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MODULE *DistributedLock*

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EXTENDS *Naturals, FiniteSets, Sequences, TLC*

**The set of clients**  
 CONSTANT *Clients*

**Client states**  
 CONSTANTS *Active, Inactive*

**Message types**  
 CONSTANT *LockRequest, LockResponse, TryLockRequest, TryLockResponse, UnlockRequest, UnlockResponse*

**An empty constant**  
 CONSTANT *Nil*

**The current lock holder**  
 VARIABLE *lock*

**The lock queue**  
 VARIABLE *queue*

**The current lock ID**  
 VARIABLE *id*

$serverVars \triangleq \langle lock, id, queue \rangle$

**Client states**  
 VARIABLE *clients*

$clientVars \triangleq \langle clients \rangle$

**Client messages**  
 VARIABLE *messages*

**Variable**  
 VARIABLE *messageCount*

$messageVars \triangleq \langle messages, messageCount \rangle$

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**The invariant checks that no client can hold more than one lock at a time**  
 $TypeInvariant \triangleq$   
 $\wedge \forall c \in DOMAIN\ clients : Cardinality(clients[c].locks) \in 0 .. 1$

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**Returns a sequence with the head removed**  
 $Pop(q) \triangleq SubSeq(q, 2, Len(q))$

Sends a message on the given client's channel  
 $Send(m, c) \triangleq$   
 $\wedge messages' = [messages \text{ EXCEPT } ![c] = Append(messages[c], m)]$   
 $\wedge messageCount' = messageCount + 1$

Removes a message from the given client's channel  
 $Accept(m, c) \triangleq$   
 $\wedge messages' = [messages \text{ EXCEPT } ![c] = Pop(messages[c])]$   
 $\wedge messageCount' = messageCount + 1$

Removes the last message and appends a message to the given client's channel  
 $Reply(m, c) \triangleq$   
 $\wedge messages' = [messages \text{ EXCEPT } ![c] = Append(Pop(messages[c]), m)]$   
 $\wedge messageCount' = messageCount + 1$

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Handles a lock request. If the lock is not currently held by another process, the lock is granted to the client. If the lock is held by a process, the request is added to a queue.

$HandleLockRequest(m, c) \triangleq$   
 $\vee \wedge lock = Nil$   
 $\wedge lock' = m @ @ ("client" :> c)$   
 $\wedge id' = id + 1$   
 $\wedge Reply([type \mapsto LockResponse, acquired \mapsto TRUE, id \mapsto id'], c)$   
 $\wedge UNCHANGED \langle queue, clientVars \rangle$   
 $\vee \wedge lock \neq Nil$   
 $\wedge queue' = Append(queue, m @ @ ("client" :> c))$   
 $\wedge Accept(m, c)$   
 $\wedge UNCHANGED \langle lock, id, clientVars \rangle$

Handles a *tryLock* request. If the lock is not currently held by another process, the lock is granted to the client. Otherwise, the request is rejected.

$HandleTryLockRequest(m, c) \triangleq$   
 $\vee \wedge lock = Nil$   
 $\wedge lock' = m @ @ ("client" :> c)$   
 $\wedge id' = id + 1$   
 $\wedge Reply([type \mapsto LockResponse, acquired \mapsto TRUE, id \mapsto id'], c)$   
 $\wedge UNCHANGED \langle queue, clientVars \rangle$   
 $\vee \wedge lock \neq Nil$   
 $\wedge Reply([type \mapsto LockResponse, acquired \mapsto FALSE], c)$   
 $\wedge UNCHANGED \langle clientVars, serverVars \rangle$

Handles an unlock request. If the lock is currently held by the given client, it will be unlocked. If any client's requests are pending in the queue, the next lock request will be removed from the queue and the lock will be granted to the requesting client.

$HandleUnlockRequest(m, c) \triangleq$   
 $\vee \wedge lock = Nil$   
 $\wedge Accept(m, c)$



Sends a lock request to the cluster with a unique *ID* for the client.

$$\begin{aligned}
\text{Lock}(c) &\triangleq \\
&\wedge \text{clients}[c].\text{state} = \text{Active} \\
&\wedge \text{Send}([type \mapsto \text{LockRequest}, id \mapsto \text{clients}[c].\text{next}], c) \\
&\wedge \text{clients}' = [\text{clients} \text{ EXCEPT } ![c].\text{next} = \text{clients}[c].\text{next} + 1] \\
&\wedge \text{UNCHANGED } \langle \text{serverVars} \rangle
\end{aligned}$$

Sends a try lock request to the cluster with a unique *ID* for the client.

$$\begin{aligned}
\text{TryLock}(c) &\triangleq \\
&\wedge \text{clients}[c].\text{state} = \text{Active} \\
&\wedge \text{Send}([type \mapsto \text{TryLockRequest}, id \mapsto \text{clients}[c].\text{next}], c) \\
&\wedge \text{clients}' = [\text{clients} \text{ EXCEPT } ![c].\text{next} = \text{clients}[c].\text{next} + 1] \\
&\wedge \text{UNCHANGED } \langle \text{serverVars} \rangle
\end{aligned}$$

Sends an unlock request to the cluster if the client is active and current holds a lock.

$$\begin{aligned}
\text{Unlock}(c) &\triangleq \\
&\wedge \text{clients}[c].\text{state} = \text{Active} \\
&\wedge \text{Cardinality}(\text{clients}[c].\text{locks}) > 0 \\
&\wedge \text{Send}([type \mapsto \text{UnlockRequest}, id \mapsto \text{CHOOSE } l \in \text{clients}[c].\text{locks} : \text{TRUE}], c) \\
&\wedge \text{clients}' = [\text{clients} \text{ EXCEPT } ![c].\text{locks} = \text{clients}[c].\text{locks} \setminus \{\text{CHOOSE } l \in \text{clients}[c].\text{locks} : \text{TRUE}\}] \\
&\wedge \text{UNCHANGED } \langle \text{serverVars} \rangle
\end{aligned}$$

Handles a lock response from the cluster. If the client's session is expired, the response is ignored. If the lock was acquired successfully, it's added to the client's lock set.

$$\begin{aligned}
\text{HandleLockResponse}(m, c) &\triangleq \\
&\wedge \vee \wedge \text{clients}[c].\text{state} = \text{Inactive} \\
&\quad \wedge \text{UNCHANGED } \langle \text{clientVars}, \text{serverVars} \rangle \\
&\vee \wedge \text{clients}[c].\text{state} = \text{Active} \\
&\quad \wedge m.\text{acquired} \\
&\quad \wedge \text{clients}' = [\text{clients} \text{ EXCEPT } ![c].\text{locks} = \text{clients}[c].\text{locks} \cup \{m.\text{id}\}] \\
&\quad \wedge \text{UNCHANGED } \langle \text{serverVars} \rangle \\
&\vee \wedge \text{clients}[c].\text{state} = \text{Active} \\
&\quad \wedge \neg m.\text{acquired} \\
&\quad \wedge \text{UNCHANGED } \langle \text{clientVars}, \text{serverVars} \rangle \\
&\wedge \text{Accept}(m, c)
\end{aligned}$$

Receives a message from/to the given client from the head of the client's message queue.

$$\begin{aligned}
\text{Receive}(c) &\triangleq \\
&\wedge \text{Len}(\text{messages}[c]) > 0 \\
&\wedge \text{LET } \text{message} \triangleq \text{Head}(\text{messages}[c]) \\
&\quad \text{IN} \\
&\quad \vee \wedge \text{message.type} = \text{LockRequest} \\
&\quad \quad \wedge \text{HandleLockRequest}(\text{message}, c)
\end{aligned}$$

$$\begin{aligned}
& \vee \wedge message.type = LockResponse \\
& \quad \wedge HandleLockResponse(message, c) \\
& \vee \wedge message.type = TryLockRequest \\
& \quad \wedge HandleTryLockRequest(message, c) \\
& \vee \wedge message.type = UnlockRequest \\
& \quad \wedge HandleUnlockRequest(message, c)
\end{aligned}$$


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Initial state predicate

$Init \triangleq$

$$\begin{aligned}
& \wedge messages = [c \in Clients \mapsto \langle \rangle] \\
& \wedge messageCount = 0 \\
& \wedge lock = Nil \\
& \wedge queue = \langle \rangle \\
& \wedge id = 0 \\
& \wedge clients = [c \in Clients \mapsto [state \mapsto Active, locks \mapsto \{\}, next \mapsto 1]]
\end{aligned}$$

Next state predicate

$Next \triangleq$

$$\begin{aligned}
& \vee \exists c \in DOMAIN \ clients : Receive(c) \\
& \vee \exists c \in DOMAIN \ clients : Lock(c) \\
& \vee \exists c \in DOMAIN \ clients : TryLock(c) \\
& \vee \exists c \in DOMAIN \ clients : Unlock(c) \\
& \vee \exists c \in DOMAIN \ clients : ExpireSession(c)
\end{aligned}$$

The specification includes the initial state predicate and the next state

$Spec \triangleq Init \wedge \Box [Next]_{\langle serverVars, clientVars, messageVars \rangle}$

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\ \* Modification History  
\ \* Last modified Sat Jan 27 01:50:12 PST 2018 by jordanhalterman  
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