ACM-ICPC TEAM REFERENCE DOCUMENT

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	3.4 Bellman Ford Algorithm	7	using namespacegnu_pbds;
	3.5 Bipartite Graph	8	typedef long long ll; typedef unsigned long long ull;
	3.6 Finding Articulation Points	8	typedef long double ld; typedef pair <int, int=""> pii;</int,>
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	3.8 Max Flow With Ford Fulkerson	9	typedef pair <double, double=""> pdd; template <typename t=""> using min_heap = priority_queue<t,< th=""></t,<></typename></double,>
			vector <t>, greater<t>>;</t></t>
	3.9 Max Flow With Dinic		template <typename t=""> using max_heap = priority_queue<t< th=""></t<></typename>
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	3.11 Min Cut		null_type, less <t>, rb_tree_tag,</t>
	3.12 Number Of Paths Of Fixed Length	11	tree_order_statistics_node_update>; template <typename k,="" typename="" v=""> using hashmap =</typename>
	3.13 Shortest Paths Of Fixed Length	11	$gp_hash_table < K, V > ;$
			template <typename a,="" b="" typename=""> ostream& operator<<(</typename>
4	Geometry	11	ostream& out, pair <a, b=""> p) { out << "(" << p.first << "," << p.second << ")"; return out;}</a,>
	4.1 2d Vector	11	template <typename t=""> ostream& operator<<(ostream& out,</typename>
	4.2 Line	11	$ \begin{array}{l} vector < T > v) \ \{ \ out << \ "["; \ for (auto\& \ x : v) \ out << \ x << \ ", _"; \ out << \ "]"; return \ out; \} $
	4.3 Convex Hull Gift Wrapping	11	template <typename t=""> ostream& operator<<(ostream& out,</typename>
	4.4 Convex Hull With Graham's Scan		$\begin{array}{l} \operatorname{set} < T > v) \ \{ \ \operatorname{out} << ``\{"; \ \operatorname{for}(\operatorname{auto\&} \ x : v) \ \operatorname{out} << \ x << ", "; \ \operatorname{out} << "\}"; \ \operatorname{return} \ \operatorname{out}; \ \} \end{array}$
	4.5 Circle Line Intersection		template <typename k,="" typename="" v=""> ostream& operator<<(</typename>
	4.6 Circle Circle Intersection		ostream& out, map <k, v=""> m) { out << "{"; for(auto& e : m) out << e.first << "u->u" << e.second << ",u"; out</k,>
			: m) out << e.nrst << ">_" << e.second << ",_"; out << ","; return out; }
	4.7 Common Tangents To Two Circles		template <typename k,="" typename="" v=""> ostream& operator<<(</typename>
	4.8 Number Of Lattice Points On Segment		ostream& out, hashmap <k, v=""> m) { out << "{"; for(auto& e : m) out << e.first << ">_" << e.second <<</k,>
	4.9 Pick's Theorem	12	", $_{\perp}$ "; out << "}"; return out; }

```
#define FAST_IO ios_base::sync_with_stdio(false); cin.tie(
       NULLA
#define TESTS(t) int NUMBER_OF_TESTS; cin >>
       NUMBER_OF_TESTS; for(int t = 1; t <= NUMBER_OF_TESTS; t++)
#define FOR(i, begin, end) for (int i = (begin) - ((begin) > (
       end)); i != (end) - ((begin) > (end)); i += 1 - 2 * ((begin)
         > (end))
#define sgn(a) ((a) > eps ? 1 : ((a) < -eps ? -1 : 0)) #define precise(x) fixed << setprecision(x) #define debug(x) cerr << ">_{\sqcup}" << x << "_{\sqcup}" << x << "
#define pb push_back
#define rnd(a, b) (uniform_int_distribution<int>((a), (b))(rng
#ifndef LOCAL
     #define cerr if(0)cout
     #define endl "\n
#endif
mt19937\ rng(chrono::steady\_clock::now().time\_since\_epoch().
       count());
clock t clock
void startTime() { __clock __ = clock();}
void timeit(string msg) {cerr << ">_" << msg << ":_" << precise(6) << ld(clock() __clock __)/
CLOCKS_PER_SEC << endl;}
const ld PI = asin(1) * 2;
const ld eps = 1e-14;
const int oo = 2e9;
const ll OO = 2e18;
const ll MOD = 10000000007;
const int MAXN = 1000000;
int main() {
    FAST IO;
    startTime();
    timeit("Finished");
    return 0;
}
```

1.2 Python Template

```
import sys
import re
from math import ceil, log, sqrt, floor

__local__run___ = False
if __local__run___:
    sys.stdin = open('input.txt', 'r')
    sys.stdout = open('output.txt', 'w')

def main():
    a = int(input())
    b = int(input())
    print(a*b)

main()
```

1.3 Compilation

```
# Simple compile
g++ -DLOCAL -O2 -o main.exe -std-c++17 -Wall -Wno-unused
-result -Wshadow main.cpp
# Debug
g++ -DLOCAL -std=c++17 -Wshadow -Wall -o main.exe main.
cpp -fsanitize=address -fsanitize=undefined -fuse-ld=gold
-D_GLIBCXX_DEBUG -g
```

1.4 Automatic Test

```
# Linux Bash
# gen, main and stupid have to be compiled beforehand
for((i=1;;++i)); do
echo $i;
```

```
./gen $i > genIn;
diff <(./main < genIn) <(./stupid < genIn) || break;
done

# Windows CMD
@echo off

FOR /L %%I IN (1,1,2147483647) DO (
echo %%I
gen.exe %%I > genIn
main.exe < genIn > mainOut
stupid.exe < genIn > stupidOut
FC mainOut stupidOut || goto :eof
)
```

1.5 Ternary Search

```
 \begin{array}{l} \mbox{double ternary\_search(double l, double r) } \{ \\ \mbox{while } (r \cdot 1 > eps) \; \{ \\ \mbox{double } m1 = l + (r \cdot l) \; / \; 3; \\ \mbox{double } m2 = r \cdot (r \cdot l) \; / \; 3; \\ \mbox{double } f1 = f(m1); \\ \mbox{double } f2 = f(m2); \\ \mbox{if } (f1 < f2) \\ \mbox{l = } m1; \\ \mbox{else} \\ \mbox{r = } m2; \\ \mbox{} \} \\ \mbox{return } f(l); \; / / \mbox{return the maximum of } f(x) \; \mbox{in } [l, \, r] \\ \mbox{} \} \end{array}
```

1.6 Big Integer

```
const int base = 100000000000:
const int base_digits = 9;
struct bigint {
    vector<int> a;
    int sign;
    int size() {
        if (a.empty()) return 0;
int ans = (a.size() - 1) * base_digits;
        int ca = a.back();
        while (ca) ans++, ca \neq 10;
        return ans;
    bigint operator (const bigint &v) {
bigint ans = 1, x = *this, y = v;
        while (!y.isZero()) {
if (y \% 2) ans *= x;
            x *= x, y /= 2;
        return ans:
    string to string() {
        stringstream ss;
        ss <\bar{<} *this;
        string s;
        ss >> s;
        return s:
    int sumof() {
        string s = to_string();
        int ans = 0;
        for (auto c : s) ans += c - '0';
        return ans;
    bigint() : sign(1) \{ \}
    bigint(long long v) {
        *this = v;
    bigint(const string &s) {
        read(s):
    void operator=(const bigint &v) {
        sign = v.sign;
        a = v.a;
    void operator=(long long v) {
        sign = 1;
        a.clear();
        if (v < 0)
```

```
\mathrm{sign} = -1, \ \mathrm{v} = -\mathrm{v};
      for (; v > 0; v = v / base)
a.push_back(v \% base);
bigint operator+(const bigint &v) const {
      if (sign == v.sign) {
            bigint res = v;
             for (int i = 0, carry = 0; i < (int)max(a.size(), v.a.
                  \begin{array}{l} \text{size())} \mid\mid \text{carry}; ++i) \mid \\ \text{if } (i==(\text{int})\text{res.a.size()}) \text{ res.a.push\_back(0)}; \\ \text{res.a[i]} += \text{carry} + (i < (\text{int})\text{a.size()} ? \text{ a[i]} : 0); \\ \text{carry} = \text{res.a[i]} >= \text{base;} \end{array}
                  if (carry) res.a[i] -= base;
             return res;
      return *this - (-v);
bigint operator-(const bigint &v) const {
      if (sign == v.sign) {
    if (abs() >= v.abs()) {
        bigint res = *this;

                  for (int i = 0, carry = 0; i < (int)v.a.size() ||
                         carry; ++i) {
res.a[i] -= carry + (i < (int)v.a.size() ? v.a[i] :
                                   0);
                         carry = res.a[i] < 0;
                        if (carry) res.a[i] += base;
                  res.trim();
                  return res;
             return -(v - *this);
      return *this + (-v);
void operator*=(int v) {
      if (v < 0) sign = -sign, v = -v;
      for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i) {
            \begin{array}{l} \mbox{if } (i == (\mbox{int}) a.size()) \ a.push\_back(0); \\ \mbox{long long cur} = a[i] * (\mbox{long long}) v + carry; \\ \mbox{carry} = (\mbox{int}) (\mbox{cur} / base); \end{array}
            a[i] = (int)(cur \% base);
      trim();
bigint operator*(int v) const {
   bigint res = *this;
      res *= v;
      return res;
void operator*=(long long v) {
      if (v < 0) sign = -sign, v = -v;

for (int i = 0, carry = 0; i < (int)a.size() || carry; ++i) {

if (i == (int)a.size()) a.push_back(0);

long long cur = a[i] * (long long)v + carry;

carry = (int)(cur / base);
             a[i] = (int)(cur \% base);
      trim();
bigint operator*(long long v) const {
      bigint res = *this;
      res *= v;
      return res;
friend pair<br/>
bigint, bigint> divmod(const bigint &a1, const
          bigint &b1) {
       int norm = base / (b1.a.back() + 1);
      bigint a = a1.abs() * norm;
bigint b = b1.abs() * norm;
      bigint q, r;
      q.a.resize(a.a.size());
for (int i = a.a.size() - 1; i >= 0; i--) {
            r *= base;
             r += a.a[i];
            \begin{array}{l} \mathrm{int} \ \mathrm{s1} = \mathrm{r.a.size}() <= \mathrm{b.a.size}() \ ? \ 0 : \mathrm{r.a[b.a.size}()]; \\ \mathrm{int} \ \mathrm{s2} = \mathrm{r.a.size}() <= \mathrm{b.a.size}() - 1 \ ? \ 0 : \mathrm{r.a[b.a.size}() \end{array}
             int d = ((long long)base * s1 + s2) / b.a.back();
            r -= b * d;
while (r < 0) r += b, --d;
             q.a[i] = d;
      q.sign = a1.sign * b1.sign;
      r.sign = a1.sign;
      q.trim();
      r.trim();
```

```
return make_pair(q, r / norm);
bigint operator/(const bigint &v) const {
     return divmod(*this, v).first;
bigint operator%(const bigint &v) const {
     return divmod(*this, v).second;
void operator/=(int v) {
     if (v < 0) \text{ sign} = -\text{sign}, v = -v;
for (int i = (int)a.size() - 1, rem = 0; i >= 0; --i) {
          long long cur = a[i] + rem * (long long)base;

a[i] = (int)(cur / v);
          rem = (int)(cur \% v);
     trim();
bigint operator/(int v) const {
   bigint res = *this;
     bigint res
     res /= v;
     return res;
int operator%(int v) const {
     if (v < 0) \dot{v} = -\dot{v};
     int m = 0;
     for (int i = a.size() - 1; i >= 0; --i) 

m = (a[i] + m * (long long)base) % v;

return m * sign;
void operator+=(const bigint &v) {
     *this = *this + v;
void operator-=(const bigint &v) {
     *this = *this - v;
void operator*=(const bigint &v) {
     *this = *this * v;
void operator/=(const bigint &v) {
    *this = *this / v;
bool operator<(const bigint &v) const {
    if (sign != v.sign) return sign < v.sign;
    if (a.size() != v.a.size())
     \begin{array}{ll} \text{If } (a.size():=\text{v.a.size()}) \\ \text{return a.size()} * \text{sign} < \text{v.a.size()} * \text{v.sign}; \\ \text{for } (\text{int } i = \text{a.size()} - 1; \ i >= 0; \ i --) \\ \text{if } (a[i] != \text{v.a[i]}) \\ \text{return } a[i] * \text{sign} < \text{v.a[i]} * \text{sign}; \end{array}
     return false;
bool operator>(const bigint &v) const { return v < *this;}
bool operator <= (const bigint &v) const {
    return !(v < *this);
bool operator>=(const bigint &v) const {
     return !(*this < v);
bool operator==(const bigint &v) const {
return !(*this < v) && !(v < *this);
bool operator!=(const bigint &v) const {
     return *this < v || v < *this;
void trim() {
   while (!a.empty() && !a.back()) a.pop_back();
     if (a.empty()) sign = 1;
bool isZero() const {
     return a.empty() || (a.size() == 1 \&\& !a[0]);
bigint operator-() const {
   bigint res = *this;
     res.sign = -sign;
     return res;
bigint abs() const {
     bigint res = *this;
res.sign *= res.sign;
     return res;
long long longValue() const {
     long long res = 0;
     for (int i = a.size() - 1; i >= 0; i--) res = res * base + a[i
     return res * sign;
```

```
friend bigint gcd<br/>(const bigint &a, const bigint &b) {
      return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint &b) {
     return a / gcd(a, b) * b;
void read(const string &s) {
      sign = 1
      a.clear();
      int pos = 0:
      while (pos < (int)s.size() && (s[pos] == '-' || s[pos] ==
            if (s[pos] == '-') sign = -sign;
            ++pos;
      for (int i = s.size() - 1; i >= pos; i -= base_digits) {
            for (int j = \max(pos, i - base\_digits + 1); j <= i; j
                 x = x * 10 + s[j] - 0;
            a.push_back(x);
      trim():
friend istream &operator>>(istream &stream, bigint &v) {
      stream >> s
      v.read(s);
      return stream;
friend ostream & operator << (ostream & stream, const bigint
      \begin{array}{l} {\rm if}\; (v.{\rm sign} == -1)\; {\rm stream} <<\, '\cdot '; \\ {\rm stream} <<\, (v.a.{\rm empty}()\;?\;0:v.a.{\rm back}()); \\ {\rm for}\; ({\rm int}\; i = ({\rm int})v.a.{\rm size}()\; -2;\; i>=0;\; -i) \end{array} 
           stream << setw(base_digits) << setfill('0') << v.a[i
      return stream;
static vector<int> convert_base(const vector<int> &a, int
      old_digits, int new_digits) {
vector<long long> p(max(old_digits, new_digits) + 1);
      p[0] = 1;
      for (int i = 1; i < (int)p.size(); i++)

p[i] = p[i - 1] * 10;
      vector<int> res;
     | long long cur = 0;
| int cur_digits = 0;
| for (int i = 0; i < (int)a.size(); i++) {
| cur += a[i] * p[cur_digits];
| cur_digits += old_digits;
           while (cur_digits >= new_digits) {
    res.push_back(int(cur % p[new_digits]));
                 cur /= p[new_digits];
cur digits -= new digits;
           }
      res.push_back((int)cur);
      while (!res.empty() && !res.back()) res.pop_back();
      return res:
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a, const vll &b) {
     int n = a.size();
      vll res(n + n);
      \begin{aligned} & \text{vii } & \text{res}(i + m), \\ & \text{if } (n <= 32) \; \{ \\ & \text{for } (\text{int } i = 0; \, i < n; \, i++) \\ & \text{for } (\text{int } j = 0; \, j < n; \, j++) \\ & \text{res}[i + j] \; += a[i] \; * \; b[j]; \end{aligned} 
            return res;
      int k = n \gg 1;
     vll a1(a.begin(), a.begin() + k);
vll a2(a.begin() + k, a.end());
vll b1(b.begin(), b.begin() + k);
      vll b2(b.begin() + k, b.end());
     vll a1b1 = karatsubaMultiply(a1, b1);
vll a2b2 = karatsubaMultiply(a2, b2);
     \begin{array}{l} {\rm for} \ ({\rm int} \ i=0; \ i < k; \ i++) \ a2[i] \ += \ a1[i]; \\ {\rm for} \ ({\rm int} \ i=0; \ i < k; \ i++) \ b2[i] \ += \ b1[i]; \end{array}
      vll r = karatsubaMultiply(a2, b2);
     \begin{array}{lll} & \text{for (int } i=0; \ i<(iint)alb1.size(); \ i++) \ r[i] -=alb1[i]; \\ & \text{for (int } i=0; \ i<(iint)a2b2.size(); \ i++) \ r[i] -=a2b2[i]; \end{array}
      for (int i = 0; i < (int)r.size(); i++) res[i + k] += r[i];
```

```
for (int i = 0; i < (int)a1b1.size(); i++) res[i] += a1b1[i
        for (int i = 0; i < (int)a2b2.size(); i++) res[i + n] +=
               a2b2[i];
        return res;
    bigint operator*(const bigint &v) const {
        vector<int> a6 = convert_base(this->a, base_digits, 6);
vector<int> b6 = convert_base(v.a, base_digits, 6);
        vll x(a6.begin(), a6.end());
vll y(b6.begin(), b6.end());
        push\_back(0);
        vll c = karatsubaMultiply(x, y);
        bigint res:
        res.sign = sign * v.sign;
        for (int i = 0, carry = 0; i < (int)c.size(); i++) {
            \begin{array}{l} long\ long\ cur = c[i] + carry; \\ res.a.push\_back((int)(cur\ \%\ 1000000)); \end{array}
             carry = (int)(cur / 1000000);
        res.a = convert base(res.a, 6, base digits);
        res.trim():
        return res;
};
```

2 Data Structures

2.1 Disjoin Set Union

```
struct DSU {
    vector<int> par;
    vector<int> sz;
    DSU(int n) {
       FOR(i, 0, n) {
           par.pb(i);
           sz.pb(1);
    }
   int find(int a) {
        return \ par[a] = par[a] == a \ ? \ a : find(par[a]); 
    bool same(int a, int b) {
       return find(a) == find(b);
    void unite(int a, int b) {
       a = find(a);
       b = find(b);
       if(sz[a] > sz[b]) swap(a, b);
       sz[b] += sz[a];
       par[a] = b;
};
```

2.2 Fenwick Tree Point Update And Range Query

```
 \begin{array}{l} struct \; Fenwick \; \{ \\ vector < ll > \; tree; \\ int \; n; \\ Fenwick() \{ \} \\ Fenwick(int \_n) \; \{ \\ n = \_n; \\ tree = vector < ll > (n+1, \, 0); \\ \} \\ void \; add(int \; i, \; ll \; val) \; \{ \; // \; arr[i] \; += \; val \\ \; for(; \; i <= \; n; \; i \; += \; i\&(-i)) \; tree[i] \; += \; val; \\ \} \\ ll \; get(int \; i) \; \{ \; // \; arr[i] \\ \; return \; sum(i, \; i); \end{array}
```

2.3 Fenwick Tree Range Update And Point Query

2.4 Fenwick Tree Range Update And Range Query

2.5 Fenwick 2D

```
 \begin{array}{l} struct \; Fenwick2D \; \{ \\ vector < vector < ll >> bit; \\ int \; n, \; m; \\ Fenwick2D (int \_n, \; int \_m) \; \{ \\ n = \_n; \; m = \_m; \\ bit = vector < vector < ll >> (n+1, \; vector < ll >(m+1, \; 0)); \\ \} \\ ll \; sum (int \; x, \; int \; y) \; \{ \\ ll \; ret = \; 0; \\ for \; (int \; i = \; x; \; i > \; 0; \; i \; -= \; i \; \& \; (-i)) \\ for \; (int \; j = \; y; \; j > \; 0; \; j \; -= \; j \; \& \; (-j)) \\ ret \; += bit[i][j]; \\ \end{array}
```

2.6 Segment Tree

```
struct SegmentTree \{
      int n;
      vector<ll> t; const ll IDENTITY = 0; // OO for min, -OO for max, ...
      ll f(ll a, ll b) {
           return a+b:
      SegmentTree(int _n) {
           n = _n; t = vector < ll > (4*n, IDENTITY);
      SegmentTree(vector<ll>& arr) {
           n = arr.size(); t = vector < ll > (4*n, IDENTITY);
            build(arr, 1, 0, n-1);
      void build(vector<ll>& arr, int v, int tl, int tr) {
           if(tl == tr) \ \{ \ t[v] = arr[tl]; \ \}
           else {
                 int tm = (tl+tr)/2;
build(arr, 2*v, tl, tm);
build(arr, 2*v+1, tm+1, tr);
                 t[v] = f(t[2*v], t[2*v+1]);
      /// \ sum(1, \, 0, \, n\text{-}1, \, l, \, r) \\ ll \ sum(int \ v, \ int \ tl, \ int \ tr, \ int \ l, \ int \ r) \ \{ \\ if(l > r) \ return \ IDENTITY; 
            if(l == tl \&\& r == tr) return t[v];
           \begin{array}{l} \mathrm{int}\;\mathrm{tm}=(\mathrm{tl}+\mathrm{tr})/2;\\ \mathrm{return}\;\mathrm{f}(\mathrm{sum}(2^*\mathrm{v},\,\mathrm{tl},\,\mathrm{tm},\,\mathrm{l},\,\mathrm{min}(\mathrm{r},\,\mathrm{tm})),\,\mathrm{sum}(2^*\mathrm{v}+1,\,\mathrm{tm}) \end{array}
                     +1, tr, \max(l, tm+1), r));
      // update(1, 0, n-1, i, v) void update(int v, int tl, int tr, int pos, ll newVal) {
           if(tl == tr) \{ t[v] = newVal; \}
                 int tm = (tl+tr)/2;
                 if (pos \le tm) update(2*v, tl, tm, pos, newVal); else update(2*v+1, tm+1, tr, pos, newVal);
                 t[v] = f(t[2*v],t[2*v+1]);
     }
};
```

2.7 Segment Tree With Lazy Propagation

```
lazy[v] = 0;
   void update(int v, int tl, int tr, int l, int r, ll addend) {
       if (l > r)
          return;
       if (l == tl \&\& tr == r) {
          t[v] += addend;
          lazy[v] += addend;
       } else {
          push(v);
          int tm = (tl + tr) / 2;
          update(v*2, tl, tm, l, min(r, tm), addend);
          }
   }
   int query(int v, int tl, int tr, int l, int r) {
       if (l > r)
          return -OO;
       if (tl == tr)
          return t[v];
       push(v);\\
      int tm = (tl + tr) / 2;
return max(query(v*2, tl, tm, l, min(r, tm)),
              query(v*2+1, tm+1, tr, max(l, tm+1), r));
};
```

2.8 Treap

```
namespace Treap {
    struct Node {
        Node *l, *r;
        ll key, prio, size;
        \begin{split} & \text{Node() \{ \}} \\ & \text{Node(ll key) : key(key), l(nullptr), r(nullptr), size(1) \{ } \\ & \text{prio = rand() $^(\text{rand()} << 15)$}; \end{split}
    };
    typedef Node* NodePtr;
    int sz(NodePtr n) {
        return n ? n->size : 0;
    void recalc(NodePtr n) {
        if (!n) return:
        n-> size = sz(n->l) + 1 + sz(n->r); // add more
               operations here as needed
    void split(NodePtr tree, ll key, NodePtr& l, NodePtr& r) {
        if (!tree) {
            l = r = nullptr;
        else if (key < tree->key) {
             split(tree->l, key, l, tree->l);
             r = tree;
        else {
             split(tree->r, key, tree->r, r);
        recalc(tree);
    void merge(NodePtr& tree, NodePtr l, NodePtr r) {
        if (!l || !r) {
tree = l ? l : r;
        else if (l->prio > r->prio) \{
             merge(l-{>}r,\; l-{>}r,\; r);
             tree = l;
             merge(r->l, l, r->l);
             tree = r;
        recalc(tree):
    void insert(NodePtr& tree, NodePtr node) {
```

```
if (!tree) {
              tree = node;
         else if (node->prio > tree->prio) {
              split(tree, node->key, node->l, node->r);
              tree = node;
              insert(node->key < tree->key ? tree->l : tree->r,
                     node);
         recalc(tree);
    void erase(NodePtr tree, ll key) {
         if (!tree) return;
         if (tree->key == key) {
              merge(tree, tree->l, tree->r);
              erase(key < tree->key ? tree->l : tree->r, key);
         recalc(tree);
    }
    void print(NodePtr t, bool newline = true) {
         if (!t) return;
         \begin{array}{l} \operatorname{print}(t->l,\,\operatorname{false});\\ \operatorname{cout} << t-> \operatorname{key} << " \llcorner "; \end{array}
         print(t->r, false);
         if (newline) cout << endl;
}
```

2.9 Implicit Treap

```
template <tvpename T>
struct Node {
Node* l, *r;
    ll prio, size, sum;
    T val;
    bool rev
    Node() {}
Node(T _
                _val) : l(nullptr), r(nullptr), val(_val), size(1), sum(
             _val), rev(false) {
         \overline{\text{prio}} = \text{rand}() \cap (\text{rand}() << 15);
    }
template <tvpename T>
struct ImplicitTreap {
    typedef Node<T>* NodePtr;
    int sz(NodePtr n) {
         return n? n->size: 0;
    fl getSum(NodePtr n) {
         return n ? n->sum : 0;
    void\ push(NodePtr\ n)\ \{
         if (n && n->rev) {
             n->rev = false;
             swap(n->l, n->r);
             if (n->1) n->1->rev ^= 1;
             if (n->r) n->r->rev = 1;
    }
    void recalc(NodePtr n) {
         if (!n) return;
         n->size = sz(n->l) + 1 + sz(n->r);
         n->sum = getSum(n->l) + n->val + getSum(n->r);
    void split(NodePtr tree, ll key, NodePtr& l, NodePtr& r) {
         push(tree):
         if (!tree) {
             l = r = nullptr;
         else if (\text{key} \le \text{sz(tree-} > l)) {
             {
m split}({
m tree-}>l, {
m key}, {
m l}, {
m tree-}>l);
             r = tree:
             {\rm split}({\rm tree-}{>}{\rm r},\,{\rm key-sz}({\rm tree-}{>}{\rm l}){\rm -}1,\,{\rm tree-}{>}{\rm r},\,{\rm r});
```

```
l = tree;
         recalc(tree);
    }
    void merge(NodePtr& tree, NodePtr l, NodePtr r) {
         push(l); push(r);
         if (!l || !r) {
            tree = 1?1:r;
         else if (l->prio > r->prio) {
             merge(l->r, l->r, r);
             tree = l;
         else {
             merge(r->l, l, r->l);
             tree = r;
         recalc(tree);
    void insert(NodePtr& tree, T val, int pos) {
        if (!tree) {
             tree = new Node<T>(val);
             return;
         NodePtr L, R;
        split(tree, pos, L, R);
merge(L, L, new Node<T>(val));
merge(tree, L, R);
        recalc(tree);
    void reverse(NodePtr tree, int l, int r) {
         NodePtr t1, t2, t3;
split(tree, l, t1, t2);
split(t2, r - l + 1, t2, t3);
         if(t2) t2->rev = true;
         merge(t2, t1, t2);
         merge(tree, t2, t3);
    void print(NodePtr t, bool newline = true) {
        push(t);
         if (!t) return;
         print(t->l, false);
         \mathrm{cout} << \mathrm{t\text{--}val} << "_{\sqcup}";
         print(t->r, false);
         if (newline) cout << endl;
    NodePtr fromArray(vector<T> v) {
         NodePtr t = nullptr;
FOR(i, 0, (int)v.size()) {
             insert(t, v[i], i);
         return t;
    ll calcSum(NodePtr t, int l, int r) {
         NodePtr L, R;
         split(t, l, L, R);
         NodePtr good;
         \operatorname{split}(R, r - l + 1, \operatorname{good}, L);
         return getSum(good);
    }
};
/* Usage: ImplicitTreap<int> t;
Node<int> tree = t.fromArray(someVector); t.reverse(tree, l, r);
```

2.10 Trie

```
nodeCount = 1;
    void insert(const string& s, int id) {
        int curr = 0:
        FOR(i, 0, (int)s.length()) {
           int c = s[i] - BASE;
           if(nextNode[curr][c] == -1) {
               nextNode[curr][c] = nodeCount++;
           curr = nextNode[curr][c];
        mark[curr] = id;
   bool exists(const string& s) {
        int curr = 0;
       FOR(i, 0, (int)s.length()) {

int c = s[i] - BASE;
           if(nextNode[curr][c] == -1) return false;
           curr = nextNode[curr][c];
        return mark[curr] != -1;
};
```

3 Graphs

3.1 Dfs With Timestamps

```
\label{eq:constraint} \begin{split} & \operatorname{vector} < \operatorname{vector} < \operatorname{int} >> \operatorname{adj}; \\ & \operatorname{vector} < \operatorname{int} > \operatorname{tIn}, \ \operatorname{tOut}, \ \operatorname{color}; \\ & \operatorname{int} \ \operatorname{dfs\_timer} = 0; \\ & \operatorname{void} \ \operatorname{dfs}(\operatorname{int} \ v) \ \{ \\ & \operatorname{tIn}[v] = \operatorname{dfs\_timer} + +; \\ & \operatorname{color}[v] = 1; \\ & \operatorname{for} \ (\operatorname{int} \ u : \operatorname{adj}[v]) \\ & \operatorname{if} \ (\operatorname{color}[u] = = 0) \\ & \operatorname{dfs}(u); \\ & \operatorname{color}[v] = 2; \\ & \operatorname{tOut}[v] = \operatorname{dfs\_timer} + +; \\ \} \end{split}
```

3.2 Lowest Common Ancestor

```
int n, l; // l == logN (usually about ~20)
vector<vector<int>> adj;
int timer:
vector<int> tin, tout:
vector<vector<int>> up;
void dfs(int v, int p)
      {\rm tin}[v] = + + {\rm timer};
      up[v][0] = p;
      // \text{wUp[v][0]} = \text{weight[v][u]}; // <- \text{path weight sum to } 2^i-\text{th}
                 ancestor
      \begin{array}{l} \text{for (int $i=1$; $i<=1$; $++$i)} \\ up[v][i] &= up[up[v][i-1]][i-1]; \\ // & wUp[v][i] &= wUp[v][i-1] + wUp[up[v][i-1]][i-1]; \end{array}
     \begin{array}{c} \text{for } (\text{int } u : \text{adj}[v]) \ \{ \\ \text{if } (u \mathrel{!}=p) \end{array}
                  dfs(u, v);
      tout[v] = ++timer;
}
bool isAncestor(int u, int v)
      \operatorname{return} \ \operatorname{tin}[u] <= \operatorname{tin}[v] \ \&\& \ \operatorname{tout}[v] <= \operatorname{tout}[u];
int lca(int u, int v)
      if (isAncestor(u, v))
```

```
\begin{tabular}{ll} return $u$; \\ if $(isAncestor(v, u))$ \\ return $v$; \\ for $(int $i=1$; $i>=0$; $--i$) $\{$ \\ if $(!isAncestor(up[u][i], v))$ \\ $u = up[u][i]$; \\ $\}$ \\ return $up[u][0]$; \\ $\}$ \\ void $preprocess(int root) $\{$ \\ tin.resize(n)$; \\ tout.resize(n)$; \\ tout.resize(n)$; \\ timer = 0$; \\ $l = ceil(log2(n))$; \\ $up.assign(n, vector<int>(l+1))$; \\ $dfs(root, root)$; } $$\}
```

3.3 Strongly Connected Components

```
vector < vector
 < int> > g, gr; // adjList and reversed adjList
vector<bool> used;
vector < int > order, component;
void dfs1 (int v) {
     used[v] = true;
     for (size_t i=0; i<g[v].size(); ++i)
           if (!used[ g[v][i] ])
dfs1 (g[v][i]);
     order.push_back (v);
}
void dfs2 (int v) {
     used[v] = true;
     \begin{array}{l} {\rm component.push\_back} \ (v); \\ {\rm for} \ ({\rm size\_t} \ i{=}0; \ i{<}{\rm gr[v].size}(); \ +{+}i) \\ {\rm if} \ (!{\rm used[} \ {\rm gr[v][i]} \ ]) \end{array}
                dfs2 (gr[v][i]);
\mathrm{int}\ \mathrm{main}()\ \{
     int n;
      // read n
     for (;;) {
           int a, b;
            // read edge a -> b
           g[a].push_back (b);
           gr[b].push_back (a);
     used.assign (n, false);
     for (int i=0; i< n; ++i)
           if\ (!used[i])
     \begin{array}{c} \operatorname{dfs1}(i);\\ \operatorname{used.assign}(n,\,\operatorname{false});\\ \operatorname{for}(\operatorname{int}i=0;\,i<\!n;\,+\!+\!i) \end{array} \{
           int v = order[n-1-i];
           if (!used[v]) {
                 dfs2(v);
                 // do something with the found component
                 component.clear(); // components are generated in
                         toposort-order
           }
     }
```

3.4 Bellman Ford Algorithm

3.5 Bipartite Graph

```
class Bipartite
Graph \{
private:
     vector<int> _left, _right;
vector<vector<int>> _adjList;
vector<int> _matchR, _matchL;
     vector < bool > used;
     bool kuhn(int v) {
         if (_used[v]) return false;
_used[v] = true;
          FOR(i, 0, (int)_adjList[v].size()) {
               \begin{array}{ll} R(t, 0, (lint)\_adjList[v].size()) \ \{\\ int \ to = \_adjList[v][i] - \_left.size();\\ if \ (\_matchR[to] == -1 \ || \_kuhn(\_matchR[to])) \ \{\\ \_matchR[to] = v;\\ \_matchL[v] = to; \end{array} 
                   return true;
          return false;
             addReverseEdges() {
     void
          FOR(i, 0, (int)_right.size()) {
    if (_matchR[i] != -1) {
                    \_adjList[\_left.size() + i].pb(\_matchR[i]);
          }
     void _dfs(int p) {
          if (_used[p]) return;
            \underline{used[p]} = true;
          for (auto x : _adjList[p]) {
               _{dfs(x);}
     vector<pii> _buildMM() {
          vector<pair<int, int> > res;
          FOR(i, 0, (int)_right.size()) {
              res.push\_back(make\_pair(\_matchR[i],\,i));
               }
          }
          return res;
public:
     void addLeft(int x) {
           left.pb(x):
          _adjList.pb({});
          _{\rm matchL.pb(-1)};
          \_used.pb(false);
     void addRight(int x) {
         _{\rm right.\widetilde{pb}(x)};
          \underline{\text{adjList.pb}}(\{\});
          _matchR.pb(-1);
          _used.pb(false);
     void addForwardEdge(int l, int r) {
           _{adjList[l].pb(r + \_left.size())}
     void addMatchEdge(int l, int r) {
          if(l != -1) _{matchL[l]} = r;

if(r != -1) _{matchR[r]} = l;
     // Maximum Matching
     vector<pii> mm() {
          _matchR = vector<int>(_right.size(), -1);
          _{\text{matchL}} = \text{vector} < \text{int} > (_{\text{left.size}}(), -1);
```

```
// \hat{} these two can be deleted if performing MM on
             already partially matched graph
       used = vector<bool>(_left.size() + _right.size(), false
             );
     bool path_found;
          fill(_used.begin(), _used.end(), false);
          path_found = false;
FOR(i, 0, (int) left.size()) {
               if (_matchL[i] < 0 && !_used[i]) {
    path_found |= _kuhn(i);
     } while (path_found);
     return _buildMM();
// Minimum Edge Cover
// Algo: Find MM, add unmatched vertices greedily.
vector<pii> mec() {
     auto ans = mm();
     FOR(i, 0, (int)_left.size()) {
          if (_matchL[i] != -1) {
               __matchE[i] := 1) {
    int ridx = x - _left.size();
    if (_matchR[ridx] == -1) {
        ans.pb({ i, ridx });
    }
                          \underline{\text{matchR}[\text{ridx}] = i};
          }
      \begin{array}{l} FOR(i,\,0,\,(int)\_left.size())\;\{\\ if\;(\_matchL[i]==-1\;\&\&\;(int)\_adjList[i].size()>0) \end{array} 
               \label{eq:int_ridx} \begin{split} &\inf \ ridx = \_adjList[i][0] \ - \_left.size(); \\ &\_matchL[i] = ridx; \end{split}
               ans.pb(\{i, ridx\});
          }
     return ans:
  / Minimum Vertex Cover
// Algo: Find MM. Run DFS from unmatched vertices from
        the left part.
// MVC is composed of unvisited LEFT and visited RIGHT
pair < vector < int >, \ vector < int >> \ mvc(bool \ runMM = true)
     if (runMM) mm();
       _addReverseEdges();
     _addreverseEdges(),
fill(_used.begin(), _used.end(), false);
FOR(i, 0, (int)_left.size()) {
          if (\underline{matchL[i]} == -1)  {
               _dfs(i);
     vector<int> left, right;
     FOR(i, 0, (int)_left.size()) {
    if (!_used[i]) left.pb(i);
     \begin{aligned} & \text{FOR}(i, \ 0, \ (int)\_right.size()) \ \{ \\ & \text{if} \ (\_used[i + (int)\_left.size()]) \ right.pb(i); \end{aligned}
     return { left,right };
// Maximal Independent Vertex Set
// Algo: Find complement of MVC.
pair<vector<int>, vector<int>> mivs(bool runMM = true)
     auto m = mvc(runMM);
     vector{<}bool{>} containsL(\_left.size(), false), containsR(
     \begin{tabular}{ll} $\tt \_right.size(), false); \\ for (auto $x : m.first) containsL[x] = true; \\ for (auto $x : m.second) containsR[x] = true; \\ \end{tabular}
     vector<int> left, right;
     FOR(i, 0, (int)\_left.size())
          if \ (!containsL[i]) \ left.pb(i); \\
     FOR(i, 0, (int)_right.size()) {
   if (!containsR[i]) right.pb(i);
     return { left, right };
```

```
};
```

3.6 Finding Articulation Points

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector < bool > visited;
vector<int> tin, fup;
int timer:
void processCutpoint(int v) {
      // problem-specific logic goes here
      // it can be called multiple times for the same v
void dfs(int v, int p = -1) {
      visited[v] = true;

tin[v] = fup[v] = timer++;

int children=0;
      for (int to : adj[v]) {
           if (to == p) continue; if (visited[to]) {
                fup[v] = min(fup[v], tin[to]);
           } else {
                 dfs(to, v);
                \begin{aligned} & \text{fup[v]} = \min(\text{fup[v]}, \text{ fup[to]}); \\ & \text{if } (\text{fup[to]} >= \text{tin[v]} \&\& \text{ p!=-1}) \\ & \text{processCutpoint(v)}; \end{aligned}
                 ++children;
      if(p == -1 \&\& children > 1)
           processCutpoint(v);
void findCutpoints() {
      visited.assign(n, false);
      tin.assign(n,\, \text{-}1);
     \begin{array}{l} \mathrm{fup.assign}(n,\,\text{-}1);\\ \mathrm{for}\;(\mathrm{int}\;i=0;\,i< n;\,++i)\;\{ \end{array}
           if (!visited[i])
                 dfs (i);
}
```

3.7 Finding Bridges

```
int n: // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<br/>bool> visited;
vector<int> tin, fup;
int timer;
void processBridge(int u, int v) {
    // do something with the found bridge
void dfs(int v, int p = -1) {
     \begin{array}{l} visited[v] = true; \\ tin[v] = fup[v] = timer++; \\ for (int to : adj[v]) \end{array} 
         if (to == p) continue;
         if (visited[to]) {
             fup[v] = min(fup[v], tin[to]);
         } else {
             dfs(to, v); 
fup[v] = min(fup[v], fup[to]);
             if (fup[to] > tin[v])
                 processBridge(v, to);
        }
    }
}
// Doesn't work with multiple edges
// But multiple edges are never bridges, so it's easy to check
```

```
 \begin{aligned} & \text{void findBridges}() \ \{ \\ & \text{timer} = 0; \\ & \text{visited.assign}(n, \text{false}); \\ & \text{tin.assign}(n, -1); \\ & \text{fup.assign}(n, -1); \\ & \text{bridges.clear}(); \\ & \text{FOR}(i, 0, n) \ \{ \\ & \text{if (!visited[i])} \\ & \text{dfs}(i); \\ \} \end{aligned}
```

3.8 Max Flow With Ford Fulkerson

```
struct Edge {
     int to, next;
     ll f, c;
     int idx, dir:
     int from:
};
int n, m;
{\tt vector}{<}{\tt Edge}{\tt > edges};
vector < int > first;
void addEdge(int a, int b, ll c, int i, int dir) {
      \begin{array}{l} edges.pb(\{\ b,\ first[a],\ 0,\ c,\ i,\ dir,\ a\ \});\\ edges.pb(\{\ a,\ first[b],\ 0,\ 0,\ i,\ dir,\ b\ \});\\ first[a] = edges.size()\ -\ 2; \end{array} 
     first[b] = edges.size() - 1;
}
void init() {
     cin >> n >> m;
edges.reserve(4 * m);
     first = vector < int > (n, -1);
     FOR(i, 0, m) {
          int a, b, c;
          cin >> a >> b >> c;
          addEdge(a, b, c, i, 1);
          addEdge(b, a, c, i, -1);
}
int cur\_time = 0;
vector<int> timestamp;
ll dfs(int v, ll flow = OO) {
      \begin{array}{l} if\ (v == n-1)\ return\ flow;\\ timestamp[v] = cur\_time;\\ for\ (int\ e = first[v];\ e != -1;\ e = edges[e].next)\ \{ \end{array} 
          if (edges[e].f < edges[e].c && timestamp[edges[e].to] !=
               cur_time) {
int pushed = dfs(edges[e].to, min(flow, edges[e].c -
               edges[e].f);
if (pushed > 0) {
  edges[e].f += pushed;
  edges[e ^ 1].f -= pushed;
                     return pushed;
          }
     return 0;
}
ll \max Flow() \{
     cur\_time = 0;
     timestamp = vector < int > (n, 0);
     ll f = 0, add;
     while (true) {
          cur\_time++;
          add = dfs(0);
          if (add > 0)
               f += add
          élse {
               break;
     return f:
```

3.9 Max Flow With Dinic

```
struct Edge \{
      int f, c;
     int to:
     pii revIdx;
     int dir;
      int idx;
};
\begin{array}{l} {\rm int}\ n,\ m;\\ {\rm vector}{<}{\rm Edge}{>}\ {\rm adjList[MAX\_N]}; \end{array}
int level[MAX_N];
void addEdge(int a, int b, int c, int i, int dir) {
     int idx = adjList[a].size();
int revIdx = adjList[b].size();
adjList[a].pb({ 0,c,b, {b, revIdx}, ,dir,i });
adjList[b].pb({ 0,0,a, {a, idx}, ,dir,i });
bool bfs(int s, int t) {
     FOR(i, 0, n) level[i] = -1;
level[s] = 0;
queue<int> Q;
      Q.push(s);
      while (!Q.empty()) {
            auto t = Q.front(); Q.pop();
            for (auto x : adjList[t]) {
                 \begin{array}{l} \text{if } (\text{level[x.to]} < 0 \text{ \&\& x.f} < \text{x.c}) \text{ } \{\\ \text{level[x.to]} = \text{level[t]} + 1; \end{array}
                        Q.push(x.to);
           }
      return level[t] >= 0;
int send(int u, int f, int t, vector<int>& edgeIdx) {
      if (u == t) return f;
     for (; edgeIdx[u] < adjList[u].size(); edgeIdx[u]++) {
    auto& e = adjList[u][edgeIdx[u]];
    if (level[e.to] == level[u] + 1 && e.f < e.c) {
        int curr_flow = min(f, e.c - e.f);
    }
                 int next_flow = send(e.to, curr_flow, t, edgeIdx);
                 if (\text{next\_flow} > 0) {
                       e.f += next_flow;
                       adjList[e.revIdx.first][e.revIdx.second].f -=
                                next\_flow;
                       return next flow;
           }
      return 0;
int maxFlow(int s, int t) {
      while (bfs(s, t)) {
            vector < int > edgeIdx(n, 0);
            while (int extra = send(s, oo, t, edgeIdx)) {
                 f += extra:
      return f;
\begin{array}{c} \mathrm{void\ init}()\ \{\\ \mathrm{cin}>>n>>m; \end{array}
      FOR(i, 0, m) {
           int a, b, c;
           \mathrm{cin} >> \mathrm{a} >> \mathrm{b} >> \mathrm{c};
            a--; b--;
           addEdge(a, b, c, i, 1);
           addEdge(b, a, c, i, -1);
     }
}
```

3.10 Max Flow With Dinic 2

```
struct FlowEdge {
   int v, u;
   long long cap, flow = 0;
```

```
FlowEdge(int\ v,\ int\ u,\ long\ long\ cap): v(v),\ u(u),\ cap(cap)
};
struct Dinic {
    const long long flow_inf = 1e18;
     vector<FlowEdge> edges
     vector<vector<int>> adj;
    int n, m = 0;
    int s, t;
    vector<int> level, ptr;
    queue<int> q;
    Dinic(int n, int s, int t): n(n), s(s), t(t) {
          adj.resize(n);
         level.resize(n);
         ptr.resize(n);
    void add_edge(int v, int u, long long cap) \{
         edges.push\_back(FlowEdge(v,\,u,\,cap));\\
         edges.push\_back(FlowEdge(u,\,v,\,0));
         adj[v].push_back(m);
adj[u].push_back(m + 1);
m += 2;
    bool bfs() {
    while (!q.empty()) {
        int v = q.front();
    }
}
              q.pop();
              for (int id : adj[v]) {
                   if (edges[id].cap - edges[id].flow < 1)
                        continue;
                   if\ (level[edges[id].u] \mathrel{!=-1})
                        continue:
                   level[edges[id].u] = level[v] + 1;
                   q.push(edges[id].u);
              }
         return level[t] != -1;
    long long dfs(int v, long long pushed) {
          if (pushed == 0)
              return 0;
         if (v == t)
              return pushed;
         for (int& cid = ptr[v]; cid < (int)adj[v].size(); cid++) {
              int id = adj[v][cid];
              int u = edges[id].u;
               \begin{array}{l} \text{if } (\operatorname{level}[v]+1 \stackrel{!}{=} \operatorname{level}[u] \mid\mid \operatorname{edges}[\operatorname{id}].\operatorname{cap-edges}[\operatorname{id}]. \\ \operatorname{flow} < 1) \end{array} 
               \begin{array}{l} continue; \\ long\ long\ tr = dfs(u,\,min(pushed,\,edges[id].cap\ \text{-} \end{array} 
                     edges[id].flow));
              if (tr == 0)
                   continue;
              edges[id].flow += tr;
              edges[id^ 1].flow -= tr;
              return tr;
         return 0;
    \begin{array}{c} long\ long\ flow()\ \{\\ long\ long\ f=0; \end{array}
          while (true) {
              fill(level.begin(), level.end(), -1);
              level[s] = 0;
              q.push(s);
              if (!bfs())
                   break:
              fill(ptr.begin(), ptr.end(), 0);
              while (long long pushed = dfs(s, flow_inf)) {
                   f += pushed;
         return f:
};
```

3.11 Min Cut

```
init();
ll f = maxFlow(); // Ford-Fulkerson
cur_time++;
dfs(0);
set<int> cc;
for (auto e : edges) {
    if (timestamp[e.from] == cur_time && timestamp[e.to] !=
        cur_time) {
        cc.insert(e.idx);
    }
}
// (# of edges in min-cut, capacity of cut)
// [indices of edges forming the cut]
cout << cc.size() << "\" "<< f << endl;
for (auto x : cc) cout << x + 1 << "\";</pre>
```

3.12 Number Of Paths Of Fixed Length

Let G be the adjacency matrix of a graph. Then $C_k = G^k$ gives a matrix, in which the value $C_k[i][j]$ gives the number of paths between i and j of length k.

3.13 Shortest Paths Of Fixed Length

Define $A \bigcirc B = C \iff C_{ij} = \min_{p=1..n} (A_{ip} + B_{pj})$. Let G be the adjacency matrix of a graph. Also, let $L_k = G \bigcirc \ldots \bigcirc G = G^{\bigcirc k}$. Then the value $L_k[i][j]$ denotes the length of the shortest path between i and j which consists of exactly k edges.

4 Geometry

4.1 2d Vector

```
template <typename T>
struct Vec {
    T x, y;
    Vec(): x(0), y(0) {}
Vec(T _x, T _y): x(_x), y(_y) {}
Vec operator+(const Vec& b) {
       return Vec < T > (x+b.x, y+b.y);
    Vec operator-(const Vec& b) {
       return Vec < T > (x-b.x, y-b.y);
    .
Vec operator*(T c) {
       return Vec(x*c, y*c);
    T operator*(const Vec& b) {
        return x*b.x + y*b.y;
    T operator (const Vec& b) {
       return x*b.y-y*b.x;
    bool operator<(const Vec& other) const {
       if(x == other.x) return y < other.y;
        return x < other.x;
    bool operator==(const Vec& other) const {
       return x==other.x && y==other.y;
    bool operator!=(const Vec& other) const {
       return !(*this == other);
    friend ostream& operator<<(ostream& out, const Vec& v) {
       return out << "(" << v.x << ",_{\sqcup}" << v.y << ")";
    friend istream& operator>>(istream& in, Vec<T>& v) {
       \mathrm{return}\ \mathrm{in} >> \mathrm{v.x} >> \mathrm{v.y};
    T norm() { // squared length return (*this)*(*this);
```

4.2 Line

```
template <typename T>
struct Line { // expressed as two vectors
    Vec<T> start, dir;
    Line() {}
    Line(Vec<T> a, Vec<T> b): start(a), dir(b-a) {}

    Vec<ld> intersect(Line l) {
        id t = ld((l.start-start)^l.dir)/(dir^l.dir);
        // For segment-segment intersection this should be in
            range [0, 1]
        Vec<ld> res(start.x, start.y);
        Vec<ld> dirld(dir.x, dir.y);
        return res + dirld*t;
    }
};
```

4.3 Convex Hull Gift Wrapping

```
vector<Vec<int>> buildConvexHull(vector<Vec<int>>& pts)
   int n = pts.size();
    sort(pts.begin(), pts.end());
   auto \operatorname{currP} = \operatorname{pts}[0]; // choose some extreme point to be on
           the hull
    vector < Vec < int >> hull;
    set<Vec<int>> used;
   hull.pb(pts[0]);
    used.insert(pts[0]);
    while(true) {
        auto candidate = pts[0]; // choose some point to be a
               candidate
        auto currDir = candidate-currP;
        vector<Vec<int>> toUpdate;
        FOR(i, 0, n) {
   if(currP == pts[i]) continue;
             // currently we have currP->candidate
// we need to find point to the left of this
             auto possibleNext = pts[i];
            auto nextDir = possibleNext - currP;
auto cross = currDir ^ nextDir;
if(candidate == currP || cross > 0) {
                 candidate = possibleNext;
                 currDir = nextDir;
             } else if(cross == 0 && nextDir.norm() > currDir.
                   \mathrm{norm}())\ \{
                 candidate = possible Next; \\
                 currDir = nextDir;
        if(used.find(candidate) != used.end()) break;
        hull.pb(candidate);
        used.insert(candidate);
currP = candidate;
    return hull;
```

4.4 Convex Hull With Graham's Scan

```
/ Takes in >= 3 points
// Returns convex hull in clockwise order
// Ignores points on the border
vector<Vec<int>> buildConvexHull(vector<Vec<int>> pts) {
   if(pts.size() <= 3) return pts;
   sort(pts.begin(), pts.end());
stack<Vec<int>> hull;
   hull.push(pts[0]);\\
   auto p = pts[0]; sort(pts.begin()+1, pts.end(), [&](Vec<int> a, Vec<int> b)
          -> bool {
        // p->a->b is a ccw turn
       int turn = sgn((a-p)^(b-a));
        //if(turn == 0) return (a-p).norm() > (b-p).norm();
            among collinear points, take the farthest one
       return turn == 1;
   hull.push(pts[1]);
FOR(i, 2, (int)pts.size()) {
       auto c = pts[i];
       if(c == hull.top()) continue;
       while(true) {
   auto a = hull.top(); hull.pop();
           auto b = hull.top();
           auto ba = a-b;
           auto ac = c-a;
           if((ba^ac)>0)\ \{
               hull.push(a);
               break:
           } else if((ba^ac) == 0) {
               if(ba*ac < 0) c = a;
               // ^ c is between b and a, so it shouldn't be
                     added to the hull
               break:
           }
       hull.push(c);
    vector<Vec<int>> hullPts;
   while(!hull.empty()) {
       hullPts.pb(hull.top());
       hull.pop();
   return hullPts;
```

4.5 Circle Line Intersection

```
double r, a, b, c; // ax+by+c=0, radius is at (0, 0) // If the center is not at (0, 0), fix the constant c to translate everything so that center is at (0, 0) double x0 = -a^*c/(a^*a+b^*b), y0 = -b^*c/(a^*a+b^*b); if (c^*c > r^*r^*(a^*a+b^*b)+eps) puts ("no_{\square}points"); else if (abs\ (c^*c - r^*r^*(a^*a+b^*b)) < eps) { puts ("1_{\square}point"); cout << x0 << '_{\square}' << y0 << '\setminus n'; } else { double d = r^*r - c^*c/(a^*a+b^*b); double mult = sqrt\ (d\ /\ (a^*a+b^*b)); double ax, ay, bx, by; ax = x0 + b * mult; bx = x0 - b * mult; ay = y0 - a * mult; by = y0 + a * mult; puts ("2_{\square}points"); cout << ax << '_{\square}' << ay << '\setminus n' << bx << '_{\square}' << by << '\n'; } }
```

4.6 Circle Circle Intersection

Let's say that the first circle is centered at (0,0) (if it's not, we can move the origin to the center of the first circle and adjust the coordinates), and the

second one is at (x_2, y_2) . Then, let's construct a line Ax + By + C = 0, where $A = -2x_2$, $B = -2y_2$, $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$. Finding the intersection between this line and the first circle will give us the answer. The only tricky case: if both circles are centered at the same point. We handle this case separately.

4.7 Common Tangents To Two Circles

```
struct pt {
    double x, y;
    pt operator- (pt p) {
         pt\ res = \{\ x\hbox{-p.x},\ y\hbox{-p.y}\ \};
         return res;
struct circle : pt {
    double r;
struct line {
    \quad \text{double $\hat{a}$, b, c;}
void tangents (pt c, double r1, double r2, vector<line> & ans) {
    double r = r2 - r1;
    double z = sqr(c.x) + sqr(c.y);
    double d = z - sqr(r); if (d < -eps) return;
    d = sqrt (abs (d));
    l.a = (c.x * r + c.y * d) / z;
l.b = (c.y * r - c.x * d) / z;
    l.c = r1;
    ans.pus\dot{h}\_back~(l);
vector<line> tangents (circle a, circle b) {
     vector<line> ans;
    for (int i=-1; i<=1; i+=2)
         for (int j=-1; j<=1; j+=2)
    tangents (b-a, a.r*i, b.r*j, ans);
for (size_t i=0; i<ans.size(); ++i)
ans[i].c -= ans[i].a * a.x + ans[i].b * a.y;
    return ans;
}
```

4.8 Number Of Lattice Points On Segment

Let's say we have a line segment from (x_1, y_1) to (x_2, y_2) . Then, the number of lattice points on this segment is given by

$$gcd(x_2-x_1,y_2-y_1)+1.$$

4.9 Pick's Theorem

We are given a lattice polygon with non-zero area. Let's denote its area by S, the number of points with integer coordinates lying strictly inside the polygon by I and the number of points lying on the sides of the polygon by B. Then:

$$S = I + \frac{B}{2} - 1.$$

5 Math

5.1 Linear Sieve

```
\begin{split} & ll \; minDiv[MAXN+1]; \\ & vector < ll > primes; \\ & void \; sieve(ll \; n) \{ \\ & \; FOR(k, \; 2, \; n+1) \{ \\ & \; minDiv[k] = k; \\ \} \\ & \; FOR(k, \; 2, \; n+1) \; \{ \\ & \; if(minDiv[k] = k) \; \{ \\ & \; primes.pb(k); \\ \} \\ & \; for(auto \; p : \; primes) \; \{ \\ & \; if(p > minDiv[k]) \; break; \\ & \; if(p > minDiv[k]) \; break; \\ & \; minDiv[p*k] = p; \\ \} \\ \} \\ & \} \end{split}
```

5.2 Extended Euclidean Algorithm

```
// ax+by=gcd(a,b)
void solveEq(ll a, ll b, ll& x, ll& y, ll& g) {
    if(b==0) {
       x = 1;
       y = 0;
        g = a;
       return;
   solveEq(b, a%b, xx, yy, g);
   x = yy;
   y = xx-yy*(a/b);
// ax + by = c
bool solveEq(ll a, ll b, ll c, ll& x, ll& y, ll& g) {
    solveEq(a, b, x, y, g);
   if(c%g!=0) return false;

x = c/g; y = c/g;
   return true;
// Finds a solution (x, y) so that x >= 0 and x is minimal
bool solveEqNonNegX(ll a, ll b, ll c, ll& x, ll &y, ll& g) {
    if(!solveEq(a, b, c, x, y, g)) return false;
   ll k = x*g/b;

x = x - k*b/g;
   y = y + k*a/g;
if(x < 0) {
       x + = b/g;
    return true:
```

All other solutions can be found like this:

$$x' = x - k \frac{b}{q}, y' = y + k \frac{a}{q}, k \in \mathbb{Z}$$

5.3 Chinese Remainder Theorem

Let's say we have some numbers m_i , which are all mutually coprime. Also, let $M = \prod_i m_i$. Then the system of congruences

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \dots \\ x \equiv a_k \pmod{m_k} \end{cases}$$

is equivalent to $x \equiv A \pmod{M}$ and there exists a unique number A satisfying $0 \le A \le M$.

Solution for two: $x \equiv a_1 \pmod{m_1}$, $x \equiv a_2 \pmod{m_2}$. Let $x = a_1 + km_1$. Substituting into the second congruence: $km_1 \equiv a_2 - a_1 \pmod{m_2}$. Then, $k = (m_1)_{m_2}^{-1}(a_2 - a_1) \pmod{m_2}$. and we can easily find x. This can be extended to multiple equations by solving them one-by-one.

If the moduli are not coprime, solve the system $y \equiv 0 \pmod{\frac{m_1}{g}}, y \equiv \frac{a_2-a_1}{g} \pmod{\frac{m_2}{g}}$ for y. Then let $x \equiv gy + a_1 \pmod{\frac{m_1m_2}{g}}$.

5.4 Euler Totient Function

5.5 Factorization With Sieve

```
// Use linear sieve to calculate minDiv
vector<pll> factorize(ll x) {
   vector<pll> res;
   ll prev = -1;
   ll cnt = 0:
   while(x != 1) {
       ll d = minDiv[x];
       if(d == prev) {
          cnt++;
       } else {
          if(prev != -1) res.pb(\{prev, cnt\});
          prev=d;
          cnt = 1:
       x /= d;
   res.pb({prev, cnt});
   return res:
```

5.6 Modular Inverse

```
\label{eq:bool invWithEuclid(ll a, ll m, ll& aInv) } \{ & ll x, y, g; \\ & if(!solveEqNonNegX(a, m, 1, x, y, g)) \ return false; \\ & aInv = x; \\ & return true; \\ \} \\ // \ Works \ only \ if \ m \ is \ prime \\ ll \ invFermat(ll \ a, \ ll \ m) \left\{ \\ & return \ pwr(a, \ m-2, \ m); \\ \} \\ // \ Works \ only \ if \ gcd(a, \ m) = 1 \\ ll \ invEuler(ll \ a, \ ll \ m) \left\{ \\ & return \ pwr(a, \ phi(m)-1, \ m); \\ \} \\ \}
```

5.7 Simpson Integration

5.8 Burnside's Lemma

Let G be a finite group that acts on a set X. For each g in G let X^g denote the set of elements in X that are fixed by g. Burnside's lemma asserts the following formula for the number of orbits:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

Example. Coloring a cube with three colors.

Let X be the set of 3^6 possible face color combinations. Let's count the sizes of the fixed sets for each of the 24 rotations:

- one 0-degree rotation which leaves all 3^6 elements of X unchanged
- six 90-degree face rotations, each of which leaves 3^3 elements of X unchanged
- three 180-degree face rotation, each of which leaves 3^4 elements of X unchanged
- eight 120-degree vertex rotations, each of which leaves 3^2 elements of X unchanged
- six 180-degree edge rotations, each of which leaves 3^3 elements of X unchanged

```
The average is then \frac{1}{24}(3^6 + 6 \cdot 3^3 + 3 \cdot 3^4 + 8 \cdot 3^2 + 6 \cdot 3^3) = 57. For n colors: \frac{1}{24}(n^6 + 3n^4 + 12n^3 + 8n^2).
```

Example. Coloring a circular stripe of n cells with two colors.

X is the set of all colored striped (it has 2^n elements), G is the group of rotations (n elements - by 0 cells, by 1 cell, ..., by (n-1) cells). Let's fix some K and find the number of stripes that are fixed by the rotation by K cells. If a stripe becomes itself after rotation by K cells, then its 1st cell must have the same color as its $(1 + K \mod n)$ -th cell, which is in turn the same as its $(1 + 2K \mod n)$ -th cell, etc., until $mK \mod n = 0$. This will happen when $m = n/\gcd(K, n)$. Therefore, we have $n/\gcd(K, n)$ cells that must all be of the same color. The same will happen when starting from the second cell and so on. Therefore, all cells are separated into qcd(K,n)groups, with each group being of one color, and that yields $2^{gcd(K,n)}$ choices. That's why the answer to the original problem is $\frac{1}{n} \sum_{k=0}^{n-1} 2^{gcd(k,n)}$.

5.9 FFT

```
name
space FFT \{
    int n;
     vector<int> r;
     vector<complex<ld>> omega;
     int logN, pwrN;
    void\ initLogN()\ \{
          logN = 0;
          pwrN = 1;
          while (pwrN < n) {
               pwrN *= 2;
               logN++;
          n = pwrN;
    }
     void initOmega() {
          FOR(i, 0, pwrN) {
              omega[i] = {cos(2 * i*PI / n), sin(2 * i*PI / n)};
    }
     void initR() {
          FOR(i, 1, pwrN) {

r[i] = r[i / 2] / 2 + ((i \& 1) << (logN - 1));
    }
     void initArrays() {
          r.clear();
          r.resize(pwrN);
          omega.clear();
          omega.resize(pwrN);
    }
     void init(int n) {
          FFT::n = n;
          initLogN();
          init Arrays():
          initOmega();
          initR();
    \label{eq:complex} \begin{tabular}{ll} void fft(complex<ld>d> a[], complex<ld>f[]) \\ \end{tabular}
          FOR(i, 0, pwrN) 
f[i] = a[r[i]];
          for (ll k = 1; k < pwrN; k *= 2) {
               for (ll i = 0; i < pwrN; i += 2 * k) {
                     \begin{array}{l} (\text{li } 1 - 0, 1 \times \text{pwith}, 1 + - 2 \times \text{k}) \, \{ \\ \text{for (ll } j = 0; \, j < k; \, j + +) \, \{ \\ \text{auto } z = \text{omega[j*n} \, / \, (2 * k)] * f[i + j + k]; \\ f[i + j + k] = f[i + j] - z; \\ f[i + j] + = z; \end{array} 
                   }
              }
         }
   }
```

5.10 FFT With Modulo

```
 \begin{aligned} & \text{bool isGenerator(ll g) } \{ \\ & \text{if } (pwr(g,\,M\,-\,1)\,!=\,1) \text{ return false;} \\ & \text{for } (ll\,\,i=\,2;\,\,i^*i\,<=\,M\,-\,1;\,\,i++)\,\, \{ \\ & \text{if } ((M\,-\,1)\,\,\%\,\,i\,\,:=\,0)\,\, \{ \\ & \text{ll } q=i; \\ & \text{if } (\text{isPrime}(q))\,\, \{ \\ & \text{ll } p=(M\,-\,1)\,/\,\,q; \\ & \text{ll } pp=pwr(g,\,p); \\ & \text{if } (pp==\,1) \text{ return false;} \\ & \} \\ & q=\,(M\,-\,1)\,/\,i; \\ & \text{if } (\text{isPrime}(q))\,\, \{ \\ & \text{ll } p=\,(M\,-\,1)\,/\,\,q; \\ & \text{ll } pp=pwr(g,\,p); \\ & \text{if } (pp==\,1) \text{ return false;} \\ & \} \\ & \} \\ & \} \\ & \} \\ & \} \\ \end{aligned}
```

```
namespace FFT {
     ll n;
      vector<ll> r;
      vector<ll> omega;
     ll\ logN,\ pwrN;
      void\ initLogN()\ \{
           logN = 0;
           pwrN = 1;
           while (pwrN < n) {
    pwrN *= 2;
                 logN++;
           n = pwrN;
      void\ initOmega()\ \{
            while (!isGenerator(g)) g++;
           ll G = 1:
            g = pwr(g, (M - 1) / pwrN);
            FOR(i, 0, pwrN) {
                omega[i] = G;
                 G \% = M;
           }
     }
      void initR() {
           FOR(i, 1, pwrN) {

r[i] = r[i / 2] / 2 + ((i \& 1) << (logN - 1));
      void\ initArrays()\ \{
           r.clear();
           r.resize(pwrN);
           omega.clear();
           \widetilde{\mathrm{omega.resize(pwrN)}};
      void\ init(ll\ n)\ \{
           FFT::n = n:
            initLogN();
           initArrays();
           initOmega();
           initR();
      \begin{array}{l} {\rm void\ fft(ll\ a[],\ ll\ f[])\ \{} \\ {\rm \ for\ (ll\ i=0;\ i< pwrN;\ i++)\ \{} \\ {\rm \ f[i]=a[r[i]];} \end{array} 
           for (ll k = 1; k < pwrN; k *= 2) {
    for (ll i = 0; i < pwrN; i += 2 * k) {
        for (ll j = 0; j < k; j++) {
            auto z = omega[j*n / (2 * k)] * f[i + j + k] %
        }
                                       M;
                             f[i + j + k] = f[i + j] - z;
                            \begin{array}{l} f[i+j+k] = f[i+j] = 2; \\ f[i+j] + = z; \\ f[i+j+k] \% = M; \\ \text{if } (f[i+j+k] < 0) \ f[i+j+k] + = M; \\ f[i+j] \% = M; \end{array}
                      }
               }
          }
     }
}
```

5.11 Big Integer Multiplication With FFT

```
\begin{split} & complex {<} ld {>} \ a[MAX\_N], \ b[MAX\_N]; \\ & complex {<} ld {>} \ fa[MAX\_N], \ fb[MAX\_N], \ fc[MAX\_N]; \\ & complex {<} ld {>} \ cc[MAX\_N]; \\ & string \ mul(string \ as, \ string \ bs) \ \{ \\ & int \ sgn1 = 1; \\ & int \ sgn2 = 1; \\ \end{split}
```

```
if (as[0] == '-') {
     sgn1 = -1;
     as = as.substr(1);
if (bs[0] == '-') {
     sgn2 = -1;
     bs = bs.substr(1);
int n = as.length() + bs.length() + 1;
FFT::init(n);
FOR(i, 0, FFT::pwrN) {
     a[i] = b[i] = fa[i] = fb[i] = fc[i] = cc[i] = 0;
FOR(i, 0, as.size()) {
a[i] = as[as.size() - 1 - i] - '0';
FOR(i, 0, bs.size()) {
   b[i] = bs[bs.size() - 1 - i] - '0';
\label{eq:fft} \begin{cases} \text{FFT:::fft}(a,\,fa);\\ \text{FFT:::fft}(b,\,fb);\\ \text{FOR}(i,\,0,\,\text{FFT::pwrN}) \ \{\\ \text{fc}[i] = fa[i] * fb[i]; \end{cases}
// turn [0,1,2,...,n-1] into [0, n-1, n-2, ..., 1] FOR(i, 1, FFT::pwrN) {
     if (i < FFT::pwrN - i) \{
          swap(fc[i], fc[FFT::pwrN - i]);
     }
FFT::fft(fc, cc);
ll carry = 0;
vector<int> v
FOR(i, 0, FFT::pwrN) {
     \begin{array}{l} \text{int num} = \text{round}(\text{cc[i].real()} \; / \; \text{FFT::pwrN)} \; + \; \text{carry;} \\ \text{v.pb(num} \; \% \; 10); \end{array}
     carry = num / 10;
while (carry > 0) {
v.pb(carry % 10);
     carry /= 10;
reverse(v.begin(), v.end());
bool start = false;
ostringstream ss;
bool allZero = true;
for (auto x : v) {
if (x \stackrel{!}{=} 0) {
          allZero = false;
           break;
if (sgn1*sgn2 < 0 \&\& !allZero) ss << "-";
for (auto x : v) {
    if (x == 0 && !start) continue;
     start = true;
     ss \ll abs(x);
if (!start) ss << 0;
return ss.str();
```

5.12 Formulas

```
\begin{array}{lll} \sum_{i=1}^{n}i & = & \frac{n(n+1)}{2}; & \sum_{i=1}^{n}i^2 & = & \frac{n(2n+1)(n+1)}{6}; \\ \sum_{i=1}^{n}i^3 & = & \frac{n^2(n+1)^2}{4}; & \sum_{i=1}^{n}i^4 & = & \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}; \\ \sum_{i=a}^{b}c^i & = & \frac{c^{b+1}-c^a}{c-1}, c & \neq & 1; & \sum_{i=1}^{n}a_1 + (i-1)d & = & \frac{n(a_1+a_n)}{2}; & \sum_{i=1}^{n}a_1r^{i-1} & = & \frac{a_1(1-r^n)}{1-r}, r \neq & 1; \\ \sum_{i=1}^{\infty}ar^{i-1} & = & \frac{a_1}{1-r}, |r| \leq 1. \end{array}
```

6 Strings

6.1 Hashing

```
struct HashedString { const ll A1 = 999999929, B1 = 1000000009, A2 = 1000000087, B2 = 1000000097;
```

```
vector<ll> A1pwrs, A2pwrs;
    vector<pll> prefixHash;
    HashedString(const string& _s) {
        init(s);
        calcHashes(s);
    void init(const string& s) {
        11 \ a2 = 1:
        FOR(i, 0, (int)s.length()+1) {
Alpwrs.pb(al);
            A2pwrs.pb(a2);

a1 = (a1*A1)\%B1;
            a2 = (a2*A2)\%B2;
        }
    void calcHashes(const string& s) {
        pll h = \{0, 0\};
        prefixHash.pb(h);
        for(char c : s) {
            ll h1 = (prefixHash.back().first*A1 + c)\%B1;
            ll h2 = (prefixHash.back().second*A2 + c)\%B2;
            prefixHash.pb(\{h1, h2\});
    pll getHash(int l, int r) {
        ll h<br/>1 = (prefixHash[r+1].first - prefixHash[l].first*A1pwrs
              [r+1-l]) % B1;
        ll h2 = (prefixHash[r+1].second - prefixHash[l].second* A2pwrs[r+1-l]) % B2;
        if(h1 < 0) h1 += B1;

if(h2 < 0) h2 += B2;
        return {h1, h2};
};
```

6.2 Prefix Function

```
\label{eq:continuous_problem} // \ pi[i] \ is the length of the longest proper prefix of the substring $$ s[0..i] \ which is also a suffix $$ // \ of this substring $$ vector<int> prefixFunction(const string& s) $$ int n = (int)s.length(); $$ vector<int> pi(n); $$ for (int i = 1; i < n; i++) $$ {$ int j = pi[i-1]; $$ while $(j > 0 \&\& s[i] != s[j])$$ {$ j = pi[j-1]; $$ if $(s[i] == s[j])$$ {$ j++; $$ pi[i] = j; $$ return pi; $$$ $$ $$ return pi; $$$ }
```

6.3 Prefix Function Automaton

```
// aut[oldPi][c] = newPi
vector<vector<int>> computeAutomaton(string s) {
     const char BASE = 'a';
     s += "#";
     int n = s.size();
     vector < int > pi = prefixFunction(s);
    vector<int> pi = prefixFunction(s);
vector<vector<int>> aut(n, vector<int>(26));
for (int i = 0; i < n; i++) {
    for (int c = 0; c < 26; c++) {
        if (i > 0 && BASE + c != s[i])
    }
                   aut[i][c] = aut[pi[i-1]][c];
                   \operatorname{aut}[i][c] = i + (\operatorname{BASE} + c == s[i]);
         }
    return aut;
vector<int> findOccurs(const string& s, const string& t) {
     auto\ aut = computeAutomaton(s);\\
     int curr = 0;
     vector<int> occurs;
     FOR(i, 0, (int)t.length()) {
         int\ c\,=\,t[i]\text{-'a'};
```

6.4 KMP

```
\label{eq:constant} \begin{tabular}{ll} // & Knuth-Morris-Pratt algorithm \\ vector & & sindOccurences (const string & s, const string & t) & \{ & int n = s.length(); \\ & int m = t.length(); \\ & string S = s + "\#" + t; \\ & auto pi = prefixFunction(S); \\ & vector & & sint > ans; \\ & FOR(i, n+1, n+m+1) & \{ & & if(pi[i] == n) & \{ & & \\ & & ans.pb(i-2*n); \\ & & \} & \} \\ & return ans; \\ \end{tabular}
```

6.5 Aho Corasick Automaton

```
// alphabet size
const int K = 70;
// the indices of each letter of the alphabet
int intVal[256];
\mathrm{void} \ \mathrm{init}(\dot{)} \ \{
   int curr = 2
    intVal[1] = 1;
    for(char c = '0'; c <= '9'; c++, curr++) intVal[(int)c] =
          curr;
   for(char\ c='A';\ c<='Z';\ c++,\ curr++)\ intVal[(int)c]=
          curr;
    for(char c = 'a'; c \le 'z'; c++, curr++) intVal[(int)c] =
          curr:
struct Vertex {
   int next[K]:
    vector<int> marks;
        this can be changed to int mark = -1, if there will be
          no duplicates
    int p = -1;
    char pch;
   int link = -1;
   int exitLink = -1:
         exitLink points to the next node on the path of suffix
          links which is marked
   int go[K];
      / ch has to be some small char
    Vertex(int \_p=-1, char ch=(char)1) : p(\_p), pch(ch) 
        fill(begin(next), end(next), -1);
        fill(begin(go), end(go), -1);
};
vector < Vertex > t(1);
void addString(string const& s, int id) {
    int v = 0
    for (char ch : s) {
        int c = intVal[(int)ch];

if (t[v].next[c] == -1) {

t[v].next[c] = t.size();
            t.emplace_back(v, ch);
        v = t[v].next[c];
    t[v].marks.pb(id);
}
int go(int v, char ch);
```

```
\begin{array}{l} \mathrm{int}\ \mathrm{getLink}(\mathrm{int}\ v)\ \{\\ \mathrm{if}\ (\mathrm{t[v].link}==-1)\ \{\\ \mathrm{if}\ (v==0\ ||\ t[v].p==0) \end{array}
            t[v].link = 0;
        else
            t[v].link = go(getLink(t[v].p), t[v].pch);
    return t[v].link;
int\ getExitLink(int\ v)\ \{
    if(t[v].exitLink != -1) return t[v].exitLink;
    int l = getLink(v);
    if(l == 0) return t[v].exitLink = 0;
    if(!t[l].marks.empty())\ return\ t[v].exitLink = l;\\
    return\ t[v].exitLink = getExitLink(l);
int go<br/>(int v, char ch) {
    int c = intVal[(int)ch];
    if (t[v].go[c] == -1) {
if (t[v].next[c] != -1)
            t[v].go[c] = t[v].next[c];
        else
            t[v].go[c] = v == 0 ? 0 : go(getLink(v), ch);
    return t[v].go[c];
void walkUp(int v, vector<int>& matches) {
    if(v == 0) return;
    if(!t[v].marks.empty())
        for(auto m : t[v].marks) matches.pb(m);
    walkUp(getExitLink(v), matches);
}
// returns the IDs of matched strings.
// Will contain duplicates if multiple matches of the same string
       are found.
vector{<}int{>}\ walk(const\ string\&\ s)\ \{
    vector<int> matches:
    int curr = 0;
    for(char c : s) {
        curr = go(curr, c);
        if(!t[curr].marks.empty()) {
            for(auto m : t[curr].marks) matches.pb(m);
        walkUp(getExitLink(curr), matches);
    return matches;
/* Usage:
 * addString(strs[i], i);
   auto matches = walk(text);
 * .. do what you need with the matches - count, check if some
       {\rm id}\ {\rm exists},\ {\rm etc}\ \dots
 * Some applications:
   - Find all matches: just use the walk function
 * - Find lexicographically smallest string of a given length that
        doesn't match any of the given strings:
 * For each node, check if it produces any matches (it either
       contains some marks or walkUp(v) returns some marks)
 * Remove all nodes which produce at least one match. Do DFS
        in the remaining graph, since none of the remaining
       nodes
  will ever produce a match and so they're safe.
 * - Find shortest string containing all given strings:
 * For each vertex store a mask that denotes the strings which
       match at this state. Start at (v = root, mask = 0),
   we need to reach a state (v, mask=2^n-1), where n is the
       number of strings in the set. Use BFS to transition
       between states
   and update the mask.
```

6.6 Suffix Array

```
vector<int> sortCyclicShifts(string const& s) {
  int n = s.size();
  const int alphabet = 256; // we assume to use the whole
   ASCII range
  vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
  for (int i = 0; i < n; i++)</pre>
```