## Unconditional Branching: Jump

Already saw conditional branch statements

Sometimes, we want to move to another section of code regardless of any condition

Can you think of an example?

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Sometimes, we want to move to another section of code regardless of any condition

Can you think of an example?
Function calls
Also helpful for loops

What about the following?

beq \$0, \$0, label

What about the following?

beq \$0, \$0, label

Works correctly, but has limited range – branch is I-type instruction with 16-bit immediate

By leaving out registers, we can save more room for address

#### Jumps

Unconditional branch is called *jump* 

# MIPS Assembly Code addi \$s0, \$0, 4 # \$s0 = 4 addi \$s1, \$0, 1 # \$s1 = 1 j target # jump to target addi \$s1, \$s1, 1 # not executed sub \$s1, \$s1, \$s0 # not executed target: add \$s1, \$s1, \$s0 # \$s1 = 1 + 4 = 5

#### Jumps

That's about it – jump is very simple instruction

Note that jump instruction uses labels, just like branch instructions

However, way jump represents labels is different

In fact, jump gets its own special instruction type

#### J-Type Instructions

Third (and final!) MIPS instruction format

#### Includes:

- 6-bit opcode (same as all instructions)
- Single 26-bit address operand

### J-type op addr 6 bits 26 bits

#### J-Type Instructions

Quick note – we will see two J-type instructions:

- j (jump)
- jal (jump and link)

Don't worry about difference for now. jal will be helpful for writing functions in assembly. Until you know you need it, just use j

**MIPS Assembly Code** 

0x0040005C .jal sum

. . .

0x004000A0 sum: add \$v0, \$a0, \$a1

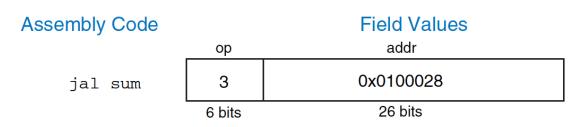
Once again, we cannot store full 32-bit address in instruction

Assume top 4 bits come from PC+4

Append 26 bits we have

Shift result by 2 bits (multiply by 4)

Result is the address of next instruction





JTA 0000 0000 0100 0000 0000 1010 0000 (0x004000A0)

26-bit addr 0000 0000 0100 0000 0000 1010 0000 (0x0100028)

0 1 0 0 0 2 8

Seems odd at first, but quite straightforward once we understand the two "tricks":

- Multiplying by 4
- Taking first 4 bits from PC + 4

Remember: reason for all this trouble is that we cannot store all 32 bits of address

#### Multiplying by 4

Instructions are word-aligned – every instruction starts at multiple of 4

#### Taking first 4 bits from PC + 4

Need 4 bits from somewhere – assume we are not jumping extremely far away

"Pseudo"-direct because we cannot fit full 32-bit address into 28 bits

Turns out not to matter – see later that all MIPS code fits within subset of 32-bit address space

Remainder of 32-bit address space reserved for data (variables, arrays, etc.)

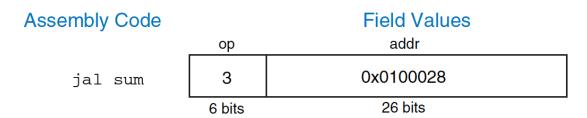
#### **MIPS Assembly Code**

0x0040005C jal

. . .

0x004000A0 sum: add \$v0, \$a0, \$a1

Essentially, we are just storing all bits of address that matter



sum



JTA 0000 0000 0100 0000 0000 1010 0000 (0x004000A0)

26-bit addr 0000 0000 0100 0000 0000 1010 0000 (0x0100028)

0 1 0 0 0 2 8

#### One more jump

There is third jump instruction in addition to j (jump) and jal (jump and link)

Final jump instruction is: jr (jump register)

Register holds 32-bit value, so no fancy addressing needed – just jump to location specified in register

#### Jump register

## MIPS Assembly Code 0x00002000 addi \$s0, \$0, 0x2010 #\$s0 = 0x2010 0x00002004 jr \$s0 # jump to 0x00002010 0x00002008 addi \$s1, \$0, 1 # not executed 0x0000200c sra \$s1, \$s1, 2 # not executed 0x00002010 lw \$s3, 44(\$s1) # executed after jr instruction

#### Jump register

You should have noticed something odd – J-type instructions do not have register as argument

Even though it is functionally a jump, jr is an R-type instruction

Remember to separate what an instruction does from how it is stored

#### Jump register

Like jal, jr is useful when writing functions

In fact, the two work together to ensure we can move to and from functions correctly

For now, just remember to end all of your functions with jr \$ra