

Program Complexity

Big-O Notation

CSCI 3700 — Data Structures and Objects

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Outline

- 1 Timing Programs
 - Bad Timing Measures
 - Counting Statements
- 2 Big-O Notation
 - Big-O Introduction

Bad Performance Measures

"Wall-clock" time

Sensitive to system load

CPU cycles

Varies from machine to machine

Machine instruction count

Varies by language and programming style

Counting Statements

Counting to ten

- Line 4:
 $1 + 11 + 10 = 22$
- Line 5: 10
- Line 7: 1
- Total: 33

```
1  int ten(int data[], int n) {  
2      int j;  
3  
4      for (j=0; j<10; j++)  
5          cout << j << endl;  
6  
7      return 0;  
8  }
```

Counting Statements

Finding the largest element

- Line 4: 1
- Line 5: $2n$
- Line 6: $n - 1$
- Line 7: $0 \dots n - 1$
- Line 9: 1
- Total: $3n + 1 \dots 4n$

```
1  int findMax(int data[], int n) {  
2      int j, m;  
3  
4      m = 0;  
5      for (j=1; j<n; j++)  
6          if (data[j] > data[m])  
7              m = j;  
8  
9      return m;  
10 }
```

Counting Statements

Sorting a list

- Line 4: $2n$
- Line 5: $n - 1$
- Line 6: $2 \sum_{i=1}^{n-1} (i + 1)$
- Line 7: $\sum_{i=1}^{n-1} i$
- Line 8: $0 \dots \frac{n^2 - n}{2}$
- Lines 10–12:
 $3(n - 1)$ total
- Total:
 $1.5n^2 + 6.5n - 6 \dots$
 $2n^2 + 6n - 6$

```
1 void sort(int data[], int n) {  
2     int s, j, m, tmp;  
3  
4     for (s=n-1; s>0; s--) {  
5         m = 0;  
6         for (j=1; j<=s; j++)  
7             if (data[j] > data[m])  
8                 m = j;  
9  
10        tmp = data[m];  
11        data[m] = data[s];  
12        data[s] = tmp;  
13    }  
14 }
```

Why Bother?

- Counting statements is hard
 - Imagine counting a large function
- Counting statements is more accurate
 - Machine and language effects minimized
 - Better counting methods exist
- Additional benefits
 - Compare functions / programs

A Better Approach

- Big- O notation
 - Easily approximates count
 - Shows trends with large data sizes
 - Allows code / algorithm comparison
- How it works
 - Concept derived from discrete mathematics
 - Simple form needed here — focus on loops

Types Of Loops

Simple loops

- Linear loop
 - Key: ++ or --
 - $O(n)$
- Logarithmic loop
 - Key: cut data in half
 - $O(\lg n)$
 - $\lg n = \log_2 n$

Nested loops

- Log-linear loop
 - Linear inside logarithmic (or vice versa)
 - $O(n \lg n)$
- Dependent quadratic loop
 - Linear inside linear
 - Inner counter depends upon outer counter
 - $O(n^2)$
- Quadratic loop
 - Linear inside linear
 - No dependency between loops
 - $O(n^2)$

Benefits and Caveats

Benefits

- Easy to determine
- Allows code / algorithm comparison
- Show trend as amount of data increases

Caveats

- Not accurate for small data size
- Not exact
- Can't distinguish two programs / algorithms with same $O(\cdot)$ value

Comparing Big-O Values

- Table shows common values
 - Fastest at top
 - Each value significantly faster than values beneath it
 - Two items, same big- O value
 - Neither significantly faster
- $O(1)$
 - $O(\log n)$
 - $O(n)$
 - $O(n \log n) = O(\lg(n!))$
 - $O(n^2)$
 - $O(n^3)$
 - $O(n^4)$
 - $O(2^n)$
 - $O(3^n)$
 - $O(n!)$
 - $O(n^n)$

Summary

- Many bad timing measures
- Counting statements better, but hard
- Big- O notation provides easy approximation