#### School of Electrical Engineering and Computing

# SENG2200/6220 PROGRAMMING LANGUAGES & PARADIGMS (S1, 2020)

## Functional Programming I

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### **Outline**

- Functional Languages
  - Features
  - History
- Scheme
  - Syntax
  - Lists
  - Predefined Functions



- Procedural Languages (e.g., Pascal, C)
  - Statement oriented
  - Built from code segments (blocks)
  - Uses variables and assignment
  - Control flow (selection and iteration) is by control structures
  - Forces sequential execution
- The main concern is efficient of execution on Von Neumann architectures



Procedural Example (Factorial)

```
if n == 0 then return 1
if n > 0 then
  prod = 1
  while n > 0 do
    prod = prod*n
    n = n-1
  return prod
```



- Functional Languages (e.g., Lisp, Scheme)
  - Function oriented
  - Control flow is by means of function calls
  - Iteration is by means of recursion
    - there are no looping constructs
  - Functions may be evaluated concurrently
- The main concern is representation of the problem
  - Tend not to be efficient on traditional architectures



Functional Example (Factorial)

```
\begin{cases}
n == 0 \mid 1 \\
\\
\text{Factl(n)} \leftarrow
\end{cases}

\begin{cases}
n > 0 \mid n * \text{Factl(n-1)}
\end{cases}
```



- Imperative Languages
  - Operators are applied to values and the results store in variables
  - These variables can be used as a source of values at any time in their life-cycle
- Functional Languages
  - Functions are applied to arguments
  - The result of a function may (will!) be an argument of another function
  - There is no concept of variable!



#### **Pure Function**

- A pure function is one with no side-effects
  - A purely functional language consists only of pure functions
  - No user input/output other than as arguments to/results from the main function
  - Not very useful (other than for scientific calculation)
  - Most practical functional languages are not pure
- If a pure function is called twice with the same arguments, the result is the same both times
  - This is called referential transparency



## **Brief History**

- LISP = LISt Processing language (1959)
- Al research needed a language that:
  - Process data in lists (rather than arrays)
  - Symbolic computation (rather than numeric)
- Syntax is based on lambda calculus
  - First interpreter just to demonstrate the usefulness of lambda notation



## **Brief History**

- Pure LISP
  - Purely functional
  - Only supported two data types: atoms and lists
- Scheme (1975)
  - A "teaching" version of LISP
  - Statically scoped
  - First-class functions



#### Scheme

- Scheme programmes are written using fully parenthesised prefix notation
  - A function in Scheme is a list
  - whose first element is the function name.
  - and whose remaining elements are the arguments
  - E.g., (+12) and (max 58)
- Scheme is (usually) an interpreted language
- Coding
  - You could use the Racket implementation from your own computer. <a href="https://download.racket-lang.org/">https://download.racket-lang.org/</a>
  - Or use online compilers, e.g.,

https://repl.it/repls/CoarseSaddlebrownComputergame



## Example

```
> ( + 4 3 2 1 )
  10
> ( * 4 3 2 1 )
  24
> ( gcd 14 21 35 )
  7
> ( quit )
```



#### **Scheme Atoms**

- Atomic Types...
- Identifiers/Symbols
  - Can start with any character except #"();0123456789
  - Can contain any characters except ( )
- Boolean constants
  - #t and #f
- Character constants
  - #\ followed by any single character.
  - *E.g.*, #\a, #\b, #\?
  - #\space and #\newline



#### **Scheme Atoms**

- String constants
  - Any sequence of characters except \ or " in "strings"
  - Escaped characters \\ and \"
- Numeric constants
  - A sequence of digits, possibly containing a decimal point . or exponent mark e or E
  - Possibly preceded by #b (binary), #o (octal), #d
     (decimal), #x (hexadecimal), #e (exact) or #i (inexact)



#### Scheme Atoms

With the exception of strings and character constants...

Scheme is case InSeNsitive (in Scheme standard R5RS)

But you should not rely on this!

Scheme is ALSO case sensitive (in Scheme standard R6RS etc)

e.g., Racket implementation



Atoms can be combined into a list

```
'( 1 2 3 4 5 )
'( red green blue )
'( "green" "eggs" "ham" )
'( 2112 "budd" #t )
```

Lists can be nested

```
'(12(34)5)
'(1(2(3(4))))
'(1(23)(4)5))
```

■ The **empty list ()** is a constant

```
'(1()(23)45)
```



- To create a list:
  - Use single-quote (indeed **quote** function): '(1 2)
  - Use list function: (list 1 2)
- By default, Scheme will attempt to evaluate any list as a function call

```
( func arg1 arg2 ... )
```

- Each arg may itself be a function call
- func may be a variable or function call (see the LAMBDA function later)
- There is no guarantee on the order of evaluation
- Arguments may even be evaluated in parallel!



The quote function suppresses evaluation of its single argument, for example

```
> ( + 2 3 )
5
> ( quote (+ 2 3) )
  (+ 2 3)
```

■ '( + 2 3 ) is shorthand for (**quote**(+ 2 3))



- The quasiquote function acts like quote
  - Shorthand `( + 2 3 )
- But can be unquoted

```
> `( 1 ( unquote ( + 2 3 ) ) )
  (1 5)
```

■ Shorthand `( 1 ,( + 2 3 ) )

```
> `( 1 ( unquote-splicing '(+ 2 3) ) )
  (1 + 2 3)
```

■ Shorthand `( 1 ,@'( + 2 3 ) )



#### **Scheme Pairs**

- A pair is written as two expressions separated by a dot
  - E.g., (1 . 2)
- A pair can be created by using cons function

```
> (cons 1 2)
(1 . 2)
```

- The first element of a pair is the car
  - E.g., (car (cons 1 2) )= 1
- The second element of a pair is the cdr
  - E.g., (cdr (cons 1 2) )= 2



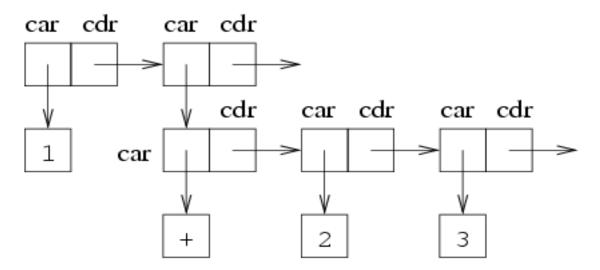
#### **Pairs and Lists**

- A list is actually a sequence of pairs
- A list can be written/represented by using nested pairs
  - (1 2 3) can be considered as(1 . (2 . (3 . () ) ) )
  - The *cdr* of the last element is the empty list ( )
  - If the last element of the list is not an empty list then it is an improper list, e.g.,



#### **Pairs and Lists**

- Considering the pairs of a list
  - car is a "pointer" to the value of the element
  - cdr is a "pointer" to the next element of the list
- For example, the list (1 + 2 3) looks like





- ( **define** var expr )
  - Evaluates expr and binds the result to var
  - This can only be used at the top-level of code
  - The scope of *var* is the remainder of the program
  - If var already exists, then it is overwritten
  - Effectively creates and intialises a global variable
- ( set! var expr )
  - Evaluates expr and stores it in an already defined var



- Functions in Scheme are first-class types
  - Functions can be assigned to variables
  - Functions can be returned by other functions
- ( lambda ( var1 ... varn ) body )
  - A function that expects *n* parameters, with parameters var1 ... varn
  - body is a sequence of one or more expressions
  - Calling the function returns the value of the last expression in body



- Scheme checks that the number of arguments matches the number of formal parameters
- ( **lambda** var body )
  - A function which expects zero-or-more arguments
  - The list of arguments is bound to var
- ( lambda ( var1 ... varn . varn+1 ) body )
  - A function which expects *n*-or-more arguments
  - The first *n* arguments are bound to *var1* ... *varn*
  - The list of remaining arguments is bound to varn+1



```
( define ( func var1 ... varn ) body )
    ⇔ (define func (
        lambda ( var1 ... varn ) body )
(define ( func var1 ... varn . varn+1 ) body )
    ⇔ (define func (
        lambda ( var1 ... varn . varn+1 ) body )
```



## Example

Swap function

```
( define ( swap a b )
  ( define temp a)
  ( set! a b )
  ( set! b temp )
  (list a b)
)
> ( swap 5 6 )
  (6 5)
```



#### The let Function

- ( let ( ( varl exprl ) ... ( varn exprn )) body )
  - Evaluates each expr and binds the result to the corresponding var
  - The scope of each *var* is the expression *body*
  - These variables may mask variables declared at higher levels
  - There is no guarantee on the order in which variables are bound



#### The let\* Function

- ( let\* ( ( varl exprl )...( varn exprn ) ) body )
  - Equivalent to let, except...
  - Expressions are evaluated and variables are bound in order, var1 through varn
  - The scope of each var is the body and all vars and exprs after it in the order



#### The letrec Function

- ( letrec ( ( varl exprl ) ... ( varn exprn ) ) body )
  - Equivalent to let, except...
  - The scope of each var is the body and all vars and exprs
  - Each expr is evaluated knowing that the vars exist, but not knowing their values
  - Used to create mutually-recursive functions



## **Examples**

```
> x
    ERROR: unbound variable: x
> ( define x 1 )
> ( let ( ( y 2 ) ( z x ) ) `( ,x ,y ,z ) )
    (1 2 1)
> ( let ( ( y 2 ) ( z y ) ) `( ,x ,y ,z ) )
    ERROR: unbound variable: y
```



## **Examples**

```
> ( let* ( ( y 2 ) ( z y ) ) `( ,x ,y ,z ) )
  (1 \ 2 \ 2)
> ( letrec (
    ( A ( lambda (x)
      (if (> x 0) (B (- x 1)) 0)
    ( B ( lambda (y)
      (if (> y 0) (A (- y 1)) ))
  ) `( ,( A 4 ) ,( A 5 ) ,( B 4 ) ,( B 5 ) )
(0 \ 1 \ 1 \ 0)
```



- ( pair? expr )
  - Is true if the expression results in a pair
- ( car pair )
  - Returns the car of the pair
- ( **cdr** pair )
  - Returns the cdr of the pair
- It is an error to use car or cdr on the empty list



- ( cons car cdr )
  - Creates a pair with the given car and cdr
- ( set-car! pair expr )
- ( set-cdr! pair expr )
  - Change the car/cdr of a pair



- ( **list** item1 ... itemn )
  - Creates a list of the given items
- ( list? expr )
  - Is true if the expression results in a list
- ( null? expr )
  - Is true if the expression results in the empty list
- ( length list )
  - Returns the length of the list



- ( append list1 ... listn )
  - Concatenates the lists into a single list
  - Newly allocates items for list1 ... listn-1 but not listn
- ( reverse list )
  - Returns the reversed list
- ( list-tail list k )
  - Returns the result of omitting the first kelements of the list
  - Same as applying cdr k times
- ( list-ref list k )
  - Returns the kth element of the list
  - Indices starting from 0
  - lacktriangle Same as (  $m{car}$  (  $m{list}$   $m{tall}$   $m{list}$   $m{k}$  ) )



- ( begin expr1 ... exprn )
  - Evaluates the expressions in order and returns the result of exprn
- ( if test then else )
  - Evaluates test
  - If test is true, then it evaluates and returns then
  - Otherwise, it evaluates and returns else
  - Only #£ is false, everything else is true



- ( **not** expr )
  - The logical not of expr
- ( and expr1 ... exprn )
  - Evaluates each expr in order until one fails, then it returns false
  - If no exprs fail then it returns true
- ( **or** expr1 ... exprn )
  - Evaluates each expr in order until one is true, then it returns true
  - If no exprs are true then it returns false



- Evaluates *tests* in order until one is true
- Evaluates the corresponding *exprs*, returning the last
- If no tests are true then evaluate the exprs of the (optional) else clause, returning the last



### Example

Factorial function

```
( define ( factorial n )
    ( cond
      ( ( = n 0 ) 1 )
      ( ( > n 0 )
        ( * n ( factorial ( - n 1 ) ) )
      ( ( < n 0 ) 'error )
```



# Example

#### Output

```
> ( factorial 0 )
  1
> ( factorial 4 )
  24
> ( factorial -1 )
  error
```



### Example

Factorial function

```
(define (compare x y)
  (cond
    ( (> x y)
      (display "x is greater than y") )
    ( (< x y)
      (display "y is greater than x") )
    (else (display "x and y are equal")
```



- ( apply func args )
  - Apply function func to the list of args and return the result
- ( map func list1 ... listn )
  - Apply function func to each list and return the list of results
  - Guarantees func will be called on the lists in order
  - Does not guarantee the order of evaluation within each list



- Other control structures (not covered)
- for-each
  - Like map guarantees argument order but does not return results
- case
  - Like a C/Java switch
- do
  - Like a C/Java for-loop better support for multivariables loops



- Scheme has 3 versions of <u>equivalence</u>...
- ( **eq?** obj1 obj2 )
  - True if obj1 and obj2 are equal simple constants or references to the same object
- ( **eqv**? *obj1 obj2* )
  - True if obj1 and obj2 are equal simple values or references to the same object
- ( equal? obj1 obj2 )
  - Recursively applies eqv? to the elements of complex types obj1 and obj2



## Examples

```
> ( equal? 1 2 )
  #f
> ( equal? (list 1) (list 2) )
  #f
> ( equal? (list 1) (list 1) )
  #t
> ( eqv? (list 1) (list 1) )
  #f
> ( eq? (list 1) (list 1))
  #f
> ( eq? '() '())
  #t
```



- Scheme has 3 versions of <u>membership</u> predicates
- ( memq obj list )
  - Returns the first sublist of list whose car eq? obj
  - If no match is found then return #£
- Similarly memv using eqv?
- Similarly member using equal?



### Examples

```
> ( memq 1 '(1) )
  (1)
> (memq 1 '(2 1 4))
  (1 \ 4)
> ( memv '(1) '(2 1 4) )
  #f
> ( memv '(1) '(2 (1) 4) )
  #f
> ( member '(1) '(2 1 4) )
  #f
> ( member '(1) '(2 (1) 4) )
  ((1) 4)
> ( memv 1 '(2 (1) 4) )
  #f
```



Type checking predicates...

```
( number? obj )
( complex? obj ) ( real? obj )
( rational? obj ) ( integer? obj )
( exact? num ) ( inexact? num )
( zero? num )
( positive? num ) ( negative? num )
( odd? num ) ( even? num )
```



Ordering predicates

```
( = num1 num2 num3 ... )
( < num1 num2 num3 ... )
( > num1 num2 num3 ... )
( <= num1 num2 num3 ... )
( >= num1 num2 num3 ... )
```



- String comparison
  - Case sensitive

```
( string=? str1 str2 )
```

Case insensitive

```
( string-ci=? str1 str2 )
```

Also < > <= >=



### **Scheme Numbers**

Useful functions on lists of numbers

```
( + num1 num2 num3 ... )
( * num1 num2 num3 ... )
( - num1 num2 num3 ... )
( / num1 num2 num3 ... )
```

- Left associative
- Also max min gcd lcm
- Remainders, rounding, trigonometry, exponentials, rectangular- and polar-coordinates



```
current-input-port
 ( current-output-port )
      Return the current input/output ports
      Default to standard input/output
• ( open-input-file filename )
  ( open-output-file filename )
      Return a port opened for input/output on the given file
  ( close-input-file port )
 ( close-output-file port )
      Close an open input/output port
```



- ( read port )
  - Returns the object whose external representation is found on *port*
- ( read-char port )
  - Removes the next character on port and returns it
- ( peek-char port )
  - Returns the next character on port without removing it
- ( eof-object? obj )
  - Return true if obj is the end-of-file object
- port defaults to current-input-port



- ( write obj port )
  - Writes a machine representation of obj to port
- ( display obj port )
  - Writes a machine representation of obj to port
  - Strings have no double-quotes or escaped characters, and characters appears as per write-char
- ( write-char char port )
  - Writes a single character to port
- ( newline port )
  - Writes a newline character to port
- port defaults to current-output-port



 Functions for converting between strings and numbers

```
( number->string number )
( number->string number radix )
( string->number string )
( string->number string radix )
```

Radix must be an exact integer. By default, radix=10.
 E.g.,

```
>( number->string 5)
   "5"
>( number->string 5 2)
   "101"
```



#### References

- R. W. Sebesta, "Concepts of Programming Languages", 9th Edn, Addison-Wesley, 2010 (Chapter 15) (also Edn.10)
- R. K. Dybvig, "The Scheme Programming Language", 3rd Edition, MIT Press, 2003. <a href="http://www.scheme.com/tspl3/">http://www.scheme.com/tspl3/</a>