

HACETTEPE UNIVERSITY COMPUTER ENGINEERING DEPARTMENT

BBM204 SOFTWARE PRACTICUM II - 2024 SPRING

Programming Assignment 1

March 22, 2024

Student name: BORAN ÇAKIL Student Number: b2210356071

1 Problem Definition

In this project we are trying to compare several sorting and searching algorithms. We have a set of value and we will use our sorting and searching algorithms on it to test how the algorithms react with different cases.

2 Solution Implementation

We will use 3 sorting and 2 searching algorithms.

2.1 Sorting Algorithm 1: Counting Sort

```
public class CountingSort {
2
       public static double[] sort(double[] array, int k) {
3
            double count[] = new double[k + 1];
4
            double output[] = new double[array.length];
5
            int size = array.length;
6
7
            for (int i = 0; i < size; i++) {</pre>
                int j = (int) array[i];
8
                count[j]++;
9
10
            for (int i = 1; i < k + 1; i++) {
11
                count[i] += count[i - 1];
12
13
            for (int i = size - 1; i >= 0; i--) {
14
                int j = (int) array[i];
15
                count[j] = count[j] - 1;
16
                output[(int) count[j]] = array[i];
17
18
            return output;
19
       }
20
21
22
```

2.2 Sorting Algorithm 2: Insertion Sort

```
public class InsertionSort {
24
       public static void sort(double[] values) {
25
            for (int j = 1; j < values.length; j++) {</pre>
26
27
                double key = values[j];
                int i = j - 1;
28
                while (i \ge 0 \&\& values[i] > key) {
29
                     values[i + 1] = values[i];
30
31
                     i--;
```

2.3 Sorting Algorithm 3: Merge Sort

```
public class MergeSort {
38
39
40
       public static double[] sort(double[] array) {
            int n = array.length;
41
            if (n <= 1) {
42
                return array;
43
            }
44
            double[] left = new double[n / 2];
45
            double[] right = new double[n - n / 2];
46
            left = sort(left);
47
            right = sort(right);
48
            return merge(left, right);
49
50
       private static double[] merge(double[] a, double[]b) {
51
            double[] c = new double[a.length + b.length];
52
53
             int i = 0, j = 0, k = 0;
54
             while (i < a.length && j < b.length) {</pre>
55
                  if (a[i] > b[j]) {
56
                      c[k++] = b[j++];
57
                  } else {
                      c[k++] = a[i++];
59
60
             }
61
62
             while (i < a.length) {</pre>
63
64
                  c[k++] = a[i++];
65
66
             while (j < b.length) {</pre>
67
                 c[k++] = b[j++];
68
             }
69
            return c;
70
       }
71
72
```

2.4 Search Algorithm 1: Linear Search

```
public class LinearSearch {
74
       public static double search(double[] array, double x) {
75
            int size = array.length;
76
            for (int i = 0; i < size; i++) {</pre>
77
                if (array[i] == x) {
78
                     return i;
79
80
81
            return -1;
82
83
        }
84
```

2.5 Search Algorithm 2: Binary Search

```
public class BinarySearch {
87
88
        public static int binarySearch(double[] array, double x) {
89
             int low = 0;
90
             int high = array.length - 1;
91
92
             while (high - low > 1) {
93
                 int mid = (high + low) / 2;
94
                 if (array[mid] < x) {</pre>
                      low = mid + 1;
96
                  } else {
97
                      high = mid;
98
99
             }
100
101
             if (array[low] == x) {
102
                 return low;
103
             } else if (array[high] == x) {
104
                 return high;
105
106
107
             return -1;
        }
109
110
111
```

3 Results, Analysis, Discussion

By the results I understood that every algorithm can be better for special cases. Like linear sort is does not waste any time if we have an array that is nearly sorted. And while testing I understood that each time we ran the program the values differ because there is so much going on at the background. Because of that when we want to measure the time precisely we have to run the algorithms multiple time and find the average time.

Running time test results for sorting algorithms are given in Table 1.

Table 1: Results of the running time tests performed for varying input sizes (in ms).

Input Size n Algorithm 500 1000 2000 4000 8000 16000 32000 64000 128000 250000 Random Input Data Timing Results in ms Insertion sort 0.0 0.0 1.6 2.0 5.6 20.2 81.5 348.3 1563.4 6765.2 $0.\overline{2}$ Merge sort 0.0 0.0 0.0 1.6 0.0 1.7 1.5 6.4 11.1 204.8Counting sort 234.0 207.1 205.0 203.2 205.8 213.0 201.8 213.2 260.4 Sorted Input Data Timing Results in ms Insertion sort 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.4 0.0 0.1 0.0 0.0 0.0 0.0 2.9 2.1 Merge sort 0.0 0.0 13.4 Counting sort 211.6 210.0 212.4 214.5 208.2 205.6 206.8 209.0 221.0 208.7 Reversely Sorted Input Data Timing Results in msInsertion sort 0.0 0.0 159.1 3190.2 12357.1 1.7 8.3 38.6 742.0Merge sort 0.0 0.0 0.0 0.0 0.1 0.1 1.6 13.5 1.9 4.8 Counting sort 224.7 199.5 205.3 211.0 209.2 210.8 202.6 207.6 211.7 225.7

Running time test results for search algorithms are given in Table 2.

Table 2: Results of the running time tests of search algorithms of varying sizes (in ns).

					mp	it size n				
Algorithm	500	1000	2000	4000	8000	16000	32000	64000	128000	250000
Linear search (random data)	747500.0	585200.0	925300.0	1963300.0	4385300.0	6227900.0	1.06722E7	1.92541E7	3.30402E7	5.20989E7
Linear search (sorted data)	261100.0	443400.0	929400.0	3240400.0	5012700.0	8823700.0	1.26335E7	2.53788E7	4.5848E7	7.44258E7
Binary search (sorted data)	571700.0	99800.0	113000.0	140500.0	166700.0	172500.0	229600.0	216100.0	272900.0	290300.0

Complexity analysis tables to complete (Table 3 and Table 4):

Table 3: Computational complexity comparison of the given algorithms.

Algorithm	Best Case	Average Case	Worst Case
Insertion sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$
Merge sort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n \log n)$
Counting Sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$
Linear Search	$\Omega(1)$	$\Theta(n)$	O(n)
Binary Search	$\Omega(n)$	$\Theta(logn)$	$O(\log n)$

Table 4: Auxiliary space complexity of the given algorithms.

Algorithm	Auxiliary Space Complexity
Insertion sort	O(1)
Merge sort	O(n)
Counting sort	O(n)
Linear Search	O(1)
Binary Search	O(1)

3 sorting algorithms used on randomly generated data.: Fig. 1.

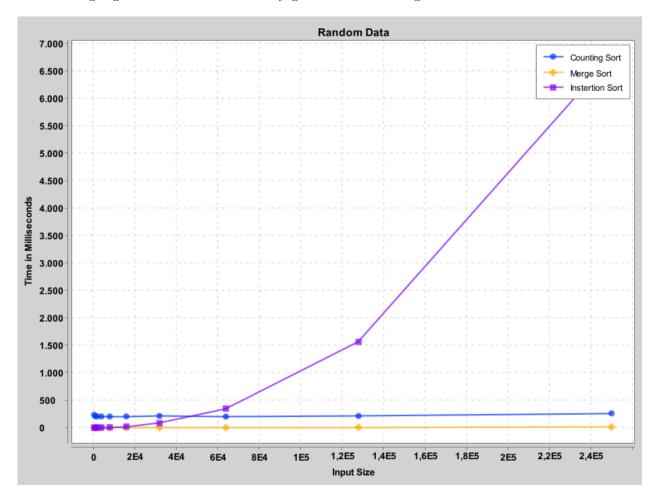


Figure 1: Sorting algorithms on randomly generated data

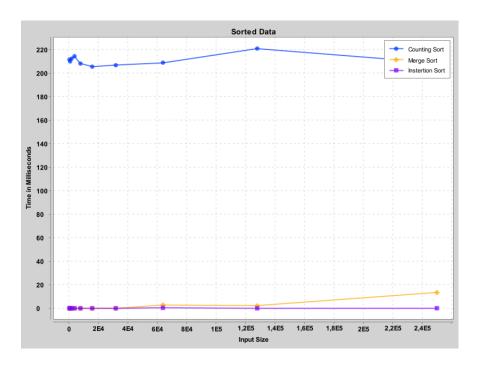


Figure 2: Sorting algorithms on sorted data

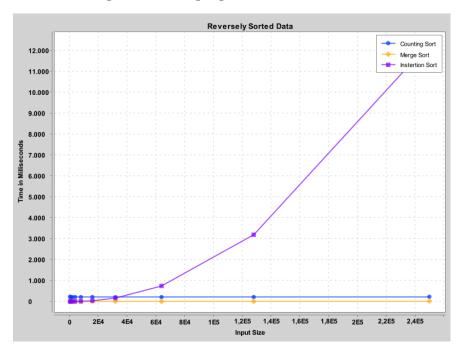


Figure 3: Sorting algorithms on reversely sorted data

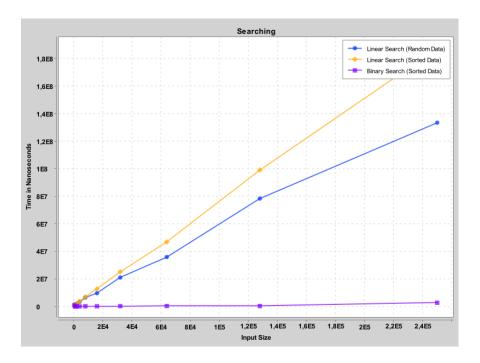


Figure 4: Searching algorithms on data

4 Notes

Algorithms react differently on different data sizes and types. It is important to use the right one for the right cause.

References

 \bullet https://www.geeksforgeeks.org