ACM-ICPC TEAM REFERENCE DOCUMENT

Vilnius University (Šimoliūnaitė, Strakšys, Strimaitis)

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1 Data Structures						
1.1 Disjoin Set Union						

```
struct DSU {
   vector<int> par;
   vector<int> sz;
   DSU(int n) {
      FOR(i, 0, n) {
          par.pb(i);
          sz.pb(1);
   int find(int a) {
      return par[a] = par[a] == a ? a : find(par[a])
   bool same(int a, int b) {
      return find(a) == find(b);
   void unite(int a, int b) {
      a = find(a);
      b = find(b);
      if(sz[a] > sz[b]) swap(a, b);
      sz[b] += sz[a];
      par[a] = b;
```

1.2 Fenwick 2D

```
struct Fenwick2D {
    vector<vector<ll>> bit;
   int n, m;
    Fenwick2D(int n, int m) {
       n = _n; m = _m;
       \text{bit} = \text{vector} < \text{vector} < \text{ll} >> (n+1, \text{vector} < \text{ll})
              >(m+1, 0);
    fl sum(int x, int y) {
       ll ret = 0;
       for (int i = x; i > 0; i -= i & (-i))
           for (int j = y; j > 0; j = j \& (-j))
               ret += bit[i][j];
       return ret;
    ll sum(int x1, int y1, int x2, int y2) {
       return sum(x2, y2) - sum(x2, y1-1) - sum(
              x1-1, y2) + sum(x1-1, y1-1);
```

```
} void add(int x, int y, ll delta) { for (int i = x; i <= n; i += i & (-i)) for (int j = y; j <= m; j += j & (-j)) bit[i][j] += delta; } } }
```

1.3 Fenwick Tree Point Update And Range Query

```
struct Fenwick {
    vector<ll> tree;
    Fenwick(){}
    Fenwick(int _n) {
        n = _n;
        tree = vector < ll > (n+1, 0);
    void add(int i, ll val) { // arr[i] += val
        for(; i \le n; i += i\&(-i)) tree[i] += val;
    ll get(int i) { // arr[i]
        return sum(i, i);
    Îl sum(int i) { // arr[1]+...+arr[i]
        ll ans = 0;
        for(; i > 0; i -= i\&(-i)) ans += tree[i];
        return ans:
    \hat{l}l \text{ sum}(\text{int } l, \text{ int } r)  {// \text{arr}[l]+...+\text{arr}[r]
        return sum(r) - sum(l-1);
```

1.4 Fenwick Tree Range Update And Point Query

```
struct Fenwick {
   vector<ll> tree;
   vector<ll> arr;
   int n;
   Fenwick(vector<ll> _arr) {
```

```
\begin{array}{l} n = \_arr.size(); \\ arr = \_arr; \\ tree = vector < ll > (n+2, 0); \\ \} \\ void add(int i, ll val) \left\{ \ // \ arr[i] \ += \ val \\ for(; i <= n; i += i\&(-i)) \ tree[i] \ += \ val; \\ \} \\ void add(int l, int r, ll val) \left\{ \ // \ arr[l..r] \ += \ val \\ add(l, val); \\ add(r+1, -val); \\ \} \\ ll \ get(int i) \left\{ \ // \ arr[i] \\ ll \ sum = \ arr[i-1]; \ // \ zero \ based \\ for(; i > 0; i -= i\&(-i)) \ sum \ += \ tree[i]; \\ return \ sum; \ // \ zero \ based \\ \} \\ \}; \end{array}
```

1.5 Fenwick Tree Range Update And Range Query

1.6 Implicit Treap

```
template <typename T>
struct Node {
   Node* l, *r;
   ll prio, size, sum;
   T val;
   bool rev;
   Node() {}
   Node(T val): l(nullptr), r(nullptr), val(val),
         size(1), sum( val), rev(false) {
      prio = rand() \cap (rand() << 15);
template <typename T>
struct ImplicitTreap {
   typedef Node<T>* NodePtr:
   int sz(NodePtr n) {
      return n ? n->size : 0;
   il getSum(NodePtr n) {
      return n ? n->sum : 0;
   void push(NodePtr n) {
      if (n && n->rev) {
          n->rev = false;
          swap(n->l, n->r);
          if (n->l) n->l->rev = 1;
          if (n->r) n->r->rev = 1;
   void recalc(NodePtr n) {
      if (!n) return;
      n-> size = sz(n->l) + 1 + sz(n->r);
      n->sum = getSum(n->l) + n->val +
             getSum(n->r):
   void split(NodePtr tree, ll key, NodePtr& l,
         NodePtr& r) {
       push(tree);
      if (!tree) {
          l = r = nullptr;
      else if (\text{key} \leq \text{sz}(\text{tree-}>l))
          split(tree->l, key, l, tree->l);
          r = tree;
      else
          split(tree->r, key-sz(tree->l)-1, tree->r,
          l = tree;
```

```
recalc(tree);
void merge(NodePtr& tree, NodePtr l, NodePtr
   push(l); push(r);
   if (!l || !r) {
       tree = 1?1:r;
   else if (l->prio > r->prio) {
       merge(l->r, l->r, r);
       tree = 1;
   else {
       merge(r->l, l, r->l);
       tree = r;
   recalc(tree);
void insert(NodePtr& tree, T val, int pos) {
   if (!tree) {
       tree = new Node < T > (val);
       return;
   NodePtr L, R;
   split(tree, pos, L, R);
   merge(L, L, new Node<T>(val));
   merge(tree, L, R);
   recalc(tree);
void reverse(NodePtr tree, int l, int r) {
   NodePtr t1, t2, t3;
   split(tree, l, t1, t2);
   split(t2, r - l + 1, t2, t3);
   if(t2) t2 - rev = true;
   merge(t2, t1, t2);
   merge(tree, t2, t3);
void print(NodePtr t, bool newline = true) {
   push(t);
   if (!t) return;
   print(t->l, false);
   cout << t->val << " ";
   print(t->r, false);
   if (newline) cout << endl;
NodePtr fromArray(vector<T> v) {
   NodePtr t = nullptr:
   FOR(i, 0, (int)v.size()) {
```

```
insert(t, v[i], i);
}
return t;
}

ll calcSum(NodePtr t, int l, int r) {
   NodePtr L, R;
   split(t, l, L, R);
   NodePtr good;
   split(R, r - 1 + 1, good, L);
   return getSum(good);
}
};
/* Usage: ImplicitTreap<int> t;
Node<int> tree = t.fromArray(someVector); t.
   reverse(tree, l, r); ...
*/
```

1.7 Segment Tree With Lazy Propagation

```
// Add to segment, get maximum of segment
struct LazySegTree {
   int n;
    vector<ll> t, lazy;
   LazySegTree(int _n) {
       n = n; t = vector < ll > (4*n, 0); lazy =
             vector < ll > (4*n, 0);
    LazySegTree(vector<ll>& arr) {
       n = n; t = \text{vector} < \text{ll} > (4*n, 0); lazy =
             \text{vector} < \text{ll} > (4*n, 0);
       build(arr, 1, 0, n-1); // same as in simple
             SegmentTree
    void push(int v) {
       t[v*2] += lazy[v];
       lazy[v*2] += lazy[v];
       t[v^*2+1] += lazy[v];
       lazy[v*2+1] += lazy[v];
       lazy[v] = 0;
   void update(int v, int tl, int tr, int l, int r, ll
         addend) {
       if (l > r)
           return;
       if (l == tl \&\& tr == r) {
           t[v] += addend;
```

```
lazy[v] += addend;
       } else -
          push(v);
          int tm = (tl + tr) / 2;
          update(v*2, tl, tm, l, min(r, tm),
                addend);
          update(v*2+1, tm+1, tr, max(l, tm+1),
                 r, addend);
          t[v] = \max(t[v*2], t[v*2+1]);
   int query(int v, int tl, int tr, int l, int r) {
       if (l > r)
          return -OO:
       if (tl == tr)
          return t[v];
       push(v);
      int tm = (tl + tr) / 2;
      return max(query(v*2, tl, tm, l, min(r, tm))
              query(v^*2+1, tm+1, tr, max(l, tm))
                    +1), r));
};
```

1.8 Segment Tree

```
struct SegmentTree {
   int n;
   vector < ll > t;
   const ll IDENTITY = 0; // OO for min, -OO
         for max, ...
   ll f(ll a, ll b) {
       return a+b;
   SegmentTree(int _n) {
      n = _n; t = vector < ll > (4*n, IDENTITY);
   SegmentTree(vector<ll>& arr) {
      n = arr.size(); t = vector < ll > (4*n,
            IDENTITY);
       build(arr, 1, 0, n-1);
   void build(vector<ll>& arr, int v, int tl, int tr)
       if(tl == tr) \{ t[v] = arr[tl]; \}
      else {
          int tm = (tl+tr)/2;
```

```
build(arr, 2*v, tl, tm);
        build(arr, 2*v+1, tm+1, tr);
        t[v] = f(t[2*v], t[2*v+1]);
// sum(1, 0, n-1, l, r)
ll sum(int v, int tl, int tr, int l, int r) {
    if(l > r) return IDENTITY;
    if (l == tl \&\& r == tr) return t[v];
    int tm = (tl+tr)/2;
    \mathrm{return}\ f(\mathrm{sum}(2^*\mathrm{v},\,\mathrm{tl},\,\mathrm{tm},\,\mathrm{l},\,\mathrm{min}(\mathrm{r},\,\mathrm{tm})),
          sum(2*v+1, tm+1, tr, max(l, tm+1),
           r));
// update(1, 0, n-1, i, v)
void update(int v, int tl, int tr, int pos, ll
      newVal) {
    if(tl == tr) \{ t[v] = newVal; \}
    else {
        int tm = (tl+tr)/2;
        if(pos <= tm) update(2*v, tl, tm, pos,
               newVal);
        else update(2*v+1, tm+1, tr, pos,
               newVal);
        t[v] = f(t[2*v],t[2*v+1]);
```

1.9 Treap

```
if (!n) return;
   n->size = sz(n->l) + 1 + sz(n->r); // add
         more operations here as needed
void split(NodePtr tree, ll key, NodePtr& l,
     NodePtr& r) {
   if (!tree) {
      l = r = nullptr;
   else if (key < tree->key) {
      split(tree->l, key, l, tree->l);
      r = tree;
      split(tree->r, key, tree->r, r);
      l = tree;
   recalc(tree);
void merge(NodePtr& tree, NodePtr l, NodePtr
   if (!l || !r) {
      tree = 1?1:r;
   else if (l->prio > r->prio) {
      merge(l->r, l->r, r);
      tree = 1;
   else {
      merge(r->l, l, r->l);
      tree = r;
   recalc(tree);
void insert(NodePtr& tree, NodePtr node) {
   if (!tree) {
      tree = node;
   else if (node->prio > tree->prio) {
      split(tree, node->key, node->l, node->r
            );
      tree = node;
   élse {
      insert(node->key < tree->key ? tree->l
            : tree->r, node);
   recalc(tree);
```

1.10 Trie

```
struct Trie {
   const int ALPHA = 26;
   const char BASE = 'a';
   vector<vector<int>> nextNode;
   vector<int> mark;
   int nodeCount:
   Trie() {
      nextNode = vector<vector<int>>(MAXN,
            vector<int>(ALPHA, -1));
      mark = vector < int > (MAXN, -1);
      nodeCount = 1;
   void insert(const string& s, int id) {
      int curr = 0;
      FOR(i, 0, (int)s.length()) {
          int c = s[i] - BASE;
          if(nextNode[curr][c] == -1) {
             nextNode[curr][c] = nodeCount++;
          curr = nextNode[curr][c];
      mark[curr] = id;
   bool exists(const string& s) {
      int curr = 0;
```

```
 \begin{array}{c} FOR(i,\,0,\,(int)s.length())\;\{\\ &int\;c=s[i]-BASE;\\ &if(nextNode[curr][c]==-1)\;return\;false;\\ &curr=nextNode[curr][c];\\ \\ \}\\ &return\;mark[curr]\;!=-1;\\ \\ \}; \end{array}
```

2 General

2.1 Automatic Test

```
# Linux Bash
# gen, main and stupid have to be compiled
     beforehand
for((i=1;;++i)); do
   echo $i;
   ./gen $i > genIn;
   diff <(./main < genIn) <(./stupid < genIn) ||
         break:
done
# Windows CMD
FOR /L %%I IN (1,1,2147483647) DO (
   echo %%I
   gen.exe %%I > genIn
   {\rm main.exe} < {\rm genIn} > {\rm mainOut}
   stupid.exe < genIn > stupidOut
   FC mainOut stupidOut || goto :eof
```

2.2 Big Integer

```
const int base = 1000000000;
const int base_digits = 9;
struct bigint {
   vector<int> a;
   int sign;
   int size() {
      if (a.empty()) return 0;
      int ans = (a.size() - 1) * base_digits;
```

```
int ca = a.back();
   while (ca) ans++, ca \neq 10;
   return ans;
bigint operator (const bigint &v) {
   bigint ans = 1, x = *this, y = v;
   while (!v.isZero()) {
       if (y \% 2) ans *= x;
       x *= x, y /= 2;
   return ans;
string to_string() {
   stringstream ss;
   ss << *this;
   string s;
   ss >> s;
   return s;
int sumof() {
   string s = to\_string();
   int ans = 0;
   for (auto c : s) ans += c - 0;
   return ans;
bigint() : sign(1) \{ \}
bigint(long long v) {
   *this = v;
bigint(const string &s) {
   read(s);
void operator=(const bigint &v) {
   sign = v.sign;
   a = v.a;
void operator=(long long v) {
   sign = 1;
   a.clear();
   if (v < 0)
       sign = -1, v = -v;
   for (; v > 0; v = v / base)
       a.push back(v % base);
bigint operator+(const bigint &v) const {
   if (sign == v.sign) {
       bigint res = v;
       for (int i = 0, carry = 0; i < (int)max(a.
             size(), v.a.size()) || carry; ++i) {
           if (i == (int)res.a.size()) res.a.
                push back(0);
           res.a[i] += carry + (i < (int)a.size())
```

? a[i] : 0);

```
carry = res.a[i] >= base;
           if (carry) res.a[i] -= base;
       return res;
   return *this - (-v);
bigint operator-(const bigint &v) const {
   if (sign == v.sign) {
       if (abs() \ge v.abs()) {
           bigint res = *this;
           for (int i = 0, carry = 0; i < (int)v.a
                 .size() \mid\mid carry; ++i) {
              res.a[i] -= carry + (i < (int)v.a.
                    size() ? v.a[i] : 0);
              carry = res.a[i] < 0;
              if (carry) res.a[i] += base;
           res.trim();
           return res;
       return -(v - *this);
   return *this + (-v);
void operator*=(int v) {
   if (v < 0) sign = -sign, v = -v;
   for (int i = 0, carry = 0; i < (int)a.size() ||
         carry; ++i) {
       if (i == (int)a.size()) a.push back(0);
       long long cur = a[i] * (long long)v +
             carry;
       carry = (int)(cur / base);
       a[i] = (int)(cur \% base);
   trim();
bigint operator*(int v) const {
   bigint res = *this;
   res *= v;
   return res;
void operator*=(long long v) {
   if (v < 0) sign = -sign, v = -v;
   for (int i = 0, carry = 0; i < (int)a.size() ||
         carry; ++i) {
       if (i == (int)a.size()) a.push_back(0);
       long long cur = a[i] * (long long)v +
             carry;
       carry = (int)(cur / base);
       a[i] = (int)(cur \% base);
   trim();
```

```
bigint operator*(long long v) const {
   bigint res = *this;
   res *= v;
   return res;
friend pair < bigint, bigint > divmod(const bigint
      &a1, const bigint &b1) {
   int norm = base / (b1.a.back() + 1);
   bigint a = a1.abs() * norm;
   bigint b = b1.abs() * norm;
   bigint q, r;
   q.a.resize(a.a.size());
   for (int i = a.a.size() - 1; i >= 0; i--) {
       r *= base;
       r += a.a[i];
       int s1 = r.a.size() \le b.a.size() ? 0 : r.a
             [b.a.size()];
       int s2 = r.a.size() \le b.a.size() - 1?0:
             r.a[b.a.size() - 1];
       int d = ((long long)base * s1 + s2) / b.
             a.back();
       r -= b * d;
       while (r < 0) r += b, --d;
       q.a[i] = d;
   q.sign = a1.sign * b1.sign;
   r.sign = a1.sign;
   q.trim();
   r.trim();
   return make_pair(q, r / norm);
bigint operator/(const bigint &v) const {
   return divmod(*this, v).first;
bigint operator%(const bigint &v) const {
   return divmod(*this, v).second;
void operator/=(int v) {
   if (v < 0) sign = -sign, v = -v;
   for (int i = (int)a.size() - 1, rem = 0; i >=
         0; --i) {
       long long cur = a[i] + rem * (long long)
             base:
       a[i] = (int)(cur / v);
      rem = (int)(cur \% v);
   trim();
bigint operator/(int v) const {
   bigint res = *this;
   res /= v;
   return res;
```

```
int operator%(int v) const {
   if (v < 0) v = -v;
   int m = 0;
   for (int i = a.size() - 1; i >= 0; --i)
      m = (a[i] + m * (long long)base) % v;
   return m * sign;
void operator+=(const bigint &v) {
   *this = *this + v;
void operator-=(const bigint &v) {
   *this = *this - v;
void operator*=(const bigint &v) {
   *this = *this * v;
void operator/=(const bigint &v) {
   *this = *this / v;
bool operator<(const bigint &v) const {
   if (sign != v.sign) return sign < v.sign;
   if (a.size() != v.a.size())
       return a.size() * sign < v.a.size() * v.
             sign;
   for (int i = a.size() - 1; i >= 0; i--)
      if (a[i] != v.a[i])
return a[i] * sign < v.a[i] * sign;
   return false;
bool operator>(const bigint &v) const {
   return v < *this;
bool operator <= (const bigint &v) const {
   return !(v < *this);
bool operator>=(const bigint &v) const {
   return !(*this < v);
bool operator==(const bigint &v) const {
   return !(*this < v) && !(v < *this);
bool operator!=(const bigint &v) const {
   return *this < v \mid \mid v < *this;
void trim() {
   while (!a.empty() && !a.back()) a.
         pop back();
   if (a.empty()) sign = 1;
bool isZero() const {
   return a.empty() || (a.size() == 1 \&\& !a[0])
```

```
bigint operator-() const {
   bigint res = *this;
   res.sign = -sign;
   return res;
bigint abs() const {
   bigint res = *this;
   res.sign *= res.sign;
   return res;
long long longValue() const {
   long long res = 0;
   for (int i = a.size() - 1; i >= 0; i--) res =
         res * base + a[i];
   return res * sign;
friend bigint gcd(const bigint &a, const bigint
   return b.isZero()? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint
   return a / gcd(a, b) * b;
void read(const string &s) {
   sign = 1;
   a.clear();
   int pos = 0;
   while (pos < (int)s.size() && (s[pos] == '-'
         || s[pos] == '+')) {
       if (s[pos] == '-') sign = -sign;
       ++pos;
   for (int i = s.size() - 1; i >= pos; i -=
         base digits) {
       int x = 0;
       for (int j = max(pos, i - base digits +
            1); i \le i; i++)
           x = x * 10 + s[i] - 0;
       a.push back(x);
   trim();
friend istream & operator >> (istream & stream,
     bigint &v) {
   string s;
   stream >> s;
   v.read(s);
   return stream;
friend ostream & operator << (ostream & stream,
      const bigint &v) {
```

```
if (v.sign == -1) stream << '-';
   stream << (v.a.empty() ? 0 : v.a.back());
   for (int i = (int)v.a.size() - 2; i >= 0; --i)
       stream << setw(base_digits) << setfill
             ('0') << v.a[i];
   return stream;
static vector<int> convert base(const vector<
     int> &a, int old digits, int new digits) {
   vector<long long> p(max(old_digits,
         new digits) +1;
   p[0] = 1;
   for (int i = 1; i < (int)p.size(); i++)
       p[i] = p[i - 1] * 10;
   vector<int> res;
   long long cur = 0;
   int cur digits = 0;
   for (int i = 0; i < (int)a.size(); i++) {
       cur += a[i] * p[cur digits];
       cur_digits += old_digits;
       while (cur_digits >= new_digits) {
           res.push back(int(cur % p[
                new_digits]));
          cur /= p[new digits];
          cur digits -= new digits;
   res.push back((int)cur);
   while (!res.empty() && !res.back()) res.
         pop back();
   return res;
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a, const
      vll &b) {
   int n = a.size():
   vll res(n + n);
   if (n <= 32) {
       for (int i = 0; i < n; i++)
          for (int j = 0; j < n; j++)
              res[i + j] += a[i] * b[j];
       return res;
   int k = n >> 1;
   vll a1(a.begin(), a.begin() + k);
   vll a2(a.begin() + k, a.end());
   vll b1(b.begin(), b.begin() + k);
   vll b2(b.begin() + k, b.end());
   vll a1b1 = karatsubaMultiply(a1, b1);
   vll a2b2 = karatsubaMultiply(a2, b2);
   for (int i = 0; i < k; i++) a2[i] += a1[i];
```

```
for (int i = 0; i < k; i++) b2[i] += b1[i];
       vll r = karatsubaMultiply(a2, b2);
       for (int i = 0; i < (int)a1b1.size(); i++) r[i]
             -= a1b1[i];
       for (int i = 0; i < (int)a2b2.size(); i++) r[i]
             -= a2b2[i];
       for (int i = 0; i < (int)r.size(); i++) res[i +
             k] += r[i];
       for (int i = 0; i < (int)a1b1.size(); i++) res
             i] += a1b1[i];
       for (int i = 0; i < (int)a2b2.size(); i++) res
            i + n] += a2b2[i];
       return res;
   bigint operator*(const bigint &v) const {
       vector<int> a6 = convert base(this->a,
             base digits, 6);
       vector<int> b6 = convert base(v.a,
             base_digits, 6);
       vll x(a6.begin(), a6.end());
       vll y(b6.begin(), b6.end());
       while (x.size() < y.size()) x.push back(0);
       while (y.size() < x.size()) y.push_back(0);
       while (x.size() & (x.size() - 1)) x.push back
             (0), y.push back(0);
       vll c = karatsubaMultiplv(x, y);
       res.sign = sign * v.sign;
       for (int i = 0, carry = 0; i < (int)c.size(); i
             ++) {
          long long cur = c[i] + carry;
          res.a.push_back((int)(cur % 1000000));
          carry = (int)(cur / 1000000);
       res.a = convert base(res.a, 6, base digits);
       res.trim();
       return res;
};
```

2.3 C++ Template

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp> //
gp_hash_table<int, int> == hash map
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
```

```
using namespace gnu pbds;
typedef long long ll;
typedef unsigned long long ull;
typedef long double ld;
typedef pair<int, int> pii;
typedef pair<ll, ll> pll;
typedef pair < double, double > pdd;
template <typename T> using min_heap =
     priority queue<T, vector<T>, greater<T
template <typename T> using max heap =
     priority_queue<T, vector<T>, less<T>>;
template <typename T> using ordered_set = tree
     <T, null type, less<T>, rb tree tag,
     tree_order_statistics_node_update>;
template <typename K, typename V> using
     hashmap = gp hash table < K, V > ;
template<typename A, typename B> ostream&
     operator << (ostream & out, pair < A, B > p) {
     out << "(" << p.first << ", " << p.second
     << ")"; return out;}
template<typename T> ostream& operator<<(
     ostream& out, vector\langle T \rangle v) { out \langle v \rangle out \langle v \rangle
     for(auto\& x : v) out << x << ", "; out <<
      "]":return out;}
template<typename T> ostream& operator<<(
     ostream& out, setT > v) { out T < v} for(
     auto& x : v) out << x << ", "; out <math><< "}";
      return out; }
template<typename K, typename V> ostream&
     operator << (ostream & out, map < K, V > m)
     { out << "{"; for(auto& e: m) out << e. first << "-> " << e.second << ", "; out
     << "}"; return out; }
template<typename K, typename V> ostream&
     operator << (ostream & out, hashmap < K, V >
     m) { out << "{"; for(auto& e : m) out <<
     e.first << " -> " << e.second << ", "; out
     << "}"; return out; }
#define FAST_IO ios_base::sync_with_stdio(
     false); cin.tie(NULL)
#define TESTS(t) int NUMBER OF TESTS; cin
      >> NUMBER_OF_TESTS; for(int t = 1; t
      <= NUMBER_OF_TESTS; t++)
#define FOR(i, begin, end) for (int i = (begin) - ((
     begin) > (end); i != (end) - ((begin) > (end)
     )); i += 1 - 2 * ((begin) > (end)))
#define sgn(a) ((a) > eps ? 1 : ((a) < -eps ? -1 : 0)
#define precise(x) fixed << setprecision(x)
```

```
#define debug(x) cerr << "> " << #x << " = "
     << x << endl;
#define pb push back
#define rnd(a, b) (uniform_int_distribution<int
     >((a), (b))(rng))
#ifndef LOCAL
   #define cerr if(0)cout
   #define endl "\n"
mt19937 rng(chrono::steady_clock::now().
     time since epoch().count());
clock_t __clock__;
void startTime() {___clock___ = clock();}
void timeit(string msg) {cerr << "> " << msg
     << ": " << precise(6) << ld(clock()-
        clock )/CLOCKS PER SEC << endl
const ld PI = asin(1) * 2;
const ld eps = 1e-14;
const int oo = 2e9;
const ll OO = 2e18;
const ll MOD = 10000000007;
const int MAXN = 1000000;
int main() {
   FAST IO:
   startTime();
   timeit("Finished");
   return 0;
```

2.4 Compilation

```
# Simple compile
g++ -DLOCAL -O2 -o main.exe -std-c++17 -Wall
-Wno-unused-result -Wshadow main.cpp
# Debug
g++ -DLOCAL -std=c++17 -Wshadow -Wall -o
main.exe main.cpp -fsanitize=address -
fsanitize=undefined -fuse-ld=gold -
D GLIBCXX DEBUG -g
```

2.5 Ternary Search

3 Geometry

3.1 2d Vector

```
template <typename T>
struct Vec {
   T x, y;
   Vec(): x(0), y(0) \{ \}
   Vec(T _x, T _y): x(_x), y(_y) {}
   Vec operator+(const Vec& b) {
      return Vec < T > (x+b.x, y+b.y);
   Vec operator-(const Vec& b) {
      return Vec<T>(x-b.x, y-b.y);
   Vec operator*(T c) {
      return Vec(x*c, y*c);
   T operator*(const Vec& b) {
      return x*b.x + y*b.y;
   T operator (const Vec& b) {
      return x*b.y-y*b.x;
   bool operator<(const Vec& other) const {
      if(x == other.x) return y < other.y;
      return x < other.x;
   bool operator==(const Vec& other) const {
      return x==other.x && y==other.y;
   bool operator!=(const Vec& other) const {
```

```
return !(*this == other);
  friend ostream& operator << (ostream& out,
       const Vec& v) {
     return out << "(" << v.x << ", " << v.v
  friend istream& operator>>(istream& in, Vec<
     return in >> v.x >> v.y;
  T norm() { // squared length
     return (*this)*(*this);
  ld len() {
     return sqrt(norm());
  ld angle(const Vec& other) { // angle between
        this and other vector
     return acosl((*this)*other/len()/other.len())
  Vec perp() {
     return Vec(-y, x);
* Cross product of 3d vectors: (ay*bz-az*by, az*
    bx-ax*bz, ax*by-ay*bx)
```

3.2 Circle Circle Intersection

Let's say that the first circle is centered at (0,0) (if it's not, we can move the origin to the center of the first circle and adjust the coordinates), and the second one is at (x_2, y_2) . Then, let's construct a line Ax + By + C = 0, where $A = -2x_2$, $B = -2y_2$, $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$. Finding the intersection between this line and the first circle will give us the answer. The only tricky case: if both circles are centered at the same point. We handle this case separately.

3.3 Circle Line Intersection

```
double r, a, b, c; // ax+by+c=0, radius is at (0, 0)
// If the center is not at (0, 0), fix the constant c
     to translate everything so that center is at
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b)
if (c*c > r*r*(a*a+b*b)+eps)
   puts ("no points");
else if (abs (c*c - r*r*(a*a+b*b)) < eps) {
   puts ("1 point");
   cout << x0 << ', ' << v0 << '\n';
   double d = r*r - c*c/(a*a+b*b);
   double mult = sqrt (d / (a*a+b*b));
   double ax, ay, bx, by;
   ax = x0 + b * mult;
   bx = x0 - b * mult;
   av = v0 - a * mult;
   by = y0 + a * mult;
   puts ("2 points");
   cout << ax << ',' , << ay << '\n' << bx <<
         ,,<< by << , ,
```

3.4 Common Tangents To Two Circles

```
struct pt {
    double x, y;

    pt operator- (pt p) {
        pt res = { x-p.x, y-p.y };
        return res;
    }
};
struct circle : pt {
        double r;
};
struct line {
        double a, b, c;
};
void tangents (pt c, double r1, double r2, vector <
        line > & ans) {
```

```
double r = r2 - r1;
   double z = sqr(c.x) + sqr(c.y);
   double d = z - sqr(r);
   if (d < -eps) return;
   d = sqrt (abs (d));
   line l;
   l.a = (c.x * r + c.y * d) / z;
   l.b = (c.y * r - c.x * d) / z;
   l.c = r1;
   ans.push_back (l);
vector<line> tangents (circle a, circle b) {
    vector<line> ans;
   for (int i=-1; i<=1; i+=2)
       for (int j=-1; j<=1; j+=2)
          tangents (b-a, a.r*i, b.r*j, ans);
    for (size t i=0; i<ans.size(); ++i)
       ans[i].c -= ans[i].a * a.x + ans[i].b * a.y;
    return ans;
```

3.5 Convex Hull Gift Wrapping

```
vector<Vec<int>> buildConvexHull(vector<Vec<
     int >> \& pts)
   int n = pts.size();
   sort(pts.begin(), pts.end());
   auto currP = pts[0]; // choose some extreme
        point to be on the hull
   vector<Vec<int>> hull;
   set<Vec<int>> used;
   hull.pb(pts[0]);
   used.insert(pts[0]);
   while(true) {
      auto candidate = pts[0]; // choose some
            point to be a candidate
      auto currDir = candidate-currP;
      vector<Vec<int>> toUpdate;
      FOR(i, 0, n) {
          if(currP == pts[i]) continue;
          // currently we have currP->candidate
          // we need to find point to the left of
          auto possibleNext = pts[i];
          auto nextDir = possibleNext - currP;
```

3.6 Convex Hull With Graham's Scan

```
// Takes in >= 3 points
// Returns convex hull in clockwise order
// Ignores points on the border
vector<Vec<int>> buildConvexHull(vector<Vec<
     int >> pts) {
   if(pts.size() <= 3) return pts;
   sort(pts.begin(), pts.end());
   stack<Vec<int>> hull;
   hull.push(pts[0]);
   auto p = pts[0];
   sort(pts.begin()+1, pts.end(), [&](Vec<int> a,
         Vec < int > b) -> bool {
      // p->a->b is a ccw turn
      int turn = sgn((a-p)^(b-a));
      //if(turn == 0) return (a-p).norm() > (b-p)
            ).norm();
       // among collinear points, take the
            farthest one
      return turn == 1;
   hull.push(pts[1]);
   FOR(i, 2, (int)pts.size()) {
      auto c = pts[i];
      if(c == hull.top()) continue;
      while(true) {
          auto a = hull.top(); hull.pop();
```

```
auto b = hull.top();
       auto ba = a-b;
       auto ac = c-a:
       if((ba^ac) > 0) {
           hull.push(a);
          break;
       else if((ba^ac) == 0) {
          if(ba*ac < 0) c = a;
          // ^ c is between b and a, so it
                shouldn't be added to the hull
          break;
   hull.push(c);
vector<Vec<int>> hullPts;
while(!hull.empty()) {
   hullPts.pb(hull.top());
   hull.pop();
return hullPts;
```

3.7 Line

```
template <typename T>
struct Line { // expressed as two vectors
    Vec<T> start, dir;
    Line() {}
    Line(Vec<T> a, Vec<T> b): start(a), dir(b-a)
    {}

    Vec<ld> intersect(Line l) {
        ld t = ld((l.start-start)^l.dir)/(dir^l.dir);
        // For segment-segment intersection this
            should be in range [0, 1]
        Vec<ld> res(start.x, start.y);
        Vec<ld> dirld(dir.x, dir.y);
        return res + dirld*t;
    }
};
```

3.8 Number Of Lattice Points On Segment

Let's say we have a line segment from (x_1, y_1) to (x_2, y_2) . Then, the number of lattice points on this segment is given by

$$gcd(x_2 - x_1, y_2 - y_1) + 1.$$

3.9 Pick's Theorem

We are given a lattice polygon with nonzero area. Let's denote its area by S, the number of points with integer coordinates lying strictly inside the polygon by I and the number of points lying on the sides of the polygon by B. Then:

$$S = I + \frac{B}{2} - 1.$$

4 Graphs

4.1 Bellman Ford Algorithm

```
struct Edge
{
    int a, b, cost;
};
int n, m, v; // v - starting vertex
vector<Edge> e;
/* Finds SSSP with negative edge weights.
```

```
* Possible optimization: check if anything changed in a relaxation step. If not - you can break early.

* To find a negative cycle: perform one more relaxation step. If anything changes - a negative cycle exists.

*/
void solve() {
    vector<int> d (n, oo);
    d[v] = 0;
    for (int i=0; i<n-1; ++i)
        for (int j=0; j<m; ++j)
        if (d[e]j].a] < oo)
            d[e[j].b] = min (d[e[j].b], d[e[j].a] + e
            [j].cost);
    // display d, for example, on the screen
```

4.2 Bipartite Graph

```
class BipartiteGraph {
private:
    vector < int > \_left, \_right;
    vector<vector<int>> _adjList;
   vector<int> _matchR, _matchL;
    vector<bool> _used;
    bool _kuhn(int v) {
       if (_used[v]) return false;
        used[v] = true;
       FOR(i, 0, (int)_adjList[v].size()) {
           int to = _adjList[v][i] - _left.size();
           if (_matchR[to] == -1 || _kuhn(
                 _matchR[to])) {
               matchR[to] = v;
               _{\text{matchL}[v]} = to;
               return true;
       return false;
    void _addReverseEdges() {
       FOR(i, 0, (int) right.size()) {
           if (\underline{\text{matchR}[i]} != -1) {
               adjList[left.size() + i].pb(
                     _matchR[i]);
```

```
void _dfs(int p) {
       if (_used[p]) return;
        \_used[p] = true;
       for (auto x : _adjList[p]) {
            _{dfs(x);
    vector<pii> buildMM() {
       vector<pair<int, int> > res;
       FOR(i, 0, (int)_right.size()) {
           if ( \operatorname{matchR}[i] != -1) {
               res.push_back(make_pair(_matchR[
                     i], i));
       return res;
public:
    void addLeft(int x) {
        _{\text{left.pb}(x)};
        adjList.pb({});
       _{\text{matchL.pb}(-1)};
        _used.pb(false);
    void addRight(int x) {
        _{right.pb(x)}
        _{adjList.pb({});}
        matchR.pb(-1);
        used.pb(false);
   void addForwardEdge(int l, int r) {
        \_adjList[l].pb(r + \_left.size());
   void addMatchEdge(int l, int r) {
       if(l != -1) \quad matchL[l] = r;
       if(r != -1) _{matchR[r]} = 1;
    // Maximum Matching
    vector<pii> mm() {
       matchR = vector < int > (right.size(), -1);
        _{\text{matchL}} = \text{vector} < \text{int} > (_{\text{left.size}}(), -1);
       // ^ these two can be deleted if performing
             MM on already partially matched
              graph
        used = vector < bool > (left.size() +
              _right.size(), false);
       bool path found;
            fill( used.begin(), used.end(), false);
            path\_found = false;
```

```
FOR(i, 0, (int) left.size()) {
           if (_{matchL[i]} < 0 \&\& !_{used[i]})  {
               path found |= kuhn(i);
   } while (path found);
   return buildMM();
// Minimum Edge Cover
// Algo: Find MM, add unmatched vertices
     greedily.
vector<pii> mec() {
   auto ans = mm();
   FOR(i, 0, (int)_left.size()) {
       if ( matchL[i] != -1) {
           for (auto x : adjList[i]) {
               int ridx = x - left.size();
               if (\underline{\text{matchR}}[\text{ridx}] == -1) {
                  ans.pb(\{ i, ridx \});
                   matchR[ridx] = i;
   FOR(i, 0, (int) left.size()) {
       if (\underline{matchL[i]} == -1 \&\& (int)\underline{adjList[i]}
             |.size() > 0)  {
           int ridx = adjList[i][0] - left.size()
           _{\text{matchL}[i]} = \text{ridx};
           ans.pb(\{i, ridx\});
   return ans;
// Minimum Vertex Cover
// Algo: Find MM. Run DFS from unmatched
     vertices from the left part.
// MVC is composed of unvisited LEFT and
     visited RIGHT vertices.
pair<vector<int>, vector<int>> mvc(bool
     runMM = true) {
   if (runMM) mm();
    addReverseEdges();
   fill(_used.begin(), _used.end(), false);
   FOR(i, 0, (int) left.size()) {
       if (\underline{\text{matchL}[i]} == -1) {
           dfs(i);
```

```
vector<int> left, right;
   FOR(i, 0, (int)_left.size()) {
       if (!_used[i]) left.pb(i);
   FOR(i, 0, (int)_right.size()) {
       if (_used[i + (int)_left.size()]) right.pb(
   return { left,right };
// Maximal Independent Vertex Set
// Algo: Find complement of MVC.
pair < vector < int >, vector < int >> mivs(bool
     runMM = true) {
   auto m = mvc(runMM);
   vector<br/>bool> containsL( left.size(), false),
         containsR( right.size(), false);
   for (auto x : m.first) containsL[x] = true;
   for (auto x : m.second) containsR[x] = true;
   vector<int> left, right;
   FOR(i, 0, (int)_left.size())
       if (!containsL[i]) left.pb(i);
   FOR(i, 0, (int)_right.size()) {
       if (!containsR[i]) right.pb(i);
   return { left, right };
```

4.3 Dfs With Timestamps

```
\label{eq:constraint} \begin{split} & \operatorname{vector} < \operatorname{vector} < \operatorname{int} > \operatorname{adj}; \\ & \operatorname{vector} < \operatorname{int} > \operatorname{tIn}, \ \operatorname{tOut}, \ \operatorname{color}; \\ & \operatorname{int} \ \operatorname{dfs\_timer} = 0; \\ & \operatorname{void} \ \operatorname{dfs}(\operatorname{int} \ v) \ \{ \\ & \operatorname{tIn}[v] = \operatorname{dfs\_timer} + +; \\ & \operatorname{color}[v] = 1; \\ & \operatorname{for} \ (\operatorname{int} \ u : \operatorname{adj}[v]) \\ & \quad \operatorname{if} \ (\operatorname{color}[u] = = 0) \\ & \quad \operatorname{dfs}(u); \\ & \operatorname{color}[v] = 2; \\ & \operatorname{tOut}[v] = \operatorname{dfs\_timer} + +; \\ \} \end{split}
```

4.4 Finding Articulation Points

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of
      graph
vector<br/>bool> visited;
vector<int> tin, fup;
int timer;
void processCutpoint(int v) {
    // problem-specific logic goes here
    // it can be called multiple times for the same
void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = fup[v] = timer++;
    int children=0;
    for (int to : adj[v]) {
       if (to == p) continue;
       if (visited[to]) {
           fup[v] = min(fup[v], tin[to]);
       } else {
           dfs(to, v);
           fup[v] = min(fup[v], fup[to]);
           if (fup[to] >= tin[v] \&\& p!=-1)
               processCutpoint(v);
           ++children;
    if(p == -1 \&\& children > 1)
       processCutpoint(v);
void findCutpoints() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    fup.assign(n, -1);
    for (int i = 0; i < n; ++i) {
       if (!visited[i])
           dfs (i);
```

4.5 Finding Bridges

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of
      graph
vector<br/>bool> visited;
vector<int> tin, fup;
int timer;
void processBridge(int u, int v) {
    // do something with the found bridge
void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = fup[v] = timer++;
   for (int to : adj[v]) {
       if (to == p) continue;
       if (visited[to]) {
           fup[v] = min(fup[v], tin[to]);
       } else {
           dfs(to, v);
           fup[v] = min(fup[v], fup[to]);
           if (fup[to] > tin[v])
               processBridge(v, to);
// Doesn't work with multiple edges
// But multiple edges are never bridges, so it's
      easy to check
void findBridges() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    fup.assign(n, -1);
   bridges.clear();
    FOR(i, 0, n) {
       if (!visited[i])
           dfs(i);
```

4.6 Lowest Common Ancestor

```
int n, l; // l == logN (usually about ~20)
vector<vector<int>> adj;
int timer;
vector<int> tin, tout;
vector<vector<int>> up;
void dfs(int v, int p)
   tin[v] = ++timer;
   up[v][0] = p;
   // wUp[v][0] = weight[v][u]; // <- path weight
         sum to 2^i-th ancestor
   for (int i = 1; i <= l; ++i)
       up[v][i] = up[up[v][i-1]][i-1];
       // wUp[v][i] = wUp[v][i-1] + wUp[up[v][i
             -1]][i-1];
   for (int u:adj[v]) {
      if (u != p)
          dfs(u, v);
   tout[v] = ++timer;
bool isAncestor(int u, int v)
   return tin[u] \le tin[v] \&\& tout[v] \le tout[u];
int lca(int u, int v)
   if (isAncestor(u, v))
       return u;
   if (isAncestor(v, u))
       return v;
   for (int i = l; i >= 0; --i) {
       if (!isAncestor(up[u][i], v))
          u = up[u][i];
   return up[u][0];
void preprocess(int root) {
   tin.resize(n);
   tout.resize(n);
```

```
\begin{array}{l} timer = 0; \\ l = ceil(log2(n)); \\ up.assign(n, vector < int > (l+1)); \\ dfs(root, root); \end{array}
```

4.7 Max Flow With Dinic 2

```
struct FlowEdge {
   int v, u;
   long long cap, flow = 0;
   FlowEdge(int v, int u, long long cap): v(v), u(
         u), cap(cap) {}
struct Dinic {
   const long long flow_inf = 1e18;
   vector<FlowEdge> edges;
   vector < vector < int >> adj;
   int n, m = 0;
   int s. t:
   vector<int> level, ptr;
   queue<int> q;
   Dinic(int n, int s, int t) : n(n), s(s), t(t) {
       adj.resize(n);
       level.resize(n);
       ptr.resize(n);
    void add_edge(int v, int u, long long cap) {
       edges.push back(FlowEdge(v, u, cap));
       edges.push_back(FlowEdge(u, v, 0));
       adj[v].push_back(m);
       adj[u].push back(m + 1);
       m += 2;
   bool bfs() {
       while (!q.empty()) {
          int v = q.front();
           q.pop();
           for (int id : adj[v]) {
               if (edges[id].cap - edges[id].flow < 1)
              if (level[edges[id].u] != -1)
                  continue;
               level[edges[id].u] = level[v] + 1;
              q.push(edges[id].u);
```

```
return level[t] != -1;
long long dfs(int v, long long pushed) {
    if (pushed == 0)
       return 0;
   if (v == t)
       return pushed;
    for (int\& cid = ptr[v]; cid < (int)adj[v].size
          (); cid++) {
       int id = adj[v][cid];
       int u = edges[id].u;
       if (level[v] + 1 != level[u] || edges[id].
              cap - edges[id].flow < 1
           continue;
       long long tr = dfs(u, min(pushed, edges))
              id].cap - edges[id].flow));
       if (tr == 0)
           continue;
       \mathrm{edges}[\mathrm{id}].\mathrm{flow} \ += \mathrm{tr};
       edges[id \hat{} 1].flow -= tr;
       return tr;
    return 0;
long long flow() {
    long long f = 0;
    while (true) {
       fill(level.begin(), level.end(), -1);
       level[s] = 0;
       q.push(s);
       if (!bfs())
            break:
       fill(ptr.begin(), ptr.end(), 0);
       while (long long pushed = dfs(s,
              flow inf)) {
           f += pushed;
    return f;
```

4.8 Max Flow With Dinic

```
struct Edge {
```

```
int f, c;
   int to:
   pii revIdx;
   int dir;
   int idx;
vector<Edge> adjList[MAX N];
int level[MAX_N];
void addEdge(int a, int b, int c, int i, int dir) {
   int idx = adjList[a].size();
   int revIdx = adjList[b].size();
   adjList[a].pb(\{ 0,c,b, \{b, revIdx\}, dir,i \});
   adjList[b].pb(\{0,0,a,\{a,idx\},dir,i\});
bool bfs(int s, int t) {
   FOR(i, 0, n) level[i] = -1;
   level[s] = 0;
   queue<int> Q;
   Q.push(s);
   while (!Q.empty()) {
       auto t = Q.front(); Q.pop();
       for (auto x : adjList[t]) {
           if (level[x.to] < 0 && x.f < x.c) {
              level[x.to] = level[t] + 1;
               Q.push(x.to);
   return level[t] >= 0;
int send(int u, int f, int t, vector<int>& edgeIdx)
   if (u == t) return f;
   for (; edgeIdx[u] < adjList[u].size(); edgeIdx[u
       auto\& e = adjList[u][edgeIdx[u]];
      if (level[e.to] == level[u] + 1 && e.f < e.c)
           int curr flow = min(f, e.c - e.f);
           int next_flow = send(e.to, curr_flow, t,
                edgeIdx);
           if (\text{next\_flow} > 0) {
               e.f += next flow;
               adjList[e.revIdx.first][e.revIdx.second\\
                    l.f -= next flow;
               return next flow;
```

```
return 0;
int maxFlow(int s, int t) {
    int f = 0;
    while (bfs(s, t)) {
        vector < int > edgeIdx(n, 0);
        while (int extra = send(s, oo, t, edgeIdx)) {
           f += extra;
    return f;
void init() {
    cin >> n >> m;
    FOR(i, 0, m) {
        int a, b, c;
        \mathrm{cin}>>\mathrm{a}>>\mathrm{b}>>\mathrm{c};
        a--; b--;
        addEdge(a, b, c, i, 1);
        addEdge(b, a, c, i, -1);
```

4.9 Max Flow With Ford Fulkerson

```
struct Edge {
    int to, next;
    ll f, c;
    int idx, dir;
    int from;
};

int n, m;
vector<Edge> edges;
vector<int> first;

void addEdge(int a, int b, ll c, int i, int dir) {
    edges.pb({ b, first[a], 0, c, i, dir, a });
    edges.pb({ a, first[b], 0, 0, i, dir, b });
    first[a] = edges.size() - 2;
    first[b] = edges.size() - 1;
}

void init() {
```

```
cin >> n >> m;
    edges.reserve(4 * m);
    first = vector < int > (n, -1);
    FOR(i, 0, m) {
       int a, b, c;
       cin >> a >> b >> c;
       a--; b--;
       addEdge(a, b, c, i, 1);
       addEdge(b, a, c, i, -1);
int cur\_time = 0;
vector<int> timestamp;
ll dfs(int v, ll flow = OO) {
    if (v == n - 1) return flow;
    timestamp[v] = cur time;
    for (int e = first[v]; e != -1; e = edges[e].next)
       if (edges[e].f < edges[e].c && timestamp[
             edges[e].to] != cur_time) {
           int pushed = dfs(edges[e].to, min(flow,
                 edges[e].c - edges[e].f));
           if (pushed > 0) {
               edges[e].f += pushed;
edges[e ^ 1].f -= pushed;
               return pushed;
    return 0;
ll maxFlow() {
    cur time = 0;
    timestamp = vector < int > (n, 0);
   ll f = 0, add;
    while (true) {
       cur time++;
       add = dfs(0);
       if (add > 0) {
           f += add;
       else {
           break;
    return f;
```

4.10 Min Cut

```
\label{eq:continuous_series} \begin{split} & \text{init}(); \\ & \text{ll } f = \text{maxFlow}(); \text{ // Ford-Fulkerson} \\ & \text{cur\_time} + +; \\ & \text{dfs}(0); \\ & \text{set} < \text{int} > \text{cc}; \\ & \text{for } (\text{auto } e : \text{edges}) \text{ } \{ \\ & \text{ if } (\text{timestamp[e.from]} == \text{cur\_time } \&\& \\ & \text{ timestamp[e.to]} \text{ != cur\_time } \} \\ & \text{ cc.insert}(e.\text{idx}); \\ & \text{ } \} \\ & \text{ // } (\# \text{ of edges in min-cut, capacity of cut)} \\ & \text{ // } [\text{indices of edges forming the cut]} \\ & \text{ cout } << \text{cc.size}() << " " << f << \text{endl}; \\ & \text{ for } (\text{auto } \text{x} : \text{cc}) \text{ cout} << \text{x} + 1 << " "; \\ \end{split}
```

4.11 Number Of Paths Of Fixed Length

Let G be the adjacency matrix of a graph. Then $C_k = G^k$ gives a matrix, in which the value $C_k[i][j]$ gives the number of paths between i and j of length k.

4.12 Shortest Paths Of A Fixed Length

Define $A \odot B = C \iff C_{ij} = \min_{p=1..n} (A_{ip} + B_{pj})$. Let G be the adjacency matrix of a graph. Also, let $L_k = G \odot \ldots \odot G = G^{\odot k}$. Then the value $L_k[i][j]$ denotes the length of the shortest path between i and j which consists of exactly k edges.

4.13 Strongly Connected Components

```
vector < vector < int> > g, gr; // adjList and
      reversed adjList
vector<br/>bool> used:
vector<int> order, component;
void dfs1 (int v) {
    used[v] = true;
    for (size_t i=0; i< g[v].size(); ++i)
       if (!used[ g[v][i] ])
           dfs1 (g[v][i]);
    order.push back (v);
void dfs2 (int v) {
    used[v] = true;
    component.push back (v);
    for (size_t i=0; i < gr[v].size(); ++i)
       if (!used[gr[v][i]])
           dfs2 (gr[v][i]);
int main() {
   int n;
    // read n
    for (;;) {
       int a, b;
       // read edge a -> b
       g[a].push_back (b);
       gr[b].push back (a);
    used.assign (n, false);
    for (int i=0; i< n; ++i)
       if (!used[i])
           dfs1 (i):
    used.assign (n, false);
    for (int i=0; i< n; ++i) {
       int v = order[n-1-i]:
       if (!used[v]) {
           dfs2 (v):
           // do something with the found
                 component
           component.clear(); // components are
                 generated in toposort-order
```

5 Math

5.1 Big Integer Multiplication With FFT

```
complex<ld> a[MAX_N], b[MAX_N];
complex<ld>fa[MAX_N], fb[MAX_N], fc[
     MAX N];
complex<ld>cc[MAX_N];
string mul(string as, string bs) {
   int sgn1 = 1;
   int sgn2 = 1;
   if (as[0] == '-') {
       sgn1 = -1;
       as = as.substr(1);
   if (bs[0] == '-') {
       sgn2 = -1;
       bs = bs.substr(1);
   int n = as.length() + bs.length() + 1;
   FFT::init(n);
   FOR(i, 0, FFT::pwrN) {
       a[i] = b[i] = fa[i] = fb[i] = fc[i] = cc[i] = 0;
   FOR(i, 0, as.size()) {
       a[i] = as[as.size() - 1 - i] - '0';
   FOR(i, 0, bs.size()) {
       b[i] = bs[bs.size() - 1 - i] - '0';
   FFT::fft(a, fa);
   FFT::fft(b, fb);
   FOR(i, 0, FFT::pwrN) {
      fc[i] = fa[i] * fb[i];
   // turn [0,1,2,...,n-1] into [0, n-1, n-2, ..., 1]
   FOR(i, 1, FFT::pwrN) {
      if (i < FFT::pwrN - i) {
          swap(fc[i], fc[FFT::pwrN - i]);
   FFT::fft(fc, cc);
   ll carry = 0;
   vector<int> v;
   FOR(i, 0, FFT::pwrN) {
       int num = round(cc[i].real() / FFT::pwrN)
             + carry;
```

```
v.pb(num % 10);
   carry = num / 10;
while (carry > 0) {
   v.pb(carry % 10);
   carry \neq 10;
reverse(v.begin(), v.end());
bool start = false;
ostringstream ss;
bool allZero = true;
for (auto x : v) {
   if (x!= 0) {
       allZero = false;
       break;
if (sgn1*sgn2 < 0 \&\& !allZero) ss << "-";
for (auto x : v) {
   if (x == 0 \&\& !start) continue;
   start = true;
   ss \ll abs(x);
if (!start) ss << 0;
return ss.str():
```

5.2 Burnside's Lemma

Let G be a finite group that acts on a set X. For each g in G let X^g denote the set of elements in X that are fixed by g. Burnside's lemma asserts the following formula for the number of orbits:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

Example. Coloring a cube with three colors.

Let X be the set of 3^6 possible face color combinations. Let's count the sizes of the fixed sets for each of the 24 rotations:

- one 0-degree rotation which leaves all 3^6 elements of X unchanged
- six 90-degree face rotations, each of which leaves 3^3 elements of X unchanged
- three 180-degree face rotation, each of which leaves 3^4 elements of X unchanged
- eight 120-degree vertex rotations, each of which leaves 3^2 elements of X unchanged
- six 180-degree edge rotations, each of which leaves 3^3 elements of X unchanged

The average is then $\frac{1}{24}(3^6 + 6 \cdot 3^3 + 3 \cdot 3^4 + 8 \cdot 3^2 + 6 \cdot 3^3) = 57$. For n colors: $\frac{1}{24}(n^6 + 3n^4 + 12n^3 + 8n^2)$.

Example. Coloring a circular stripe of n cells with two colors.

X is the set of all colored striped (it has 2^n elements), G is the group of rotations (n elements - by 0 cells, by 1 cell, ..., by (n-1) cells). Let's fix some K and find the number of stripes that are fixed by the rotation by K cells. If a stripe becomes itself after rotation by K cells, then its 1st cell must have the same color as its $(1+K \mod n)$ -th cell, which is in turn the same as its $(1+2K \mod n)$ -th cell, etc., until $mK \mod n = 0$. This will happen when m = n/gcd(K, n). Therefore, we

have $n/\gcd(K,n)$ cells that must all be of the same color. The same will happen when starting from the second cell and so on. Therefore, all cells are separated into $\gcd(K,n)$ groups, with each group being of one color, and that yields $2^{\gcd(K,n)}$ choices. That's why the answer to the original problem is $\frac{1}{n}\sum_{k=0}^{n-1}2^{\gcd(k,n)}$.

5.3 Chinese Remainder Theorem

Let's say we have some numbers m_i , which are all mutually coprime. Also, let $M = \prod_i m_i$. Then the system of congruences

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \dots \\ x \equiv a_k \pmod{m_k} \end{cases}$$

is equivalent to $x \equiv A \pmod{M}$ and there exists a unique number A satisfying 0 < A < M.

Solution for two: $x \equiv a_1 \pmod{m_1}$, $x \equiv a_2 \pmod{m_2}$. Let $x = a_1 + km_1$. Substituting into the second congruence: $km_1 \equiv a_2 - a_1 \pmod{m_2}$. Then, $k = (m_1)_{m_2}^{-1}(a_2 - a_1) \pmod{m_2}$. and we can easily find x. This can be extended to multiple equations by solving them one-by-one.

If the moduli are not coprime, solve the system $y \equiv 0 \pmod{\frac{m_1}{g}}, y \equiv \frac{a_2-a_1}{g} \pmod{\frac{m_2}{g}}$ for y. Then let $x \equiv gy + a_1 \pmod{\frac{m_1m_2}{g}}$.

5.4 Euler Totient Function

5.5 Extended Euclidean Algorithm

```
// ax+by=gcd(a,b)
void solveEq(ll a, ll b, ll& x, ll& y, ll& g) {
   if(b==0) {
      x = 1;
      y = 0;
      g = a;
      return;
   ĺl xx, yy;
   solveEq(b, a%b, xx, yy, g);
   x = yy;
   y = xx-yy*(a/b);
// ax+bv=c
bool solveEq(ll a, ll b, ll c, ll& x, ll& y, ll& g) {
   solveEq(a, b, x, y, g);
   if(c\%g != 0) return false;
   x *= c/g; y *= c/g;
   return true;
// Finds a solution (x, y) so that x >= 0 and x is
bool solveEqNonNegX(ll a, ll b, ll c, ll& x, ll &y, ll
     & g) {
   if(!solveEq(a, b, c, x, y, g)) return false;
   ll k = x*g/b;
   x = x - k*b/g;
   y = y + k*a/g;
   if(x < 0) {
      x += b/g;
```

```
y -= a/g;
}
return true;
}
```

All other solutions can be found like this:

$$x' = x - k\frac{b}{g}, y' = y + k\frac{a}{g}, k \in \mathbb{Z}$$

5.6 Factorization With Sieve

```
// Use linear sieve to calculate minDiv
vector<pll> factorize(ll x) {
   vector<pll> res;
   ll prev = -1;
   ll cnt = 0;
   while(x != 1) {
       ll d = minDiv[x];
       if(d == prev) {
          cnt++;
       } else {
          if(prev != -1) res.pb(\{prev, cnt\});
          prev = d;
          cnt = 1;
       x /= d;
   res.pb({prev, cnt});
   return res;
```

5.7 FFT With Modulo

```
\begin{split} & bool \; is Generator(ll \; g) \; \{ \\ & \; if \; (pwr(g, \; M-1) \; != 1) \; return \; false; \\ & for \; (ll \; i=2; \; i^*i <= M-1; \; i++) \; \{ \\ & \; if \; ((M-1) \; \% \; i==0) \; \{ \\ & \; ll \; q=i; \\ & \; if \; (is Prime(q)) \; \{ \\ & \; ll \; p=(M-1) \; / \; q; \\ & \; ll \; pp=pwr(g, \; p); \end{split}
```

```
if (pp == 1) return false;
          q = (M - 1) / i;
          if (isPrime(q)) {
             ll p = (M - 1) / q;
             ll pp = pwr(g, p);
             if (pp == 1) return false;
   return true;
namespace FFT {
   ll n;
   vector<ll> r;
   vector<ll> omega;
   ll logN, pwrN;
   void initLogN() {
      logN = 0;
       pwrN = 1;
       while (pwrN < n) {
          pwrN *= 2;
          logN++;
      n = pwrN;
   void initOmega() {
      ll g = 2;
       while (!isGenerator(g)) g++;
      ll G = 1;
      g = pwr(g, (M - 1) / pwrN);
      FOR(i, 0, pwrN) {
          omega[i] = G;
          G *= g;
          G \% = M;
   void initR() {
      r[0] = 0;
      FOR(i, 1, pwrN) {
          r[i] = \hat{r[i / 2]} / 2 + ((i \& 1) << (logN -
               1));
   void initArrays() {
      r.clear();
      r.resize(pwrN);
```

omega.clear();

```
omega.resize(pwrN);
void init(ll n) {
   FFT::n = n;
   initLogN();
   initArrays();
   initOmega();
   initR();
void fft(ll a[], ll f[]) {
   for (ll i = 0; i < pwrN; i++) {
       f[i] = a[r[i]];
   for (ll k = 1; k < pwrN; k *= 2) {
       for (ll i = 0; i < pwrN; i += 2 * k) {
           for (ll j = 0; j < k; j++) {
              auto z = omega[j*n / (2 * k)] * f
                    [i + j + k] \% M;
               f[i + j + k] = f[i + j] - z;
              f[i + j] += z;
              f[i + j + k] \% = M;
              if (f[i + j + k] < 0) f[i + j + k]
                    += M;
              f[i + j] \% = M;
```

5.8 FFT

```
namespace FFT {
   int n;
   vector<int> r;
   vector<complex<ld>>> omega;
   int logN, pwrN;

void initLogN() {
     logN = 0;
     pwrN = 1;
     while (pwrN < n) {
        pwrN *= 2;
        logN++;
     }
     n = pwrN;
}</pre>
```

```
void initOmega() {
   FOR(i, 0, pwrN) {
      omega[i] = { cos(2 * i*PI / n), sin(2 * i)
             *PI / n) };
void initR() {
   r[0] = 0;
   FOR(i, 1, pwrN) {
       r[i] = \hat{r}[i / 2] / 2 + ((i \& 1) << (logN -
            1));
void initArrays() {
   r.clear();
   r.resize(pwrN);
   omega.clear();
   omega.resize(pwrN);
void init(int n) {
   FFT::n = n;
   initLogN();
   initArrays();
   initOmega();
   initR();
void fft(complex<ld> a[], complex<ld> f[]) {
   FOR(i, 0, pwrN) {
       f[i] = a[r[i]];
   for (ll k = 1; k < pwrN; k *= 2) {
       for (ll i = 0; i < pwrN; i += 2 * k) {
           for (ll j = 0; j < k; j++) {
              auto z = \text{omega}[j*n / (2 * k)] * f
                    [i + j + k];
              f[i + j + k] = f[i + j] - z;
              f[i + j] += z;
        }
}
```

5.9 Formulas

```
\begin{array}{lll} \sum_{i=1}^n i & = & \frac{n(n+1)}{2}; & \sum_{i=1}^n i^2 & = \\ \frac{n(2n+1)(n+1)}{6}; & \sum_{i=1}^n i^3 & = & \frac{n^2(n+1)^2}{4}; \\ \sum_{i=1}^n i^4 & = & \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}; \\ \sum_{i=a}^b c^i & = & \frac{c^{b+1}-c^a}{c-1}, c \neq 1; & \sum_{i=1}^n a_1 + \\ (i-1)d & = & \frac{n(a_1+a_n)}{2}; & \sum_{i=1}^n a_1 r^{i-1} = \\ \frac{a_1(1-r^n)}{1-r}, r \neq 1; & \sum_{i=1}^\infty ar^{i-1} = \frac{a_1}{1-r}, |r| \leq 1. \end{array}
```

5.10 Linear Sieve

```
\begin{split} & ll \; minDiv[MAXN+1]; \\ & vector < ll > primes; \\ & void \; sieve(ll \; n) \{ \\ & \; FOR(k, \; 2, \; n+1) \{ \\ & \; minDiv[k] = k; \\ \} \\ & \; FOR(k, \; 2, \; n+1) \; \{ \\ & \; if(minDiv[k] = k) \; \{ \\ & \; primes.pb(k); \\ \} \\ & \; for(auto \; p : primes) \; \{ \\ & \; if(p > minDiv[k]) \; break; \\ & \; if(p > minDiv[k]) \; break; \\ & \; minDiv[p*k] = p; \\ \} \\ & \} \\ \} \end{split}
```

5.11 Modular Inverse

```
bool invWithEuclid(ll a, ll m, ll& aInv) {
    ll x, y, g;
    if(!solveEqNonNegX(a, m, 1, x, y, g)) return
        false;
    aInv = x;
    return true;
}
```

```
// Works only if m is prime
ll invFermat(ll a, ll m) {
    return pwr(a, m-2, m);
}
// Works only if gcd(a, m) = 1
ll invEuler(ll a, ll m) {
    return pwr(a, phi(m)-1, m);
}
```

5.12 Simpson Integration

```
 \begin{array}{l} {\rm const\ int\ N} = 1000\ ^*\ 1000;\ //\ number\ of\ steps\ (\\ {\rm already\ multiplied\ by\ 2}) \\ \\ {\rm double\ simpsonIntegration}({\rm double\ a,\ double\ b}) \{\\ {\rm double\ h} = (b\ -\ a)\ /\ N;\\ {\rm double\ s} = f(a)\ +\ f(b);\ //\ a = x\_0\ and\ b = \\ x\_2n\\ {\rm for\ (int\ i=1;\ i<=N-1;++i)\ \{}\\ {\rm double\ x=a+h\ *\ i;}\\ {\rm s\ +=f(x)\ *\ ((i\ \&\ 1)\ ?\ 4:2);}\\ \}\\ {\rm s\ *=h\ /\ 3;}\\ {\rm return\ s;} \\ \} \\ \end{array}
```

6 Strings

6.1 Aho Corasick Automaton

```
\label{eq:constint} \begin{split} // & \text{ alphabet size} \\ & \text{const int } K = 70; \\ // & \text{ the indices of each letter of the alphabet} \\ & \text{ int intVal}[256]; \\ & \text{void init()} \; \{ \\ & \text{ int curr} = 2; \\ & \text{ intVal}[1] = 1; \\ & \text{ for}(\text{char } c = '0'; \ c <= '9'; \ c++, \ curr++) \ \text{ intVal} \\ & & [(\text{int})c] = \text{curr}; \\ & \text{ for}(\text{char } c = 'A'; \ c <= 'Z'; \ c++, \ curr++) \\ & & \text{ intVal}[(\text{int})c] = \text{ curr}; \\ \end{split}
```

```
for(char c = 'a'; c \le 'z'; c++, curr++) intVal
         [(int)c] = curr;
struct Vertex {
    int next[K];
    vector<int> marks;
    // ^ this can be changed to int mark = -1, if
         there will be no duplicates
    int p = -1;
    char pch;
    int link = -1;
    int exitLink = -1;
    // ^ exitLink points to the next node on the
          path of suffix links which is marked
    int go[K];
    // ch has to be some small char
    Vertex(int p=-1, char ch=(char)1): p(p),
         pch(ch) {
       fill(begin(next), end(next), -1);
       fill(begin(go), end(go), -1);
};
vector < Vertex > t(1);
void addString(string const& s, int id) {
    int v = 0;
    for (char ch : s) {
       int c = intVal[(int)ch];
       if (t[v].next[c] == -1)
           t[v].next[c] = t.size();
           t.emplace_back(v, ch);
       v = t[v].next[c];
    t[v].marks.pb(id);
int go(int v, char ch);
int getLink(int v) {
    if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0)
           t[v].link = 0;
           t[v].link = go(getLink(t[v].p),\, t[v].pch);\\
    return t[v].link;
int getExitLink(int v) {
```

```
if(t[v].exitLink != -1) return t[v].exitLink;
   int l = getLink(v);
   if(l == 0) return t[v].exitLink = 0;
   if(!t[l].marks.empty()) return t[v].exitLink = l;
   return t[v].exitLink = getExitLink(1);
int go(int v, char ch) {
   int c = intVal[(int)ch];
   if (t[v].go[c] = -1) {
       if (t[v].next[c] != -1)
          t[v].go[c] = t[v].next[c];
       else
          t[v].go[c] = v == 0 ? 0 : go(getLink(v),
                ch);
   return t[v].go[c];
void walkUp(int v, vector<int>& matches) {
   if(v == 0) return;
   if(!t[v].marks.empty()) {
       for(auto m: t[v].marks) matches.pb(m);
   walkUp(getExitLink(v), matches);
// returns the IDs of matched strings.
// Will contain duplicates if multiple matches of
     the same string are found.
vector<int> walk(const string& s) {
   vector<int> matches;
   int curr = 0;
   for(char c : s) {
       curr = go(curr, c);
       if(!t[curr].marks.empty()) {
          for(auto m : t[curr].marks) matches.pb(
       walkUp(getExitLink(curr), matches);
   return matches;
/* Usage:
 * addŠtring(strs[i], i);
* auto matches = walk(text);
* .. do what you need with the matches - count,
      check if some id exists, etc ..
* Some applications:
* - Find all matches: just use the walk function
* - Find lexicographically smallest string of a
      given length that doesn't match any of the
      given strings:
```

```
* For each node, check if it produces any matches (it either contains some marks or walkUp(v) returns some marks).

* Remove all nodes which produce at least one match. Do DFS in the remaining graph, since none of the remaining nodes

* will ever produce a match and so they're safe.

* - Find shortest string containing all given strings:

* For each vertex store a mask that denotes the strings which match at this state. Start at (v = root, mask = 0),

* we need to reach a state (v, mask=2^n-1), where n is the number of strings in the set. Use BFS to transition between states

* and update the mask.
```

6.2 Hashing

```
struct HashedString {
   = 1000000087, B2 = 1000000097;
   vector<ll> A1pwrs, A2pwrs;
   vector<pll> prefixHash;
   HashedString(const string& _s) {
      init(\underline{\hspace{0.1cm}}s);
      calcHashes(\_s);
   void init(const string& s) {
      11 a1 = 1;
      11 \ a2 = 1:
      FOR(i, 0, (int)s.length()+1) {
         A1pwrs.pb(a1);
         A2pwrs.pb(a2);
         a1 = (a1*A1)\%B1;
         a2 = (a2*A2)\%B2;
   void calcHashes(const string& s) {
      pll h = \{0, 0\};
      prefixHash.pb(h);
      for(char c : s) {
         ll h1 = (prefixHash.back().first*A1 + c)
         ll h2 = (prefixHash.back().second*A2 +
         prefixHash.pb(\{h1, h2\});
```

```
 \begin{cases} & \text{pll getHash(int l, int r) } \{ \\ & \text{ll h1 = (prefixHash[r+1].first - prefixHash[l l l.first*Alpwrs[r+1-l]) \% B1;} \\ & \text{ll h2 = (prefixHash[r+1].second - prefixHash[l].second*A2pwrs[r+1-l]) \% B2;} \\ & \text{B2;} \\ & \text{if(h1 < 0) h1 += B1;} \\ & \text{if(h2 < 0) h2 += B2;} \\ & \text{return } \{\text{h1, h2}\}; \\ \} \}; \end{cases}
```

6.3 KMP

6.4 Prefix Function Automaton

```
 \begin{array}{c} if \ (i>0 \ \&\& \ BASE + c \ != s[i]) \\ aut[i][c] = aut[pi[i-1]][c]; \\ else \\ aut[i][c] = i + (BASE + c == s[i]); \\ \} \\ return \ aut; \\ \} \\ vector < int > findOccurs (const \ string\& \ s, \ const \ string\& \ t) \ \{ \\ auto \ aut = computeAutomaton(s); \\ int \ curr = 0; \\ vector < int > occurs; \\ FOR(i, 0, (int)t.length()) \ \{ \\ int \ c = t[i]-'a'; \\ curr = aut[curr][c]; \\ if (curr == (int)s.length()) \ \{ \\ occurs.pb(i-s.length()+1); \\ \} \\ return \ occurs; \\ \} \end{array}
```

6.5 Prefix Function

```
\label{eq:continuous_problem} \begin{tabular}{ll} // & problem probl
```

6.6 Suffix Array

```
vector<int> sortCyclicShifts(string const& s) {
    int n = s.size();
    const int alphabet = 256; // we assume to use
          the whole ASCII range
    vector<int> p(n), c(n), cnt(max(alphabet, n),
    for (int i = 0; i < n; i++)
       cnt[s[i]]++;
    for (int i = 1; i < alphabet; i++)
       \operatorname{cnt}[i] += \operatorname{cnt}[i-1];
    for (int i = 0; i < n; i++)
       p[-cnt[s[i]]] = i;
    c[p[0]] = 0;
   int classes = 1;
    for (int i = 1; i < n; i++) {
       if (s[p[i]] != s[p[i-1]])
           classes++;
       c[p[i]] = classes - 1;
    vector < int > pn(n), cn(n);
    for (int h = 0; (1 << h) < n; ++h) {
       for (int i = 0; i < n; i++) {
           pn[i] = p[i] - (1 << h);
           if (pn[i] < 0)
                pn[i] += n;
       fill(cnt.begin(), cnt.begin() + classes, 0);
       for (int i = 0; i < n; i++)
        \cot[c[pn[i]]] ++; 
for (int i = 1; i < classes; i++)
           \operatorname{cnt}[i] += \operatorname{cnt}[i-1];
       for (int i = n-1; i >= 0; i--)
           p[-cnt[c[pn[i]]]] = pn[i];
       \operatorname{cn}[p[0]] = 0;
       classes = 1;
       for (int i = 1; i < n; i++) {
           pair < int, int > cur = \{c[p[i]], c[(p[i] + (1
                   << h)) % n]};
           pair < int, int > prev = \{c[p[i-1]], c[(p[i-1])\}
                   + (1 << h)) \% n];
           if (cur != prev)
                ++classes;
           cn[p[i]] = classes - 1;
       c.swap(cn);
    return p;
vector<int> constructSuffixArray(string s) {
    s += "$"; // <- this must be smaller than any
          character in s
    vector<int> sorted shifts = sortCyclicShifts(s)
```

```
sorted_shifts.erase(sorted_shifts.begin());
return sorted_shifts;
```