ACM-ICPC TEAM REFERENCE DOCUMENT

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	3.3 Strongly Connected Components	7	#include <bits stdc++.h=""></bits>	
	3.4 Bellman Ford Algorithm	8	#include <ext assoc_container.hpp="" pb_ds=""> // gp_hash_table<int, int=""> == hash map</int,></ext>	
	3.5 Bipartite Graph	8	#include <ext pb_ds="" tree_policy.hpp=""></ext>	
	3.6 Finding Articulation Points	9	using namespace std; using namespacegnu_pbds;	
	3.7 Finding Bridges	9	typedef long long ll; typedef unsigned long long ull;	
	3.8 Max Flow With Ford Fulkerson		typeder unsigned long long un, typedef long double ld;	
		9	typedef pair <int, int=""> pii; typedef pair<ll, ll=""> pll;</ll,></int,>	
	3.9 Max Flow With Dinic	10	typedef pair <double, double=""> pdd;</double,>	
	3.10 Max Flow With Dinic 2	10	template <typename t=""> using min_heap = priority_queue < T, vector <t>, greater <t>>;</t></t></typename>	
	3.11 Min Cut	11	template <typename t=""> using max_heap = priority_queue<</typename>	
	3.12 Number Of Paths Of Fixed Length .	11	T, vector <t>, less<t>>; template <typename t=""> using ordered_set = tree<t,< td=""></t,<></typename></t></t>	
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4	Geometry	11	template <typename k,="" typename="" v=""> using hashmap = gp_hash_table<k, v="">;</k,></typename>	
	4.1 2d Vector	11		
	4.2 Line	11	template <typename a,="" b="" typename=""> ostream& operator<<(ostream& out, pair<a, b=""> p) { out << "(" << p.first</a,></typename>	
	4.3 Convex Hull Gift Wrapping	11	<pre><< "," << p.second << ")"; return out;}</pre>	
	4.4 Convex Hull With Graham's Scan .	12	template <typename t=""> ostream& operator<<$($ostream& out vector<t> v) { out << "["; for(auto& x : v) out << x</t></typename>	
	4.5 Circle Line Intersection	12	<< ",_"; out << "]";return out;}	
	4.6 Circle Circle Intersection	12	template <typename t=""> ostream& operator<<(ostream& out set<t> v) { out << "{"; for(auto& $x: v$) out << x</t></typename>	
	4.7 Common Tangents To Two Circles .	12	<= ",∟"; out << "}"; return out; }	
	4.8 Number Of Lattice Points On Segment		template <typename k,="" typename="" v=""> ostream& operator<<(ostream& out, map<k, v=""> m) { out << "{"; for(auto&</k,></typename>	
	4.9 Pick's Theorem		e: m) out << e.first << ">_" << e.second << ",_"; out << "}"; return out; }	
			1	

```
template<typename K, typename V> ostream& operator<<(
       ostream& out, hashmap<K, V> m) { out << "{\"; for(auto& e : m) out << e.first << "\_->\_" << e.second <<< ",\_"; out << "\}"; return out; }
\#define\ FAST\_IO\ ios\_base::sync\_with\_stdio(false);\ cin.tie(
       NULL)
 #define TESTS(t) int NUMBER_OF_TESTS; cin >>
NUMBER_OF_TESTS; for(int t = 1; t <= NUMBER_OF_TESTS; t++) #define FOR(i, begin, end) for (int i = (begin) - ((begin) > (
       end)); i != (end) - ((begin) > (end)); i += 1 - 2 * ((
       begin) > (end))
#define sgn(a) ((a) > eps ? 1 : ((a) < -eps ? -1 : 0))
#define precise(x) fixed << set
precision(x) #define debug(x) cerr << ">^{"} << #x << "
^{"} << x <<
        endl;
#define pb push_back
#define rnd(a, b) (uniform_int_distribution<int>((a), (b))(
       rng))
 #ifndef LOCAL
     #define cerr if(0)cout
     #define endl "\n'
#endif
mt19937 rng(chrono::steady_clock::now().time_since_epoch()
const ld eps = 1e-14;
\begin{array}{l} {\rm const\ int\ oo} = 2e9; \\ {\rm const\ ll\ OO} = 2e18; \\ {\rm const\ ll\ MOD} = 1000000007; \end{array}
const int MAXN = 1000000;
int main() {
    FAST_IO;
    startTime();
    timeit("Finished");
    return 0;
```

1.2 Python Template

```
import sys
import re
from math import ceil, log, sqrt, floor

__local_run__ = False
if __local_run__:
    sys.stdin = open('input.txt', 'r')
    sys.stdout = open('output.txt', 'w')

def main():
    a = int(input())
    b = int(input())
    print(a*b)

main()
```

1.3 Compilation

```
# Simple compile
g++ -DLOCAL -O2 -o main.exe -std-c++17 -Wall -Wno-
unused-result -Wshadow main.cpp
# Debug
g++ -DLOCAL -std=c++17 -Wshadow -Wall -o main.exe
main.cpp -fsanitize=address -fsanitize=undefined -fuse-
ld=gold -D_GLIBCXX_DEBUG -g
```

1.4 Automatic Test

```
\# Linux Bash \# gen, main and stupid have to be compiled beforehand for((i=1;;++i)); do echo $i; ./gen $i > genIn; diff <(./main < genIn) <(./stupid < genIn) || break; done \# Windows CMD (@echo off FOR /L %%I IN (1,1,2147483647) DO (echo %%I gen.exe %%I > genIn main.exe < genIn > mainOut stupid.exe < genIn > stupidOut FC mainOut stupidOut || goto :eof )
```

1.5 Ternary Search

```
 \begin{array}{l} \mbox{double ternary\_search(double l, double r) } \{ \\ \mbox{while } (r - l > eps) \, \{ \\ \mbox{double } m1 = l + (r - l) \, / \, 3; \\ \mbox{double } m2 = r - (r - l) \, / \, 3; \\ \mbox{double } f2 = f(m1); \\ \mbox{double } f2 = f(m2); \\ \mbox{if } (f1 < f2) \\ \mbox{l} = m1; \\ \mbox{else} \\ \mbox{r} = m2; \\ \mbox{} \} \\ \mbox{return } f(l); \, / / \mbox{return the maximum of } f(x) \mbox{ in } [l, \, r] \\ \mbox{} \} \end{array}
```

1.6 Big Integer

```
const int base = 100000000000;
const\ int\ base\_digits = 9;
struct bigint {
     vector<int> a:
    int sign;
    int size() {
         if (a.empty()) return 0;
int ans = (a.size() - 1) * base_digits;
int ca = a.back();
         while (ca) ans++, ca \neq 10;
         return ans;
    bigint operator^(const bigint &v) {
         while (!y.isZero()) {
    if (y % 2) ans *= x;
        x *= x, y /= 2;
         return ans;
    string to_string() {
         stringstream ss;
ss << *this;
         string s;
         ss >> s;
         return s;
    int sumof() {
         string s = to\_string();
         int ans = 0;
         for (auto c : s) ans += c - 0;
         return ans;
    \operatorname{bigint}():\operatorname{sign}(1) {}
    bigint(long long v) {
         *{\rm this}={\rm v};
    bigint(const string &s) {
```

```
read(s);
void operator=(const bigint &v) {
       sign = v.sign;
       a = v.a:
void operator=(long long v) {
       sign = 1;
       a.clear();
       if (v < 0)
             \operatorname{sign} \stackrel{\cdot}{=} -1, \ v = -v;
       for (; v > 0; v = v / base)
a.push_back(v % base);
bigint operator+(const bigint &v) const {
       if\ (sign == v.sign)\ \{
               bigint res = v;
               for (int i = 0, carry = 0; i < (int)max(a.size(), v.a.
                      \begin{array}{l} \operatorname{res.a[i]} += \operatorname{carry} + (\mathrm{i} < (\operatorname{int}) \operatorname{a.size}(\underline{)} ? \ \operatorname{a[i]} : 0); \\ \operatorname{carry} = \operatorname{res.a[i]} >= \operatorname{base}; \\ \operatorname{if} \ (\operatorname{carry}) \ \operatorname{res.a[i]} -= \operatorname{base}; \end{array}
               return res;
       return *this - (-v);
bigint operator-(const bigint &v) const {
       if\ (sign == v.sign)\ \{
              if (abs() \ge v.abs()) {
                      bigint res = *this;
                      for (int i = 0, carry = 0; i < (int)v.a.size() ||
                              \begin{array}{l} {\rm carry; \ ++i) \ \{} \\ {\rm res.a[i] \ -= \ carry \ + \ (i < (int)v.a.size() \ ? \ v.a[i]} \end{array}
                              \begin{array}{c} \text{res.a[i]} = \text{carry} + (i) \\ \text{j} : 0); \\ \text{carry} = \text{res.a[i]} < 0; \end{array}
                             if (carry) res.a[i] += base;
                      res.trim();
                      return res;
               return -(v - *this);
       return *this + (-v);
void operator*=(int v) {
       if (v < 0) sign = -sign, v = -v;
for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i)
              \label{eq:continuous_state} \begin{array}{l} \overset{\  \, \mathsf{i}}{\mathrm{i}} = = (\mathrm{int}) \mathrm{a.size}()) \ \mathrm{a.push\_back}(0); \\ \mathrm{long} \ \mathrm{long} \ \mathrm{cur} = \mathrm{a}[\mathrm{i}] \ ^* \ (\mathrm{long} \ \mathrm{long}) \mathrm{v} + \mathrm{carry}; \\ \mathrm{carry} = (\mathrm{int}) (\mathrm{cur} \ / \ \mathrm{base}); \\ \mathrm{a}[\mathrm{i}] = (\mathrm{int}) (\mathrm{cur} \ \% \ \mathrm{base}); \end{array}
       trim();
bigint operator*(int v) const {
  bigint res = *this;
  res *= v;
       return res;
void operator*=(long long v) {
       if (v < 0) sign = -sign, v = -v;
       for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i)
               \begin{array}{l} if \ (i == (int)a.size()) \ a.push\_back(0); \\ long \ long \ cur = a[i] \ * (long \ long)v + carry; \\ carry = (int)(cur \ / \ base); \\ a[i] = (int)(cur \ \% \ base); \end{array} 
       trim();
bigint operator*(long long v) const {
   bigint res = *this;
       res^* = v;
       return res;
friend pair < bigint, bigint > divmod(const bigint &a1,
       const bigint &b1) {
int norm = base / (b1.a.back() + 1);
bigint a = a1.abs() * norm;
bigint b = b1.abs() * norm;
       bigint q, r;
       \begin{array}{l} {\rm q.a.resize(a.a.size());} \\ {\rm for~(int~i=a.a.size()-1;~i>=0;~i--)~\{} \end{array}
```

```
r *= base
         r \mathrel{+}= a.a[i];
         \begin{array}{l} \text{int s1} = \text{r.a.size}() <= \text{b.a.size}() ? 0 : \text{r.a[b.a.size}()]; \\ \text{int s2} = \text{r.a.size}() <= \text{b.a.size}() - 1 ? 0 : \text{r.a[b.a.size}) \end{array}
                 () - 1];
         int d = ((long long)base * s1 + s2) / b.a.back();
r -= b * d;
         while (r < 0) r += b, --d;
         q.a[i]=d;
    q.sign = a1.sign * b1.sign;
    r.sign = a1.sign;
    r.trim();
    return make_pair(q, r / norm);
bigint operator/(const bigint &v) const {
    return divmod(*this, v).first;
bigint operator%(const bigint &v) const {
    return divmod(*this, v).second;
void operator/=(int v) {
if (v < 0) sign = -sign, v = -v;
    for (int i = (int)a.size() - 1, rem = 0; i >= 0; --i) {
         long long cur = a[i] + rem * (long long)base;
         a[i] = (int)(cur / v);
         rem = (int)(cur \% v);
    trim();
bigint operator/(int v) const {
    bigint res = *this;
    res /= v;
    return res:
int operator%(int v) const {
    if (v < 0) \dot{v} = -\dot{v};
     int m = 0;
    for (int i = a.size() - 1; i >= 0; --i)

m = (a[i] + m * (long long)base) % v;

return m * sign;
void operator+=(const bigint &v) {
     *this = *this + v;
void operator==(const bigint &v) {
   *this = *this - v;
void operator*=(const bigint &v) {
     *this = *this * v;
void operator/=(const bigint &v) {
    *this = *this / v;
bool operator < (const bigint &v) const {
    if (sign != v.sign) return sign < v.sign;
    if (a.size() != v.a.size()) return a.size() * sign < v.a.size() * v.sign; for (int i = a.size() - 1; i >= 0; i-) if (a[i] != v.a[i]) return a[i] * sign < v.a[i] * sign;
    return false;
bool operator>(const bigint &v) const {
   return v < *this;</pre>
bool operator <= (const bigint &v) const {
    return !(v < *this);
bool operator>=(const bigint &v) const {
    return !(*this < v);
bool operator==(const bigint &v) const {
    return !(*this < v) && !(v < *this);
bool operator!=(const bigint &v) const {
    return *this < v || v < *this;
void trim() {
    while (!a.empty() && !a.back()) a.pop_back();
    if (a.empty()) sign = 1;
bool isZero() const {
     return \stackrel{\cdot}{a.empty()}|| \; (a.size() == 1 \; \&\& \; !a[0]); \\
```

```
bigint operator-() const {
   bigint res = *this;
     res.sign = -sign;
     return res;
bigint abs() const {
     bigint res = *this;
     res.sign *= res.sign;
     return res;
long long longValue() const {
     long long res = 0;
     for (int i = a.size() - 1; i >= 0; i--) res = res * base +
           a[i];
     return res * sign;
friend bigint gcd(const bigint &a, const bigint &b) {
return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint &b) {
return a / gcd(a, b) * b;
void read(const string &s) {
    sign = 1:
     a.clear();
     while (pos < (int)s.size() && (s[pos] == '-' || s[pos] == '+')) {
  if (s[pos] == '-') sign = -sign;
         ++pos;
     for (int i = s.size() - 1; i \ge pos; i -= base\_digits) {
         for (int j = max(pos, i - base\_digits + 1); j \le i; j
              x = x * 10 + s[j] - 0;
         a.push_back(x);
     trim();
friend istream & operator >> (istream & stream, bigint &v)
     string s;
     stream >> s;
     v.read(s);
     return stream;
friend ostream & operator << (ostream & stream, const
     bigint &v) {
if (v.sign == -1) stream << '-';
     \begin{array}{l} \text{stream} << (\text{v.a.empty}() ? 0 : \text{v.a.back}()); \\ \text{for (int } i = (\text{int})\text{v.a.size}() - 2; i >= 0; --i) \end{array}
         {\rm stream} << {\rm setw}({\rm base\_digits}) << {\rm setfill}('0') << {\rm v}.
                a[i];
    return stream;
static vector<int> convert_base(const vector<int> &a,
       int old_digits, int new_digits) {
     vector<long long> p(max(old_digits, new_digits) +
            1);
     p[0] = 1;
     for (int i = 1; i < (int)p.size(); i++)

p[i] = p[i - 1] * 10;
     vector<int> res;
     long long cur = 0;
    for cur_digits = 0;

for (int i = 0; i < (int)a.size(); i++) {

    cur += a[i] * p[cur_digits];

    cur_digits += old_digits;
         while (cur_digits >= new_digits) {
    res.push_back(int(cur % p[new_digits]));
              cur \mathrel{/}= p[new\_digits];
              cur\_digits -= new\_digits;
         }
     res.push_back((int)cur);
     while (!res.empty() && !res.back()) res.pop_back();
     return res:
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a, const vll &b) {
     int n = a.size();
     vll res(n + n);
     if (n <= 32) {
         for (int i = 0; i < n; i++)
              for (int j = 0; j < n; j++)
```

```
res[i + j] += a[i] * b[j];
                  return res
            int k = n >> 1;
            vll a1(a.begin(), a.begin() + k);
vll a2(a.begin() + k, a.end());
            vll b1(b.begin(), b.begin() + k);
            vll b2(b.begin() + k, b.end());
            vll a1b1 = karatsubaMultiply(a1, b1):
            vll a2b2 = karatsubaMultiply(a2, b2);
            \begin{array}{l} {\rm for} \ ({\rm int} \ i=0; \ i< k; \ i++) \ a2[i] \ += \ a1[i]; \\ {\rm for} \ ({\rm int} \ i=0; \ i< k; \ i++) \ b2[i] \ += \ b1[i]; \end{array}
            \begin{array}{l} vll \ r = karatsubaMultiply(a2, b2); \\ for \ (int \ i = 0; \ i < (int)a1b1.size(); \ i++) \ r[i] \ -= \ a1b1[i]; \\ for \ (int \ i = 0; \ i < (int)a2b2.size(); \ i++) \ r[i] \ -= \ a2b2[i]; \end{array}
            \begin{array}{l} {\rm for} \ ({\rm int} \ i=0; \ i<({\rm int})r.{\rm size}(); \ i++) \ {\rm res}[i+k] += r[i]; \\ {\rm for} \ ({\rm int} \ i=0; \ i<({\rm int})a1b1.{\rm size}(); \ i++) \ {\rm res}[i] \ += a1b1 \end{array}
                      [i];
            for (int i = 0; i < (int)a2b2.size(); i++) res[i + n] +=
                      a2b2[i];
            return res;
      bigint operator*(const bigint &v) const {
            vector<int> a6 = convert_base(this->a, base_digits,
                      6):
            vector < int > b6 = convert_base(v.a, base_digits, 6);
            vll x(a6.begin(), a6.end());
vll y(b6.begin(), b6.end());
            \label{eq:while (x.size() < y.size()) x.push_back(0);} while (y.size() < x.size()) y.push_back(0); while (x.size() & (x.size() - 1)) x.push_back(0), y. \\
                      push_back(0);
            vll c = karatsubaMultiply(x, y);
            bigint res;
            res.sign = sign * v.sign;
            for (int i = 0, carry = 0; i < (int)c.size(); i++) {
long long cur = c[i] + carry;
res.a.push_back((int)(cur % 1000000));
                   carry = (int)(cur / 1000000);
            res.a = convert_base(res.a, 6, base_digits);
            res.trim();
            return res;
      }
};
```

2 Data Structures

2.1 Disjoin Set Union

```
struct DSU {
      vector<int> par;
      {\tt vector}{<} {\tt int}{>} \ {\tt sz};
      DSU(int n) {
             FOR(i, 0, n) {
                    par.pb(i);
                     \operatorname{sz.pb}(1);
      }
      int find(int a) {
             return par[a] = par[a] == a ? a : find(par[a]);
      \begin{array}{c} bool \ same(int \ a, \ int \ b) \ \{\\ return \ find(a) == find(b); \end{array}
       void unite(int a, int b) {
             a = find(a);
              b = find(b);
             \begin{array}{l} -1 & \text{if } (S); \\ \text{if } (\operatorname{sz}[a] > \operatorname{sz}[b]) & \operatorname{swap}(a, b); \\ \operatorname{sz}[b] & += \operatorname{sz}[a]; \end{array}
             par[a] = b;
};
```

2.2 Fenwick Tree Point Update And Range Query

```
struct Fenwick {
      vector<ll> tree;
     int n;
     Fenwick(){}
     Fenwick(\overset{.}{int}\ \_n)\ \{
           n = n;
           tree = vector < ll > (n+1, 0);
      \begin{array}{l} \text{void add(int } i, \ ll \ val) \ \{ \ // \ arr[i] \ += \ val \\ \text{for(; } i <= n; \ i \ += \ i\&(-i)) \ tree[i] \ += \ val; \end{array} 
     il get(int i) { // arr[i]
           return sum(i, i);
     ll sum(int i) { // arr[1]+...+arr[i]
           \ln ans = 0:
           for(; i > 0; i \rightarrow i\&(-i)) ans += tree[i];
           return ans;
     \hat{l}l \text{ sum}(\text{int } l, \text{ int } r)  {// \text{arr}[l]+...+\text{arr}[r]
           return sum(r) - sum(l-1);
};
```

2.3 Fenwick Tree Range Update And Point Query

2.4 Fenwick Tree Range Update And Range Query

```
struct RangedFenwick {
    Fenwick F1, F2; // support range query and point update
    RangedFenwick(int _n) {
        F1 = Fenwick(_n+1);
        F2 = Fenwick(_n+1);
    }
    void add(int l, int r, ll v) { // arr[l..r] += v
        F1.add(l, v);
        F1.add(r+1, -v);
        F2.add(l, v*(l-1));
        F2.add(l, v*[l-1]);
        F2.add(r+1, -v*r);
    }
    ll sum(int i) { // arr[l..i]
        return F1.sum(i)*i-F2.sum(i);
    }
    ll sum(int l, int r) { // arr[l..r]
        return sum(r)-sum(l-1);
    }
};
```

2.5 Fenwick 2D

2.6 Segment Tree

```
struct SegmentTree {
    int n:
    vector<ll> t;
    const ll IDENTITY = 0; // OO for min, -OO for max, ...
    ll f(ll a, ll b) {
         return a+b;
    SegmentTree(int _n) {
n = _n; t = vector < ll > (4*n, IDENTITY);
    SegmentTree(vector<ll>& arr) {
         n = arr.size(); t = vector < ll > (4*n, IDENTITY);
         build(arr, 1, 0, n-1);
    void build(vector<ll>& arr, int v, int tl, int tr) {
         if(tl == tr) \ \{ \ t[v] = arr[tl]; \ \}
             int tm = (tl+tr)/2;
             \begin{array}{l} build(arr,\,2^*v,\,tl,\,tm);\\ build(arr,\,2^*v{+}1,\,tm{+}1,\,tr);\\ t[v]\,=\,f(t[2^*v],\,t[2^*v{+}1]); \end{array}
         }
     // sum(1, 0, n-1, l, r)
    ll sum(int v, int tl, int tr, int l, int r) {
         if (l > r) \ return \ IDENTITY; \\
         tm+1, tr, max(l, tm+1), r));
    // update(1, 0, n-1, i, v) void update(int v, int tl, int tr, int pos, ll newVal) { if(tl == tr) { t[v] = newVal; }}
         else {
             int tm = (tl+tr)/2;
             if(pos <= tm) update(2*v, tl, tm, pos, newVal);
             else update(2*v+1, tm+1, tr, pos, newVal);

t[v] = f(t[2*v],t[2*v+1]);
        }
    }
};
```

2.7 Segment Tree With Lazy Propagation

```
// Add to segment, get maximum of segment
struct LazySegTree {
   int n;
    vector<ll> t, lazy;
   LazySegTree(int \_n) \ \{
       n = n; t = vector < ll > (4*n, 0); lazy = vector < ll > (4*
             n, 0);
   LazySegTree(vector<ll>& arr) {
       n = n; t = \text{vector} < \text{ll} > (4*n, 0); lazy = vector < ll>(4*
             n, 0);
       build
(arr, 1, 0, n-1); // same as in simple Segment
Tree
   void push(int v) {
       d push(int v) {
t[v^*2] += lazy[v];
lazy[v^*2] += lazy[v];
t[v^*2+1] += lazy[v];
lazy[v^*2+1] += lazy[v];
       lazy[v] = 0;
   void update(int v, int tl, int tr, int l, int r, ll addend) {
       if (l > r)
           return:
       if (1 == tl \&\& tr == r) \{ t[v] += addend; \}
           lazy[v] += addend;
       } else {
           int tm = (tl + tr) / 2;
           }
   }
   int query(int v, int tl, int tr, int l, int r) \{
       if (l > r)
           return -OO;
       if (tl == tr)
           return t[v];
       push(v);
       };
```

2.8 Treap

```
namespace Treap {
   struct Node {
Node *l, *r;
        ll key, prio, size;
        Node() {} 
Node(ll key) : key(key), l(nullptr), r(nullptr), size(1) { 
prio = rand() ^ (rand() << 15);
        }
   };
   typedef Node* NodePtr;
   int sz(NodePtr n) {
       return n ? n->size : 0;
    void recalc(NodePtr n) {
        if (!n) return;
        n-size = sz(n-sl) + 1 + sz(n-sr); // add more
              operations here as needed
    void split(NodePtr tree, ll key, NodePtr& l, NodePtr& r)
        if (!tree) {
            l = r = nullptr;
        else if (key < tree->key) {
```

```
split(tree->l, key, l, tree->l);
            r = tree;
       else {
            split(tree->r, key, tree->r, r);
            l = tree;
       recalc(tree);
   void merge(NodePtr& tree, NodePtr l, NodePtr r) {
       if (!l || !r) {
tree = 1 ? l : r;
        else if (l->prio > r->prio) {
            merge(l->r, l->r, r);
            tree = 1;
       else {
            merge(r->l, l, r->l);
            tree = r;
        recalc(tree);
   }
    void insert(NodePtr& tree, NodePtr node) {
       if (!tree) \{
           tree = node;
       else if (node->prio > tree->prio) {
            split(tree, node->key, node->l, node->r);
        else {
           insert(node->key < tree->key ? tree->l : tree->r,
                  node);
        recalc(tree);
   void erase(NodePtr tree, ll key) {
       if (!tree) return;
if (tree->key == key) {
            merge(tree, tree->l, tree->r);
            \dot{\rm erase}(\rm key < \rm tree-> key ? \rm tree-> l : \rm tree-> r, \, key);
       recalc(tree);
   }
   void print(NodePtr t, bool newline = true) {
       if (!t) return;
       print(t->l, false);

cout << t->key << "\left";
       print(t->r, false);
        if (newline) cout << endl;
}
```

2.9 Implicit Treap

```
}
void\ push(NodePtr\ n)\ \{
    if (n && n->rev) {
n->rev = false;
         swap(n->l, n->r);
         if (n->l) n->l->rev = 1;
         if (n->r) n->r->rev ^= 1;
}
void recalc(NodePtr n) {
    if (!n) return;
    n-size = sz(n-sl) + 1 + sz(n-sr);
    n\text{-}{>}sum = getSum(n\text{-}{>}l) \, + \, n\text{-}{>}val \, + \, getSum(n\text{-}{>}r);
void split(NodePtr tree, ll key, NodePtr& l, NodePtr& r)
    push(tree);
    if (!tree) {
        l = r = nullptr;
    else if (key \leq sz(tree->l)) {
        split(tree->l, key, l, tree->l);
    else {
         \begin{array}{l} \mbox{split}(\mbox{tree-}{>}\mbox{r},\;\mbox{key-sz}(\mbox{tree-}{>}\mbox{l})\mbox{-}1,\;\mbox{tree-}{>}\mbox{r},\;\mbox{r});\\ \mbox{l} = \mbox{tree}; \end{array} 
    recalc(tree);
void merge(NodePtr& tree, NodePtr l, NodePtr r) {
    push(l); push(r);
if (!l || !r) {
         tree = \hat{l} ? 1 : r;
    else if (l->prio > r->prio) {
         \mathrm{merge}(l{-}{>}r,\; l{-}{>}r,\; r);
         tree = 1:
    élse {
         merge(r->l, l, r->l);
         tree = r;
    recalc(tree);
void insert
(NodePtr& tree, T val, int pos) {
    if (!tree) {
        tree = new Node < T > (val);
         return;
     NodePtr L, R;
    split(tree, pos, L, R);
    merge(L, L, new Node<T>(val));
    merge(tree, L, R);
    recalc(tree);
void reverse(NodePtr tree, int l, int r) {
    NodePtr t1, t2, t3;
     split(tree, l, t1, t2);
    split(t2, r - l + 1, t2, t3);

if(t2) t2->rev = true;
    merge(t2, t1, t2);
    merge(tree, t2, t3);
void print(NodePtr t, bool newline = true) {
    push(t);
    if (!t) return;
    print(t->l, false);
    cout << t->val << "";
    print(t->r, false);
    if (newline) cout << endl;
NodePtr fromArray(vector<T> v) {
     NodePtr t = nullptr;
    FOR(i, 0, (int)v.size()) {
         \mathrm{insert}(t,\,v[i],\,i);
    return t:
```

```
}

ll calcSum(NodePtr t, int l, int r) {
    NodePtr L, R;
    split(t, l, L, R);
    NodePtr good;
    split(R, r - l + 1, good, L);
    return getSum(good);
}

};
/* Usage: ImplicitTreap<int> t;
Node<int> tree = t.fromArray(someVector); t.reverse(tree, l, r); ...
*/
```

2.10 Trie

```
struct Trie {
   const int ALPHA = 26;
   const char BASE = 'a';
   vector<vector<int>> nextNode;
   vector<int> mark;
   int nodeCount;
   Trie() {
       nextNode = vector<vector<int>>(MAXN, vector<int
            >(ALPHA, -1));
       mark = vector < int > (MAXN, -1);
       nodeCount = 1;
   void insert(const string& s, int id) {
       int curr = 0:
       FOR(i, 0, (int)s.length()) {
          int c = s[i] - BASE;
           if(nextNode[curr][c] == -1) {
              nextNode[curr][c] = nodeCount++;
          curr = nextNode[curr][c];
       mark[curr] = id;
   bool exists(const string& s) {
       int curr = 0:
       FOR(i, 0, (int)s.length()) {

int c = s[i] - BASE;
           if(nextNode[curr][c] == -1) return false;
           curr = nextNode[curr][c];
       return mark[curr] != -1;
};
```

3 Graphs

3.1 Dfs With Timestamps

```
\label{eq:control_vector} \begin{split} & \operatorname{vector} < \operatorname{vector} < \operatorname{int} >> \operatorname{adj}; \\ & \operatorname{vector} < \operatorname{int} > \operatorname{tIn}, \ \operatorname{tOut}, \ \operatorname{color}; \\ & \operatorname{int} \ \operatorname{dfs\_timer} = 0; \\ & \operatorname{void} \ \operatorname{dfs}(\operatorname{int} \ v) \ \{ \\ & \operatorname{tIn}[v] = \operatorname{dfs\_timer} + +; \\ & \operatorname{color}[v] = 1; \\ & \operatorname{for} \ (\operatorname{int} \ u : \operatorname{adj}[v]) \\ & \operatorname{if} \ (\operatorname{color}[u] == 0) \\ & \operatorname{dfs}(u); \\ & \operatorname{color}[v] = 2; \\ & \operatorname{tOut}[v] = \operatorname{dfs\_timer} + +; \\ \} \end{split}
```

3.2 Lowest Common Ancestor

```
int n, l; // l == logN (usually about ~20)
vector<vector<int>> adj;
int timer:
vector<int> tin, tout:
vector<vector<int>> up;
void dfs(int v, int p)
    {\rm tin}[v] = + + {\rm timer};
    up[v][0] = p;
    // \text{wUp[v][0]} = \text{weight[v][u]}; // <- \text{path weight sum to } 2^i
            th ancestor
    \begin{array}{l} \text{for (int $i=1$; $i<=1$; $++$i)} \\ up[v][i] &= up[up[v][i-1]][i-1]; \\ // & wUp[v][i] &= wUp[v][i-1] + wUp[up[v][i-1]][i-1]; \end{array}
    for (int u : adj[v]) {
         if (u != p)
              dfs(u, v);
    tout[v] = ++timer;
}
bool isAncestor(int u, int v)
    \operatorname{return} \ \operatorname{tin}[u] <= \operatorname{tin}[v] \ \&\& \ \operatorname{tout}[v] <= \operatorname{tout}[u];
int lca(int u, int v)
    if (isAncestor(u, v))
          return u;
    if (isAncestor(v, u))
          return v;
    for (int i = 1; i >= 0; --i) {
         if \; (!isAncestor(up[u][i], \; v)) \\
              u = up[u][i];
    return up[u][0];
}
void preprocess(int root) \{
    tin.resize(n);
    tout.resize(n);
    timer=0;
    l = ceil(log2(n));
    up.assign(n, vector<int>(l + 1));
    dfs(root, root);
```

3.3 Strongly Connected Components

```
vector < vector<int> > g, gr; // adjList and reversed adjList
vector<bool> used;
vector<int> order, component;
void dfs1 (int v) {
   used[v] = true;
    for (size_t i=0; i<g[v].size(); ++i)
        if (!used[ g[v][i] ])
           dfs1 (g[v][i])
   order.push_back (v);
}
void dfs2 (int v) {
    used[v] = true;
    component.push_back (v);
   for (size_t i=0; i<gr[v].size(); ++i) if (!used[ gr[v][i] ])
            dfs2 (gr[v][i]);
}
\mathrm{int}\ \mathrm{main}()\ \{
   int n;
    // read n
   for (;;) {
        int a, b;
        // read edge a -> b
        g[a].push_back (b);
```

```
gr[b].push_back (a);
}

used.assign (n, false);
for (int i=0; i<n; ++i)
    if (!used[i])
        dfs1 (i);
used.assign (n, false);
for (int i=0; i<n; ++i) {
    int v = order[n-1-i];
    if (!used[v]) {
        dfs2 (v);
        // do something with the found component
        component.clear(); // components are generated in
        toposort-order
    }
}</pre>
```

3.4 Bellman Ford Algorithm

```
struct Edge
{
    int a, b, cost;
};

int n, m, v; // v - starting vertex
vector<Edge> e;

/* Finds SSSP with negative edge weights.
    * Possible optimization: check if anything changed in a
        relaxation step. If not - you can break early.
    * To find a negative cycle: perform one more relaxation step.
    If anything changes - a negative cycle exists.

*/
void solve() {
    vector<int> d (n, oo);
    d[v] = 0;
    for (int i=0; i<n-1; ++i)
        for (int j=0; j<m; ++j)
        if (d[e[j].a] < oo)
        d[e[j].b] = min (d[e[j].b], d[e[j].a] + e[j].cost);
    // display d, for example, on the screen
}</pre>
```

3.5 Bipartite Graph

```
class BipartiteGraph {
private:
    {\tt vector}{<} {\tt int}{>} \_{\tt left}, \_{\tt right};
    vector<vector<int>> _adjList;
vector<int> _matchR, _matchL;
     vector<bool> _used;
    bool \underline{\phantom{a}} kuhn(int v) {
         if (_used[v]) return false;
         \_used[v] = true;
FOR(i, 0, (int)_adjList[v].size()) {
              int to = _adjList[v][i] - _left.size();
              if (\_matchR[to] == -1 \mid\mid \_kuhn(\_matchR[to])) \; \{
                   _{\text{matchR[to]}} = v;
                    matchL[v] = to;
                   return true;
              }
         return false;
             _addReverseEdges() {
    void
         FOR(i, 0, (int)_right.size()) {
    if (_matchR[i] != -1) {
                   \_adjList[\_left.size() + i].pb(\_matchR[i]);
         }
    void _dfs(int p) {
         if (_used[p]) return;
_used[p] = true;
         for (auto x : _adjList[p]) {
```

```
_{dfs(x);
         }
    vector<pii> _buildMM() {
         vector<pair<int, int> > res;
FOR(i, 0, (int)_right.size()) {
             if (_matchR[i] != -1) {
                  res.push\_back(make\_pair(\_matchR[i],\ i));
         }
         return res;
public:
    void addLeft(int x) {
         _left.pb(x);
_adjList.pb({});
_matchL.pb(-1);
         _used.pb(false);
    void addRight(int x) {
        _right.pb(x);
_adjList.pb({});
           matchR.pb(-1);
         _used.pb(false);
    void addForwardEdge(int l, int r) \{
         \_adjList[l].pb(r + \_left.size());
    void addMatchEdge(int l, int r) {
        if(l!=-1) _matchL[l] = r;
if(r!=-1) _matchR[r] = l;
    // Maximum Matching
    vector<pii> mm() {
         _matchR = vector<int>(_right.size(), -1);
_matchL = vector<int>(_left.size(), -1);
             ^ these two can be deleted if performing MM on
                already partially matched graph
         \_used = vector < bool > (\_left.size() + \_right.size(),
                false);
         bool path_found;
             fill(_used.begin(), _used.end(), false);
             \begin{array}{l} path\_found = false; \\ FOR(i,\ 0,\ (int)\_left.size())\ \{\\ if\ (\_matchL[i] < 0\ \&\&\ !\_used[i])\ \{ \end{array}
                      path\_found = _kuhn(i);
         } while (path_found);
         return _buildMM();
    }
       Minimum Edge Cover
    // Algo: Find MM, add unmatched vertices greedily.
    (anto x : __atplise();

int ridx = x - _left.size();

if (_matchR[ridx] == -1) {

    ans.pb({ i, ridx });

    _matchR[ridx] = i;
                      }
                 }
             }
          \begin{array}{l} FOR(i,\,0,\,(int)\_left.size()) \; \{\\ if \; (\_matchL[i] == -1 \;\&\& \; (int)\_adjList[i].size() \; > \end{array} 
                   int ridx =
                  ans.pb(\{ i, ridx \});
             }
         return ans;
        Minimum Vertex Cover
    // Algo: Find MM. Run DFS from unmatched vertices
           from the left part.
```

```
// MVC is composed of unvisited LEFT and visited
           RIGHT vertices.
    pair<vector<int>, vector<int>> mvc(bool runMM =
           true) {
        if (runMM) mm();
         _addReverseEdges();
         \begin{array}{ll} & \text{fill(\_used.begin(), \_used.end(), false);} \\ & \text{FOR(i, 0, (int)\_left.size()) } \\ & \text{if } (\_matchL[i] == -1) \end{array} 
                  _{\mathrm{dfs}(i)};
         vector<int> left, right;
         FOR(i, 0, (int)\_left.size()) {
             if (!_used[i]) left.pb(i);
         FOR(i, 0, (int)_right.size()) {
             \begin{array}{l} \text{if } (\_\text{used}[i+(\text{int})\_\text{left.size}()]) \text{ right.pb}(i); \\ \end{array}
         return { left,right };
    // Maximal Independent Vertex Set
    // Algo: Find complement of MVC.
    pair<vector<int>, vector<int>> mivs(bool runMM =
           true) {
         auto m = mvc(runMM);
         vector<br/>bool> containsL(_left.size(), false), containsR(
        _right.size(), false);
for (auto x : m.first) containsL[x] = true;
         for (auto x : m.second) contains R[x] = true;
         vector<int> left, right;
         FOR(i, 0, (int)\_left.size())
             if (!containsL[i]) left.pb(i);
         FOR(i, 0, (int)_right.size()) {
             if (!containsR[i]) right.pb(i);
         return { left, right };
};
```

3.6 Finding Articulation Points

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> tin, fup;
int timer;
void processCutpoint(int v) {
      // problem-specific logic goes here
// it can be called multiple times for the same v
 \begin{array}{l} \mbox{void dfs(int $v$, int $p=-1$) \{} \\ \mbox{visited[$v$] = true;} \\ \mbox{tin[$v$] = fup[$v$] = timer++;} \end{array} 
      int children=0;
      for (int to : adj[v]) {
            if (to == p) continue;
           if (visited[to]) {
                 fup[v] = \min(fup[v],\, tin[to]);
            } else {
                  dfs(to, v); 
fup[v] = min(fup[v], fup[to]);
                   \begin{array}{l} \operatorname{int}(\operatorname{dip}[t], \operatorname{dip}[t]), \\ \operatorname{if}(\operatorname{fup}[t]) >= \operatorname{tin}[v] & & \operatorname{p!}=-1) \\ \operatorname{processCutpoint}(v); \\ \end{array} 
                  ++children;
           }
      if(p == -1 \&\& children > 1)
           processCutpoint(v);
void findCutpoints() {
      timer = 0;
      visited.assign(n, false);
      tin.assign(n, -1);
      fup.assign(n, -1);
```

```
 \begin{array}{c} \text{for (int } i = 0; \ i < n; \ ++i) \ \{ \\ & \text{if (!visited[i])} \\ & \text{dfs (i);} \\ \} \\ \} \end{array}
```

3.7 Finding Bridges

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<br/>bool> visited;
vector<int> tin, fup;
int timer:
void processBridge(int u, int v) {
     // do something with the found bridge
void dfs(int v, int p = -1) {
    visited[v] = true;
tin[v] = fup[v] = timer++;
    for (int to : adj[v]) {
          if (to == p) continue;
          if (visited[to]) {
               fup[v] = \min(fup[v], \, tin[to]);
          } else {
              \begin{aligned} &\operatorname{dfs}(\operatorname{to}, \, \operatorname{v}); \\ &\operatorname{fup}[\operatorname{v}] = \min(\operatorname{fup}[\operatorname{v}], \, \operatorname{fup}[\operatorname{to}]); \end{aligned}
               if (fup[to] > tin[v])
processBridge(v, to);
    }
}
// Doesn't work with multiple edges
// But multiple edges are never bridges, so it's easy to check
void findBridges() {
    timer = 0:
     visited.assign(n, false);
    tin.assign(n, -1);
     fup.assign(n, -1);
     bridges.clear();
    FOR(i, 0, n)
          if (!visited[i])
              dfs(i);
}
```

3.8 Max Flow With Ford Fulkerson

```
struct Edge {
      int to, next;
      ll f, c;
      int idx dir:
      int from;
};
vector<Edge> edges;
{\tt vector}{<} {\tt int}{>} \; {\tt first};
void addEdge(int a, int b, ll c, int i, int dir) {
     datablege(int a, int b, in c, int i, int d
edges.pb({ b, first[a], 0, c, i, dir, a });
edges.pb({ a, first[b], 0, 0, i, dir, b });
first[a] = edges.size() - 2;
first[b] = edges.size() - 1;
}
void init() {
      cin >> n >> m;
edges.reserve(4 * m);
      \begin{array}{l} first = vector < int > (n, -1); \\ FOR(i, 0, m) \end{array} 
             int a, b, c;
             cin >> a >> b >> c;
             addEdge(a, b, c, i, 1);
```

```
addEdge(b, a, c, i, -1);
}
int cur time = 0:
vector<int> timestamp;
ll dfs(int v, ll flow = OO)
      \begin{array}{ll} i\dot{r}\,(v==n-1)\; return\; flow;\\ timestamp[v]=cur\_time;\\ for\;(int\;e=first[v];\;e\;!=-1;\;e=edges[e].next)\; \{ \end{array} 
          if (edges[e].f < edges[e].c && timestamp[edges[e].to] != cur_time) {
               int pushed = dfs(edges[e].to, min(flow, edges[e].c -
               edges[e].f);
if (pushed > 0) {
  edges[e].f += pushed;
  edges[e ^ 1].f -= pushed;
                    return pushed;
          }
     return 0;
}
ll \max Flow()  {
     cur\_time = 0;
     timestamp = vector{<}int{>}(n,\,0);
     ll f = 0, add;
     while (true) {
         cur\_time++; add = dfs(0);
          if (add > 0)
               f += add;
          élse {
               break;
     return f;
}
```

3.9 Max Flow With Dinic

```
struct Edge {
       int f, c;
       int to:
       pii revIdx;
       int dir;
       int idx;
};
int n, m;
vector<Edge> adjList[MAX_N];
int level[MAX_N];
void addEdge(int a, int b, int c, int i, int dir) {
      int idx = adjList[a].size();
int revIdx = adjList[b].size();
adjList[a].pb({ 0,c,b, {b, revIdx} ,dir,i });
adjList[b].pb({ 0,0,a, {a, idx} ,dir,i });
bool bfs(int s, int t) {
       \begin{split} & FOR(i,\;0,\;n)\;level[i] = \text{-1}; \\ & level[s] = 0; \\ & queue{<}int{>}\;Q; \end{split}
       \hat{Q}.push(s);
       while (!Q.empty()) {
              auto t = Q.front(); Q.pop();
             \begin{array}{l} \text{auto } t = \lozenge, \text{mone}(), \ \circledast, \text{res}(), \\ \text{for } (\text{auto } x : \text{adjList[t]}) \ \{ \\ \text{if } (\text{level[x.to]} < 0 \ \&\& \ x.f < x.c) \ \{ \\ \text{level[x.to]} = \text{level[t]} + 1; \end{array}
                            Q.push(x.to);
             }
       return level[t] >= 0;
}
int send(int u, int f, int t, vector<int>& edgeIdx) {
       if (u == t) return f;
```

```
for (; \operatorname{edgeIdx}[u] < \operatorname{adjList}[u].\operatorname{size}(); \operatorname{edgeIdx}[u]++) {
           (, eugerdx[u] \ adjinst[u].size(), eugerdx[u]++) {
    auto& e = adjList[u][edgeIdx[u]];
    if (level[e.to] == level[u] + 1 && e.f < e.c) {
        int curr_flow = min(f, e.c - e.f);
        int next_flow = send(e.to, curr_flow, t, edgeIdx);
    }
                  if (\text{next\_flow} > 0) {
                        e.f += next_flow;
                        adjList[e.revIdx.first][e.revIdx.second].f \mathrel{\textit{-}}=
                                 next_flow;
                        return next_flow;
                  }
           }
      return 0;
}
int maxFlow(int s, int t) {
     int f = 0;
      while (bfs(s, t)) {
            vector < int > edgeIdx(n, 0);
            while (int extra = send(s, oo, t, edgeIdx)) {
                 f += extra;
            }
     return f;
\mathrm{void}\ \mathrm{init}()\ \{
     cin >> n >> m;

FOR(i, 0, m) {
           int a, b, c;
            cin >> a >> b >> c;
            addEdge(a,\,b,\,c,\,i,\,1);
            addEdge(b, a, c, i, -1);
}
```

3.10 Max Flow With Dinic 2

```
struct FlowEdge {
   int v, u;
    long long cap, flow = 0;
    FlowEdge(int v, int u, long long cap): v(v), u(u), cap(cap
struct Dinic {
   const long long flow_inf = 1e18;
    vector<FlowEdge> edges;
    vector < vector < int >> adj
   int n, m = 0;
   int s, t;
   vector<int> level, ptr;
   queue{<}int{>}\ q;
   Dinic(int\ n,\ int\ s,\ int\ t):n(n),\,s(s),\,t(t)\ \{
       adj.resize(n);
       level.resize(n);
       ptr.resize(n);
   }
   void add_edge(int v, int u, long long cap) {
       edges.push\_back(FlowEdge(v,\,u,\,cap));
       edges.push\_back(FlowEdge(u,\,v,\,0));
       adj[v].push\_back(m);
       adj[u].push\_back(m + 1);
       m += 2;
   bool bfs() {
       while (!q.empty()) {
           int v = q.front();
           q.pop();
           for (int id : adj[v]) {
               if (edges[id].cap - edges[id].flow < 1)
                  continue;
               if (level[edges[id].u] != -1)
                  continue;
               level[edges[id].u] = level[v] + 1;
               q.push(edges[id].u);
```

```
return level[t] != -1;
   }
   long long dfs(int v, long long pushed) \{
        if (pushed == 0)
            return 0;
        if (v == t)
           return pushed;
       for\ (int\&\ cid=ptr[v];\ cid<(int)adj[v].size();\ cid++)
            int id = adj[v][cid];
            int u = edges[id].u;
            if (level[v] + 1 != level[u] || edges[id].cap - edges[id]
                  ].flow < 1)
                continue;
            long\ long\ tr = dfs(u,\,min(pushed,\,edges[id].cap\ \text{-}
                 edges[id].flow));
            if (tr == 0)
                continue;
            edges[id].flow += tr;
            return tr;
        return 0;
   \log\,\log\,\mathrm{flow}()\,\,\{
       long long f = 0; while (true) {
            fill(level.begin(), level.end(), -1);
            level[s] = 0;
            q.push(s);
            if (!bfs())
                break:
            fill(ptr.begin(), ptr.end(), 0);
while (long long pushed = dfs(s, flow_inf)) {
                f += pushed;
       return f;
   }
};
```

3.11 Min Cut

```
\label{eq:continuity} \begin{array}{l} \operatorname{init}();\\ \operatorname{ll}\ f = \operatorname{maxFlow}();\ //\ \operatorname{Ford-Fulkerson}\\ \operatorname{cur\_time} + +;\\ \operatorname{dfs}(0);\\ \operatorname{set} < \operatorname{int} > \operatorname{cc};\\ \operatorname{for}\ (\operatorname{auto}\ e : \operatorname{edges})\ \{\\ \operatorname{if}\ (\operatorname{timestamp}[\operatorname{e.from}] == \operatorname{cur\_time}\ \&\&\ \operatorname{timestamp}[\operatorname{e.to}]\ !=\\ \operatorname{cur\_time})\ \{\\ \operatorname{cc.insert}(\operatorname{e.idx});\\ \}\\ \}\\ //\ (\#\ \operatorname{of}\ \operatorname{edges}\ \operatorname{in}\ \operatorname{min-cut},\ \operatorname{capacity}\ \operatorname{of}\ \operatorname{cut})\\ //\ [\operatorname{indices}\ \operatorname{of}\ \operatorname{edges}\ \operatorname{forming}\ \operatorname{the}\ \operatorname{cut}]\\ \operatorname{cout}\ << \operatorname{cc.size}()\ << \ ``\ ''\ << f\ << \operatorname{endl};\\ \operatorname{for}\ (\operatorname{auto}\ x : \operatorname{cc})\ \operatorname{cout}\ << x + 1 << \ `'\ '';\\ \end{array}
```

3.12 Number Of Paths Of Fixed Length

Let G be the adjacency matrix of a graph. Then $C_k = G^k$ gives a matrix, in which the value $C_k[i][j]$ gives the number of paths between i and j of length k.

3.13 Shortest Paths Of Fixed Length

Define $A \odot B = C \iff C_{ij} = min_{p=1..n}(A_{ip} + B_{pj})$. Let G be the adjacency matrix of a graph. Also, let $L_k = G \odot ... \odot G = G^{\odot k}$. Then the

value $L_k[i][j]$ denotes the length of the shortest path between i and j which consists of exactly k edges.

4 Geometry

4.1 2d Vector

```
template <typename T>
struct Vec {
    T x, y;
    Vec(): x(0), y(0) {}

Vec(T _x, T _y): x(_x), y(_y) {}

Vec operator+(const Vec& b) {
       return Vec < T > (x+b.x, y+b.y);
    Vec operator-(const Vec& b) {
return Vec<T>(x-b.x, y-b.y);
    \acute{\text{V}}ec operator*(T c) {
       return Vec(x*c, y*c);
    T operator*(const Vec& b) {
       \mathrm{return}\ x^*\mathrm{b.x} + y^*\mathrm{b.y};
    T operator^(const Vec& b) {
       return x*b.y-y*b.x;
    bool operator<(const Vec& other) const {
       if(\hat{x} == other.x) return y < other.y;
        return x < other.x;
    bool operator==(const Vec& other) const {
        return x==other.x && y==other.y;
    bool operator!=(const Vec& other) const {
       return !(*this == other);
    friend ostream& operator<<(ostream& out, const Vec& v)
        return out << "(" << v.x << "," << v.y << ")";
   friend istream& operator>>(istream& in, Vec<T>& v) {
       return in >> v.x >> v.v;
   T norm() { // squared length
return (*this)*(*this);
    ld len() {
       return sqrt(norm());
    ld angle(const Vec& other) { // angle between this and
          other vector
        return\ acosl((*this)*other/len()/other.len());\\
    Vec perp() {
       return Vec(-y, x);
   Cross product of 3d vectors: (ay*bz-az*by, az*bx-ax*bz, ax
      *by-ay*bx)
```

4.2 Line

```
template <typename T>
struct Line { // expressed as two vectors
    Vec<T> start, dir;
    Line() {}
    Line(Vec<T> a, Vec<T> b): start(a), dir(b-a) {}

    Vec<ld> intersect(Line l) {
        ld t = ld((l.start-start)^l.dir)/(dir^l.dir);
        // For segment-segment intersection this should be in
        range [0, 1]
        Vec<ld> res(start.x, start.y);
```

```
\begin{array}{c} {\rm Vec < ld > \ dirld(dir.x, \ dir.y);} \\ {\rm return \ res} \ + \ dirld*t;} \\ \} \\ \}; \end{array}
```

4.3 Convex Hull Gift Wrapping

```
vector<Vec<int>> buildConvexHull(vector<Vec<int>>&
      pts) {
    int n = pts.size();
   sort(pts.begin(), pts.end());
auto currP = pts[0]; // choose some extreme point to be
           on the hull
    {\tt vector}{<}{\tt Vec}{<}{\tt int}{\gt}{\gt}~{\tt hull};
    \begin{array}{l} \mathrm{set}{<}\mathrm{Vec}{<}\mathrm{int}{>}{>}\ \mathrm{used};\\ \mathrm{hull.pb}(\mathrm{pts}[0]); \end{array}
    used.insert(pts[0]);
    while(true) {
        auto candidate = pts[0]; // choose some point to be a
               candidate
        auto currDir = candidate-currP:
         vector<Vec<int>> toUpdate;
        FOR(i, 0, n) {
             if(currP == pts[i]) continue;
             // currently we have currP->candidate
// we need to find point to the left of this
             auto possibleNext = pts[i];
             auto nextDir = possibleNext - currP;
             auto cross = currDir ^ nextDir;
if(candidate == currP || cross > 0) {
                 candidate = possibleNext;
                  currDir = nextDir;
             } else if(cross == 0 && nextDir.norm() > currDir.
                   norm()) {
                  candidate = possibleNext;
                 currDir = nextDir;
        if(used.find(candidate) != used.end()) break;
        hull.pb(candidate);
        used.insert(candidate);
        currP = candidate;
    return hull;
```

4.4 Convex Hull With Graham's Scan

```
// Takes in >= 3 points
// Returns convex hull in clockwise order
// Ignores points on the border
vector<Vec<int>> buildConvexHull(vector<Vec<int>> pts)
    if(pts.size() <= 3) return pts;
    sort(pts.begin(), pts.end());
    stack < Vec < int >> hull;
    \begin{array}{l} hull.push(pts[0]);\\ auto\ p=pts[0]; \end{array}
    sort(pts.begin()+1, pts.end(), [&](Vec<int> a, Vec<int>
           b) -> bool {
         // p->a->b is a ccw turn
        int turn = sgn((a-p)^(b-a));
        \label{eq:continuous} $$//if(turn == 0) \ return (a-p).norm() > (b-p).norm(); $$// \ among collinear points, take the farthest one
        return turn == 1:
    hull.push(pts[1]);
    FOR(i, 2, (int)pts.size()) {
        auto c = pts[i];
if(c == hull.top()) continue;
        while(true) {
   auto a = hull.top(); hull.pop();
             auto b = hull.top();
             auto ba = a-b;
```

```
auto ac = c-a;
       if((ba^ac) > 0)
          hull.push(a);
          break:
       else if((ba^ac) == 0) 
          if(ba*ac < 0) c = a;
              c is between b and a, so it shouldn't be
                added to the hull
          break:
       }
   hull.push(c);
vector<Vec<int>> hullPts;
while(!hull.empty()) {
   hullPts.pb(hull.top());
   hull.pop();
return hullPts;
```

4.5 Circle Line Intersection

```
double r, a, b, c; // ax+by+c=0, radius is at (0, 0) // If the center is not at (0, 0), fix the constant c to translate everything so that center is at (0, 0) double x0 = a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b); if (c*c > r*r*(a*a+b*b)+eps) puts ("no_points"); else if (abs (c*c - r*r*(a*a+b*b)) < eps) { puts ("1_point"); cout << x0 << '\_' << y0 << '\n'; } else { double d = r*r - c*c/(a*a+b*b); double mult = sqrt (d / (a*a+b*b)); double ax, ay, bx, by; ax = x0 + b * mult; bx = x0 - b * mult; bx = x0 - b * mult; by = y0 - a * mult; puts ("2_points"); cout << ax << '\_' << ay << '\n' << bx << '\_' << by << '\n'; <> cout << ax << '\_' << ay << '\n' << bx << '\_' << by << '\n'; <> by << '\n'; <> cout << ax << '\_' << by << '\n'; <> by << '\n'; <> cout << ax << '\_' << cout << bx << '\_' << cout << cout << ax << '\_' << cout << c
```

4.6 Circle Circle Intersection

Let's say that the first circle is centered at (0,0) (if it's not, we can move the origin to the center of the first circle and adjust the coordinates), and the second one is at (x_2, y_2) . Then, let's construct a line Ax + By + C = 0, where $A = -2x_2, B = -2y_2, C = x_2^2 + y_2^2 + r_1^2 - r_2^2$. Finding the intersection between this line and the first circle will give us the answer. The only tricky case: if both circles are centered at the same point. We handle this case separately.

4.7 Common Tangents To Two Circles

```
struct pt {
    double x, y;

    pt operator- (pt p) {
        pt res = { x-p.x, y-p.y };
        return res;
    }
};
struct circle : pt {
```

```
double r;
struct line {
    double a, b, c;
void tangents (pt c, double r1, double r2, vector<line> & ans)
    double r = r2 - r1;
    double z = sqr(c.x) + sqr(c.y);
double d = z - sqr(r);
    if (d < -eps) return;
    d = sqrt (abs (d));
    l.a = (c.x * r + c.y * d) / z;
    l.b = (c.y * r - c.x * d) / z;
    1.c = r1:
    ans.push_back (1);
vector<line> tangents (circle a, circle b) {
    vector<line> ans;
    for (int i=-1; i<=1; i+=2)
        for (int j=-1; j<=1; j+=2)
tangents (b-a, a.r*i, b.r*j, ans);
    for (size_t i=0; i<ans.size(); ++i)

ans[i].c -= ans[i].a * a.x + ans[i].b * a.y;
```

4.8 Number Of Lattice Points On Segment

Let's say we have a line segment from (x_1, y_1) to (x_2, y_2) . Then, the number of lattice points on this segment is given by

$$gcd(x_2-x_1,y_2-y_1)+1.$$

4.9 Pick's Theorem

We are given a lattice polygon with non-zero area. Let's denote its area by S, the number of points with integer coordinates lying strictly inside the polygon by I and the number of points lying on the sides of the polygon by B. Then:

$$S = I + \frac{B}{2} - 1.$$

5 Math

5.1 Linear Sieve

```
\begin{split} & ll \ minDiv[MAXN+1]; \\ & vector < ll > primes; \\ & void \ sieve(ll \ n) \{ \\ & \ FOR(k, \, 2, \, n+1) \{ \\ & \ minDiv[k] = k; \\ \} \\ & FOR(k, \, 2, \, n+1) \ \{ \\ & \ if(minDiv[k] = k) \ \{ \\ & \ primes.pb(k); \\ \} \\ & \ for(auto \ p : primes) \ \{ \\ & \ if(p > minDiv[k]) \ break; \\ & \ if(p > k > n) \ break; \\ & \ minDiv[p*k] = p; \end{split}
```

```
}
}
```

5.2 Extended Euclidean Algorithm

```
// ax+by=gcd(a,b)
void solve
Eq(ll a, ll b, ll& x, ll& y, ll& g) {
    if(b==0) {
         x = 1;
          y = 0;
          g = a;
          return;
    ll xx, yy;
    solveEq(b, a%b, xx, yy, g);
    x = yy;
    y = xx-yy*(a/b);
// ax + by = c
bool solve
Eq(ll a, ll b, ll c, ll& x, ll& y, ll& g) {
    solveEq(a, b, x, y, g);

if(c\%g != 0) return false;

x *= c/g; y *= c/g;
     return true;
]/ // Finds a solution (x, y) so that x>=0 and x is minimal bool solveEqNonNegX(ll a, ll b, ll c, ll& x, ll &y, ll& g) {
     if(!solveEq(a,\;b,\;c,\;x,\;y,\;g)) \ return \ false;\\
    ll k = x*g/b;

x = x - k*b/g;
    y = y + k*a/g;

if(x < 0) \{

x += b/g;
         y = a/g;
     return true;
```

All other solutions can be found like this:

$$x' = x - k \frac{b}{g}, y' = y + k \frac{a}{g}, k \in \mathbb{Z}$$

5.3 Chinese Remainder Theorem

Let's say we have some numbers m_i , which are all mutually coprime. Also, let $M = \prod_i m_i$. Then the system of congruences

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \dots \\ x \equiv a_k \pmod{m_k} \end{cases}$$

is equivalent to $x \equiv A \pmod{M}$ and there exists a unique number A satisfying $0 \le A \le M$.

Solution for two: $x \equiv a_1 \pmod{m_1}$, $x \equiv a_2 \pmod{m_2}$. Let $x = a_1 + km_1$. Substituting into the second congruence: $km_1 \equiv a_2 - a_1 \pmod{m_2}$. Then, $k = (m_1)_{m_2}^{-1}(a_2 - a_1) \pmod{m_2}$. and we can easily find x. This can be extended to multiple equations by solving them one-by-one.

If the moduli are not coprime, solve the system $y\equiv 0\pmod{\frac{m_1}{g}}, y\equiv \frac{a_2-a_1}{g}\pmod{\frac{m_2}{g}}$ for y. Then let $x\equiv gy+a_1\pmod{\frac{m_1m_2}{g}}$.

5.4 Euler Totient Function

5.5 Factorization With Sieve

```
// Use linear sieve to calculate minDiv
vector<pll> factorize(ll x) {
    vector<pll> res;
    ll prev = -1;
    ll cnt = 0;
    while(x != 1) {
        ll d = minDiv[x];
        if(d == prev) {
            cnt++;
        } else {
            if(prev != -1) res.pb({prev, cnt});
            prev = d;
            cnt = 1;
        }
        x /= d;
    }
    res.pb({prev, cnt});
    return res;
}
```

5.6 Modular Inverse

5.7 Simpson Integration

```
 \begin{array}{l} {\rm const\ int\ N=1000\ ^*\ 1000;\ //\ number\ of\ steps\ (already\ multiplied\ by\ 2)} \\ \\ {\rm double\ simpsonIntegration(double\ a,\ double\ b)\{} \\ {\rm double\ h=(b-a)\ /\ N;} \\ {\rm double\ s=f(a)+f(b);\ //\ a=x\_0\ and\ b=x\_2n} \\ {\rm for\ (int\ i=1;\ i<=N-1;\ ++i)\ \{} \\ {\rm double\ x=a+h\ ^*\ i;} \\ {\rm s+=f(x)\ ^*\ ((i\ \&\ 1)\ ?\ 4:2);} \\ {\rm \}} \\ {\rm s\ ^*=h\ /\ 3;} \\ {\rm return\ s;} \\ \\ \end{array} \right\}
```

5.8 Burnside's Lemma

Let G be a finite group that acts on a set X. For each g in G let X^g denote the set of elements in X that are fixed by g. Burnside's lemma asserts the following formula for the number of orbits:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

Example. Coloring a cube with three colors.

Let X be the set of 3^6 possible face color combinations. Let's count the sizes of the fixed sets for each of the 24 rotations:

- one 0-degree rotation which leaves all 3^6 elements of X unchanged
- six 90-degree face rotations, each of which leaves 3^3 elements of X unchanged
- three 180-degree face rotation, each of which leaves 3^4 elements of X unchanged
- eight 120-degree vertex rotations, each of which leaves 3^2 elements of X unchanged
- six 180-degree edge rotations, each of which leaves 3³ elements of X unchanged

The average is then $\frac{1}{24}(3^6 + 6 \cdot 3^3 + 3 \cdot 3^4 + 8 \cdot 3^2 + 6 \cdot 3^3) = 57$. For n colors: $\frac{1}{24}(n^6 + 3n^4 + 12n^3 + 8n^2)$.

Example. Coloring a circular stripe of n cells with two colors.

X is the set of all colored striped (it has 2^n elements), G is the group of rotations (n elements by 0 cells, by 1 cell, ..., by (n-1) cells). Let's fix some K and find the number of stripes that are fixed by the rotation by K cells. If a stripe becomes itself after rotation by K cells, then its 1st cell must have the same color as its $(1+K \mod n)$ th cell, which is in turn the same as its (1 + 2K) $\mod n$)-th cell, etc., until $mK \mod n = 0$. This will happen when m = n/gcd(K, n). Therefore, we have $n/\gcd(K,n)$ cells that must all be of the same color. The same will happen when starting from the second cell and so on. Therefore, all cells are separated into gcd(K, n) groups, with each group being of one color, and that yields $2^{gcd(K,n)}$ choices. That's why the answer to the original problem is $\frac{1}{n} \sum_{k=0}^{n-1} 2^{\gcd(k,n)}.$

5.9 FFT

```
namespace FFT {
   int n;
   vector<int> r;
   vector<complex<ld>> omega;
```

```
int logN, pwrN;
 void initLogN() {
                 log N = 0;

pwr N = 1;
                   while (pwrN < n) {
                                    pwrN *= 2;
                                     logN++;
                   n = pwrN:
 void initOmega() {
                 FOR(i, 0, pwrN) {
                                   omega[i] = { \cos(2 * i*PI / n), \sin(2 * i*PI / n) };
void initR() {
                 FOR(i, 1, pwrN) {

r[i] = r[i / 2] / 2 + ((i \& 1) << (logN - 1));
 void initArrays() {
                  r.clear();
                  r.resize(pwrN);
                  omega.clear();
                  omega.resize(pwrN);
void init(int n) {
                  FFT::n = n;
                 initLogN():
                  initArrays():
                  initOmega();
                   initR();
\label{eq:complex} \mbox{void } \mbox{ftt(complex<ld>a[], complex<ld>f[]) } \{ \mbox{} \mbox{
                  \begin{array}{c} FOR(i,\,0,\,pwrN) \ \{\\ f[i] = a[r[i]]; \end{array} 
                   for (ll k = 1; k < pwrN; k *= 2) {
                                    \begin{array}{l} \text{(ii } k=1; \ k < pwrin; \ k \ ^{=} 2) \ \{ \\ \text{for (ll } i=0; \ i < pwrin; \ i+=2 \ ^{*} \ k) \ \{ \\ \text{for (ll } j=0; \ j < k; \ j++) \ \{ \\ \text{auto } z = omega[j^{*}n \ / \ (2 \ ^{*} \ k)] \ ^{*} \ f[i+j+k]; \\ f[i+j+k] = f[i+j] - z; \\ f[i+j] +=z; \end{array}
                                 }
               }
```

5.10 FFT With Modulo

```
\begin{array}{l} bool \ is Generator(ll \ g) \ \{\\ if \ (pwr(g, \ M-1) \ != 1) \ return \ false;\\ for \ (ll \ i = 2; \ i^*i < = M-1; \ i++) \ \{ \end{array}
           if ((M - 1) \% i == 0) {
                 ll q = i;
                 if (isPrime(q)) {
                      ll p = (M - 1) / q;
                      ll pp = pwr(g, p);
if (pp == 1) return false;
                 q = (M - 1) / i;
                 if (isPrime(q)) {
                      ll p = (\widetilde{M} - 1) / q;
                      ll pp = pwr(g, p);
                      if (pp == 1) return false;
           }
      return true;
namespace FFT {
     ll n:
      vector<ll> r;
```

```
vector<ll> omega;
ll logN, pwrN;
 {\rm void~initLogN}()~\{
       logN = 0:
       pwrN = 1;
        while (pwrN < n) {
             pwrN *= 2;
             logN++;
       n = pwrN;
}
{\bf void~initOmega()}~\{
       ll g = 2; while (!isGenerator(g)) g++;
       ll G = 1;
        g = pwr(g, (M - 1) / pwrN);
       FOR(i, 0, pwrN) {
    omega[i] = G;
    G *= g;
    G %= M;
       }
}
 void\ initR()\ \{
        \begin{array}{l} r[0] = 0; \\ FOR(i, \, 1, \, pwrN) \ \{ \\ r[i] = r[i \, / \, 2] \ / \ 2 + ((i \, \& \, 1) << (logN \, - \, 1)); \end{array} 
}
{\rm void~initArrays}()~\{
       r.clear();
       r.resize(pwrN);
       omega.clear();
       omega.resize(pwrN);
 void init(ll n) {
       FFT::n = n;
initLogN();
       initArrays();
       initOmega();
       initR();
 \begin{array}{c} {\rm void\ fft}(ll\ a[],\ ll\ f[])\ \{ \\ {\rm \ for\ } (ll\ i\ =\ 0;\ i\ <\ pwrN;\ i++)\ \{ \end{array} 
             \dot{f}[i] = a[r[i]];
      for (ll k = 1; k < pwrN; k *= 2) {
  for (ll i = 0; i < pwrN; i += 2 * k) {
    for (ll j = 0; j < k; j++) {
      auto z = omega[j*n / (2 * k)] * f[i + j + k]
      ^{\circ} M.
                                     % M;
                           f[i+j+k] = f[i+j] - z;
                          \begin{array}{l} f[i+j+k] = f[i+j] - 2, \\ f[i+j] + = z; \\ f[i+j+k] \% = M; \\ \text{if } (f[i+j+k] < 0) \ f[i+j+k] + = M; \\ f[i+j] \% = M; \end{array}
     } }
}
```

$\begin{array}{ccc} 5.11 & Big & Integer & Multiplication \\ & With \ FFT \end{array}$

}

```
 \begin{array}{l} complex < ld > a[MAX\_N], \ b[MAX\_N]; \\ complex < ld > fa[MAX\_N], \ fb[MAX\_N], \ fc[MAX\_N]; \\ complex < ld > cc[MAX\_N]; \\ string \ mul(string \ as, \ string \ bs) \ \{ \\ int \ sgn1 = 1; \\ int \ sgn2 = 1; \\ if \ (as[0] == `-') \ \{ \\ sgn1 = -1; \\ as = as.substr(1); \\ \} \\ if \ (bs[0] == `-') \ \{ \end{array}
```

```
bs = bs.substr(1);
int n = as.length() + bs.length() + 1;
FFT::init(n);
FOR(i, 0, FFT::pwrN) {
    a[i] = b[i] = fa[i] = fb[i] = fc[i] = cc[i] = 0;
FOR(i, 0, as.size()) {
    a[i] = as[as.size() - 1 - i] - '0';
FOR(i, 0, bs.size()) {
    b[i] = bs[bs.size() - 1 - i] - '0';
FFT::fft(a, fa);
FFT::fft(b, fb);
FOR(i, 0, FFT::pwrN) {
fc[i] = fa[i] * fb[i];
// turn [0,1,2,...,n-1] into [0, n-1, n-2, ..., 1]

FOR(i, 1, FFT::pwrN) {

    if (i < FFT::pwrN - i) {
        swap(fc[i], fc[FFT::pwrN - i]);
    }
FFT::fft(fc, cc);
ll carry = 0;
vector<int> v;
FOR(i, 0, FFT::pwrN)
    int num = round(cc[i].real() / FFT::pwrN) + carry;
    v.pb(num % 10);
    carry = num / 10;
while (carry > 0) {
    v.pb(carry % 10);
    carry /= 10;
reverse(v.begin(), v.end());
bool start = false;
ostringstream ss;
bool allZero = true;
for (auto x : v) {
    if (x != 0) {
         allZero = false;
         break;
    }
if (sgn1*sgn2 < 0 \&\& !allZero) ss << "-"; for (auto x : v) {
    if (x == 0 &&!start) continue;
    start = true;
    ss \ll abs(x);
if (!start) ss << 0;
return ss.str();
```

5.12 Formulas

```
\sum_{i=1}^{n} i = \frac{n(n+1)}{2}; \quad \sum_{i=1}^{n} i^2 = \frac{n(2n+1)(n+1)}{6};
\sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4}; \quad \sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30};
\sum_{i=a}^{b} c^i = \frac{c^{b+1}-c^a}{c-1}, c \neq 1; \quad \sum_{i=1}^{n} a_1 + (i-1)d = \frac{n(a_1+a_n)}{2}; \quad \sum_{i=1}^{n} a_1r^{i-1} = \frac{a_1(1-r^n)}{1-r}, r \neq 1;
\sum_{i=1}^{\infty} ar^{i-1} = \frac{a_1}{1-r}, |r| \leq 1.
```

6 Strings

6.1 Hashing

```
init(\underline{\hspace{1em}}s);
          calcHashes(_s);
     void init(const string& s) {
          ll \ a1 = 1:
          11 \text{ a} 2 = 1;
          FOR(i,\,0,\,(int)s.length(){+}1)~\{
               A1pwrs.pb(a1);
               A2pwrs.pb(a2);

a1 = (a1*A1)\%B1;
               a2 = (a2*A2)\%B2;
          }
     void calcHashes(const string& s) {
          pll h = \{0, 0\};
prefixHash.pb(h);
          for(char c : s) {

ll h1 = (prefixHash.back().first*A1 + c)%B1;
               ll h2 = (prefixHash.back().second*A2 + c)\%B2;
               prefixHash.pb({h1, h2});
      \begin{array}{l} \text{pll getHash(int l, int r) } \\ \text{ll h1} = (\text{prefixHash[r+1].first - prefixHash[l].first*} \\ \text{A1pwrs[r+1-l])} \% \text{ B1;} \\ \end{array} 
          ll h2 = (prefixHash[r+1].second - prefixHash[l].second*
                  A2pwrs[r+1-l]) % B2;
          if(h1 < 0) h1 += B1;

if(h2 < 0) h2 += B2;
          return {h1, h2};
};
```

6.2 Prefix Function

```
\label{eq:continuous_problem} // \ pi[i] \ is the length of the longest proper prefix of the substring s[0..i] which is also a suffix <math display="block">\ // \ of \ this \ substring \ vector<int> \ prefixFunction(const string& s) \ \{ \ int \ n = (int)s.length(); \ vector<int> \ pi(n); \ for \ (int \ i = 1; \ i < n; \ i++) \ \{ \ int \ j = pi[i-1]; \ while \ (j>0 \&\& \ s[i] \ != \ s[j]) \ j = pi[j-1]; \ if \ (s[i] = = \ s[j]) \ j++; \ pi[i] = \ j; \ \} \ return \ pi; \ \}
```

6.3 Prefix Function Automaton

6.4 KMP

```
\label{eq:cont_string} // \ Knuth-Morris-Pratt algorithm \\ vector<int> \ findOccurences(const string& s, const string& t) \\ \left\{ \begin{array}{l} int \ n = s.length(); \\ int \ m = t.length(); \\ string \ S = s + "\#" + t; \\ auto \ pi = prefixFunction(S); \\ vector<int> \ ans; \\ FOR(i, n+1, n+m+1) \ \left\{ \begin{array}{l} if(pi[i] = n) \ \left\{ \\ ans.pb(i-2*n); \\ \right\} \\ return \ ans; \\ \end{array} \right\}
```

6.5 Aho Corasick Automaton

```
// alphabet size
const int K = 70:
// the indices of each letter of the alphabet
int intVal[256];
void init() {
   int curr = 2;
   intVal[1] = 1;
   for(char c = '0'; c <= '9'; c++, curr++) intVal[(int)c] =
         curr;
   for(char\ c='A';\ c<='Z';\ c++,\ curr++)\ intVal[(int)c]=
         curr;
   for(char\ c='a';\ c<='z';\ c++,\ curr++)\ intVal[(int)c]=
}
struct Vertex {
   int next[K];
    vector<int> marks;
   // ^ this can be changed to int mark = -1, if there will be
         no duplicates
   int p = -1;
   char pch;
   int link = -1;
   int exitLink = -1;
        exitLink points to the next node on the path of suffix
          links which is marked
   int go[K];
      ch has to be some small char
    Vertex(int _p=-1, char ch=(char)1) : p(_p), pch(ch) {
       fill(begin(next), end(next), -1);
       fill(begin(go), end(go), -1);
vector < Vertex > t(1);
void addString(string const& s, int id) {
   int v = 0;
   for (char ch : s) {
       int c = intVal[(int)ch];
       if (t[v].next[c] == -1) {

t[v].next[c] = t.size();
           t.emplace_back(v, ch);
       v = t[v].next[c];
   t[v].marks.pb(id);
```

```
int go(int v, char ch);
int\ getLink(int\ v)\ \{
   if (t[v].link == -1) {
if (v == 0 || t[v].p == 0)
            t[v].link = 0;
            t[v].link = go(getLink(t[v].p), \ t[v].pch);\\
   return t[v].link;
}
int getExitLink(int v) {
    if(t[v].exitLink != -1) return t[v].exitLink;
   int l = getLink(v);

if(l == 0) return t[v].exitLink = 0;

if(lt[l].marks.empty()) return t[v].exitLink = l;
   return t[v].exitLink = getExitLink(l);
int go(int v, char ch) {
   int c = intVal[(int)ch];

if (t[v].go[c] == -1) {

    if (t[v].next[c] != -1)

        t[v].go[c] = t[v].next[c];
            t[v].go[c] = v == 0 ? 0 : go(getLink(v), \, ch); \\
   return t[v].go[c];
}
void walk
Up(int v, vector<int>& matches) {
    if(v == 0) return;
   if(!t[v].marks.empty())
        for(auto m : t[v].marks) matches.pb(m);
    walkUp(getExitLink(v), matches);
}
   returns the IDs of matched strings.
   Will contain duplicates if multiple matches of the same
      string are found.
vector<int> walk(const string& s) {
    vector<int> matches;
    int curr = 0;
    for(char c : s) {
        curr = go(curr, c);
        if(!t[curr].marks.empty())  {
            for(auto m: t[curr].marks) matches.pb(m);
        walkUp(getExitLink(curr), matches);
   return matches;
/* Usage:
  * addString(strs[i], i);
 * auto matches = walk(text);
 \ast .. do what you need with the matches - count, check if
       some id exists, etc ..
 * Some applications:
  - Find all matches: just use the walk function
 * - Find lexicographically smallest string of a given length
       that doesn't match any of the given strings:
 * For each node, check if it produces any matches (it either
       contains some marks or walkUp(v) returns some marks)
 * Remove all nodes which produce at least one match. Do
       DFS in the remaining graph, since none of the
       remaining nodes
 * will ever produce a match and so they're safe.
  - Find shortest string containing all given strings:
* For each vertex store a mask that denotes the strings which
 match at this state. Start at (v = root, mask = 0), * we need to reach a state (v, mask=2^n-1), where n is the
       number of strings in the set. Use BFS to transition
       between states
 * and update the mask
```

6.6 Suffix Array

```
int n = s.size();
     const int alphabet = 256; // we assume to use the whole
             ASCII range
     \label{eq:continuous} \begin{array}{l} vector < int > p(n), \; c(n), \; cnt(max(alphabet, \; n), \; 0); \\ for \; (int \; \underline{i} \; \underline{=} \; 0; \; i < n; \; i++) \end{array}
     cnt[s[i]]++;
for (int i = 1; i < alphabet; i++)
     cnt[i] += cnt[i-1];
for (int i = 0; i < n; i++)
     p[-cnt[s[i]]] = i;

c[p[0]] = 0;
     int classes = 1;
     for (int i = 1; i < n; i++) {
          if (s[p[i]] != s[p[i-1]])
               classes++;
          c[p[i]] = classes - 1;
     vector < int > pn(n), cn(n);
     for (int h = 0; (1 << h) < n; ++h) {
for (int i = 0; i < n; i++) {
               pn[i] = p[i] - (1 << h);
if (pn[i] < 0)
                     pn[i] += n;
          fill(cnt.begin(), cnt.begin() + classes, 0);
          for (int i = 0; i < n; i++)
          cnt[c[pn[i]]]++;

for (int i = 1; i < classes; i++)

cnt[i] += cnt[i-1];

for (int i = n-1; i >= 0; i--)
               p[--cnt[c[pn[i]]]] = pn[i];
          \operatorname{cn}[p[0]] = 0;
          classes = 1;
          \begin{array}{ll} {\rm for} \ ({\rm int} \ i = 1; \ i < n; \ i++) \ \{ \\ {\rm pair}<{\rm int}, \ {\rm int}> \ {\rm cur} = \{ c[p[i]], \ c[(p[i] + (1 << h)) \ \% \end{array}
               pair < int, int > prev = \{c[p[i-1]], c[(p[i-1] + (1 <<
                       h)) % n]};
               if (cur != prev)
                      ++classes;
               cn[p[i]] = classes - 1;
          c.swap(cn);
     return p;
vector<int> constructSuffixArray(string s) {
     s += "\$"; // <- this must be smaller than any character
            in s
     vector<int> sorted_shifts = sortCyclicShifts(s);
     sorted_shifts.erase(sorted_shifts.begin());
     return sorted_shifts;
```