

CS561_HW3_Boya_Zhou

Question1:

	Partition1	Partition2	Partition3
Query1	1	2	3
Query2	4	5	6
Query3	7	8	9

There is **m** machine.

- Query1, Partition1
 - Parallel group by city: Server i partitions chunks R_i using a hash function $(\text{Customer.city} \bmod P)$ on city (there is p city), get $R_{i1}, R_{i2}, \dots, R_{i,p-1}$.
 - Server i send partition R_{ij} to server j (assume j less than m)
 - Server j computes $\text{AVG}(\text{age})$ on $R_{0j}, R_{1j}, \dots, R_{p-1,j}$
 - Send the result from server j to other server
- Query1, Partition2:
 - Since the data is evenly distributed to all nodes, still need to exchange data.
 - The rest is as same as 1
- Query1, Partition3:
 - No need to do parallel group by
 - Calculate the $\text{avg}(\text{age})$ just on the local machine
- Query2, Partition1
 - Parallel duplicate elimination: Duplicate elimination the age on local machine on Customer.age since is a range-based partition, no need to exchange the data
 - Each sever deal its own data on its own server.
- Query2, Partition 2
 - Since all the records is evenly distributed on each server.
 - Each server partition chunks R_i using a hash function $(\text{Customer.city} \bmod P)$ on city, get $R_{i1}, R_{i2}, \dots, R_{i,p-1}$
 - Server j send partition R_{ij} to other server
 - Right now, each city is only in one server. Server j eliminate duplicate (age duplication) on $R_{0j}, R_{1j}, \dots, R_{p-1,j}$
- Query2, Partition3
 - Since it is based on city hash-based partition, each server perform duplication eliminate on its own server when considering age.
- Query3, Partition 1
 - Since the partition is range-based partition on age. Only the nodes contain age > 30 will work on data.
 - Then select city="Boston" or city="New York"
- Query3, Partition 2
 - Since the data is well distributed on all servers, each server will perform the selection and projection on its own server.
- Query3, Partition 3

- a. Only the server contain city=Boston and city = New York do selection on age > 30

Question2:

1. First question
 - Send only S.y to R's location
 - Select R.x1 on R
 - Do the join based on Y columns in R's location after selection
 - Send the records of R that will join (without duplicates) to S's location
 - Perform the final join in S's location
2. Second question
 - Select on S.z1
 - Send all the records in S fit the selection to R's location
 - Perform the final join based on column y in R's location.
3. Third question
 - We assume here are two relations : R(cusID, age, salary), S(cusID, phoneNum, gender), cusID is primary key, two relations have same number of primary key and the two relations are almost the same large, we assume there is C customers.
 - When using semi-join, the cost is S.cusID to R. The whole relation R to S.
 - When just send S to R, the cost is whole relation S to R.
 - Since the two relations are almost the same large, so the semi-join is less efficient here.
 - Beside this scenario, if column y is large, semi-join will be less effective, too.

Question 3

1. First question
 - The purpose of two-phase commit is coordinates all the processes that participate in a distributed atomic transaction on whether to commit or abort the transaction
 - Why it is widely is because The protocol achieves its goal even in many cases of temporary system failure.
2. Second question
 - Since the coordinator receive "ready T" from all sites, it will go to the Phase 2, after coordinator send "commit T" to all sites.
 - i. After just the site recover, it will frozen in ready state. It will re-contact the coordinator, if can not, it will contact other participants, if there is any site in abort of commit state, this site will do the same thing. Abort T or Commit T.
3. Third question
 - Right now when coordinator recover, it is in prepare state, it will either
 - i. Resend "prepare T" to other sites and wait for response.
 - ii. Put "abort T" in its own log, send global "abort T" to all participants.
 - iii. If after a long time, coordinator recover, maybe the participants already selected a new coordinator to replace it.