

2-0 @ A circuit switched network will be suitable for this Scenario since the application generates data for a long period of time where data is continuously generated. Since the data generation characteristics is known, network resources can be reserved in advance. In addition, the overhead cost of setting up and tearing down connections frequently will introduce additional processing load and delay.

Data generation pattern:

K time unit K time unit

(b) Assume $\sum_{i=1}^{X} N_i < R_{link}$ Where N_i is the data generated in every burst. $KTU \leq Isec$ $X = \frac{1}{KTU} = \frac{1}{2} = \frac{1}{KTU} = \frac{1}{2} = \frac{1}{1} = \frac{1}{2} = \frac{1}{1} = \frac$

Assum If the above condition exist i.e. total data generated by the application in a time unit is less than the transmission rate of the link, then no queue will build up hence, no congestion control is necessary.

2-1:0 First bit is transmitted at t=0, where the last bit will be transmitted at t = ditrans

| Host | Host | Last bit | L= dprop Last bit t = dtrans

At t = drans, the lost bit will leave the host A.

@ dprop < dtrans, at time t = dtrans

First bit has reached host B, host A is still transmitting
bits

(a) $S = 2.5 \times 10^8 \text{ m/sec}$, L = 120 bits, R = 56 kbps d = ?Given dprop = dtrans $dprop = \frac{d}{s} \Rightarrow d = S * dprop = S * \frac{L}{R}$ $= 2.5 \times 10^8 \times \frac{120}{56 \times 10^3} = 2.5 \times 10^8 \times 2.14 \times 10^{-3}$ = 535,000 meter= 535 km

2-2 a A circuit Awitched link provide a connection between a pair of users for the duration of a call, irrespective of active call time. In this case call active period is 20% of the total call time. Data is transmitted continuously @ 1 Mbps whereas the link data rate is 2 Mbps.

No. of circuit Awitched users is given by:

When two or less than two users transmit at the same time there will be no queuing delay $\Sigma R_{\nu} < R$. When three or more users transmit at the same time, i.e. $\Sigma R_{\nu} > R$, there will be queuing delay.

@ Since each user is transmitting for 20% of the call time, the probability a user transmit is 0.2



(d) To solve this problem we have to use the Bionomial distribution which is given by

$$Pr(x=k) = {n \choose k} P^{k} (1-b)^{n-k}$$

where n is the population size and k is no of certive users. p is the probability of transmission.

$$P_{r}(x=3) = \begin{pmatrix} 3 \\ 3 \end{pmatrix} p^{3}(1-p)^{3-3}$$

$$= \frac{3!}{3!(3-3)!} (0.2)^{3} (1-0.2)^{0}$$

$$= \frac{3!(3-3)!}{1 \times 0.008 \times 1} = 0.008$$

Prob. that all three transmitters will transmit simultaneously is 0.008

$$\frac{2-3}{2}$$
 (a) $R = 3$ Mbps, $R_u = 150$ kbps, $p = 0.1$

$$N_{circuit} = \frac{3 \times 10^6}{150 \times 10^3} = 20$$

Probability a given user is transmitting

Probability a given use
$$(1-k)^{K}$$

P(x=k) = $\binom{n}{k}$ $\binom{k}{k}$ $\binom{(1-k)^{K}}{(1-k)^{K}}$
 $n = 50$, $k = 5$, $k = 0.1$
 $p(x=5) = \binom{50}{5} (0.1)^{5} (1-0.1)^{50-5}$
 $= \frac{50!}{5! \cdot 45!} (0.1)^{5} (0.9)^{45}$
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$$= \frac{3.041 \times 10^{64}}{120 \times 1.196 \times 10^{56}} \times 1 \times 10^{-5} \times 8.727 \times 10^{-3}$$

$$\frac{2-4}{2}$$
 (a) $N_{\text{circuit}} = \frac{1 \times 10^9}{100 \times 10^3} = 10,000$

(b) Probability that more than N users out of M users are transmitting.

$$P_r\left(X \geqslant M\right) = \sum_{n=N+1}^{M} \binom{M}{n} p^n (1-p)^{M-n}$$

For n=5, and M=10, p=0.1

$$P_r(x > 5) = \sum_{n=6}^{10} {M \choose n} p^n (1-p)^{m-n}$$

$$\Rightarrow Pr(x \ge 5) = \frac{10!}{6!(10-6)!} (60.1)^6 (0.9)^6 + \frac{10!}{7!(10-7)!} (0.1)^7 (0.9)^{10-7}$$

$$+\frac{10!}{8! 2!0-8!!} (0.1)^{8} (0.9)^{10-8} + \frac{10!}{9!(10-9)!} (0.1)^{9} (0.9)^{10-9}$$

$$+\frac{10!}{10!(10-10)!} (0.1)^{10} (0.9)^{10-10}$$

$$= \frac{3628800}{720 \times 24} \times 1 \times 10^{-6} \times 0.6561 + \frac{3628800}{5040 \times 6} \times 1 \times 10^{-7} \times 0.729$$

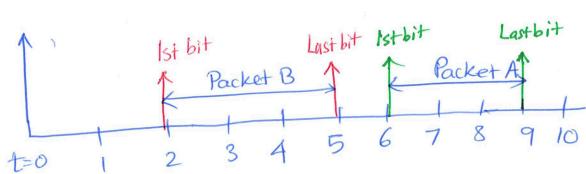
2-5: Host B to router A has lower propagation delay

$$\Rightarrow d_{P,B} = 2 \times 10^{-3} + \frac{1500 \times 8}{4 \times 10^{6}} = 2 \times 10^{-3} + 3 \times 10^{-3}$$

$$=5\times10^{-3}$$
 sec $=5$ ms

Host A packet arrival delay dp,A = 6×10-3+ 1500×8 (5) ⇒ dP,A = 9×10-3 sec = 9 ms

Timing diagram:



the above diagram we see that packet from host B will arrive ahead of packet A, hence, no queuing delay occur.

2-6:
$$1 \text{ KB} = 1024 \text{ byte in binary, } 1000 \text{ Byte (decimal)}$$

$$1 \text{ MB} = 1 \text{ KB} \times 1 \text{ KB} = 1,048576 \text{ byte in binary}$$

$$1 \text{ MB} = 106 \text{ Byte in decimal}$$

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1 GB = 1 KB × 1 MB = 1024 × 1048 576 = 1.07374 × 109 bytes IGB = 109 in decimal

the binary conversion

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dfile = dprop +
$$\frac{L}{R} = \frac{150 \times 10^3}{3 \times 10^8} + \frac{1.03374 \times 10^9 \times 8}{2 \times 10^6}$$

= $5 \times 10^{-4} + 4294.96 \approx 4295 \text{ sec}$ (Ignore Ist dezm)

= $71.58 \text{ min} = 1.193 \text{ hour}$

(b)
$$d_{file} = (d_{prop} + \frac{L}{R_1}) + d_{journey} + (d_{prop} + \frac{L}{R_2})$$

We ignore the propagation delays

$$\Rightarrow d_{file} = \frac{L}{R_1} + \frac{d}{S_b} + \frac{L}{R_2} \quad | \quad S_b \text{ speed of bas}$$

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=)
$$dfile = 85.89 + \frac{150 \times 10^{3} + 8.589}{16.66} + 8.589 = 9098.078ec$$

= $85.89 + 9003.6 + 8.589 = 9098.078ec$
= 2.527 hour

$$\frac{2-7!}{d} = 32 \text{ kbps}$$
, $T_p = 20 \text{ ms}$, $H = 18 \text{ byte}$
 $d = 10 \text{ km}$, $R = 10 \text{ Mbps}$

Voice packets are generated every 20 ms. No. of bits accumulated in a packet is given by Npacket = Roader X Tp

= 32×103 × 20×10-3 = 640 bits = 80 bytes

Voice packet:

$$P_{packet} = P_{prop} + \frac{L}{R} = \frac{10 \times 10^3}{3 \times 10^8} + \frac{(80 + 18) \times 8}{10 \times 10^6}$$