## COMP2270/6270 – Theory of Computation Sixth week

## School of Electrical Engineering & Computing The University of Newcastle

**Exercise 1**) Define the function  $twice(L) = \{w : \exists x \in L \ (x \text{ can be written as } c_1c_2 \dots c_n, \text{ for some } n \ge 1, \text{ where each } c_i \in \Sigma_L, \text{ and } w = c_1c_1c_2c_2 \dots c_nc_n\}$ .

- a) Let  $L = (1 \cup 0)*1$ . Write a regular expression for *twice(L)*.
- b) Are the regular languages closed under twice? Prove your answer.

Exercise 2) For each of the following claims, state whether it is *True* or *False*. Prove your answer.:

- a) The union of an infinite number of regular languages must be regular.
- b) The union of an infinite number of regular languages is never regular.
- c) If  $L_1$  and  $L_2$  are regular languages and  $L_1 \subseteq L \subseteq L_2$ , then L must be regular.
- d) The intersection of two nonregular languages must not be regular.
- e) The intersection of an infinite number of regular languages must be regular.
- f) If L is a language that is not regular, then  $L^*$  is not regular.
- g) If  $L^*$  is regular, then L is regular.
- h) Every subset of a regular language is regular.

**Exercise 3)** For each of the following languages L, state whether L is regular or not and prove your answer:

- a)  $\{w \in \{a, b, c\}^* : \text{ in each prefix } x \text{ of } w, \#_a(x) = \#_b(x) = \#_c(x)\}\}.$
- b)  $\{w \in \{a, b, c\}^* : \exists \text{ some prefix } x \text{ of } w (\#_a(x) = \#_b(x) = \#_c(x))\}.$
- c)  $\{w \in \{a, b, c\}^* : \exists \text{ some prefix } x \text{ of } w (x \neq \epsilon \text{ and } \#_a(x) = \#_b(x) = \#_c(x))\}.$

Exercise 4) Define the following two languages:

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L_a = \{ w \in \{ a, b \}^* : \text{ in each prefix } x \text{ of } w, \#_a(x) \ge \#_b(x) \}.

L_b = \{ w \in \{ a, b \}^* : \text{ in each prefix } x \text{ of } w, \#_b(x) \ge \#_a(x) \}.
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- a) Let  $L_1 = L_a \cap L_b$ . Is  $L_1$  regular? Prove your answer.
- b) Let  $L_2 = L_a \cup L_b$ . Is  $L_2$  regular? Prove your answer.

**Exercise 5**) For each of the following languages L, state whether L is regular or not and prove your answer:

- a)  $\{uww^{R}v: u, v, w \in \{a, b\}^{+}\}.$
- b)  $\{xyzy^{R}x : x, y, z \in \{a, b\}^{+}\}.$

**Exercise 6**) Let  $\Sigma = \{a, b\}$ . For the languages that are defined by each of the following grammars, do each of the following:

- *i*. List five strings that are in *L*.
- *ii*. List five strings that are not in *L*.
- iii. Describe L concisely. You can use regular expressions, expressions using variables (e.g.,  $a^nb^n$ , or set

theoretic expressions (e.g.,  $\{x: ...\}$ )

- iv. Indicate whether or not L is regular. Prove your answer.
- a)  $S \rightarrow aS \mid Sb \mid \epsilon$
- b)  $S \rightarrow aSa \mid bSb \mid a \mid b$

**Exercise 7**) Consider the following context free grammar *G*:

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S \rightarrow aSa

S \rightarrow T

S \rightarrow \varepsilon

T \rightarrow bT
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 $T \to cT$  $T \to \varepsilon$ 

One of these rules is redundant and could be removed without altering L(G). Which one?

**Exercise 8**) Show a context-free grammar for each of the following languages *L*:

- a) BalDelim =  $\{w : \text{ where } w \text{ is a string of delimeters: } (, ), [, ], \{, \}, \text{ that are properly balanced} \}.$
- b)  $\{a^ib^j: 2i \neq 3j+1\}.$
- c)  $\{w \in \{a, b\}^* : \#_a(w) = 2 \#_b(w)\}.$

## **REFERENCES**

[1] Elaine Rich, Automata Computatibility and Complexity: Theory and Applications, Pearson, Prentice Hall, 2008.