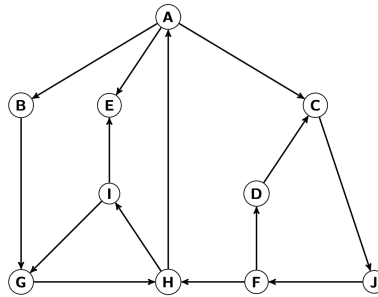


## Week 12: Lab

Module: Graphs

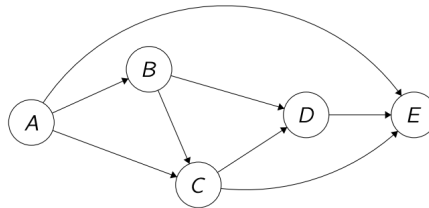
COLLABORATION LEVEL 0 (NO RESTRICTIONS). OPEN NOTES.

1. Consider the graph below:



- What are all the strongly connected components?
- Perform DFS on the graph above starting from vertex A. Assume that the adjacency lists are ordered in alphabetical order. As you go, label each vertex with the start and finish time. Draw the DFS-tree and highlight the tree edges on the graph.
- Perform BFS on the graph above starting from vertex A. Again, assume the adjacency lists are ordered in alphabetical order. Draw the BFS-tree and highlight the tree edges on the graph.

2. Consider the graph below:



- Without using any specific algorithm, come up with a topological order for the graph. How many different orders can you come up with?
- Run DFS on the whole graph starting at vertex C, and consider the edges in the adjacency lists in alphabetical order. Recall that when you run DFS on the graph, if it stops but has not yet explored the whole graph, it will start again at the next unexplored (white) vertex.
  - What do you get when you order the vertices by ascending start time?
  - What do you get when you order the vertices by descending finish time?

(c) Now run DFS starting at vertex C, but consider the edges in the adjacency lists in reverse alphabetical order.

- What do you get when you order the vertices by ascending start time?
  - What do you get when you order the vertices by descending finish time?
3. Explain why the following algorithm does not necessarily produce a topological order: Run BFS, and label the vertices by increasing distance to their respective sources.

Note: To prove that a certain algorithm does not work, it is sufficient to show a counterexample.

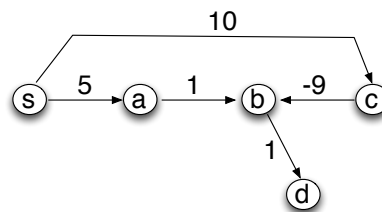
4. Consider a directed graph  $G$ , and assume that instead of shortest paths we want to compute *longest paths*. Longest paths are defined in the natural way, i.e. the longest path from  $u$  to  $v$  is the path of maximum weight among all possible paths from  $u$  to  $v$ . Note that if the graph contains a positive cycle, then longest paths are not well defined (for the same reason that shortest paths are not well defined when the graph has a negative cycle). So what we mean is the *longest simple path*, (a path is called *simple* if it contains no vertex more than once).

Show that the *longest simple path* problem does not have optimal substructure by coming up with a small graph that provides a counterexample. Note: Finding longest (simple) paths is a classical *hard* problem, and it is known to be NP-complete.

5. **Shortest paths on DAGs:** Consider a directed acyclic graph  $G$ . Unlike the problems so far, in which we considered that the edges are “equal” and we counted the weight of a path as the number of edges on the path, in this problem we’ll assume that the graph is *weighted*. What that means is that each edge  $(u, v)$  has a weight which we denote by  $w_{uv}$ . The weight represents the cost of traversing that edge (for e.g. the weight could be the distance between the vertices). The weight of a path in a weighted graph is defined as the sum of the weights of the edges along the path.

So: given a weighted DAG and two arbitrary vertices  $u$  and  $v$ , the goal is to compute the **shortest path** from  $u$  to  $v$ . Note that if the graph was not weighted or if all edges weights were equal, we could simply run BFS from  $u$ ; but BFS does not guarantee shortest paths when there are weights.

Explore an algorithm for finding the shortest path from  $u$  to  $v$ . You can assume, for simplicity, that  $u$  is a vertex without any incoming edges.



Hint: Argue that shortest paths have optimal substructure, choose a subproblem, express the problem recursively and from here go to a recursive dynamic programming solution.