

# P r e s e n t a t i o n   o n

# **GALOIS FIELD**

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# INTRODUCTION

In mathematics, *Galois theory*, originally introduced by **Évariste Galois**, provides a connection between field theory and group theory.



**Évariste Galois (1811 - 1832)**

Pronunciation: eh-vah-reest gah-lwah

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$



$$V = \frac{4}{3} \pi r^3$$

# ÉVARISTE GALOIS LIFE

Not a brilliant mathematician, but he was also a political activist.

In 1830, the **July Revolution (7 August 1830)** overthrew King Charles X, replacing him with **Louis-Philippe**, the so-called "Citizen King."

At first, Galois believed this change **would bring true democracy** to France, but he soon realized that the new monarchy still **avored the wealthy and powerful**.

- Galois became a **committed Republican**, advocating for a people's government.
- He joined the revolutionary organization ***Société des Amis du Peuple*** (***Society of the Friends of the People***).
- He was **arrested several times** for his outspoken criticism of the monarchy.
- Galois even **missed one of his mathematics exams** to take part in a political protest!
- **In 1831, he was imprisoned** for allegedly plotting against the government and possessing weapons.

After his release from prison, Galois remained deeply involved in revolutionary circles.

**In May 1832, he was killed in a mysterious duel at just 20 years old.**



Revolution in France: an eyewitness account - archive, 1830 - theguardian.com  
<https://www.theguardian.com/world/2020/aug/07/revolution-in-france-eyewitness-account-1830>  
[https://en.wikipedia.org/wiki/July\\_Revolution](https://en.wikipedia.org/wiki/July_Revolution)

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$



$$V = \frac{4}{3} \pi r^3$$

Group  $\rightarrow$  Ring  $\rightarrow$  Field  $\rightarrow$  Galois field

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

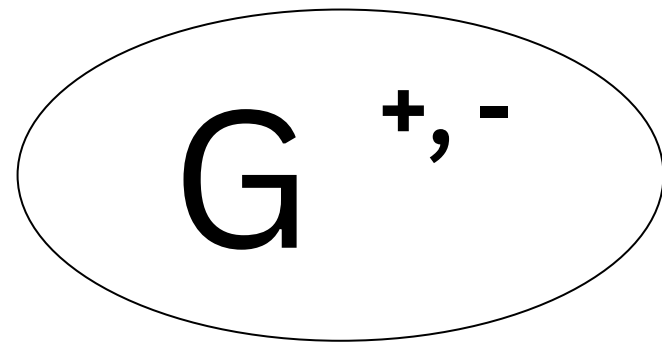
$$y = mx + b$$



$$V = \frac{4}{3} \pi r^3$$

# 1. GROUP THEORY

A group is a set of elements  $G$  together with an operation  $\circ$  which combine two elements of  $G$ .



- The group operation  $\circ$  is **closed**. That is, for all  $a, b \in G$ , it holds that  $a \circ b = c \in G$ .
- The group operation is **associative**. That is,  $a \circ (b \circ c) = (a \circ b) \circ c$  for all  $a, b, c \in G$ .
- There is an element  $1 \in G$ , called the **neutral element** (or identity element), such that  $a \circ 1 = 1 \circ a = a$  for all  $a \in G$ .
- For each  $a \in G$  there exists an element  $a^{-1} \in G$ , called the **inverse** of  $a$ , such that  $a \circ a^{-1} = a^{-1} \circ a = 1$ .
- A group  $G$  is **abelian (or commutative)** if, furthermore,  $a \circ b = b \circ a$  for all  $a, b \in G$ .

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$

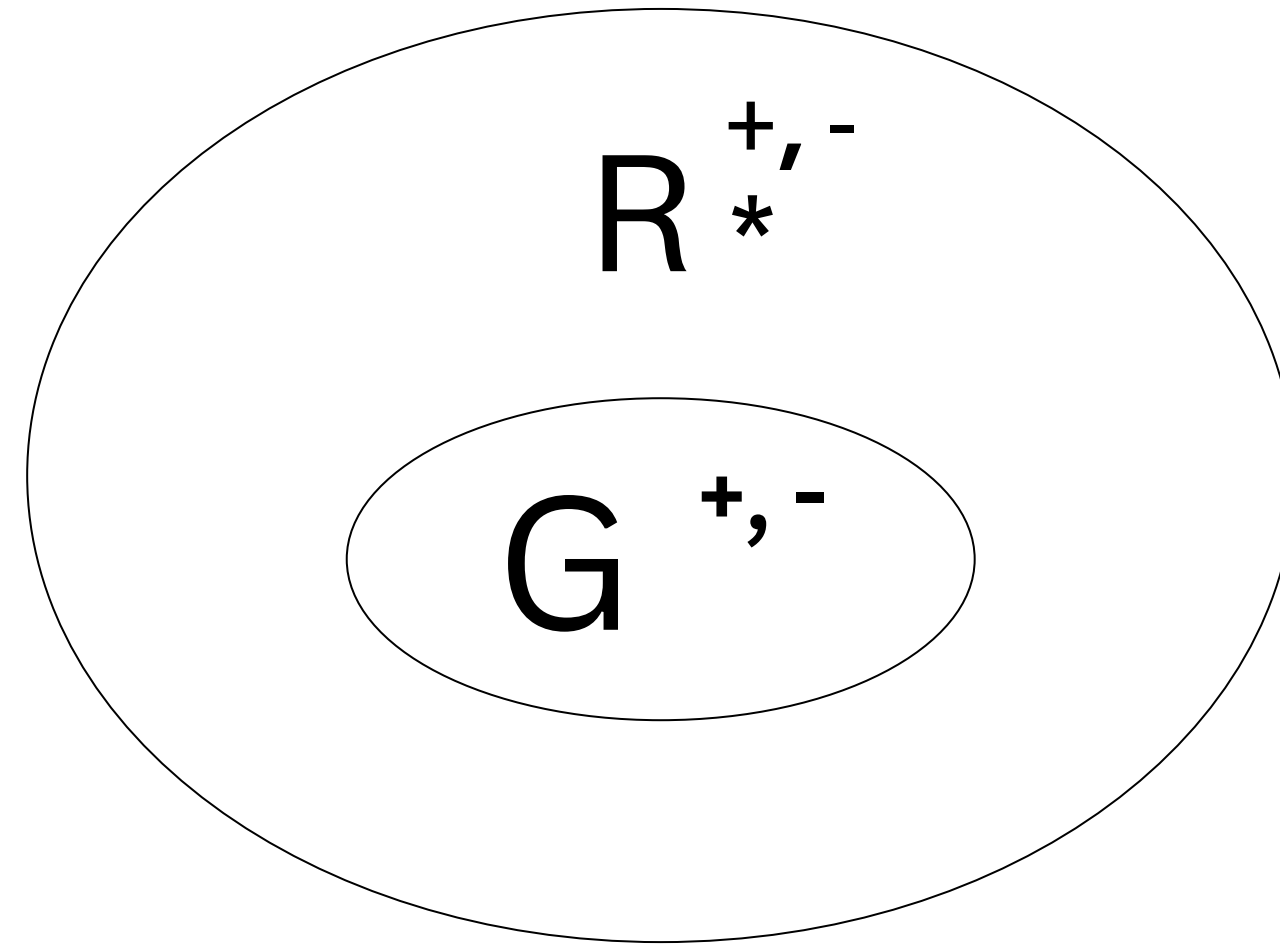


$$V = \frac{4}{3} \pi r^3$$

# 2.RING

A ring is a set that has two operations, usually written as:

- addition (+)
- multiplication ( $\cdot$ )



$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$

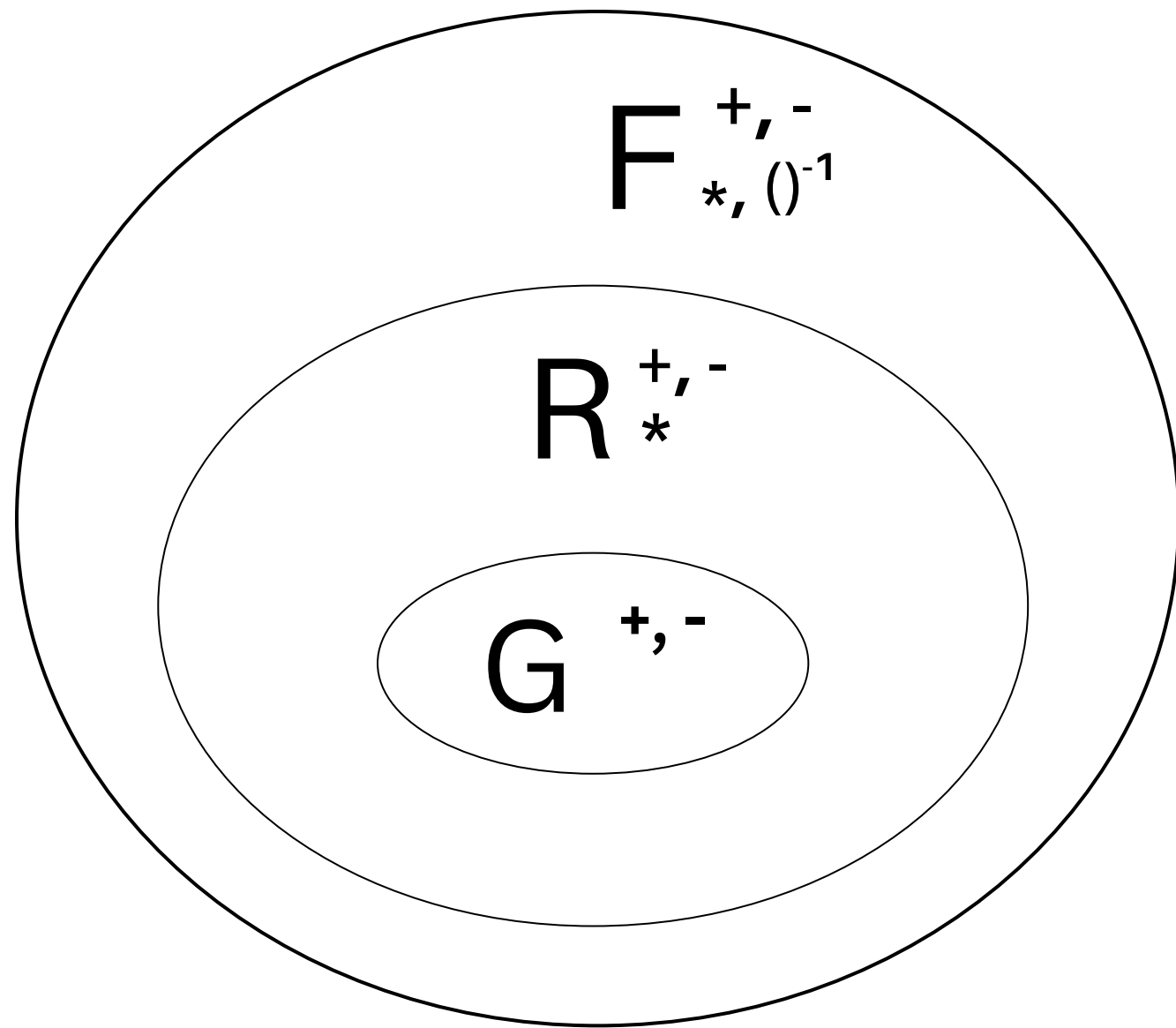


$$V = \frac{4}{3} \pi r^3$$



# 3.FIELD

In order to have all four basic arithmetic operations (i.e.. addition, subtraction, multiplication, division) in one structure, we need a set which contains an additive and a multiplicative group. This is what we call a field.



A field is a set of numbers in which we can add, subtract, multiply and divide.

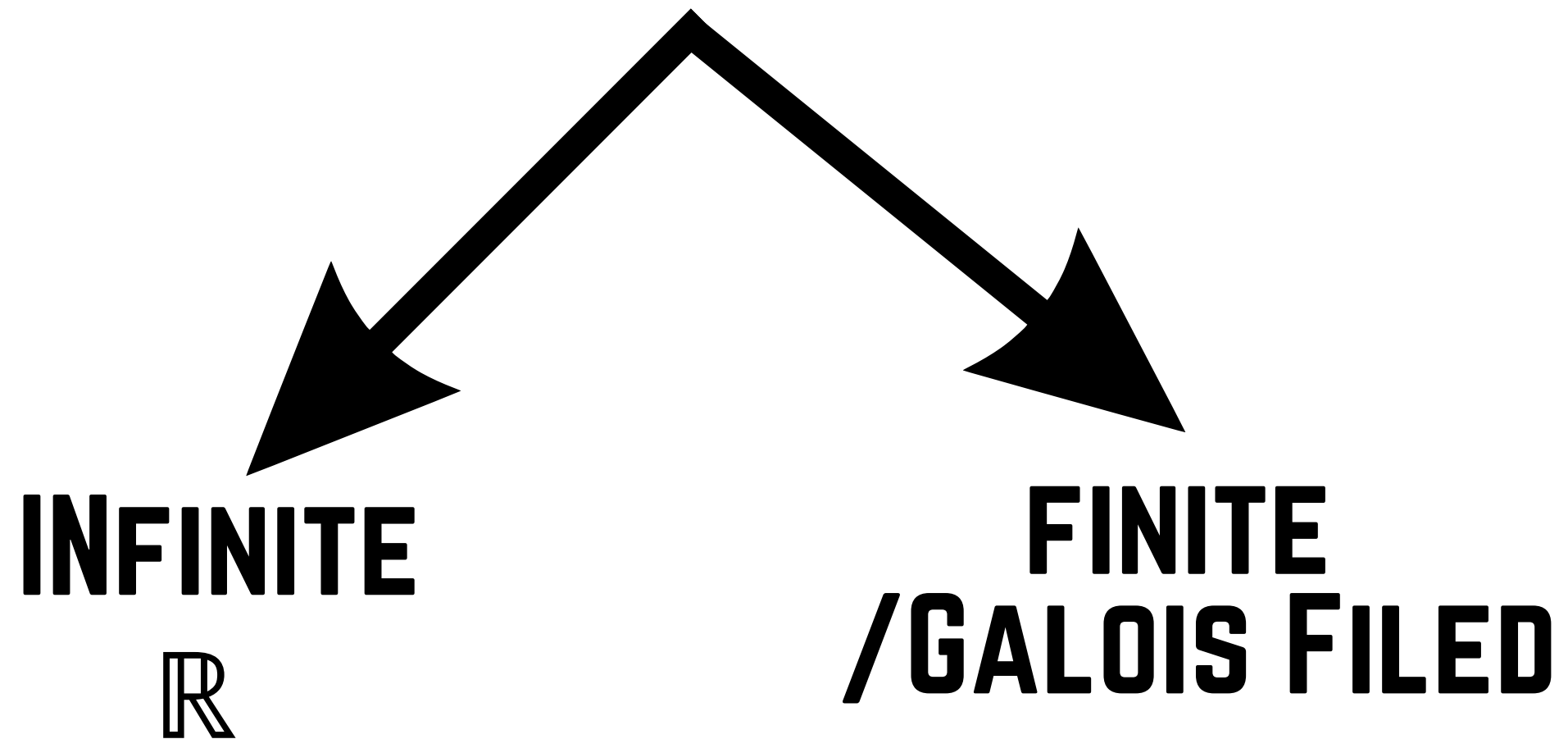
$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

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# FIELD TYPES



$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$



$$V = \frac{4}{3} \pi r^3$$



# FINITE FIELD

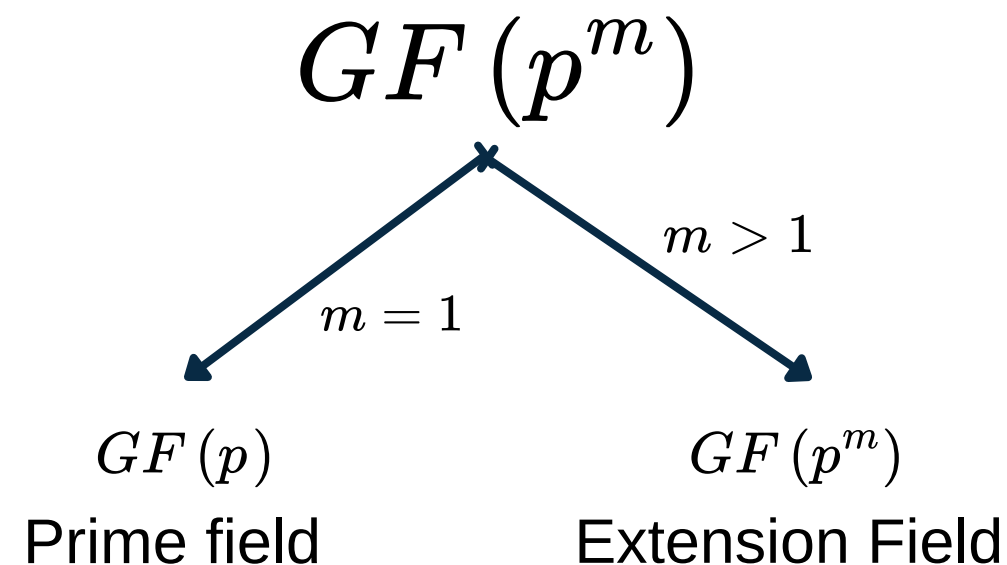
A finite field is a field with a finite field order also called a *Galois field*.

## The Existence Theorem

A finite field of order  $q$  exists if and only if  $q = p^n$  for some prime number  $p$  and positive integer  $n$ .

### Example:

- (1) There is a Galois field with 13 elements :  $GF(13^1)$
- (2) There is a finite field with 81 elements :  $GF(3^4)$
- (3) There is a finite field with 256 elements :  $GF(2^8)$
- (4) There is not a finite field with 12 elements :  $GF(2^2 \times 3) \neq p^n$



$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

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# PRIME FIELD

The elements of prime field of  $GF(p)$  are the

$$GF(p) = \{0, 1, \dots, P - 1\}$$

Example:

$$GF(2) = \{0, 1\}$$

$$GF(7) = \{0, 1, 2, 3, 4, 5, 6\}$$

**1. Addition & Subtraction**

$$GF(7) = \{0, 1, 2, 3, 4, 5, 6\}$$

$$c = 5 + 6 = 11 \equiv 4 \pmod{7}$$

**2. Multiplication**

$$c = 5 \times 6 = 30 \equiv 2 \pmod{7}$$

**3. Inversion**

$$a \in GF(p)$$

The inverse  $a^{-1}$  must satisfy

$$a \cdot a^{-1} \equiv 1 \pmod{p}$$

(Using Extended Euclidian Algorithm )

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$



$$V = \frac{4}{3} \pi r^3$$

# EXTENSION FIELD $(2^m)$

The element of  $GF(2^m)$  are polynomials

$$a_{m-1}x^{m-1} + a_{m-2}x^{m-2} + \dots + a_1x + a_0 = A(x) \in GF(2^m)$$

$$a_i \in GF(2) = \{0, 1\}$$

Example:

$$GF(2^3), m = 3$$

$$A(x) = a_2x^2 + a_1x + a_0$$

$$GF(2^3) = \{0, 1, x, x + 1, x^2, x^2 + 1, x^2 + x, x^2 + x + 1\}$$



a2	a1	a0
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$



$$V = \frac{4}{3} \pi r^3$$

# EXTENSION FIELD ( $2^m$ )

## 1. Addition and Subtraction

$$A(x), B(x) \in GF(2^m)$$

$$C(x) = A(x) + B(x) = \sum_{i=0}^{m-1} c_i x^i \quad c_i = (a_i + b_i) \bmod 2$$

$$C(x) = A(x) - B(x) = \sum_{i=0}^{m-1} c_i x^i \quad c_i = a_i - b_i = (a_i + b_i) \bmod 2$$

Addition  $\bmod 2 =$  Subtraction  $\bmod 2$

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$


$$V = \frac{4}{3} \pi r^3$$

# EXTENSION FIELD ( $2^m$ )

## Example

$$A(x) = x^2 + x + 1, B(x) = x^2 + 1$$

$$C(x) = A(x) + B(x) = (2x^2 + x + 2) \bmod 2 = x$$

## 2. Multiplication

$$C(x) = A(x) \times B(x) \bmod 2 = C'(x)$$

Reduce  $C'(x)$  modulo a polynomial that "behaves like a prime".

Which is called irreducible polynomial  $P(x)$

$$P(x) = \sum_{i=0}^m P_i x^i, P_i \in GF(2)$$

$$C(x) = A(x) \cdot B(x) \bmod P(x)$$

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$



$$V = \frac{4}{3} \pi r^3$$

# EXTENSION FIELD ( $2^m$ )

$$A(x) = x^2 + x + 1, B(x) = x + 1 \in GF(2^3)$$

$$C(x) = A(x) \cdot B(x) \text{ modulo } 2 = x^4 + x^3 + x + 1 = C'(x)$$

$$P(x) = \sum_{i=0}^m P_i x^i = P_0 x^0 + P_1 x^1 + P_2 x^2 + P_3 x^3, P_i \in GF(2)$$
$$= x^3 + x + 1$$

Divide  $C'(x)$  by  $P(x)$  and find the remainder

$$C(x) = x^2 + x = A(x) \cdot B(x) \text{ mod } P(x)$$

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$


$$V = \frac{4}{3} \pi r^3$$

# EXTENSION FIELD ( $2^m$ )

$$A(x) = x^2 + x + 1, B(x) = x + 1 \in GF(2^3)$$

$$C(x) = A(x) \cdot B(x) \text{ modulo } 2 = x^4 + x^3 + x + 1 = C'(x)$$

$$P(x) = \sum_{i=0}^m P_i x^i = P_0 x^0 + P_1 x^1 + P_2 x^2 + P_3 x^3, P_i \in GF(2)$$
$$= x^3 + x + 1$$

Divide  $C'(x)$  by  $P(x)$  and find the remainder

$$C(x) = x^2 + x = A(x) \cdot B(x) \text{ mod } P(x)$$

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

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# GALOIS FIELD IN AES

## SubBytes (Substitution Step)

- SubBytes is the first step in each AES round.
- It replaces each byte in the 4×4 matrix with a new byte using an S-box (Substitution box).
- The S-box is a fixed table that gives one output byte for each input byte.
- This makes the data non-linear and adds confusion, which means the relationship between the key and ciphertext becomes harder to guess.

## S-Box

- How it's created: The values are carefully constructed using mathematical operations in the finite field  $GF(2^8)$  including:
  - Taking the multiplicative inverse of each byte (except 0, which maps to 0).
  - Applying an affine transformation over  $GF(2)$  (simple bitwise operations).

byte  $\rightarrow GF(2^8)$  inverse  $\rightarrow$  affine  $\rightarrow$  S-box

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$y = mx + b$$



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# GALOIS FIELD IN AES

## MixColumn

MixColumn is a crucial part of the Diffusion Layer in AES.

To provide diffusion by mixing the data within each column of the state matrix. A change in a single input bit will affect all output bits after this operation.

### Calculate : A3 \* 02

$$A3(hex) = 10100011(binary) = x^7 + x^5 + x + 1$$

$$02(hex) = 00000010 = x^1$$

$$\begin{pmatrix} C_0 \\ C_1 \\ C_2 \\ C_3 \end{pmatrix} = \begin{pmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{pmatrix} \times \begin{pmatrix} B_0 \\ B_5 \\ B_{10} \\ B_{15} \end{pmatrix}$$

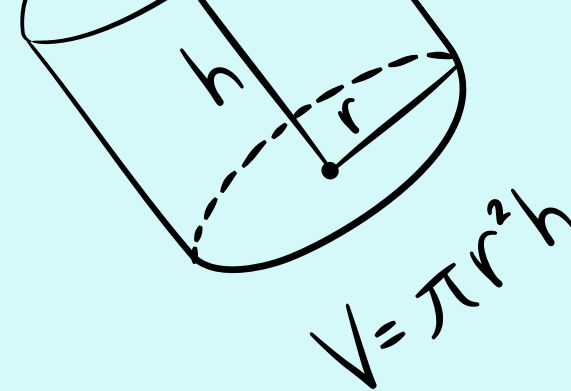
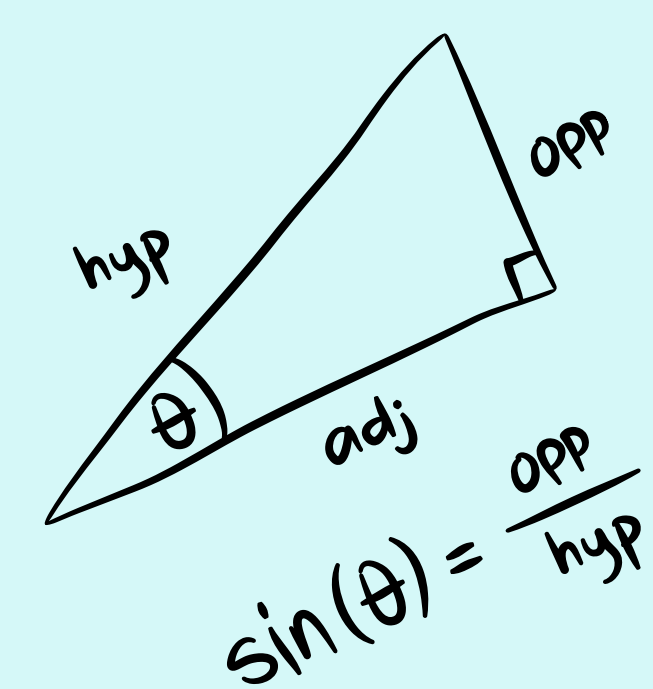
B\_0...B\_3 are the 4 bytes of the input column.  
C\_0...C\_3 are the 4 bytes of the transformed output column.  
01, 02, 03 are hexadecimal values representing elements in GF(2<sup>8</sup>).

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

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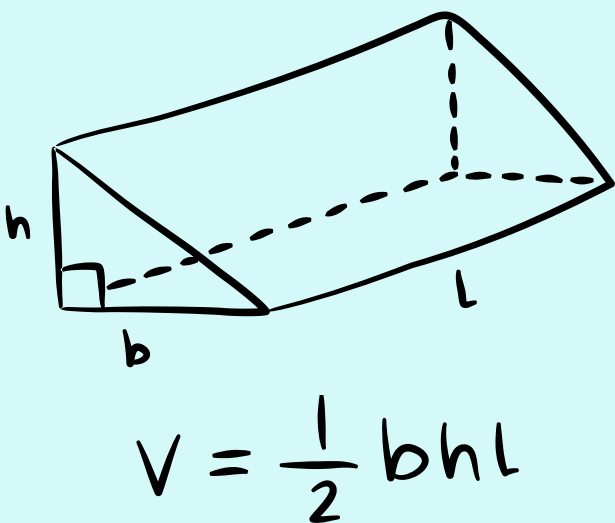


$$x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$a = \frac{V_f - V_i}{t}$$

**THANK YOU!**

$$mx + b$$



$$\frac{x}{a} + \frac{y}{b} = 1$$

$$ax^2 + bx + c = 0$$

