Implementation:

Syscalls:

I added these inside the SyscallHandler class. The syscall handler enum is used to make system calls using inline assembly in the syscall functions in userspace.

Like: asm("int \$0x80" :: "a" (SYSCALL::FORK));

I included a pointer to the task manager to call some task manager functions when an interrupt that needs those functions is occurred.

```
232 uint32_t sysGetPid()
233 {
234         asm("int $0x80" :: "a" (SYSCALL::GETPID));
235 }
236
237 void sysExit()
238         asm("int $0x80" :: "a" (SYSCALL::EXIT));
240 }
```

I added these functions in the kernel.cpp file. (sysPrintf was already in there, I just changed the call input from 4 to SYSCALL::PRINT.)

Here, these functions cause 0x80 (syscall) interrupts by putting the desired call into eax register and some other arguments to other registers if necessary. These interrupts are handled by the SyscallHandler class, if the system call need to return a value, the eax register is set directly in the handler.

```
23 uint32_t SyscallHandler::HandleInterrupt(uint32_t esp)
        CPUState* cpu = (CPUState*)esp;
        switch(cpu \rightarrow eax)
            case SYSCALL:: FORK:
                 printf("syscall fork called\n");
                 tm→Fork(cpu);
            case SYSCALL::WAITPID:
                 printf("syscall waitpid called\n");
                 cpu \rightarrow eax = tm \rightarrow WaitPid(cpu \rightarrow ebx);
                 esp = (uint32_t)tm→Schedule((CPUState*)esp); // c
            case SYSCALL :: EXECVE:
                 printf("syscall execve called\n");
                 esp = (uint32_t) tm→ChangeCurrentTask(cpu→ebx);
            case SYSCALL:: PRINT:
                 printf((char*)cpu \rightarrow ebx);
                 break;
            case SYSCALL:: GETPID:
                 cpu \rightarrow eax = tm \rightarrow GetPid();
                 break;
            case SYSCALL :: EXIT:
                 tm→Exit();
54
                 esp = (uint32_t)tm→Schedule((CPUState*) esp);
                 break;
            default:
                 break:
        return esp;
63 }
```

This is the HandleInterrupt function in the SyscallHandler class. I called the necessary functions in the taskmanager for each wanted system call (taken from cpu->eax), and then I put the return value to the return register (cpu->eax) if that system call has to return a value.

Exception here is the Fork call. Because the child and parent process after fork has to return different values, I set them in the Fork function in the taskmanager directly.

Wait and Exit syscalls call scheduler explicitly because those tasks won't be running after the call and we need to decide the next task.

TaskManager:

Task:

I added variables for state, pid, parent pid, the pid of the child that this task is waiting for. I added a default constructor explicitly because I hold the tasks in the TaskManager directly (not pointers to them.) I did this to make sure that the tasks are in the stack.

```
class TaskManager
{

private:

Task tasks[256];

int numTasks;

int currentTask;

int newPid;

GlobalDescriptorTable* gdt;

Task* TaskFromPid(common::uint32_t pid);
```

I added a pid counter named newPid to set the pid of a new task.

## Task Manager Functions:

#### Schedule:

```
99 CPUState* TaskManager:: Schedule(CPUState* cpustαte)
        printf("Schedular is called\n");
        // PrintTasks(tasks, numTasks);
        if(numTasks ≤ 0)
            return cpustate;
        if(currentTask ≥ 0)
            tasks[currentTask].cpustate = cpustate;
        // stop the current process.
        if((tasks[currentTask].state ≠ Task::BLOCKED) && (tasks[currentTask].state ≠ Task::TERMINATED))
            tasks[currentTask].state = Task::READY; // maybe some checks here.
        // (the quanta is the time between hardware timer interrupts)
        // Searches the next ready process.
119
120
121
122
123
124
125
126
            ++currentTask;
            if(currentTask ≥ numTasks)
                currentTask %= numTasks;
        }while(tasks[currentTask].state # Task::READY || tasks[currentTask].waiting_child_pid # -1);
        // I added the second check in while loop to pass the processes that are waiting for its child to finish.
        tasks[currentTask].state = Task::RUNNING;
        return tasks[currentTask].cpustate;
```

The schedule function first saves the current cpu state to the cpu state of the currently running task and make currently running task ready. I added some checks to make sure that scheduler won't unintentionally make a blocked or terminated process ready. After this, it searches for the next ready (and not waiting for a child) task in a loop. When it founds that task, it sets that tasks state to running and returns the cpustate of that task like older version.

```
179 uint32_t TaskManager::GetPid(){
        return tasks[currentTask].pid;
181 }
182
183 uint32_t TaskManager::GetParrentPid(){
        return tasks[currentTask].p_pid;
185 }
187 uint32_t TaskManager::WaitPid(uint32_t pid){
        // todo: check this before in the syscall to preve
        if(TaskFromPid(pid) \rightarrow state = Task::TERMINATED)
            return -1;
        tasks[currentTask].waiting_child_pid = pid;
        return pid;
196 void TaskManager::Exit(){
        Task* current = &tasks[currentTask];
        Task* parent = TaskFromPid(current→p_pid);
        if(current \rightarrow p_pid \neq 0){
            Task* parent = TaskFromPid(current→p_pid);
            if(parent = nullptr){
                 printf("Can't find parent in exit.\n");
            if(parent \rightarrow waiting\_child\_pid = current \rightarrow pid)
                 parent→waiting_child_pid = -1;
        current→state = Task::TERMINATED;
212 }
214 Task* TaskManager::TaskFromPid(uint32_t pid){
        for(int i = 0; i < numTasks; ++i){</pre>
             Task* taskP = &tasks[i];
             if(taskP \rightarrow pid = pid)
                 return taskP;
        return nullptr;
221 }
```

The getpid, getParrentPid functions are simply return the pid and p\_pid.

WaitPid function sets the waitingPid of the current task to the pid from the syscall argument that is set in syscallhandler if that task is not terminated already.

The TaskFromPid returns the pointer to the task that its pid is pid in the argument.

Exit function first checks whether exited processes parent is waiting for exited process to finish, if so, it sets the waiting pid of the exited processes parent to -1 (I used -1 as that process is not waiting for a child to finish.) and then it sets the exited processes state as terminated.

#### Fork:

```
221 void TaskManager::Fork(CPUState* cpuStαte){
         if(numTasks \ge 256){
              cpuState \rightarrow eax = -1;
         Task * parent = &tasks[currentTask];
         Task * child = &tasks[numTasks];
         child→stαte = Task::READY;
         child \rightarrow p_pid = parent \rightarrow pid;
         child \rightarrow pid = newPid \leftrightarrow ;
         child→waiting_child_pid = -1;
234
         for (int i = 0; i < sizeof(parent \rightarrow stack); i \leftrightarrow)
              child→stack[i]=parent→stack[i];
240
         vint32_t parentCpuStackOffset = (vint32_t)cpuState - (vint32_t) (parent +> stack);
241
         child \rightarrow cpustate = (CPUState*)(((uint32_t) (child \rightarrow stack)) + parentCpuStackOffset);
243
         uint32_t sp0ffset = (uint32_t)(parent→stack) - (uint32_t) (cpuState);
         child→cpustate→esp = (uint32_t) (child→stack) + sp0ffset;
245
         child \rightarrow cpustate \rightarrow eax = 0;
         cpuState \rightarrow eax = child \rightarrow pid;
         numTasks++;
249 }
```

My fork function is almost as same as the shared code, except I fixed some bugs. Here I add a child process to the end of the tasks array and assign its state, parent pid, pid and waiting\_child\_pid. Then copy the stack. And calculate the cpu state and assign it also.

But after this the esp register of the child is same with the parent (that value is copied when we copy the stack) so we need to assign it to the correct value which is an offset from child stack. I calculated this offset by parent stack – current esp (since the stack pointer is incremented by decreasing it.) And I used the current cpuState as current esp because they are essentially same (in syscallhandler the current esp is set as cpustate).

Second change I did is, I used the return value register (eax) to set the fork return values of the child and parent correctly.

## Helper Additions:

I added a sleep function, it just counts in a nested for loop, I tuned the values for my computer so that argument 1 approximately counts for 1 second but these values may need to be changed if sleep takes longer or shorter time.

I added a simple function to print decimal values.

I added this check to slow down the scheduler.

1 / [SCH\_SLW\_DWN] hardware interrupt calls schedule function, on remaining interrupts, the kernel sleeps for very short amount of time.

In my code, value of SCH\_SLW\_DWN is defined as 99.

## Collatz Sequence function:

```
void collatzSequence() {
        int32_t r = random();
        if(r < 0)
        if(r = 0)
        int n = r \% 100;
        if (n \leq 0) {
            sysPrintf("Input must be a positive integer.\n");
        sysPrintf("Collatz sequence for ");
        sysPrintfDec(n);
sysPrintf(" is: ");
        while (n \neq 1) {
            sysPrintfDec(n);
            sysPrintf(" ");
if (n % 2 = 0) {
            n = n / 2;
} else {
        sysPrintf("1\n");
        sleep(0.5);
        sysExit();
454 }
```

## Long Running Program:

The programs are implemented as in the pdf, Collatz sequence program calculated the collatz sequence of a random number.

```
71 int32_t random(){
72    int32_t time;
73    asm("rdtsc": "=A"(time));
74
75    time *= time;
76    time #= 100;
77    return (int32_t)(time % 10000);
78 }
```

Random function is a very simple function that takes the time from rdtsc and returns it after some random calculations.

```
481 void initialTask(){
        int32_t pids[6];
        for(int32_t i = 0; i < 6; i++) {
             if(i\%2 = 0){
                 pids[i] = sysFork();
                 if(pids[i] = 0){
                     sysPrintf("Forked and executing collatz..\n");
                     sysPrintf("Pid is: ");
                     sysPrintfDec(sysGetPid());
                     sysPrintf("\n");
sysExecve(collatzSequence);
                     sleep(0.5);
                 }
             else{
                 pids[i] = sysFork();
                 if(pids[i] = 0){
                     sysPrintf("Forked and executing long running program..\n");
sysPrintf("Pid is: ");
                     sysPrintfDec(sysGetPid());
                     sysPrintf("\n");
                     sysExecve(longRunningProgram);
                     sleep(0.5);
        for(int32_t i = 0; i < 6; ++i){
             sysWaitPid(pids[i]);
        while(true){
             sysPrintf("Processes have finished. \n");
             sleep(0.5);
518 }
```

I added this initial task that loads each program 3 times. After that, it waits for each of them, and then it prints "Processes have finished" in a while loop. My waitpid implementation takes a single process to wait. So in order to wait all of them, I called all of them in a loop. Since sysWaitPid return immediately if the task that we are trying to wait is done already, this didn't cause any problems.

Tests:

#### Fork and GetPid:

I made a simple program that forks 2 times (second fork is called only in the child). Then all these processes print their process id using sysGetPid call and the return values from fork. The output is as expected:

```
Hello
INTERRUPT FROM AMD am79c973
AMD am79c973 DATA SENT
AMD am79c973 INIT DONE
Schedular is called
I'm the process to test fork, I am going to fork and my child will fork too.
We will be printing our pids and forkReturn values.
syscall fork called
Hello, I'm the parent process, and my pid is: 1
I'm going to print parent from now on.
Parent pid: 1, forkReturnParent: 2
Parent pid: 2, forkReturnParent: 0
Schedular is called
Hello, I'm the child process, and my pid is: 2
I'm going to fork again and print child from now on.
syscall fork called
Child pid: 2, forkReturnParent: 0, forkReturnChild: 3
Child pid: 2, forkReturnParent: 0, forkReturnChild: 3
Child pid: 2, forkReturnParent: 0, forkReturnChild: 3
```

```
Child pid: 2, forkReturnParent: 0, forkReturnChild: 3
Schedular is called
Hello, I'm the grand child process, and my pid is: 3
I'm going to print grand child from now on.
Grand Child pid: 3, forkReturnParent: 0, forkReturnChild: 0
Schedular is called
Parent pid: 1, forkReturnParent: 2
```

```
Child pid: 2, forkReturnParent: 0, forkReturnChild: 3
Schedular is called
Grand Child pid: 3, forkReturnParent: 0, forkReturnChild: 0
Schedular is called
Parent pid: 1, forkReturnParent: 2
```

```
Execve:

In this test, I made a program like:

taskTestExecve:

forkReturn = sysfork()

if forkReturn == 0:

execve(taskPrintExecve)

while(true):

print("original")

taskPrintExecve:

while(true):

print("from execve")

The output is as expected:

Kernel Main is started.
INTERRUPT FROM AMD am79c973
AMD am79c973 DATA SENT
AMD am79c973 INIT DONE
Schedular is called
I'm the process to test execve. I'm going to for ecuting something else.
syscall fork called
```

```
Kernel Main is started.
INTERBUFT FROM AMD am/9c973
AMD am/9c973 DATA SENT
AMD am/9c973 INIT DONE
Schedular is called
I'm the process to test execue. I'm going to fork and my child is going to be executing something else.
syscall fork called
Original
Schedular is called
From execue
```

# My output:

With process table (with printing the process table, I can't show some parts of the output because even when I slowed the schedular and video record the output, when the printed task is big, it skips some of the output.):

```
:1, waiting pid: -1, state:1,

Porked and executing long running program..

Pid is: 7

syscall execue called

Process id: 7

result: 1884769809

Schedular is called

task:0: pid:1, p_pid:0, waiting pid: -1, state:1, task:1: pid:2, p_pid:1, waiting pid: +1, state:3, task:2: pid:3, p_pid:1, waiting pid: -1, state:3, task:3: pid:4, p_pid:1, waiting pid: -1, state:3, task:3: pid:6, p_pid:1, waiting pid: -1, state:3, task:6: pid:7, p_pid:1, waiting pid: -1, state:3, task:3: pid:4, p_pid:1, waiting pid: -1, state:3, task:6: pid:7, p_pid:1, waiting pid: -1, state:3, task:6: pid:1, p_pid:1, waiting pid: -1, state:3, ta
```

Without printing the process table (I put this output to show the results more clearly):

```
INTERRUPT FROM AMD am79c973
AMD am79c973 DATA SENT
AMD am79c973 INIT DONE
Schedular is called
syscall fork called
Schedular is called
Forked and executing collatz..
Pid is: 2
syscall execve called
Collatz sequence for 69 is: 69 208 104 52 26 13 40 20 10 5 16 8 4 2 1
Schedular is called
Forked and executing long running program..
Fid is: 3
syscall execve called
Process id: 3
```

```
result: 1884769809
Schedular is called

Forked and executing collatz..
Pid is: 4
syscall execve called
Collatz sequence for 56 is: 56 28 14 7 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2
1
Schedular is called

Forked and executing long running program..
Pid is: 5
syscall execve called
Process id: 5
Schedular is called

Forked and executing collatz..
Pid is: 6
syscall execve called
Collatz sequence for 69 is: 69 208 104 52 26 13 40 20 10 5 16 8 4 2 1
```

```
Process id: 7
result: 1884769809
Schedular is called
syscall waitpid called
result: 1884769809
Schedular is called
syscall waitpid called
syscall waitpid called
Schedular is called
syscall waitpid called
Schedular is called
Processes have finished.
```