## ACM-ICPC TEAM REFERENCE DOCUMENT

Vilnius University (Šimoliūnaitė, Strakšys, Strimaitis)

Contents					
1	Data	a Structures	1		
	1.1	Disjoin Set Union	1		
	1.2	Fenwick 2D	1		
	1.3	Fenwick Tree Point Update And			
		Range Query	1		
	1.4	Fenwick Tree Range Update And			
		Point Query	1		
	1.5	Fenwick Tree Range Update And Range Query	1		
	1.6	Implicit Treap	2		
	1.7	Segment Tree With Lazy Propagation	2		
	1.8	Segment Tree	3		
	1.9	Treap	3		
	1.10	Trie	3		
	1.10		•		
2	Gen	eral	4		
	2.1	Automatic Test	4		
	2.2	Big Integer	4		
	2.3	C++ Template	6		
	2.4	Compilation	6		
	2.5	Ternary Search	6		
3	Cas	metry	7		
3	3.1	0.1.77	7		
	3.2	2d Vector	7		
	$\frac{3.2}{3.3}$	Circle Line Intersection	7		
	3.4	Common Tangents To Two Circles .	7		
	3.5	Convex Hull Gift Wrapping	8		
	3.6	Convex Hull With Graham's Scan .	8		
	3.7	Line	8		
	3.8	Number Of Lattice Points On Segment	8		
	3.9	Pick's Theorem	8		
	3.0		Ü		
4	Gra	phs	8		
	4.1	Bellman Ford Algorithm	8		
	4.2	Bipartite Graph	9		
	4.3	Dfs With Timestamps	10		
	4.4	Finding Articulation Points	10		
	4.5	Finding Bridges	10		
	4.6	Lowest Common Ancestor	10		
	4.7	Max Flow With Dinic 2	11		
	4.8	Max Flow With Dinic	11		
	4.9	Max Flow With Ford Fulkerson	12		
	4.10		12		
	4.11		12		
	4.12	Q	12		
	4.13	Strongly Connected Components	12		

5	Mat	h	13
	5.1	Big Integer Multiplication With FFT	13
	5.2	Burnside's Lemma	13
	5.3	Chinese Remainder Theorem	14
	5.4	Euler Totient Function	14
	5.5	Extended Euclidean Algorithm	14
	5.6	Factorization With Sieve	14
	5.7	FFT With Modulo	14
	5.8	FFT	15
	5.9	Formulas	15
	5.10		15
	5.11		15
	5.12	Simpson Integration	16
6	Strings		
	6.1	Aho Corasick Automaton	16
	6.2	Hashing	16
	6.3	KMP	17
	6.4	Prefix Function Automaton	17
	6.5	Prefix Function	17
	6.6	Suffix Array	17

### 1 Data Structures

## 1.1 Disjoin Set Union

```
struct DSU {
    vector < int > par;
    vector < int > sz;

DSU(int n) {
        FOR(i, 0, n) {
            par.pb(i);
            sz.pb(1);
        }
}

int find(int a) {
        return par[a] = par[a] == a ? a : find(par[a]);
}

bool same(int a, int b) {
        return find(a) == find(b);
}

void unite(int a, int b) {
        a = find(a);
        b = find(b);
        if(sz[a] > sz[b]) swap(a, b);
        sz[b] += sz[a];
        par[a] = b;
};
```

#### 1.2 Fenwick 2D

```
 \begin{array}{l} struct \; Fenwick2D \; \{ \\ vector < vector < ll >> bit; \\ int \; n, \; m; \\ Fenwick2D (int \_n, \; int \_m) \; \{ \\ n = \_n; \; m = \_m; \\ bit = vector < vector < ll >> (n+1, \; vector < ll >(m+1, \; 0)); \\ \} \\ ll \; sum (int \; x, \; int \; y) \; \{ \\ ll \; ret = \; 0; \\ for \; (int \; i = \; x; \; i \; > \; 0; \; i \; -= \; i \; \& \; (-i)) \\ for \; (int \; j = \; y; \; j \; > \; 0; \; j \; -= \; j \; \& \; (-j)) \\ ret \; + = bit[i][j]; \\ return \; ret; \\ \} \\ ll \; sum (int \; x1, \; int \; y1, \; int \; x2, \; int \; y2) \; \{ \\ return \; sum (x2, \; y2) \; - \; sum (x2, \; y1-1) \; - \; sum (x1-1, \; y2) \; + \\ sum (x1-1, \; y1-1); \\ \} \\ void \; add (int \; x, \; int \; y, \; ll \; delta) \; \{ \\ for \; (int \; i = \; x; \; i \; < \; n; \; i \; += \; i \; \& \; (-i)) \\ for \; (int \; j = \; y; \; j \; < \; m; \; j \; += \; j \; \& \; (-j)) \\ bit[i][j] \; += \; delta; \\ \} \\ \}; \end{array}
```

## 1.3 Fenwick Tree Point Update And Range Query

## 1.4 Fenwick Tree Range Update And Point Query

## 1.5 Fenwick Tree Range Update And Range Query

## 1.6 Implicit Treap

```
template <typename T>
struct Node {
Node* l, *r;
    ll prio, size, sum;
    T val:
    bool rev
    Node() {}
Node(T _val) : l(nullptr), r(nullptr), val(_val), size(1),
        sum(_val), rev(false) {
prio = rand() ^ (rand() << 15);
template <typename T>
struct ImplicitTreap {
    typedef Node<T>* NodePtr;
    int sz(NodePtr n) {
        return n ? n->size : 0;
    il getSum(NodePtr n) {
        return n ? n->sum : 0;
    void push(NodePtr n) {
        if (n && n->rev) {
            n->rev = false;
            swap(n->l, n->r);
            if (n->1) n->1->rev ^= 1;
            if (n->r) n->r->rev = 1;
    }
    void recalc(NodePtr n) {
        if (!n) return;
        n->size = sz(n->l) + 1 + sz(n->r);
        n\text{-}{>}sum = getSum(n\text{-}{>}l) + n\text{-}{>}val + getSum(n\text{-}{>}r);
    void split(NodePtr tree, ll key, NodePtr& l, NodePtr& r)
        push(tree);
        if (!tree) {
            l = r = nullptr;
        else if (key \leq sz(tree->l)) {
            split(tree->l, key, l, tree->l);
        else {
            split(tree->r, key-sz(tree->l)-1, tree->r, r);
            l = tree;
        recalc(tree);
```

```
void merge(NodePtr& tree, NodePtr l, NodePtr r) {
         push(l); push(r);
        if (!l || !r) {
tree = l ? l : r;
        else if (l->prio > r->prio) {
             merge(l{-}{>}r,\; l{-}{>}r,\; r);
             tree = 1;
        else {
             merge(r->l, l, r->l);
        recalc(tree);
    }
    void insert(NodePtr& tree, T val, int pos) {
        if (!tree) {
             tree = new Node<T>(val);
             return;
        NodePtr L, R;
        split(tree, pos, L, R);
        merge(L, L, new Node<T>(val));
        merge(tree, L, R);
        recalc(tree);
    void reverse(NodePtr tree, int l, int r) {
        NodePtr t1, t2, t3;
        split(tree, l, t1, t2);
         split(t2, r - l + 1, t2, t3);
        if(t2) t2->rev = true;
        merge(t2, t1, t2);
        merge(tree, t2, t3);
    void print(NodePtr t, bool newline = true) {
        push(t);
        if (!t) return;
        \begin{array}{l} \operatorname{print}(t->l,\,\operatorname{false});\\ \operatorname{cout} << t->\operatorname{val} << "{\scriptscriptstyle \sqcup}"; \end{array}
        print(t->r, false);
        if (newline) cout << endl;
    NodePtr fromArray(vector<T> v) {
        NodePtr t = nullptr;
         FOR(i, 0, (int)v.size()) {
             insert(t, v[i], i);
        return t;
    }
    ll calcSum(NodePtr t, int l, int r) {
         NodePtr L, R;
         split(t, l, L, R);
        NodePtr good;
split(R, r - l + 1, good, L);
        return getSum(good);
};
/* Usage: ImplicitTreap<int> t;
Node<int> tree = t.fromArray(someVector); t.reverse(tree, l,
      r); ...
```

# 1.7 Segment Tree With Lazy Propagation

```
build(arr, 1, 0, n-1); // same as in simple SegmentTree
   lazy[v*2+1] += lazy[v];
       lazy[v] = 0;
   void update(int v, int tl, int tr, int l, int r, ll addend) {
       if (l > r)
          return;
       if (l == tl^{'} \&\& tr == r) {
           t[v] += addend;
          lazy[v] += addend;
       } else {
           push(v);
           int tm = (tl + tr) / 2;
           update(v*2, tl, tm, l, min(r, tm), addend);
           update(v*2+1, tm+1, tr, max(l, tm+1), r, addend);
           t[v] = max(t[v*2], t[v*2+1]);
       }
   }
   int query(int v, int tl, int tr, int l, int r) {
           return -OO;
       if (tl == tr)
          return t[v];
       push(v);
       int tm = (tl + tr) / 2;
       return max(query(v*2, tl, tm, l, min(r, tm)),
              query(v^*2{+}1,\; tm{+}1,\; tr,\; max(l,\; tm{+}1),\; r));
};
```

## 1.8 Segment Tree

```
struct SegmentTree {
   int n;
   vector<|l> t:
   const ll IDENTITY = 0; // OO for min, -OO for max, ...
   ll f(ll a, ll b) {
       return a+b;
   SegmentTree(int _n) {
       n = n; t = vector < ll > (4*n, IDENTITY);
   SegmentTree(vector<ll>& arr) {
       n = arr.size(); t = vector < ll > (4*n, IDENTITY);
       build(arr, 1, 0, n-1);
   void build(vector<ll>& arr, int v, int tl, int tr) {
       if(tl == tr) \ \{ \ t[v] = arr[tl]; \ \}
       else {
          int tm = (tl+tr)/2;
           build(arr, 2*v, tl, tm);
build(arr, 2*v+1, tm+1, tr);
          t[v] = f(t[2*v], t[2*v+1]);
       }
    // sum(1, 0, n-1, l, r)
   ll sum(int v, int tl, int tr, int l, int r) {
       if(l > r) return IDENTITY;
       tm+1, tr, max(l, tm+1), r));
   // update(1, 0, n-1, i, v)
   void update(int v, int tl, int tr, int pos, ll newVal) {
       if(tl == tr) \ \{ \ t[v] = newVal; \ \}
       else {
           int tm = (tl+tr)/2;
           if(pos <= tm) update(2*v, tl, tm, pos, newVal);
           else update(2*v+1, tm+1, tr, pos, newVal);
           t[v] = f(t[2*v],t[2*v+1]);
       }
   }
};
```

### 1.9 Treap

```
namespace Treap {
    struct Node {
Node *l, *r;
        ll key, prio, size;
Node() {}
        Node(ll key) : key(key), l(nullptr), r(nullptr), size(1) { prio = rand() \widehat{} (rand() << 15);
    };
    typedef Node* NodePtr;
    int\ sz(NodePtr\ n)\ \{
        return n ? n->size : 0;
    void recalc(NodePtr n) {
        if (!n) return;
        n-size = sz(n-sl) + 1 + sz(n-sr); // add more
              operations here as needed
    }
    void split(NodePtr tree, ll key, NodePtr& l, NodePtr& r)
        if (!tree) {
            l = r = nullptr;
        else if (key < tree->key) {
            split(tree->l, key, l, tree->l);
            r = tree:
        else {
            split(tree->r, key, tree->r, r);
            l = tree;
        recalc(tree);
    void merge(NodePtr& tree, NodePtr l, NodePtr r) {
        if (!l || !r) {
            \text{tree} \stackrel{\cdot}{=} \stackrel{\cdot}{l} ? l : r;
        else if (l->prio > r->prio) {
            merge(l->r, l->r, r);
            tree = 1;
        else {
            merge(r->l,\;l,\;r->l);
            tree = r:
        recalc(tree);
    void insert(NodePtr& tree, NodePtr node) {
        if (!tree) {
            tree = node;
        else if (node->prio > tree->prio) {
            {\rm split}({\rm tree,\ node\text{-}>}{\rm key},\ {\rm node\text{-}>}{\rm l},\ {\rm node\text{-}>}{\rm r});
            tree = node:
        else {
            insert(node->key < tree->key ? tree->l : tree->r,
        recalc(tree);
    void erase(NodePtr tree, ll key) {
        if (!tree) return;
        if (tree->key == key) {
            merge(tree, tree->l, tree->r);
            erase(key < tree->key ? tree->l : tree->r, key);
        recalc(tree);
    void print(NodePtr t, bool newline = true) {
        if (!t) return;
        print(t->l, false);
        cout << t->key << "";
```

```
\begin{array}{c} \operatorname{print}(t\text{--}r,\,\operatorname{false});\\ \operatorname{if}\,\left(\operatorname{newline}\right)\,\operatorname{cout}<<\operatorname{endl};\\ \end{array}\}
```

#### 1.10 Trie

```
struct Trie {
   const int ALPHA = 26;
   const char BASE = 'a';
    vector<vector<int>> nextNode;
    vector<int> mark;
   int nodeCount:
   Trie() {
       nextNode = vector<vector<int>>(MAXN, vector<int
             >(ALPHA, -1));
       mark = vector < int > (MAXN, -1);
       {\bf nodeCount}=1;
   void insert(const string& s, int id) {
       int curr = 0:
       FOR(i, 0, (int)s.length()) {
           c = s[i] - BASE;

c = s[i] - BASE;

c = s[i] - BASE;
               nextNode[curr][c] = nodeCount++;
           curr = nextNode[curr][c];
       mark[curr] = id;
   bool exists
(const string& s) {
       int curr = 0;
       FOR(i, 0, (int)s.length()) {
           int c = s[i] - BASE;

if(nextNode[curr][c] == -1) return false;
           curr = nextNode[curr][c];
       return mark[curr] != -1;
};
```

#### 2 General

#### 2.1 Automatic Test

```
# Linux Bash
# gen, main and stupid have to be compiled beforehand
for((i=1;;++i)); do
    echo $i;
    ./gen $i > genIn;
    diff <(./main < genIn) <(./stupid < genIn) || break;
done

# Windows CMD
@echo off
FOR /L %%I IN (1,1,2147483647) DO (
    echo %%I
    gen.exe %%I > genIn
    main.exe < genIn > mainOut
    stupid.exe < genIn > stupidOut
    FC mainOut stupidOut || goto :eof
)
```

#### 2.2 Big Integer

```
const int base = 1000000000;
const int base_digits = 9;
struct bigint {
  vector<int> a;
  int sign;
  int size() {
```

```
if (a.empty()) return 0;
     int ans = (a.size() - 1)* base_digits;
     int ca = a.back();
     while (ca) ans++, ca \neq 10;
     return ans:
bigint operator^(const bigint &v) {
     while (!y.isZero()) {
    if (y % 2) ans *= x;
        x *= x, y /= 2;
     return ans;
string to_string() {
     stringstream ss;
ss << *this;
     string s;
     ss >> s;
     return s;
int sumof() {
     string s = to_string();
int ans = 0;
     for (auto c : s) ans += c - 0;
     return ans;
bigint(): sign(1) \{ \}
bigint(long long v) {
    *this = v;
bigint(const string &s) {
    read(s);
void operator=(const bigint &v) {
     sign = v.sign;
     a = v.a;
void operator=(long long v) {
     sign = 1;
     a.clear();
     if (v < 0)
          sign = -1, v = -v;
     for (; v > 0; v = v / base)
a.push_back(v % base);
bigint operator+(const bigint &v) const {
     if (sign == v.sign) {
          bigint res = v;
          for (int i = 0, carry = 0; i < (int)max(a.size(), v.a.
                 size()) || carry; ++i) {
               if\ (i == (int)res.a.size()) \ res.a.push\_back(0); \\
               \begin{array}{l} res.a[i] += carry + (i < (int)a.size() ? \ a[i] : 0); \\ carry = res.a[i] >= base; \\ if \ (carry) \ res.a[i] -= base; \\ \end{array} 
          return res:
     return *this - (-v);
bigint operator-(const bigint &v) const {
     if (sign == v.sign) {
          if(abs() >= v.abs()) {
               bigint res = *this;
               for (int i = 0, carry = 0; i < (int)v.a.size() |
                   carry; ++i) {
res.a[i] -= carry + (i < (int)v.a.size() ? v.a[i] : 0);
                    carry = res.a[i] < 0;
                    if (carry) res.a[i] += base;
               res.trim();
               return res:
          return -(v - *this);
     return *this + (-v);
void operator*=(int v) {
     if (v < 0) sign = -sign, v = -v;
     for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i)
         \label{eq:continuous_state} \begin{array}{l} \overset{1}{\text{if }} i = (int)a.size()) \; a.push\_back(0); \\ long \; long \; cur = a[i] * (long \; long)v + carry; \\ carry = (int)(cur / base); \\ a[i] = (int)(cur \% \; base); \end{array}
```

```
trim();
bigint operator*(int v) const {
   bigint res = *this;
     res *= v;
     return res;
void operator*=(long long v) {
    if (v < 0) sign = -sign, v = -v;
for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i)
           \begin{array}{l} i \ \ (i == (int)a.size()) \ a.push\_back(0); \\ long \ long \ cur = a[i] \ * (long \ long)v + carry; \\ carry = (int)(cur \ / \ base); \\ a[i] = (int)(cur \ \% \ base); \end{array} 
     trim();
bigint operator*(long long v) const {
   bigint res = *this;
   res *= v;
     return res;
friend pair<br/>bigint, bigint> divmod(const bigint &a1,
        const bigint &b1) {
     \begin{array}{ll} \mathrm{int\ norm} = \mathrm{base\ /\ (b1.a.back()\ +\ 1);} \\ \mathrm{bigint\ a} = \mathrm{a1.abs()\ *\ norm;} \\ \mathrm{bigint\ b} = \mathrm{b1.abs()\ *\ norm;} \\ \end{array}
     bigint q, r;
     q.a.resize(a.a.size());
     for (int i = a.a.size() - 1; i >= 0; i--) {
          r *= base;
          r += a.a[i];
          int s1 = r.a.size() \le b.a.size() ? 0 : r.a[b.a.size()];
          int s2 = r.a.size() \le b.a.size() - 1 ? 0 : r.a[b.a.size()]
                 () - 1];
          int d = ((long long)base * s1 + s2) / b.a.back(); r -= b * d;
          q.sign = a1.sign * b1.sign;
     r.sign = a1.sign;
     q.trim();
     r.trim();
     return make_pair(q, r / norm);
bigint operator/(const bigint &v) const {
     return divmod(*this, v).first;
bigint operator%(const bigint &v) const {
     return divmod(*this, v).second;
void operator/=(int v) {
     if (v < 0) sign = -sign, v = -v;
     for (int i = (int)a.size() - 1, rem = 0; i >= 0; --i) {
          long long cur = a[i] + rem * (long long)base;
a[i] = (int)(cur / v);
rem = (int)(cur % v);
     trim();
bigint operator/(int v) const {
     bigint res = *this;
     res /= v;
return res;
int operator%(int v) const {
     if(v < 0) v = -v;
     int m = 0;
      \begin{array}{ll} {\rm for~(int~i=a.size()~-1;~i>=0;--i)} \\ {\rm m=(a[i]+m~*~(long~long)base)~\%~v;} \\ {\rm return~m~*sign;} \end{array} 
void operator+=(const bigint &v) {
     *this = *this + v;
void operator=(const bigint &v) {
     *this = *this - v;
void operator*=(const bigint &v) {
    *this = *this * v;
void operator/=(const bigint &v) {
    *this = *this / v;
```

```
bool operator < (const bigint &v) const {
    if (sign != v.sign) return sign < v.sign;
if (a.size() != v.a.size())</pre>
    \begin{array}{ll} \text{If } (\text{a.size}():=\text{v.a.size}()) \\ \text{return a.size}() * \text{sign} < \text{v.a.size}() * \text{v.sign}; \\ \text{for } (\text{int } i = \text{a.size}() - 1; i >= 0; i--) \end{array}
        \begin{array}{l} \text{if } (a[i] \mathrel{!=} v.a[i]) \\ \text{return } a[i] \mathrel{*} sign < v.a[i] \mathrel{*} sign; \end{array}
    return false;
bool operator>(const bigint &v) const {
    return v < *this;
bool operator<=(const bigint &v) const {
    return !(v < *this);
bool operator>=(const bigint &v) const {
    return !(*this < v);
bool operator==(const bigint &v) const {
    return !(*this < v) \&\& !(v < *this);
bool operator!=(const bigint &v) const {
    return *this < v || v < *this;
     while (!a.empty() && !a.back()) a.pop_back();
    if (a.empty()) sign = 1;
bool isZero() const {
    return a.empty() || (a.size() == 1 && !a[0]);
bigint operator-() const {
    bigint res = *this;
res.sign = -sign;
    return res;
bigint abs() const {
    bigint res = *this;
res.sign *= res.sign;
    return res;
long long longValue() const {
     long long res = 0;
    for (int i = a.size() - 1; i >= 0; i--) res = res * base +
           a[i];
    return res * sign;
friend bigint gcd(const bigint &a, const bigint &b) {
    return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint &b) {
    return a / \gcd(a, b) * b;
void read(const string &s) {
    sign = 1;
    while (pos < (int)s.size() && (s[pos] == '-' || s[pos] == '+')) {
  if (s[pos] == '-') sign = -sign;
    for (int i = s.size() - 1; i \ge pos; i -= base\_digits) {
         for (int j = max(pos, i - base_digits + 1); j <= i; j
             x = x * 10 + s[j] - 0;
         a.push_back(x);
    trim();
friend istream & operator >> (istream & stream, bigint & v)
    string s;
    stream >> s;
    v.read(s);
    return stream:
friend ostream & operator << (ostream & stream, const
       {\rm bigint}\ \&v)\ \{
    if (v.sign == -1) stream << '-'; stream << (v.a.empty() ? 0 : v.a.back());
    for (int i = (int)v.a.size() - 2; i >= 0; --i)
         stream << setw(base\_digits) << setfill('0') << v.
                a[i];
```

```
static vector<int> convert_base(const vector<int> &a,
           int old_digits, int new_digits) {
vector<long long> p(max(old_digits, new_digits) +
           p[0] = 1;
          p[i] - 1;
for (int i = 1; i < (int)p.size(); i++)
p[i] = p[i - 1] * 10;
vector<int> res;
long long cur = 0;
           int cur\_digits = 0;
          int cur_digits = 0;
for (int i = 0; i < (int)a.size(); i++) {
    cur += a[i] * p[cur_digits];
    cur_digits += old_digits;
    while (cur_digits >= new_digits) {
        res.push_back(int(cur % p[new_digits]));
        cur /= p[new_digits];
    }
                      cur_digits -= new_digits;
                }
          res.push_back((int)cur);
while (!res.empty() && !res.back()) res.pop_back();
           return res:
     typedef vector<long long> vll;
     static vll karatsubaMultiply(const vll &a, const vll &b) {
          int n = a.size();
           vll res(n + n);
          if (n \le 32) {
for (int i = 0; i < n; i++)
                     for (int j = 0; j < n; j++)

res[i + j] += a[i] * b[j];
                return res;
           int k = n \gg 1:
          vll a1(a.begin(), a.begin() + k);
vll a2(a.begin() + k, a.end());
vll b1(b.begin(), b.begin() + k);
           vll b2(b.begin() + k, b.end());
           vll a1b1 = karatsubaMultiply(a1, b1):
           vll a2b2 = karatsubaMultiply(a2, b2);
           for \; (int \; i = 0; \; i < k; \; i{+}{+}) \; a2[i] \; {+}{=} \; a1[i];
          for (int i = 0; i < k; i++) b2[i] += b1[i];
           vll r = karatsubaMultiply(a2, b2):
          for (int i = 0; i < (int)a1b1.size(); i++) r[i] -= a1b1[i];
           for (int i = 0; i < (int)a2b2.size(); i++) r[i] -= a2b2[i];
          \begin{array}{l} for \; (int \; i=0; \; i < (int)r.size(); \; i++) \; res[i+k] \; += r[i]; \\ for \; (int \; i=0; \; i < (int)alb1.size(); \; i++) \; res[i] \; += \; alb1 \end{array}
           for (int i = 0; i < (int)a2b2.size(); i++) res[i + n] +=
                   a2b2[i];
           return res:
     bigint operator*(const bigint &v) const {
           vector<int> a6 = convert_base(this->a, base_digits,
                   6);
           vector < int > b6 = convert_base(v.a, base_digits, 6);
          vll x(a6.begin(), a6.end());
vll y(b6.begin(), b6.end());
          \label{eq:while (x.size() < y.size()) x.push_back(0);} while (y.size() < x.size()) y.push_back(0); while (x.size() & (x.size() - 1)) x.push_back(0), y. \\
                   push_back(0);
           vll c = karatsubaMultiply(x, y);
          bigint res;
           res.sign = sign * v.sign;
          for (int i = 0, carry = 0; i < (int)c.size(); i++) { long long cur = c[i] + carry; res.a.push_back((int)(cur % 1000000));
                carry = (int)(cur / 1000000);
           res.a = convert_base(res.a, 6, base_digits);
           res.trim();
           return res:
};
```

## 2.3 C++ Template

```
#include <bits/stdc++.h>
 #include <ext/pb_ds/assoc_container.hpp> //
gp_hash_table<int, int> == hash map
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
 using namespace ___gnu_pbds;
 typedef long long ll;
typedef unsigned long long ull;
typedef long double ld;
typedef pair<int, int> pii;
 typedef pair<ll, ll> pll;
 typedef pair<double, double> pdd;
 template <typename T> using min_heap = priority_queue<
T, vector<T>, greater<T>>;
template <typename T> using max_heap = priority_queue<
T, vector<T>, less<T>>;
template <typename T> using ordered_set = tree<T,
              null_type, less<T>, rb_tree_tag,
               tree_order_statistics_node_update>;
 template <typename K, typename V> using hashmap =
              gp\_hash\_table{<}K,\;V{>};
template<typename A, typename B> ostream& operator<<( ostream& out, pair<A, B> p) { out << "(" << p.first << ",\square" << p.second << ")"; return out;}
set <T> v) { out << "{"; for (auto& x : v) out << x << ",..."; out << "}"; return out; }
template<typename K, typename V> ostream& operator<<( ostream& out, map<K, V> m) { out << "{"; for(auto& e:m) out << e.first << "_->_" << e.second << ",_"; out << "}"; return out; }
template<typename K, typename V> ostream& operator<<(
              ostream& out, hashmap< K, V> m) { out << "{"; for( auto& e: m) out << e.first << "_{-}>_{-}" << e.second << ",_{-}"; out << "}"; return out; }
 #define FAST IO ios base::sync with stdio(false); cin.tie(
 #define TESTS(t) int NUMBER_OF_TESTS; cin >>
              NUMBER_OF_TESTS; for(int t = 1; t <= NUMBER_OF_TESTS; t++)
 #define FOR(i, begin, end) for (int i = (begin) - ((begin) > (end)); i != (end) - ((begin) > (end)); i += 1 - 2 * ((begin) > (end)))
 #define sgn(a) ((a) > eps ? 1 : ((a) < -eps ? -1 : 0))
#define debug(x) cerr << ">^{\circ}' < #x << "^{\sqcup}" << x <<
                endl:
 #define pb push_back
 \#define \ \bar{r}nd(\bar{a}, \ b) \ (uniform\_int\_distribution < int > ((a), \ (b))(a)
              rng))
 #ifndef LOCAL
         #define cerr if(0)cout
         #define endl "\n'
 #endif
\verb|mt19937 rng| (chrono::steady\_clock::now().time\_since\_epoch()|\\
             .count());
 \begin{array}{lll} clock\_t & \underline{clock}\_; \\ void \ startTime() \ \{\underline{\phantom{clock}}\_ = clock(); \} \\ void \ timeit(string \ msg) \ \{cerr <<">>\_" << msg << ":\_" << msg << ":_"" << msg << ":_" << constraints of the precise of the precis
 clock_t ___clock
const ld eps = 1e-14;
const int oo = 2e9;
const ll OO = 2e18;
const ll MOD = 1000000007;
const int MAXN = 1000000;
int main() {
    FAST_IO;
         startTime();
        timeit("Finished"):
        return 0:
```

### 2.4 Compilation

```
# Simple compile
g++ -DLOCAL -O2 -o main.exe -std-c++17 -Wall -Wno-
unused-result -Wshadow main.cpp
# Debug
g++ -DLOCAL -std=c++17 -Wshadow -Wall -o main.exe
main.cpp -fsanitize=address -fsanitize=undefined -fuse-
ld=gold -D_GLIBCXX_DEBUG -g
```

### 2.5 Ternary Search

## 3 Geometry

#### 3.1 2d Vector

```
template <typename T>
struct Vec {
    T x, y;
     \begin{array}{l} {\rm Vec}() \colon x(0), \, y(0) \, \, \{\} \\ {\rm Vec}(T \, \_x, \, T \, \_y) \colon x(\_x), \, y(\_y) \, \, \{\} \\ {\rm Vec \, \, operator+(const \, Vec\& \, b) \, \, \{} \end{array} 
        return Vec < T > (x+b.x, y+b.y);
    Vec operator-(const Vec& b) {
        \operatorname{return} \ \operatorname{Vec} < T > (x-b.x, \ y-b.y);
    Vec operator*(T c) {
        return Vec(x*c, y*c);
    T operator*(const Vec& b) {
        return x*b.x + y*b.y;
    T operator^(const Vec& b) {
        return x*b.y-y*b.x;
    bool operator<(const Vec& other) const {
        if(x == other.x) return y < other.y;
        return x < other.x;
    bool operator==(const Vec& other) const {
        return x==other.x && y==other.y;
    bool operator!=(const Vec& other) const {
        return !(*this == other);
    friend ostream& operator << (ostream& out, const Vec& v)
        return out << "(" << v.x << "," << v.y << ")";
    friend istream& operator>>(istream& in, Vec<T>& v) {
        return in >> v.x >> v.y;
    T norm() { // squared length
return (*this)*(*this);
    ld len() {
        return sqrt(norm());
    ld angle(const Vec& other) { // angle between this and
        return acosl((*this)*other/len()/other.len());
```

```
}
Vec perp() {
    return Vec(-y, x);
}
};
/* Cross product of 3d vectors: (ay*bz-az*by, az*bx-ax*bz, ax
    *by-ay*bx)
*/
```

#### 3.2 Circle Circle Intersection

Let's say that the first circle is centered at (0,0) (if it's not, we can move the origin to the center of the first circle and adjust the coordinates), and the second one is at  $(x_2, y_2)$ . Then, let's construct a line Ax + By + C = 0, where  $A = -2x_2, B = -2y_2, C = x_2^2 + y_2^2 + r_1^2 - r_2^2$ . Finding the intersection between this line and the first circle will give us the answer. The only tricky case: if both circles are centered at the same point. We handle this case separately.

#### 3.3 Circle Line Intersection

```
double r, a, b, c; // ax+by+c=0, radius is at (0, 0) // If the center is not at (0, 0), fix the constant c to translate everything so that center is at (0, 0) double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b); if (c*c > r*r*(a*a+b*b)+eps) puts ("no_points"); else if (abs (c*c - r*r*(a*a+b*b)) < eps) { puts ("1_point"); cout << x0 << '_' << y0 << '\n'; } else { double d = r*r - c*c/(a*a+b*b); double mult = sqrt (d / (a*a+b*b)); double ax, ay, bx, by; ax = x0 + b * mult; bx = x0 - b * mult; by = y0 - a * mult; by = y0 + a * mult; puts ("2_points"); cout << ax << '_' << ay << '\n'; << by << '\n'; << by
```

## 3.4 Common Tangents To Two Circles

```
struct pt {
    double x, y;

    pt operator- (pt p) {
        pt res = { x-p.x, y-p.y };
        return res;
    }
};
struct circle : pt {
    double r;
};
struct line {
    double a, b, c;
};
void tangents (pt c, double r1, double r2, vector<line> & ans)
    {
        double z = sqr(c.x) + sqr(c.y);
        double d = z - sqr(r);
        if (d < -eps) return;
        d = sqrt (abs (d));
    }
}</pre>
```

```
line l;  
l.a = (c.x * r + c.y * d) / z;  
l.b = (c.y * r - c.x * d) / z;  
l.c = r1;  
ans.push_back (l);  
} vectorline> tangents (circle a, circle b) { vectorline> ans; for (int i=-1; i<=1; i+=2)  
for (int j=-1; j<=1; j+=2)  
tangents (b-a, a.r*i, b.r*j, ans);  
for (size_t i=0; i<ans.size(); ++i)  
ans[i].c -= ans[i].a * a.x + ans[i].b * a.y;  
return ans;  
}
```

## 3.5 Convex Hull Gift Wrapping

```
vector<Vec<int>> buildConvexHull(vector<Vec<int>>&
     pts) {
   int n = pts.size();
   sort(pts.begin(), pts.end());
auto currP = pts[0]; // choose some extreme point to be
         on the hull
   vector<Vec<int>> hull;
   set < Vec < int >> used;
   hull.pb(pts[0]);
   used.insert(pts[0]);
   while(true) {
       auto candidate = pts[0]; // choose some point to be a
             candidate
       auto currDir = candidate-currP;
        vector<Vec<int>> toUpdate;
       FOR(i, 0, n) {
   if(currP == pts[i]) continue;
           // currently we have currP->candidate
            // we need to find point to the left of this
           auto possibleNext = pts[i];
           auto nextDir = possibleNext - currP;
auto cross = currDir ^ nextDir;
if(candidate == currP || cross > 0) {
                candidate = possibleNext;
                currDir = nextDir;
            } else if(cross == 0 && nextDir.norm() > currDir.
                 norm()) {
                candidate = possibleNext;
               currDir = nextDir;
       if(used.find(candidate) != used.end()) break;
       hull.pb(candidate)
       used.insert(candidate):
       currP = candidate;
   return hull;
```

## 3.6 Convex Hull With Graham's Scan

```
// Takes in >= 3 points
// Returns convex hull in clockwise order
// Ignores points on the border
vector<Vec<int>> buildConvexHull(vector<Vec<int>> pts)
{
    if(pts.size() <= 3) return pts;
    sort(pts.begin(), pts.end());
    stack<Vec<int>> hull;
    hull.push(pts[0]);
    auto p = pts[0];
    sort(pts.begin()+1, pts.end(), [&](Vec<int> a, Vec<int>
        b) -> bool {
        // p->a-> b is a ccw turn
        int turn = sgn((a-p)^(b-a));
        //if(turn == 0) return (a-p).norm() > (b-p).norm();
```

```
^ among collinear points, take the farthest one
    return turn == 1;
});
hull.push(pts[1]);
FOR(i, 2, (int)pts.size()) { auto c = pts[i];
    if(c == hull.top()) continue;
    while(true) {
        auto a = hull.top(); hull.pop();
auto b = hull.top();
        auto ba = a-b;
        auto ac = c-a;
        if((ba^ac) > 0)
            hull.push(a);
            break;
        } else if((ba^ac) == 0) {
 if(ba*ac < 0) c = a;
                c is between b and a, so it shouldn't be
                   added to the hull
            break;
        }
    hull.push(c);
vector<Vec<int>> hullPts;
while(!hull.empty())  {
    hullPts.pb(hull.top());
    hull.pop();
return hullPts;
```

#### 3.7 Line

```
 \begin{array}{l} template < typename \ T> \\ struct \ Line \{ \ // \ expressed \ as \ two \ vectors \\ Vec < T> \ start, \ dir; \\ Line() \ \{ \} \\ Line(Vec < T> \ a, \ Vec < T> \ b): \ start(a), \ dir(b-a) \ \{ \} \\ \\ Vec < ld> \ intersect(Line \ l) \ \{ \\ \ ld \ t = \ ld((l.start-start)^l.dir)/(dir^l.dir); \\ \ // \ For \ segment-segment \ intersection \ this \ should \ be \ in \ range \ [0, \ l] \\ Vec < \ ld> \ res(start.x, \ start.y); \\ Vec < \ ld> \ dirld(dir.x, \ dir.y); \\ \ return \ res + \ dirld^*t; \\ \} \\ \}; \end{array}
```

## 3.8 Number Of Lattice Points On Segment

Let's say we have a line segment from  $(x_1, y_1)$  to  $(x_2, y_2)$ . Then, the number of lattice points on this segment is given by

$$gcd(x_2 - x_1, y_2 - y_1) + 1.$$

#### 3.9 Pick's Theorem

We are given a lattice polygon with non-zero area. Let's denote its area by S, the number of points with integer coordinates lying strictly inside the polygon by I and the number of points lying on the sides of the polygon by B. Then:

$$S = I + \frac{B}{2} - 1.$$

## 4 Graphs

## 4.1 Bellman Ford Algorithm

```
struct Edge
    int a, b, cost;
};
int n, m, v; // v - starting vertex
vector<Edge> e;
/* Finds SSSP with negative edge weights.
* Possible optimization: check if anything changed in a
       relaxation step. If not - you can break early.
 * To find a negative cycle: perform one more relaxation step.
       If anything changes - a negative cycle exists.
void solve() {
    vector <int> d (n, oo);
    d[v]=0;
    for (int i=0; i< n-1; ++i)
        for (int j=0; j<m; ++j)
            if (d[e[j].a] < oo)

d[e[j].b] = min (d[e[j].b], d[e[j].a] + e[j].cost);
    // display d, for example, on the screen
```

## 4.2 Bipartite Graph

```
class BipartiteGraph {
private:
     {\tt vector}{<} {\tt int}{\gt} \_{\tt left}, \_{\tt right};
     vector < vector < int >> \_adjList;
     vector<int> _matchR, _matchL;
     vector<bool> used;
     bool _kuhn(int v) {
          if (_used[v]) return false;
             used[v] = true;
          FOR(i, 0, (int)_adjList[v].size()) {
                \begin{array}{l} \mathrm{int} \ to = \_\mathrm{adjList[v][i]} - \_\mathrm{left.size}(); \\ \mathrm{if} \ (\_\mathrm{matchR[to]} == -1 \ || \ \_\mathrm{kuhn}(\_\mathrm{matchR[to]})) \ \{ \\ \ \_\mathrm{matchR[to]} = v; \\ \end{array} 
                     _matchL[v] = to;
                     return true;
               }
          return false:
              _addReverseEdges() {
          FOR(i, 0, (int)_right.size()) {
               if (_matchR[i] != -1) {
                    \_adjList[\_left.size() + i].pb(\_matchR[i]);
               }
          }
             _dfs(int p) {
     void
          if (_used[p]) return;
          _used[p] = true;
for (auto x : _adjList[p]) {
                _{dfs(x)}
     vector<pii> _buildMM() {
          vector<pair<int, int> > res
          \begin{aligned} FOR(i,\,0,\,(int)\_right.size()) \; \{ \\ if \; (\_matchR[i] \; != -1) \; \{ \end{aligned}
                    res.push\_back(make\_pair(\_matchR[i],\ i));
               }
          }
          return res;
public:
     void addLeft(int x) {
            _{\rm left.pb(x)};
          \_adjList.pb(\{\});
```

```
_{\text{matchL.pb}(-1)};
    _used.pb(false);
void\ addRight(int\ x)\ \{
    _{\rm right.\widetilde{pb}(x)};
    \underline{\underline{\text{adjList.pb}}}(\{\});
     \underline{\phantom{a}} matchR.pb(-1);
     _used.pb(false);
void addForwardEdge(int l, int r) {
     \_adjList[l].pb(r + \_left.size());
void addMatchEdge(int l, int r) {
    if(l != -1) _{matchL[l]} = r;
    if(r != -1) _{matchR[r]} = l;
// Maximum Matching
vector<pii> mm() {
    _matchR = vector<int>(_right.size(), -1);
_matchL = vector<int>(_left.size(), -1);
         these two can be deleted if performing MM on
    already partially matched graph
_used = vector<br/>bool>(_left.size() + _right.size(),
           false);
    bool path_found;
    \begin{array}{c} \text{do } \{ \\ \text{fill(\_used.begin(), \_used.end(), false);} \end{array}
         path_found = false;
         FOR(i, 0, (int)_left.size()) {
    if (_matchL[i] < 0 && !_used[i]) {
                  path\_found \mid = \_kuhn(i);
    } while (path_found);
    return _buildMM();
// Minimum Edge Cover
// Algo: Find MM, add unmatched vertices greedily.
vector<pii> mec()
    auto ans = mm();
    FOR(i, 0, (int)\_left.size()) {
         if (_matchL[i] != -1) {
             for (auto x : _adjList[i]) {
    int ridx = x - _left.size();
    if (_matchR[ridx] == -1) {
                      ans.pb({ i, ridx });
                       _matchR[ridx] = i;
             }
         }
     FOR(i, 0, (int)_left.size()) {
         if (\underline{matchL[i]} == -1 \&\& (int)\underline{adjList[i].size()} >
               nt ridx = _adjList[i][0] - _left.size();
_matchL[i] = ridx;
             int ridx =
             ans.pb(\{i, ridx\});
         }
    return ans;
}
   Minimum Vertex Cover
// Algo: Find MM. Run DFS from unmatched vertices
       from the left part.
    MVC is composed of unvisited LEFT and visited
       RIGHT vertices.
pair<vector<int>, vector<int>> mvc(bool runMM =
       true) {
    if (runMM) mm();
      _addReverseEdges();
    fill(_used.begin(), _used.end(), false);
FOR(i, 0, (int)_left.size()) {
         if (\underline{\phantom{a}}matchL[i] == -1) \{
             __dfs(i);
     vector<int> left, right;
    FOR(i, 0, (int)\_left.size()) {
         if (!\_used[i]) left.pb(i);\\
     FOR(i, 0, (int)\_right.size())  {
```

```
if (\_used[i + (int)\_left.size()]) right.pb(i);
        return { left,right };
    }
      Maximal Independent Vertex Set
    // Algo: Find complement of MVC.
    pair<vector<int>, vector<int>> mivs(bool runMM =
          true) {
        auto m = mvc(runMM);
        vector<br/>bool> containsL(_left.size(), false), containsR(
               _right.size(), false);
        for (auto x : m.first) containsL[x] = true;
        for (auto x : m.second) contains R[x] = true;
        vector<int> left, right;
        FOR(i, 0, (int)_left.size())
           if (!containsL[i]) left.pb(i);
       FOR(i, 0, (int)\_right.size())  { if (!containsR[i]) right.pb(i);
        return { left, right };
};
```

## 4.3 Dfs With Timestamps

```
\label{eq:control_vector} \begin{split} & \operatorname{vector} < \operatorname{vector} < \operatorname{int} > \operatorname{tIn}, \ \operatorname{tOut}, \ \operatorname{color}; \\ & \operatorname{vector} < \operatorname{int} > \operatorname{tIn}, \ \operatorname{tOut}, \ \operatorname{color}; \\ & \operatorname{tinf} = 0; \\ & \operatorname{void} \ \operatorname{dfs}(\operatorname{int} \ v) \ \{ \\ & \operatorname{tIn[v]} = \operatorname{dfs\_timer} + +; \\ & \operatorname{color[v]} = 1; \\ & \operatorname{for} \ (\operatorname{int} \ u : \operatorname{adj[v]}) \\ & \operatorname{if} \ (\operatorname{color[u]} = 0) \\ & \operatorname{dfs}(u); \\ & \operatorname{color[v]} = 2; \\ & \operatorname{tOut[v]} = \operatorname{dfs\_timer} + +; \\ \} \end{split}
```

## 4.4 Finding Articulation Points

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector < bool > visited;
vector<int> tin, fup;
int timer:
void processCutpoint(int v) {
    // problem-specific logic goes here
    // it can be called multiple times for the same v
void dfs(int v, int p = -1) {
    visited[v] = true;

tin[v] = fup[v] = timer++;
    int children=0;
    for (int to : adj[v]) {
        if (to == p) continue; if (visited[to]) {
            fup[v] = min(fup[v], tin[to]);
        } else {
            dfs(to, v);
fup[v] = min(fup[v], fup[to]);
if (fup[to] >= tin[v] && p!=-1)
                processCutpoint(v);
             ++children;
    if(p == -1 \&\& children > 1)
        processCutpoint(v);
}
void findCutpoints() {
    timer = 0;
```

```
\label{eq:visited.assign(n, false);} \begin{aligned} & visited.assign(n, false); \\ & tin.assign(n, -1); \\ & fup.assign(n, -1); \\ & for \; (int \; i = 0; \; i < n; \; ++i) \; \{ \\ & \; if \; (!visited[i]) \\ & \; dfs \; (i); \\ & \} \end{aligned}
```

## 4.5 Finding Bridges

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> tin, fup;
void\ processBridge(int\ u,\ int\ v)\ \{
    // do something with the found bridge
void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = fup[v] = timer++;
for (int to : adj[v]) {
  if (to == p) continue;
        if (visited[to]) {
            fup[v] = min(fup[v], tin[to]);
        } else {
            dfs(to, v);
fup[v] = min(fup[v], fup[to]);
            if (fup[to] > tin[v])
processBridge(v, to);
        }
    }
// Doesn't work with multiple edges
// But multiple edges are never bridges, so it's easy to check
void findBridges() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);\\
    fup.assign(n, -1);\\
    bridges.clear();
    FOR(i, 0, n)
        if (!visited[i])
            dfs(i);
}
```

#### 4.6 Lowest Common Ancestor

```
\begin{split} &\inf \ n, \ l; \ // \ l == \log N \ (usually \ about \ ^20) \\ &\operatorname{vector} < \operatorname{vector} < \operatorname{int} >> \operatorname{adj}; \\ &\operatorname{int \ timer}; \\ &\operatorname{vector} < \operatorname{int} > \operatorname{tin}, \ \operatorname{tout}; \\ &\operatorname{vector} < \operatorname{vector} < \operatorname{int} >> \operatorname{up}; \\ &\operatorname{void \ dfs}(\operatorname{int \ v}, \ \operatorname{int \ p}) \\ &\left\{ & \operatorname{tin}[v] = + + \operatorname{timer}; \\ &\operatorname{up}[v][0] = p; \\ &// \ w \operatorname{Up}[v][0] = \operatorname{weight}[v][u]; \ // <- \ \operatorname{path \ weight \ sum \ to \ 2^{-i}} \\ &\operatorname{th \ ancestor} \\ &\operatorname{for \ (int \ i \ = 1; \ i <= l; \ ++i)} \\ &\operatorname{up}[v][i] = \operatorname{up}[\operatorname{up}[v][i-1]][i-1]; \\ &// \ w \operatorname{Up}[v][i] = \operatorname{wUp}[v][i-1]] + \ w \operatorname{Up}[\operatorname{up}[v][i-1]][i-1]; \\ &\operatorname{for \ (int \ u \ : adj[v]) \ \{} \\ &\operatorname{if \ (u \ != p)} \\ &\operatorname{dfs}(u, v); \\ &\} \\ &\operatorname{tout}[v] = + + \operatorname{timer}; \\ \end{split}
```

```
bool isAncestor(int u, int v)
    \operatorname{return} \ \operatorname{tin}[u] <= \operatorname{tin}[v] \ \&\& \ \operatorname{tout}[v] <= \operatorname{tout}[u];
int lca(int u, int v)
    if\ (is Ancestor(u,\ v))
         return u:
    if (isAncestor(v, u))
         return v;
    for (int i = 1; i >= 0; --i) {
         if \; (!isAncestor(up[u][i], \; v)) \\
             u = up[u][i];
    return up[u][0];
void preprocess(int root) {
    tin.resize(n);
    tout.resize(n);
    timer = 0;
    l = ceil(log2(n));
    up.assign(n, vector<int>(l + 1));
    dfs(root, root);
```

#### 4.7 Max Flow With Dinic 2

```
struct FlowEdge {
    int v, u;
    long long cap, flow = 0;
    FlowEdge(int\ v,\ int\ u,\ long\ long\ cap): v(v),\ u(u),\ cap(cap)
};
struct Dinic {
    const long long flow_inf = 1e18;
    vector<FlowEdge> edges;
    vector<vector<int>> adj;
    int n, m = 0;
    int s, t;
    vector<int> level, ptr;
    queue<int> q;
    Dinic(int\ n,\ int\ s,\ int\ t):n(n),\,s(s),\,t(t)\ \{
       adj.resize(n);
        level.resize(n);
       ptr.resize(n);
    adj[v].push_back(m);
        adj[u].push\_back(m+1);
   \begin{array}{c} \text{bool bfs() } \{\\ \text{while (!q.empty()) } \{\\ \text{int } \mathbf{v} = \text{q.front()}; \end{array}
            q.pop();
            for (int id : adj[v]) {
               if (edges[id].cap - edges[id].flow < 1)
                    continue;
                if (level[edges[id].u] != -1)
                    continue;
                level[edges[id].u] = level[v] + 1;
                q.push(edges[id].u);
            }
       return level[t] != -1;
    long long dfs(int v, long long pushed) \{
       if (pushed == 0)
            return 0;
        if (v == t)
           return pushed;
        for (int& cid = ptr[v]; cid < (int)adj[v].size(); cid++)
```

```
int id = adj[v][cid];
               int u = edges[id].u;
               \label{eq:condition} \text{if } (\text{level}[v] + 1 \stackrel{!}{!} = \text{level}[u] \mid\mid \text{edges}[\text{id}].\text{cap - edges}[\text{id}]
                       ].flow < 1)
                    continue:
               long long tr = dfs(u, min(pushed, edges[id].cap -
                      edges[id].flow));
               if (tr == 0)
                    continue;
               edges[id].flow += tr;
edges[id ^ 1].flow -= tr;
               return tr;
          return 0;
    \begin{array}{c} long\ long\ flow()\ \{\\ long\ long\ f=0; \end{array}
          while (true) {
               fill(level.begin(), level.end(), -1);
               level[s] = 0;
               q.push(s);
               if (!bfs())
                    break:
               fill(ptr.begin(), ptr.end(), 0);
               while (long long pushed = dfs(s, flow_inf)) {
                    f += pushed;
          return f;
};
```

#### 4.8 Max Flow With Dinic

```
struct Edge {
     int f, c;
     int to;
     pii revIdx;
     int dir:
     int idx;
};
int n, m;
vector<Edge> adjList[MAX_N];
int\ level[MAX\_N];
void addEdge(int a, int b, int c, int i, int dir) {
     int idx = adjList[a].size();
     int revIdx = adjList[b].size();
adjList[a].pb({ 0,c,b, {b, revIdx} ,dir,i });
adjList[b].pb({ 0,0,a, {a, idx} ,dir,i });
}
bool bfs(int s, int t) {
     FOR(i, 0, n) level[i] = -1;
     level[s] = 0;
     queue<int> Q;
     Q.push(s):
     while (!Q.empty()) {
          auto t = Q.front(); Q.pop();
for (auto x : adjList[t]) {
                 \begin{array}{l} \text{if } (\text{level}[x.\text{to}] < 0 \&\& \text{ x.f} < \text{x.c}) \ \{\\ \text{level}[x.\text{to}] = \text{level}[t] + 1; \end{array} 
                      Q.push(x.to);
          }
     return level[t] >= 0;
}
int send(int u, int f, int t, vector<int>& edgeIdx) {
     if (u == t) return f;
     for (; edgeIdx[u] < adjList[u].size(); edgeIdx[u]++) {
           auto\& e = adjList[u][edgeIdx[u]];
            \begin{split} &\text{if (level[e.to] == level[u] + 1 \&\& e.f < e.c) \{} \\ &\text{int curr\_flow = min(f, e.c - e.f);} \\ &\text{int next\_flow = send(e.to, curr\_flow, t, edgeIdx);} \\ \end{aligned} 
                if (\text{next\_flow} > 0) {
                      e.f += next_flow;
                      adjList[e.revIdx.first][e.revIdx.second].f -=
                              next_flow;
```

```
return next flow:
         }
    return 0:
}
int maxFlow(int s, int t) {
    int f = 0
    while (bfs(s, t)) {
vector < int > edgeIdx(n, 0);
         while (int extra = send(s, oo, t, edgeIdx)) {
              f += extra;
    return f:
}
\begin{array}{l} \mathrm{void\ init}()\ \{\\ \mathrm{cin} >> n >> m; \end{array}
    FOR(i,\,0,\,m)\,\,\{
         int a, b, c;
         cin >> a >> b >> c;
a--; b--;
         addEdge(a, b, c, i, 1);
         addEdge(b, a, c, i, -1);
}
```

#### 4.9 Max Flow With Ford Fulkerson

```
struct Edge \{
     int to, next;
     ll f, c;
     int idx. dir:
     int from;
}:
vector < Edge > edges;
vector<int> first;
void addEdge(int a, int b, ll c, int i, int dir) {
     edges.pb({ b, first[a], 0, c, i, dir, a });
edges.pb({ a, first[b], 0, 0, i, dir, b });
     first[a] = edges.size() - 2;

first[b] = edges.size() - 1;
void init() {
     cin >> n >> m;
edges.reserve(4 * m);
     first = vector < int > (n, -1);
     FOR(i, 0, m) {
          int a, b, c;
          cin >> a >> b >> c;
          a--; b--;
          addEdge(a, b, c, i, 1);
          addEdge(b, a, c, i, -1);
     }
}
int cur\_time = 0;
vector<int> timestamp;
\begin{array}{l} ll~dfs(int~v,~ll~flow=OO)~\{\\ if~(v==n-1)~return~flow;\\ timestamp[v]=cur\_time; \end{array}
     for (int e = first[v]; e != -1; e = edges[e].next) {
          if (edges[e].f < edges[e].c && timestamp[edges[e].to] !=
               cur_time) {
int pushed = dfs(edges[e].to, min(flow, edges[e].c -
               edges[e].f);

if (pushed > 0) {

edges[e].f += pushed;

edges[e ^ 1].f -= pushed;
                    return pushed;
               }
          }
     return 0;
```

```
ll maxFlow() {
    cur_time = 0;
    timestamp = vector<int>(n, 0);
    ll f = 0, add;
    while (true) {
        cur_time++;
        add = dfs(0);
        if (add > 0) {
            f += add;
        }
        else {
                break;
        }
    }
    return f;
}
```

#### 4.10 Min Cut

```
\begin{split} & \operatorname{init}(); \\ & \text{ll } f = \operatorname{maxFlow}(); \ / / \text{ Ford-Fulkerson} \\ & \operatorname{cur\_time} + +; \\ & \operatorname{dfs}(0); \\ & \operatorname{set} < \operatorname{int} > \operatorname{cc}; \\ & \operatorname{for } (\operatorname{auto } e : \operatorname{edges}) \ \{ \\ & \operatorname{if } (\operatorname{timestamp}[e.\operatorname{from}] == \operatorname{cur\_time} \ \& \ \operatorname{timestamp}[e.\operatorname{to}] \ != \\ & \operatorname{cur\_time} \ \{ \\ & \operatorname{cc.insert}(e.\operatorname{idx}); \\ & \} \\ & / / \ \# \ \text{of } \operatorname{edges} \ \text{in } \operatorname{min-cut}, \operatorname{capacity} \ \text{of } \operatorname{cut}) \\ & / / \ [\operatorname{indices} \ \text{of } \operatorname{edges} \ \text{forming } \operatorname{the } \operatorname{cut}] \\ & \operatorname{cout} << \operatorname{cc.size}() << ``\_'' << f << \operatorname{endl}; \\ & \operatorname{for } (\operatorname{auto} \ x : \operatorname{cc}) \operatorname{cout} << x + 1 << ``\_''; \\ \end{split}
```

## 4.11 Number Of Paths Of Fixed Length

Let G be the adjacency matrix of a graph. Then  $C_k = G^k$  gives a matrix, in which the value  $C_k[i][j]$  gives the number of paths between i and j of length k.

## 4.12 Shortest Paths Of A Fixed Length

Define  $A \odot B = C \iff C_{ij} = \min_{p=1..n} (A_{ip} + B_{pj})$ . Let G be the adjacency matrix of a graph. Also, let  $L_k = G \odot ... \odot G = G^{\odot k}$ . Then the value  $L_k[i][j]$  denotes the length of the shortest path between i and j which consists of exactly k edges.

## 4.13 Strongly Connected Components

```
\label{eq:continents} \begin{split} & \operatorname{vector} < \operatorname{vector} < \operatorname{int} > > \operatorname{g}, \operatorname{gr}; \; // \operatorname{adjList} \operatorname{and} \operatorname{reversed} \operatorname{adjList} \\ & \operatorname{vector} < \operatorname{int} > \operatorname{order}, \operatorname{component}; \\ & \operatorname{void} \operatorname{dfs1} \; (\operatorname{int} \; v) \; \{ \\ & \operatorname{used}[v] = \operatorname{true}; \\ & \operatorname{for} \; (\operatorname{size}_t \; i = 0; \; i < \operatorname{g}[v].\operatorname{size}(); \; + + \mathrm{i}) \\ & \quad \operatorname{if} \; (!\operatorname{used}[\; \operatorname{g}[v][i] \; ]) \\ & \quad \operatorname{dfs1} \; (\operatorname{g}[v][i]); \\ & \quad \operatorname{order.push\_back} \; (v); \end{split}
```

```
void dfs2 (int v) {
     used[v] = true; \\
    \begin{array}{l} {\rm component.push\_back} \ (v); \\ {\rm for} \ ({\rm size\_t} \ i{=}0; \ i{<}{\rm gr[v].size()}; \ +{+}i) \\ {\rm if} \ (!{\rm used[} \ {\rm gr[v][i]} \ ]) \end{array}
               dfs2 (gr[v][i]);
\mathrm{int}\ \mathrm{main}()\ \{
     int n;
     // read n
     for (;;) {
          int a, b;
          // read edge a -> b
          g[a].push\_back\ (b);
          gr[b].push_back (a);
     used.assign (n, false);
     for (int i=0; i< n; ++i)
          if (!used[i])
              dfs1 (i);
     used.assign (n, false);
     for (int i=0; i < n; ++i) {
          int v = order[n-1-i];
          if (!used[v]) {
               dfs2(v);
               // do something with the found component
               component.clear(); // components are generated in
                       toposort-order
}
```

## 5 Math

## 5.1 Big Integer Multiplication With FFT

```
\begin{array}{l} {\rm complex\!<\!ld\!>\,a[MAX\_N],\;b[MAX\_N];} \\ {\rm complex\!<\!ld\!>\,fa[MAX\_N],\;fb[MAX\_N],\;fc[MAX\_N];} \end{array}
complex<ld> cc[MAX_N];
string mul(string as, string bs) {
     int sgn1 = 1:
     int sgn2 = 1;
     if (as[0] == '-') {
          sgn1 = -1;
          as = as.substr(1);
     if (bs[0] == '-') {
          sgn2 = -1;
          bs = bs.substr(1);
     int\ n = as.length() \, + \, bs.length() \, + \, 1;
    FFT::init(n);

FOR(i, 0, FFT::pwrN) {
          a[i] = b[i] = fa[i] = fb[i] = fc[i] = cc[i] = 0;
     FOR(i, 0, as.size()) {
a[i] = as[as.size() - 1 - i] - '0';
     FOR(i, 0, bs.size()) {
   b[i] = bs[bs.size() - 1 - i] - '0';
     FFT::fft(a, fa);
     FFT::fft(b, fb);
FOR(i, 0, FFT::pwrN) {
fc[i] = fa[i] * fb[i];
      // turn [0,1,2,...,n-1] into [0, n-1, n-2, ..., 1]
     FOR(i, 1, FFT::pwrN) {
    if (i < FFT::pwrN - i) {
               swap(fc[i], fc[FFT::pwrN - i]);
          }
     FFT::fft(fc, cc);
     ll carry = 0;
     vector<int> v;
```

```
FOR(i, 0, FFT::pwrN) {
    int num = round(cc[i].real() / FFT::pwrN) + carry; v.pb(num % 10);
    carry = num / 10;
while (carry > 0) {
v.pb(carry % 10);
    carry /= 10;
reverse(v.begin(), v.end());
bool start = false;
ostringstream ss;
bool allZero = true;
for (auto x : v) {
    if (x != 0) {
        allZero = false;
        break;
if (sgn1*sgn2 < 0 \&\& !allZero) ss << "-";
for (auto x : v) {
if (x == 0 && !start) continue;
    start = true;
    ss \ll abs(x):
if (!start) ss << 0;
return ss.str();
```

### 5.2 Burnside's Lemma

Let G be a finite group that acts on a set X. For each g in G let  $X^g$  denote the set of elements in X that are fixed by g. Burnside's lemma asserts the following formula for the number of orbits:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

### Example. Coloring a cube with three colors.

Let X be the set of  $3^6$  possible face color combinations. Let's count the sizes of the fixed sets for each of the 24 rotations:

- one 0-degree rotation which leaves all  $3^6$  elements of X unchanged
- six 90-degree face rotations, each of which leaves  $3^3$  elements of X unchanged
- three 180-degree face rotation, each of which leaves  $3^4$  elements of X unchanged
- eight 120-degree vertex rotations, each of which leaves  $3^2$  elements of X unchanged
- six 180-degree edge rotations, each of which leaves  $3^3$  elements of X unchanged

The average is then  $\frac{1}{24}(3^6 + 6 \cdot 3^3 + 3 \cdot 3^4 + 8 \cdot 3^2 + 6 \cdot 3^3) = 57$ . For n colors:  $\frac{1}{24}(n^6 + 3n^4 + 12n^3 + 8n^2)$ .

## Example. Coloring a circular stripe of n cells with two colors.

X is the set of all colored striped (it has  $2^n$  elements), G is the group of rotations (n elements - by 0 cells, by 1 cell, ..., by (n-1) cells). Let's fix

some K and find the number of stripes that are fixed by the rotation by K cells. If a stripe becomes itself after rotation by K cells, then its 1st cell must have the same color as its  $(1+K \mod n)$ -th cell, which is in turn the same as its  $(1+2K \mod n)$ -th cell, etc., until  $mK \mod n = 0$ . This will happen when  $m = n/\gcd(K,n)$ . Therefore, we have  $n/\gcd(K,n)$  cells that must all be of the same color. The same will happen when starting from the second cell and so on. Therefore, all cells are separated into  $\gcd(K,n)$  groups, with each group being of one color, and that yields  $2^{\gcd(K,n)}$  choices. That's why the answer to the original problem is  $\frac{1}{n} \sum_{k=0}^{n-1} 2^{\gcd(k,n)}$ .

#### 5.3 Chinese Remainder Theorem

Let's say we have some numbers  $m_i$ , which are all mutually coprime. Also, let  $M = \prod_i m_i$ . Then the system of congruences

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \dots \\ x \equiv a_k \pmod{m_k} \end{cases}$$

is equivalent to  $x \equiv A \pmod{M}$  and there exists a unique number A satisfying  $0 \le A \le M$ .

Solution for two:  $x \equiv a_1 \pmod{m_1}$ ,  $x \equiv a_2 \pmod{m_2}$ . Let  $x = a_1 + km_1$ . Substituting into the second congruence:  $km_1 \equiv a_2 - a_1 \pmod{m_2}$ . Then,  $k = (m_1)_{m_2}^{-1}(a_2 - a_1) \pmod{m_2}$ . and we can easily find x. This can be extended to multiple equations by solving them one-by-one.

If the moduli are not coprime, solve the system  $y\equiv 0\pmod{\frac{m_1}{g}}, y\equiv \frac{a_2-a_1}{g}\pmod{\frac{m_2}{g}}$  for y. Then let  $x\equiv gy+a_1\pmod{\frac{m_1m_2}{q}}$ .

#### 5.4 Euler Totient Function

#### 5.5 Extended Euclidean Algorithm

```
solveEq(b, a%b, xx, yy, g);
   x = yy;
   y = xx-yy*(a/b);
// ax+by=c
bool solve
Eq(ll a, ll b, ll c, ll& x, ll& y, ll& g) {
    solveEq(a, b, x, y, g);
   if(c\%g != 0) return false;

x *= c/g; y *= c/g;
   return true;
// Finds a solution (x, y) so that x >= 0 and x is minimal
bool solveEqNonNegX(ll a, ll b, ll c, ll& x, ll &y, ll& g) {
    if(!solveEq(a,\;b,\;c,\;x,\;y,\;g)) \ return \ false;\\
   ll k = x*g/b;

x = x - k*b/g;
    y = y + k*a/g;
   if(x < 0) {
       x + = b/g;
        y = a/g;
   return true;
}
```

All other solutions can be found like this:

$$x' = x - k \frac{b}{g}, y' = y + k \frac{a}{g}, k \in \mathbb{Z}$$

#### 5.6 Factorization With Sieve

```
// Use linear sieve to calculate minDiv
vector<pll> factorize(ll x) {
    vector<pll> res;
    ll prev = -1;
    ll cnt = 0;
    while(x != 1) {
        ll d = minDiv[x];
        if(d == prev) {
            cnt++;
        } else {
            if(prev != -1) res.pb({prev, cnt});
            prev = d;
            cnt = 1;
        }
        x /= d;
    }
    res.pb({prev, cnt});
    return res;
}
```

## 5.7 FFT With Modulo

```
 \begin{array}{lll} bool \ is Generator(ll \ g) \ \{ \\ if \ (pwr(g, \ M-1) \ != 1) \ return \ false; \\ for \ (ll \ i = 2; \ i^*i <= M-1; \ i++) \ \{ \\ if \ ((M-1) \ \% \ i == 0) \ \{ \\ ll \ q = \ i; \\ if \ (isPrime(q)) \ \{ \\ ll \ p = \ (M-1) \ / \ q; \\ ll \ pp = pwr(g, \ p); \\ if \ (pp == 1) \ return \ false; \\ \} \\ q = \ (M-1) \ / \ i; \\ if \ (isPrime(q)) \ \{ \\ ll \ p = \ (M-1) \ / \ q; \\ ll \ pp = pwr(g, \ p); \\ if \ (pp == 1) \ return \ false; \\ \} \\ \} \\ return \ true; \\ \} \\ namespace \ FFT \ \{ \end{array}
```

```
vector < ll > r;
vector<ll> omega;
ll\ logN,\ pwrN;
void initLogN() {

log N = 0; 

pwrN = 1;

      while (pwrN < n) {
    pwrN *= 2;
            logN++;
      n = pwrN;
void initOmega() {
      ll g = 2;
      while (!isGenerator(g)) g++;
      ll G = 1;
      g = pwr(g, (M - 1) / pwrN);
FOR(i, 0, pwrN) {
           omega[i] = G;
G *= g;
G %= M;
      }
void initR() {
      FOR(i, 1, pwrN) {

r[i] = r[i / 2] / 2 + ((i \& 1) << (logN - 1));
{\rm void~initArrays}()~\{
      r.clear();
      r.resize(pwrN);
      omega.clear();
      omega.resize(pwrN);
void\ init(ll\ n)\ \{
      FFT::n = n;
      initLogN();
      initArrays();
      initOmega();
      initR();
 \begin{array}{l} {\rm void\ fft}({\rm ll\ a[],\ ll\ f[])\ \{} \\ {\rm for\ (ll\ i=0;\ i< pwrN;\ i++)\ \{} \end{array} 
           f[i] = a[r[i]];
      for (ll k = 1; k < pwrN; k *= 2) {
  for (ll i = 0; i < pwrN; i += 2 * k) {
    for (ll j = 0; j < k; j++) {
                       auto z = omega[j*n / (2 * k)] * f[i + j + k] % M;
                       \begin{array}{l} f[i\ +\ j\ +\ k] = f[i\ +\ j]\ -\ z;\\ f[i\ +\ j]\ +=\ z;\\ f[i\ +\ j\ +\ k]\ \%=\ M;\\ if\ (f[i\ +\ j\ +\ k]\ <0)\ f[i\ +\ j\ +\ k]\ +=\ M;\\ f[i\ +\ j]\ \%=\ M; \end{array}
                }
          }
     }
}
```

#### 5.8 FFT

```
namespace FFT {
  int n;
  vector<int> r;
  vector<complex<ld>> omega;
  int logN, pwrN;

  void initLogN() {
    logN = 0;
    pwrN = 1;
    while (pwrN < n) {
        pwrN *= 2;
    }
}</pre>
```

```
logN++;
          n = pwrN;
     void initOmega() {
          FOR(i, 0, pwrN) {
              omega[i] = { \cos(2 * i*PI / n), \sin(2 * i*PI / n) };
     }
     void initR() {
          FOR(i, 1, pwrN) {
              r[i] = r[i / 2] / 2 + ((i \& 1) << (logN - 1));
     }
     void initArrays() {
          r.clear();
          r.resize(pwrN);
          omega.clear();
          omega.resize(pwrN);
     void\ init(int\ n)\ \{
          FFT::n\,=\,n;
          initLogN();
          initArrays():
          initOmega();\\
          initR();
     void fft(complex<ld> a[], complex<ld> f[]) {
          FOR(i, 0, pwrN) {
               f[i] = a[r[i]];
          for (ll k = 1; k < pwrN; k *= 2) {
for (ll i = 0; i < pwrN; i += 2 * k) {
                     \begin{array}{l} (\text{li } 1-0), \ 1 < \text{pwith}, \ i-2 & \text{if} \ \\ \text{for } (\text{ll } j=0); \ j < k; \ j++) \ \{ \\ \text{auto } z = \text{omega}[j^*n \ / \ (2 * k)] * f[i+j+k]; \\ f[i+j+k] = f[i+j] - z; \\ f[i+j] +=z; \end{array} 
                   }
          }
}
```

#### 5.9 Formulas

```
\begin{array}{lll} \sum_{i=1}^n i & = & \frac{n(n+1)}{2}; & \sum_{i=1}^n i^2 & = & \frac{n(2n+1)(n+1)}{6}; \\ \sum_{i=1}^n i^3 & = & \frac{n^2(n+1)^2}{4}; & \sum_{i=1}^n i^4 & = & \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}; \\ \sum_{i=a}^b c^i & = & \frac{c^{b+1}-c^a}{c-1}, c & \neq & 1; & \sum_{i=1}^n a_1 + (i-1)d \\ & = & \frac{n(a_1+a_n)}{2}; & \sum_{i=1}^n a_1r^{i-1} & = & \frac{a_1(1-r^n)}{1-r}, r \neq & 1; \\ \sum_{i=1}^\infty ar^{i-1} & = & \frac{a_1}{1-r}, |r| \leq 1. \end{array}
```

#### 5.10 Linear Sieve

```
\label{eq:local_problem} \begin{split} & ll \ minDiv[MAXN+1]; \\ & vector < ll > primes; \\ & void \ sieve(ll \ n) \{ \\ & FOR(k, \ 2, \ n+1) \{ \\ & \ minDiv[k] = k; \\ \} \\ & FOR(k, \ 2, \ n+1) \ \{ \\ & \ if(minDiv[k] = k) \ \{ \\ & \ primes.pb(k); \\ \} \\ & for(auto \ p : primes) \ \{ \\ & \ if(p > minDiv[k]) \ break; \\ & \ if(p > minDiv[p^*k] = p; \\ \} \\ \} \\ & \} \end{split}
```

#### 5.11 Modular Inverse

```
\label{eq:bool invWithEuclid(ll a, ll m, ll& aInv) } \left\{ \begin{array}{l} ll \ x, \ y, \ g; \\ if(!solveEqNonNegX(a, \ m, \ 1, \ x, \ y, \ g)) \ return \ false; \\ aInv = x; \\ return \ true; \\ \right\} \\ // \ Works \ only \ if \ m \ is \ prime \\ ll \ invFermat(ll \ a, \ ll \ m) \ \left\{ \\ return \ pwr(a, \ m-2, \ m); \\ \right\} \\ // \ \ Works \ only \ if \ gcd(a, \ m) = 1 \\ ll \ invEuler(ll \ a, \ ll \ m) \ \left\{ \\ return \ pwr(a, \ phi(m)-1, \ m); \\ \right\}
```

### 5.12 Simpson Integration

```
 \begin{array}{l} {\rm const\ int\ N=1000\ ^*\ 1000;\ //\ number\ of\ steps\ (already\ multiplied\ by\ 2)} \\ \\ {\rm double\ simpsonIntegration(double\ a,\ double\ b)\{} \\ {\rm double\ h=(b-a)\ /\ N;} \\ {\rm double\ s=f(a)+f(b);\ //\ a=x\_0\ and\ b=x\_2n} \\ {\rm for\ (int\ i=1;\ i<=N-1;\ ++i)\ \{} \\ {\rm double\ x=a+h\ ^*\ i;} \\ {\rm s\ +=f(x)\ ^*\ ((i\ \&\ 1)\ ?\ 4:2);} \\ {\rm s\ ^*=h\ /\ 3;} \\ {\rm return\ s;} \\ {\rm \}} \\ \\ \end{array}
```

## 6 Strings

#### 6.1 Aho Corasick Automaton

```
// alphabet size
const int K = 70:
// the indices of each letter of the alphabet
int intVal[256];
void init() {
    int curr = 2
    \mathrm{intVal}[1] \, = \, 1;
    for(char c = '0'; c <= '9'; c++, curr++) intVal[(int)c] =
          curr;
    for(char c = 'A'; c <= 'Z'; c++, curr++) intVal[(int)c] =
    for(char c = 'a'; c \le 'z'; c++, curr++) intVal[(int)c] =
          curr:
}
struct Vertex {
    int next[K];
    vector<int> marks;
    // \hat{} this can be changed to int mark = -1, if there will be
         no duplicates
    int p = -1;
    char pch;
    int exitLink = -1;
         exitLink points to the next node on the path of suffix
           links which is marked
    int go[K];
      / ch has to be some small char
    Vertex(int \_p=-1, char ch=(char)1) : p(\_p), pch(ch) \{
       {\rm fill}({\rm begin}({\rm next}),\,{\rm end}({\rm next}),\,{\text{\rm -1}});
        fill(begin(go), end(go), -1);
};
vector < Vertex > t(1);
```

```
void addString(string const& s, int id) {
    int v = 0;
    for (char ch : s) {
        int c = intVal[(int)ch];
        if (t[v].next[c] == -1) {
             t[v].next[c] = t.size();
             t.emplace_back(v, ch);
         v = t[v].next[c];
    t[v].marks.pb(id);
}
int go(int v, char ch);
int getLink(int v) {
    if (t[v].link == -1) {
        if (v == 0 || t[v].p == 0)
            t[v].link = 0;
             t[v].link = go(getLink(t[v].p), \ t[v].pch); \\
    return t[v].link:
}
\begin{array}{l} \mathrm{int} \ \mathrm{getExitLink(int} \ v) \ \{ \\ \mathrm{if}(\mathrm{t[v].exitLink} \ !=-1) \ \mathrm{return} \ \mathrm{t[v].exitLink}; \end{array}
    int l = getLink(v);
if(l == 0) return t[v].exitLink = 0;
     \begin{aligned} & \text{if}(!t[l].marks.empty()) \ & \text{return} \ t[v].exitLink = l; \\ & \text{return} \ t[v].exitLink = getExitLink(l); \end{aligned} 
\mathrm{int}\ \mathrm{go}(\mathrm{int}\ \mathrm{v},\,\mathrm{char}\ \mathrm{ch})\ \{
    int c = intVal[(int)ch];
    if (t[v].go[c] == -1) {
        if (t[v].next[c] != -1)
            t[v].go[c] = t[v].next[c];
             t[v].go[c] = v == 0 ? 0 : go(getLink(v), ch);
    return t[v].go[c];
}
void walkUp(int v, vector<int>& matches) {
    if(v == 0) return;
    if(!t[v].marks.empty()) {
   for(auto m : t[v].marks) matches.pb(m);
     walkUp(getExitLink(v), matches);
}
   returns the IDs of matched strings.
   Will contain duplicates if multiple matches of the same
      string are found.
vector<int> walk(const string& s) {
    vector<int> matches;
    int curr = 0;
    for(char c : s) {
        curr = go(curr, c);
        if(!t[curr].marks.empty()) {
             for(auto m : t[curr].marks) matches.pb(m);
         walkUp(getExitLink(curr), matches);
    return matches:
/* Usage:
 * addString(strs[i], i);
 * auto matches = walk(text);
 st .. do what you need with the matches - count, check if
       some id exists, etc ..
 * Some applications:
 * - Find all matches: just use the walk function
 * - Find lexicographically smallest string of a given length
        that doesn't match any of the given strings:
 * For each node, check if it produces any matches (it either
        contains some marks or walkUp(v) returns some marks)
 * Remove all nodes which produce at least one match. Do
        DFS in the remaining graph, since none of the
        remaining nodes
 * will ever produce a match and so they're safe.
 * - Find shortest string containing all given strings:
```

```
* For each vertex store a mask that denotes the strings which match at this state. Start at (v = root, mask = 0),

* we need to reach a state (v, mask=2^n-1), where n is the number of strings in the set. Use BFS to transition between states

* and update the mask.

*/
```

### 6.2 Hashing

```
struct HashedString { const ll A1 = 999999929, B1 = 1000000009, A2 =
             1000000087, B2 = 1000000097;
     vector<ll> A1pwrs, A2pwrs;
     vector<pll> prefixHash;
     HashedString(const\ string\&\ \_s)\ \{
          init(\underline{\hspace{1em}}s);
          calcHashes(s);
     void init(const string& s) {
          11 \ a1 = 1;
          ll a2 = 1
          FOR(i,\,0,\,(int)s.length(){+}1)~\{
               A1pwrs.pb(a1);
               A2pwrs.pb(a2);
               a1 = (a1*A1)\%B1;
               a2 = (a2*A2)\%B2;
          }
     void calcHashes(const string& s) {
          pll h = \{0, 0\};
          prefixHash.pb(h);
               ll h1 = (prefixHash.back().first*A1 + c)%B1;
ll h2 = (prefixHash.back().second*A2 + c)%B2;
               \overrightarrow{prefixHash.pb}(\{h1,\;h2\});
          }
      \begin{array}{l} \label{eq:continuous_policy} \text{pll getHash(int l, int r) } \{ \\ \label{eq:continuous_policy} \text{ll h1} = (\text{prefixHash[r+1].first - prefixHash[l].first*} \end{array} 
                  A1pwrs[r+1-l]) \% B1;
          ll h2 = (prefixHash[r+1].second - prefixHash[l].second* A2pwrs[r+1-l]) % B2;
          if(h1 < 0) h1 += B1;

if(h2 < 0) h2 += B2;
          return {h1, h2};
};
```

#### 6.3 KMP

```
// Knuth-Morris-Pratt algorithm
vector<int> findOccurences(const string& s, const string& t)
    {
    int n = s.length();
    int m = t.length();
    string S = s + "#" + t;
    auto pi = prefixFunction(S);
    vector<int> ans;
    FOR(i, n+1, n+m+1) {
        if(pi[i] == n) {
            ans.pb(i-2*n);
        }
    }
    return ans;
}
```

#### 6.4 Prefix Function Automaton

```
// aut[oldPi][c] = newPi
vector<vector<int>> computeAutomaton(string s) {
const char BASE = 'a';
s += "#";
int n = s.size();
```

```
vector<int> pi = prefixFunction(s);
    vector<vector<int>> aut(n, vector<int>(26));
    for (int i = 0; i < n; i++) \{
        for (int c = 0; c < 26; c++) {
    if (i > 0 && BASE + c != s[i])
                  \operatorname{aut}[i][c] = \operatorname{aut}[\operatorname{pi}[i-1]][c];
                  \operatorname{aut}[i][c] = i + (BASE + c == s[i]);
        }
    return aut;
vector<int> findOccurs(const string& s, const string& t) {
    auto aut = computeAutomaton(s);
    \quad \text{int curr} = 0; \\
    vector<int> occurs;
    FOR(i, 0, (int)t.length()) {
    int c = t[i]-'a';
         curr = aut[curr][c];
         if(curr == (int)s.length()) {
             occurs.pb(i-s.length()+1);
    return occurs:
}
```

## 6.5 Prefix Function

```
\label{eq:continuous_problem} // \ pi[i] \ is the length of the longest proper prefix of the substring $s[0..i]$ which is also a suffix // of this substring wector<int> prefixFunction(const string& $s$) { int $n = (int)s.length(); } vector<int> pi(n); for (int $i = 1; i < n; i++) { int $j = pi[i-1]; } while ($j > 0 && $s[i] != s[j]) $ j = pi[j-1]; if ($s[i] == s[j]) $ j++; pi[i] = j; } return $pi; } \\
```

#### 6.6 Suffix Array

```
{\tt vector}{<} {\tt int}{\gt} \ {\tt sortCyclicShifts} ({\tt string} \ {\tt const} \& \ {\tt s}) \ \{
     int n = s.size();
     const int alphabet = 256; // we assume to use the whole
               ASCIÎ range
       vector < int > p(n), c(n), cnt(max(alphabet, n), 0);
      for (int i = 0; i < n; i++)
            \operatorname{cnt}[s[i]]++;
     for (int i = 1; i < alphabet; i++)
     cnt[i] += cnt[i-1];
for (int i = 0; i < n; i++)
     p[-cnt[s[i]]] = i;

c[p[0]] = 0;
     int classes = 1;
     for (int i = 1; i < n; i++) {
    if (s[p[i]] != s[p[i-1]])
                  classes++;
            c[p[i]] = classes - 1;
     } vector<int> pn(n), cn(n); for (int h = 0; (1 << h) < n; ++h) { for (int i = 0; i < n; i++) { pn[i] = p[i] - (1 << h); if (pn[i] < 0) pn[i] += n;
             \begin{array}{l} \text{fill(cnt.begin(), cnt.begin() + classes, 0);} \\ \text{for (int } i = 0; \ i < n; \ i++) \end{array} 
                  \operatorname{cnt}[\operatorname{c[pn[i]]}]++;
            for (int i = 1; i < classes; i++)
                  \operatorname{cnt}[i] += \operatorname{cnt}[i-1];
            for (int i = n-1; i > = 0; i--)
```

```
p[-cnt[c[pn[i]]]] = pn[i];
       \operatorname{cn}[p[0]] = 0;
       classes = 1;
       for (int i = 1; i < n; i++) {
    pair<int, int> cur = {c[p[i]], c[(p[i] + (1 << h)) %
                n]};
           pair < int, int > prev = \{c[p[i-1]], c[(p[i-1] + (1 <<
               h)) % n]};
           if (cur != prev)
              ++classes:
           cn[p[i]] = classes - 1;
       c.swap(cn);
   return p;
in s
   vector < int > sorted\_shifts = sortCyclicShifts(s);
   sorted\_shifts.erase(sorted\_shifts.begin());
   return\ sorted\_shifts;
```