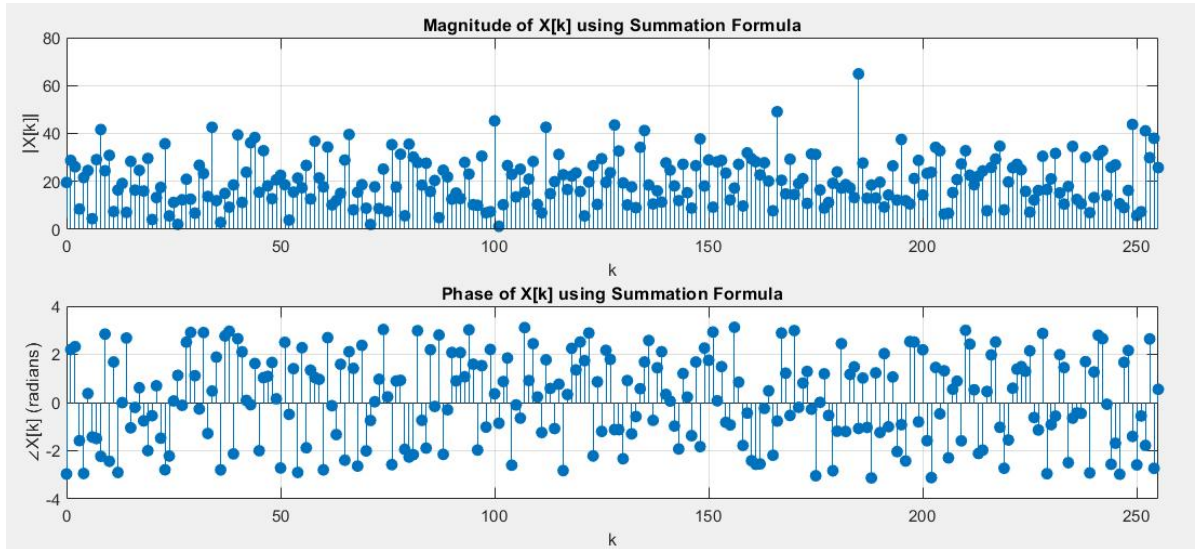
steps\_a\_h(x);

Real and imaginary parts of  $x[n]$  of length  $N = 256$  are plotted:

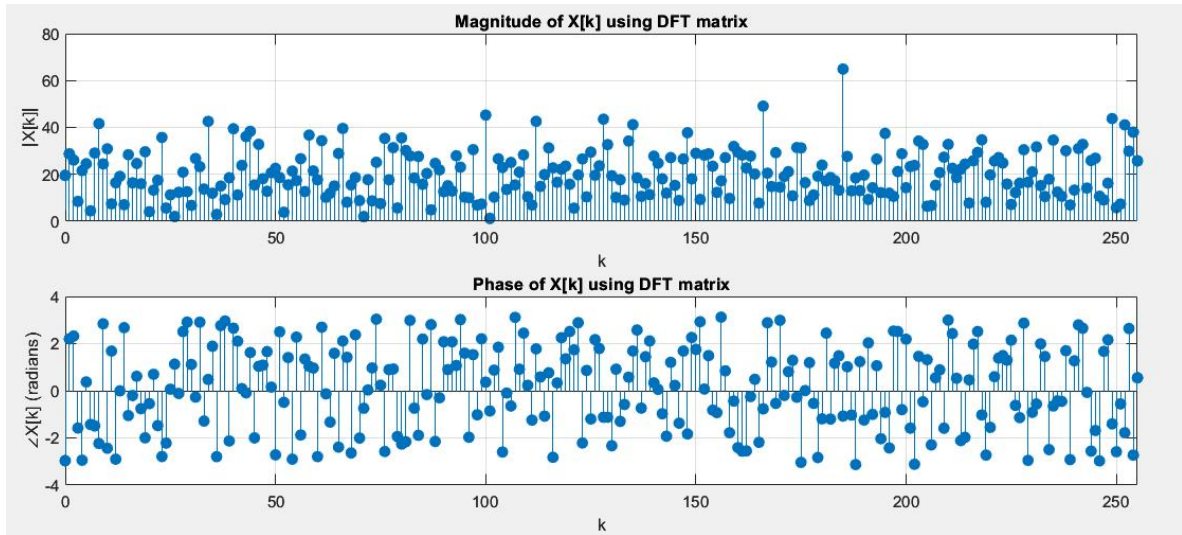


**Fig.8** Real and Imaginary parts of  $x[n]$ ,  $N= 256$

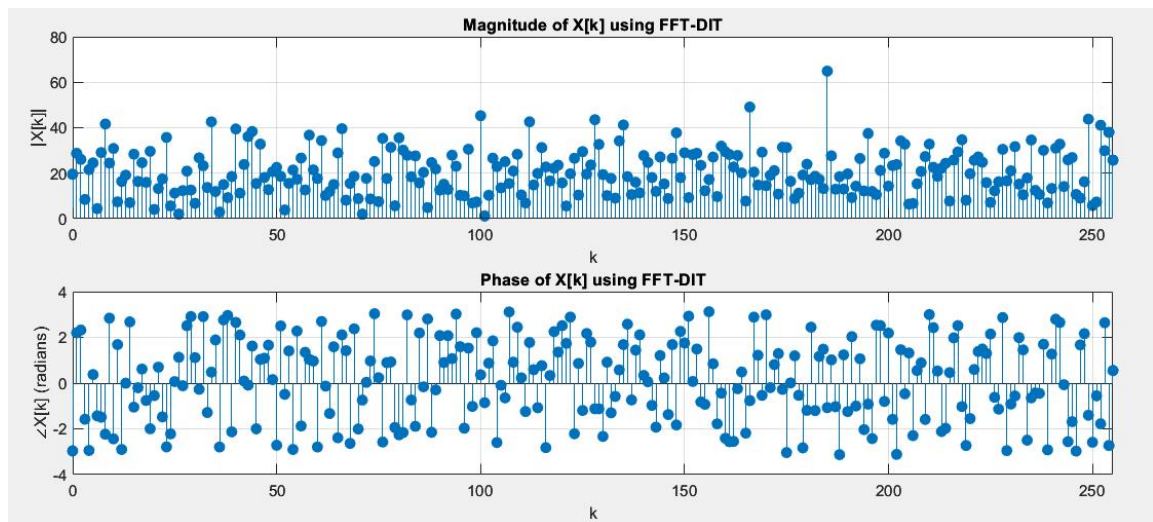
The magnitude and phase graph of the different algorithms:



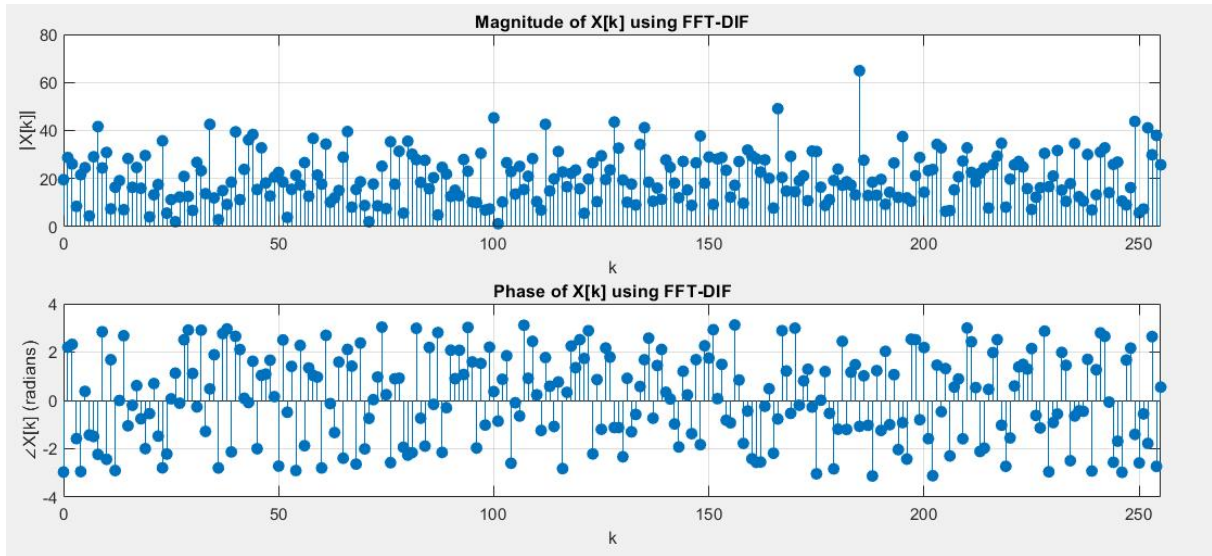
**Fig.9** DFT using summation formula,  $N = 256$



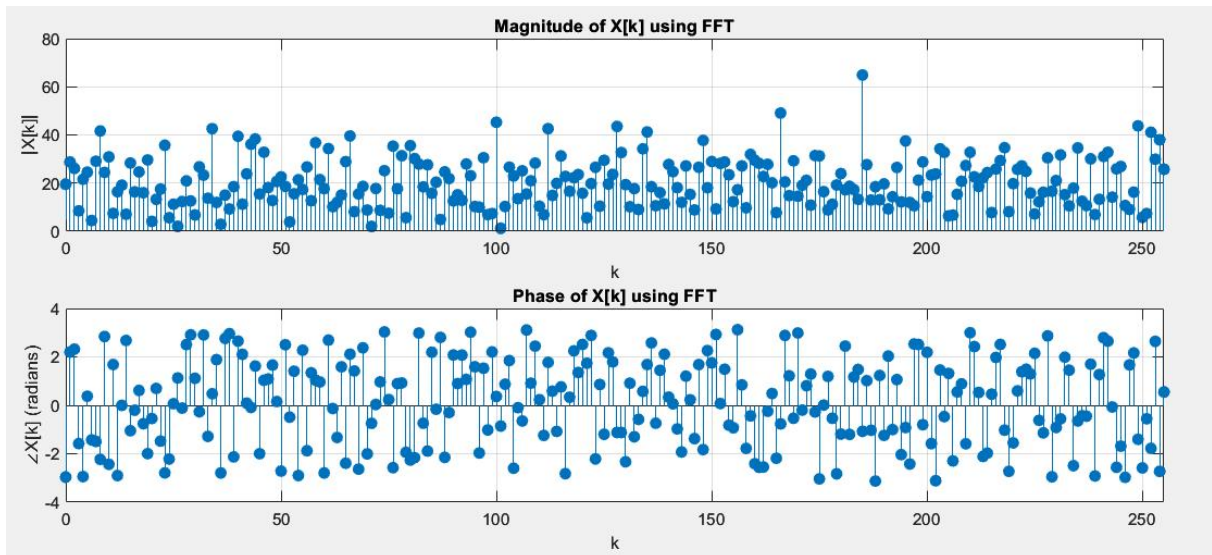
**Fig.10** DFT using DFT matrix,  $N = 256$



**Fig.11** DFT using FFT-DIT,  $N = 256$



**Fig.12** DFT using FFT-DIF,  $N = 256$



**Fig.13** DFT using FFT,  $N = 256$

Comparing Figures 9-13 it is visible that each algorithm produces equivalent results.

Differences of the norms of different algorithms:

Summation Formula:	1.5703e-11
DFT matrix:	1.1515e-10
FFT-DIT:	6.6354e-13
FFT-DIF:	1.4041e-13
fft:	0

**Fig.14** Difference of the norms between the algorithms,  $N = 256$

From Fig.14 it is evident that the difference of the norms are very close to zero, therefore, all the algorithms generates the same result.



j)

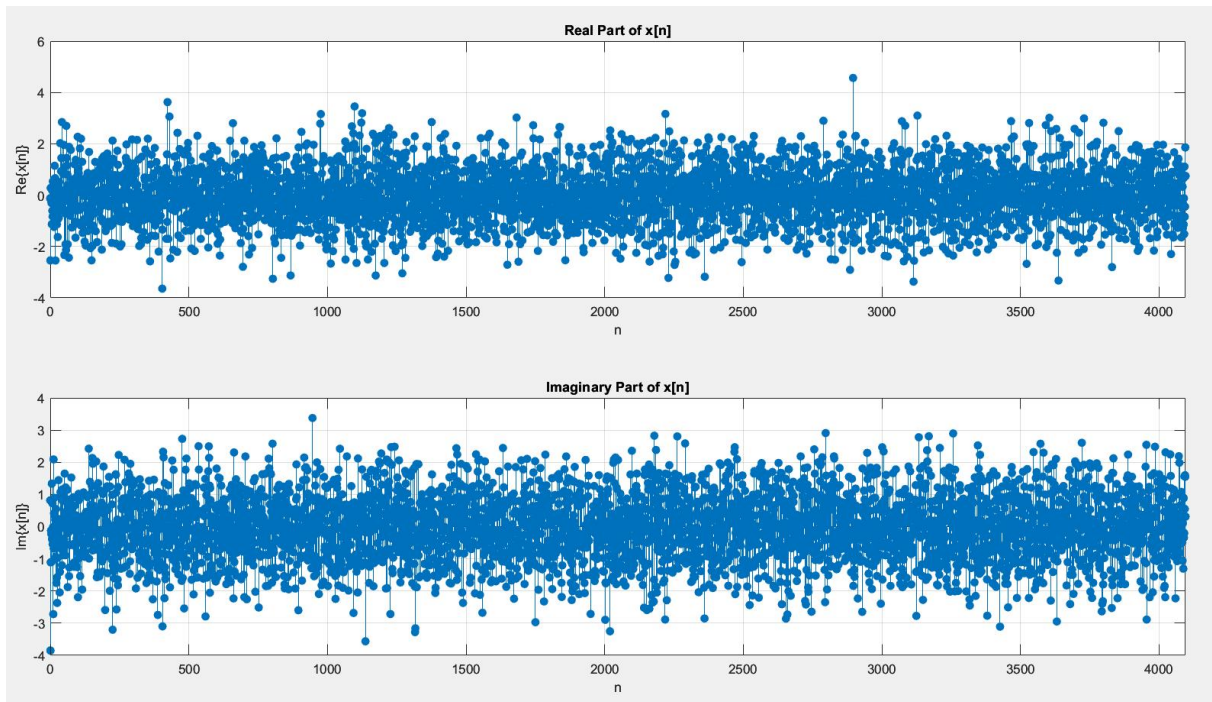
Now, steps (a) - (h) repeated for a array of length with  $N = 2^{12}$ . The vector is crated using the provided code snippet as follows:

```
rng(2,"twister")
N = 2^12;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;
```

And the steps (a)-(h) are repeated using the above mentioned custom function **steps\_a\_h(x)** (see Appendix A):

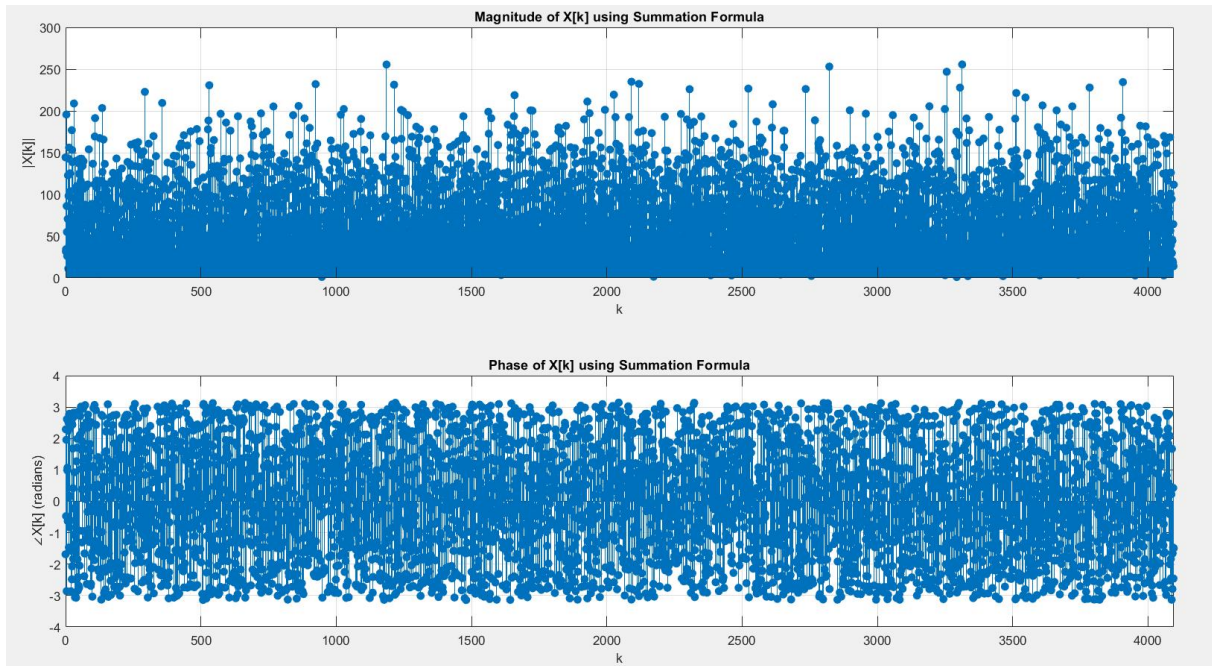
```
steps_a_h(x);
```

Real and imaginary parts of  $x[n]$  of length  $N = 2^{12}$  are plotted:

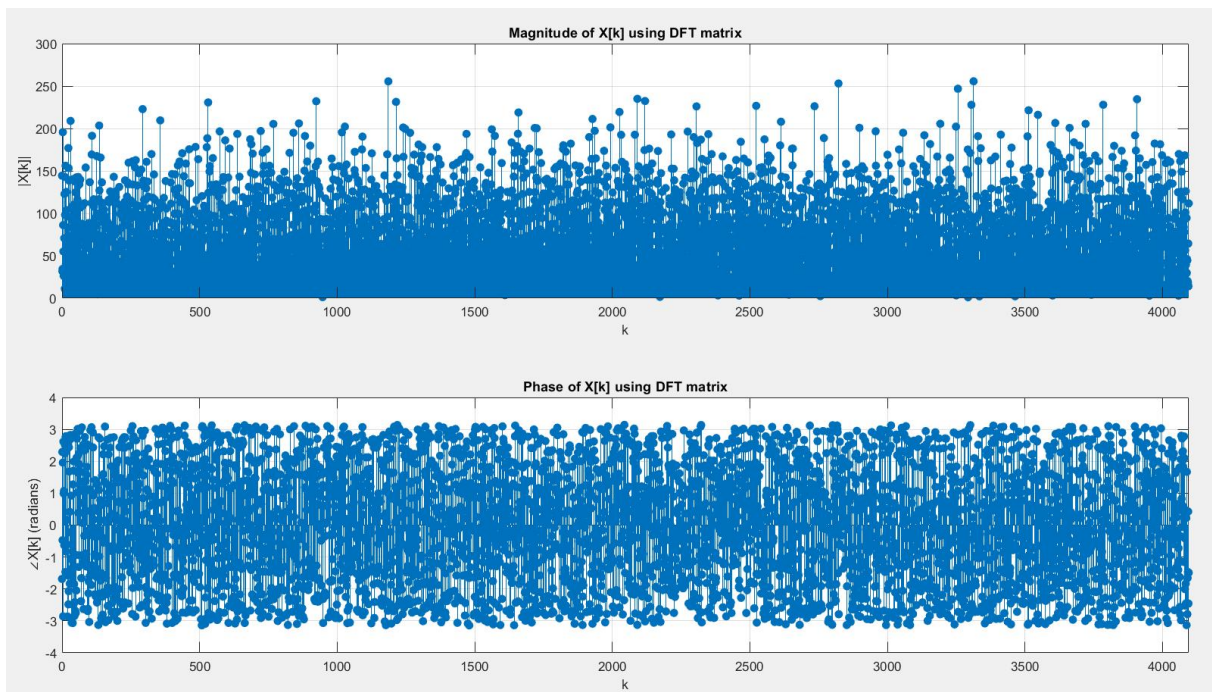


**Fig.15** Real and Imaginary parts of  $x[n]$ ,  $N = 2^{12}$

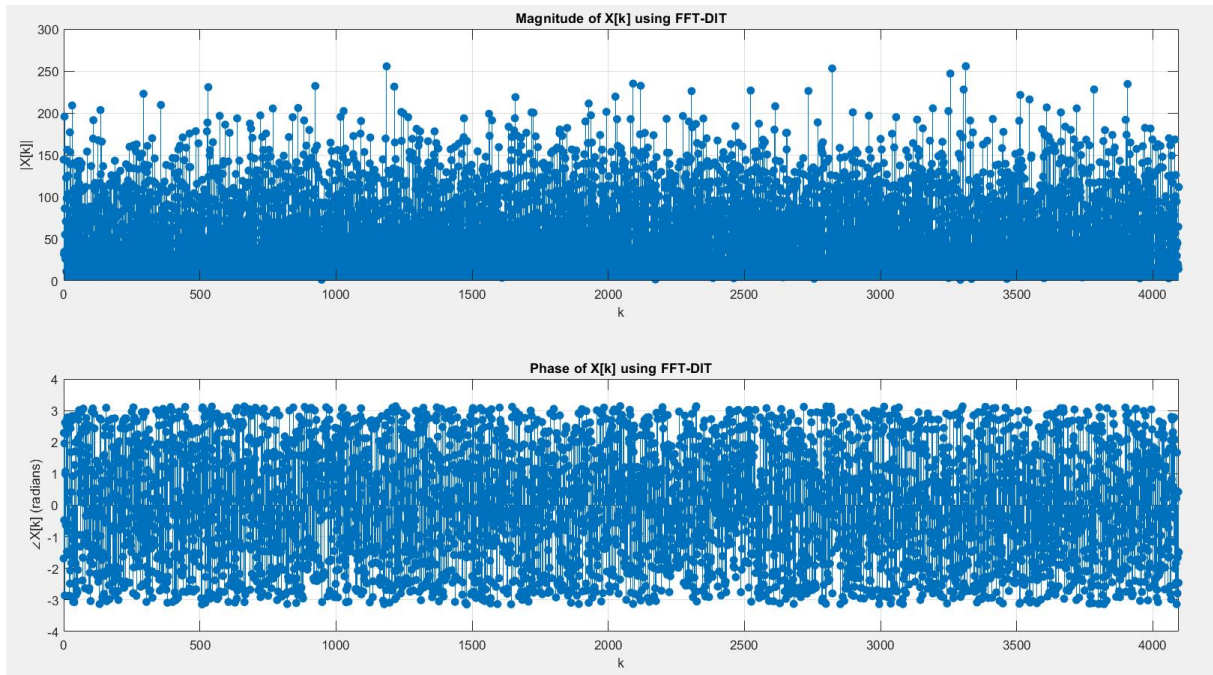
The magnitude and phase graph of the different algorithms:



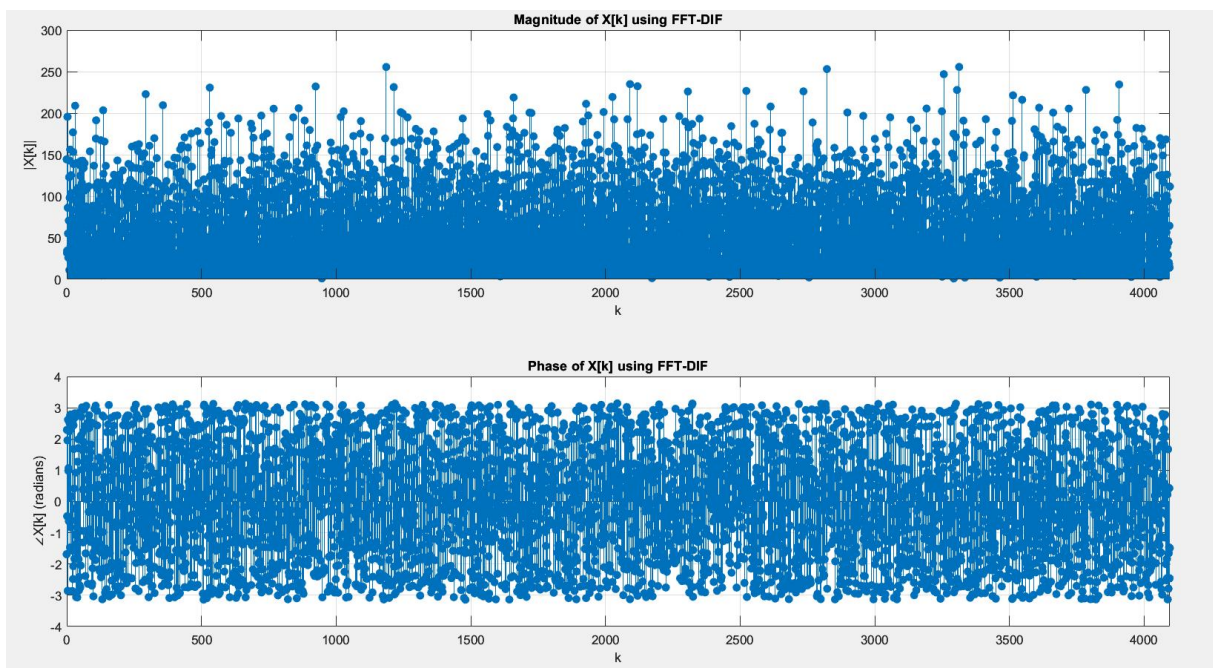
**Fig.16** DFT using summation formula,  $N = 2^{12}$



**Fig.17** DFT using DFT matrix,  $N = 2^{12}$

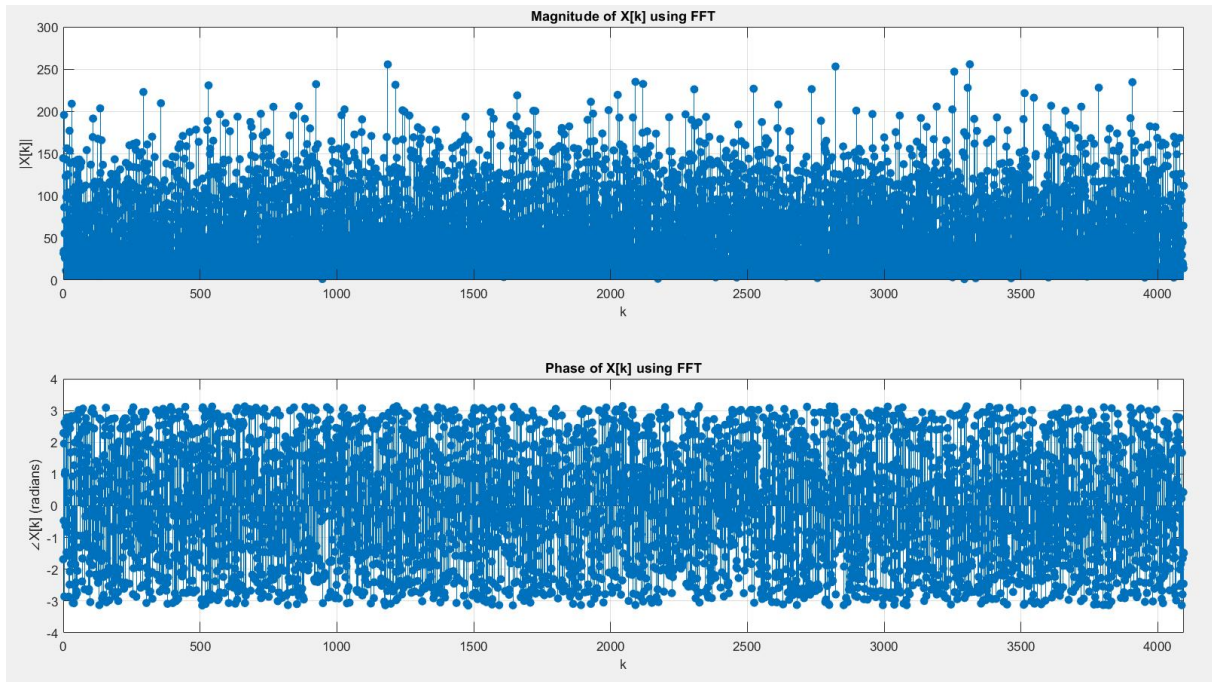


**Fig.18** DFT using FFT-DIT,  $N = 2^{12}$



**Fig.19** DFT using FFT-DIF,  $N = 2^{12}$





**Fig.20** DFT using FFT,  $N = 2^{12}$

Comparing Figures 16-20 it is visible that each algorithm produces equivalent results.

Differences of the norms of different algorithms:

```
Summation Formula:  3.9066e-09
DFT matrix:         1.1446e-06
FFT-DIT:            1.7455e-10
FFT-DIF:            3.0277e-12
fft:                 0
```

**Fig.21** Difference of the norms between the algorithms,  $N = 2^{12}$

From Fig.21 it is evident that the difference of the norms are very close to zero, therefore, all the algorithms generates the same result.

**k)**

Time it takes for each algorithm is calculated using *tic* and *toc* commands. The results are gathered in the Table 1.

	$N = 32$	$N = 256$	$N = 2^{12}$
Part (b)	0.001658 s	0.018429 s	4.120413 s
Part (d)	0.001597 s	0.034487 s	8.044837 s
Part (e)	0.000797 s	0.004292 s	0.019670 s
Part (f)	0.016621 s	0.004007 s	0.032825 s
Part (g)	0.000464 s	0.002475 s	0.000225 s

**Table 1:** Time it takes to compute different DFT approaches on signals of varying length

Table 1 shows the time taken by five different DFT approaches for three signal lengths ( $N = 32, 256$ , and  $4096$ ). As expected, the time increases with larger signal sizes, and there is a

significant variation in time across different methods. The direct DFT summation (part b) and DFT matrix (part d) exhibit a much slower computation time for larger arrays compared to the FFT-based approaches. Specifically, the FFT-DIT (part e) and FFT-DIF (part f) methods perform significantly faster as the signal length increases, highlighting the advantages of these algorithms in handling large data sets. The built-in MATLAB FFT (part g) shows the fastest computation time, aligning with theoretical expectations that FFT-based methods reduce complexity from  $O(N^2)$  in the direct summation to  $O(N \log N)$ . These findings confirm the efficiency of FFT algorithms, particularly as the signal length grows.

## Q2)

In this part, we are focusing on a uncommon scenario where we are computing FFT of a signal length  $N = 3^v$  instead of  $N = 2^v$ . In particular, 9-pt signal array. Whole code of this question is provided in **Appendix B**.

### a)

Following similar steps with part (f) where we computed FFT of a signal with length  $N = 2^v$  using the decimation-in-frequency algorithm.  
N point DFT defined as:

$$X[k] = \sum_{n=0}^{N-1} x_n W_N^{kn}$$

where:

$$W_N = e^{-i \frac{2\pi}{N}}$$

So in order to find  $X[k]$ , let  $k = 3r$ , where  $r = 0, \dots, N/3-1$

$$X[3r] = \sum_{n=0}^{N-1} x_n W_N^{(3r)n}$$

Split this over 3 sums with  $N/3$  intervals.

$$X[3r] = \sum_{n=0}^{N/3-1} x[n] W_N^{3rn} + \sum_{n=0}^{N/3-1} x[n + N/3] W_N^{3r(n+N/3)} + \sum_{n=0}^{N/3-1} x[n + 2N/3] W_N^{3r(n+2N/3)}$$

$$W_N^{3rn} = W_{N/3}^{rn}$$

$$W_N^{3r(n+N/3)} = W_N^{3rn} \cdot W_N^{Nr} = W_{N/3}^{rn}, \quad W_N^{Nr} = 1$$

$$W_N^{3r(n+2N/3)} = W_N^{3rn} \cdot W_N^{2Nr} = W_{N/3}^{rn}, \quad W_N^{2Nr} = 1$$

$$X[3r] = \sum_{n=0}^{N/3-1} (x[n] + x[n + N/3] + x[n + 2N/3]) W_{N/3}^{rn}$$

For  $k = 3r + 1$  where  $r = 0, \dots, N/3-1$

$$X[3r + 1] = \sum_{n=0}^{N-1} x_n W_N^{(3r+1)n}$$

Split this over 3 sums with  $N/3$  intervals.



$$X[3r+1] = \sum_{n=0}^{N/3-1} x[n] W_N^{(3r+1)n} + \sum_{n=0}^{N/3-1} x[n+N/3] W_N^{(3r+1)(n+N/3)} + \sum_{n=0}^{N/3-1} x[n+2N/3] W_N^{(3r+1)(n+2N/3)}$$

$$W_N^{(3r+1)n} = W_{N/3}^{rn} \cdot W_N^n$$

$$W_N^{(3r+1)(n+N/3)} = W_N^{3rn} \cdot W_N^n \cdot W_N^{Nr} \cdot W_N^{N/3} = W_{N/3}^{rn} \cdot W_N^n \cdot W_N^{N/3}, \quad W_N^{Nr} = 1$$

$$W_N^{(3r+1)(n+2N/3)} = W_N^{3rn} \cdot W_N^n \cdot W_N^{2Nr} \cdot W_N^{2N/3} = W_{N/3}^{rn} \cdot W_N^n \cdot W_N^{2N/3}, \quad W_N^{2Nr} = 1$$

$$X[3r+1] = \sum_{n=0}^{N/3-1} (x[n] + x[n+N/3] W_N^{N/3} + x[n+2N/3] W_N^{2N/3}) W_{N/3}^{rn} \cdot W_N^n$$

For  $k = 3r + 2$  where  $r = 0, \dots, N/3-1$

$$X[3r+2] = \sum_{n=0}^{N-1} x_n W_N^{(3r+2)n}$$

Split this over 3 sums with  $N/3$  intervals.

$$X[3r+2] = \sum_{n=0}^{N/3-1} x[n] W_N^{(3r+2)n} + \sum_{n=0}^{N/3-1} x[n+N/3] W_N^{(3r+2)(n+N/3)} + \sum_{n=0}^{N/3-1} x[n+2N/3] W_N^{(3r+2)(n+2N/3)}$$

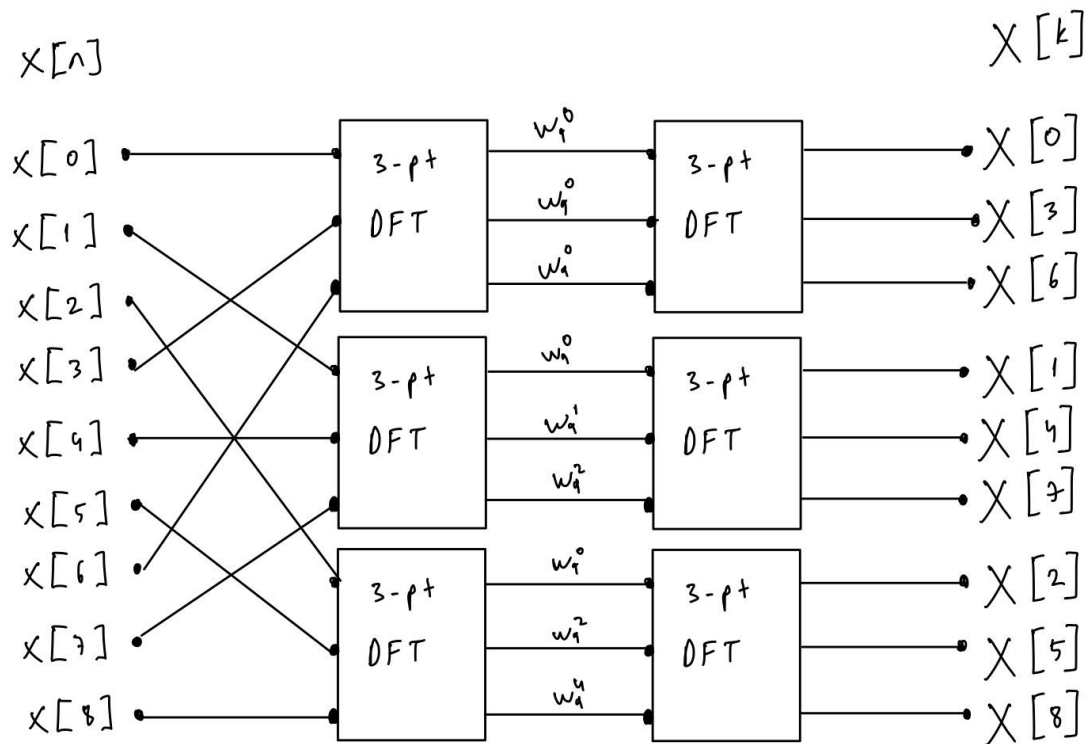
$$W_N^{(3r+2)n} = W_{N/3}^{rn} \cdot W_{N/2}^n$$

$$W_N^{(3r+2)(n+N/3)} = W_N^{3rn} \cdot W_{N/2}^n \cdot W_N^{Nr} \cdot W_N^{2N/3} = W_{N/3}^{rn} \cdot W_N^n \cdot W_N^{2N/3}, \quad W_N^{Nr} = 1$$

$$W_N^{(3r+2)(n+2N/3)} = W_N^{3rn} \cdot W_N^{2n} \cdot W_N^{2Nr} \cdot W_N^{4N/3} = W_{N/3}^{rn} \cdot W_{N/2}^n \cdot W_N^{4N/3}, \quad W_N^{2Nr} = 1$$

$$X[3r+2] = \sum_{n=0}^{N/3-1} (x[n] + x[n+N/3] W_N^{2N/3} + x[n+2N/3] W_N^{4N/3}) W_{N/3}^{rn} \cdot W_N^{2n}$$

b)



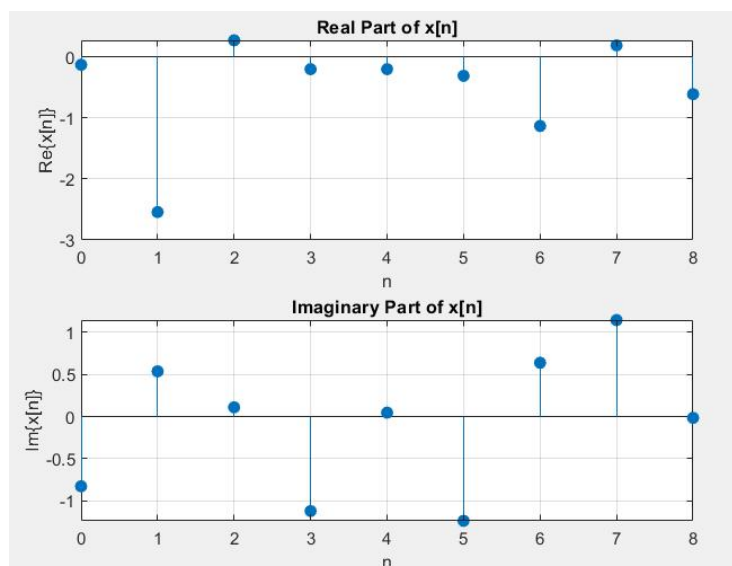
**Fig.21** Sample flow graph for the 9-pt decimation-in-frequency FFT algorithm.

c)

First, a complex array of length  $N = 9$  is created with the given code snippet:

```
rng(2, "twister")
N = 9;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;
```

Real and imaginary parts of  $x[n]$  are plotted:



**Fig.22** Real and Imaginary parts of  $x[n]$ ,  $N=9$ 

In order to compute DFT of array  $x$  using the DFT definition in summation form, the same function defined in **Q1) (b)**, **DFT\_summation(x)**, has been called as in the following code snippet:

```
t1 = tic;
Xsum = DFT_summation(x);
elapsed_time = toc(t1);
fprintf('Elapsed time for DFT_summation is %.6f seconds.\n', elapsed_time);
plotDFT(Xsum, 'Summation Formula')
```

**d)**

In order to implment the 9-pt FFT using decimation-in-frequency algorithm derived in part (a) the following function is defined:

```
function X = nineptFFT(x)
N = length(x);
X = zeros(1, N);
for r = 0:(N/3 - 1)
for p = 0:2
temp = 0;
for n = 0:(N/3 - 1)
inner_sum = 0;
for l = 0:2
inner_sum = inner_sum + x(n + l*N/3 + 1) * exp(-1j*2*pi*p*l/3);
end
temp = temp + inner_sum * exp(-1j*2*pi*n*r/(N/3)) * exp(-1j*2*pi*p*n/N);
end
X(3*r + p + 1) = temp;
end
end
end
```

The DFT of  $x[n]$  found with the following code snippet:

```
t2 = tic;
X_9pt_fft = nineptFFT(x);
elapsed_time = toc(t2);
fprintf('Elapsed time for 9-pt FFT is %.6f seconds.\n', elapsed_time);
plotDFT(X_9pt_fft, '9-pt FFT')
```

**e)**

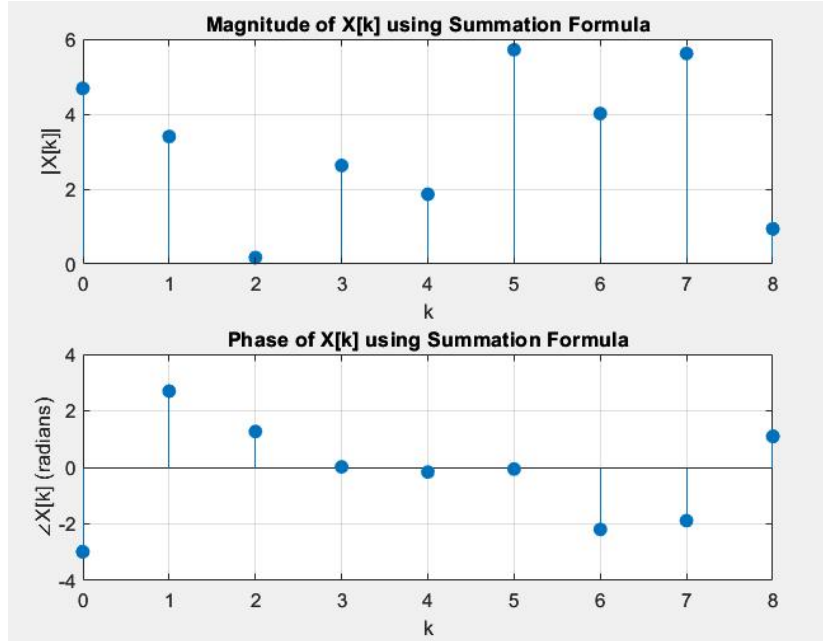
In order to compute DFT of  $x[n]$  using the MATLAB's built-in FFT command, the `fft` command is called as follows:

```
t3 = tic;
X_fft = fft(x);
elapsed_time = toc(t3);
fprintf('Elapsed time for fft is %.6f seconds.\n', elapsed_time);
plotDFT(X_fft, 'FFT')
```

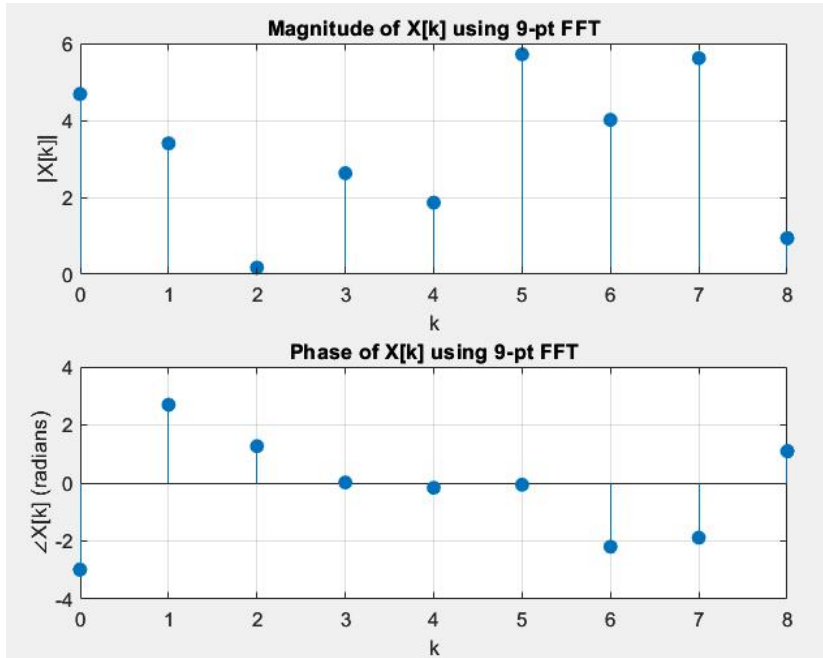
**f)**

In order to compare the results norm difference of the algorithms are computed following the same steps in **Q1** (h).

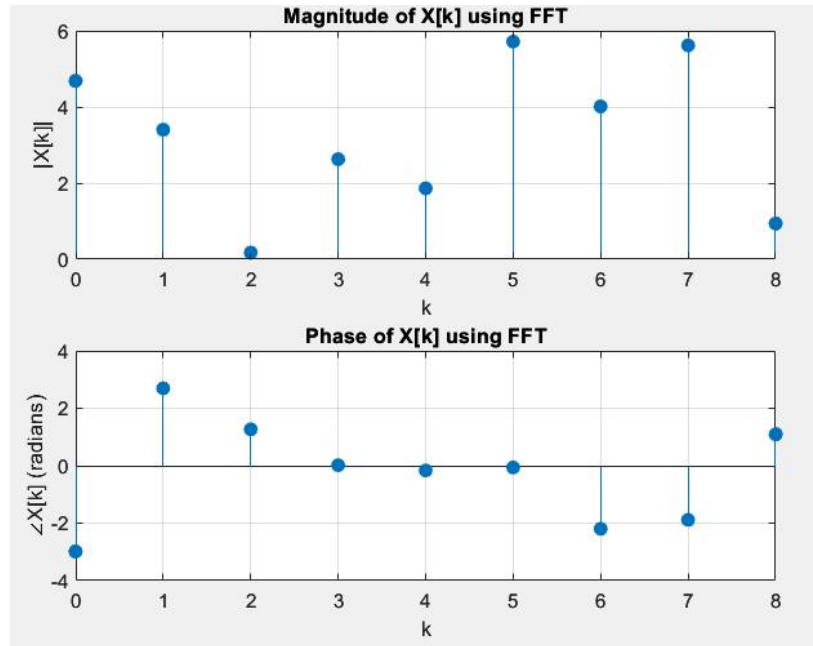
The magnitude and phase graph of the different algorithms:



**Fig.23** DFT using summation formula,  $N = 9$



**Fig.24** DFT using 9-pt FFT DIF,  $N = 9$



**Fig.25** DFT using FFT,  $N = 9$

Comparing Figures 23-25 it is visible that each algorithm produces equivalent results.

Differences of the norms of different algorithms:

```
Summation Formula: 1.2612e-14
DFT matrix:        6.9869e-15
fft:                0
```

**Fig.26** Difference of the norms between the algorithms,  $N = 9$

From Fig.26 it is evident that the difference of the norms are very close to zero, therefore, all the algorithms generates the same result.

**g)**

Time it takes for each algorithm is calculated using *tic* and *toc* commands. The results are gathered in the Table 2.

Part	Time
Part (c)	0.000041 s
Part (d)	0.000074 s
Part (e)	0.000013 s

**Table 2:** Time it takes to compute different DFT approaches,  $N = 9$

Table 2 shows the time taken by three different DFT approaches for DFT of length  $N = 9$ . As expected, there is a significant variation in time across different methods. These times indicate that the MATLAB built-in FFT function is still the fastest, as expected, with the custom 9-pt FFT slightly slower than both the summation method and the built-in function. Despite the minimal time differences, the overall performance aligns with expectations: built-in FFTs are highly optimized, while custom methods like the 9-pt FFT have higher overheads. 9-pt FFT DIF algorithm is slower than the direct summation because array length is too small  $N = 9$ . In

**Q1)** (k) it is shown that for  $N = 32$  FFT DIF algorithm is slightly slower than the direct summation but for  $N = 2^{12}$  FFT DIF algorithm is significantly faster. Therefore, for a larger array like  $N = 3^8$  the result could have been significantly faster.

## Appendices

### Appendix A - Q1 whole code

```

close all
clear
clc

rng(2,"twister")
N = 32;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;
steps_a_h(x);

rng(2,"twister")
N = 256;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;
steps_a_h(x);

rng(2,"twister")
N = 2^12;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;
steps_a_h(x);

function steps_a_h(x)
N = length(x);
fprintf('N = %d \n', N);
n = 0:N-1;

figure;
subplot(2,1,1);
stem(n, real(x), 'filled');
title('Real Part of x[n]');
xlabel('n');
ylabel('Re\{x[n]\}');
xlim([0 N-1]);
grid on;

subplot(2,1,2);
stem(n, imag(x), 'filled');
title('Imaginary Part of x[n]');
xlabel('n');
ylabel('Im\{x[n]\}');
xlim([0 N-1]);
grid on;

t1 = tic;
Xsum = DFT_summation(x);
elapsed_time = toc(t1);
fprintf('Elapsed time for DFT_summation is %.6f seconds.\n', elapsed_time);
plotDFT(Xsum, 'Summation Formula')

t2 = tic;
X_mat = DFT_matrix(x);

```

```

elapsed_time = toc(t2);
fprintf('Elapsed time for DFT_matrix is %.6f seconds.\n', elapsed_time);
plotDFT(X_mat, 'DFT matrix')

t3 = tic;
X_DIT = FFT_DIT(x);
elapsed_time = toc(t3);
fprintf('Elapsed time for FFT_DIT is %.6f seconds.\n', elapsed_time);
plotDFT(X_DIT, 'FFT-DIT')

t4 = tic;
X_DIF = bitrevorder(FFT_DIF(x));
elapsed_time = toc(t4);
fprintf('Elapsed time for FFT_DIF is %.6f seconds.\n', elapsed_time);
plotDFT(X_DIF, 'FFT-DIF')

t5 = tic;
X_fft = fft(x);
elapsed_time = toc(t5);
fprintf('Elapsed time for fft is %.6f seconds.\n', elapsed_time);
plotDFT(X_fft, 'FFT')
fprintf('\n');

disp(['Summation Formula: ', num2str(norm(X_fft - Xsum))]);
disp(['DFT matrix: ', num2str(norm(X_fft - X_mat))]);
disp(['FFT-DIT: ', num2str(norm(X_fft - X_DIT))]);
disp(['FFT-DIF: ', num2str(norm(X_fft - X_DIF))]);
disp(['fft: ', num2str(norm(X_fft - X_fft))]);
fprintf('\n');
end

function plotDFT(X, str)
N = length(X);
n = 0:N-1;

figure;
subplot(2,1,1);
stem(n, abs(X), 'filled');
title(['Magnitude of X[k] using ', str]);
xlabel('k');
ylabel('|X[k]|');
xlim([0 N-1]);
grid on;

subplot(2,1,2);
stem(n, angle(X), 'filled');
title(['Phase of X[k] using ', str]);
xlabel('k');
ylabel('∠X[k] (radians)');
xlim([0 N-1]);
grid on;
end

function X = FFT_DIF(x)
N = length(x);
if N == 1
X = x;
else
X = x;

```



```

% Butterfly stage
for k = 1:N/2
    temp = X(k);
    X(k) = temp + X(k + N/2);
    X(k + N/2) = (temp - X(k + N/2)) * exp(-1i * 2 * pi * (k - 1) / N);
end
% Recursive stage
X(1:N/2) = FFT_DIF(X(1:N/2));
X(N/2+1:N) = FFT_DIF(X(N/2+1:N));
end
end

```

```

function X = FFT_DIT(x)
N = length(x);
if N == 1
    X = x;
else
    % Divide
    x_even = x(1:2:end);
    x_odd = x(2:2:end);
    % Conquer
    X_even = FFT_DIT(x_even);
    X_odd = FFT_DIT(x_odd);
    % Combine
    WN = exp(-1i * 2 * pi / N);
    W = 1;
    X = zeros(1, N);
    for k = 1:N/2
        X(k) = X_even(k) + W * X_odd(k);
        X(k + N/2) = X_even(k) - W * X_odd(k);
        W = W * WN;
    end
end
end

```

```

function X = DFT_matrix(x)
N = length(x);
WN = exp(-1i * 2 * pi / N);
DFT_mat = zeros(N, N);
for k = 0:N-1
    for n = 0:N-1
        DFT_mat(k+1, n+1) = WN^(k * n);
    end
end
X = (DFT_mat * x.').';
end

```

```

function X = DFT_summation(x)
N = length(x);
X = zeros(1, N);
for k = 0:N-1
    for n = 0:N-1
        X(k+1) = X(k+1) + x(n+1) * exp(-1i * 2 * pi * k * n / N);
    end
end
end

```

## Appendix B - Q2 whole code

```

close all
clear
clc

rng(2, "twister")
N = 9;
real_part = randn(1,N);
imag_part = randn(1,N);
x = real_part + 1i*imag_part;

fprintf('N = %d \n', N);
n = 0:N-1;

figure;
subplot(2,1,1);
stem(n, real(x), 'filled');
title('Real Part of x[n]');
xlabel('n');
ylabel('Re\{x[n]\}');
xlim([0 N-1]);
grid on;

subplot(2,1,2);
stem(n, imag(x), 'filled');
title('Imaginary Part of x[n]');
xlabel('n');
ylabel('Im\{x[n]\}');
xlim([0 N-1]);
grid on;

t1 = tic;
Xsum = DFT_summation(x);
elapsed_time = toc(t1);
fprintf('Elapsed time for DFT_summation is %.6f seconds.\n', elapsed_time);
plotDFT(Xsum, 'Summation Formula')

t2 = tic;
X_9pt_fft = nineptFFT(x);
elapsed_time = toc(t2);
fprintf('Elapsed time for 9-pt FFT is %.6f seconds.\n', elapsed_time);
plotDFT(X_9pt_fft, '9-pt FFT')

t3 = tic;
X_fft = fft(x);
elapsed_time = toc(t3);
fprintf('Elapsed time for fft is %.6f seconds.\n', elapsed_time);
plotDFT(X_fft, 'FFT')

fprintf('\n');
disp(['Summation Formula: ', num2str(norm(X_fft - Xsum))]);
disp(['DFT matrix: ', num2str(norm(X_fft - X_9pt_fft))]);
disp(['fft: ', num2str(norm(X_fft - X_fft))]);
fprintf('\n');

function X = nineptFFT(x)
N = length(x);
X = zeros(1, N);

```

```

for r = 0:(N/3 - 1)
for p = 0:2
temp = 0;
for n = 0:(N/3 - 1)
inner_sum = 0;
for l = 0:2
inner_sum = inner_sum + x(n + l*N/3 + 1) * exp(-1j*2*pi*p*l/3);
end
temp = temp + inner_sum * exp(-1j*2*pi*n*r/(N/3)) * exp(-1j*2*pi*p*n/N);
end
X(3*r + p + 1) = temp;
end
end
end

```

```

function X = DFT_summation(x)
N = length(x);
X = zeros(1, N);
for k = 0:N-1
for n = 0:N-1
X(k+1) = X(k+1) + x(n+1) * exp(-1i * 2 * pi * k * n / N);
end
end
end

```

```

function plotDFT(X, str)
N = length(X);
n = 0:N-1;

figure;
subplot(2,1,1);
stem(n, abs(X), 'filled');
title(['Magnitude of X[k] using ', str]);
xlabel('k');
ylabel('|X[k]|');
xlim([0 N-1]);
grid on;

subplot(2,1,2);
stem(n, angle(X), 'filled');
title(['Phase of X[k] using ', str]);
xlabel('k');
ylabel('∠X[k] (radians)');
xlim([0 N-1]);
grid on;
end

```