

# **Bora Sahin, 11/01/2015**

## **Backgammon Client / Server Application Design**

### **Abstract**

The backgammon application design document for client and server was written as part of the term project effort, in partial fulfillment of SWE544 Internet Programming course requirements at Bogazici University.

It explains client-server communication, client and server architecture and design.

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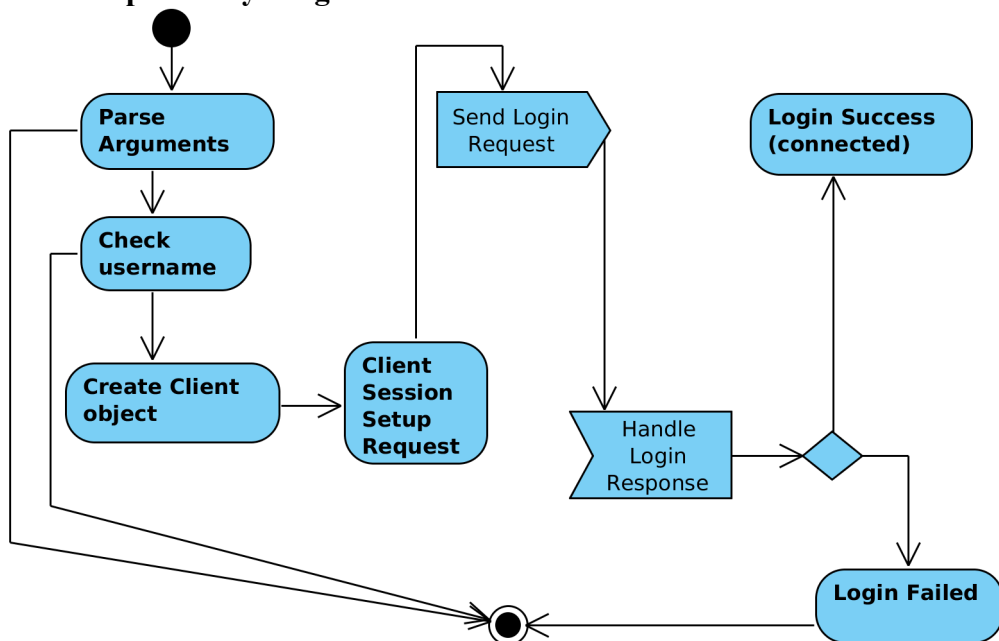
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## 1. Client / Server Communication

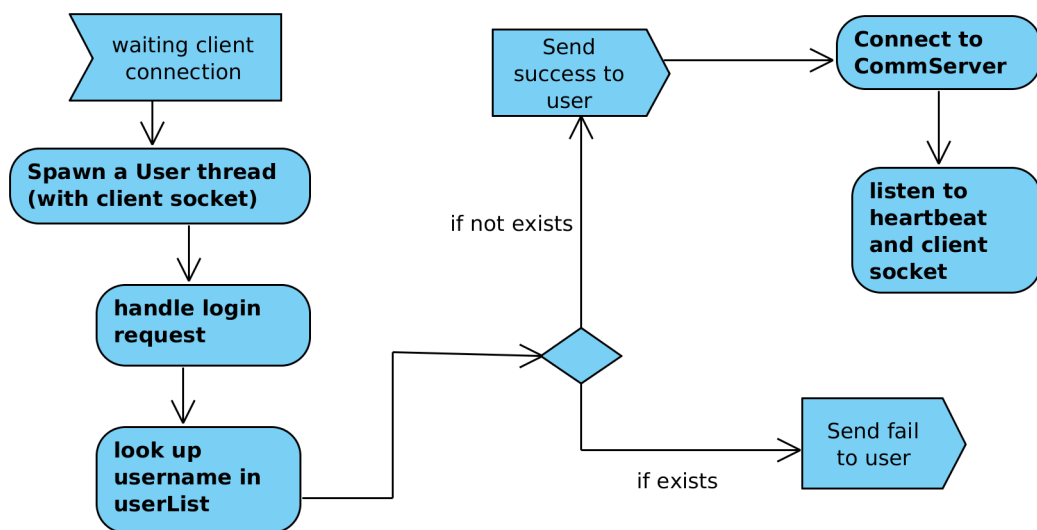
### 1.1. Session Setup

Client Messages	Server Messages
CLOGIN(username) ->	
	Server receives CLOGIN(username) and checks if that username is used to log in to the system
	<- SLRSPS(success)
	<- SLRSPS(fail)

#### Client Session Setup Activity Diagram



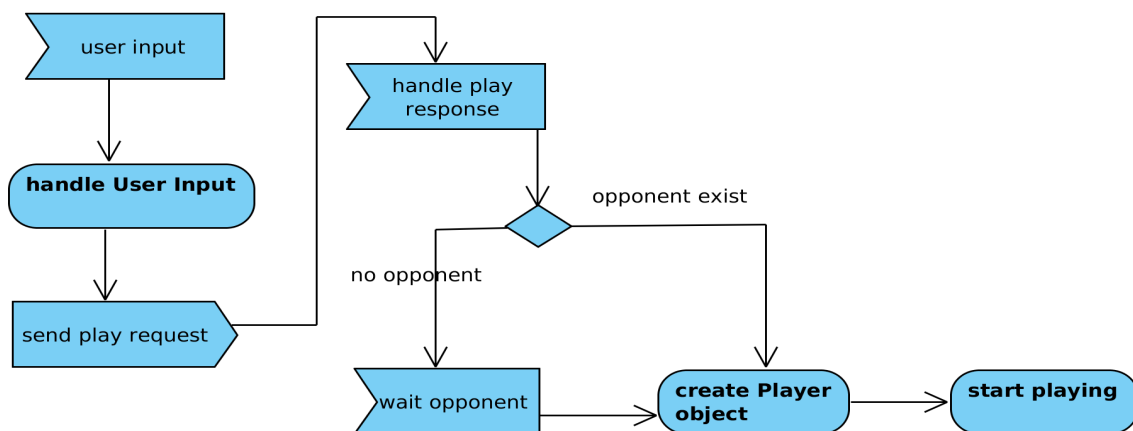
#### Server Session Setup Activity Diagram



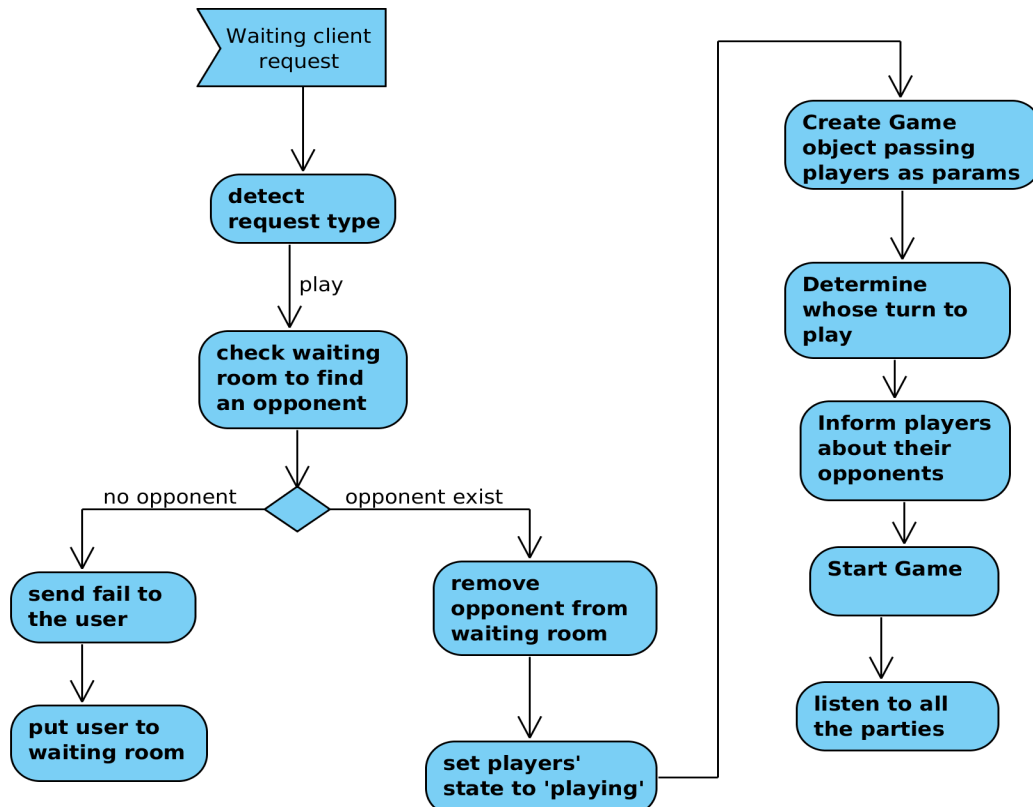
## 1.2. Play Request

Client Messages	Server Messages
CREQST(play)	
	Server receives CREQST(play) and checks if there is someone in the waiting room wanting to play
	<- SREQRP(play, success)
	<- SREQRP(play, fail)

### Client Play Request Activity Diagram



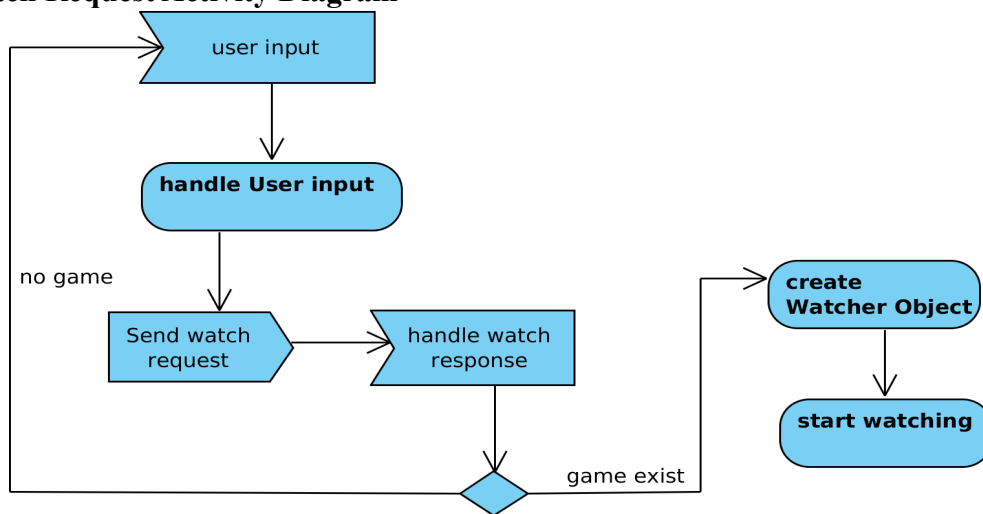
### Server Play Request Activity Diagram



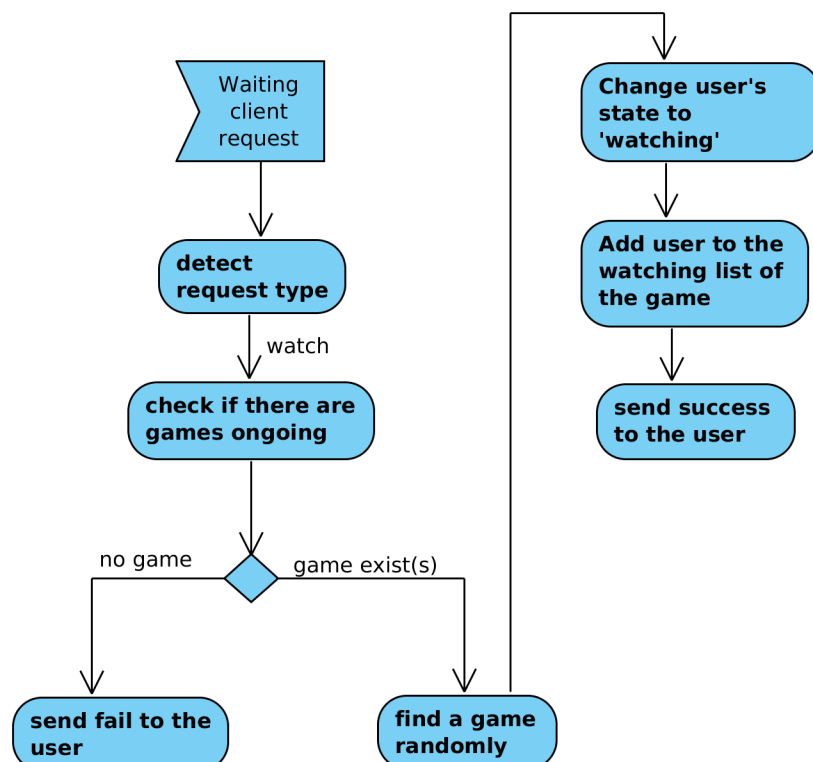
### 1.3. Watch Request

Client Messages	Server Messages
CREQST(watch)	
	Server receives CREQST(watch) and checks if there is a match or are matches. If there is, then attaches watcher to a game randomly
	<- SREQRP(watch, success)
	<- SREQRP(watch, fail)

#### Client Watch Request Activity Diagram



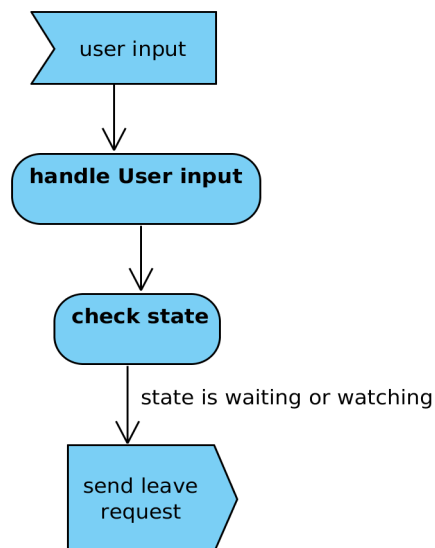
#### Server Watch Request Activity Diagram



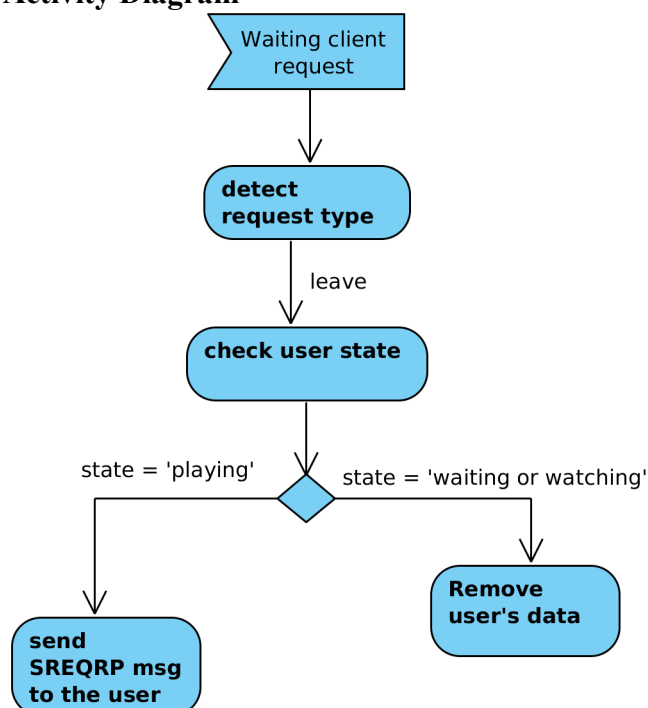
## 1.4. Leave Request

Client Messages	Server Messages
CREQST(leave)	
	Server receives CREQST(leave) and checks the state of the user. If it is either 'waiting' or 'watching' then remove its data and does nothing. If its state is 'playing', then SVRNOK is sent to the client
	nothing
	<- SVRNOK()

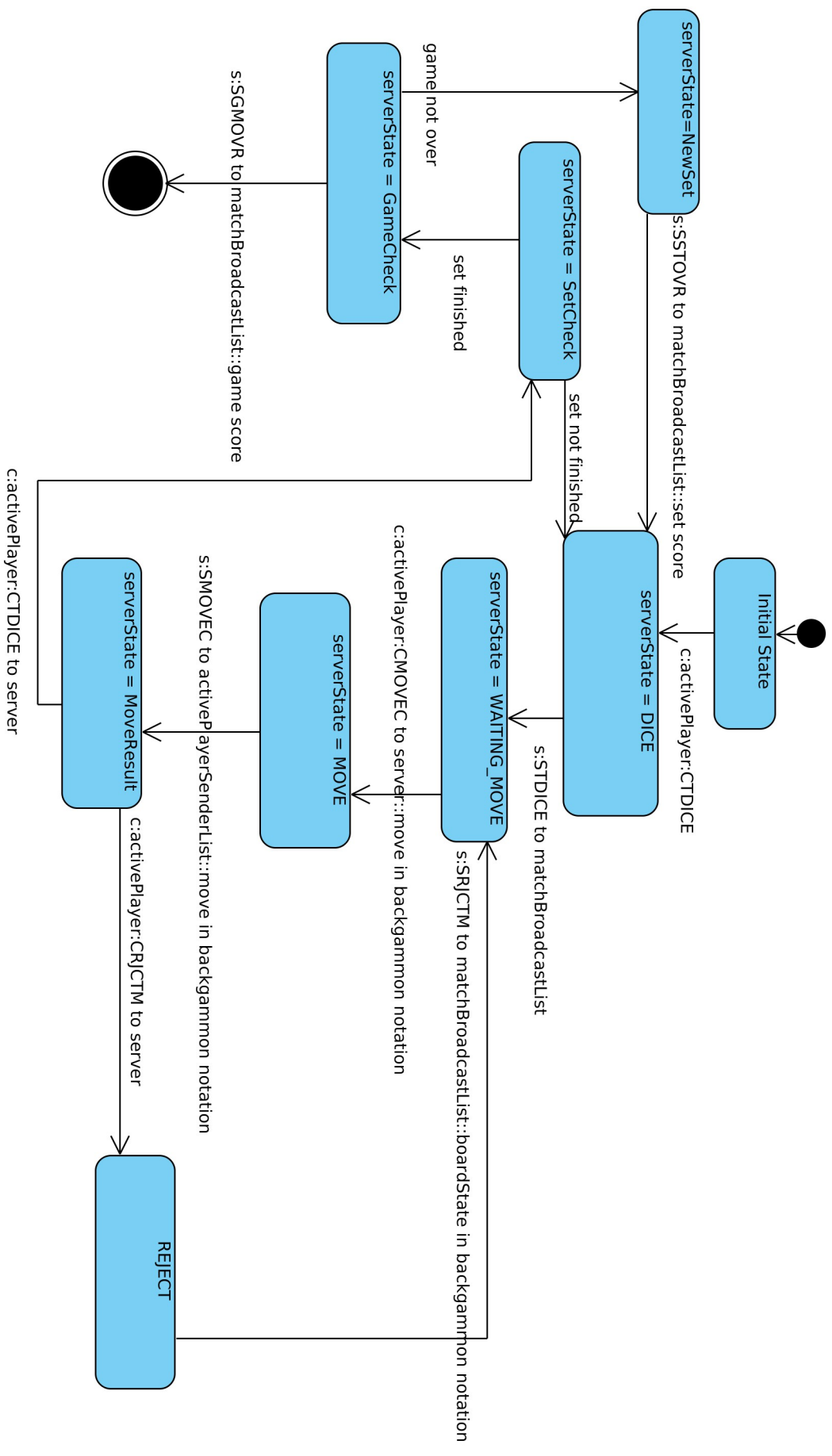
### Client Leave Request Activity Diagram



### Server Leave Request Activity Diagram



## 2. Playing State Machine and Messaging



Above state machine diagram represents the situation in a game occurring in the server, while also showing messages exchanged by server and players.

When two players who want to play backgammon are in the server, a game is started. The first thing which is done by the server is to determine whose turn to play and color for each player. Therefore, server throws single dice on behalf of players and informs them about the result (SREQRP). They are shown in previous sections. Before commenting furthermore, a few things need to be cleared:

- The player who wins to the challenge has the right to request throw dice from the server and is marked as active player
- All the moves are represented in backgammon notation
- sockList is a dictionary and any participant be it either player or watcher is added to it. As values of dictionary, socketType and userObject is hold.
- Passive player is the one who is waiting his/her opponent's move
- c: refers to a message sent by the client
- s: refers to a message sent by the server

Active player starts to the game by requesting throw dice from the server, which is represented by CTDICE message (c:ActivePlayer:CTDICE). Upon receiving this message, server enters into 'DICE' state and throws dice, which are broadcasted to every socket in sockList (s:STDICE) and brings the server state to 'WAITING\_MOVE'. Active player which is said to have the right to make a move decides what to play and tells it to server (c:ActivePlayer:CMOVEC). When server receives CMOVEC from a player, rebrands it as SMOVEC and then sends to activePlayerSenderList (s:SMOVEC). At this point activePlayer is changed to other player. Opponent gets SMOVEC and may have two reservation: Accept or Reject. If opponent accepts the move, it sends CTDICE request to the server (c:CTDICE); if not, then sends CRJCTM (c:CRJCTM). It can be described as below:

#### Scenario 1

Player1	Server	Player2
CTDICE ->	Player1:active, Player2:passive	
	Player1:active, Player2:passive <- STDICE ->	
Receives STDICE		Receives STDICE
CMOVEC ->		
	Receives CMOVEC and rebrands it as SMOVEC Player1:passive, Player2:active SMOVEC ->	
		Receives SMOVEC, then decides its fate
		ACCEPT
		<- CTDICE
	<- STDICE ->	

## Scenario 2

Player1	Server	Player2
CTDICE ->	Player1:active, Player2:passive	
	Player1:active, Player2:passive <- STDICE ->	
Receives STDICE		Receives STDICE
CMOVEC ->		
	Receives CMOVEC and rebrands it as SMOVEC Player1:passive, Player2:active SMOVEC ->	
		Receives SMOVEC, then decides its fate
		REJECT
		<- CRJCTM
	Server receives CRJCTM and rebrands it as SRJCTM and sends it to matchBroadcastList Player1:active, Player2:passive <- SRJCTM ->	
Receives SRJCTM		Receives SRJCTM
Needs to make another move		

After a move is accepted by an opponent, it is told to the server by sending CTDICE message. As server keeps some backgammon logic, namely to tell if a game is won or not, it should check it. A game consists of sets and a game finishes if one player reaches 5 set point. When a set is won, server informs all the parties with a message called SSTOVR and then goes on to check if the game is finished or not, and if finished, then informs all the parties with a message called SGMVOR. At this point, sessions of all the parties are closed. SSTOVR is just an informational message. Its only impact is that the player who won the set has the right to request throw dice from the server.

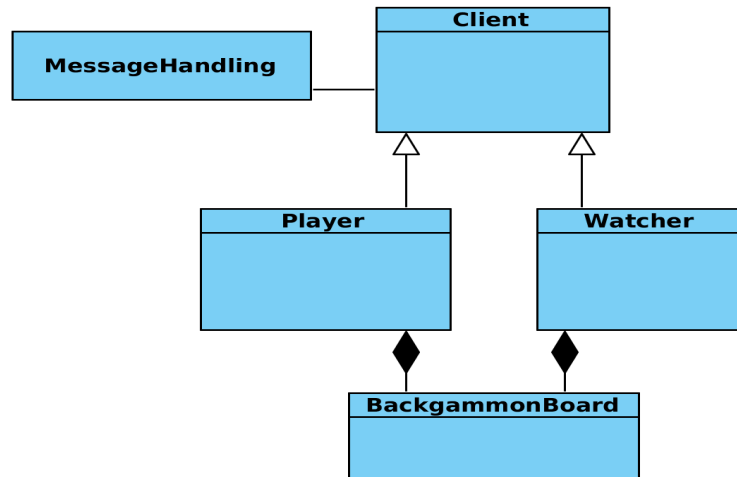
If a player loses his/her connection to the server, which means two back-to-back PING messages missed by the client, then server informs all the parties that a player lost his/her connection by sending STEARD message. All the data related to users and game are removed from the server and it is expected that clients receiving this message should also act upon appropriately.



### 3. Client Architecture and Design

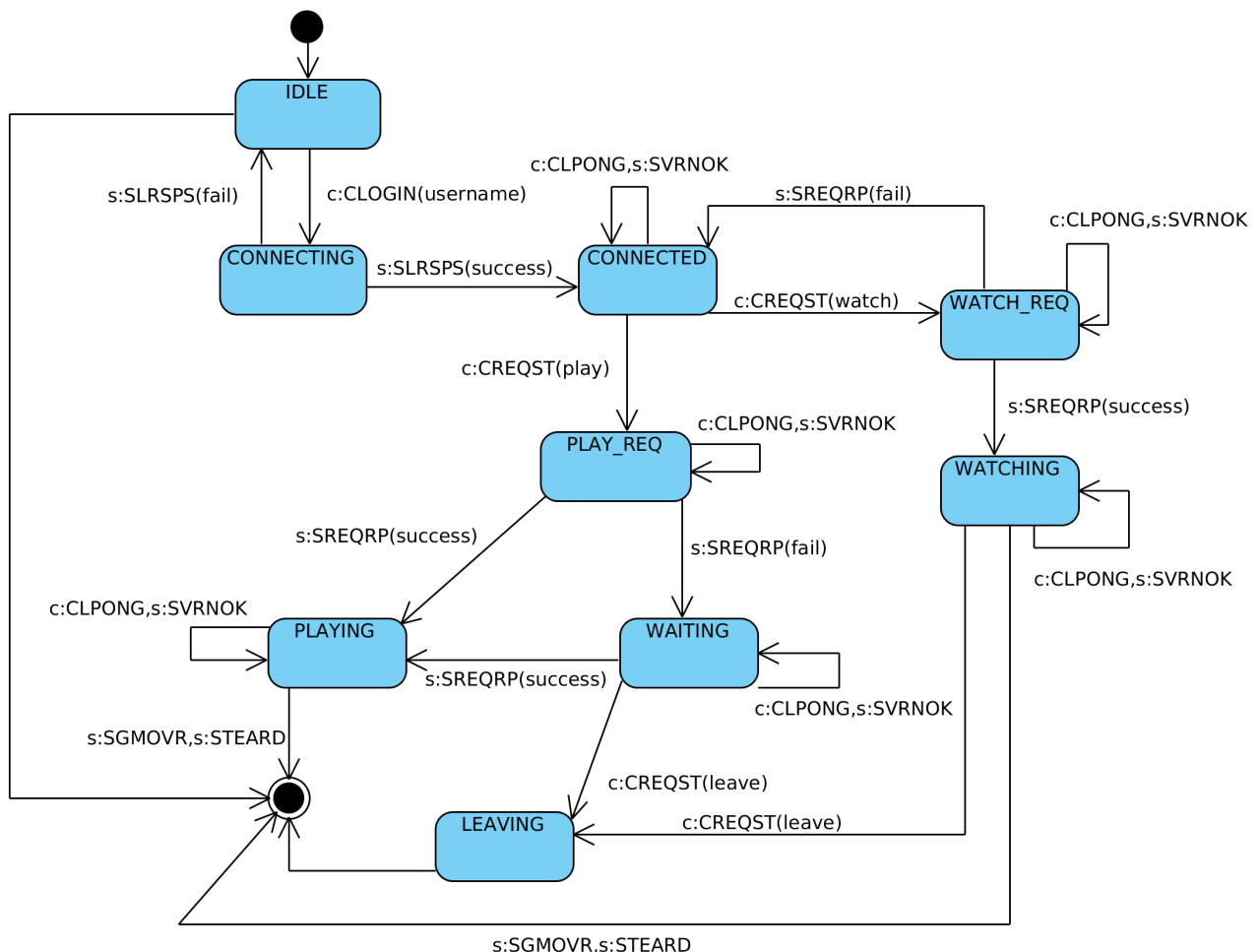
#### 3.1. Client Object Model

Client class represents a client be it either a player or a watcher.



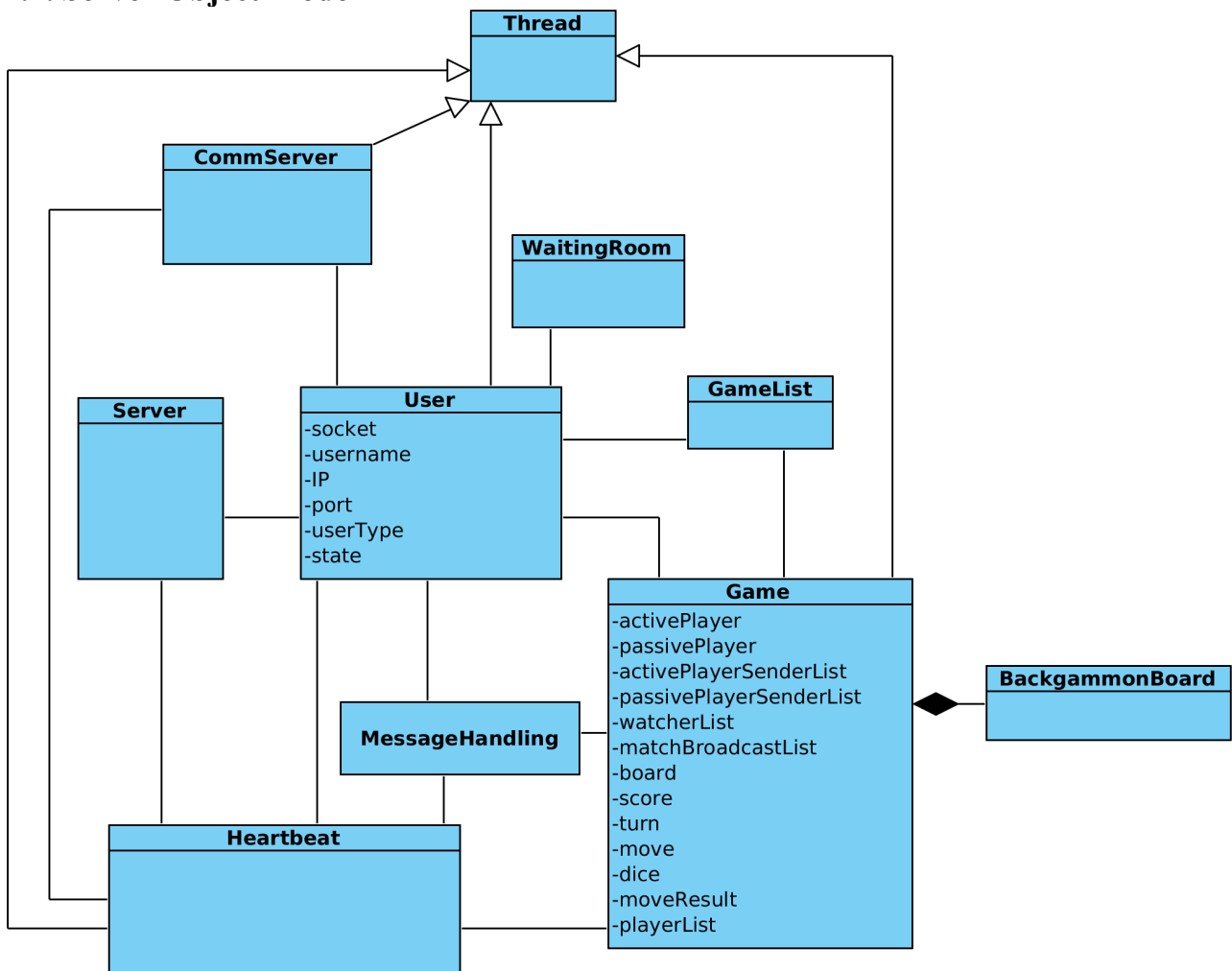
MessageHandling and BackgammonBoard are used both by client and server. Client creates a backgammon board only its request as either play or watch is accepted by the Server. Client has two endpoints to listen to after its login request is satisfied by the Server: (1) messages sent by the Server (2) input entered by the User. Python select.poll() or something similar will be used for this purpose.

#### 2.2. Client State Machine



## 4. Server Architecture and Design

### 4.1. Server Object Model



**Server** is the main thread. Its sole purpose is to receive a connection request from a client, create a new **User** object (a new thread) and give its handling to it.

**User** handles 'play' and 'watch' requests in **connected** state and 'leave' request in **waiting** state. A unix domain socket is used between **Heartbeat** and **User** object. **User** connects to **CommServer** before entering into connected state. **CommServer** sends the client socket to **Heartbeat**. **User** listens to both client socket and unix domain socket it obtained from **Comm** server by using `select.poll()`. If CLPONG message is received from the client, it calls the ping method of **Heartbeat** object. If play request is satisfied, a **Game** object is created, all the information were passed to **Game** and **User** object wait for an event by **Game** object. Python Event will be used for this purpose.

**Heartbeat** object runs in its own thread and started by **Server**. Every 30 seconds (heartbeat time), it sends all the clients SVPING message, which are registered to it. However, it does not collect response messages, instead **User** or **Game** thread collects the responses and informs **Heartbeat**. If **Heartbeat** object notices that a client is not responsive for more than 2 back-to-back PING, then it notifies **User** or **Game** thread, which are listening all the ports (a user connected to the server have 2 sockets: client socket and unix domain socket) by sending a message through the unix domain socket.

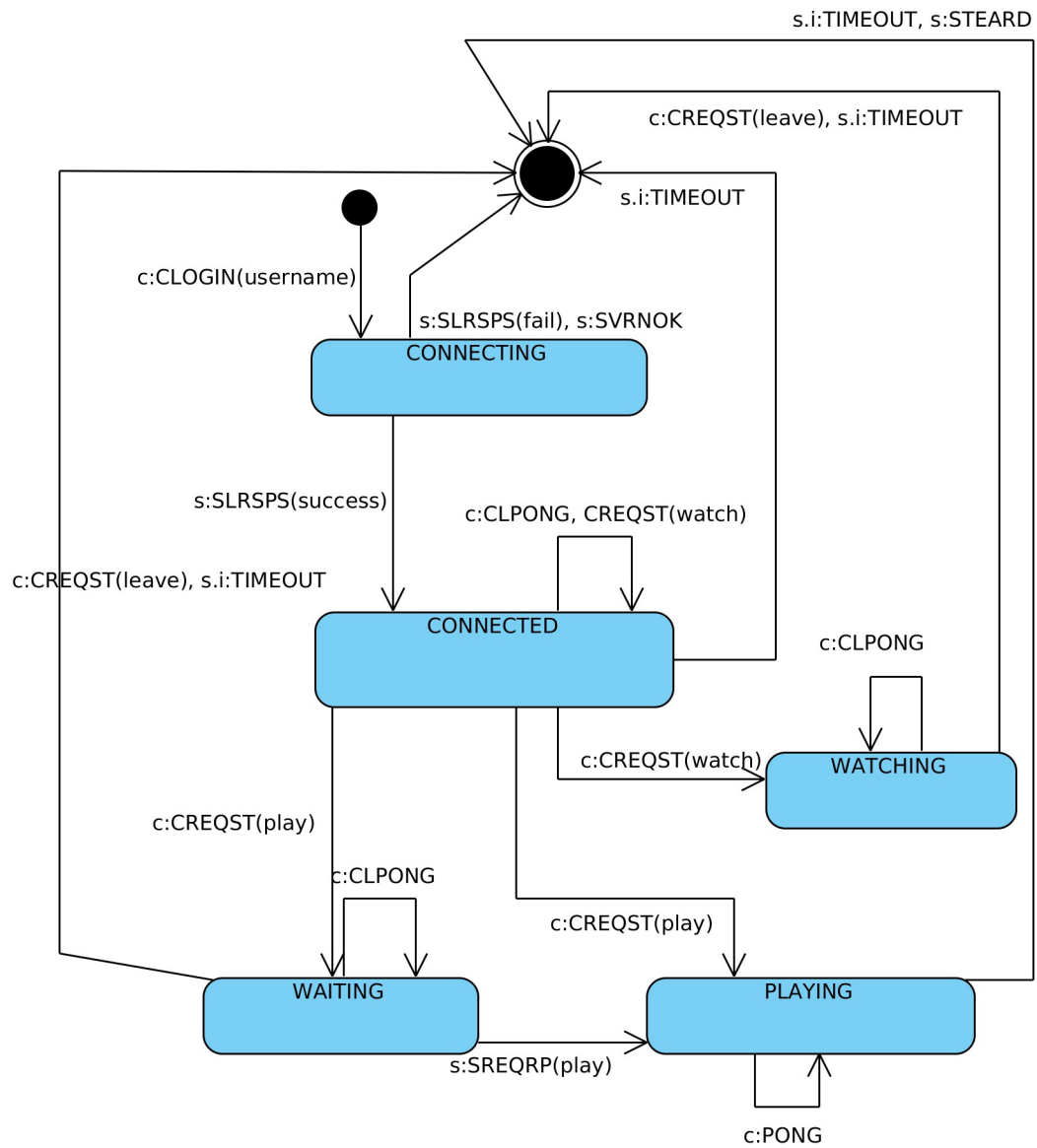
**CommServer** is used to obtain a unix domain socket. Every user whose login is successful connects to CommServer to have a unix domain pair, which may going to be used by Heartbeat object to inform the user that his/her counterpart in the client side is dead.

**Game** also runs in its own thread. If a user wants to play a game and there is a user in the waiting room, then User object creates a Game object and passes all the information to it. Game object listens to player sockets (client socket + unix domain socket for Heartbeat) and watcher sockets. Actually "Playing State Machine" in the server side is controlled by Game object. User object handles most port of "Server State Machine"

**WaitingRoom** object is used to host clients who want to play a game but couldn't find a opponent. If an opponent comes to the server and wants to play a game consults to waitingRoom and if it finds someone, the removes it from the waitingRoom and starts a game.

**GameList** object is used to hold all the games currently on-going in the server. When a watcher comes, it consults gameList to ask for if there is/are game(s). If there is one, a game is randomly chosen, otherwise a fail message is returned to the client.

## 4.2. Server State Machine



## **References**

1. SWE544\_RFC\_1 document
2. SWE544\_RFC\_2 document
3. SWE544 lectures
4. Visual Paradigm, [www.visual-paradigm.com](http://www.visual-paradigm.com)