



Cloud Infrastructure Week#5

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January 2024

AGENDA – WEEK#5

- Traditional Network
 - Today's Network
 - OSI Reference Model & Encapsulation
 - Switches & L2 problems & ARP and types
 - TCP/IP & Transport Layer
 - IPv4 Addressing
 - VLANs and Trunks
 - Redundancy & HA (High-Availability) Solutions
 - Routing Protocols



Today's Network

NETWORK & DEVICES / COMPONENTS

Communication is almost as important to us as our reliance on air, water, food, and shelter. In today's world, with networks, we are connected like never before.

END DEVICES:

An end device is where a message originates from or where it is received. Data originates with an end device, flows through the network, and arrives at an end device.

INTERMEDIARY NETWORK DEVICES:

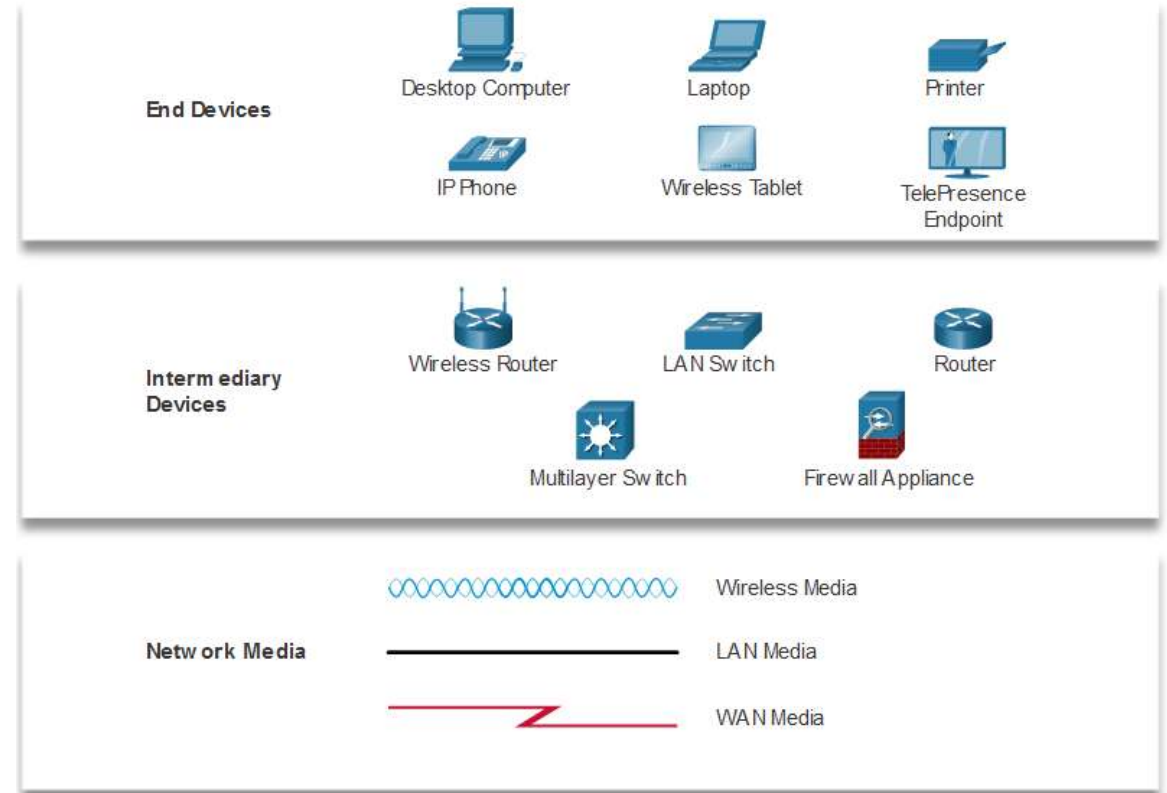
An intermediary device interconnects end devices. Examples include switches, wireless access points, routers, and firewalls.

Management of data as it flows through a network is also the role of an intermediary device, including:

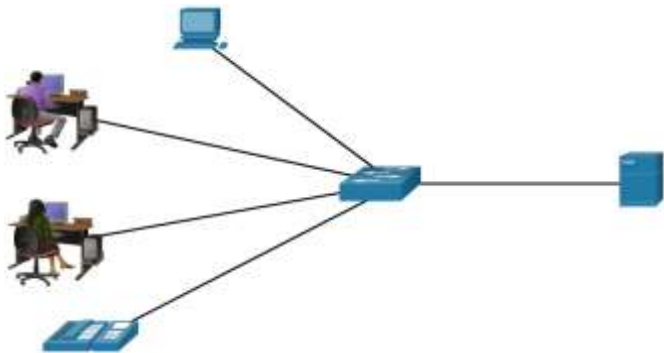
- Regenerate and retransmit data signals.
- Maintain information about what pathways exist in the network.
- Notify other devices of errors and communication failures.

NETWORK MEDIA:

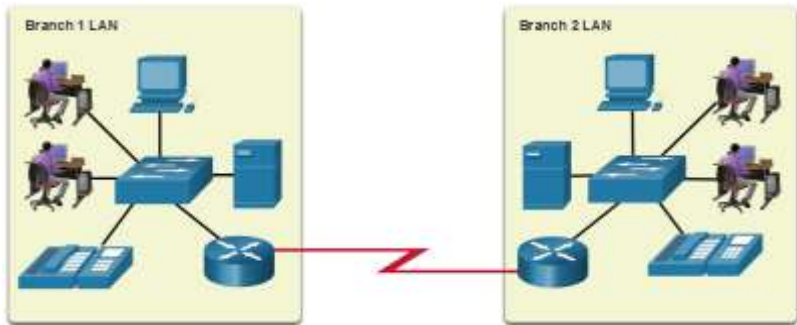
Communication across a network is carried through a medium which allows a message to travel from source to destination.



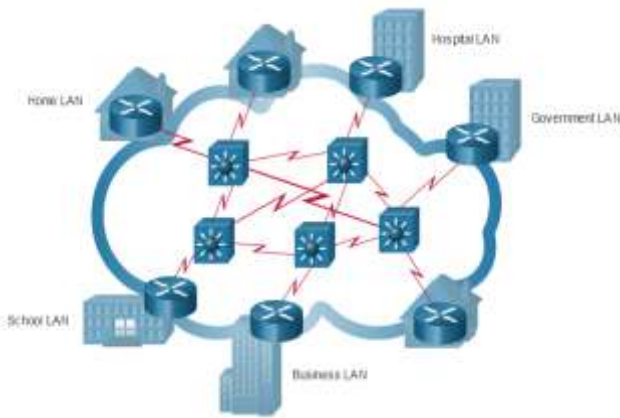
A **LAN** is a network infrastructure that spans a small geographical area.



A **WAN** is a network infrastructure that spans a wide geographical area.



The **internet** is a worldwide collection of interconnected LANs and WANs.



| LAN | WAN |
|--|--|
| Interconnect end devices in a limited area. | Interconnect LANs over wide geographical areas. |
| Administered by a single organization or individual. | Typically administered by one or more service providers. |
| Provide high-speed bandwidth to internal devices. | Typically provide slower speed links between LANs. |

- **LANs** are connected to each other using WANs.
- **WANs** may use copper wires, fiber optic cables, and wireless transmissions.

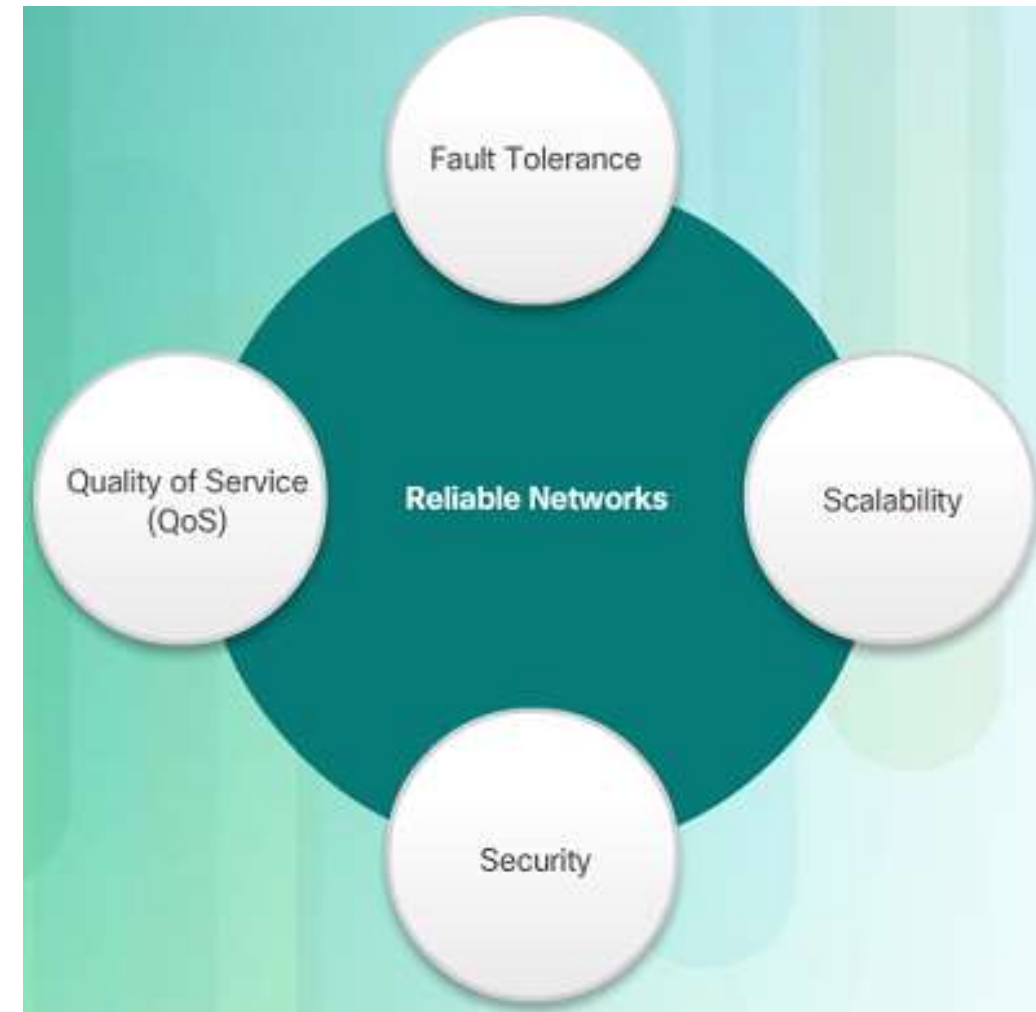
The internet is not owned by any individual or group. The following groups were developed to help maintain structure on the internet:

- IETF
- ICANN
- IAB

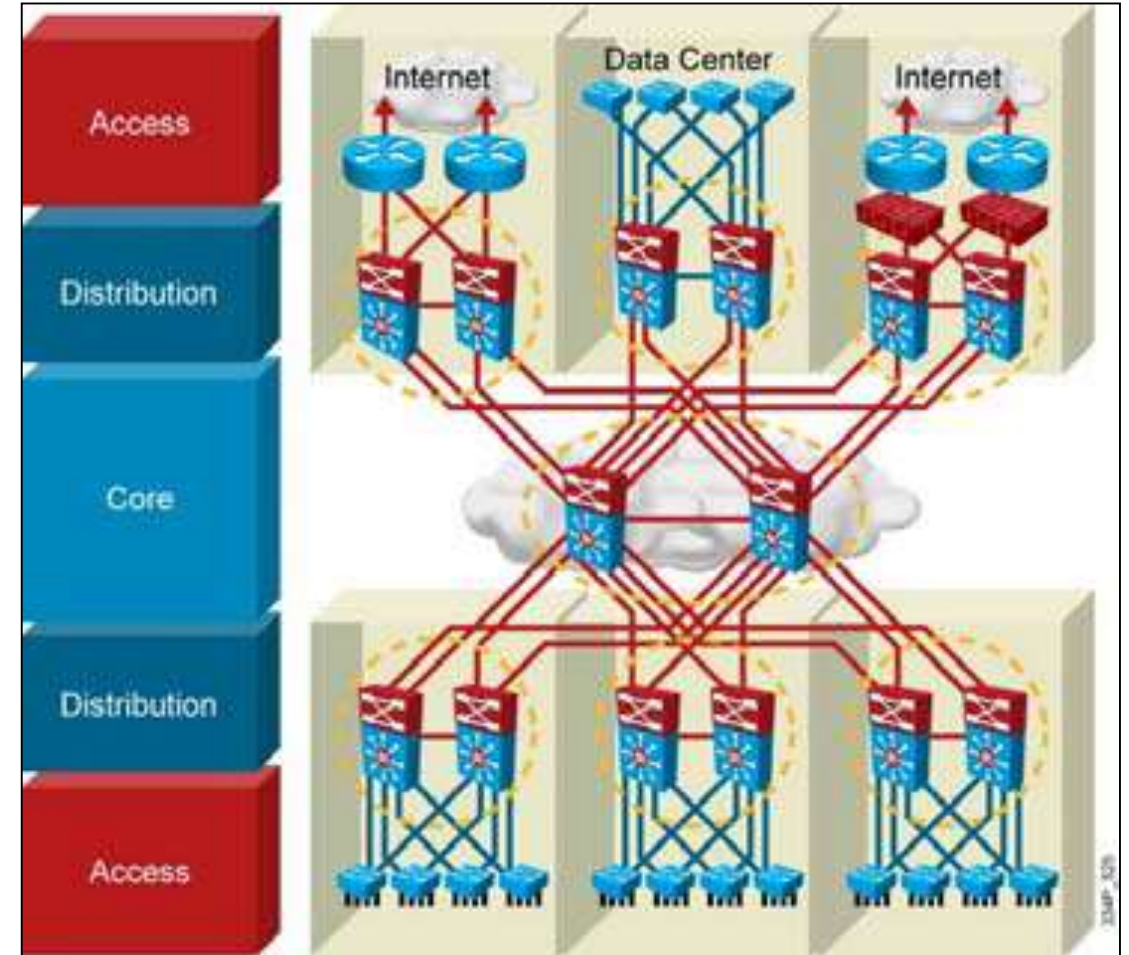
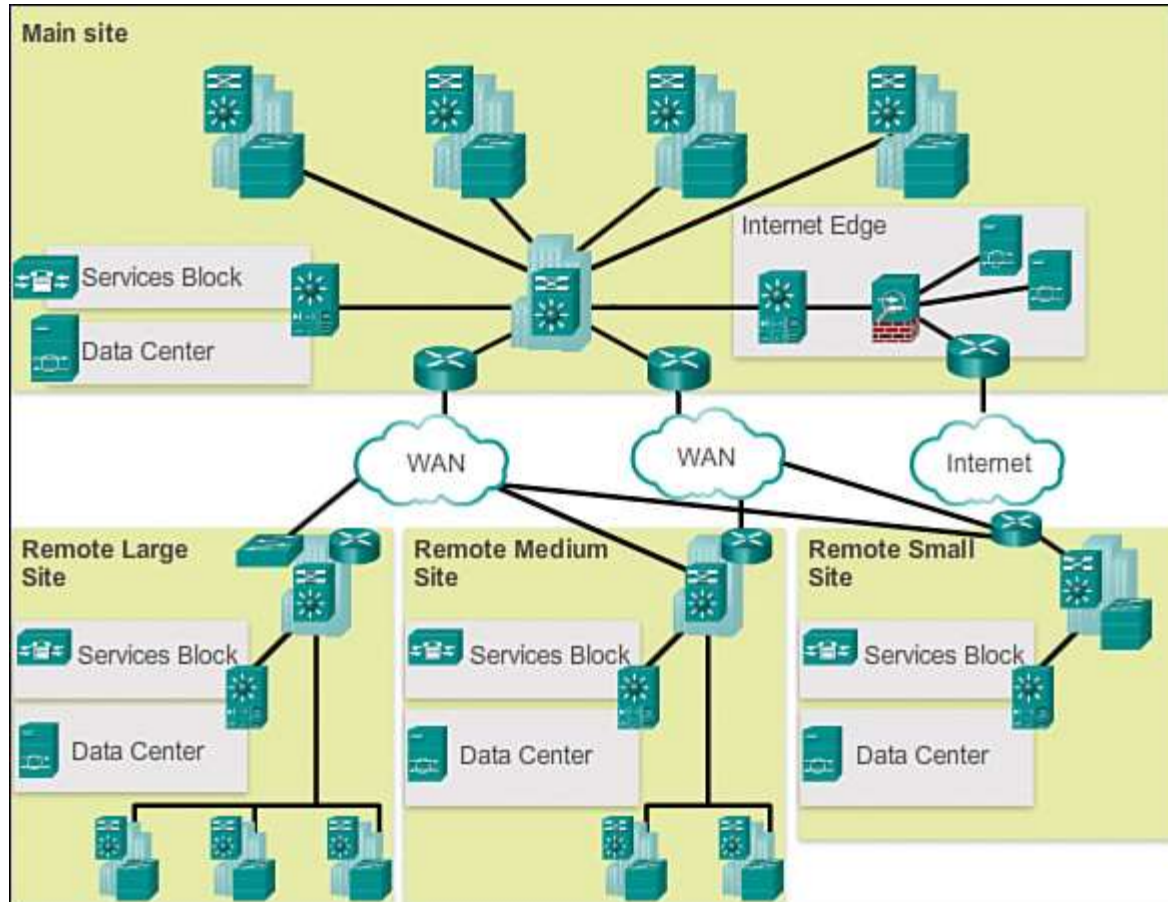
- What is the demarcation point?
- What is intranet, extranet and internet?

NETWORK ARCHITECTURE: COMMON TERMS

- **Fault Tolerance:** A fault tolerant network limits the impact of a failure by limiting the number of affected devices. Redundancy is used for Fault-tolerant networks.
 - What is a single point of failure and what are the redundancy types?
- **Scalability:** A scalable network can expand quickly and easily to support new users and applications without impacting the performance of services to existing users. Similarly, it needs to be shrunk when needed.
- **QoS:** Quality of Service (QoS) is the primary mechanism used to ensure reliable delivery of content for all users.
 - What are the packets in the network and how they are prioritized with QoS?
- **Network Security:** Huge concept, ideally, there are common terms in Network to secure the device physically and the data with encryption / network security devices.
 - What is encryption?
 - What are network security devices?

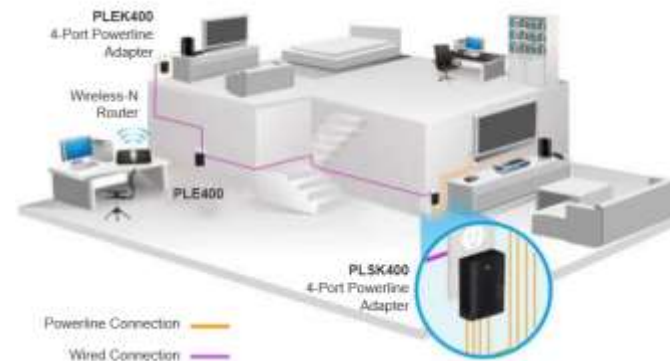


NETWORK ARCHITECTURE: HIERARCHICAL NETWORK MODEL



NETWORK TRENDS

- BYOD (Bring-Your-Own-Device)
- Online Collaboration
- Video Communication
- Cloud Computing
- Smart Home Technologies
- Powerline Networking
- Wireless Broadband Service

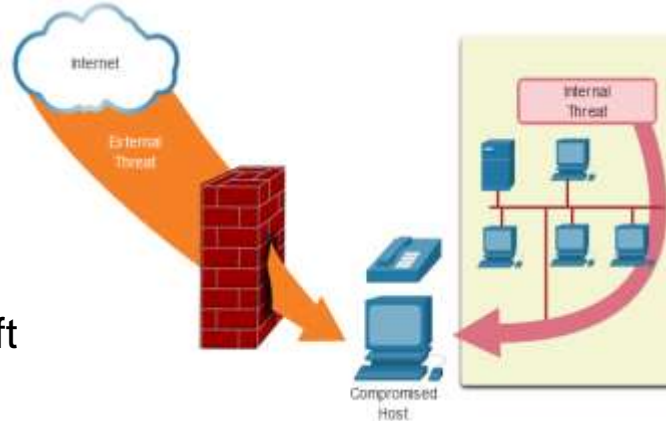


SECURITY TRENDS

Threats:

External Threats:

- Viruses, worms, and Trojan horses
- Spyware and adware
- Zero-day attacks
- Threat Actor attacks
- Denial of service attacks
- Data interception and theft
- Identity theft

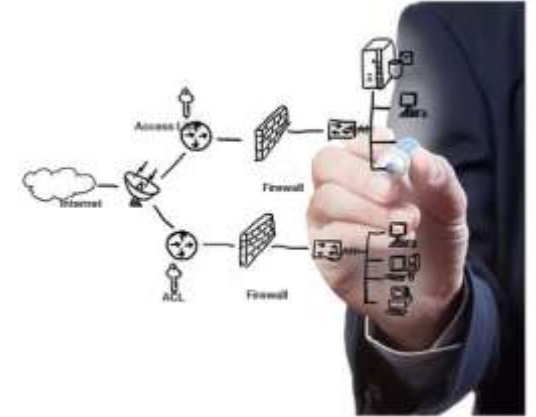


Internal Threats:

- lost or stolen devices
- accidental misuse by employees
- malicious employees

Solutions:

- Defense Strategies
 - Physical Security
 - Device Security
 - Network Security
 - Access Control
 - Permission Control
 - Encryption (Data)
- Antivirus / Antimalware / Antispyware
- Patching / Upgrades
- Firewall / ACLs (Access Control Lists)
- IDS & IPS Devices
- Proxies (DDoS Protection)
- VPNs (Virtual Private Networks)



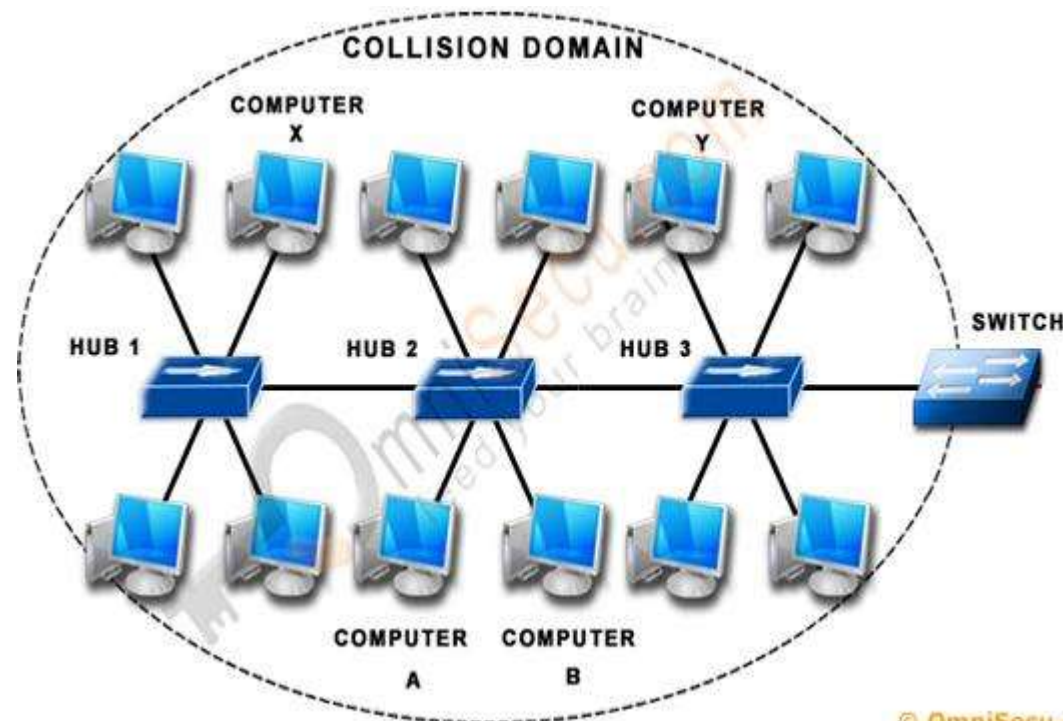
OSI Reference Model

COLLISION AND BROADCAST DOMAIN

A collision domain is a section of a network where data packets can collide with one another when being sent on a shared medium or through repeaters, in particular, when using early versions of Ethernet.

- *A network collision occurs when more than one device attempts to send a packet on a network segment at the same time.*

A broadcast domain is a logical division of a computer network, in which all nodes can reach each other by broadcast at the data link layer. A broadcast domain can be within the same LAN segment, or it can be bridged to other LAN segments.



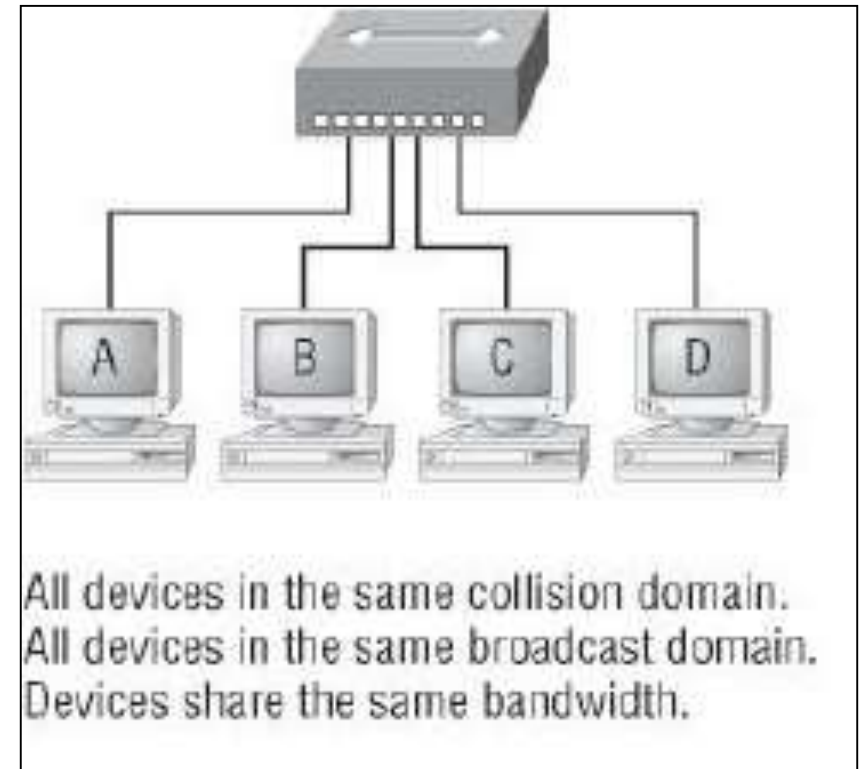
OSI LAYER-1: PHYSICAL

Jobs: Sends and receive **bits**, provides connectivity

Devices: Hub, Repeater, Cable, Connectors (RJ-45, DB-9, etc)

The term Bandwidth is the capacity at which a medium can carry data.

| Unit of Bandwidth | Abbreviation | Equivalence |
|---------------------|--------------|--|
| Bits per second | bps | 1 bps = fundamental unit of bandwidth |
| Kilobits per second | Kbps | 1 Kbps = 1,000 bps = 10^3 bps |
| Megabits per second | Mbps | 1 Mbps = 1,000,000 bps = 10^6 bps |
| Gigabits per second | Gbps | 1 Gbps = 1,000,000,000 bps = 10^9 bps |
| Terabits per second | Tbps | 1 Tbps = 1,000,000,000,000 bps = 10^{12} bps |



OSI LAYER-2: DATA LINK

Jobs:

Provides unique MAC address

Take bits and encapsulate a **frame**

Devices: Switch, Bridge

Protocols:

LAN: IEEE's 802.2, 802.3, 802.5; ANSI's FDDI

WAN: ATM, FR (Frame-Relay), PPP, HDLC, SDLC, X.25

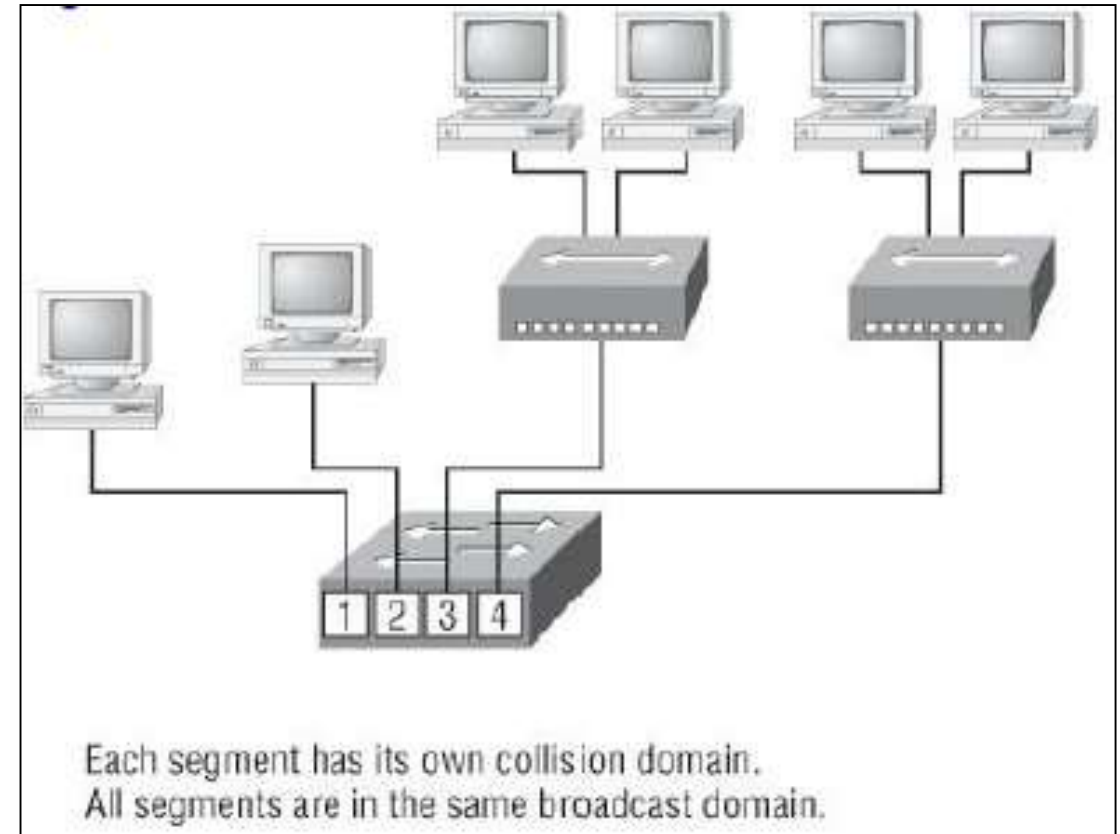
Address Types:

Broadcast: Every device on a segment

Multicast: A group of devices on a segment

Unicast: A single device on a segment

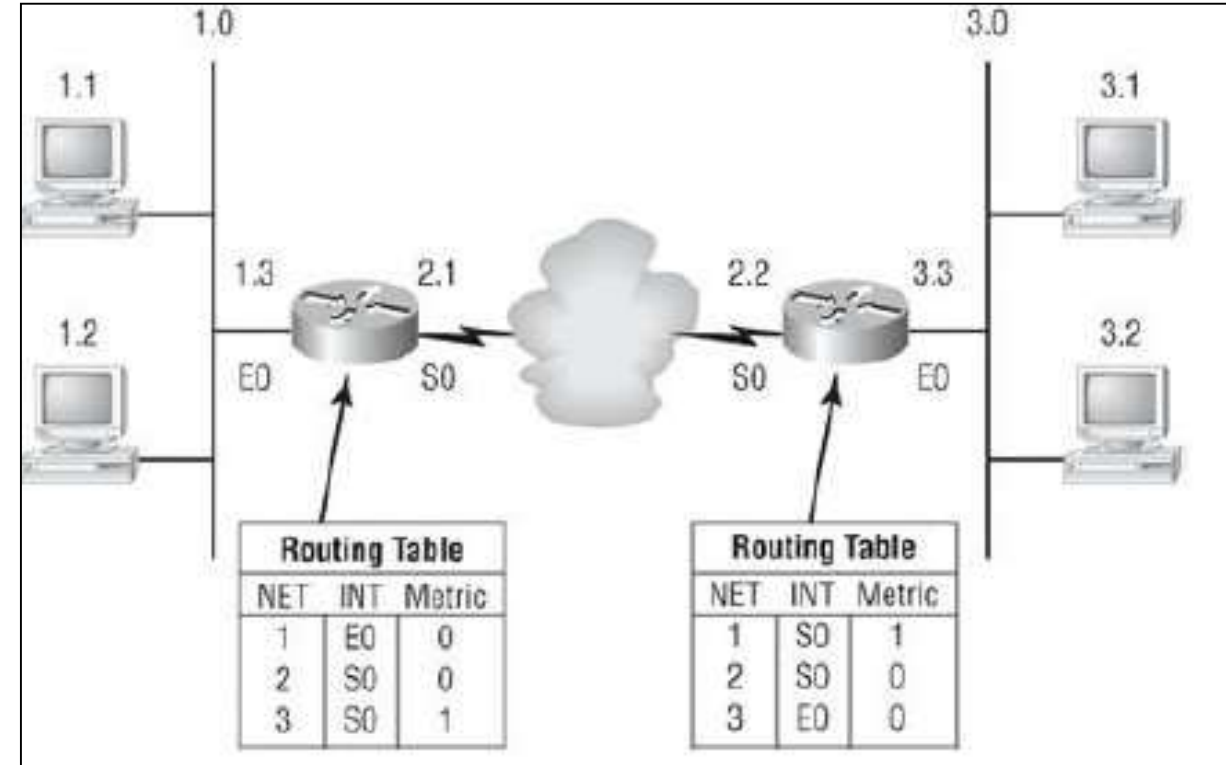
Troubleshoot: ARP, Address Tables, Protocol Analyzers



OSI LAYER-3: NETWORK

Jobs:

- Encapsulates a **packet**
- Defines IP Addresses
- Find paths to destinations
- Creates broadcast domains
- Provides QoS, ACL, Encryption
- Connects different data link types together

Devices: Router, L3 Switch**Protocols:** TCP/IP, IPX, AppleTalk**Troubleshoot:** Traceroute, ping, ARP, Routing Table, Sniffing

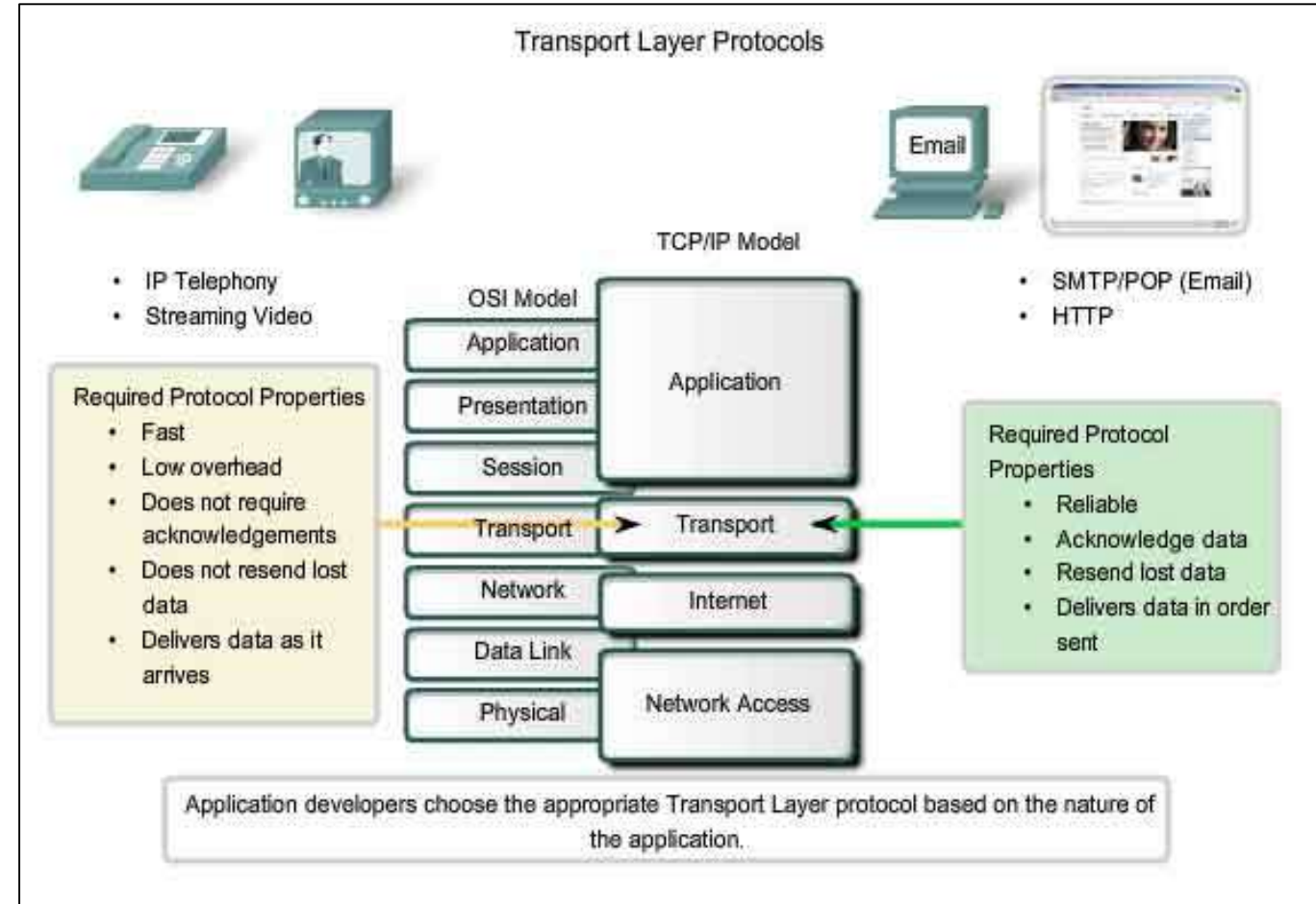
OSI LAYER-4: TRANSPORT

Jobs:

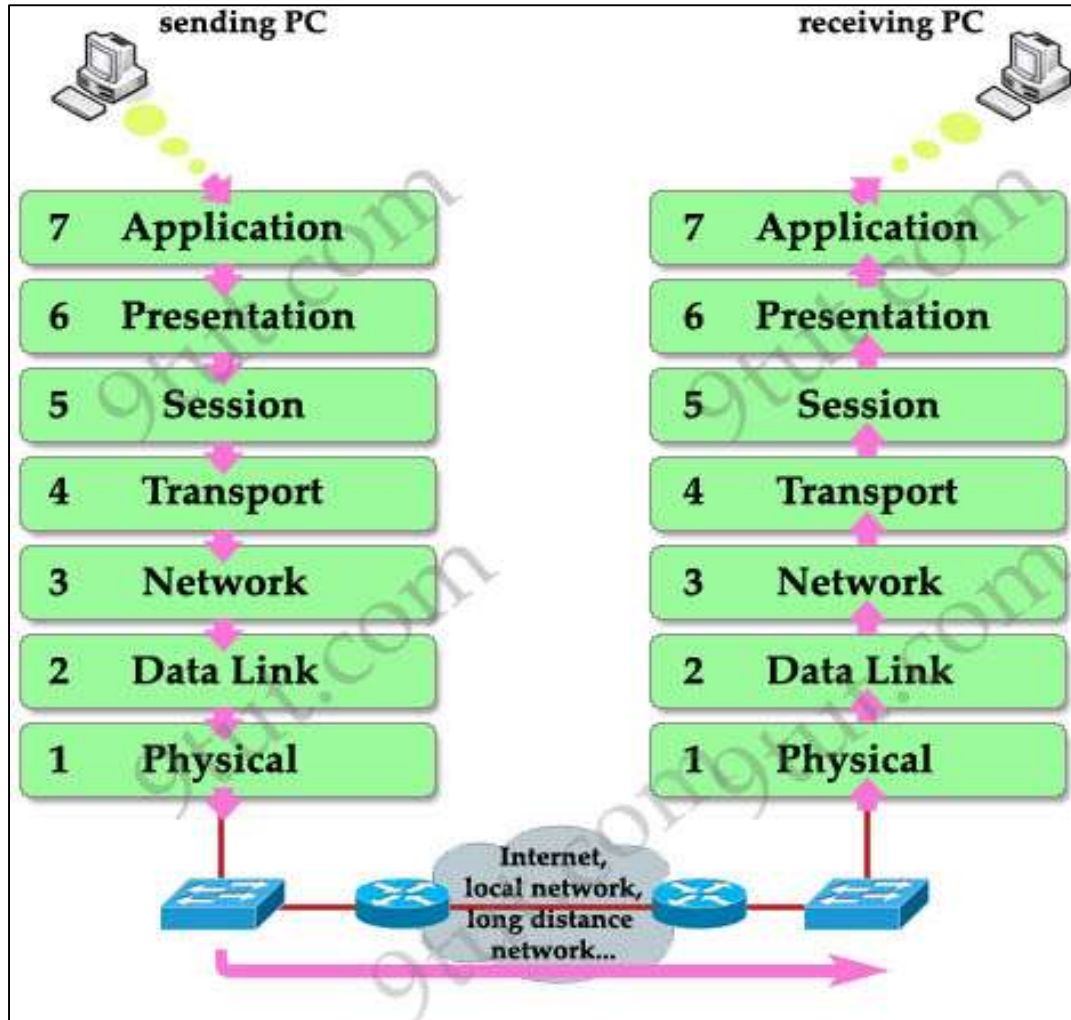
- Encapsulates a **segment**
- Connection management (setup / tear down)
- Reliable / unreliable delivery of data
- Flow control (buffering, windowing, congestion avoidance)
- Multiplexing

Devices: Firewall, Multilayer Switch

Protocols: TCP, UDP

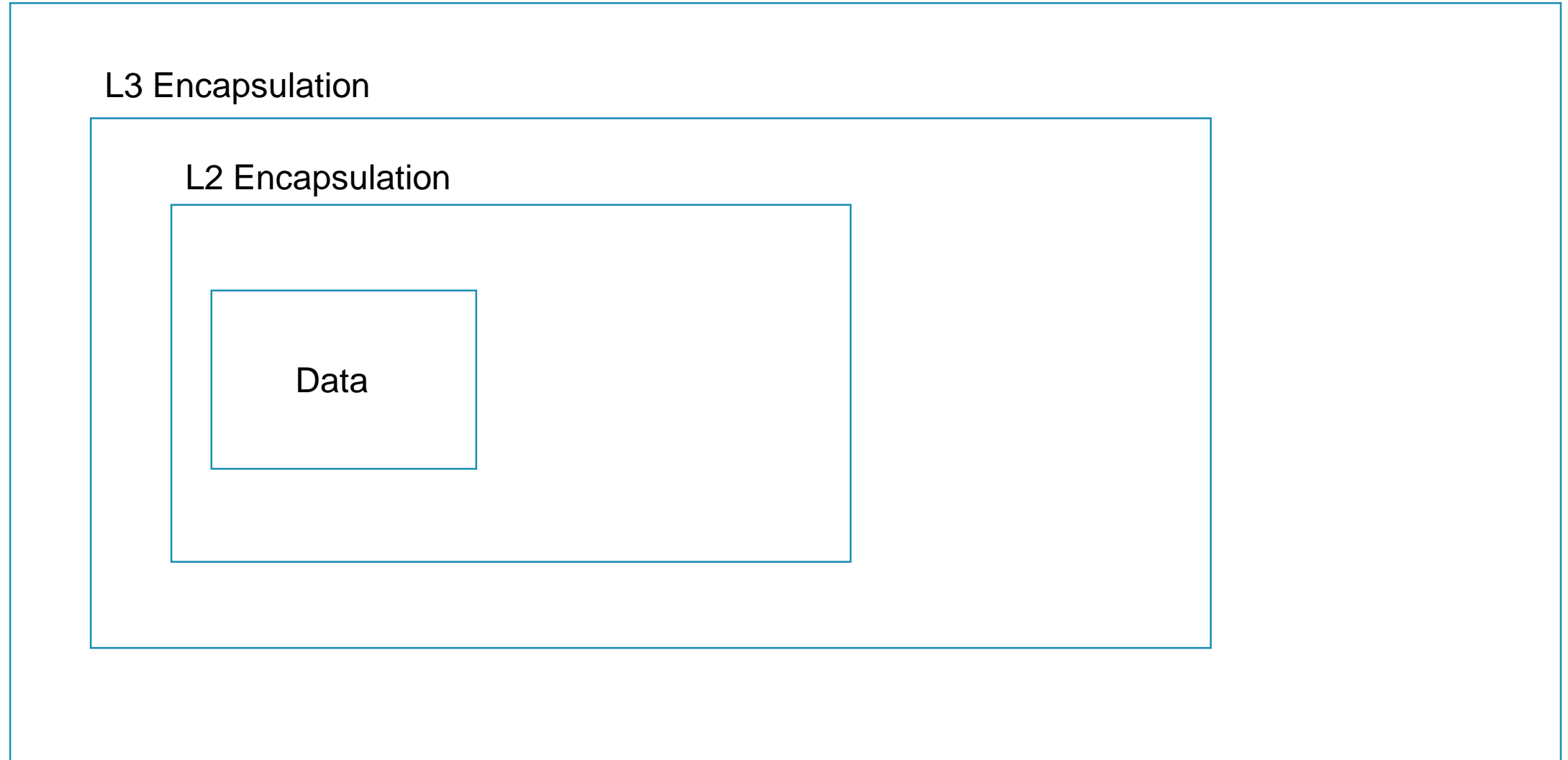


OSI REFERENCE MODEL AS A WHOLE

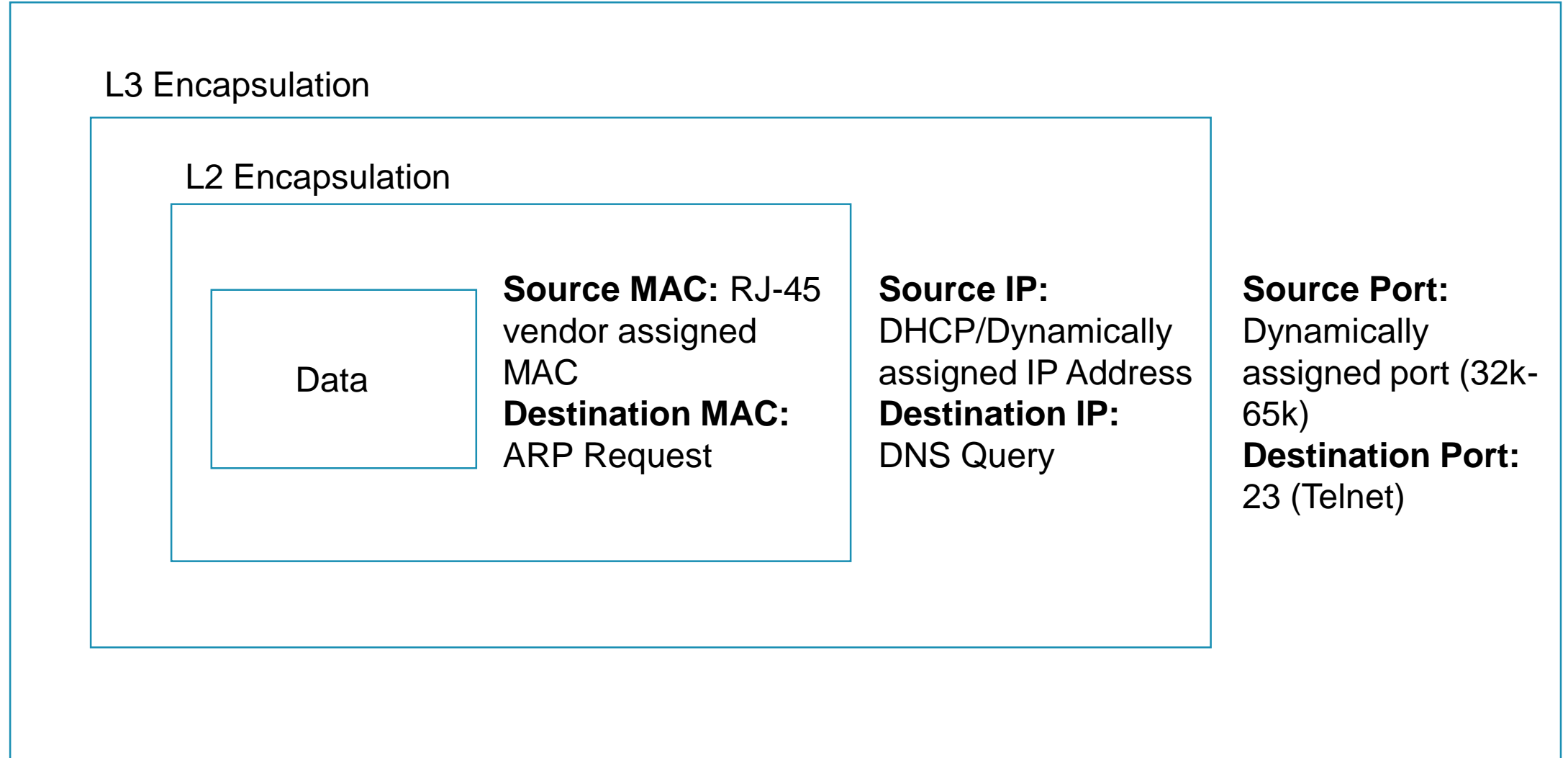


| | |
|--------------|--|
| Application | • File, print, message, database, and application services |
| Presentation | • Data encryption, compression, and translation services |
| Session | • Dialog control |
| Transport | • End-to-end connection |
| Network | • Routing |
| Data Link | • Framing |
| Physical | • Physical topology |

L4 Encapsulation



L4 Encapsulation





Switches & L2 Problems

SWITCH FUNCTIONS

Uses (ASICs – Application Specific Integrated Circuits) **hardware switching** (Max. Speed 400M fps)

Supports **full-duplex**

Switching Methods are;

Store and Forward: Put entire frame to buffer, check FCS via CRC, if passes forward, otherwise drop

Cut-Through: Read first 14-bytes and start forwarding. Faster but may cause collisions

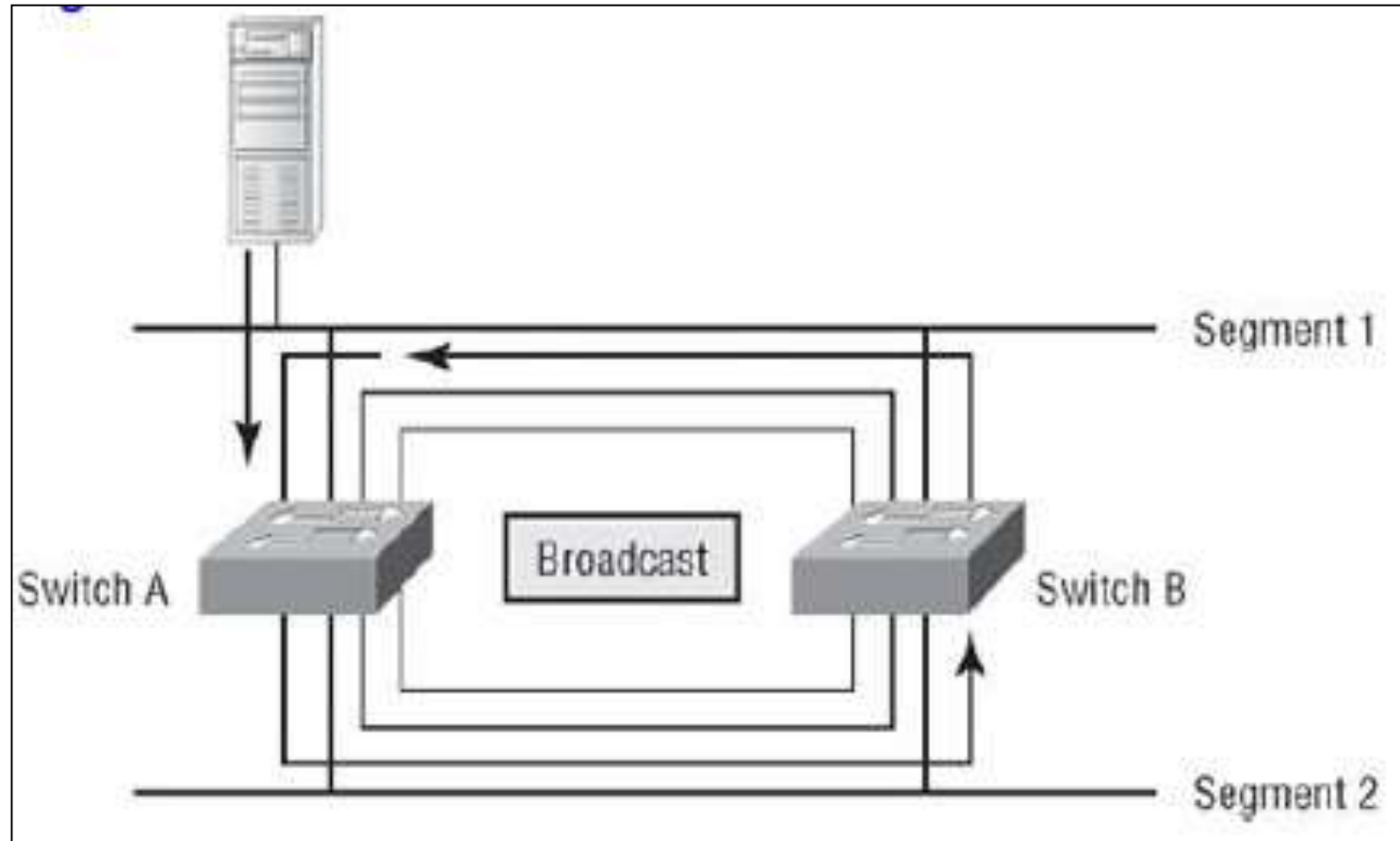
Fragment-Free: Read first 64-bytes and start forwarding

Uses **STP** to overcome L2 loops



L2 PROBLEM & RESOLUTION

- L2 loop is the main L2 problem
- The resolution is the STP

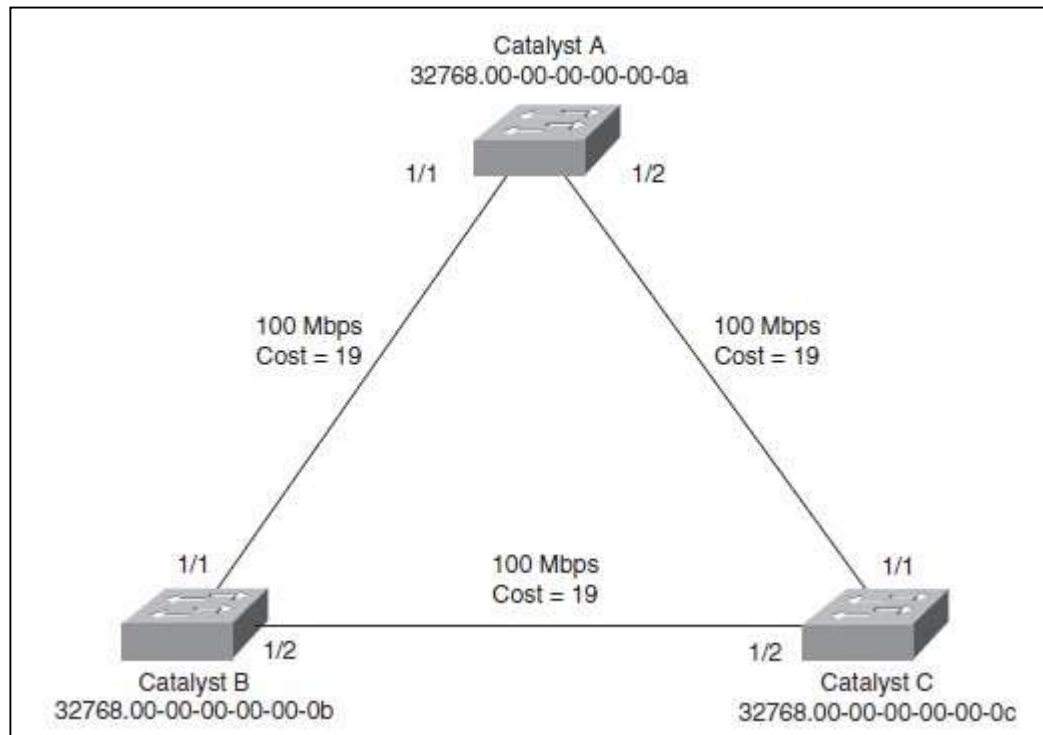


STP (SPANNING TREE PROTOCOL)

STP is a L2 loop prevention method.

Rules;

1. Elect the root bridge (Lowest Priority, if equals, lowest MAC Address)
2. Determine Root bridge's ports as Designated Ports (FWD)
3. Determine Root Ports (FWD); other bridge's the nearest cost'ed ports to the Root Bridge
4. Remind that each segment should have one Designated Port (FWD)
5. Block the non-designated ports in each segment to prevent L2 loops



| Link Bandwidth | Old STP Cost | New STP Cost |
|----------------|--------------|--------------|
| 4 Mbps | 250 | 250 |
| 10 Mbps | 100 | 100 |
| 16 Mbps | 63 | 62 |
| 45 Mbps | 22 | 39 |
| 100 Mbps | 10 | 19 |
| 155 Mbps | 6 | 14 |
| 622 Mbps | 2 | 6 |
| 1 Gbps | 1 | 4 |
| 10 Gbps | 0 | 2 |

STP PORT STATES

- CST or Common STP (IEEE 802.1d) has 50 secs convergence time, which is not acceptable in today's networks. BPDU Hello Timer = 2secs.
- During calculation of the STP, all ports are disabled and the following port states are observed.
- If ports are access ports, then PortFast feature should be used for faster convergence.
- Only the root bridge can send the topology change information.

| STP State | The Port Can... | The Port Cannot... | Duration |
|------------|--|---|---|
| Disabled | N/A | Send or receive data | N/A |
| Blocking | Receive BPDUs | Send or receive data or learn MAC addresses | Indefinite if loop has been detected |
| Listening | Send and receive BPDUs | Send or receive data or learn MAC addresses | Forward Delay timer (15 seconds) |
| Learning | Send and receive BPDUs and learn MAC addresses | Send or receive data | Forward Delay timer (15 seconds) |
| Forwarding | Send and receive BPDUs, learn MAC addresses, and send and receive data | | Indefinite as long as port is up and loop is not detected |

STP ENHANCEMENTS; RSTP

Since STP convergence time is too long, new technologies are implemented as RSTP (Rapid STP) and MSTP (Multi-STP) for faster convergence and Per VLAN capabilities.

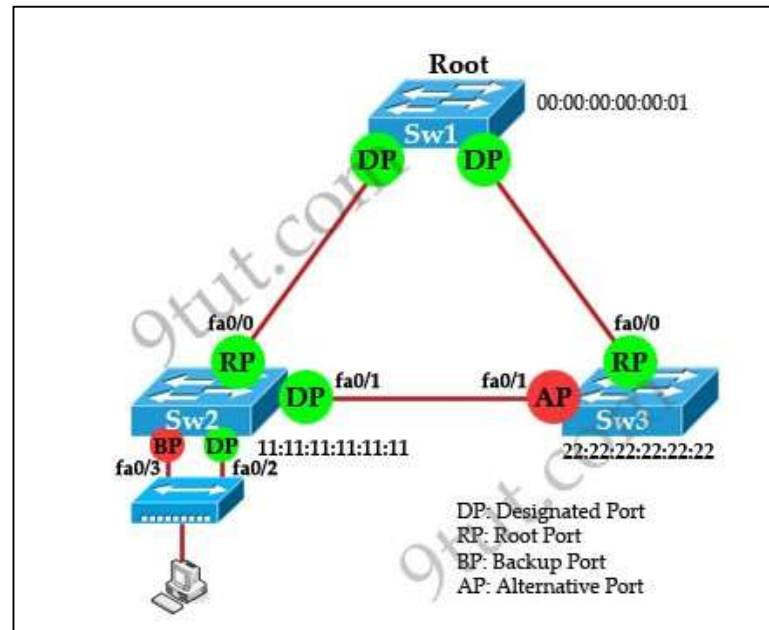
RSTP is an IEEE 802.1w protocol that speeds up the convergence time to 6 seconds only with the following new port states:

Discarding - Incoming frames simply are dropped; no MAC addresses are learned. (This state combines the 802.1D Disabled, Blocking, and Listening states because all three did not effectively forward anything. The Listening state is not needed because RSTP quickly can negotiate a state change without listening for BPDUs first.)

Learning - Incoming frames are dropped, but MAC addresses are learned.

Forwarding - Incoming frames are forwarded according to MAC addresses that have been (and are being) learned.

EdgePort term is replaced with PortFast for RSTP. Full-duplex switch-to-switch ports are defined as “Point-to-Point Port”.



STP ENHANCEMENTS; MST

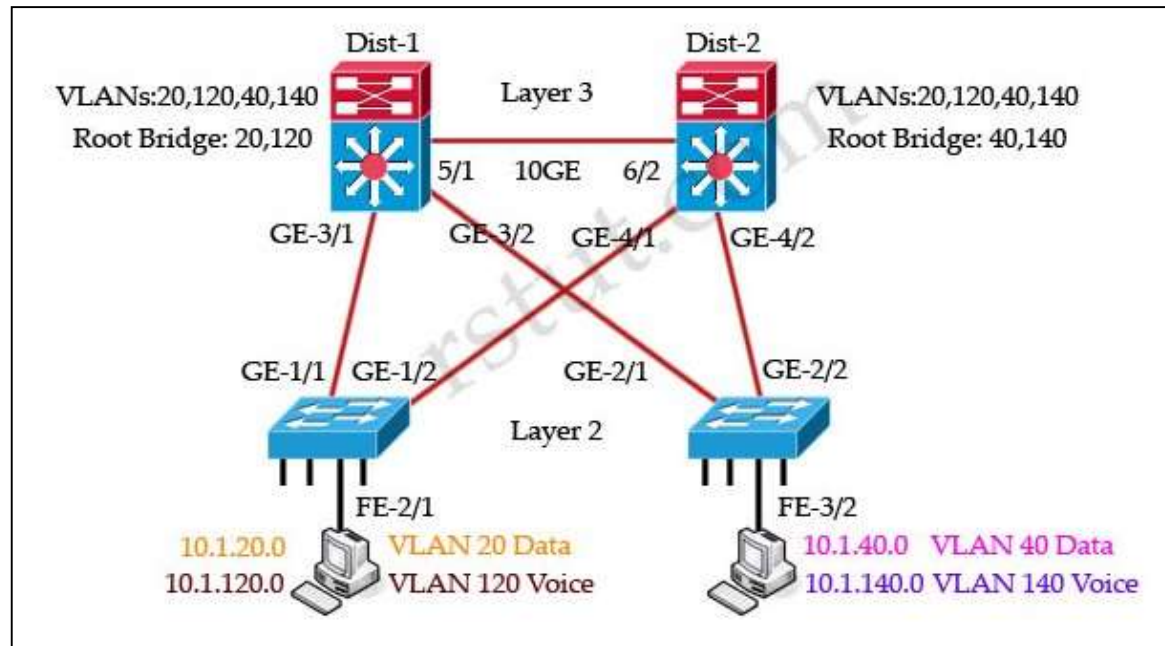
MST (IEEE 802.1s) is built on the concept of mapping one or more VLANs to a single STP instance. Multiple instances of STP can be used (hence the name MST), with each instance supporting a different group of VLANs.

MST is working with regions. If two switches have the same set of attributes, they belong to the same MST region. If not, they belong to two independent regions.

- MST configuration name (32 characters)

- MST configuration revision number (0 to 65535)

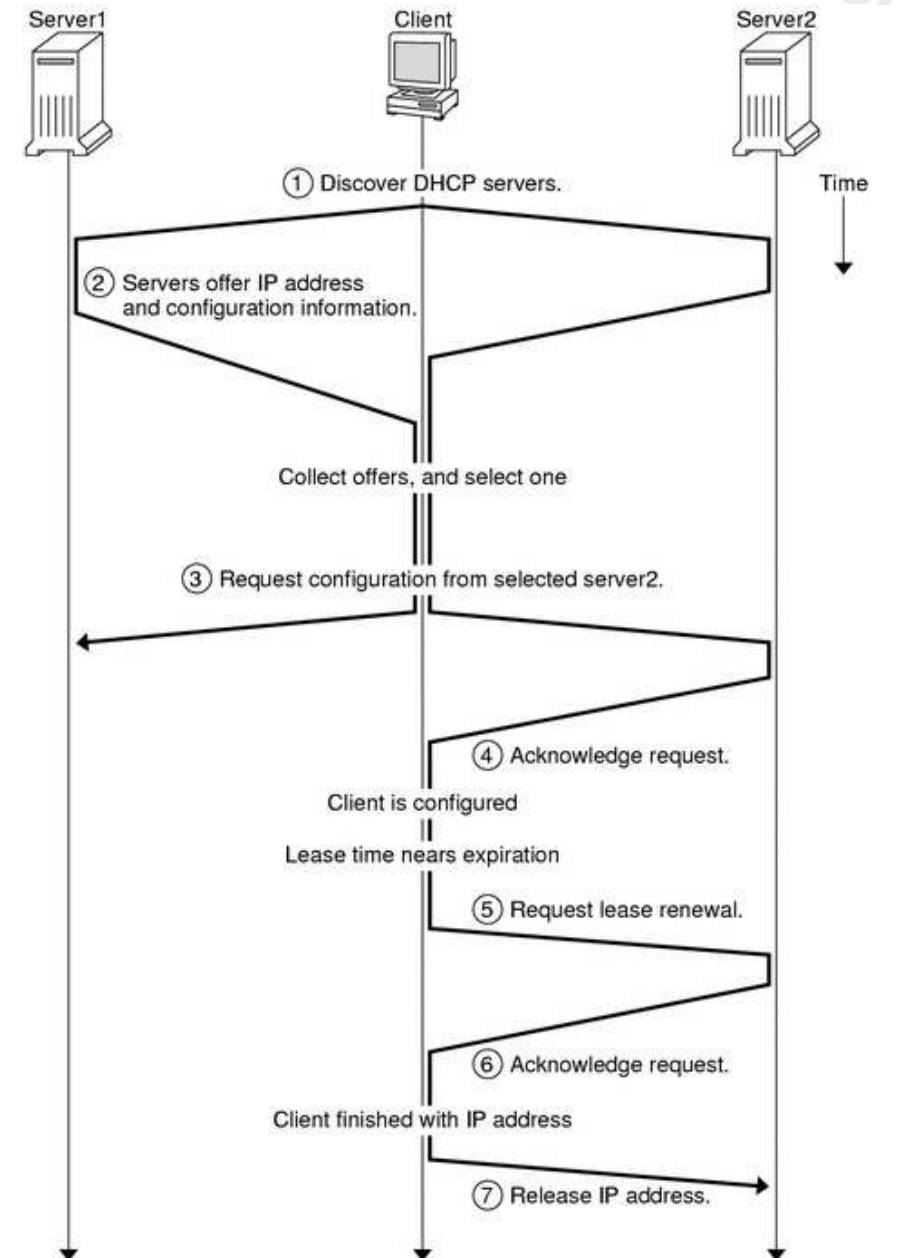
- MST instance-to-VLAN mapping table (4096 entries)



DHCP SEQUENCE

DHCP follows the below sequence:

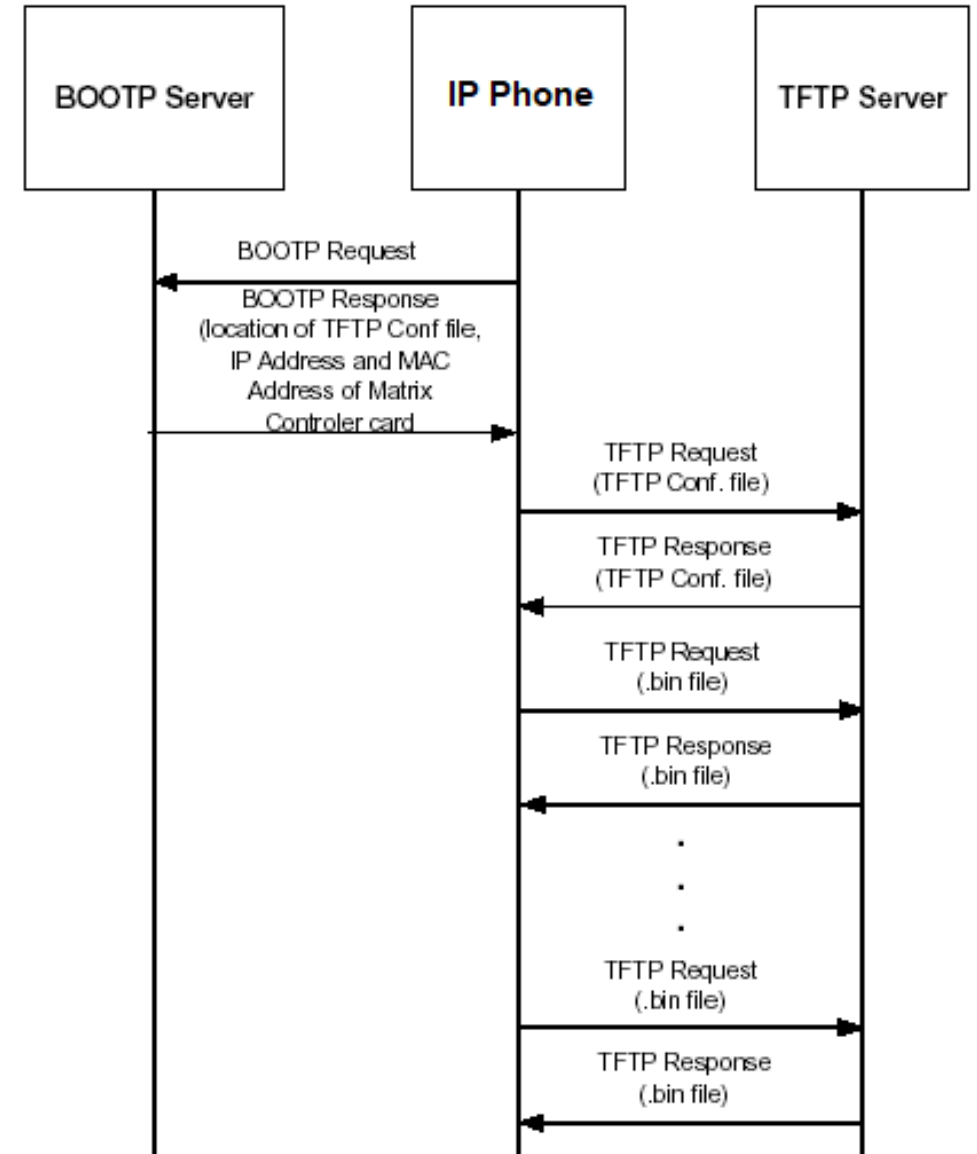
1. Client generates DHCPDISCOVER (multicast)
2. All DHCP Servers send DHCPOFFER (unicast) with the following information;
 - IP Address
 - Subnet Mask
 - Default Gateway IP Address
3. Client accepts the first arrived offer with DHCPREQUEST (DHCPDECLINE)
4. Server responds DHCPACK (DHCPNACK)



BOOTP SEQUENCE

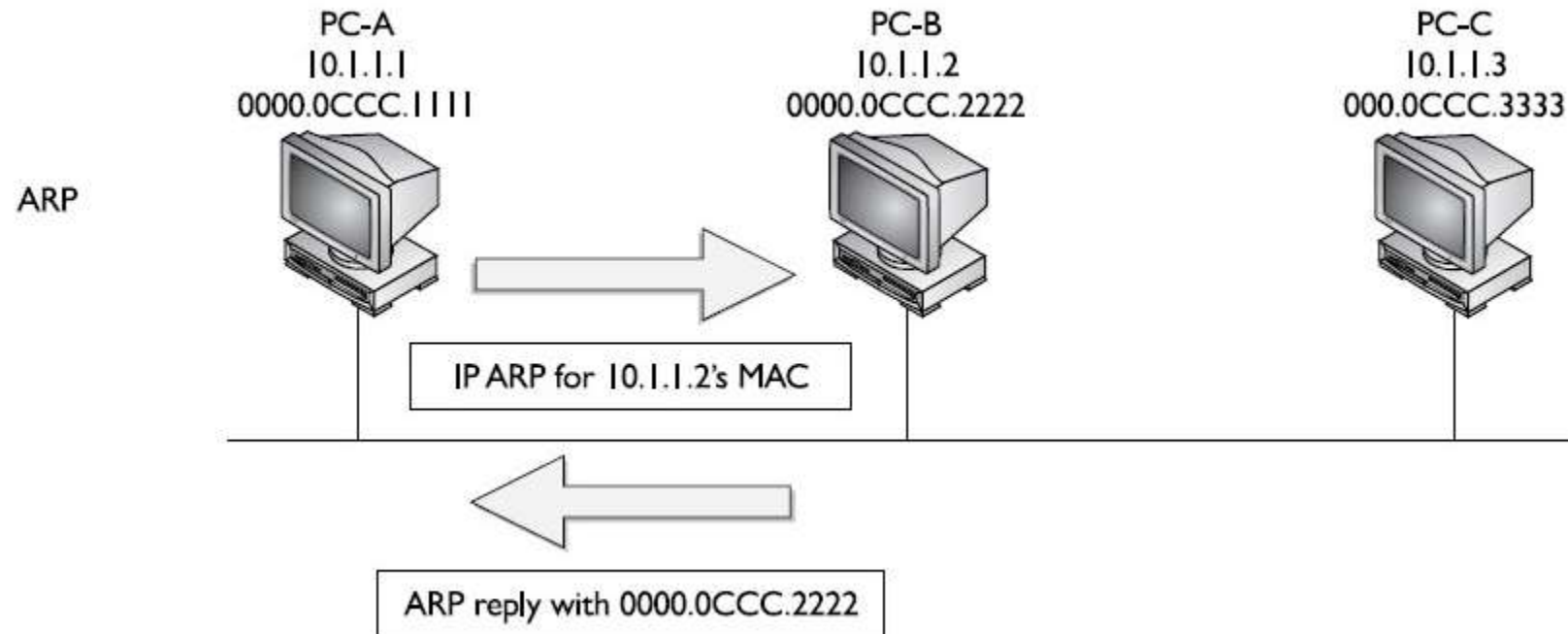
BootP follows the below sequence:

1. Client generates BootP_Request (Source IP: 0.0.0.0, Destination IP: 255.255.255.255, UDP Port: 67)
2. If Client and Server are not in the same subnet, then router performs BootP_Relay task between subnets and increments TTL
3. Server returns BootP_Reply (via ARP, UDP Port: 68) with the following information;
 - IP Address
 - Subnet Mask
 - Default Gateway Address
 - TFTP Server IP Address
 - The load path in the TFTP Server
4. Client requests boot file via TFTP (FTP over UDP)
5. Server provides boot file via TFTP
6. Client install the boot file and starts running

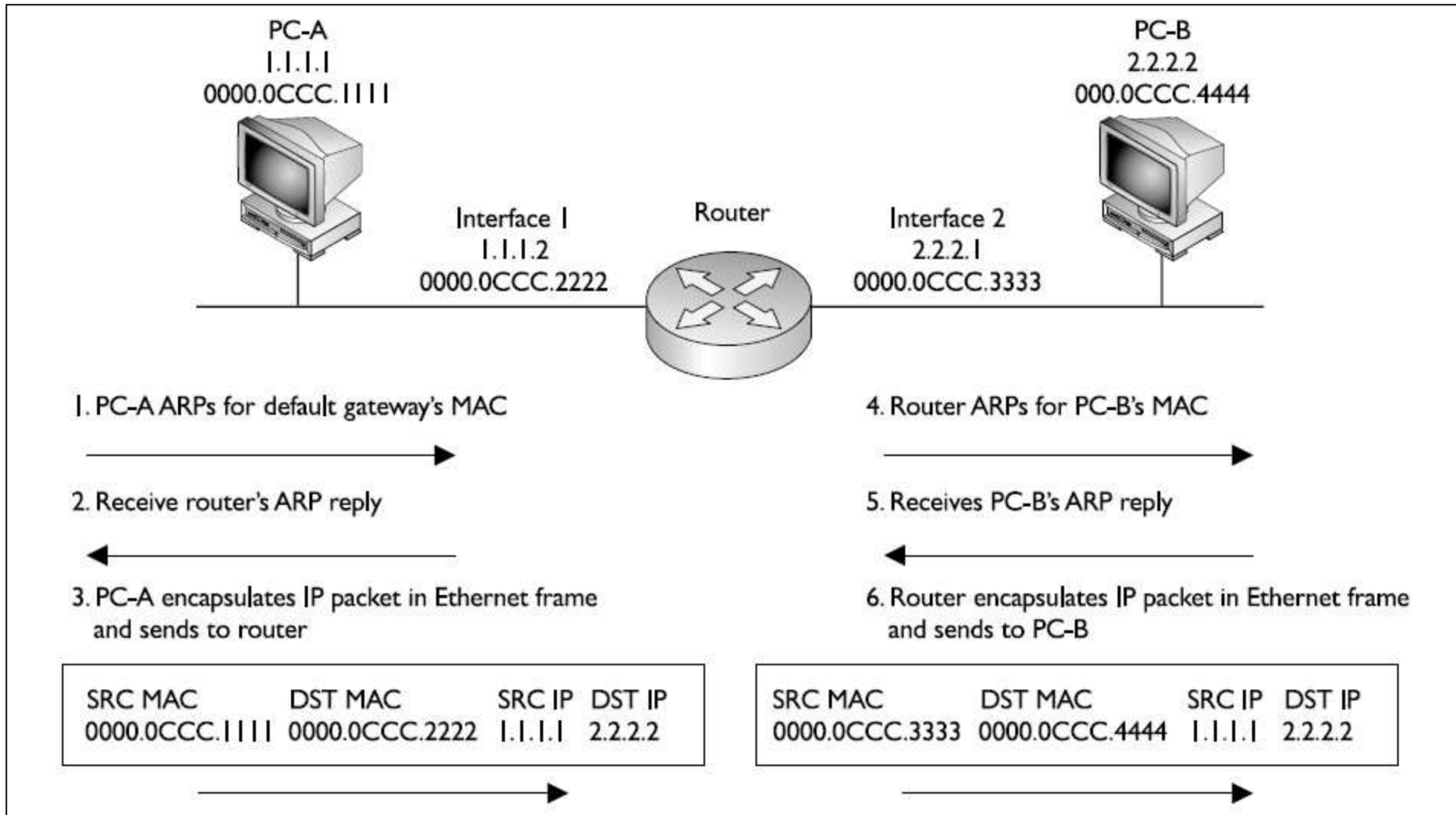


ARP - ADDRESS RESOLUTION PROTOCOL

ARP is a Layer-2 protocol, which is used to learn a device's MAC address.

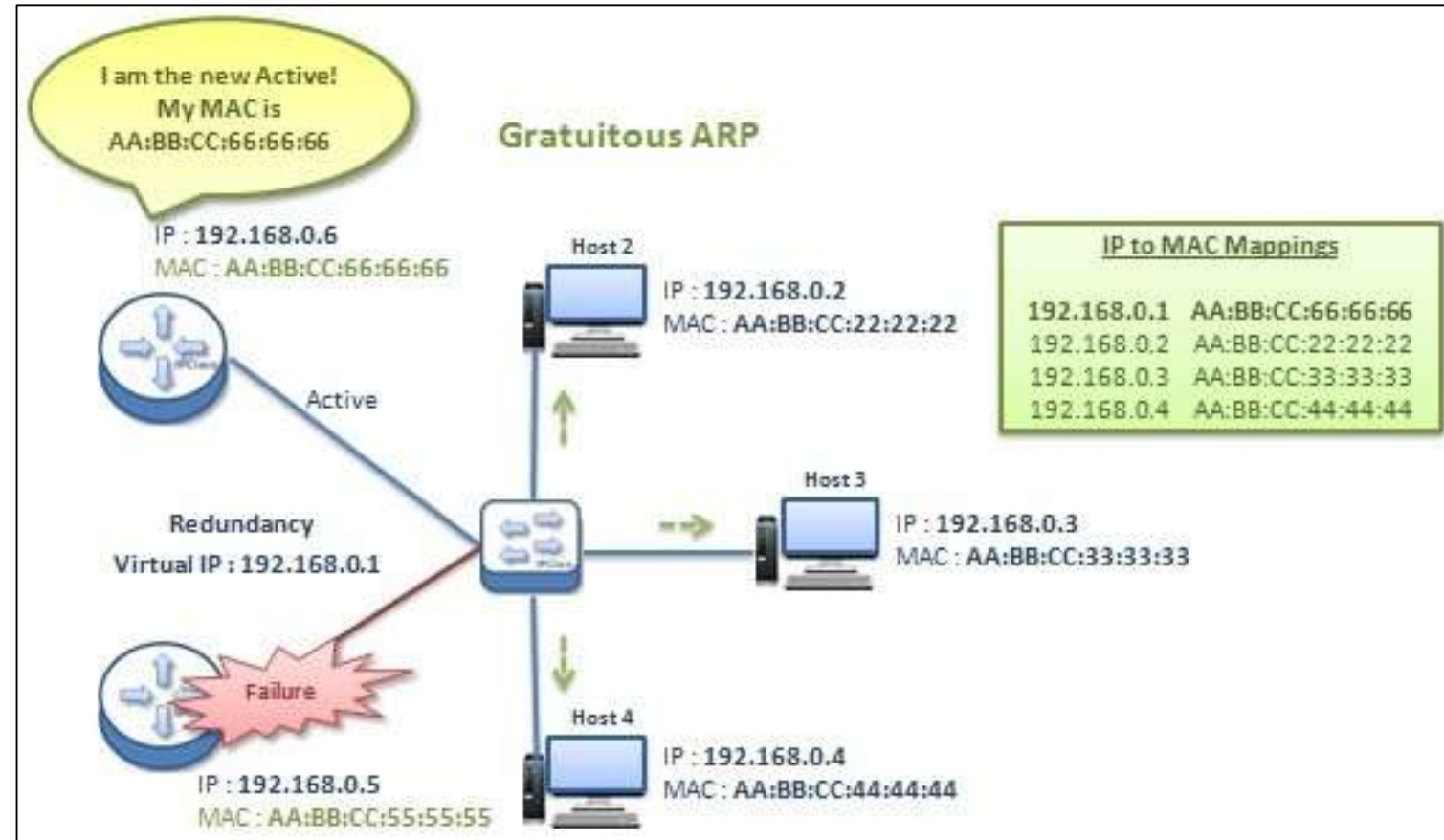
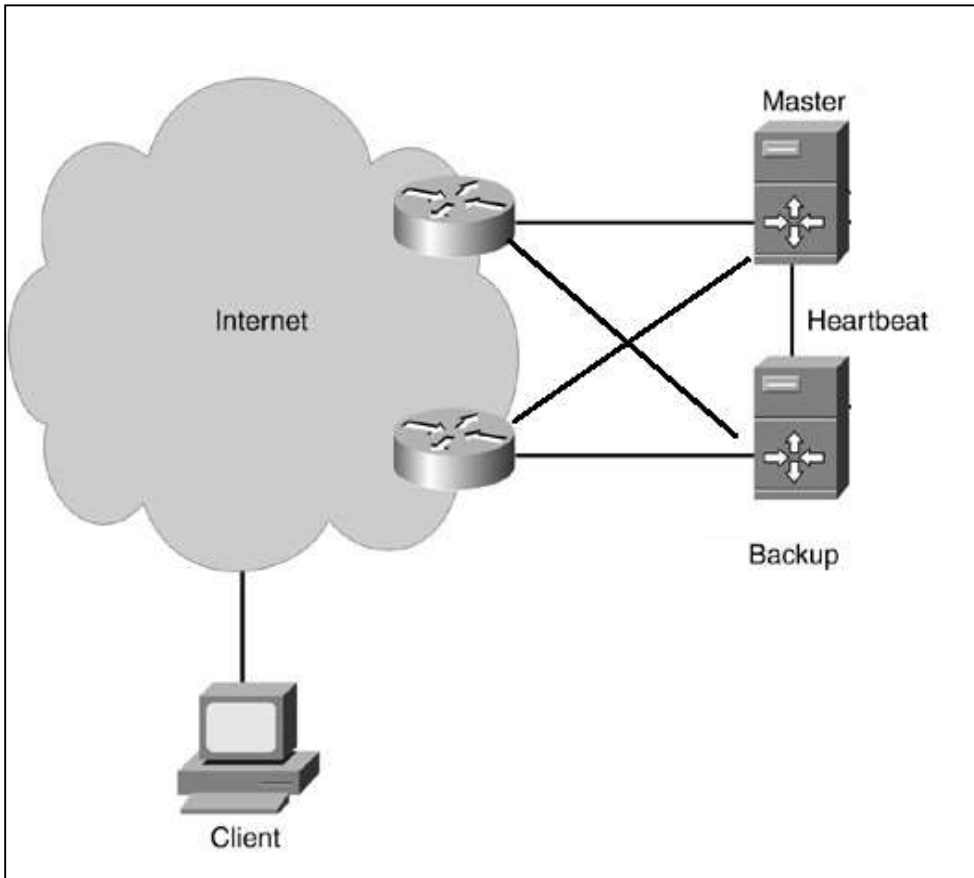


ARP DETAILED



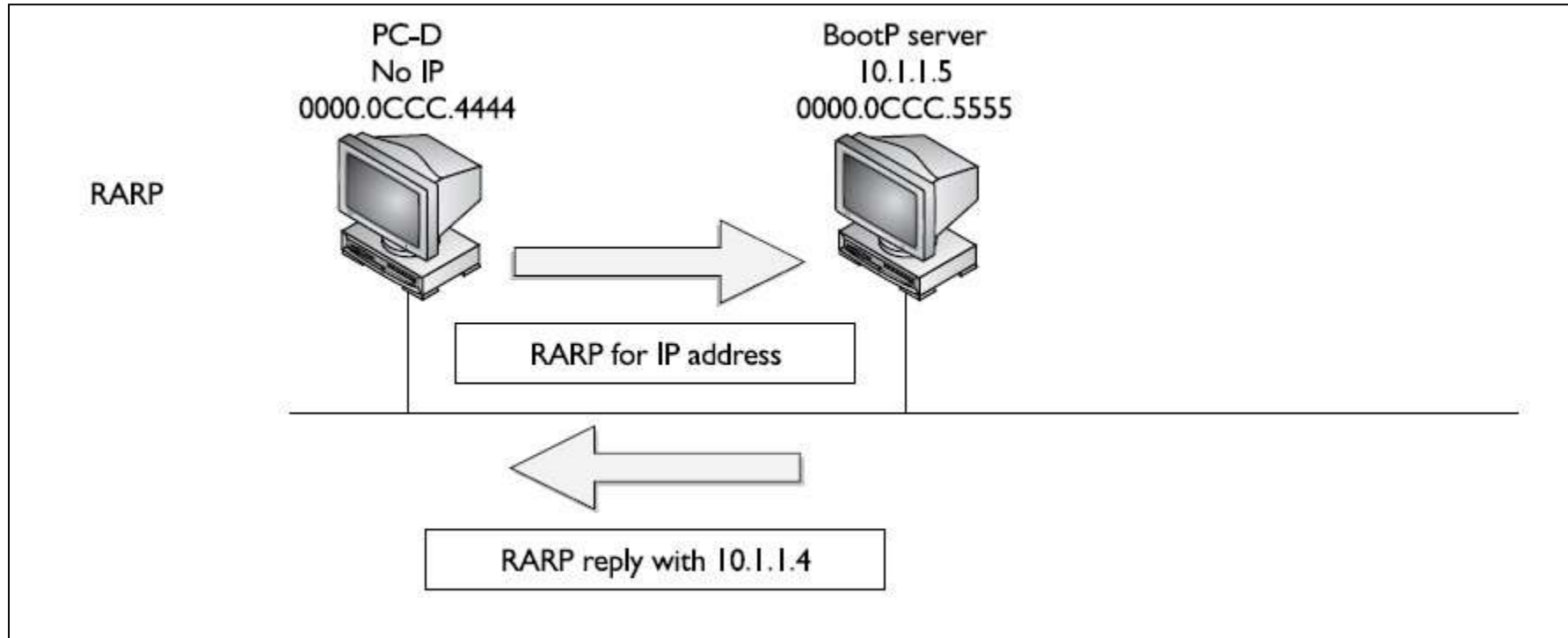
GARP

Gratuitous ARP is a kind of ARP reply without asking of ARP request.



RARP

Reverse ARP is used to learn the IP Address during DHCP / BootP operations



ICMP

ICMP is used to send error and control information between TCP/IP devices at the Internet layer (RFC 792)

Message types; Address Reply, Address Request, **Destination Unreachable**, **Echo**, **Echo Reply**, Information Reply, Information Request, Parameter Problem, Redirect, Subnet Mask Request, Time Exceeded, Timestamp, and Timestamp Reply.

```
C:\Users\emutlu>arp -a
```

Interface: 10.254.127.20 --- 0x9 - WI-FI

| Internet Address | Physical Address | Type |
|------------------|-------------------|---------|
| 10.254.127.1 | 00-00-0c-07-ac-7f | dynamic |
| 10.254.127.127 | ff-ff-ff-ff-ff-ff | static |
| 224.0.0.22 | 01-00-5e-00-00-16 | static |
| 224.0.0.251 | 01-00-5e-00-00-fb | static |
| 224.0.0.252 | 01-00-5e-00-00-fc | static |
| 239.255.255.250 | 01-00-5e-7f-ff-fa | static |
| 255.255.255.255 | ff-ff-ff-ff-ff-ff | static |

Interface: 10.254.116.38 --- 0xb -- ETHERNET

| Internet Address | Physical Address | Type |
|------------------|-------------------|---------|
| 10.254.116.1 | b4-0c-25-e0-40-16 | dynamic |
| 10.254.116.23 | 30-e1-71-83-ca-b6 | dynamic |
| 10.254.116.24 | f8-b4-6a-96-bc-ca | dynamic |
| 10.254.116.29 | 48-ba-4e-f7-48-43 | dynamic |
| 10.254.116.35 | bc-e9-2f-cc-10-ab | dynamic |
| 10.254.116.255 | ff-ff-ff-ff-ff-ff | static |
| 224.0.0.2 | 01-00-5e-00-00-02 | static |
| 224.0.0.22 | 01-00-5e-00-00-16 | static |
| 224.0.0.251 | 01-00-5e-00-00-fb | static |
| 224.0.0.252 | 01-00-5e-00-00-fc | static |
| 239.255.255.250 | 01-00-5e-7f-ff-fa | static |
| 255.255.255.255 | ff-ff-ff-ff-ff-ff | static |

```
C:\Users\emutlu>ping ribboncommunications.com
```

```
Pinging ribboncommunications.com [23.185.0.4] with 32 bytes of data:
```

```
Reply from 23.185.0.4: bytes=32 time=45ms TTL=54
```

```
Reply from 23.185.0.4: bytes=32 time=45ms TTL=54
```

```
Reply from 23.185.0.4: bytes=32 time=45ms TTL=54
```

```
Reply from 23.185.0.4: bytes=32 time=45ms TTL=54
```

```
Ping statistics for 23.185.0.4:
```

```
Packets: Sent = 4, Received = 4, Lost = 0 (0% loss),
```

```
Approximate round trip times in milli-seconds:
```

```
Minimum = 45ms, Maximum = 45ms, Average = 45ms
```

```
C:\Users\emutlu>tracert ribboncommunications.com
```

```
Tracing route to ribboncommunications.com [23.185.0.4]
```

```
over a maximum of 30 hops:
```

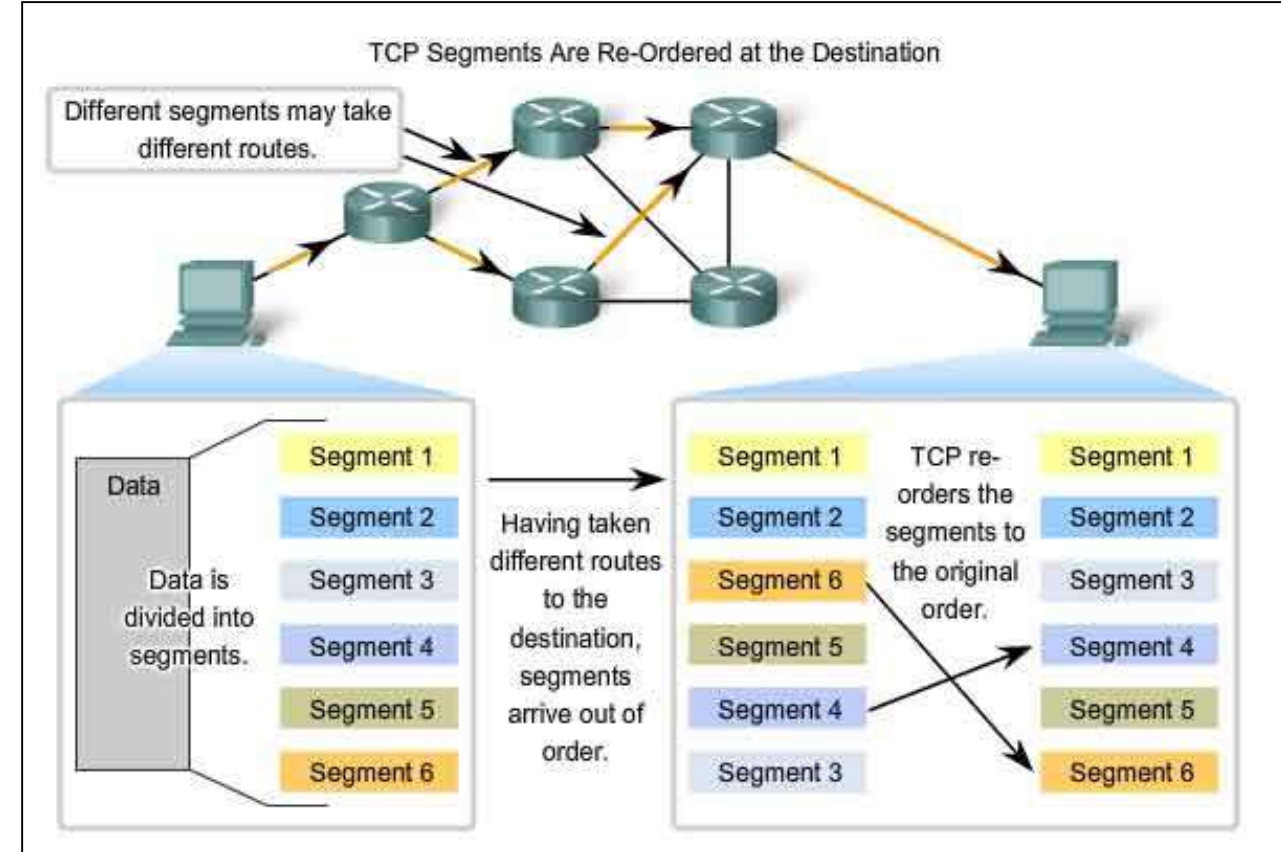
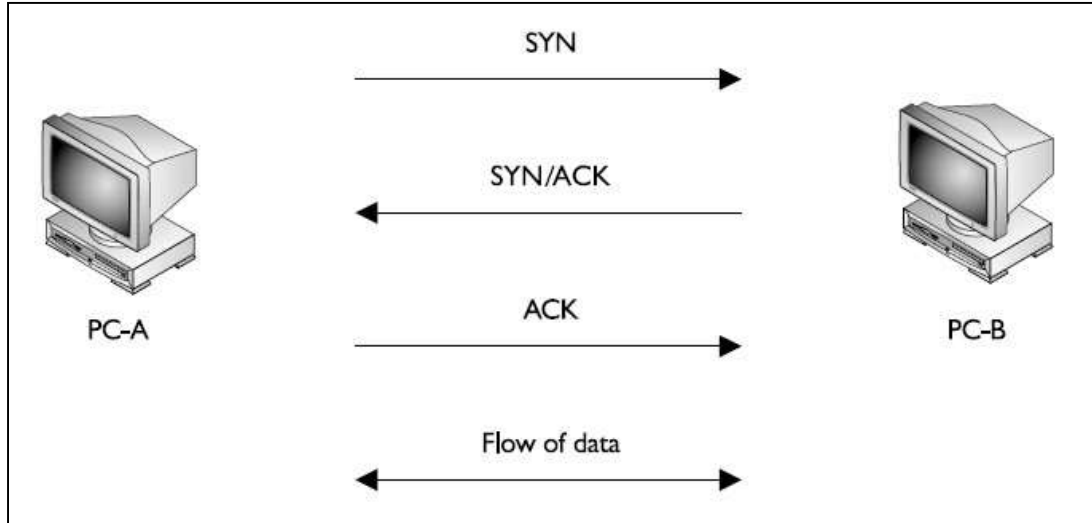
| | | | | |
|---|-------|-------|-------|--|
| 1 | <1 ms | <1 ms | <1 ms | 10.254.116.1 |
| 2 | 1 ms | 1 ms | 1 ms | netasfw.netas.lab.nortel.com |
| 3 | 2 ms | 1 ms | 1 ms | intrtr1.netas.lab.nortel.com |
| 4 | 3 ms | 16 ms | 3 ms | host-213-74-185-149.superonline.net [213.74.185.149] |
| 5 | * | * | * | Request timed out. |
| 6 | * | * | * | Request timed out. |

TCP/IP and Transport Layer

TCP PROTOCOL

TCP (Transmission Control Protocol) is a connection-oriented, reliable protocol

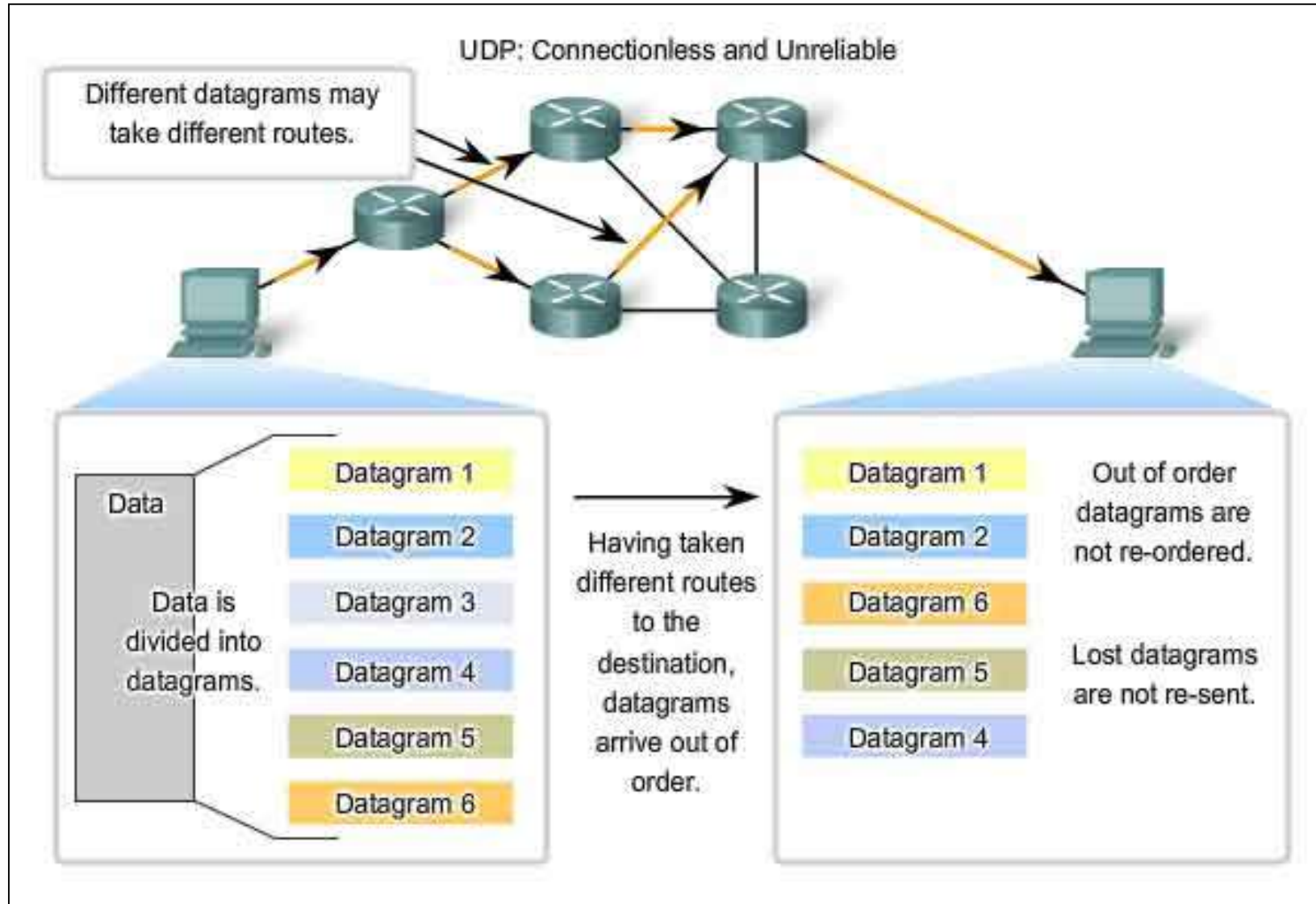
Applications: FTP, HTTP, Telnet, SSH, SMTP, POP3 for TCP

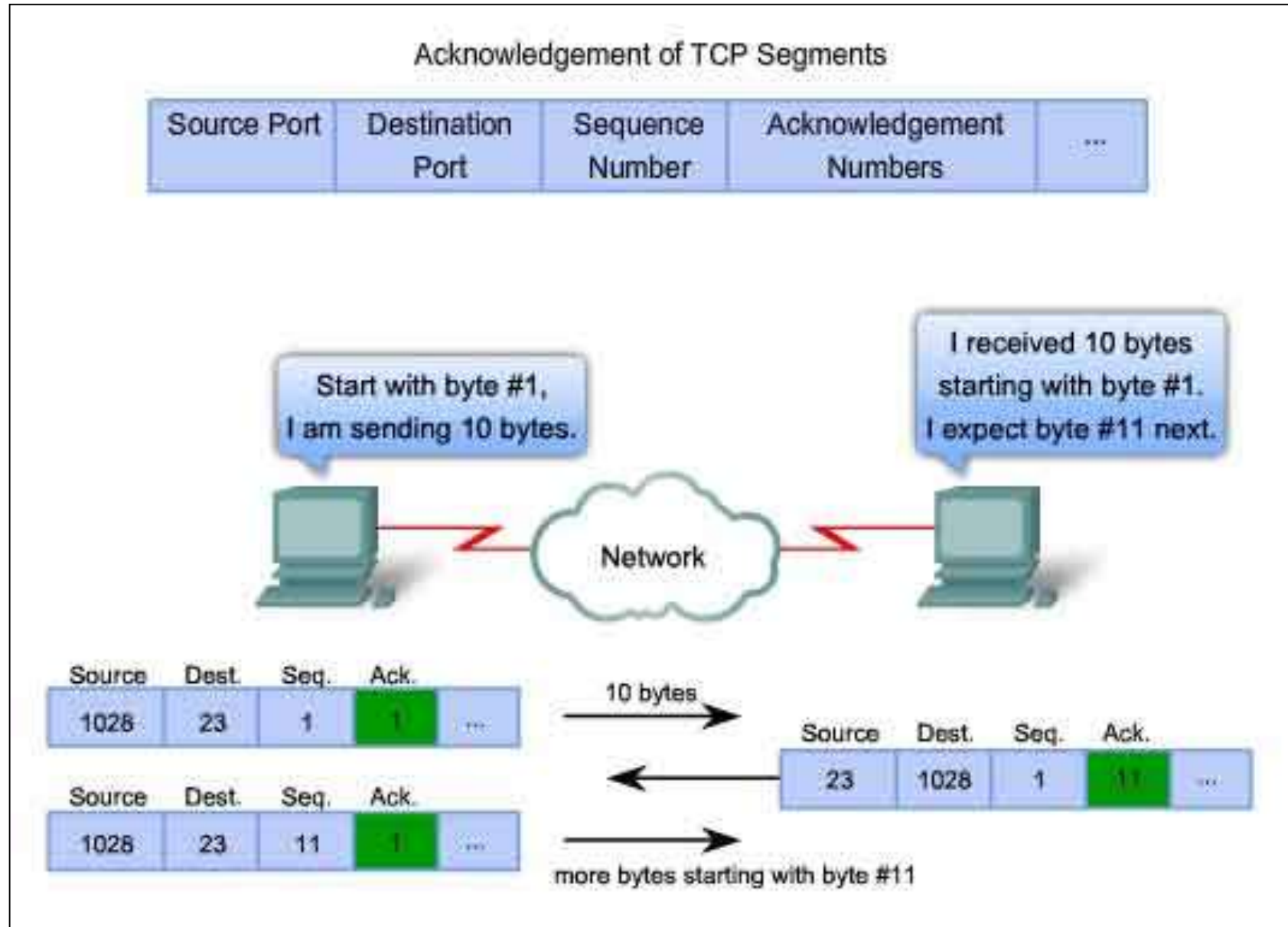


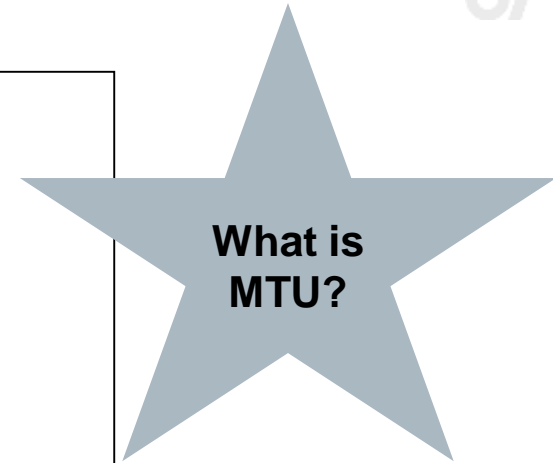
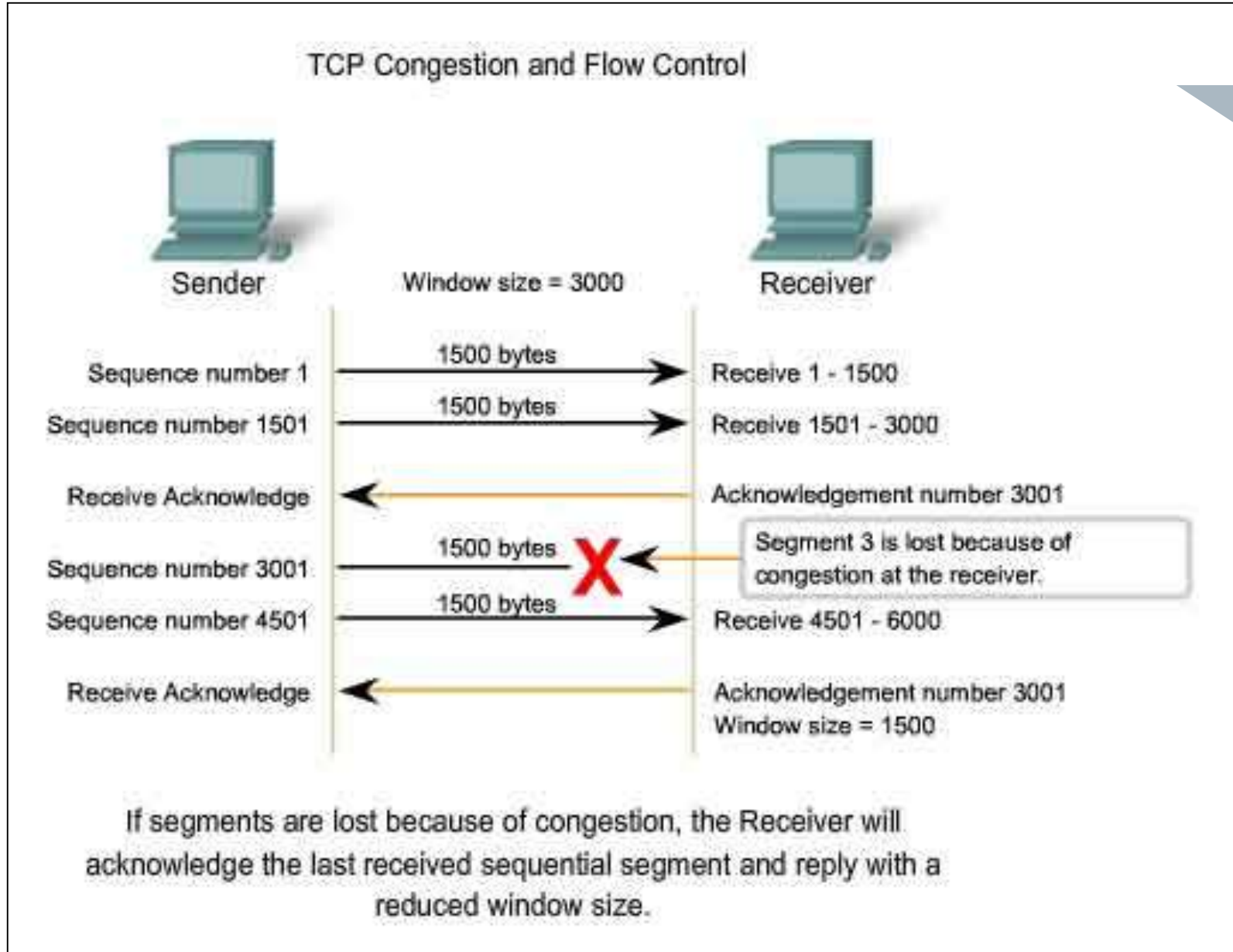
UDP PROTOCOL

UDP (User Datagram Protocol) is a connectionless, unreliable protocol

Applications: DNS, TFTP, SNMP, RIP for UDP







COMMON PORT NUMBERS

- Total 65536 (16-bits) of ports are available
- 0-1023 are reserved by IANA
- 1024-49151 are assigned by IANA to applications
- Operating System can assign ports in range 49152-to-65535
- Well-known ports are; HTTP:80, POP3:110, SMTP:25, SNMP:161, DNS:53, RIP:520, TFTP:69, SSH:22, Telnet:23, FTP:21

| | | | |
|---------------------|-----------------------|------------------------|------------------------|
| 7 Echo | 554 RTSP | 2745 Bagle.H | 6891-6901 Windows Live |
| 19 Chargen | 546-547 DHCPv6 | 2967 Symantec AV | 6970 Quicktime |
| 20-21 FTP | 560 rmonitor | 3050 Interbase DB | 7212 GhostSurf |
| 22 SSH/SCP | 563 NNTP over SSL | 3074 XBOX Live | 7648-7649 CU-SeeMe |
| 23 Telnet | 587 SMTP | 3124 HTTP Proxy | 8000 Internet Radio |
| 25 SMTP | 591 FileMaker | 3127 MyDoom | 8080 HTTP Proxy |
| 42 WINS Replication | 593 Microsoft DCOM | 3128 HTTP Proxy | 8086-8087 Kaspersky AV |
| 43 WHOIS | 631 Internet Printing | 3222 GLBP | 8118 Privoxy |
| 49 TACACS | 636 LDAP over SSL | 3260 iSCSI Target | 8200 VMware Server |
| 53 DNS | 639 MSDP (PIM) | 3306 MySQL | 8500 Adobe ColdFusion |
| 67-68 DHCP/BOOTP | 646 LDP (MPLS) | 3389 Terminal Server | 8767 TeamSpeak |
| 69 TFTP | 691 MS Exchange | 3689 iTunes | 8866 Bagle.B |
| 70 Gopher | 860 iSCSI | 3690 Subversion | 9100 HP JetDirect |
| 79 Finger | 873 rsync | 3724 World of Warcraft | 9101-9103 Bacula |
| 80 HTTP | 902 VMware Server | 3784-3785 Ventrilo | 9119 MXit |
| 88 Kerberos | 989-990 FTP over SSL | 4333 mSQL | 9800 WebDAV |
| 102 MS Exchange | 993 IMAP4 over SSL | 4444 Blaster | 9898 Dabber |
| 110 POP3 | 995 POP3 over SSL | 4664 Google Desktop | 9988 Rbot/Spybot |



IPv4 Addressing

IPV4 ADDRESS CLASSES & TYPES

Class A: 8-bits network, 24-bits host number. 1 to 126, 0 means all, 127 is reserved for testing

Class B: 16-bits network, 16-bits host number. 128 to 191 (Binary: 1000 0000 - 1011 1111)

Class C: 24-bits network, 8-bits host number. 192 to 223 (Binary: 1100 0000 - 1110 1111)

Class D: Multicasting. 224 to 239 (Binary: 1110 0000 - 1110 1111)

Class E: Reserved. 240 to 254. 255 is reserved for broadcasting.

Private IP Addresses;

Class A: 10.0.0.0 - 10.255.255.255 (1 Class A Network)

Class B: 172.16.0.0 - 172.31.255.255 (16 Class B Network)

Class C: 192.168.0.0 - 192.168.255.255 (256 Class C Network)

IPv4 Address Types:

Reserved: Broadcast and Class D&E IPs.

Network ID: VLAN IP

Host ID: Host's IP

Directed Broadcast: Routable addresses by routers

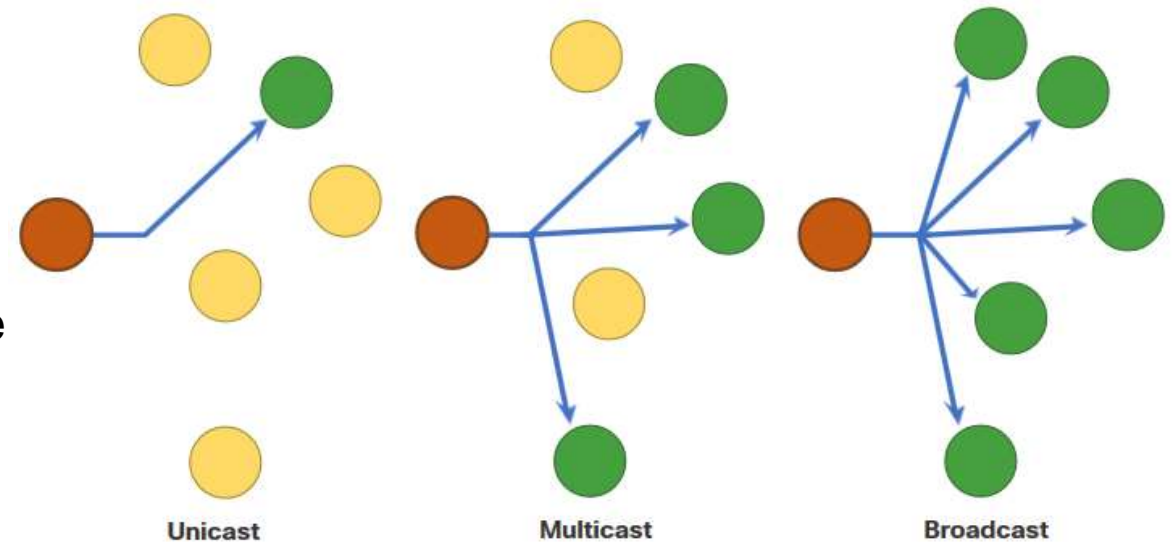
Local Broadcast: 255.255.255.255 - router will not route

Loopback: 127.0.0.1

Auto-configured: DHCP assigned IPs

Public: Internet accessible IPs

Private: Local Network IPs.



IPV4 ADDRESS CLASSES & TYPES

Five Different Classes of IPv4 Addresses

| Class | First Octet decimal (range) | First Octet binary (range) | IP range | Subnet Mask | Hosts per Network ID | # of networks |
|---------------------------|-----------------------------|----------------------------|---------------------------|---------------|----------------------|---------------|
| Class A | 0 — 127 | 0 XXXXXXXX | 0.0.0.0-127.255.255.255 | 255.0.0.0 | $2^{24} - 2$ | 2^7 |
| Class B | 128 — 191 | 10 XXXXXX | 128.0.0.0-191.255.255.255 | 255.255.0.0 | $2^{16} - 2$ | 2^{14} |
| Class C | 192 — 223 | 110 XXXXX | 192.0.0.0-223.255.255.255 | 255.255.255.0 | $2^8 - 2$ | 2^{21} |
| Class D (Multicast) | 224 — 239 | 1110 XXXX | 224.0.0.0-239.255.255.255 | | | |
| Class E (Experimental) | 240 — 255 | 1111 XXXX | 240.0.0.0-255.255.255.255 | | | |

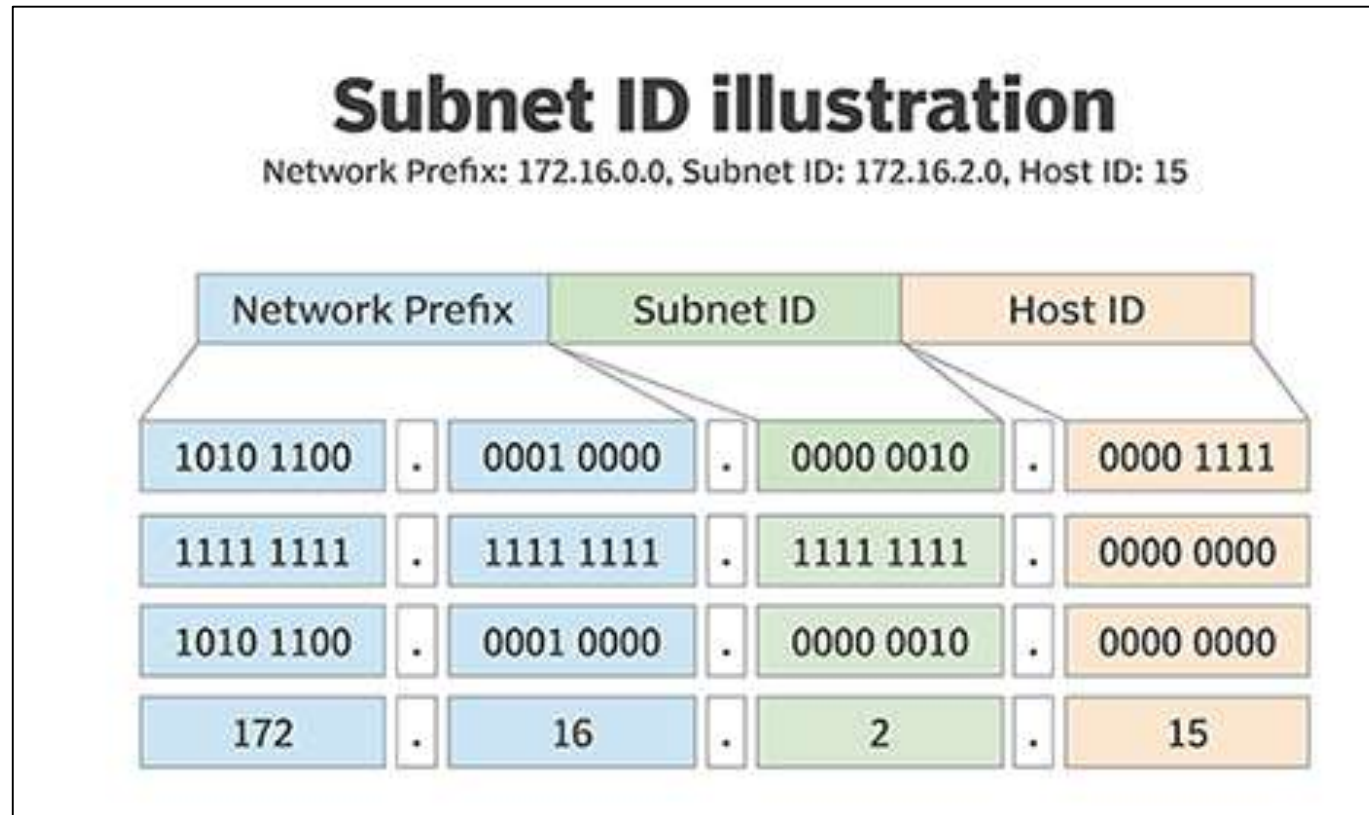
SUBNETTING

Subnetting is to take some of “Higher order host bits” in a network number and use them to create more network. These networks are called as **subnets**.

Subnet mask is a 32-bits long number which decides the length of the network.

The 0's and 1's in the subnet mask should be contiguous.

Valid subnet masks are; 0, 128, 192, 224, 252, 254, 255.



IP ADDRESS PLANNING

1. Determine the network and host requirements (Max # of hosts required / exist, Max # of segments required / exist)
2. Satisfy Host and Network Requirements
(#_of_Networks: 2^S , where S =subnet bits, #_of_hosts= 2^H-2 , where H =host bits, Total_#_of_Host_Bits= $S+H$)
3. Determine the Subnet Mask ($A=8$, $B=16$, $C=24 + S$, e.g. for a C-Class Network $24+S$)
4. Determine the Network Addresses
5. Determine the Directed Broadcast Addresses
6. Determine the Host Addresses

Examples:

1. You are given a Class C network (192.168.1.0) and you have four segments in your network, where the largest segment has 50 hosts. What subnet mask should you use and what is the layout of your addresses?
2. You are given a Class B network (172.16.0.0) and you have 490 segments in your network, where the largest segment needs 112 host addresses. What subnet mask should you use and what is the layout of your addresses?
3. You are given a Class A network (10.0.0.0) and you have 9000 segments in your network, where the largest segment needs 560 host addresses. What subnet mask should you use and what is the layout of your addresses?

| SUBNETTING TABLE | | | | | | | | |
|------------------|------|-----|-----|-----|-----|-----|-----|-----|
| | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |
| Bits borrowed | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Subnet Mask | 128 | 192 | 224 | 240 | 248 | 252 | n/a | n/a |
| /Mask | /25 | /26 | /27 | /28 | /29 | /30 | n/a | n/a |
| Wildcard Masks | .127 | .63 | .31 | .15 | .7 | .3 | n/a | n/a |
| Networks* | 2 | 4 | 8 | 16 | 32 | 64 | n/a | n/a |
| Hosts | 126 | 62 | 30 | 14 | 6 | 2 | n/a | n/a |

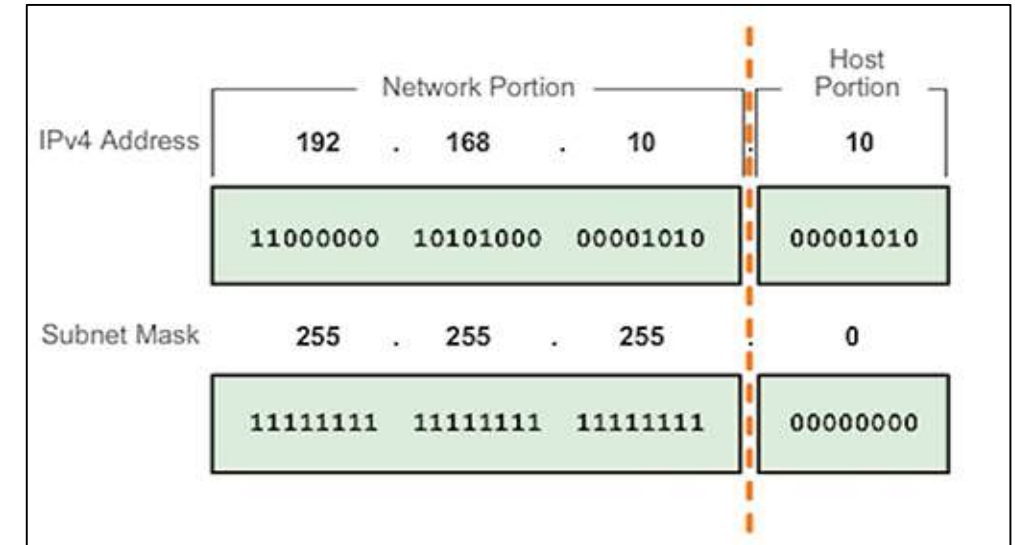
DETERMINING IP ADDRESS COMPONENTS

1. An IP Address and Subnet Mask required
2. Examine the subnet mask and find interesting OCTET
3. Subtract the interesting OCTET from 256
4. Write down network addresses
5. Write down broadcast addresses
6. Write down host addresses

Examples:

Determine the network and address type:

1. 192.168.1.37/27
2. 172.16.5.0/23
3. 192.16.1.63/29





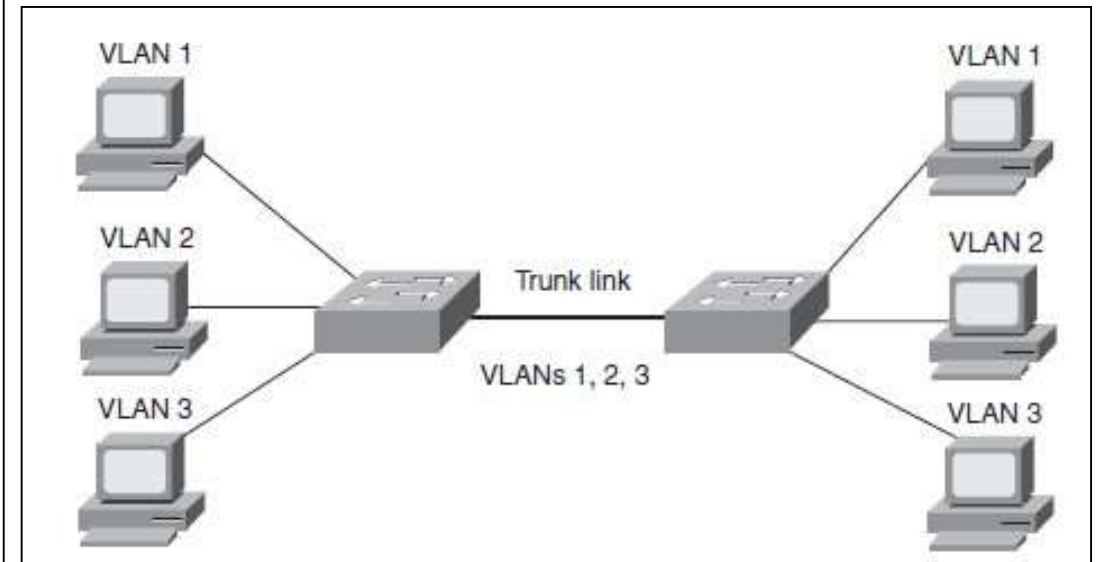
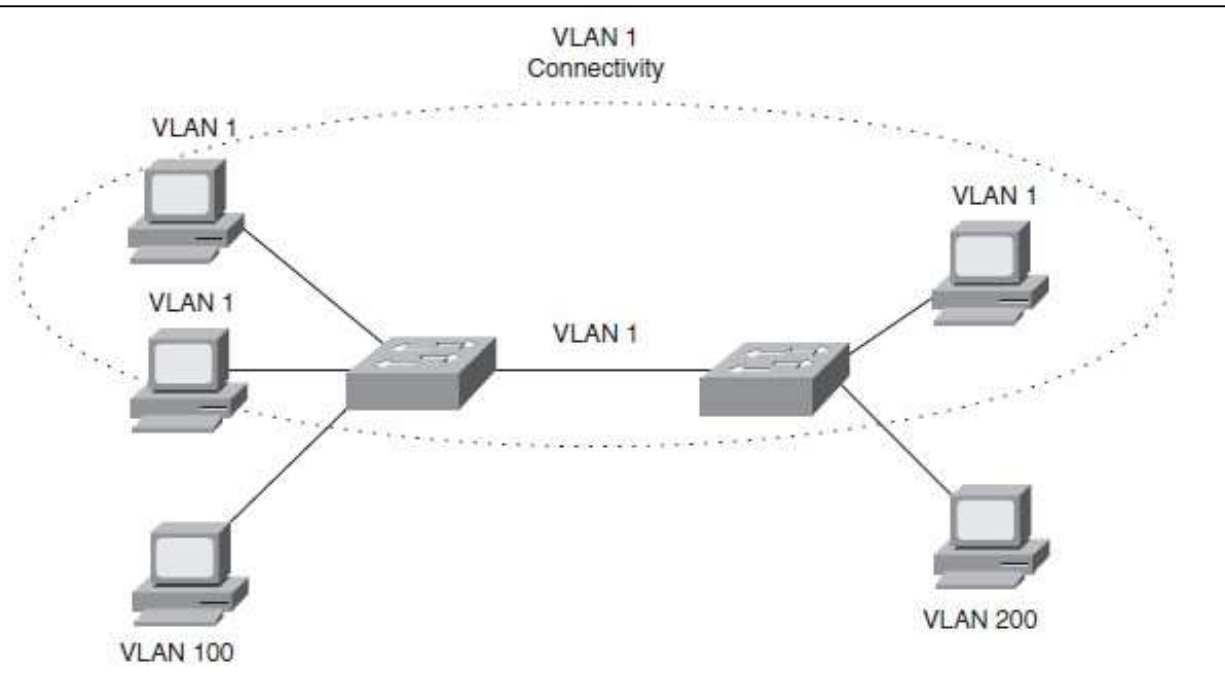
VLANs & Trunks

ACCESS AND TRUNK PORTS

VLAN (Virtual LAN) is a logical separation, which is used to create broadcast domains in Layer-2.

Access Ports are assigned to hosts and allowed to pass only one VLAN.

Trunk Ports are interfaces between switches and can be configured to pass multiple VLANs.

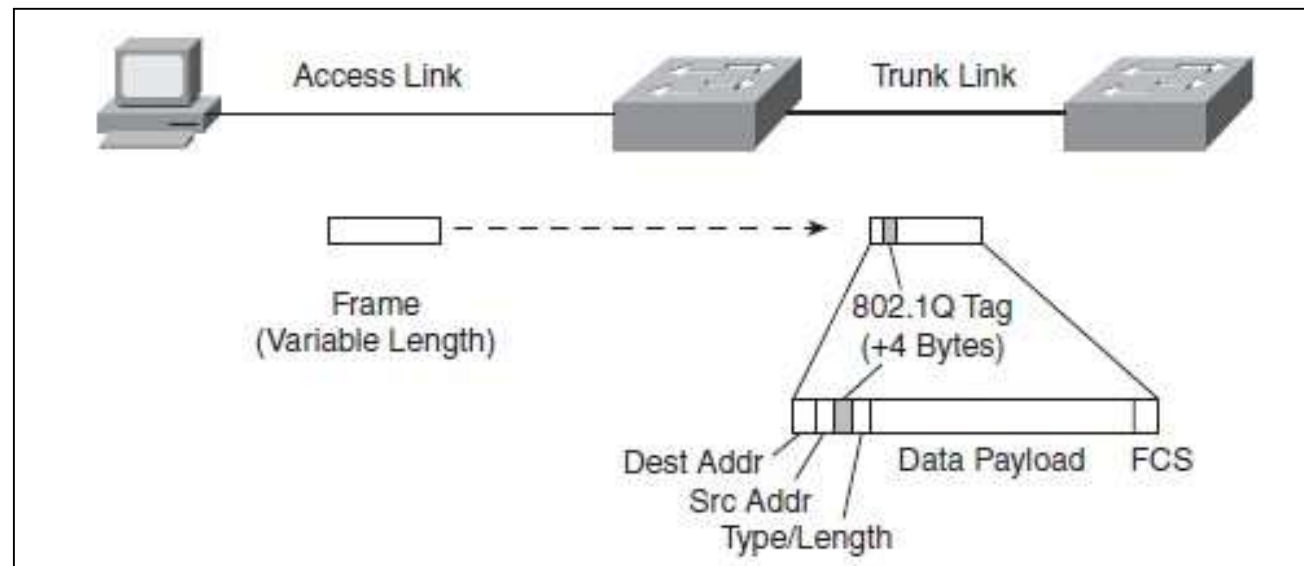


802.1Q ENCAPSULATION

IEEE standard for marking packets (a.k.a frame identification).

If frames must be transported out another trunk link, the **VLAN identifier** is added back into the frame header. Otherwise, if frames are destined out an access (non-trunk) link, the switch removes the VLAN identifier before transmitting the frames to the end station. Therefore, all traces of VLAN association are hidden from the end station.

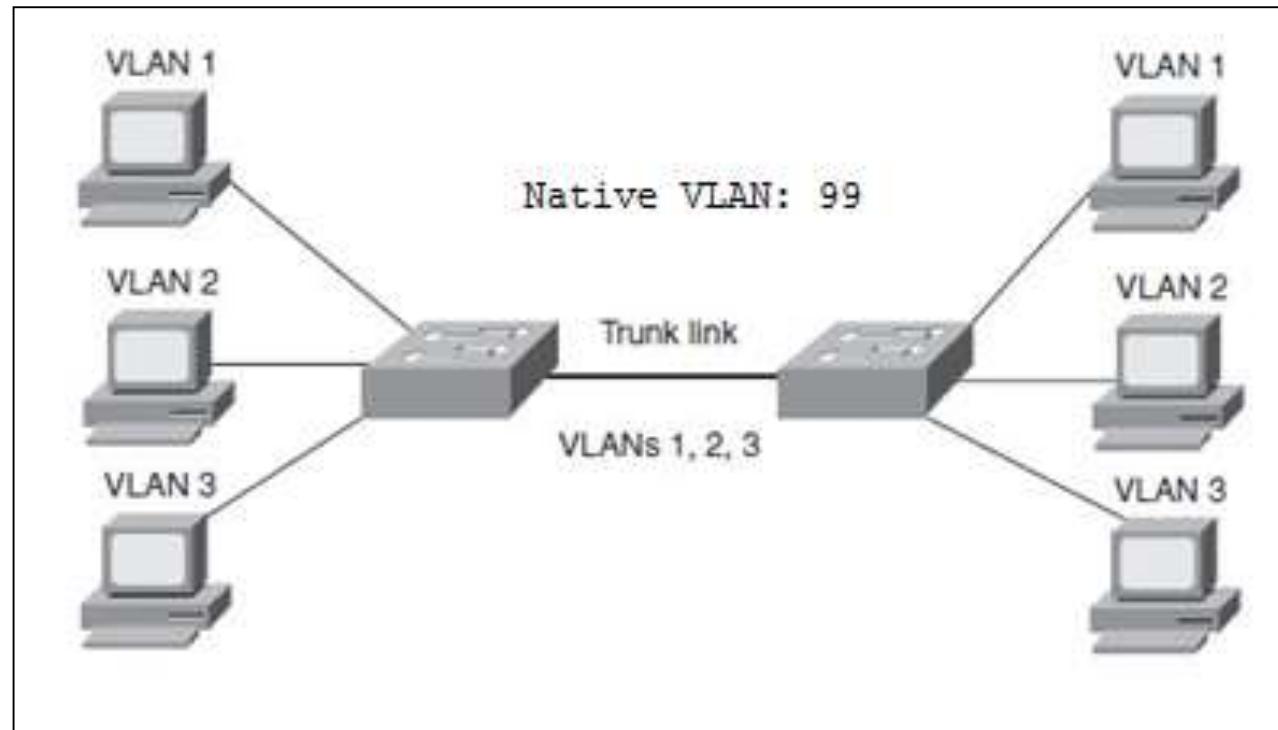
The first **two bytes** are used as a **Tag Protocol Identifier (TPID)** and always have a value of 0x8100 to signify an 802.1Q tag. The remaining two bytes are used as a Tag Control Information (TCI) field. The TCI information contains a **three-bit Priority field**, which is used to implement class-of-service (CoS) functions in the accompanying 802.1Q/802.1p prioritization standard. **One bit of the TCI** is a Canonical Format Indicator (CFI), flagging whether the MAC addresses are in Ethernet or Token Ring format. The last **12 bits** are used as a **VLAN identifier (VID)** to indicate the source VLAN for the frame. The VID can have values from 0 to 4095, but VLANs 0, 1, and 4095 are reserved.



NATIVE VLAN

Frames belonging to **Native VLAN** are *not* encapsulated with any tagging information. If an end station is connected to an **802.1Q** trunk link, the end station can receive and understand only the native VLAN frames. This provides a simple way to offer full trunk encapsulation to the devices that can understand it, while giving normal-access stations some inherent connectivity over the trunk.

- If there is a Native VLAN mismatch between switches, then switch communication via BPDU may have problems!

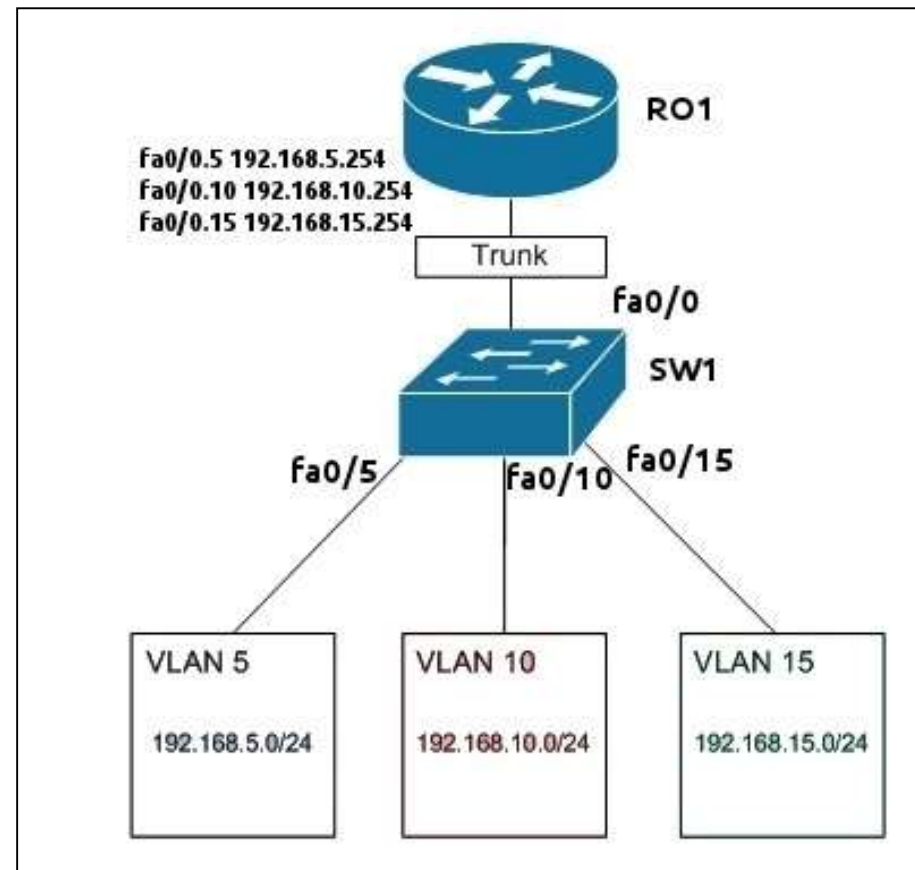


INTERVLAN COMMUNICATION

VLANs require a **L3 device** such as Router or L3 switch to communicate with each other.

Router specific **sub-interface configuration** is called as “**router-on-a-stick**”.

If **L3 switches** are used, then an **SVI** needs to be defined for each VLAN.





Redundancy & HA Solutions

MLT (A.K.A ETHERCHANNEL)

Bundling **parallel links into a single, logical link** to increase the bandwidth and overcome STP.

Provides **redundancy** with several **bundled physical links**. If one of the links within the bundle fails, traffic sent through that link automatically is moved to an adjacent link. Failover occurs in less than a few milliseconds and is transparent to the end user.

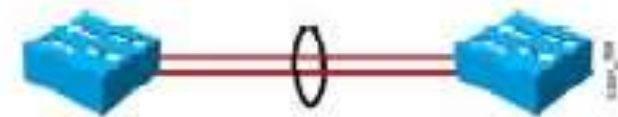
Ports should be on the same Ethernet media type and speed as well as the same STP settings.

Generally, all bundled ports first must belong to the same VLAN. If used as a trunk, bundled ports must be in **trunking mode**, have the same native VLAN, and pass the same set of VLANs.

Load balancing method:

| method Value | Hash Input | Hash Operation |
|--------------|-----------------------------------|----------------|
| src-ip | Source IP address | bits |
| dst-ip | Destination IP address | bits |
| src-dst-ip | Source and destination IP address | XOR |
| src-mac | Source MAC address | bits |
| dst-mac | Destination MAC address | bits |
| src-dst-mac | Source and destination MAC | XOR |
| src-port | Source port number | bits |
| dst-port | Destination port number | bits |
| src-dst-port | Source and destination port | XOR |

- Same speed and duplex.
- Same mode (access or trunk).
- Same native and allowed VLANs on trunk ports.
- Same access VLAN on access ports.
- Configure these parameters on the port-channel interface.



LACP (LINK AGGREGATION CONTROL PROTOCOL)

LACP is an IEEE 802.3ad standard, **MLT control protocol**.

LACP packets are exchanged between switches over MLT-capable ports.

Ports are selected and become active according to their *port priority* value (a 2-byte priority followed by a 2-byte port number), where a low value indicates a higher priority.

The lowest port priorities as active MLT and the other links are placed in a standby state and will be enabled in the MLT if one of the active links goes down.

| Mode | Negotiation Packets Sent? | Characteristics |
|-------------|---------------------------|---------------------------------|
| LACP | | |
| On | No | All ports channeling |
| Passive | Yes | Waits to channel until asked |
| Active | Yes | Actively asks to form a channel |

HIGH AVAILABILITY AND VRRP (VIRTUAL ROUTER REDUNDANCY PROTOCOL)

VRRP provides **L3 redundancy** on multilayer switches.

VRRP provides one redundant gateway address from a group of routers. The active router is called the **master router**, whereas all others are in the **backup state**.

The highest prioritized router becomes master, the default priority value is 100.

VRRP sends its advertisements to the multicast destination address 224.0.0.18 (All VRRP Routers), using IP protocol 112.

A virtual IP & MAC Address is created for each VRRP group. The virtual IP Address can be assigned to a VRRP Cluster. Total 1000 IP Addresses can be defined as VRRP Virtual IP Address.

The virtual router MAC address is of the form 0000.5e00.01**xx**, where xx is a two-digit hex VRRP group number.

VRRP advertisements are sent at **1-second intervals**. Backup routers optionally can learn the advertisement interval from the master router. The failover time is the **3 times of Adv. time**.

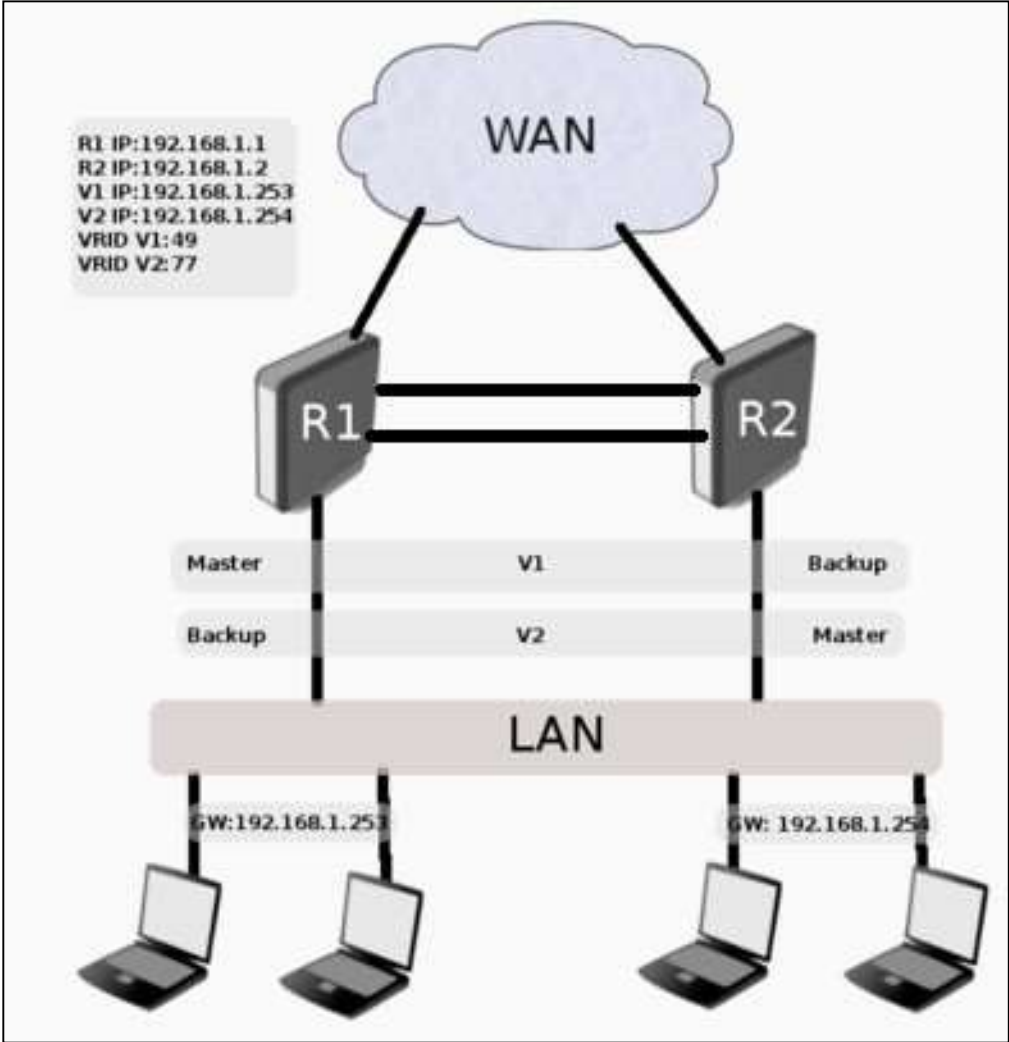
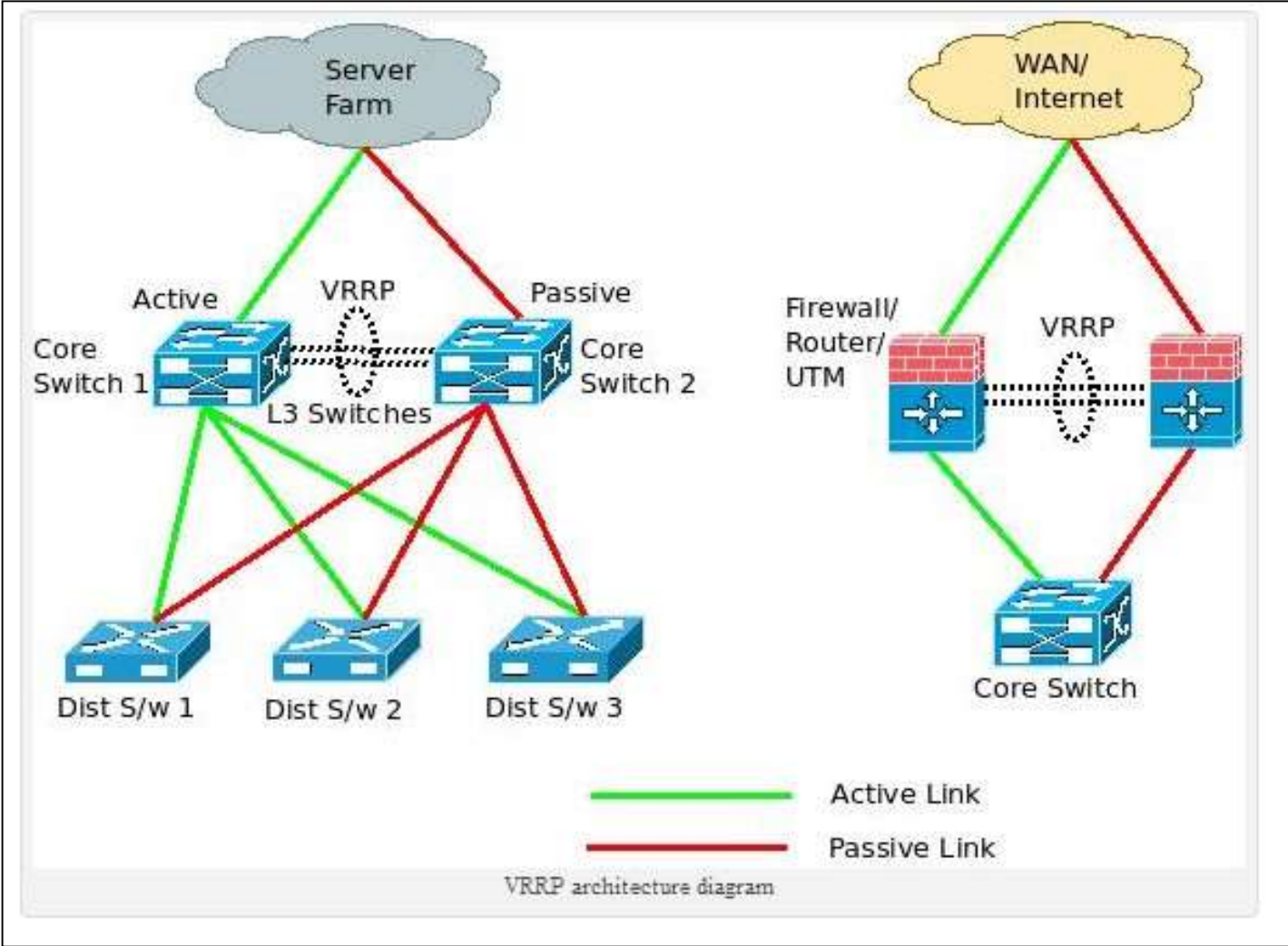
By default, all VRRP routers are configured to **preempt** the current master router if their priorities are greater.

VRRP can **track objects** to determine if the critical paths still be reachable, otherwise performs a failover.

Load balancing can be performed per VRRP Group basis.

VRRP and STP should complete each other with design perspective.

HIGH AVAILABILITY AND VRRP (VIRTUAL ROUTER REDUNDANCY PROTOCOL)



STP ENHANCEMENTS; MST

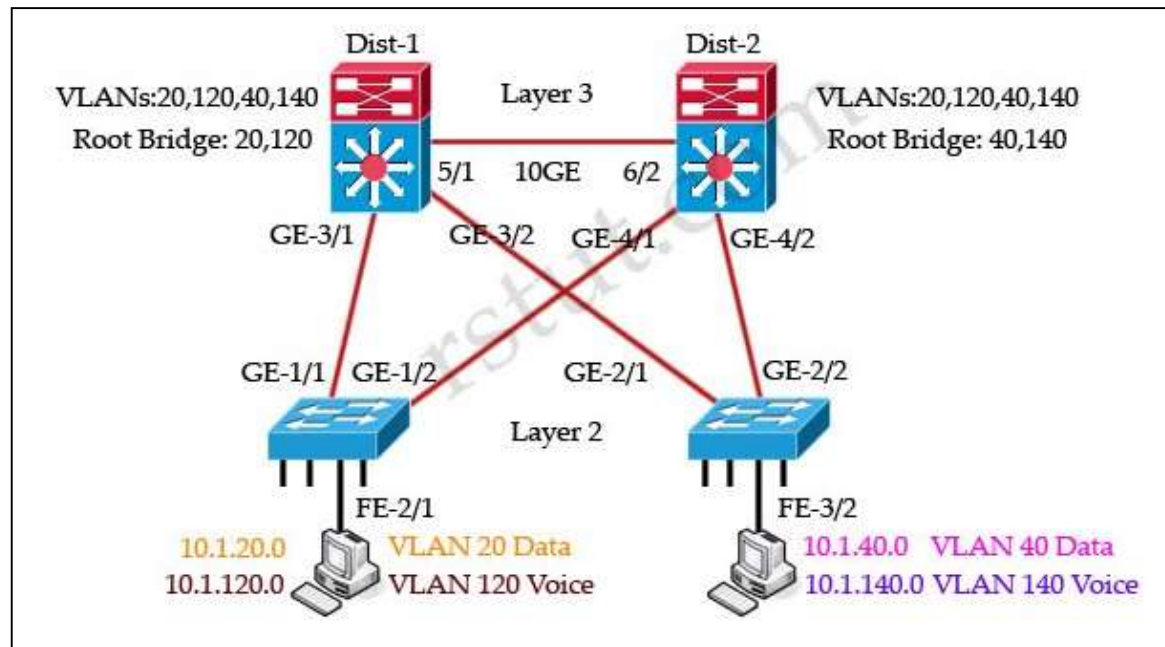
MST (IEEE 802.1s) is built on the concept of mapping one or more VLANs to a single STP instance. Multiple instances of STP can be used (hence the name MST), with each instance supporting a different group of VLANs.

MST is working with regions. If two switches have the same set of attributes, they belong to the same MST region. If not, they belong to two independent regions.

MST configuration name (32 characters)

MST configuration revision number (0 to 65535)

MST instance-to-VLAN mapping table (4096 entries)





Routing Protocols

STATIC ROUTING

A static route is a manually configured route on your router. Static routes are typically used in smaller networks and when few networks or subnets exist, or with WAN links that have little available bandwidth.

With a network that has hundreds of routes, static routes are not scalable, since you would have to configure each route and any redundant paths for that route on each router.

```
HQ#
HQ#show ip route
Codes: L - local, C - connected, S - static, R - RIP, M - mobile, B - BGP
       D - EIGRP, EX - EIGRP external, O - OSPF, IA - OSPF inter area
       N1 - OSPF NSSA external type 1, N2 - OSPF NSSA external type 2
       E1 - OSPF external type 1, E2 - OSPF external type 2, E - EGP
       i - IS-IS, L1 - IS-IS level-1, L2 - IS-IS level-2, ia - IS-IS inter area
       * - candidate default, U - per-user static route, o - ODR
       P - periodic downloaded static route

Gateway of last resort is 10.0.2.2 to network 0.0.0.0

    10.0.0.0/8 is variably subnetted, 7 subnets, 2 masks
C       10.0.2.0/24 is directly connected, GigabitEthernet0/2
L       10.0.2.1/32 is directly connected, GigabitEthernet0/2
C       10.0.3.0/24 is directly connected, GigabitEthernet0/0
L       10.0.3.1/32 is directly connected, GigabitEthernet0/0
C       10.10.0.0/24 is directly connected, Loopback0
L       10.10.0.1/32 is directly connected, Loopback0
S       10.11.0.0/24 [1/0] via 10.0.3.2
S*    0.0.0.0/0 [2/0] via 10.0.2.2
HQ#
```

RIPv1

RIPv1 uses **local broadcasts** (255.255.255.255) to share routing information.

These **updates are periodic** in nature, occurring, by default, every 30 seconds, with a hold-down period of 180 seconds.

RIP use **hop count** as a metric.

RIPv1 is a **classful** protocol.

RIP supports up to **six equal-cost paths** to a single destination, where all six paths can be placed in the routing table and the router can **load-balance** across them. The default is actually four paths, but this can be increased up to a maximum of six. RIP will not load-balance across *unequal-cost* paths.

```
Router#show ip route
Codes: C - connected, S - static, I - IGRP, R - RIP, M - mobile, B - BGP
       D - EIGRP, EX - EIGRP external, O - OSPF, IA - OSPF inter area
       N1 - OSPF NSSA external type 1, N2 - OSPF NSSA external type 2
       E1 - OSPF external type 1, E2 - OSPF external type 2, E - EGP
       i - IS-IS, L1 - IS-IS level-1, L2 - IS-IS level-2, ia - IS-IS inter area
       * - candidate default, U - per-user static route, o - ODR
       P - periodic downloaded static route

Gateway of last resort is not set

C    192.168.1.0/24 is directly connected, FastEthernet0/0
R    192.168.2.0/24 [120/1] via 200.1.1.2, 00:00:15, FastEthernet0/1
R    192.168.3.0/24 [120/1] via 200.1.1.3, 00:00:03, FastEthernet0/1
C    200.1.1.0/24 is directly connected, FastEthernet0/1
```

LAN2

LAN3

RIPV2

Classless, hybrid

Uses **multicast** IP Address (224.0.0.9) for routing updates with triggered updates to speed up convergence

Metric is the "**hop-count**" with the maximum limit of 15

Auto-summarization is enabled by default

Supports **authentication** (Clear-text or hashed-key)

Supports **equal load balancing** as the same with RIPv1

Protocol Number: 17

```
R1#show ip route
Codes: C - connected, S - static, I - IGRP, R - RIP, M - mobile, B - BGP
       D - EIGRP, EX - EIGRP external, O - OSPF, IA - OSPF inter area
       N1 - OSPF NSSA external type 1, N2 - OSPF NSSA external type 2
       E1 - OSPF external type 1, E2 - OSPF external type 2, E - EGP
       i - IS-IS, L1 - IS-IS level-1, L2 - IS-IS level-2, ia - IS-IS inter area
       * - candidate default, U - per-user static route, o - ODR
       P - periodic downloaded static route

Gateway of last resort is not set

    172.30.0.0/16 is variably subnetted, 6 subnets, 2 masks
R       172.30.200.32/28 [120/2] via 209.165.200.229, 00:00:01, Serial0/0/0
R       172.30.200.16/28 [120/2] via 209.165.200.229, 00:00:01, Serial0/0/0
C 172.30.1.0/24 is directly connected, FastEthernet 0/0
C 172.30.2.0/24 is directly connected, FastEthernet 0/1
R       172.30.100.0/24 [120/2] via 209.165.200.229, 00:00:01, Serial0/0/0
R       172.30.110.0/24 [120/2] via 209.165.200.229, 00:00:01, Serial0/0/0
```

OSPF

Open standard, **classless**, **link-state** protocol

Uses **SPF** (Dijkstra / Shortest Path First) algorithm for a **loop free** topology

Uses **incremental**, **triggered**, **multicast** [224.0.0.5 (OSPF All Routers), 224.0.0.6 (OSPF DRs)] LSAs to update the routing table

Uses **inverse of the bandwidth** as a cost value ($10^8/\text{BW}$ where BW: 56Kbps Serial: 1785, 64Kbps Serial: 1652, T1: 64, Ethernet: 10, Fast Eth.: 1, FDDI: 1)

Uses **load balancing** up to 16 equal-cost paths

AS (Autonomous System) numbers are used to separate OSPF areas

Protocol Number: 89

Requires more memory, CPU. Complex to troubleshoot. Large network disadvantage (use BGP instead)

```
R1#show ip route | begin Gateway
Gateway of last resort is 0.0.0.0 to network 0.0.0.0

    172.16.0.0/24 is subnetted, 6 subnets
O       172.16.144.0 [110/2] via 10.1.124.4, 00:52:01, FastEthernet1/0
O       172.16.133.0 [110/65] via 10.1.13.3, 00:38:18, Serial0/1
O       172.16.104.0 [110/2] via 10.1.124.4, 00:52:01, FastEthernet1/0
C       172.16.101.0 is directly connected, Loopback0
O       172.16.102.0 [110/129] via 10.1.13.3, 00:38:18, Serial0/1
O       172.16.103.0 [110/65] via 10.1.13.3, 00:38:18, Serial0/1
    10.0.0.0/24 is subnetted, 3 subnets
C       10.1.13.0 is directly connected, Serial0/1
O       10.1.23.0 [110/128] via 10.1.13.3, 00:38:18, Serial0/1
C       10.1.124.0 is directly connected, FastEthernet1/0
S*    0.0.0.0/0 is directly connected, Serial0/1
R1#
```

OSPF

Operation:

Unique Router ID: The active loopback with the highest IP address or the active physical interface with the highest IP address.

Neighboring: LSA Hello message sends every 10secs. Dead Time is 4 times of Hello (40 secs). If a router does not send Hello LSA within the dead time, then declared as dead.

The following should match for the routers to become neighbors:

- The area number
- Hello and Dead timer intervals
- OSPF password (if configured)
- Area Stub Flag
- MTU Size

During learning phase (exchange process), the following states are seen;

- Down State: Not exchanged any LSA with any router
- Init State: LSA Hello is received and added to neighbor database, unidirectional communication
- Two-way State: LSA Reply from the destination and update the neighbor database. DR / BDR election starts.

Operation:**DR / BDR Election: Client / Server implementation to OSPF**

- The router with the highest Router ID becomes DR, second becomes BDR (Priority 0 means, not involve to election)
- DR + BDR called as DROTHERs. Each broadcast segment has the DROTHERs except WAN PPP.
- Routers send LSAs to DR directly (via multicast 224.0.0.6) and DR sends to all (via multicast 224.0.0.5). PPP sends to 224.0.0.5 only.

Routing Information Sharing: The following states are seen after the DR / BDR election:

- Exstart State: The DR is selected and started the exchange of routing database, others stays as slaves
- Exchange State: DR sends DBD/DDP (Database Description Packets) which contains the link state type, ID of advertising router, cost of advertising router and the sequence number of the link. Slave responds with LSACK and compares the DBD with its own database.
- Loading State: If the DR has more up-to-date database than slave, the slave will respond to DR's original DBD with a LSR (Link State Request). The DR sends a LSU (Link State Update) and slave generates LSACK again. If the slave has more up-to-date information, it'll repeat the exchange and loading states.
- Full State: DR and others are synch'ed.

BGP

BGP is an **exterior gateway protocol** (EGP), which means that it performs routing between multiple autonomous systems or domains and exchanges routing and reachability information with other BGP systems.

BGP uses 13-steps “**Best Path Selection**” algorithm to choose the best path.

Load balancing is supported but not recommended.

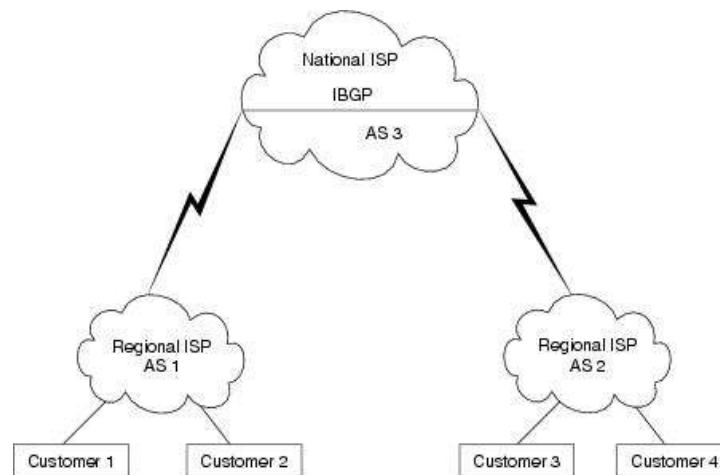
BGP is used on **Core layer routers**. Designed for **very-large scaled networks**, but convergence time is slower than OSPF.

BGP neighbors **exchange full routing information** when the TCP connection between neighbors is first established. When changes to the routing table are detected, the BGP routers send to their neighbors **only those routes that have changed**. BGP routers **do not send periodic routing updates**, and BGP routing updates advertise only the optimal path to a destination network.

If a full BGP table is requested (not recommended for small routers), you can see the **whole internetwork**.

BGP is **controlled globally**, and you should confirm that the IP Address is a member of **YOUR** network, before you announce the network through BGP!

BGP is controlled by ISPs with **ACLs**.



**BGP Best Path
Selection Algorithm**

COMPARISON OF ROUTING PROTOCOLS

| Name | Type | Update | Metric | VLSM | Summary |
|--------------|------------------|---------------|---------------|-------------|-----------------------|
| RIPv1 | DV | 30 sec | Hops | No | Automatic |
| RIPv2 | DV | 30 sec | Hops | Yes | Automatic |
| IGRP | DV | 90 sec | Composite | No | Automatic |
| EIGRP | Advanc ed. DV | triggered | Composite | Yes | Automatic + Manual |
| OSPF | LS | triggered | Cost | Yes | Manual |
| IS-IS | LS | triggered | Cost | Yes | Automatic |
| BGP | DV | triggered | N/A | N/A | Manual |

HOMEWORK

- Provide an architectural view of an Enterprise Network that contains;
 - A server farm to host Web Servers, E-Mail servers, CRM and ERP data.
 - Office users (3 floors, 150 users in each floor) that will use a VoIP phone and a PC at their desks.
 - An approach to access from home to allow “Work from Home” concept.
 - A lab network to access lab area.
 - Customers could reach to Web Servers.
 - Key points:
 - Ensure that local redundancy is maintained and there is no single point of failure in design.
 - Ensure that the security focus is maintained.
 - Ensure that all activities are monitored and logged for data regulations.
 - Ensure that minimum cost rule is followed.
 - Provide your answers with drawings and explanations of key points.

Thank You



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