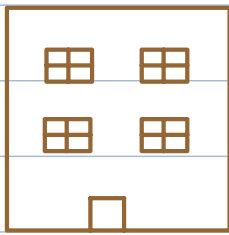


Hotel



← Vishal

301

Key
Unique

Room	Available
1	✓
2	✗
3	✓
⋮	
301	✓
⋮	

Value

, capacity, AC, etc
Multiple

Notebook

Hash Map

Q → Store population of every country →
country name → population
string → long

Q → Store # states of every country →
country name → # states
string → int

Q → Store name of all states of every country →
country name → list of states
string → list of string

Q → Store population of each state in every country →
country name → (state → population)
string → HM < string, int >

< Key, Null > → < Key > Hash Set

✓ Hash Set / Hash Map → Store Keys in random order.
Hashing $TC = O(1)$

Tree Set / Tree Map → Maintain Keys in Sorted Order

(use balanced BST internally) $TC = O(\log N)$

Linked Hash Map / Linked Hash Set → Maintain Keys in order of insertion.

Operations

To update → insert the same key again.

HashMap

- 1) $insert(key, value) / put(key, value)$
- 2) $size()$
- 3) $delete(key) / remove(key)$
- 4) $get(key)$
- 5) $containsKey(key)$
- ⋮

$TC = O(1)$

HashSet

- 1) $insert(key) / add(key)$
- 2) $size()$
- 3) $delete(key) / remove(key)$
- 4) $contains(key)$
- ⋮

$TC = O(1)$

Q → Given N elements & Q queries.

Find frequency of an element for multiple queries.

(Ans)

$A = [2 \quad \underline{6} \quad 3 \quad \underline{6} \quad 3 \quad \textcircled{8} \quad 2 \quad \underline{6}]$								<u>Query</u>	<u>Frequency</u>
0	1	2	3	4	5	6	7	6	3

Brute force → \forall query, travel & find frequency.

$$TC = \underline{O(Q * N)}$$

$$SC = \underline{O(1)}$$

8	1
5	0

Frequency Array → $F[i] \rightarrow$ frequency of i
 ↑ ↑
 key value

```
for i → 0 to (N-1) {  
    | F[A[i]] ++  
}
```

// Find freq. of $x \rightarrow F[x]$

$$SC = \underline{O(\text{Range of } A[i])}$$

 ✓ ✓ ✓ ✓ ✓ ✓ ✓ ✓
 0 1 2 3 4 5 6 7
A = [2 1 3 3 0 3 2 2]
 0 1 2 3
F = [1 1 3 3]

Hash Map → $\langle \text{int}, \text{int} \rangle$

$A[i] \rightarrow$ frequency of $A[i]$

```
for i → 0 to (N-1) {  
    | if (hm.containsKey(A[i])) {  
        | v = hm.get(A[i])  
        | hm.put(A[i], v+1)  
    } else {  
        | hm.put(A[i], 1)  
    }  
}
```

```

for i → 0 to (Q-1) {
    x = B[i] // B → queries
    if (hm.containsKey(x))
        print(hm.get(x))
    else
        print(0)
}

```

$hm.getOrDefault(x, 0)$

$TC = O(N + Q)$ $SC = O(N)$

Q → Given N integers, find first non-repeating value.
 Return -1 if all elements are repeating.

N = 6 A = [1⁰ 3¹ 5² 1³ 2⁴ 3⁵]
 → ✓
 Ans = 5

N = 5 A = [2 3 2 4 4] Ans = 3
 ✓

Bruteforce → ∀ elements check if that element
 is repeating in future. $TC = O(N^2)$ $SC = O(1)$

Sol → 1) store frequency ∀ A[i].
 2)

```

for i → 0 to (N-1) {
    if (hm.get(A[i]) == 1)
        return A[i]
}

```

$TC = O(N)$ $SC = O(N)$

Q → Find the count of distinct elements in A[].

$A = [3 \ 5 \ 6 \ 5 \ 6]$ $\{3, 5, 6\}$

Ans = 3

$A = [5 \ 5 \ 5]$ Ans = 1

```
for i → 0 to (N-1) {  
    hs.add(A[i])  
}
```

return hs.size()

TC = $O(N)$

SC = $O(N)$

Q → Given an integer array, check if there exist a subarray with 0 sum.

$A = [2 \ 2 \ 1 \ -3 \ 4 \ 3]$

Ans = true

Bruteforce → \forall subarray, calculate sum & check if sum = 0.

$$\frac{N * (N+1)}{2}$$

$$TC = O(N^3) \rightarrow O(N^2)$$

$$SC = O(1)$$

subarray sum → prefix sum

subarray sum l — r $\Rightarrow P[r] - P[l-1]$ // $l > 0$

$$P[r] - P[l-1] = 0$$

$$\text{if } (l == 0) \rightarrow P[r]$$

$$\Rightarrow \boxed{P[r] = P[l-1]} \checkmark$$

$$\boxed{P[r] = 0} \checkmark$$

$P[0] = A[0]$

$P[i] = A[0] + A[1] + \dots + A[i]$

if ($P[0] == 0$) return true

hs.add($P[0]$)

for $i \rightarrow 1$ to $(N-1)$ {

$P[i] = P[i-1] + A[i]$

if ($P[i] == 0$) return true

hs.add($P[i]$)

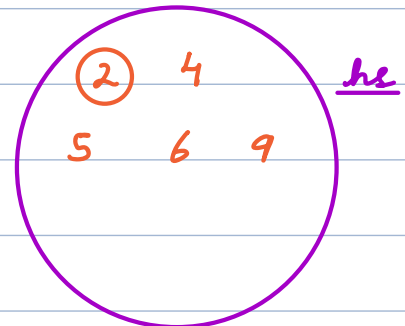
}

if (hs.size() == N) return false

else return true

TC = $O(N)$ SC = $O(N)$

$A = [2, 2, 1, -3, 4, 3]$
 $P = 2, 4, 5, 2, 6, 9$



contest — learner id → not attempted / score
Key Value
-1 int