

Networks (2IRR20)

Wireless and Mobile (07)

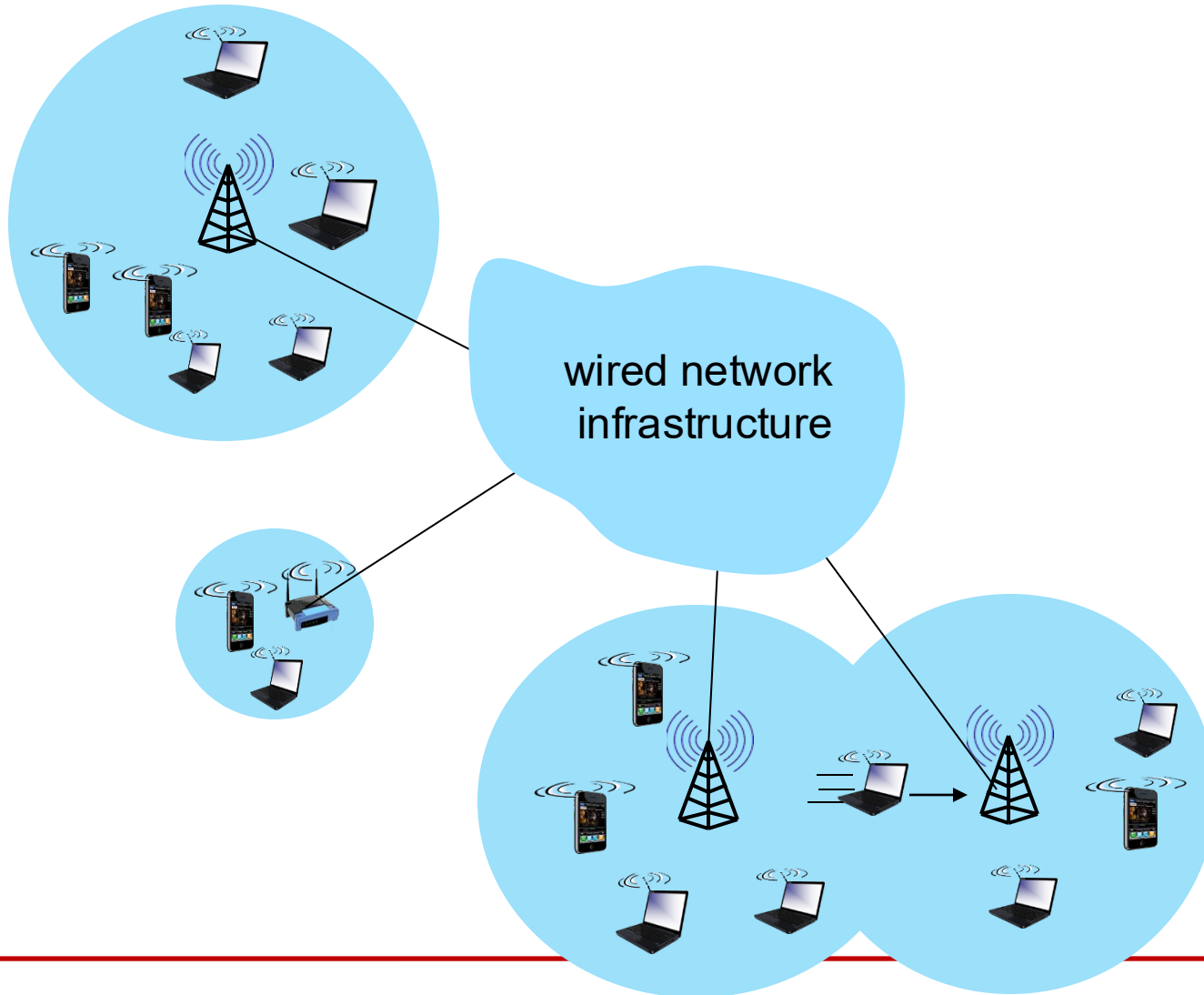
Dr. Tanir Ozcelebi

This slide set

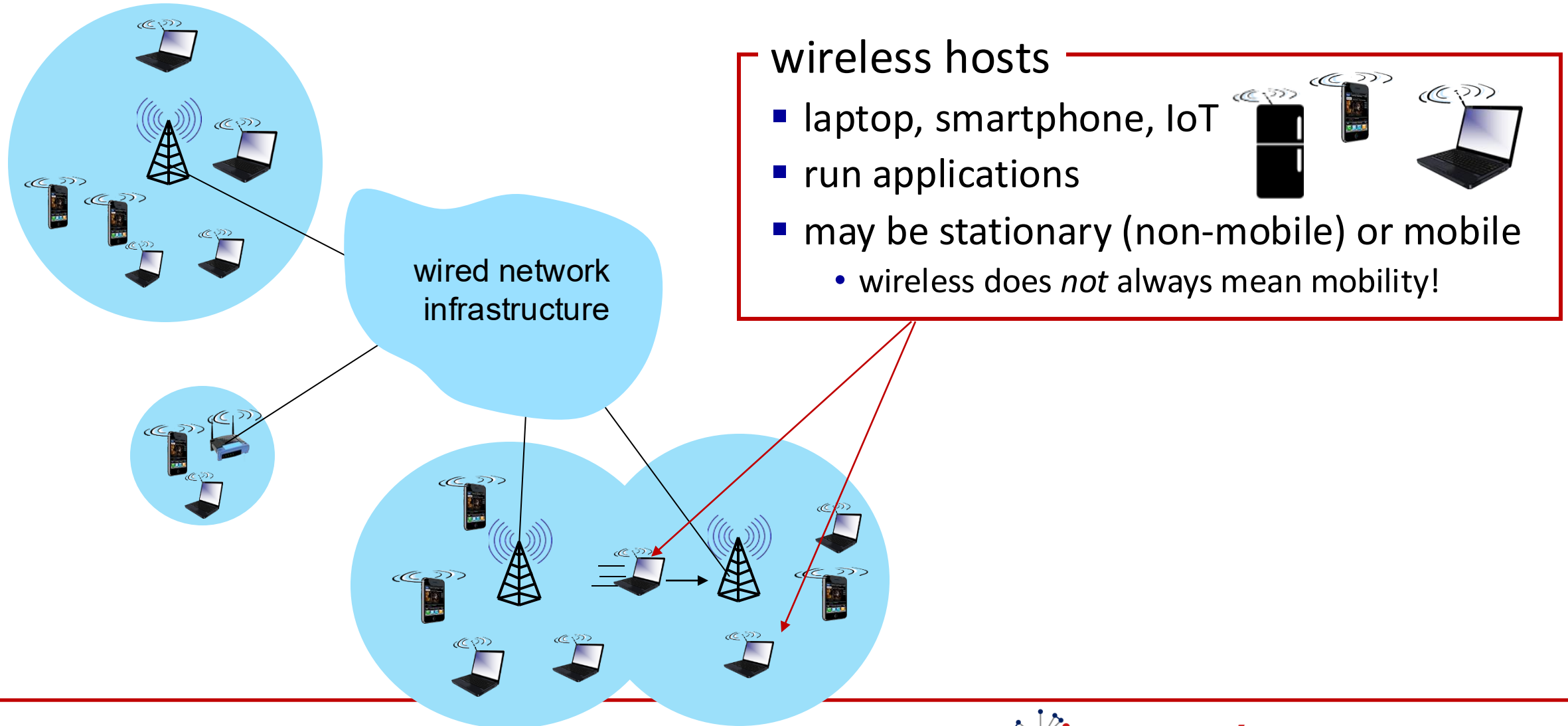
- Wireless
 - Wireless Links and network characteristics
 - WiFi: 802.11 wireless LANs
- Mobile
 - Mobility management
 - Mobile IP
 - Mobility: impact on higher-layer protocols

Introduction

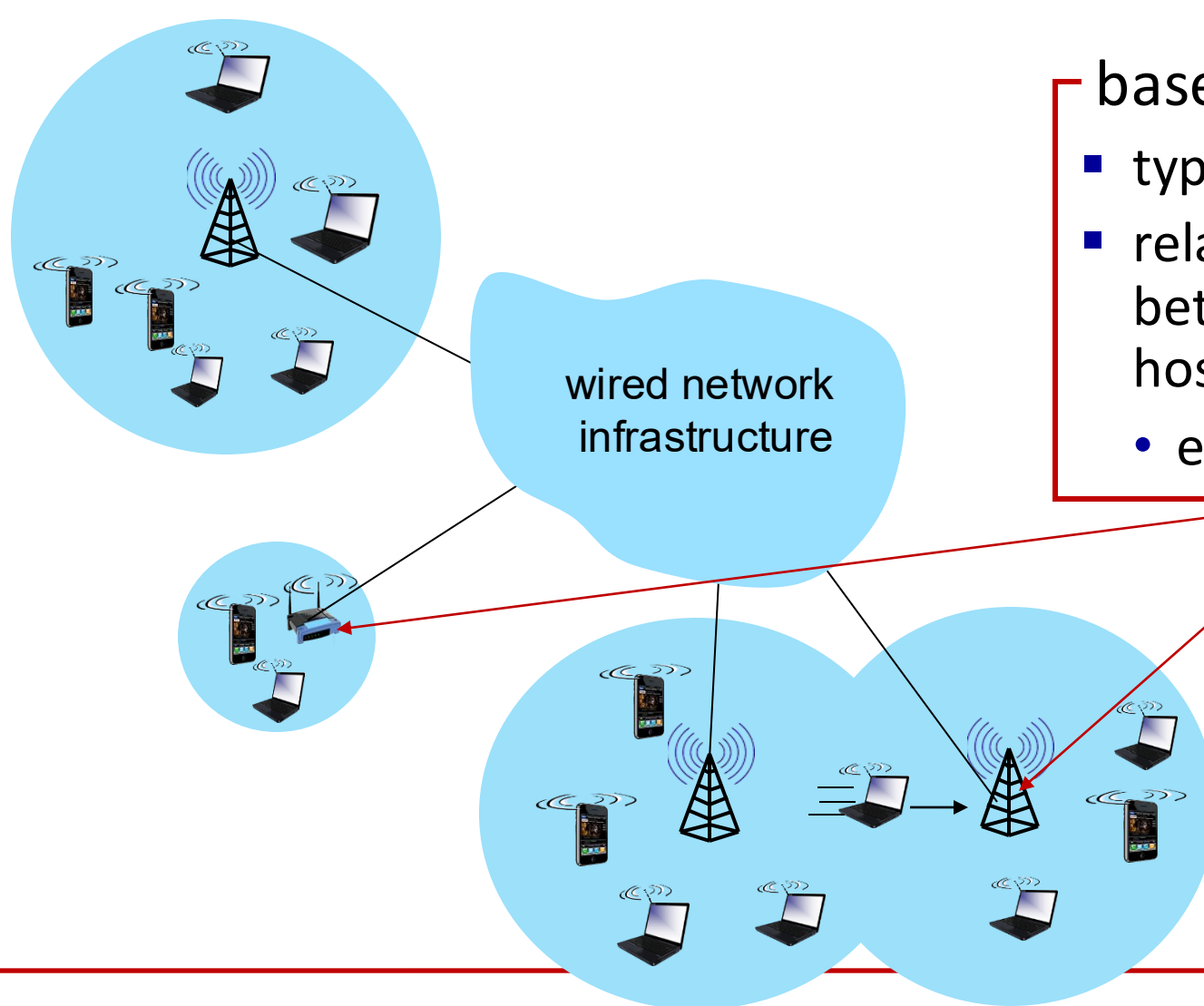
Elements of a wireless network



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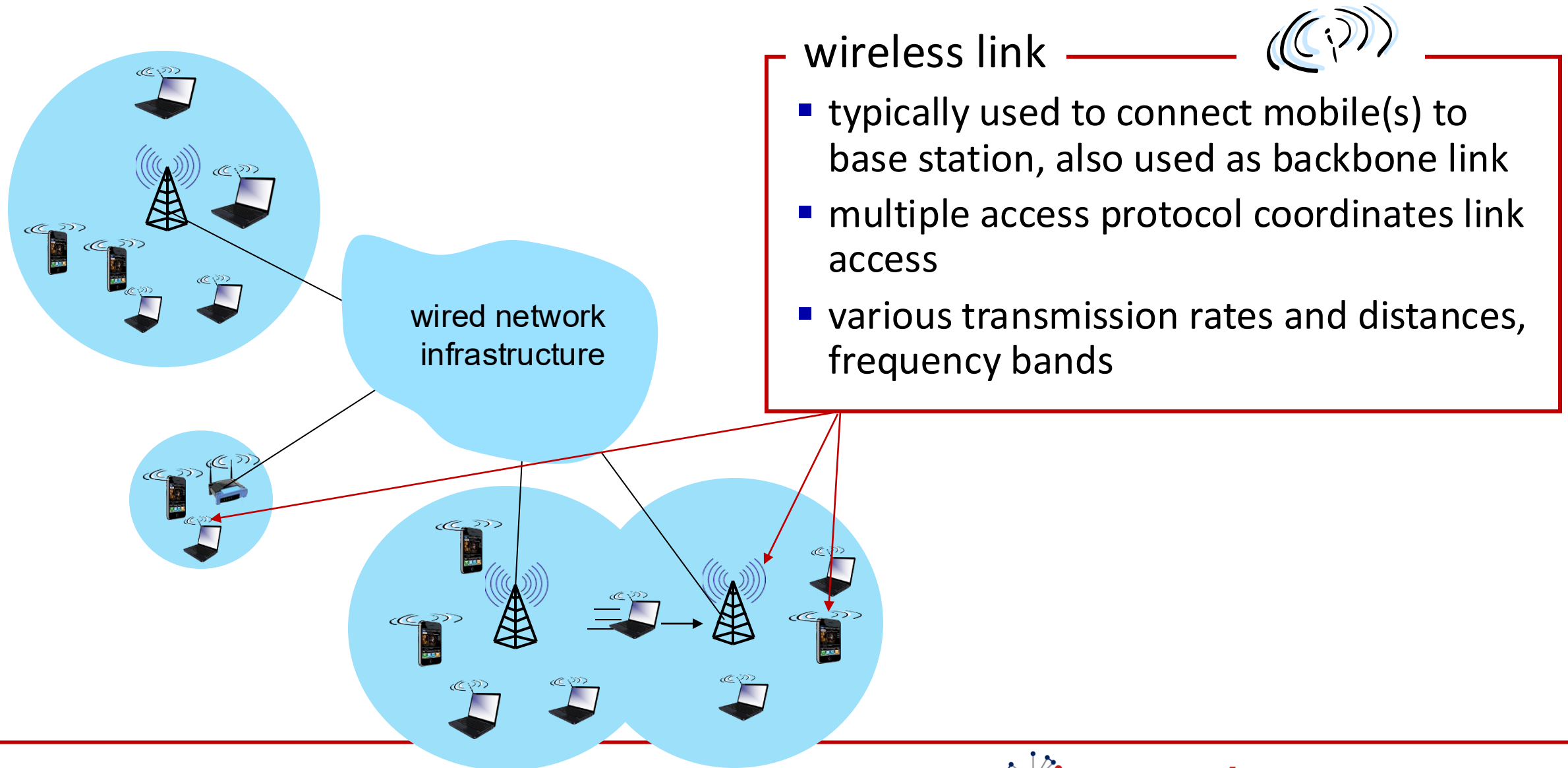


base station

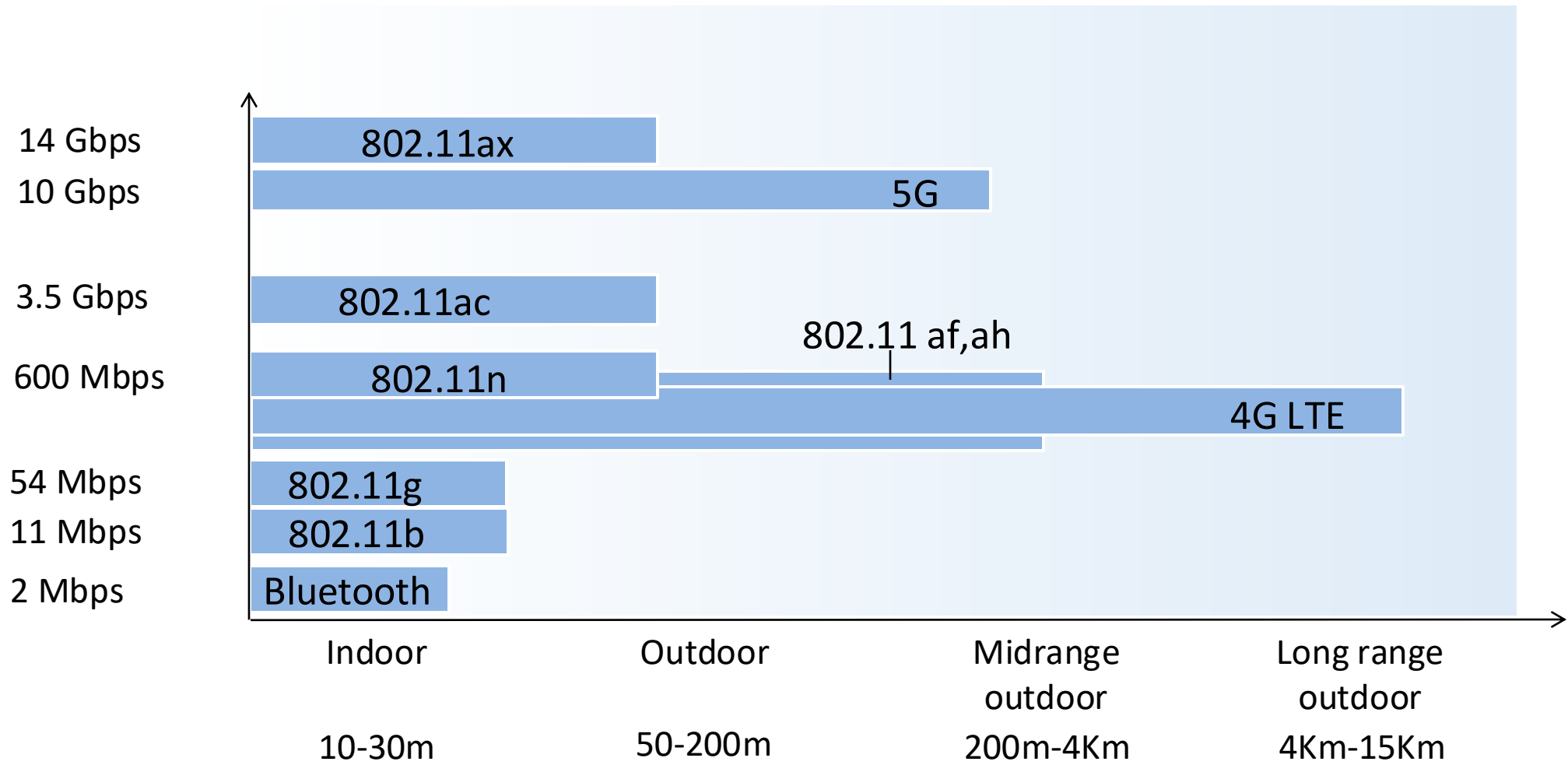


- typically connected to wired network
- relay - responsible for sending packets between wired network and wireless host(s) in its "area"
 - e.g., cell towers, 802.11 access points

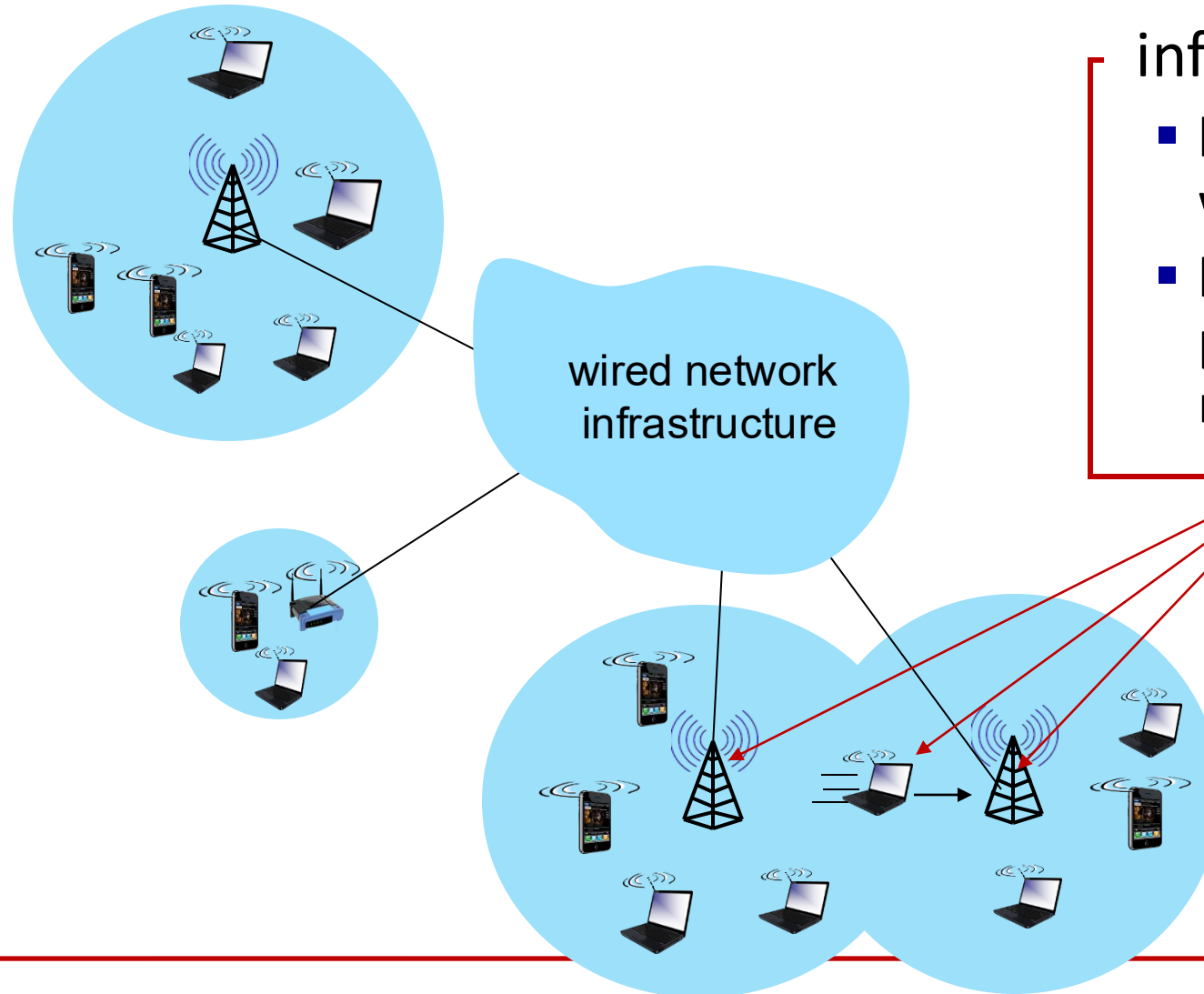
Elements of a wireless network



Characteristics of selected wireless links



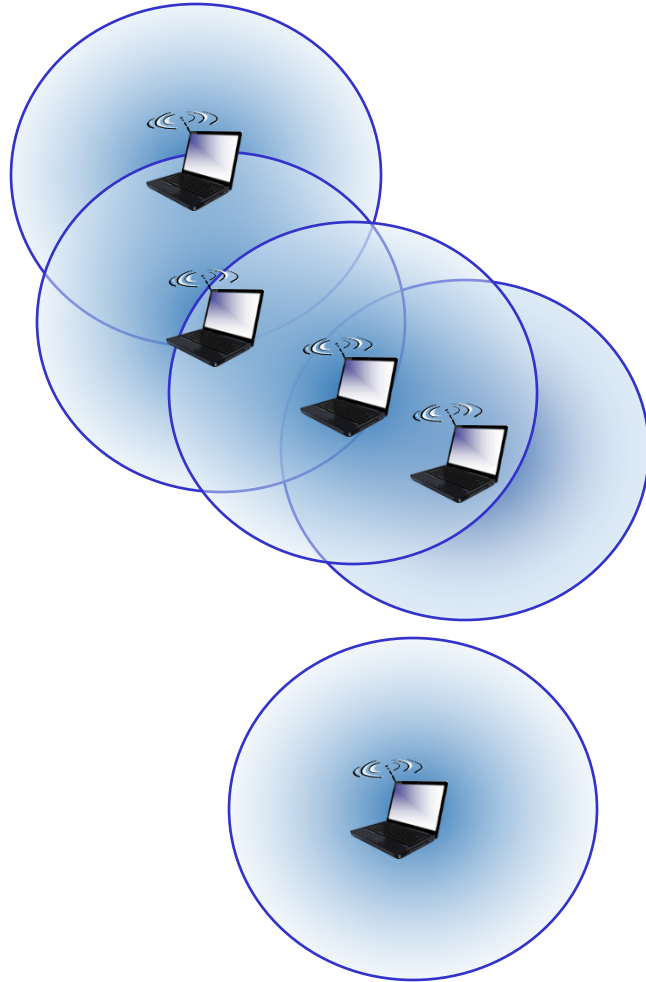
Elements of a wireless network



infrastructure mode

- base station connects mobiles into wired network
- handoff: mobile changes base station providing connection into wired network

Elements of a wireless network



ad hoc mode

- no base stations
- nodes can only transmit to other nodes within link coverage
- nodes organize themselves into a network: route among themselves

Wireless network taxonomy

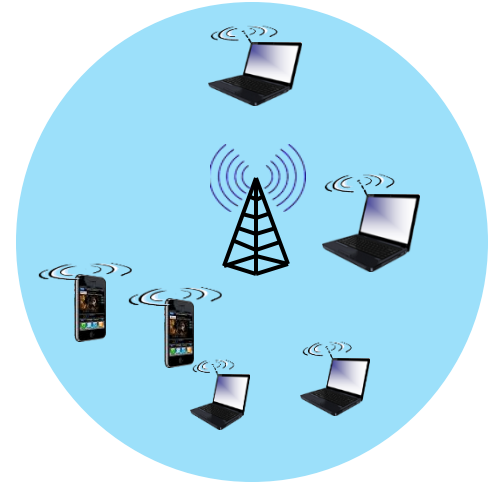
| | single hop | multiple hops |
|-------------------------------|--|--|
| infrastructure (e.g., APs) | host connects to base station (WiFi, cellular) which connects to larger Internet | host may have to relay through several wireless nodes to connect to larger Internet: <i>mesh net</i> |
| <i>no infrastructure</i> | no base station, no connection to larger Internet (Bluetooth, ad hoc nets) | no base station, no connection to larger Internet. May have to relay to reach other a given wireless node MANET, VANET |

Wireless links and network characteristics

Wireless link characteristics (1)

important differences from wired link

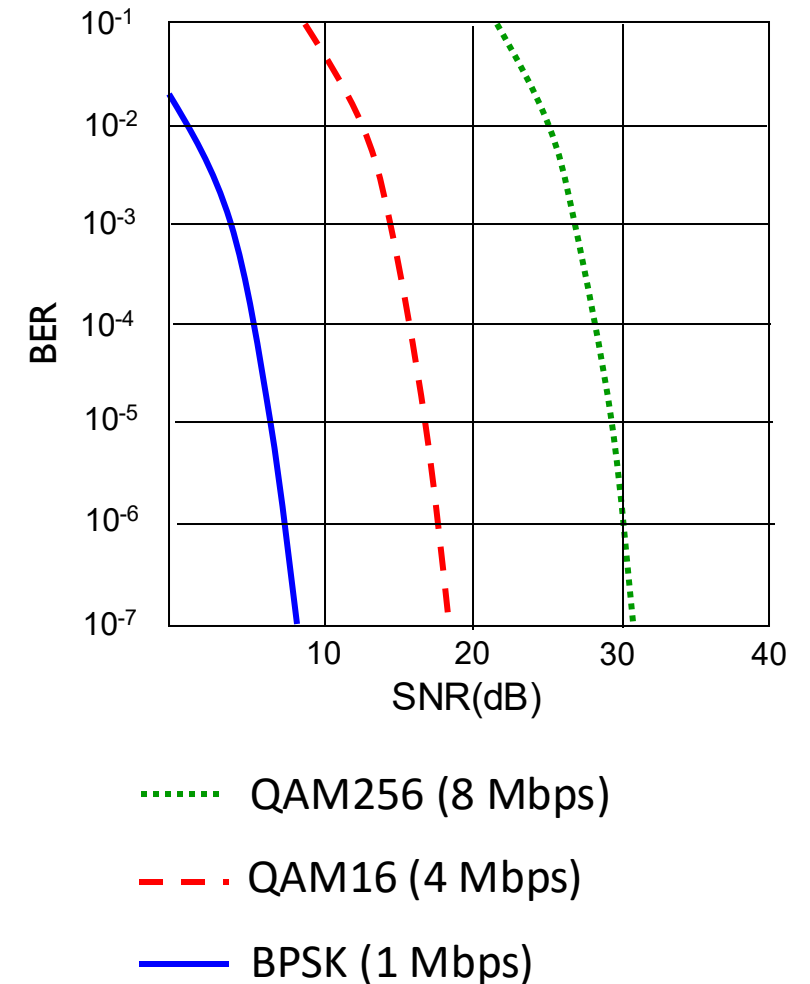
- **decreased signal strength:** radio signal attenuates as it propagates through matter (path loss)
- **interference from other sources:** wireless network frequencies (e.g., 2.4 GHz) shared by many devices (e.g., WiFi, cellular, motors): interference
- **multipath propagation:** radio signal reflects off objects ground, arriving at destination at slightly different times



.... make communication across (even a point to point) wireless link much more “difficult”

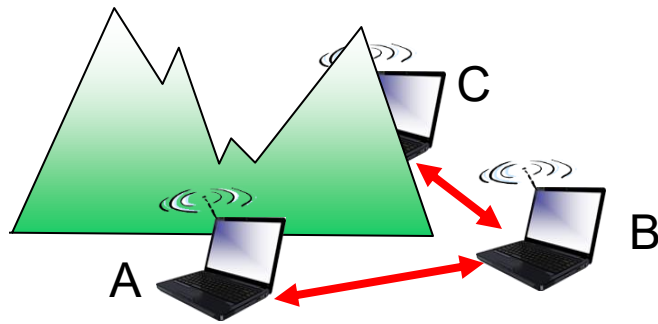
Wireless link characteristics (2)

- SNR: signal-to-noise ratio
 - larger SNR – easier to extract signal from noise (a “good thing”)
- SNR versus BER tradeoffs
 - *given physical layer*: increase power -> increase SNR->decrease BER
 - *given SNR*: choose physical layer that meets BER requirement, giving highest throughput
 - SNR may change with mobility: dynamically adapt physical layer (modulation technique, rate)



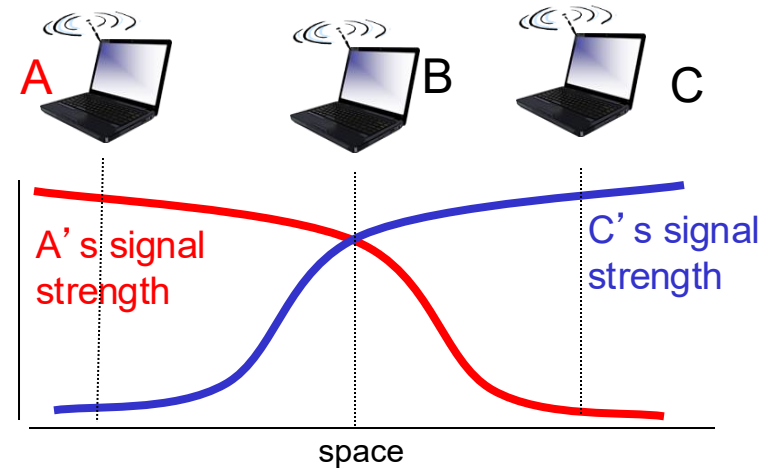
Wireless link characteristics (3)

Multiple wireless senders, receivers create additional problems (beyond multiple access):



Hidden terminal problem

- B, A hear each other
- B, C hear each other
- A, C can not hear each other means A, C unaware of their interference at B



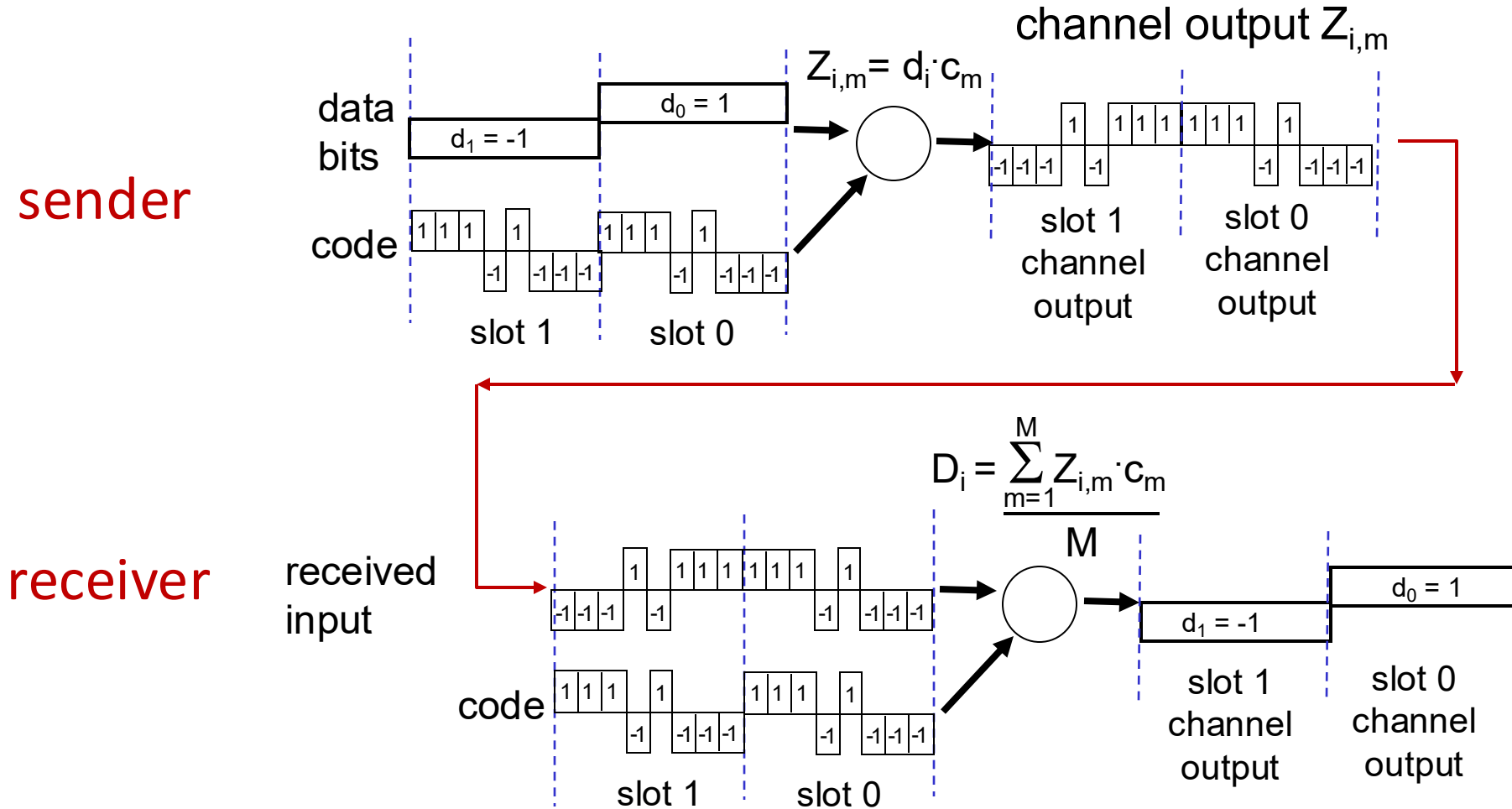
Signal attenuation:

- B, A hear each other
- B, C hear each other
- A, C can not hear each other interfering at B

Code Division Multiple Access (CDMA)

- unique “code” assigned to each user; i.e., code set partitioning
 - all users share same frequency, but each user has own “chipping” sequence (i.e., code) to encode data
 - allows multiple users to “coexist” and transmit simultaneously with minimal interference (if codes are “orthogonal”)
- **encoding:** inner product: (original data) \times (chipping sequence)
- **decoding:** summed inner-product: (encoded data) \times (chipping sequence)

CDMA encode/decode

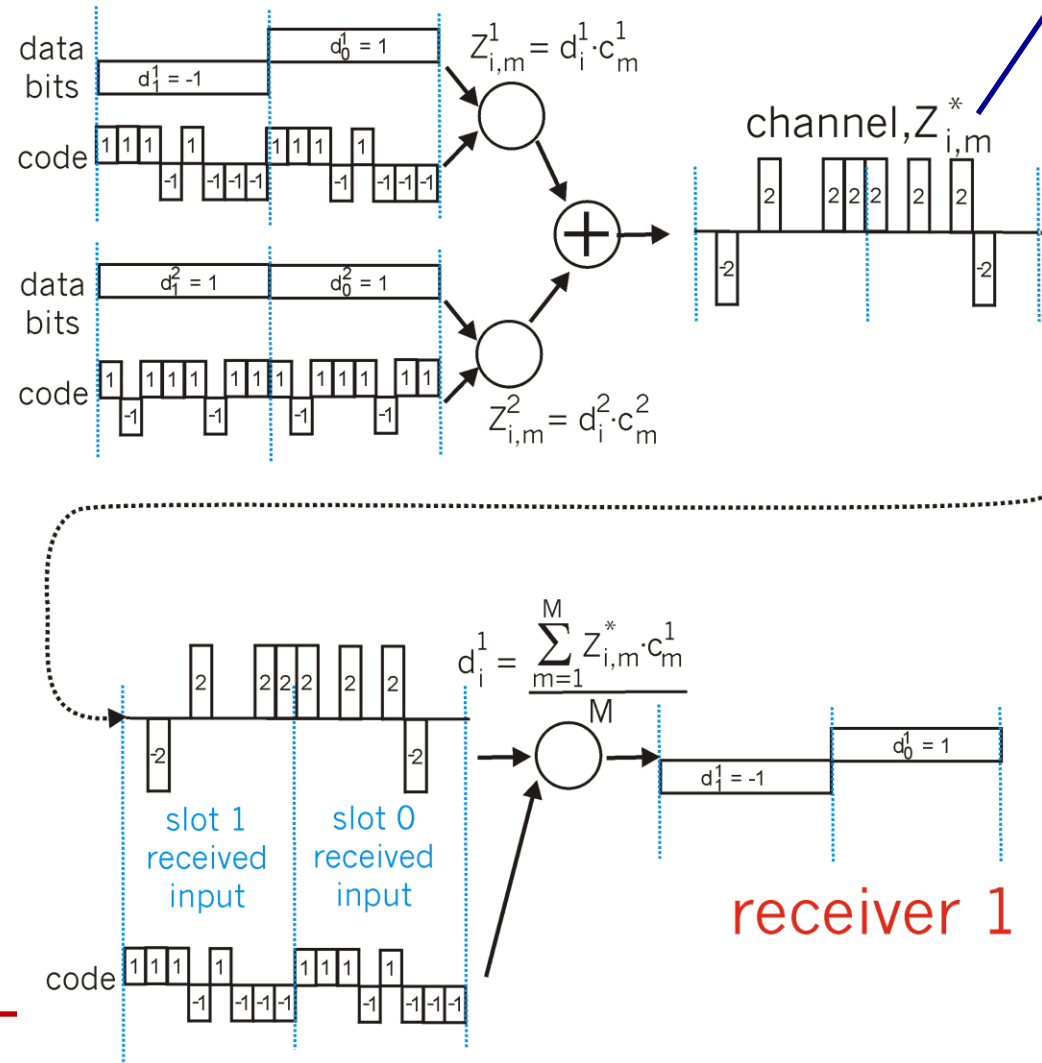


... but this isn't really useful yet!

CDMA: two-sender interference

Sender 1

Sender 2



channel sums together transmissions by sender 1 and 2

using same code as sender 1, receiver recovers sender 1's original data from summed channel data!

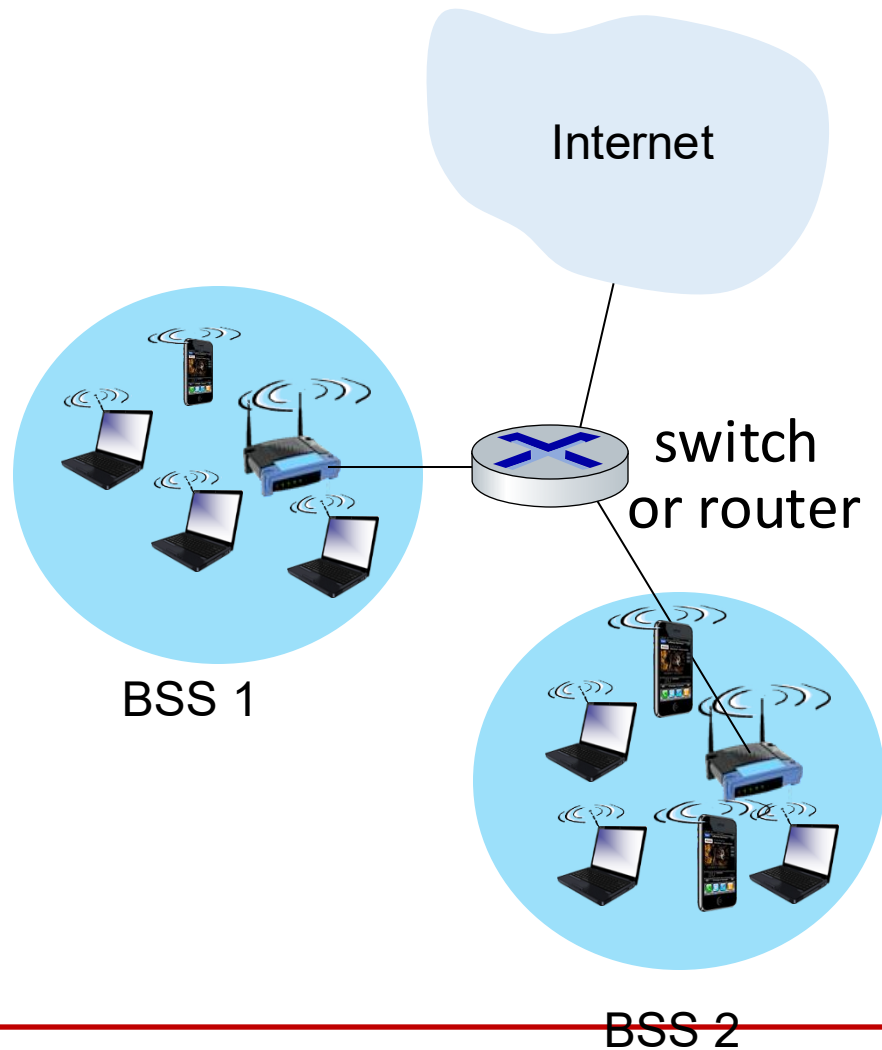
... now *that's* useful!

WiFi: 802.11 wireless LANs

IEEE 802.11 Wireless LAN

| IEEE 802.11 standard | Year | Max data rate | Range | Frequency |
|----------------------|------------|---------------|-------|------------------------------|
| 802.11b | 1999 | 11 Mbps | 30 m | 2.4 Ghz |
| 802.11g | 2003 | 54 Mbps | 30m | 2.4 Ghz |
| 802.11n (WiFi 4) | 2009 | 600 | 70m | 2.4, 5 Ghz |
| 802.11ac (WiFi 5) | 2013 | 3.47Gpbs | 70m | 5 Ghz |
| 802.11af | 2014 | 35 – 560 Mbps | 1 Km | unused TV bands (54-790 MHz) |
| 802.11ah | 2017 | 347Mbps | 1 Km | 900 Mhz |
| 802.11ax (WiFi 6) | 2020 | 14 Gbps | 70m | 2.4, 5 Ghz |
| 802.11be (Wifi 7) | 2024 (exp) | 23 Gbps | 70m | 2.4, 5 , 6 Ghz |
| 802.11bn (Wifi 8) | 2028 (exp) | 100 Gbps | 70m | 2.4, 5, 6 Ghz |

802.11 LAN architecture



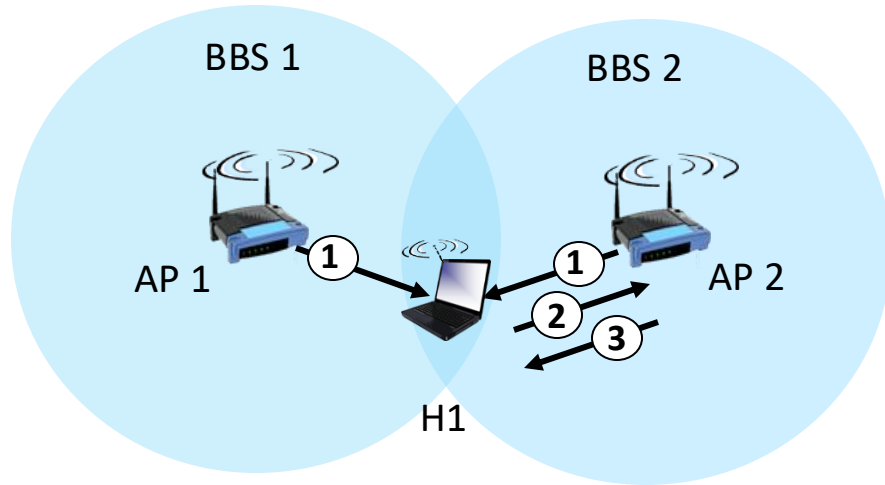
- wireless host communicates with base station
 - base station = access point (AP)
- Basic Service Set (BSS) (aka “cell”) in infrastructure mode contains:
 - wireless hosts
 - access point (AP): base station
 - ad hoc mode: hosts only

802.11: Channels, association

- spectrum divided into channels at different frequencies
 - AP admin chooses frequency for AP
 - interference possible: channel can be same as that chosen by neighboring AP!
- arriving host: must **associate** with an AP
 - scans channels, listening for *beacon frames* containing AP's name (SSID) and MAC address
 - selects AP to associate with
 - then may perform authentication [Chapter 8]
 - then typically run DHCP to get IP address in AP's subnet

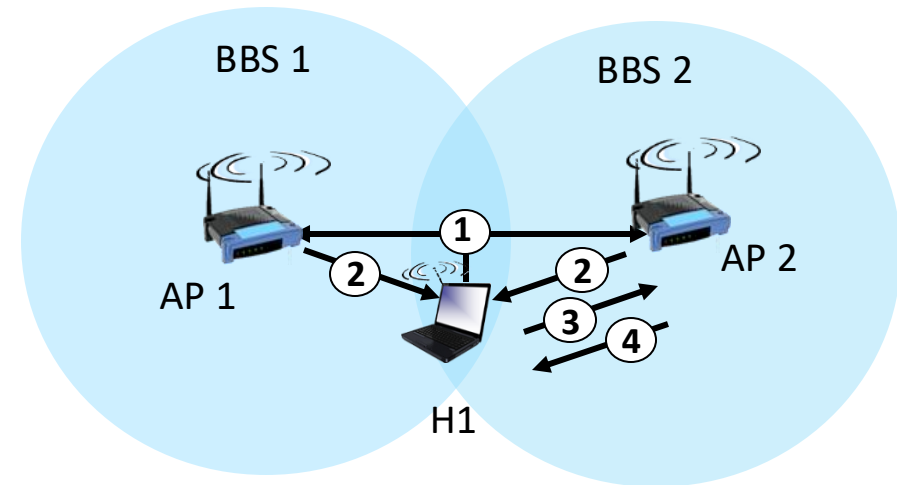


802.11: passive/active scanning



passive scanning:

- (1) beacon frames sent from APs
- (2) association Request frame sent: H1 to selected AP
- (3) association Response frame sent from selected AP to H1

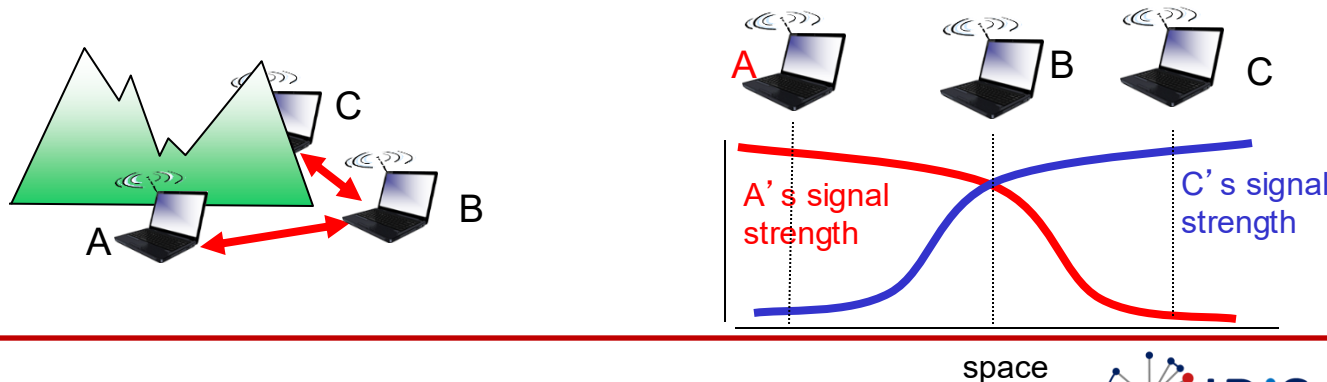


active scanning:

- (1) Probe Request frame broadcast from H1
- (2) Probe Response frames sent from APs
- (3) Association Request frame sent: H1 to selected AP
- (4) Association Response frame sent from selected AP to H1

IEEE 802.11: multiple access

- avoid collisions: 2⁺ nodes transmitting at same time
- 802.11: CSMA - sense before transmitting
 - don't collide with detected ongoing transmission by another node
- 802.11: *no* collision detection!
 - difficult to sense collisions: high transmitting signal, weak received signal due to fading
 - can't sense all collisions in any case: hidden terminal, fading
 - goal: *avoid collisions*: CSMA/CollisionAvoidance



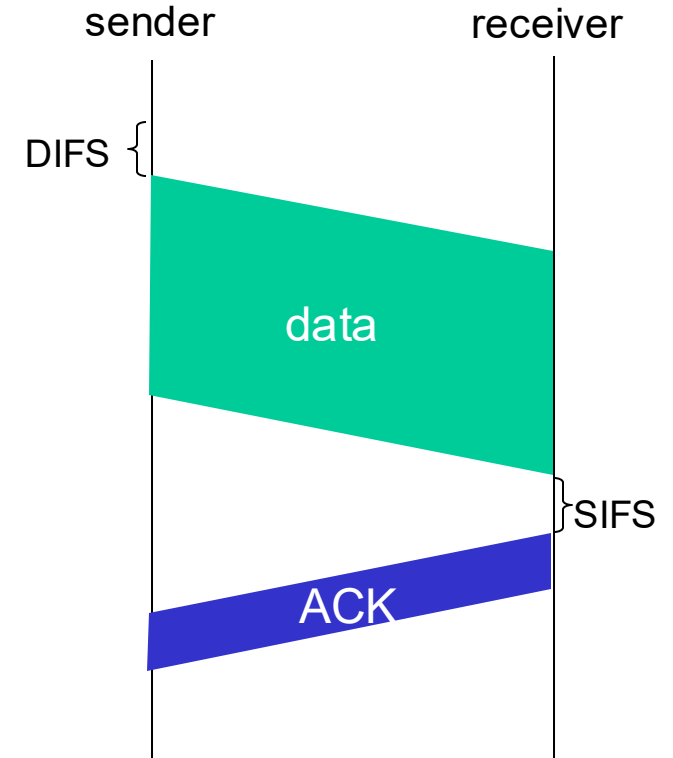
IEEE 802.11 MAC Protocol: CSMA/CA

802.11 sender

- 1 if sense channel idle for **DIFS** then
transmit entire frame (no CD)
- 2 if sense channel busy then
start random backoff time
timer counts down while channel idle
transmit when timer expires
if no ACK, increase random backoff interval, repeat 2

802.11 receiver

- if frame received OK
return ACK after **SIFS** (ACK needed due to hidden terminal problem)

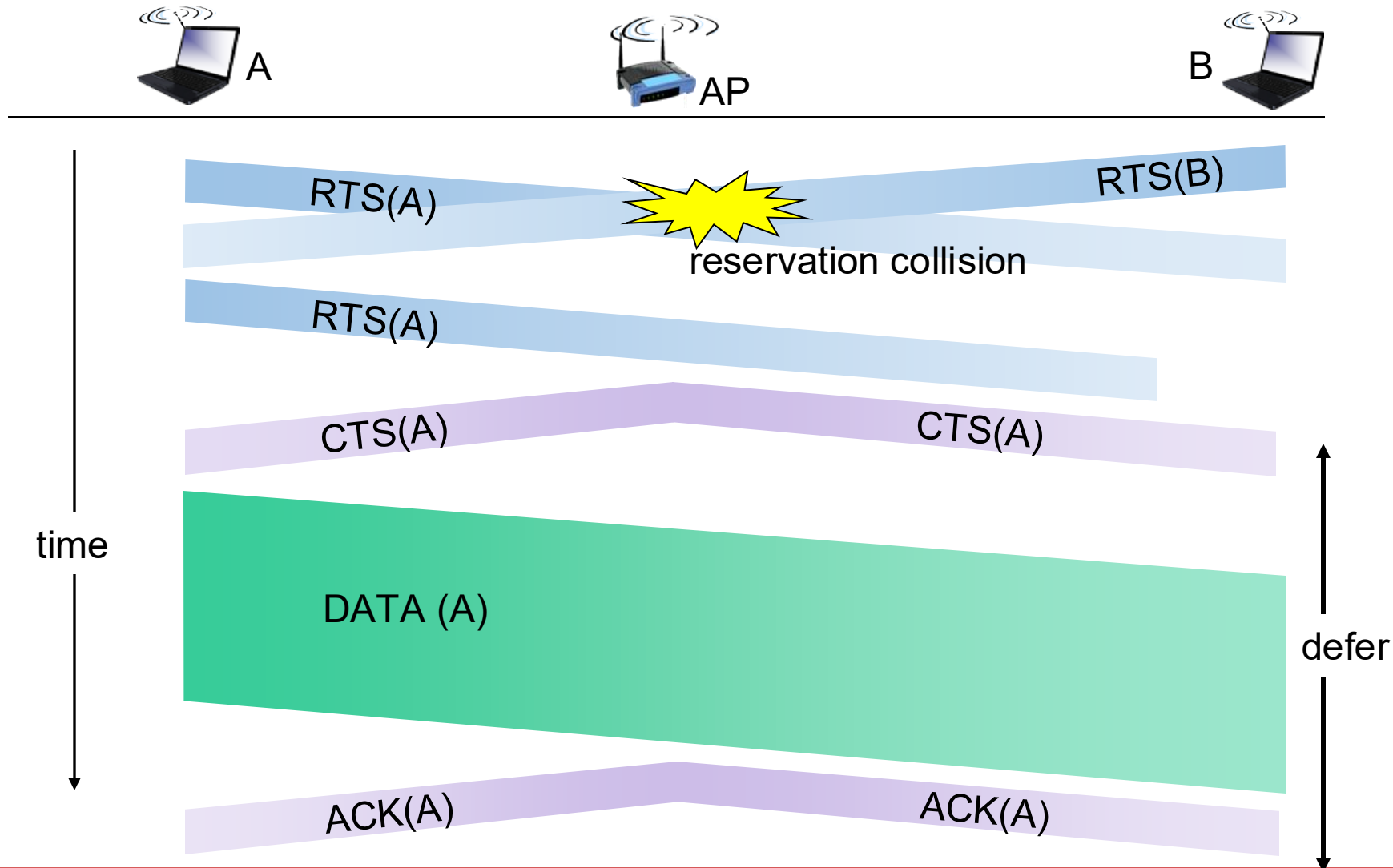


Avoiding collisions (more)

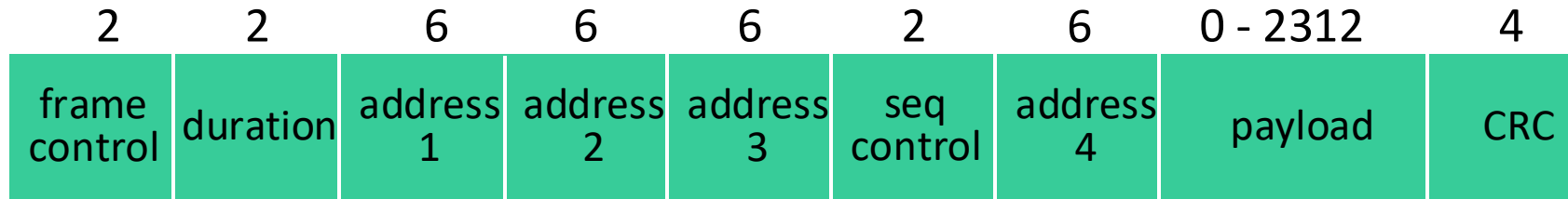
idea: sender “reserves” channel use for data frames using small reservation packets

- sender first transmits *small* request-to-send (RTS) packet to BS using CSMA
 - RTSs may still collide with each other (but they’re short)
- BS broadcasts clear-to-send CTS in response to RTS
- CTS heard by all nodes
 - sender transmits data frame
 - other stations defer transmissions

Collision Avoidance: RTS-CTS exchange



802.11 frame: addressing



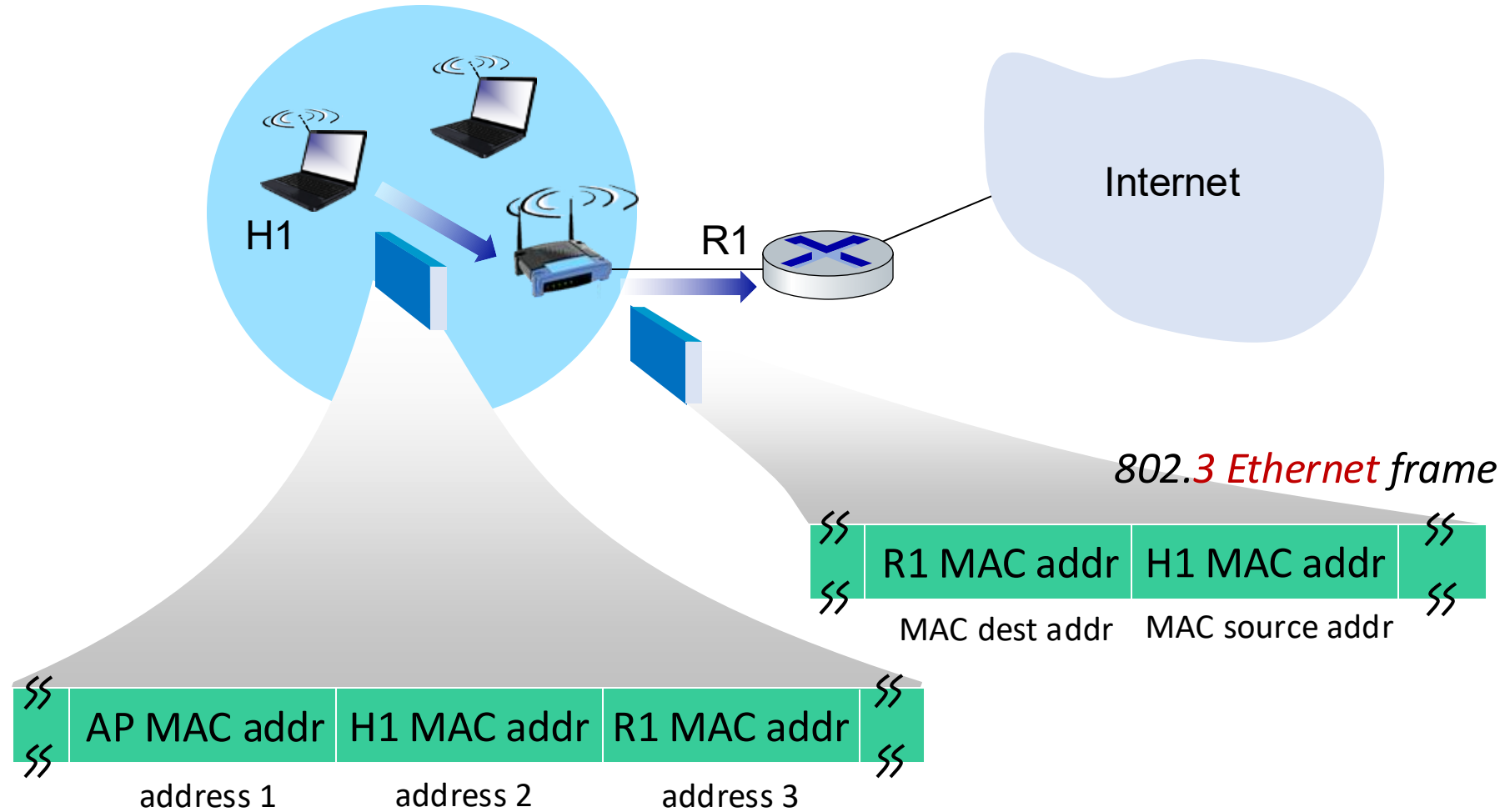
Address 1: MAC address of wireless host or AP to receive this frame

Address 2: MAC address of wireless host or AP transmitting this frame

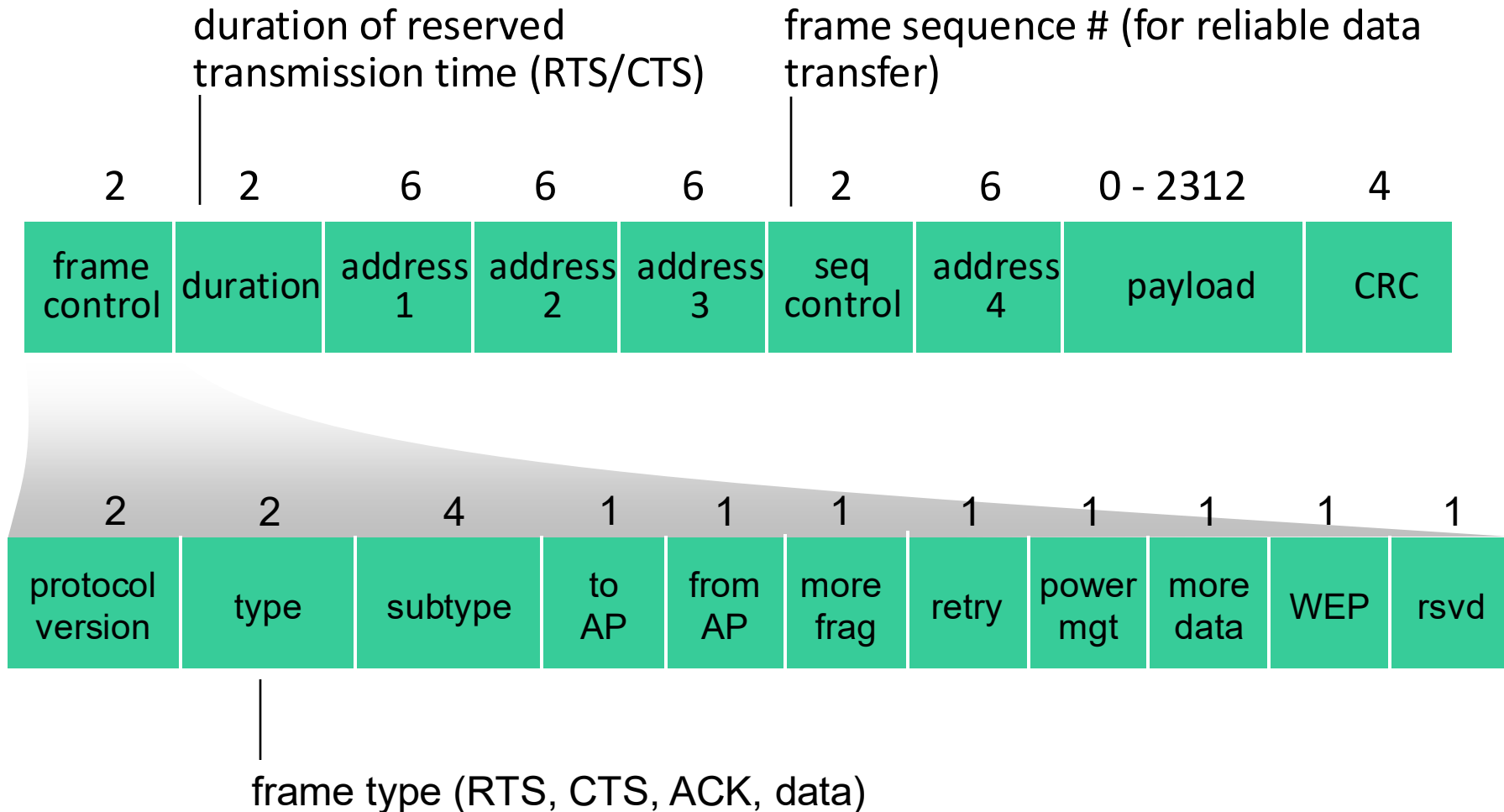
Address 3: MAC address of router interface to which AP is attached

Address 4: used only in ad hoc mode

802.11 frame: addressing

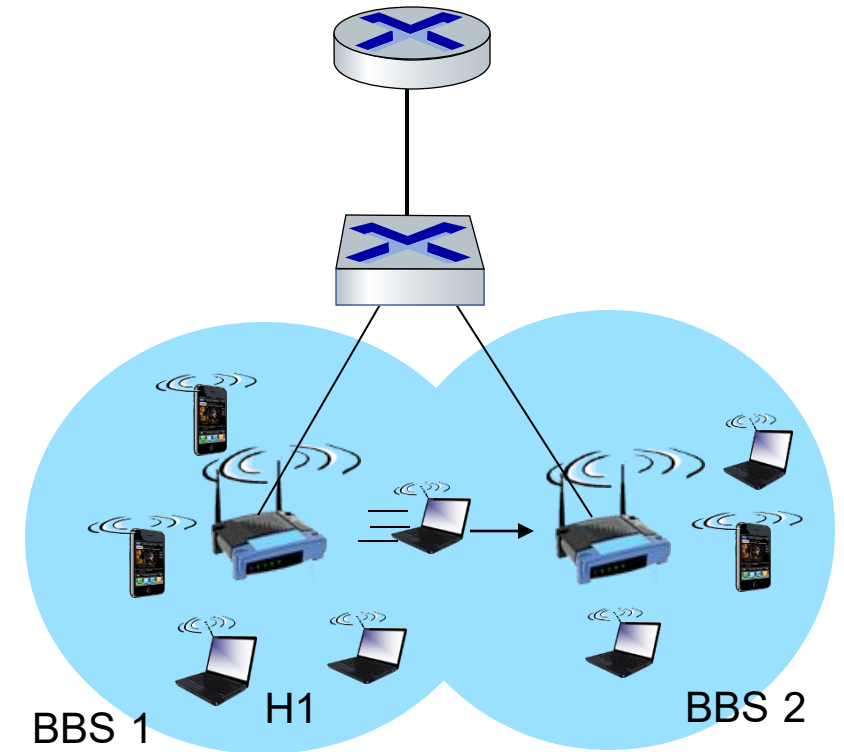


802.11 frame: addressing



802.11: mobility within same subnet

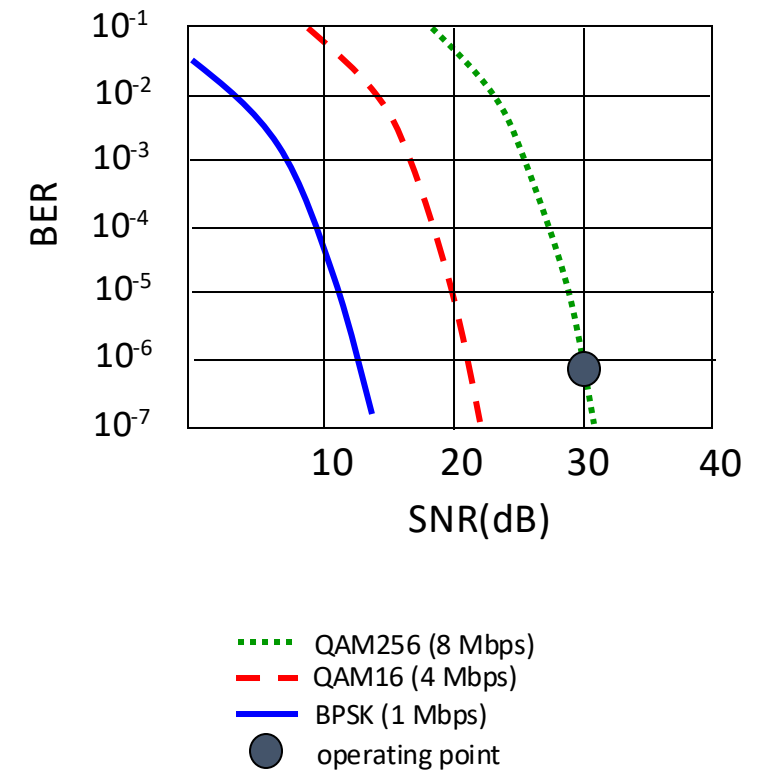
- H1 remains in same IP subnet: IP address can remain same
- switch: which AP is associated with H1?
 - Low mobility: Self-learning will do the job
→ switch will see frame from H1 and “learn again” which switch port can be used to reach H1.
 - High mobility: We need more
→ New AP could broadcast a frame with source address MAC_{H1} .



802.11: advanced capabilities

Rate adaptation

- base station, mobile dynamically change transmission rate (physical layer modulation technique) as mobile moves, SNR varies
 1. SNR decreases, BER increase as node moves away from base station
 2. When BER becomes too high, switch to lower transmission rate to reduce BER



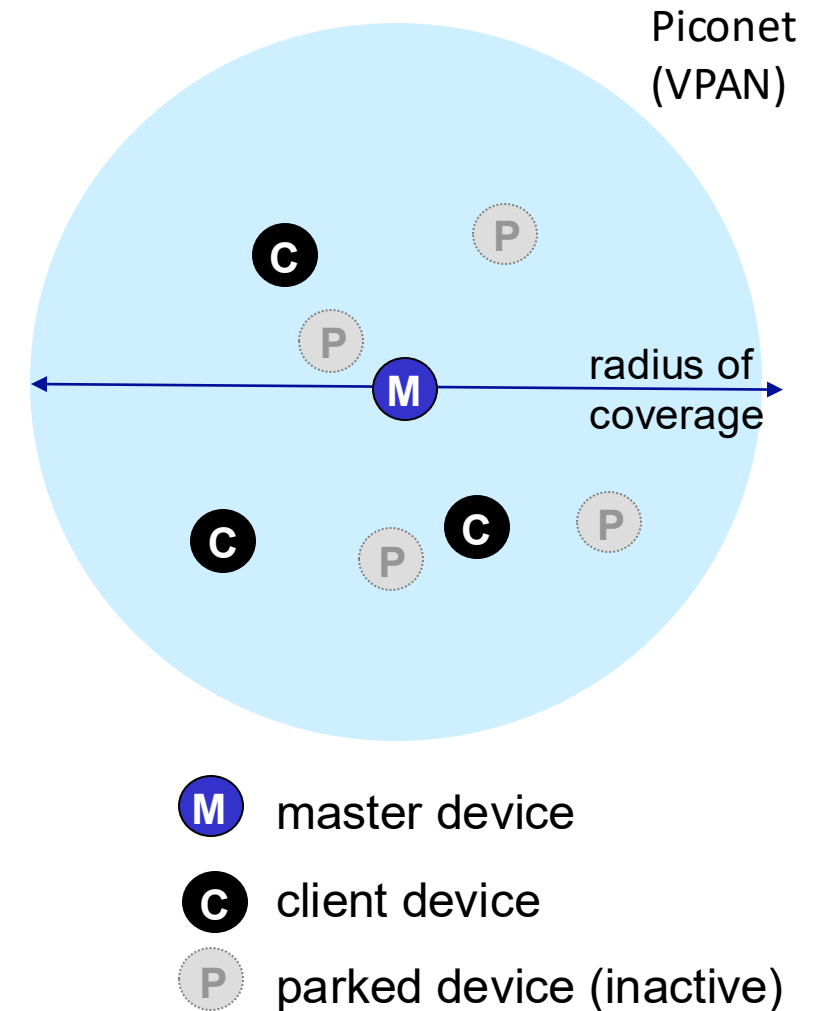
802.11: advanced capabilities

power management

- node-to-AP: “I am going to sleep until next beacon frame”
 - AP knows not to transmit frames to this node
 - node wakes up before next **beacon frame** (~every 100ms)
- beacon frame: contains list of mobiles with AP-to-mobile frames waiting to be sent
 - node will stay awake if AP-to-mobile frames are to be sent; otherwise sleep again until next beacon frame (asleep 99% of the time!)

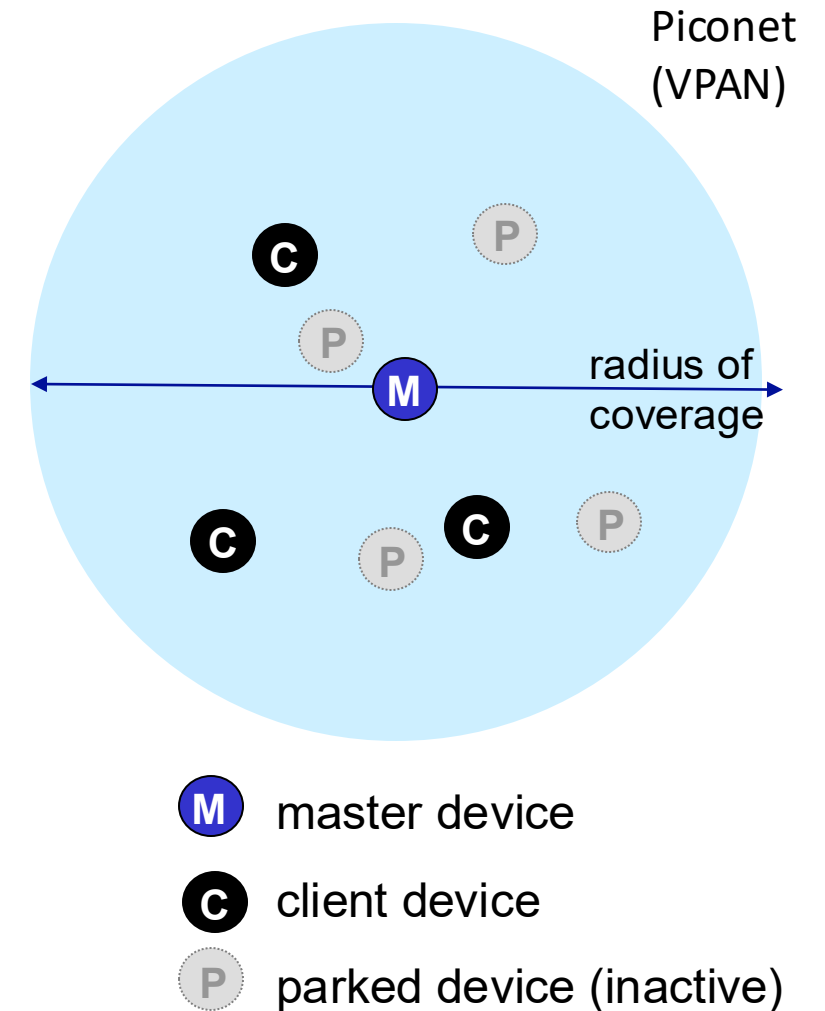
Personal area networks: Bluetooth

- less than 10 m diameter
 - max 8 active nodes
plus some “parked” (sleeping) ones
- ad hoc: no infrastructure
- 2.4-2.5 GHz ISM radio band, up to 3 Mbps
- master controller / clients devices:
 - master polls clients, grants requests for client transmissions
 - master determines which nodes are active, their clocks, the frequency sequence to use etc etc.



Personal area networks: Bluetooth

- TDM, 625 μ sec sec. slot
- FDM: sender uses 79 frequency channels in known, pseudo-random order slot-to-slot (spread spectrum)
 - other devices/equipment not in piconet (VPAN) only interfere in some slots
- **parked mode:** clients can “go to sleep” (park) and later wakeup (to preserve battery) → master determines



Mobility management

What is mobility?

- spectrum of mobility, from the **network** perspective:

no mobility

high mobility



device moves
between
networks, but
powers down
while moving

device moves
within same AP in
one provider
network

device moves
among APs in
one provider
network

device moves
among multiple
provider networks,
while maintaining
ongoing
connections

We're interested in these!

Mobility approaches

- **let network (routers) handle it:**
 - routers advertise well-known name, address (e.g., permanent 32-bit IP address), or number (e.g., cell #) of visiting mobile node via usual routing table exchange.
 - Internet routing could do this already *with no* changes! Routing tables can indicate where each mobile is located via longest prefix match!

Mobility approaches

- let network (routers) handle it:
 - routers advertise well-known address (e.g., permanent 32-bit IP address), or number (e.g., cell ID) of mobile node via usual routing table exchange
 - Internet routing could do this *no changes!* Routing tables can indicate where each mobile is located via longest prefix match!
- **let end-systems handle it:** functionality at the “edge”
 - *indirect routing*: communication from correspondent to mobile goes through mobile’s home network, then forwarded to remote mobile
 - *direct routing*: correspondent gets foreign address of mobile, sends directly to mobile

not
scalable
to billions of
mobiles

Contacting a mobile friend:

Consider friend frequently changing locations, how do you find him/her?

- search all phone books?
- expect her to let you know where he/she is?
- call his/her parents?
- social media!

The importance of having a “home”:

- a definitive source of information about you
- a place where people can find out where you are



Home network, visited network: general IP

[permanent “home” of mobile]

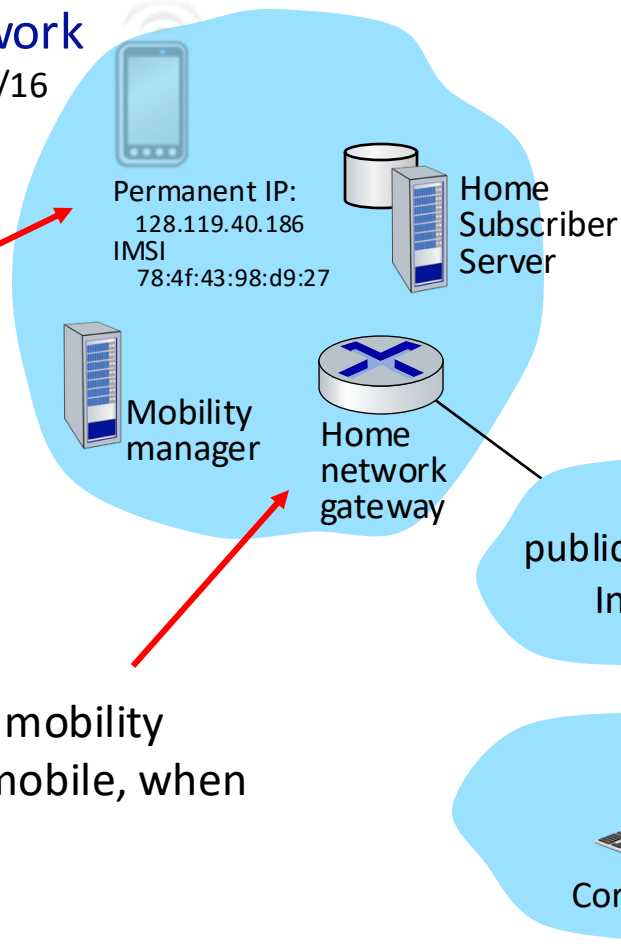
Home Network

e.g., 128.119/16

Permanent address:

address in home network,
can always be used to
reach mobile.

entity that will perform mobility
functions on behalf of mobile, when
mobile is remote



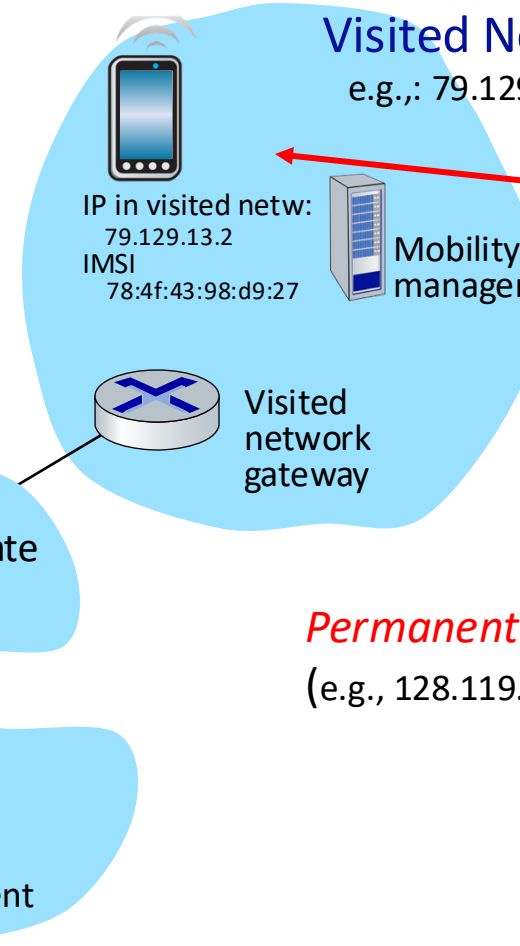
Visited Network

e.g., 79.129/16

Address in visited
network.
(e.g., 79.129.13.2)

*Known as care-of-address
(CoA) in mobile IP.*

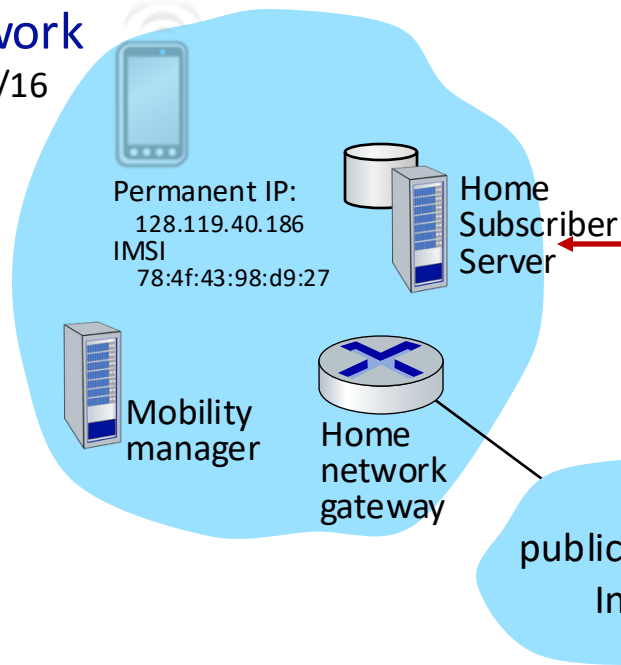
Permanent address: remains constant
(e.g., 128.119.40.186)



Registration: home needs to know where you are!

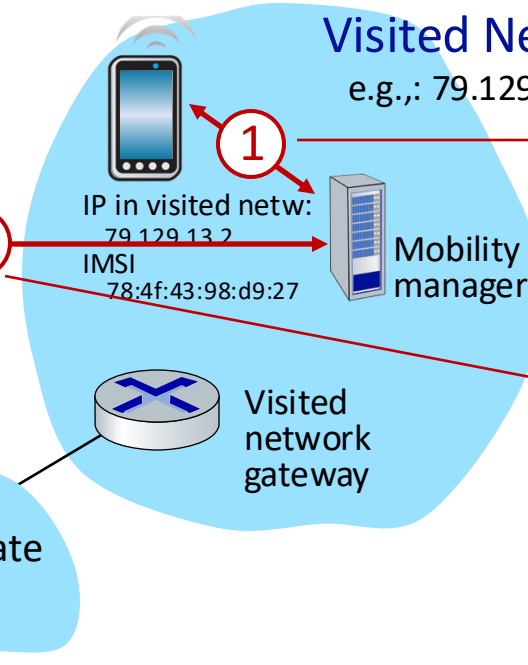
Home Network

e.g.,: 128.119/16



Visited Network

e.g.,: 79.129/16



2

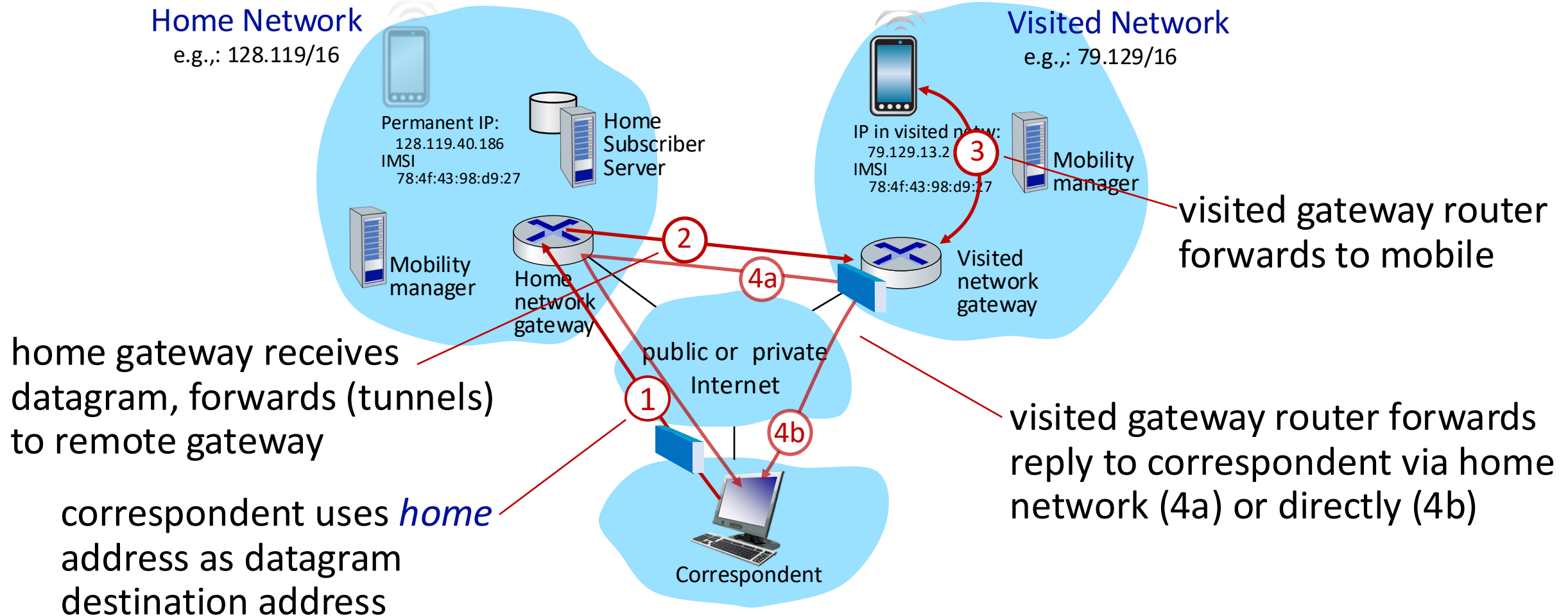
mobile *associates* with visited mobility manager

visited mobility manager *registers* mobile's location with home HSS

end result:

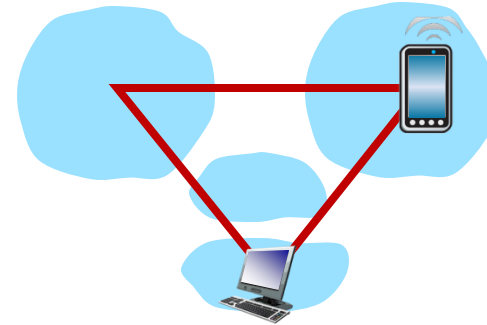
- visited mobility manager knows about mobile
- home HSS knows location of mobile

Mobility with indirect routing

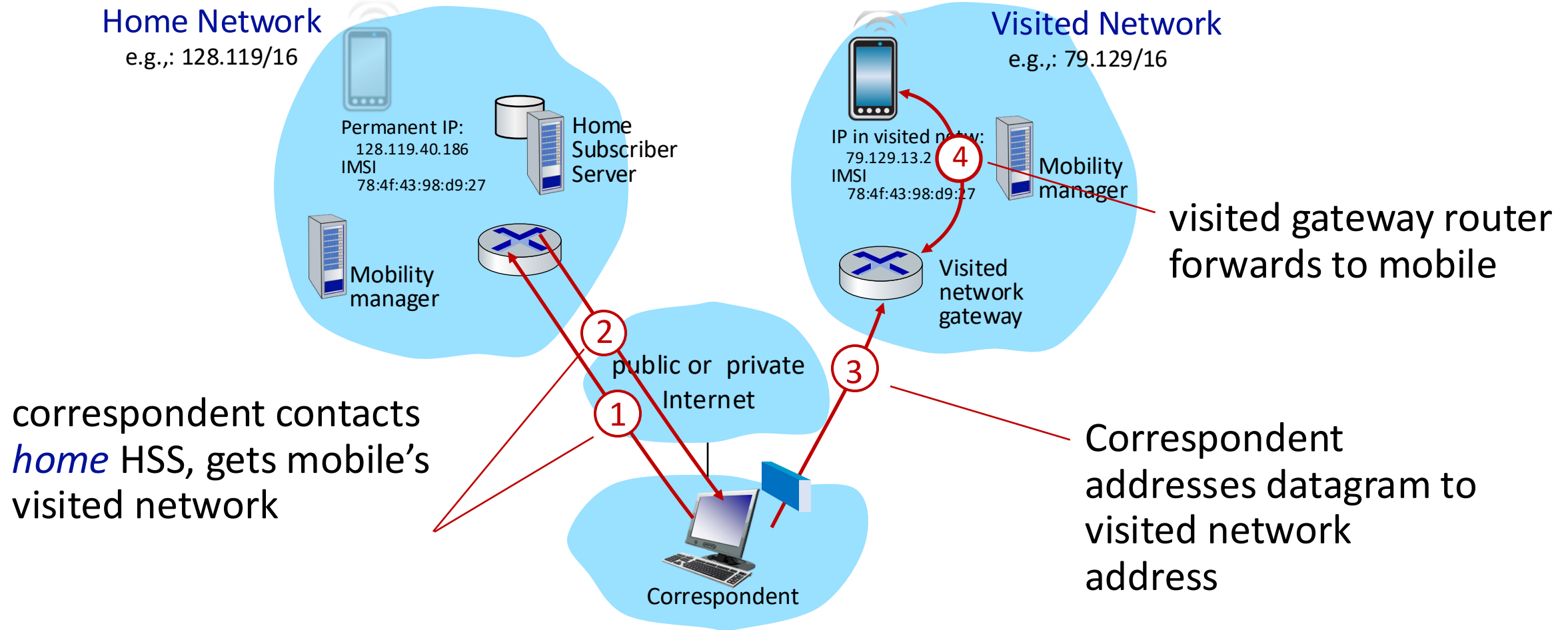


Mobility with indirect routing: comments

- triangle routing:
 - inefficient when correspondent and mobile are in same network
- mobile moves among visited networks: transparent to correspondent!
 - registers in new visited network
 - new visited network registers with home HSS
 - datagrams continue to be forwarded from home network to mobile in new network
 - *on-going (e.g., TCP) connections between correspondent and mobile can be maintained!*



Mobility with direct routing



Mobility with direct routing: comments

- overcomes triangle routing inefficiencies
- *non-transparent to correspondent*: correspondent must get care-of-address from home agent
- what if mobile changes visited network?
 - can be handled, but with additional complexity

Mobile IP

- mobile IP architecture standardized ~20 years ago [RFC 5944]
 - did not see wide deployment/use
 - perhaps WiFi for Internet, and 2G/3G phones for voice were “good enough” at the time
- mobile IP architecture:
 - indirect routing to node (via home network) using tunnels
 - mobile IP home agent: combined roles of 4G HSS and home P-GW
 - mobile IP foreign agent: combined roles of 4G MME and S-GW
 - protocols for agent discovery in visited network, registration of visited location in home network via ICMP extensions

Wireless, mobility: impact on higher layer protocols

- logically, impact *should* be minimal ...
 - best effort service model remains unchanged
 - TCP and UDP can (and do) run over wireless, mobile
- ... but performance-wise:
 - packet loss/delay due to bit-errors (discarded packets, delays for link-layer retransmissions), and handover loss
 - TCP interprets loss as congestion, will decrease congestion window unnecessarily
 - delay impairments for real-time traffic
 - bandwidth a scarce resource for wireless links