CIS 452 - Operating Systems Concepts Nathan Bowman Images taken from Silberschatz book

Paging

Multiprogramming setups discussed so far:

- contiguous memory allocation
- segmentation

Paging is another way of managing memory that does not require contiguous allocation for a process

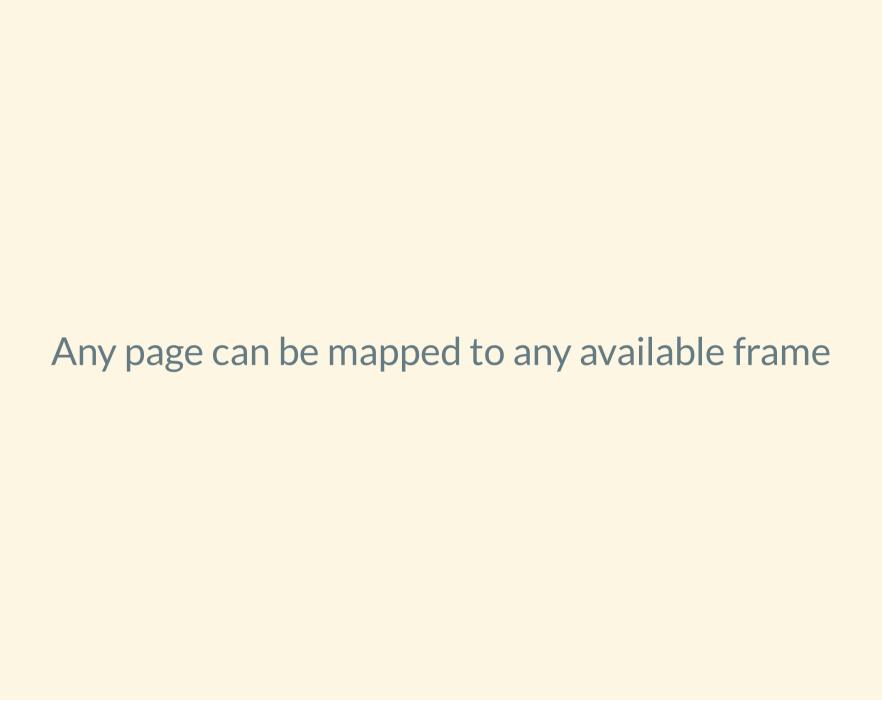
Paging

Physical memory and logical memory are split into fixed-size blocks

Blocks of physical memory are frames

Blocks of logical memory are pages

Pages and frames are the same size

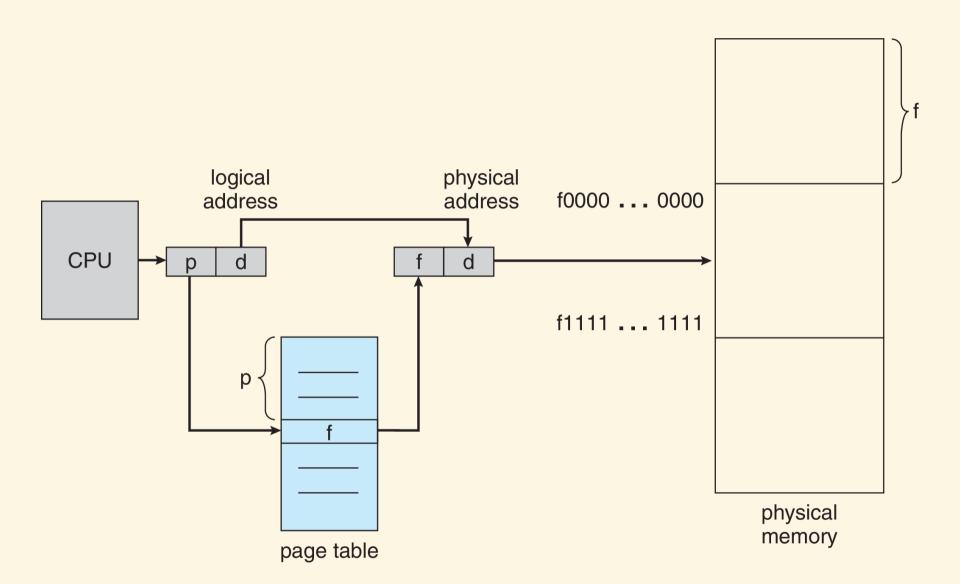


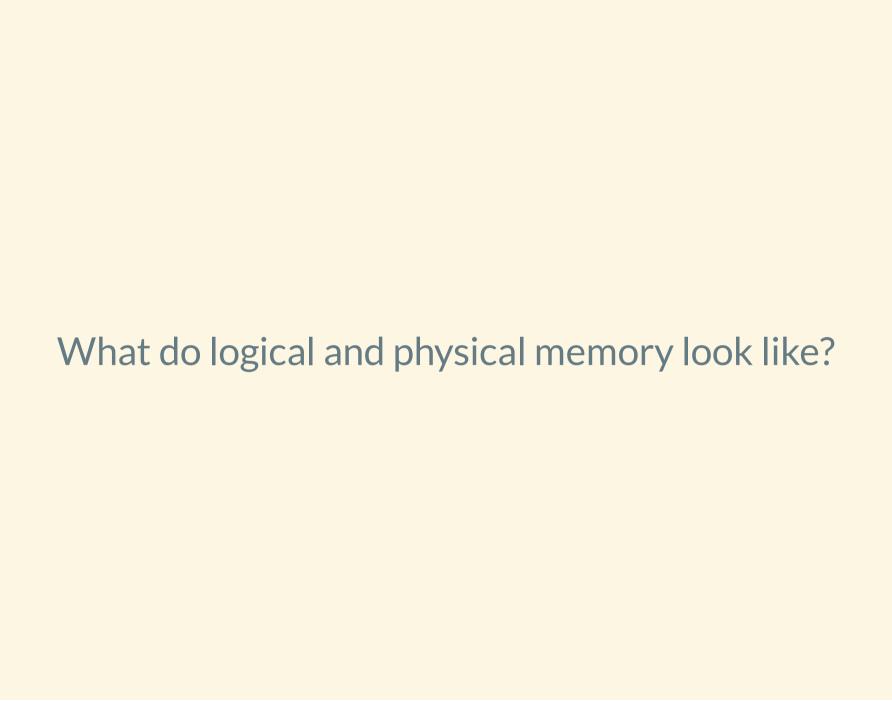
How does this memory translation happen?

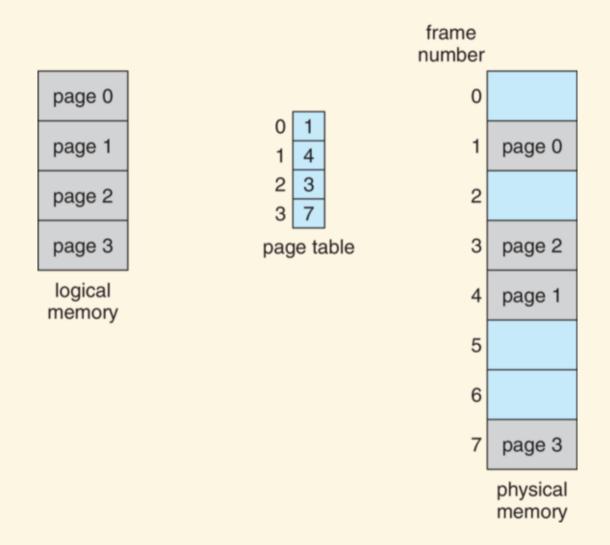
Every address is split into

- page number (p)
- page offset (d)

OS keeps **page table** mapping page numbers to frame numbers







Notice that (binary) address is broken up into two sets of bits

For a memory of $128 (= 2^7)$ bytes, e.g.,

```
101 0000
--- ---
p d
```

So, both number of pages and size of page will be powers of two

(This splitting of addresses may remind you of cache mapping from 351)

Some general rules:

If logical memory holds 2^m bytes, logical address will have m bits

If page size is 2ⁿ bytes, page offset requires n bits

Remaining (m - n) bits specify page number, and logical memory has 2^(m - n) pages

page number

page offset

p	d

m-n

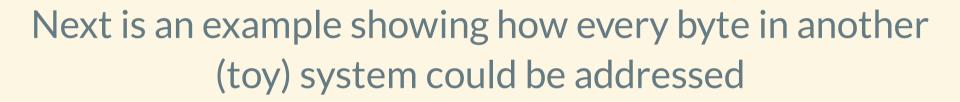
n

Assume 128 bytes of memory again Page size is 4 bytes Split up the address:

1010000

```
10100 00
----- --
p d
```

Page 20, offset 0



0	а		
1	b		
2	С		
2 3 4	d		
4	е		
5	f		
6	g		
7_	h i j k		
8	į į		
9	j		
10	k		
11			
12	m		
13	n		
14	0		
_15	р		
aical memo			

0	5			
1	6			
2	1			
3	2			
page table				

logical memory

i j k l
Впор
a b c d
e f g h

physical memory

Note that logical memory and physical memory can have different sizes

Translation is completely transparent to process

Process once again believes it has memory from [0, max] and has no idea that pages are actually scattered throughout memory