

# Implementing Conditional Branches

# I-Type Instructions

Previously made circuit that worked for I-type instructions, then combined with R-type circuit

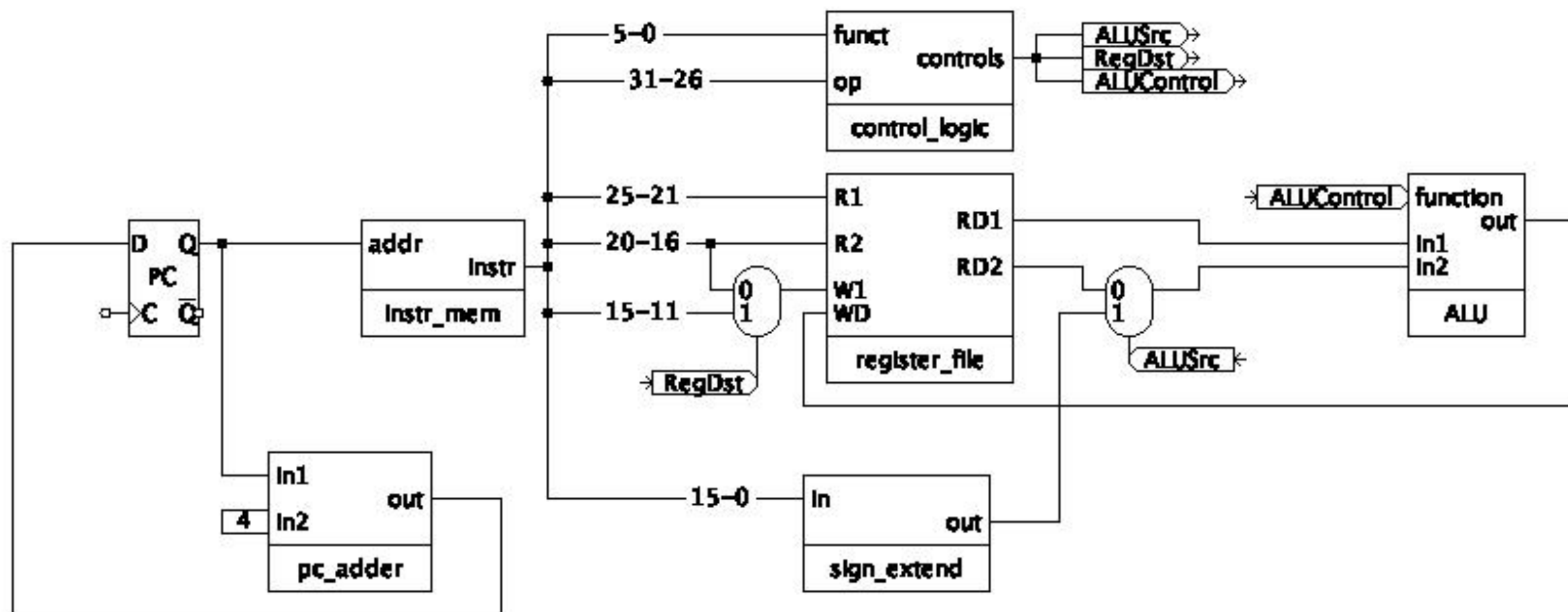
In reality, our circuit was only for subset of I-type instructions

Even though branches seem very different functionally, they are still I-type because of how they are stored

# I-Type Instructions

We follow similar process again:

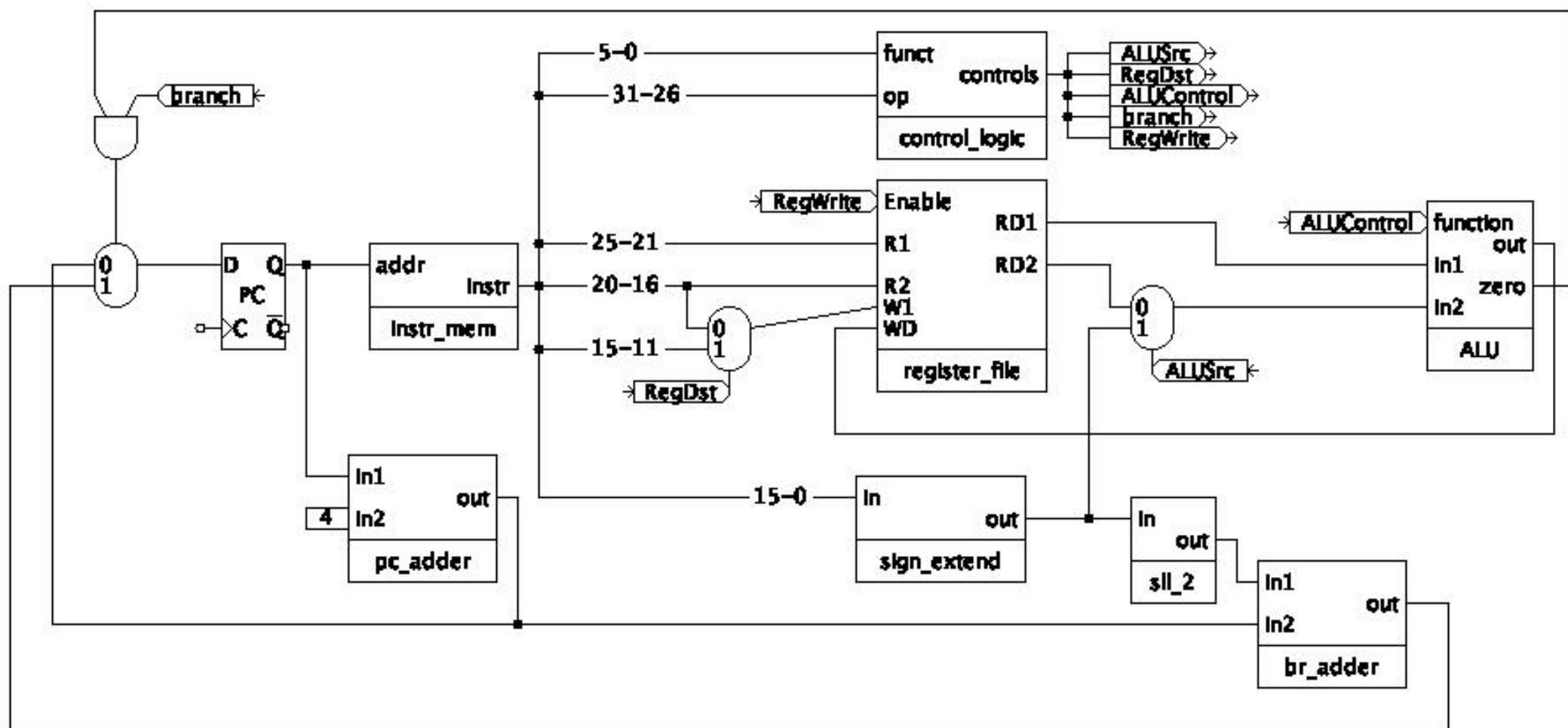
- Add capability for branches to circuit
- Use muxes to ensure backward compatibility



# Adding branches

In the big picture, there are three ways branch is different from, e.g., `addi`. Branch instruction must

- Compute PC value of branch target
- Determine whether condition is satisfied
- *Not* write to general-purpose registers



# Modifications to hardware

register\_file:

- Enable bit added – if input is 1, registers writable; otherwise, registers do not change
- Needed because we cannot “cut wires” for different inputs – WD will have *something* going into it even for branch instruction
- Register must be able to ignore bogus inputs – no register is modified during branch instruction

ALU output zero:

- Answering yes/no question: was output 0?
- Single bit – 1 for yes, 0 for no
- Feels backwards at first – 1 if output 0, 0 if output anything nonzero
- Used for checking branch condition

# New control logic

branch:

- Whether instruction is branch – 1 for yes, 0 for no
- Branch taken only if branch is 1 *and* operands are equal

RegWrite:

- Whether register should be written to – 1 for yes, 0 for no
- Do not want to accidentally overwrite register during branch



