Assignment 3 – Bo Henderson and Burde Prerana Kamath

Question 1

The doubly linked list was implemented with four functions. The FIFO operations which comprises of enqueue and dequeue was implemented followed by a priority enqueue function and a display function.

The structure consists of process id and the key value. The priority enqueue function was implemented using the key value as the priority. It also consisted of next and previous pointers that were local to every node and three global pointers that was current, front and rear.

Xinu uses arrays as the base structure whereas the program uses structures as the base structure. Since XINU uses relative pointers and has an implicit data structure such as the queuetab array, it takes less memory space and processor time. Since our implementation consists of structures using pointers there are some downfalls such as:

- 1 Pointers always take up extra memory space.
- 2 Since arrays have a better cache locality, they perform in a shorter time than our implementation and hence, their performance and memory requirements is better.
- 3 Linked lists cannot be accessed randomly. As seen in the implementation, they need to be accessed using the front pointer always which doesn't hold true for arrays since they have already been assigned a particular amount of memory space since its static memory allocation and can be accessed through their indices.

Pseudocode of the program:

The structure of the program is of the type:-

```
struct node
{
       int pid; //This is the process id
       int key; //This is the key value according to which priority is assigned
       //This is to maintain the connection among nodes. Every node will know the previous
node and the node after it.
       struct node *next;
       struct node *prev;
};
struct node *current = NULL; //This is to know where the current position of the queue is
present.
struct node *rear = NULL; //This is to know the last node in the queue.
struct node *front = NULL; //This is to know the first node in the queue.
The FIFO functions implemented were enqueue and dequeue.
Pseudocode for enqueue:
1. Algorithm – FIFO Enqueue
2. Input – Process ID, Key Value
3. Output – A node has been entered in the queue
4. *** pid – processid, key – keyvalue
5. ***link – new link to enqueue
6. struct node *link ← (struct node*) malloc (sizeof(struct node))
7. if front = NULL
                                   //If the queue is empty
       front ← rear ← link
8.
9. endif
10. else
                                   //Pointers of the rear should point to the new link
11.
       rear ← next ← link
```

```
12. link ← prev ← rear13. end else
```

14. rear ← link //The new addition is the rear node now

Pseudocode for dequeue:

- 1. Algorithm FIFO Dequeue
- 2. Input Process ID, Key Value
- 3. Output A node has been deleted from the queue
- 4. *** pid processid, key keyvalue
- 5. ***temp the node to dequeue
- 6. ***tempo temporary variable
- 7. struct node* temp
- 8. if front = NULL //The queue is empty
- 9. temp = NULL
- 10. end if
- 11. else if front -> next = NULL //The queue has only one node
- 12. struct node *tempo
- 13. $temp \leftarrow tempo$
- 14. front ← NULL
- 15. rear ← NULL
- 16. end elseif
- 17. else //When there is more than one node in the queue
- 18. struct node* tempo ← front
- 19. temp ← tempo
- 20. front ← front -> next
- 21. front -> prev ← NULL
- 22. end else
- 23. return temp //To print the information of the deleted node

Pseudocode for priority enqueue:

```
1. Algorithm - Priority Enqueue
2. Input – Process ID, Key Value
3. Output – A node has been entered in the queue
4. *** pid – processid, key – keyvalue
5. ***link – new link to enqueue
6. struct node *link ← (struct node*) malloc (sizeof(struct node))
7. if front = NULL
                                    //If the queue is empty
       front ← rear ← link
8.
9. endif
10. else
11.
       if link -> key > find -> key
                                    //The node has most priority
12.
              front -> prev ← link
13.
              link -> next ← front
14.
              front ← link
15.
       end if
16.
       else
17.
              while
                                    //Find the position per priority
                      if find -> next ← NULL
18.
19.
                             break
20.
                      endif
21.
                      find ← find -> next
22.
              end while
              if find -> next ← NULL and link -> key < find -> key
23.
                      rear -> next ← link
24.
25.
                      link -> prev ← rear
                      rear ← link
26.
```

27.

endif

```
28. else
```

29. find -> next -> prev \leftarrow link

30. $link \rightarrow prev \leftarrow find \rightarrow prev$

31. link -> next ← find

32. find -> prev \leftarrow link

33. end else

34. end else

35. end else

Question 2:

In XINU, a valid queue id is defined in the queue.h file.

Valid queue id's are:

- It must be positive

It must lie between 0 and NQENT – 2 (both inclusive)

NQENT = NPROC + 4 + NSEM + NSEM

NPROC – Number of processes.

NSEM – Number of semaphores.

These are hardcoded values in XINU.

The code changes were made in the getitem.c file present in the system folder of xinu.

The changes were made using the function isbadqid(q) where isbadqid was a function defined in the queue.h file present in the include folder of the Xinu. The functions getfirst and getlast had the check included.

An appropriate check for getitem was made using the pid where the function isbadpid(pid) was also defined in the include folder queue.h.

These checks were included in the getitem.c

A short function was added in each of these functions.

```
if(isbadqid(q))
{
printf(" The queue id is not valid");
return SYSERR;
}
```