

Name: _____

StudentNo: _____

The University of Newcastle
School of Electrical Engineering and Computer Science

COMP3260/6360 Data Security

Midterm Test 1

20 March 2019

Test duration: 55 min

100 marks

In order to score marks, you must show all the workings!

STUDENT NUMBER: _____

STUDENT NAME: _____

PROGRAM ENROLLED: _____

<i>Question 1</i>	<i>Question 2</i>	<i>Question 3</i>	<i>Question 4</i>	<i>Question 5</i>	<i>TOTAL</i>

$$\lg 26! \approx 88.4$$

$$\lg 25! \approx 83.7$$

$$\lg 3 \approx 1.58$$

$$\lg 26 \approx 4.7$$

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1. (20 marks) Suppose that there are 5 possible messages, A, B, C, D and E, with probabilities $p(A) = p(B) = p(C) = p(D) = 1/8$ and $p(E) = 1/2$.
- What is the expected number of bits needed to encode these messages in optimal encoding?
 - Give an example of an optimal encoding.
 - Calculate the average number of bits needed to encode the message using your encoding.

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2. (20 marks) True or false?

- a. Every integer in the range $[1, 971]$ has a multiplicative inverse modulo 972.
- b. Every integer in the range $[0, 18]$ has a multiplicative inverse modulo 19.
- c. Every integer in the range $[1, 34]$ except 5 and 7 has a multiplicative inverse modulo 35.
- d. Equation $3x \bmod 15 = 12$ has no solutions.
- e. Computing in $GF(2^n)$ is less efficient than computing in $GF(p)$, as it is easier to work with integers than polynomials.
- f. There is an efficient algorithm for factoring large numbers, as to find factors of n , we only need to check if it is divisible by all prime numbers less than square root of n , thus the algorithm is sub-linear.
- g. There is an efficient algorithm for finding a greatest common divisor of any two integers.
- h. There is no efficient algorithm for fast exponentiation.
- i. 100 and 111 are multiplicative inverses in $GF(2^3)$ with irreducible polynomial $p(x) = x^3 + x^2 + 1$.
- j. 101 and 110 are additive inverses in $GF(2^3)$ with irreducible polynomial $p(x) = x^3 + x^2 + 1$.

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3. Explain the following terms.
- (a) (**8 marks**) Euler's Totient Function (also provide formula)
 - (b) (**6 marks**) Steganography (also give an example)
 - (c) (**6 marks**) Absolute Rate of Language

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4. (20 mark) Let $a=100$. If $GF(2^3)$ with irreducible polynomial $p(x)=x^3+x^2+1$, use Euler's theorem to find a^{-1} and then verify that $a \times a^{-1} \bmod p(x)=1$.

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5. (20 marks) Find a solution to the equation $7x \bmod 40 = 1$ in the following 3 ways. Note that you must show all the workings and/or trace the algorithm in order to score marks.

a) **Euler's Theorem** (by fast exponentiation): $a^{\phi(n)} \bmod n = 1$, where $\gcd(a,n)=1$

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b) Chinese Remainder Theorem: Let d_1, \dots, d_t be pairwise relatively prime, and let $n = d_1 \times d_2 \times \dots \times d_t$. Then the system of equations $(x \bmod d_i) = x_i$ ($i = 1, \dots, t$) has a common solution x in the range $[0, n-1]$. The common solution is

$$x = \sum_{i=1}^t \frac{n}{d_i} y_i x_i \bmod n$$

where y_i is a solution of $(n/d_i) y_i \bmod d_i = 1$, $i = 1, \dots, t$.

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c) Extended Euclid's algorithm:

```
Algorithm inv(a,n)
begin
   $g_0 := n; g_1 := a; u_0 = 1; v_0 := 0; u_1 := 0; v_1 := 1; i := 1;$ 
  while  $g_i \neq 0$  do “ $g_i = u_i \times n + v_i \times a$ ”
    begin
       $y := g_{i-1} \text{ div } g_i; g_{i+1} := g_{i-1} - y \times g_i;$ 
       $u_{i+1} := u_{i-1} - y \times u_i; v_{i+1} := v_{i-1} - y \times v_i;$ 
       $i := i + 1$ 
    end
  end
   $x := v_i - 1;$ 
  if  $x \geq 0$  then inv := x else inv := x+n
end
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