

WQ7.

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$$T(n) = \begin{cases} 0 & n=1 \\ n-1+T(n-1) & n>1 \end{cases}$$

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$$T(n) = n-1 + T(n-1) \quad ①$$

$$= n-1 + (n-2 + T(n-2))$$

$$= (n-1) + (n-2) + T(n-2) \quad ②$$

$$= (n-1) + (n-2) + (n-3) + T(n-3) \quad ③$$

⋮

$$i\text{-th row} = (n-1) + (n-2) + \dots + (n-i) + T(n-i) \quad \textcircled{④} \quad i\text{-th row.}$$

$$= i \cdot n - (1+2+3+\dots+i) + T(n-i)$$

$$= i \cdot n - \frac{(1+i)i}{2} + T(n-i)$$

$$\text{row } r \quad \text{we conclude} \quad = r \cdot n - \frac{(1+r)r}{2} + T(n-r) \quad | \quad n-r=1 \quad r=n-1 \quad \textcircled{⑤}$$

$$= (n-1)n - \frac{(1+n-1)(n-1)}{2} + 0$$

$$= n(n-1) - \frac{n(n-1)}{2} = n(n-1)(1-\frac{1}{2}) = \frac{n(n-1)}{2}$$

$$\text{Guess: } T(n) = \frac{n(n-1)}{2}$$

proof by induction:

$$\text{Base case: } n=1 \quad T(1) = \frac{1 \times 0}{2} = 0 \rightarrow \text{true.}$$

$$\text{Inductive step: Assume that } T(k) = \frac{k(k-1)}{2}$$

By recurrence relation:

$$T(k+1) = k + T(k) = k + \frac{k(k-1)}{2}$$

$$= \frac{2k + k^2 - k}{2} = \frac{k^2 + k}{2} = \frac{k(k+1)}{2}$$

$$= \frac{(k+1)(k+1-1)}{2}$$

Thus, it also holds for $k+1$

$$\text{Therefore, } T(n) = \frac{n(n-1)}{2} \text{ holds for any } n \geq 1$$

$$\text{Namely, } T(n) = \frac{n(n-1)}{2} \in O(n^2)$$

