

# 1.8 Ben Pugh

```
1. for(i=0; i < n)
    for(k=0; k < n)
        for(j=0; j < n)
```

So each time the j loop runs it does n.  
 Each time k runs it does n (n of i)  $\leq n^2$   
 Each time i runs it does n (not k)  $= n(n^2)$   
 Overall  $T(n) = O(n^3)$   $= \frac{n^3}{4^3}$

2. (same as above but analyze)

```
for(i=0; i < n)
    for(k=i; k < n)
        for(j=k; j < n)
```

Starting from the outer most we have n iterations. moving to our middle loop we have:  $0, 1, 2, \dots, n-1 = \frac{n(n-1)}{2}$  total. Moving inner loop we once again have  $0, 1, 2, \dots, n-1$  for a total of  $\frac{n(n-1)}{2} \cdot \frac{n-1}{2}$   $\frac{n^3}{4}$  this

to the loop is  $O(n^3)$

```
for(i < n)
    for(k=i+1; k < n)
        for(j=k+2; j < n, k++)
```

Assuming this code is correct, the code is an infinite loop because the innermost loop will never resolve.

Assuming you meant j++ in the middle, then it would be  $\frac{n(n-1)}{2} \cdot \frac{n-2}{2}$  (instead of k++)

The outermost runs n times. The middle runs n-1 but you can think of the overall formula like a sum. Then the innermost runs similar to middle but it is at most n-3 when k=1 & i=0. As k increases the formula still holds so overall  $O(n^3)$