

# Exam Revision

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# Introduction

- ▶ Date and Time: 13th May 2020, 9am to 6pm.
- ▶ Place: Brightspace
- ▶ Type: Open Book Exam
- ▶ Contents:
  - ▶ Cryptographic Tools
  - ▶ User Authentication
  - ▶ Access Control
  - ▶ Malicious Software
  - ▶ Distributed Denial of Service Attacks
  - ▶ Intrusion Detection
  - ▶ Firewalls and Intrusion Prevention
- ▶ **NO PLAGIARISM!**

# Cryptographic Tools

- ▶ Confidentiality with symmetric encryption
- ▶ Message authentication and hash functions
- ▶ Random and pseudorandom numbers
- ▶ Public-key encryption
- ▶ Digital signatures and key management

# User Authentication

- ▶ Digital user authentication principles
- ▶ Password-based authentication
- ▶ Remote user authentication
- ▶ Security issues for user authentication

# Access Control

- ▶ Subjects, objects, and access rights
- ▶ Discretionary access control
- ▶ UNIX file access control
- ▶ Role-based access control
- ▶ Attribute-based access control
- ▶ Identity, credential, and access management

# Malicious Software

- ▶ Types of malicious software (malware)
- ▶ Advanced persistent threat
- ▶ Propagation-vulnerability exploit-worms
- ▶ Payload-stealth-backdoors
- ▶ Propagation-social engineering-span E-mail
- ▶ Payload-system corruption
- ▶ Countermeasures

# Distributed Denial of Service Attacks

- ▶ Distributed denial-of-service attacks
- ▶ Application-based bandwidth attacks
- ▶ Reflector and amplifier attacks●
- ▶ Denial-of-service attacks
- ▶ Flooding attacks
- ▶ Responding to a denial-of-service attack

# Intrusion Detection

- ▶ Intruders
- ▶ Intrusion detection
- ▶ Analysis approaches
- ▶ Distributed or hybrid intrusion detection
- ▶ Honeypot



# Firewalls and Intrusion Prevention

- ▶ The need for firewalls
- ▶ Firewall characteristics and access policy
- ▶ Types of firewalls
- ▶ Firewall basing
- ▶ Intrusion prevention systems

# Example Exam Question

## 8.4 Consider the following login protocol.

user knows password P  
user knows Hash function  $H(.)$  and has a mobile calculator  
user gives login name N to machine  
machine generates random number R  
machine gives R to user  
user computes  $X := \text{Hash}(P) \text{ XOR } \text{Hash}(R)$   
user gives X to machine  
machine uses N to obtain P from password table  
machine computes  $Y := \text{Hash}(P) \text{ XOR } \text{Hash}(R)$   
if  $X=Y$  then machine allows login

- a. Explain what is wrong with it and how can it be broken.
- b. Show a simple way to strengthen this protocol against your attack.

## 8.4 a. Adversary

- Sees the messages: N, R, X
- Computes  $\text{Hash}(R)$
- Computes  $X \text{ XOR } \text{Hash}(R) = \text{Hash}(P)$

Later on:

- Adversary Requests Login, submits N.
- Machine generates random number  $R'$ .
- Adversary computes  $\text{Hash}(R')$ .
- Adversary computes  $y = \text{Hash}(P) \text{ XOR } \text{Hash}(R')$ .
- Adversary submits y, and logs in as user.

- b. To strengthen, simply require the protocol to compute  $\text{Hash}(R \text{ XOR } P)$  instead of  $\text{Hash}(R) \text{ XOR } \text{Hash}(P)$