



Names, Scopes, and Bindings

ORGANIZATION OF PROGRAMMING LANGUAGES
JUCHEOL MOON

Name, Scope, and Binding

- A name is exactly what you think it is
 - Most names are identifiers
- A binding is an association between two things, such as a name and the thing it names
- The scope of a binding is the part of the program (textually) in which the binding is active

Binding Time

- | | |
|--|-----------------------------------|
| ▪ language design time
<u>3</u> | 1. static data in memory |
| ▪ language implementation time
<u>4</u> | 2. function (method) names |
| ▪ program writing time
<u>2</u> | 3. possible data type |
| ▪ compile time
<u>1</u> | 4. maximum size of heap and stack |

Binding Time

- | | |
|-------------------------|---------------------------------------|
| ▪ link time
<u>3</u> | 1. physical addresses to instructions |
| ▪ load time
<u>1</u> | 2. dynamic data in memory |
| ▪ run time
<u>2</u> | 3. object in another module (library) |
- The terms static and dynamic are generally used to refer to things bound before run time and at run time, respectively

Binding Time

- In general, early binding times are associated with greater (efficiency / flexibility)
- Later binding times are associated with greater (efficiency / flexibility)
- Compiled languages tend to have (early / later) binding times
- Interpreted languages tend to have (early / later) binding times

Lifetime

- Key events
 - creation of objects
 - creation of bindings
 - references to variables (which use bindings)
 - (temporary) deactivation of bindings
 - reactivation of bindings
 - destruction of bindings
 - destruction of objects

Lifetime

- The period of time from creation to destruction is called the **LIFETIME** of a binding
- If object outlives binding it's garbage
- If binding outlives object it's a dangling reference
- The textual region of the program in which the binding is *active* is its scope

```
void garbage() {
    int* a = new int(11);
    cout << "a << endl;
    int b = 22;
    a = &b;
    cout << "a << endl;
    delete a;
}

int* dangling() {
    int* c = new int(33);
    return c;
}
```

Diagram illustrating memory management:

- Variable `a` points to a memory location containing `11`. This location is marked as **out of scope** (garbage).
- Variable `b` points to a memory location containing `22`. This location is marked as **in scope**.

Storage Management

- Storage allocation mechanisms
 - Static
 - Stack
 - Heap
- Static allocation for
 - arguments
 - local variables
 - temporaries

Recursive Subroutine

- Factorial Function
- ```
int fact (int n)
{
 if (n == 1) return 1;
 else return n * fact(n - 1);
}
```

## Recursive Subroutine

\$a0=3 and the instruction "0x341000 jal 0x424000"

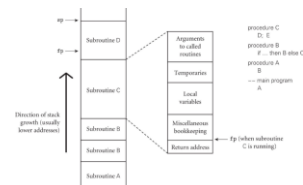
| PC       | Code                    |
|----------|-------------------------|
| 0x424000 | addi \$sp, \$sp, -8     |
| 0x424004 | sw \$ra, 4(\$sp)        |
| 0x424008 | sw \$a0, 0(\$sp)        |
| 0x424012 | lbp \$a0, \$1, 0x424032 |
| 0x424020 | addi \$v0, \$zero, 1    |
| 0x424024 | addi \$sp, \$sp, 8      |
| 0x424028 | jr \$ra                 |
| 0x424032 | addi \$a0, \$a0, -1     |
| 0x424036 | (jal) 0x424000          |
| 0x424040 | lw \$a0, 0(\$sp)        |
| 0x424044 | lw \$ra, 4(\$sp)        |
| 0x424048 | addi \$sp, \$sp, 8      |
| 0x424052 | mul \$v0, \$a0, \$v0    |
| 0x424056 | jr \$ra                 |

Handwritten notes and stack diagram:

- `$ra = 0x341004` (initial value)
- `$a0 = 0x424040` (initial value)
- Stack grows downwards (increasing address):
  - Address 0x424040: 1
  - Address 0x424044: 2
  - Address 0x424048: 3
  - Address 0x424052: 24004
- `$v0 = 6` (result of multiplication)
- Stack pointer `$sp` points to the top of the stack (0x424040).

## Storage Management

- Contents of a stack frame
  - Assigned at compile time
  - arguments
  - local vars.
  - temporaries
  - return addr.



## Storage Management

- Maintenance of stack is responsibility of calling subroutines
  - How can we (compilers) write time efficient and space efficient assembly codes?
    - ↑ using stack
    - ↓ using stack

## Leaf Procedure Example

### •MIPS code:

```
leaf_example:
 addi $sp, $sp, -4
 sw $s0, 0($sp)
 add $t0, $a0, $a1
 add $t1, $a2, $a3
 sub $s0, $t0, $t1
 add $v0, $s0, $zero
 lw $s0, 0($sp)
 addi $sp, $sp, 4
 jr $ra
```

•C code:

```
int leaf_example (int g, h, i, j) {
 int f;
 f = (g + h) - (i + j);
 return f;
}
```

•Arguments g, ..., j in \$a0, ..., \$a3  
 •f in \$s0 (hence, need to save \$s0 on stack)  
 •Result in \$v0

## Name Declarations

### •Some constructs must first be introduced by explicit declarations

▪`int javaNum;`

### •Some constructs can be introduced by implicit declarations

▪`pyNum = someNum + 10`

## Name Scope

### •Once a name is declared, how long is that declaration valid?

- file ?
- program ?
- function ?

### •C/C++ scoping example

- Declarations are valid in block that they are declared
- Declarations not in a block are global, unless the `static` keywords is used, in which case the declaration is valid in that file only

## C++ example

```
#include <iostream>
using namespace std;
string x = "Global";

int main()
{
 {
 string x = "Block 1";
 cout << x << endl;
 }
 {
 string x = "Block 2";
 cout << x << endl;
 }
}
```

<< `string x = "what?";`

## Scope Rules

- A scope is a program section of maximal size in which no bindings change, or at least in which no re-declaration are permitted.
- In most languages with subroutines, we OPEN a new scope on subroutine entry:
  - create bindings for new local variables,
  - deactivate bindings for global variables that are re-declared (these variable are said to have a "hole" in their scope)

## Scope Rules

### •On subroutine exit:

- destroy bindings for local variables
- reactivate bindings for global variables that were deactivated