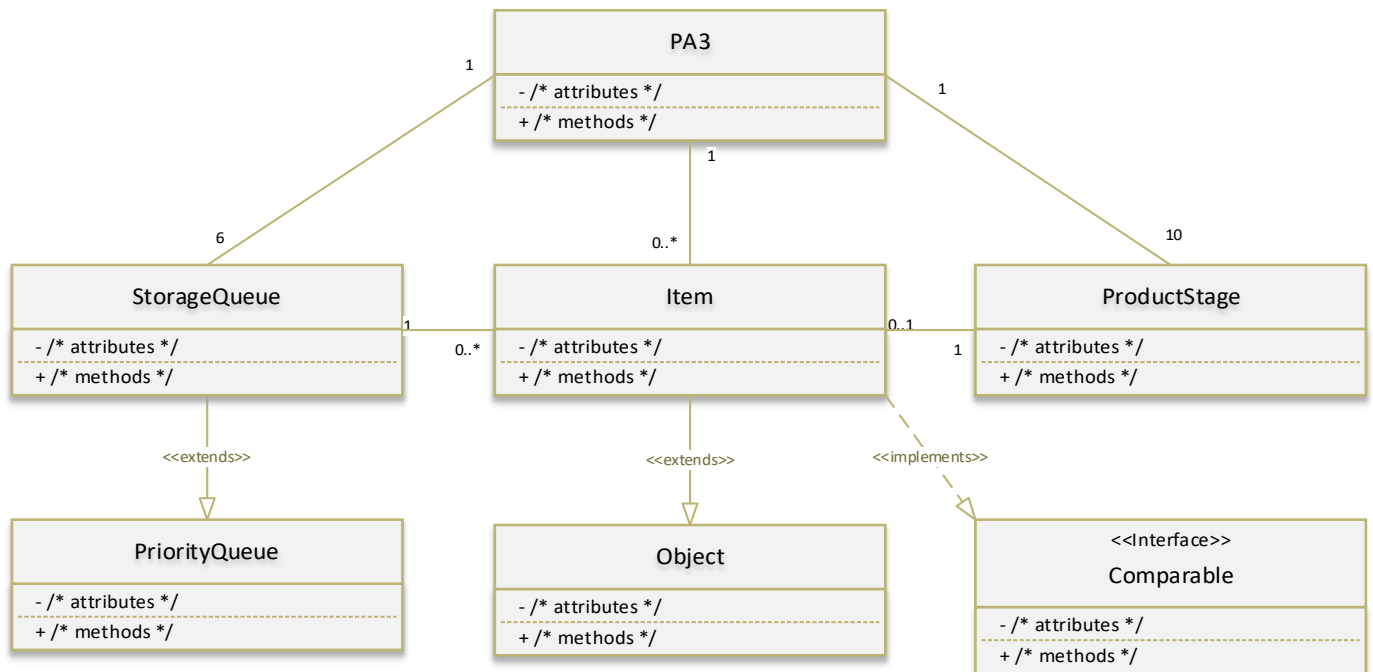

SENG2200 ASSIGNMENT 3 REPORT

1. UML diagram for Assignment 3:



2. The use of Inheritance/Polymorphism

Inheritance played a crucial role in this assessment, despite the relatively small number of classes used. The `StorageQueue` class extended the standard Java `PriorityQueue`, and the `Item` class extended the generic `Object` class while implementing the `Comparable` interface, greatly reducing the overall workload. These two classes allow the developer to very easily create, manipulate and queue a newly-defined object.

In the final product the `StorageQueue` class is extremely small – just one attribute and one method were added. The only thing that the super class `PriorityQueue` lacked was a name, which was important especially for the output.

This deviated from the initial outline where the `StorageQueue` class was used for `Starve` and `Block` checks as well. It quickly became apparent that those checks were better used in the `ProductStage` class.

I made the decision to extend `Object` for the `Item` class in the initial outline. `Object` has a lot of useful generic functionality that allows it to synergise well with the `PriorityQueue` class, and the `StorageQueue` class by extension. Like the `StorageQueue` class, the `Item` class adds attributes and methods needed for the assignment output – the more practical functionality is handled by the `Object` class.

The `Item` class did not implement the `Comparable` interface in the initial outline. An unforeseen side-effect of extending `PriorityQueue` and `Object` appeared to be that the interface was mandatory to prevent compiler-time errors, despite adding no comparing functionality. This effectively means that, although the `Item` class implements the interface, `Comparable` goes unused in the program.

3. Updating functionality for different production line topology

Several changes need to be made to the program to allow for a different path. Most significantly, the stages and queues are hardcoded. This is unavoidable – the layout of the stages and queues does not follow any obvious pattern and attempting to define the production line procedurally does not add any more flexibility. In the program the queues are defined in class `PA3`, lines 16-18, and the stages are defined in class `PA3`, lines 15 & 21-46. These need to be updated to change the production line.

Next is the unique IDs. Currently there are 2 variables that control IDs: `idNumberA` and `idNumberB`, found in class `PA3`, lines 50-51. `idNumberA` creates unique IDs for items created in stage 0a and `idNumberB` for those made in stage 0b. This is fine for any program that has 2 potential starting stages, but suppose it had 50. An array of unique ID seeds would be

much more useful in that situation. On the other hand, `idNumberB` is useless for a production line with only 1 starting stage.

Only a few small changes to the conditionals in the do-while loop in class [PA3](#), lines [53-146](#) are really necessary. Besides that, the `StarveTimeByStage` method in class [PA3](#) only allows for production lines with splits of up to 2 stages. This could be easily updated by using a local integer to track how many stages are in a “group”.

There are solutions to most of these problems that can be solved in a few lines of code each. The only problem with no simple procedural solution is the first one – unless the stages of the production line follow a predictable pattern (e.g. `0a/0b -> 1 -> 2a/2b -> ... -> na/nb`), they must be hardcoded with each change.

One point not raised in the first paragraph is that, since the inter-storage queues link one “group” of stages to another (for example, queue `Q01` might link `0a/0b/0c` to `1a/1b/1c/1d`), queues can be defined procedurally, unlike stages.

4. Updating functionality for more complex production lines, beyond the “straight-line” used in the program

The program would not be designed significantly differently if these changes were made. Each item has a full record of the stages it went through from start to finish. Suppose the end product is a plane if the item starts at `R0` and a car if it starts at `T1`. This information is in the attribute `stages`, making it simple to express through discrete-event simulation.

A few questions arise though. Will all completed items go to the same storage or is there a separate storage for each type of item? For the former, there are no changes. If the latter, how many possible different items can be created? Using the example production line, there are likely 3 different types of item, 1 for each origin point (`R0`, `S0`, `T1`). Although this can be hard-coded with an array of 3 `ArrayList`s, a better solution would be an `ArrayList` of `ArrayList`s – check the prefix of the `stages` attribute for items as they are completed, create a new `ArrayList` whenever an item with a different origin is detected and push the item to the appropriate `ArrayList` based on the prefix. Without additional information, this assumes a lot about how the production line is structured and what the assessed requirements of the program are.

The main change that will need to be made to the program seems small but is deceptively problematic – each stage may have more than 1 previous inter-storage queue. A major part of the Starving functionality is checking the size of the previous queue. In the current program this is simple; there is only 1 previous queue so only that 1 queue needs to be checked. Applying FCFS would be best, to give all potential previous queues a fair share.

The *prevQueue* attribute in the ProductStage class would need to be changed to an array or ListArray, as well as a method and attributes used to determine the next queue.