Chapter 6 The Link Layer and LANs

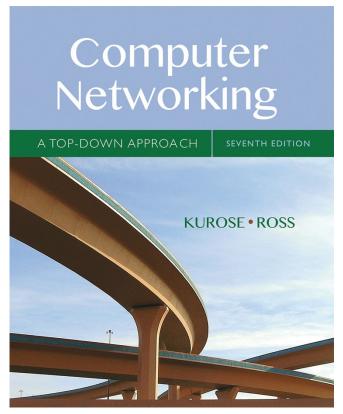
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Computer Networking: A Top Down Approach

7th edition
Jim Kurose, Keith Ross
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Minor modifications made to original slides by Nathan Bowman

Link layer, LANs: outline

- 6. I introduction, services
- 6.2 error detection, correction
- 6.3 multiple access protocols
- 6.4 LANs
 - addressing, ARP
 - Ethernet
 - switches
 - VLANS

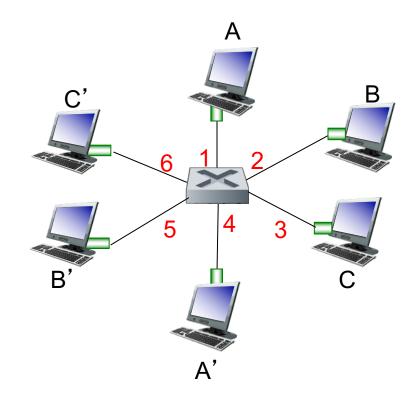
- 6.5 link virtualization: MPLS
- 6.6 data center networking
- 6.7 a day in the life of a web request

Ethernet switch

- link-layer device: takes an active role
 - store, forward Ethernet frames
 - examine incoming frame's MAC address, selectively forward frame to one-or-more outgoing links when frame is to be forwarded on segment, uses CSMA/CD to access segment
- transparent
 - hosts are unaware of presence of switches
- plug-and-play, self-learning
 - switches do not need to be configured

Switch: multiple simultaneous transmissions

- hosts have dedicated, direct connection to switch
- switches buffer packets
- Ethernet protocol used on each incoming link, but no collisions; full duplex
 - each link is its own collision domain
- switching: A-to-A' and B-to-B' can transmit simultaneously, without collisions



switch with six interfaces (1,2,3,4,5,6)

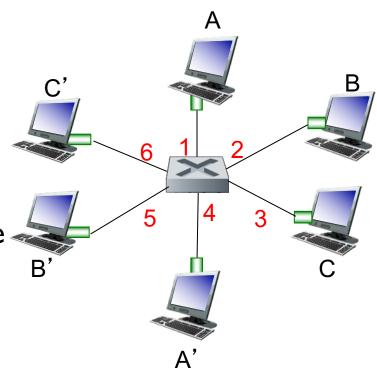
Switch forwarding table

Q: how does switch know A' reachable via interface 4, B' reachable via interface 5?

- A: each switch has a switch table, each entry:
 - (MAC address of host, interface to reach host, time stamp)
 - looks like a routing table!

Q: how are entries created, maintained in switch table?

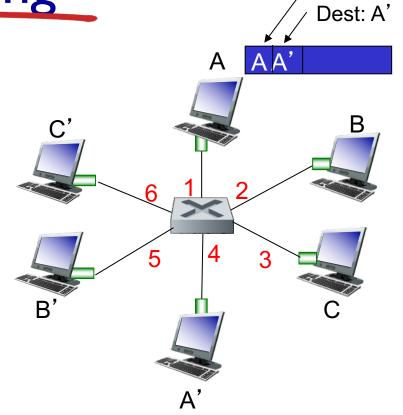
something like a routing protocol?



switch with six interfaces (1,2,3,4,5,6)

Switch: self-learning

- switch learns which hosts can be reached through which interfaces
 - when frame received, switch "learns" location of sender: incoming LAN segment
 - records sender/location pair in switch table



MAC addr	interface	TTL	
Α	1	60	

Switch table (initially empty)

Source: A

Switch: frame filtering/forwarding

when frame received at switch:

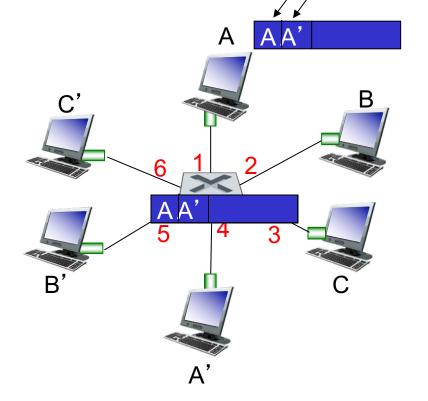
- I. record incoming link, MAC address of sending host
- 2. index switch table using MAC destination address

```
3. if entry found for destination then {
if destination on segment from which frame arrived then drop frame
else forward frame on interface indicated by entry }
else flood /* forward on all interfaces except arriving interface */
```

Self-learning, forwarding: example

Source: A Dest: A'

- frame destination, A', location unknown: flood
- destination A location known: selectively send on just one link



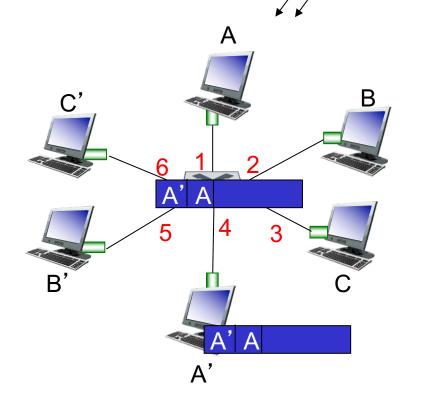
MAC addr	interface	TTL
Α	1	60
Α'	4	60

switch table (initially empty)

Self-learning, forwarding: example

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/ Dest: A'

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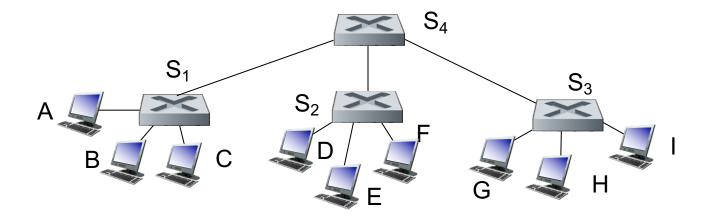


MAC addr	interface	TTL
Α	1	60
Α'	4	60

switch table (initially empty)

Interconnecting switches

self-learning switches can be connected together:

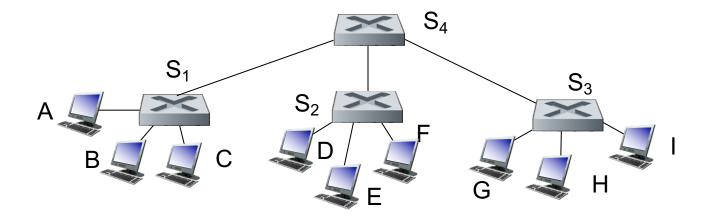


Q: sending from A to G - how does S_1 know to forward frame destined to G via S_4 and S_3 ?

A: self learning! (works exactly the same as in single-switch case!)

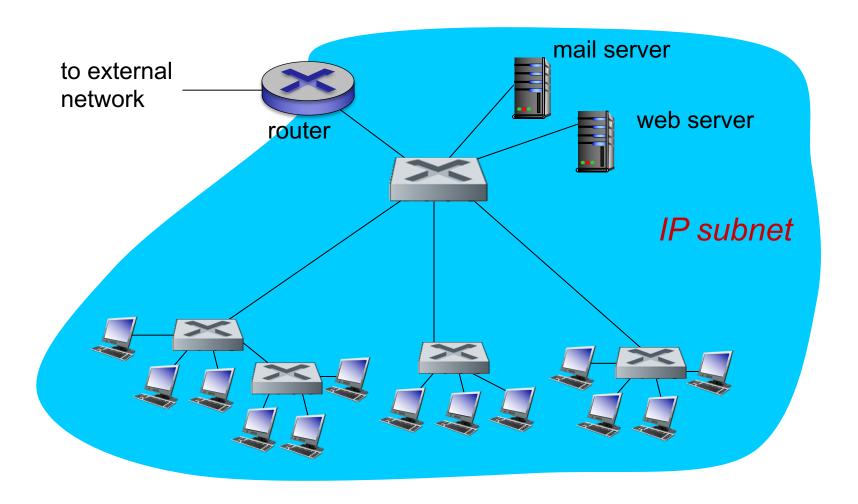
Self-learning multi-switch example

Suppose C sends frame to I, I responds to C



• Q: show switch tables and packet forwarding in S_1 , S_2 , S_3 , S_4

Institutional network



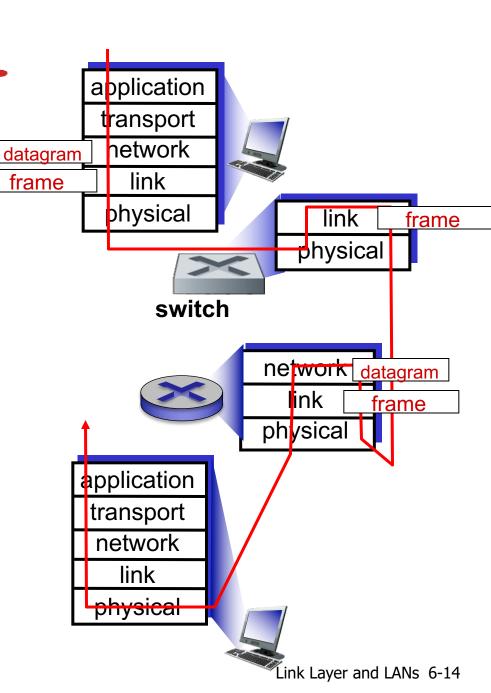
Switches vs. routers

both are store-and-forward:

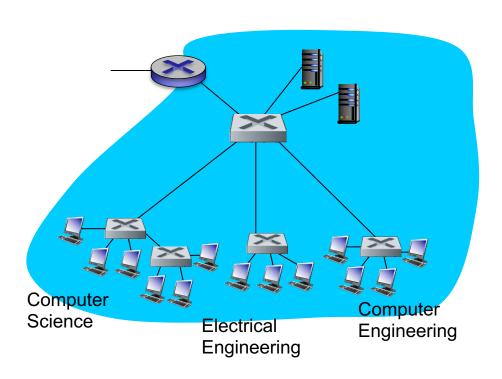
- routers: network-layer devices (examine networklayer headers)
- switches: link-layer devices (examine link-layer headers)

both have forwarding tables:

- routers: compute tables using routing algorithms, IP addresses
- switches: learn forwarding table using flooding, learning, MAC addresses



VLANs: motivation



consider:

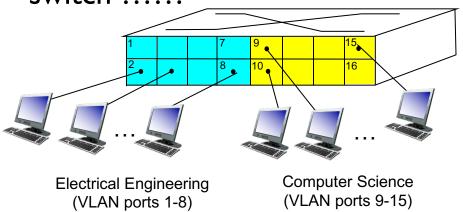
- CS user moves office to EE, but wants connect to CS switch?
- single broadcast domain:
 - all layer-2 broadcast traffic (ARP, DHCP, unknown location of destination MAC address) must cross entire LAN
 - security/privacy, efficiency issues

VLANs

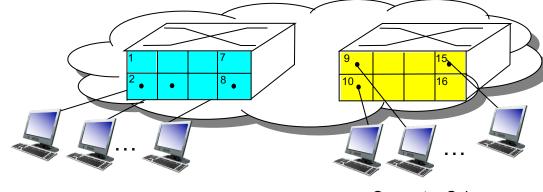
Virtual Local Area Network

switch(es) supporting VLAN capabilities can be configured to define multiple *virtual* LANS over single physical LAN infrastructure.

port-based VLAN: switch ports grouped (by switch management software) so that single physical switch



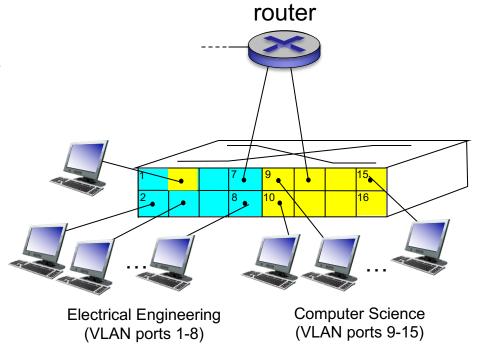
... operates as multiple virtual switches



Computer Science (VLAN ports 9-16)

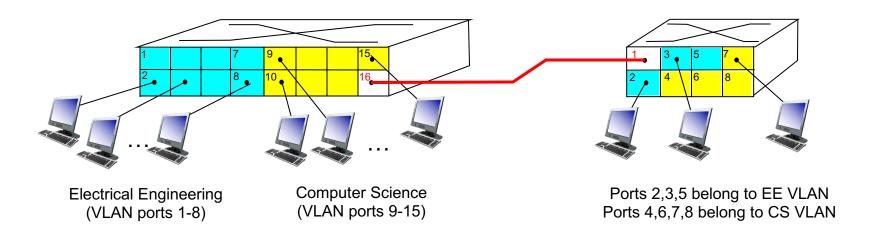
Port-based VLAN

- traffic isolation: frames to/from ports 1-8 can only reach ports 1-8
 - can also define VLAN based on MAC addresses of endpoints, rather than switch port
- dynamic membership: ports can be dynamically assigned among VLANs



- forwarding between VLANS: done via routing (just as with separate switches)
 - in practice vendors sell combined switches plus routers

VLANS spanning multiple switches



- trunk port: carries frames between VLANS defined over multiple physical switches
 - frames forwarded within VLAN between switches can't be vanilla 802.1 frames (must carry VLAN ID info)
 - 802. Iq protocol adds/removed additional header fields for frames forwarded between trunk ports

802. I Q VLAN frame format

