

Functional Programming

Type Classes — Overloading in Haskell

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Overloading

Remember previous classes?

We were able to use

- equality `==` and ordering `<` with many different types
- arithmetic operations with many different types

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Parametric polymorphism

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Haskell integrates overloading with polymorphism

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- Types contain type variables and **constraints**

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Examples

```
1  -- elem x xs : is x an element of list xs?  
2  -- type a must have equality  
3  elem :: Eq a => a -> [a] -> Bool  
4  -- insert x xs : insert x into sorted list xs  
5  -- type a must have comparison  
6  insert :: Ord a => a -> [a] -> [a]  
7  -- square x : compute the square of x  
8  -- type a has numeric operations  
9  square :: Num a => a -> a
```

Type classes

- Each constraint mentions a **type class** like `Eq`, `Ord`, `Num`, ...
- A type class specifies a set of operations for a type e.g. `Eq` requires `==` and `/=`
- Type classes form a hierarchy e.g. `Ord a => Eq a`
- Many classes are predefined, but you can roll your own

Classes and Instances

- A *class declaration* specifies a signature (i.e., the class members and their types)

```
1 class Num a where  
2   (+), (*), (-) :: a -> a -> a  
3   negate, abs, signum :: a -> a  
4   fromInteger :: Integer -> a
```

- An *instance declaration* specifies that a type belongs to a class by giving definitions for all class members

```
1 instance Num Int where ...  
2 instance Num Integer where ...  
3 instance Num Double where ...  
4 instance Num Float where ...
```

- This info can be obtained from GHCi by

```
1 :i Num
```

Equality

The type class Eq

```
1 class Eq a where  
2   (==), (/=) :: a -> a -> Bool  
3   x /= y = not (x == y) -- default definition
```

An instance must only provide (**==**).

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Instances of Eq

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1 instance Eq Int -- Defined in 'GHC.Classes'  
2 instance Eq Float -- Defined in 'GHC.Classes'  
3 instance Eq Double -- Defined in 'GHC.Classes'  
4 instance Eq Char -- Defined in 'GHC.Classes'  
5 instance Eq Bool -- Defined in 'GHC.Classes'  
6 { -- and many more -- }
```

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Tacit assumption

Equality is a congruence relation.

Defining Eq for pairs

When are two pairs equal?

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Solution

```
1 instance (Eq a, Eq b) => Eq (a, b) where
2   (a1, b1) == (a2, b2) = a1 == a2 && b1 == b2
```

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When are two pairs equal?

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YES: on types, NO: on values

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Solution

```
1 instance Eq a => Eq [a] where
2   [] == [] = True
3   (x:xs) == (y:ys) = x == y && xs == ys
4   _ == _ = False
```

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```

Is this definition recursive?

YES: no types, YES: on values

The equality `xs == ys`.

Handwriting vs deriving an instance

Remember the Hearts game

```
1 data Color = Black | Red  
2   deriving (Show)
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Define your own equality

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1 instance Eq Color where  
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```

Same result as deriving Eq

```
1 data Color = Black | Red  
2   deriving (Show, Eq)
```

Further useful classes

Show and Read

```
1 class Show a where
2   show :: a -> String
3   {- ... -}
4
5 class Read a where
6   read :: String -> a
7   {- ... -}
```

- Predefined for most built-in types
- Derivable for most datatype definitions

The Ord class (derivable)

```
1 class Eq a => Ord a where  
2   compare :: a -> a -> Ordering  
3   (<) :: a -> a -> Bool  
4   (<=) :: a -> a -> Bool  
5   (>) :: a -> a -> Bool  
6   (>=) :: a -> a -> Bool  
7   max :: a -> a -> a  
8   min :: a -> a -> a  
9  
10 data Ordering = LT | EQ | GT  -- Defined in 'GHC.Types'
```

More classes for you to investigate

- Enum (derivable)
 - Bounded (derivable)
-

Ambiguity

Some combinations of overloaded functions can lead to ambiguity

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1 f x = read (show x)
2 g x = show (read x)
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1 f x = read (show x)
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Solution

```
1 f :: (Read a, Show b) => b -> a
2 g :: String -> String
```

Further pitfalls / features

- Definitions without arguments and without type signatures are not overloaded (monomorphism restriction)
- Numeric literals are overloaded at type `Num a => a`
- Haskell has a **defaulting** mechanism that resolves violations of the monomorphism restriction
- Caveat: GHCi behaves differently than code in a file

Type classes

- provide a signature for an abstract data type
- instances provide implementations at unrelated types
- many classes are predefined and derivable
- pervasively used in Haskell / some pitfalls