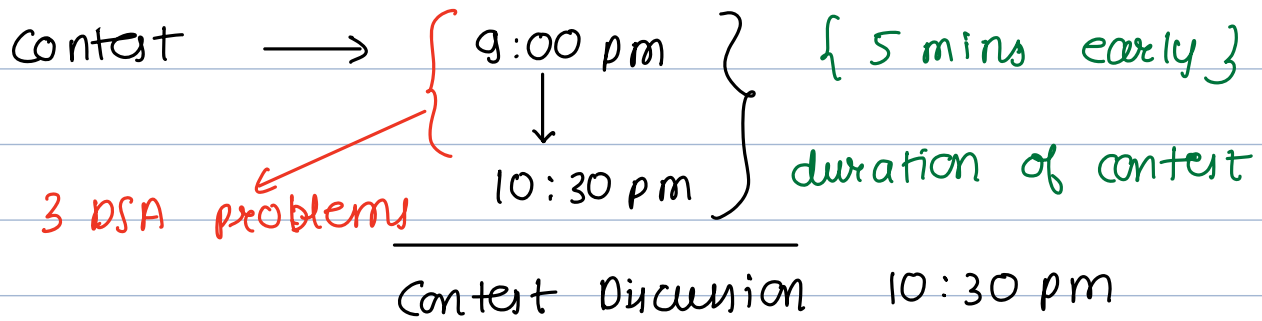
Recursion 2

Content

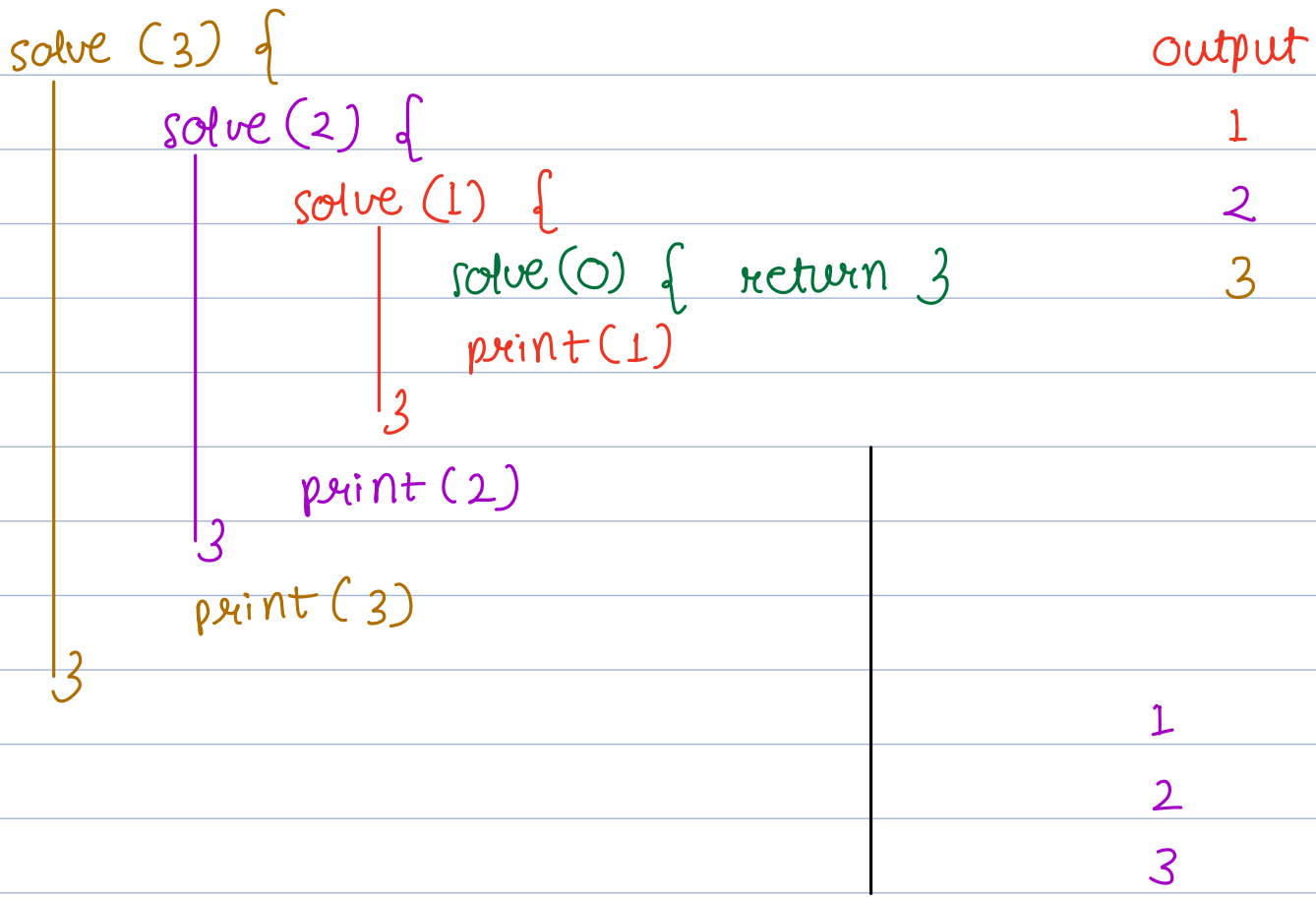
- Quizzes on Recursion
- Tower of Hanoi
- Balanced Parenthesis.



Contest Discussion 10:30 pm

What is the output of the following code for N = 3?

```
void solve(int N){
    if(N == 0)
        return;
    solve(N-1);
    print(N);
}
```



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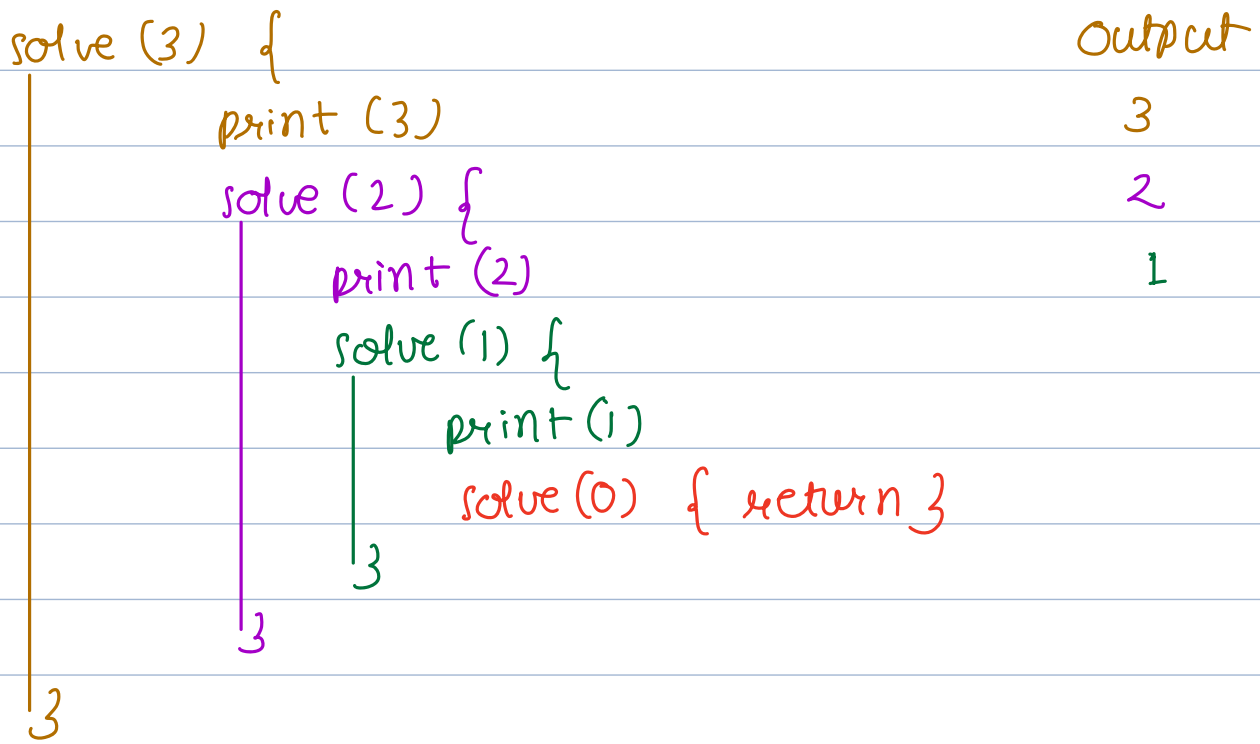


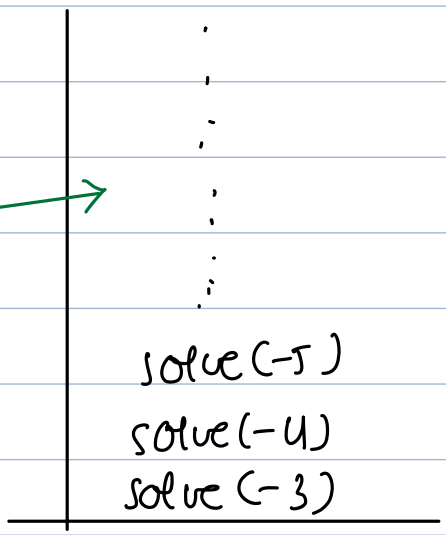
What is the output of the following code for N = -3?

```
void solve(int N){
    if(N == 0)
        return;
    print(N);
    solve(N-1);
}
```

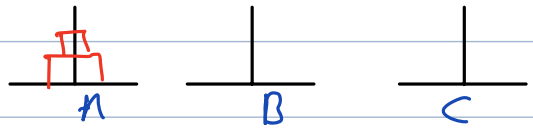
solve (-3) {
print (-3)
solve (-4) {
print (-4)
solve (-5) {
print (-5)

output
-3
-4
-5

stackoverflow



Tower of Hanoi ***



There are n disks placed on a tower A
{different sizes}

Move all disks from tower A to C using B

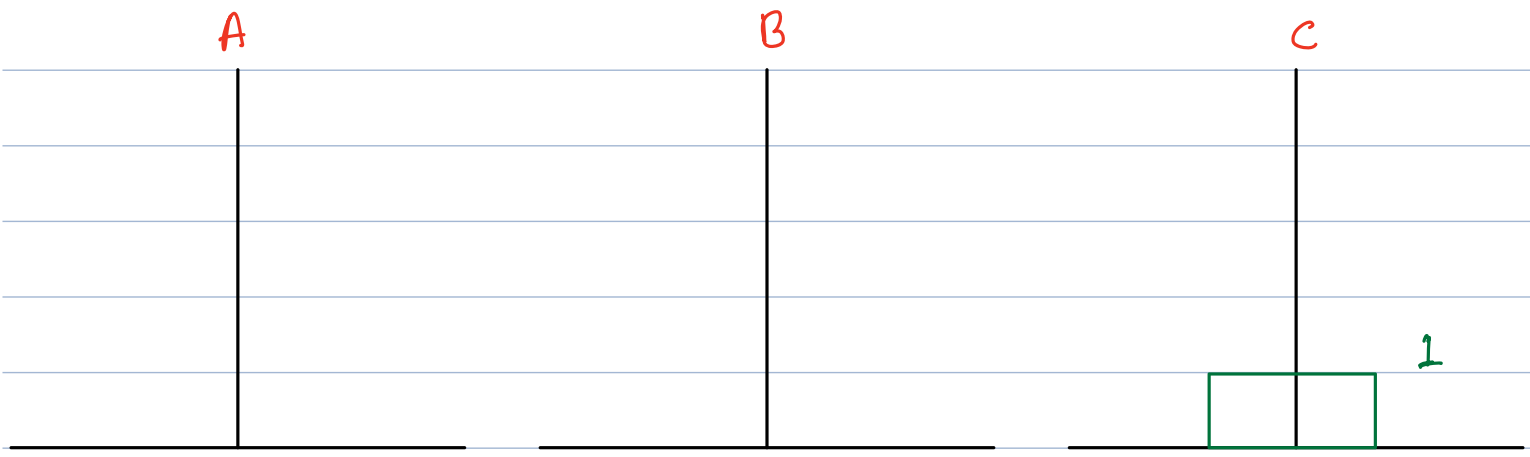
Constraints

— Only 1 disk can be moved at a time
— Larger disk cannot be placed on a small disk at any step.

Print the movement of disks from A to C in minimum steps.

Example

N = 1



Disk 1 A ⟶ C

N = 2

TOH (2, A, B, C)

Initial State

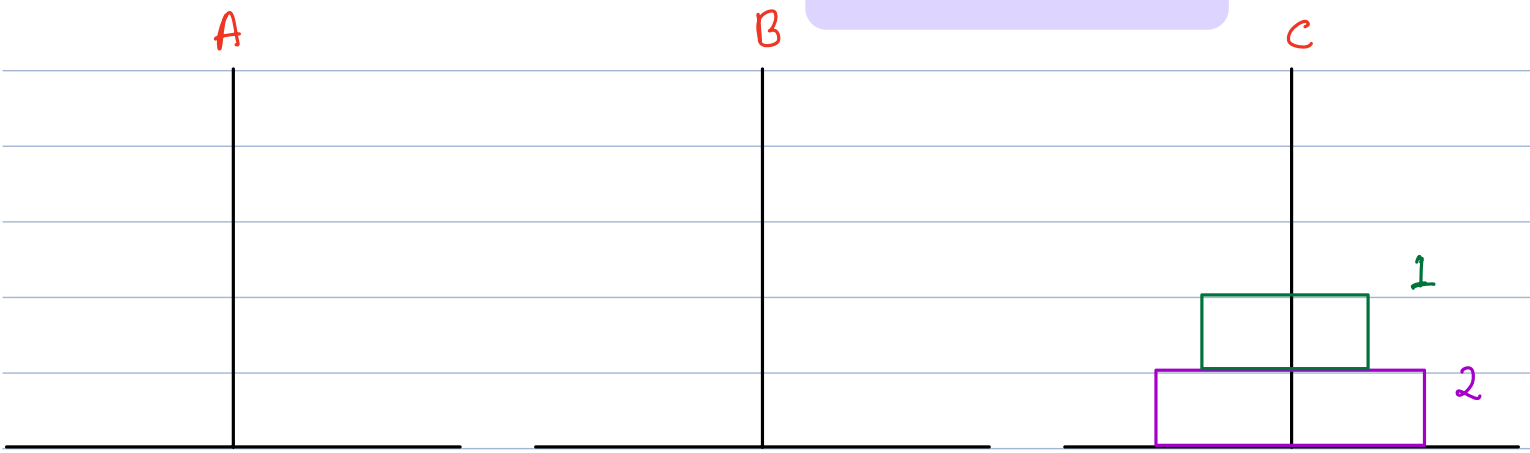
steps ⟶

D 1 A ⟶ B

D 2 A ⟶ C

D 1 B ⟶ C

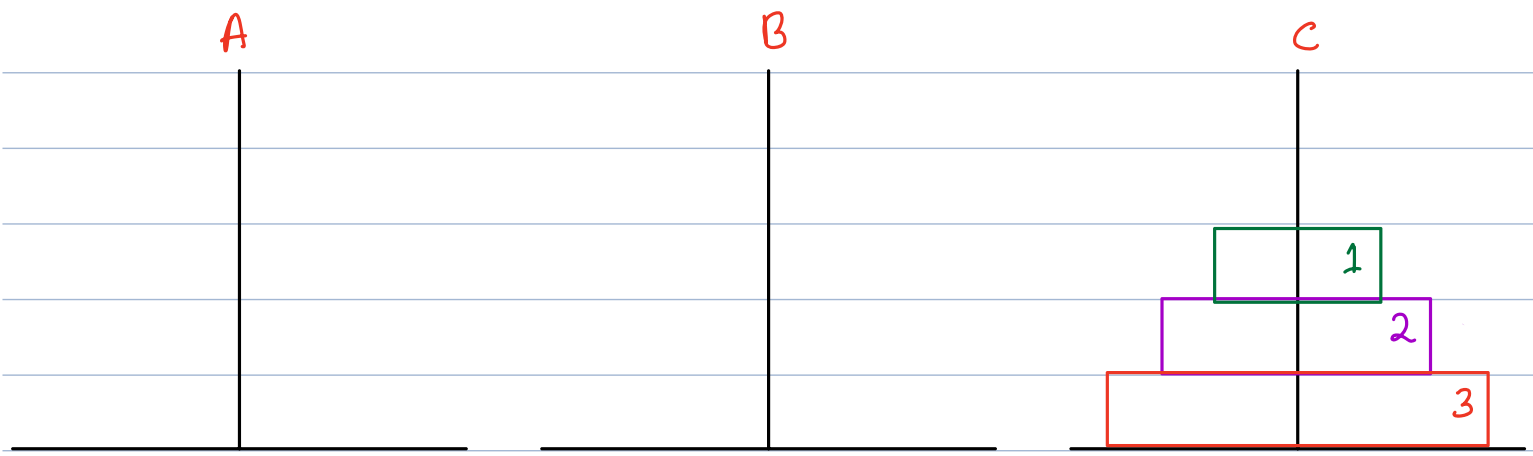
Final State



Assumption ⟶ TOH (N, A, B, C)

⟶ Move N disks from tower A to tower C using a spare tower B

N = 3



We moved 2 disks from tower A to tower B using C as spare tower.

spare

from to

TOH (2 , A , C , B)

D3 A ⟶ C

TOH (2 , B , A , C)

Steps

D1 A ⟶ C

D2 A ⟶ B

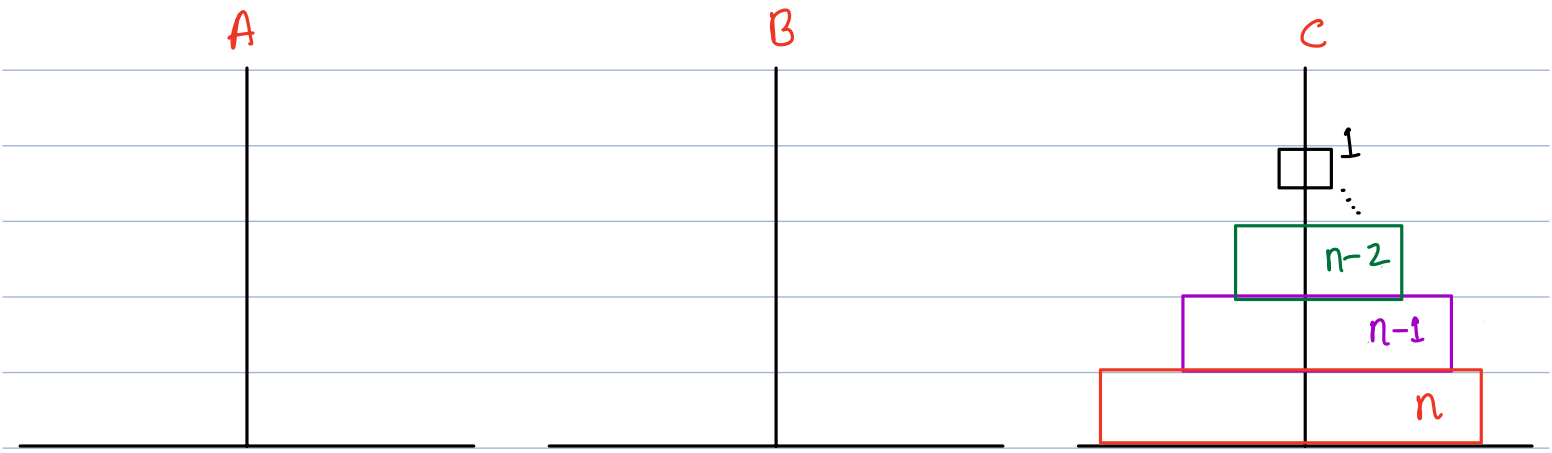
D1 C ⟶ B

D3 A ⟶ C

D1 B ⟶ A

D2 B ⟶ C

D1 A ⟶ C



fr Sp to

TOH (N , A , B , C) step 1

// Move N-1 disks from A to B using spare C step 2

TOH (N-1 , A , C , B)

print (D-N , A ⟶ C)

// Move N-1 disks from B to C using spare A

TOH (N-1 , B , A , C)

Base case

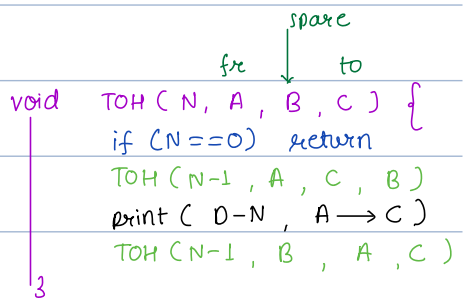
if N == 0 return ;

Pseudo code

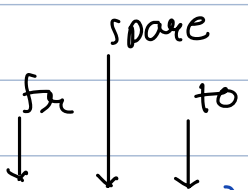
spare
fr to

```
void  TOH ( N , A , B , C ) {
        if (N==0) return
        TOH (N-1 , A , C , B )
        print ( D-N , A ⟶ C )
        TOH (N-1 , B , A , C )
}
```

N	steps	
1	1	2-1
2	3	4-1
3	7	8-1
4	7 1 7 = 15	16-1
5		32-1
⋮		
n	$2^n - 1$	



N = 2




```
TOH (2, A, B, C) {
    TOH ( 1 , A , C , B ) {
        TOH ( 0 , A , B , C ) { return }
        print ( 1 , A ⟶ B )
        TOH ( 0 , C , A , B ) { return }
    }
    print ( 2 , A ⟶ C )
```

```
TOH(1, B, A, C) {
    TOH(0, B, C, A) { return }
    print(1  B ⟶ C)
    TOH(0, A, B, C) { return }
}
```

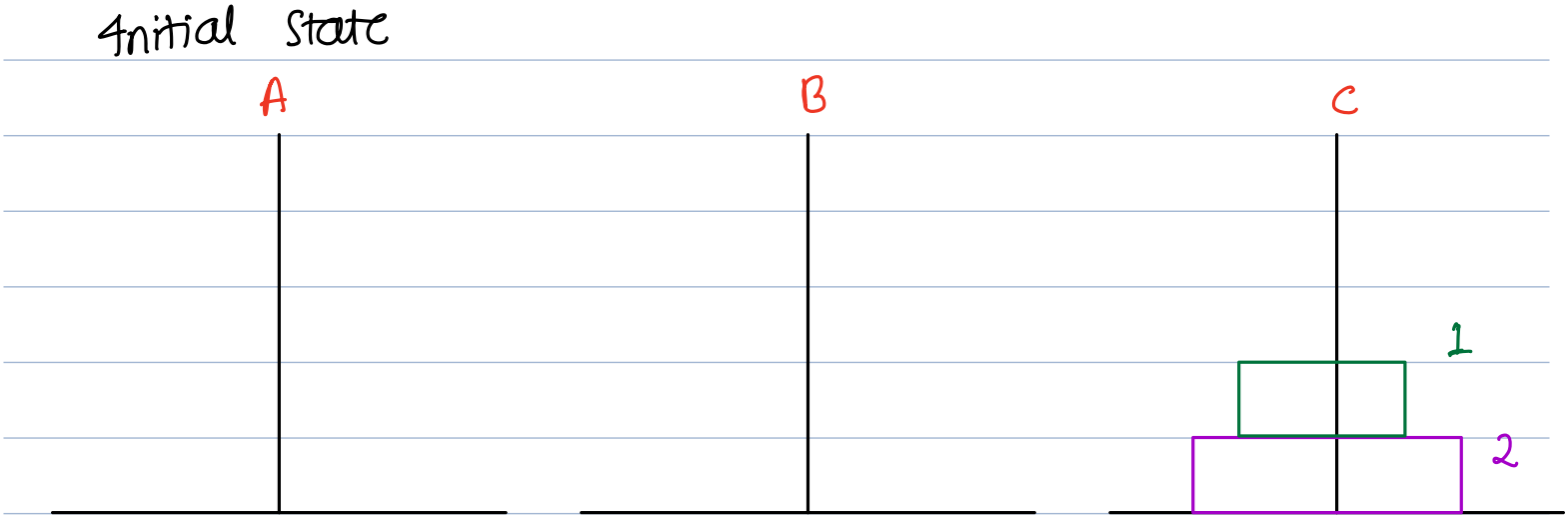
output

1 A ⟶ B

2, A ⟶ C

1 B ⟶ C

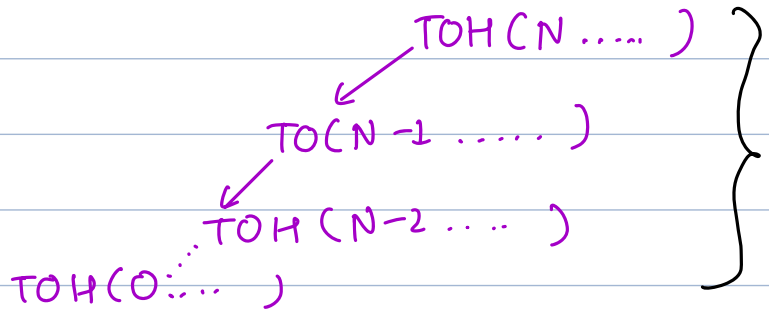
TOH(2, A, B, C)



Break : 22 : 34

TC : $O(2^N)$

SC : O(N) ⟶ recursive stack space.



Valid Parenthesis

Print all valid parenthesis of length 2*N for a given N

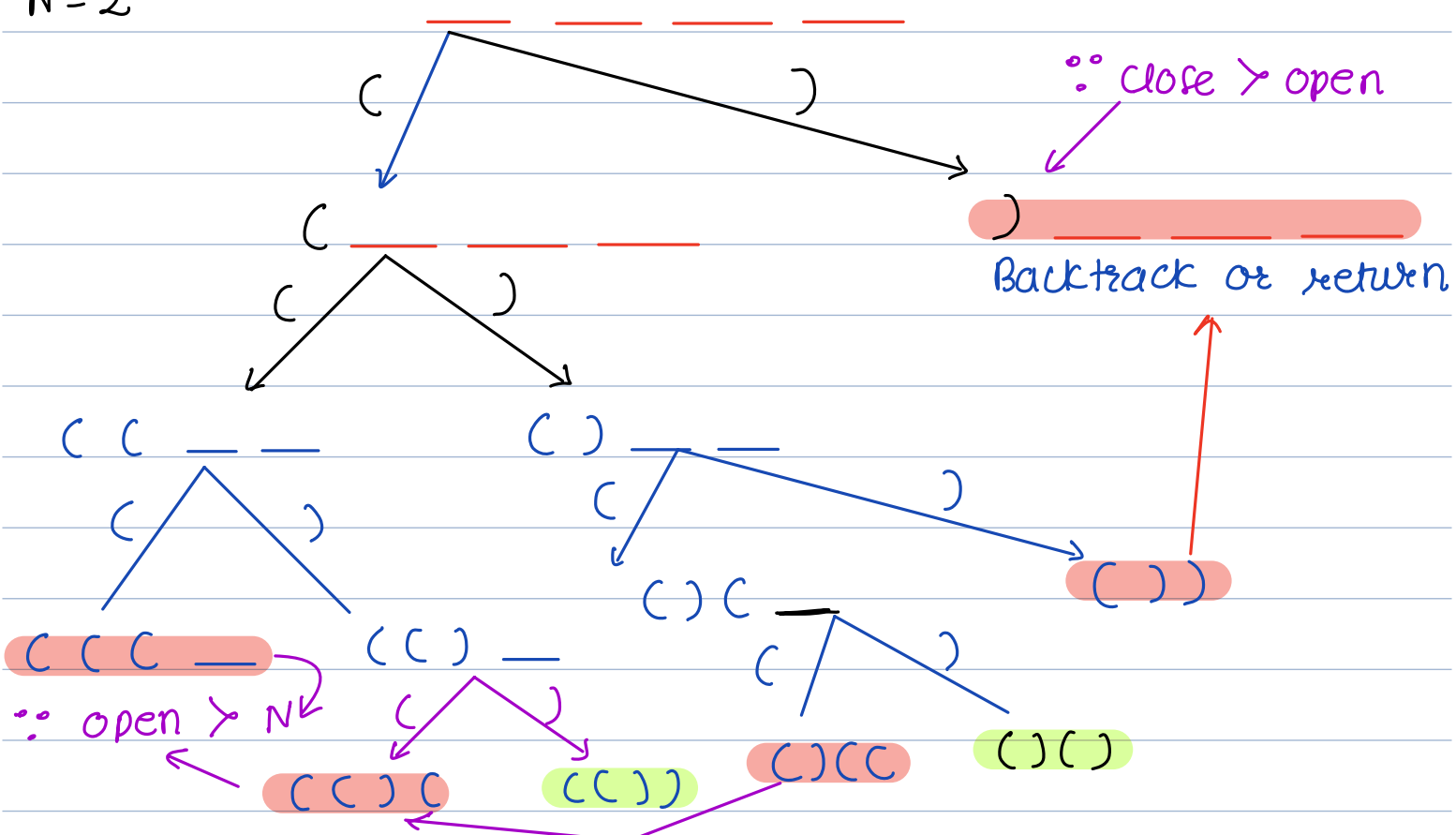
at any instant → open >= close } cond for valid parenthesis

open == close == N

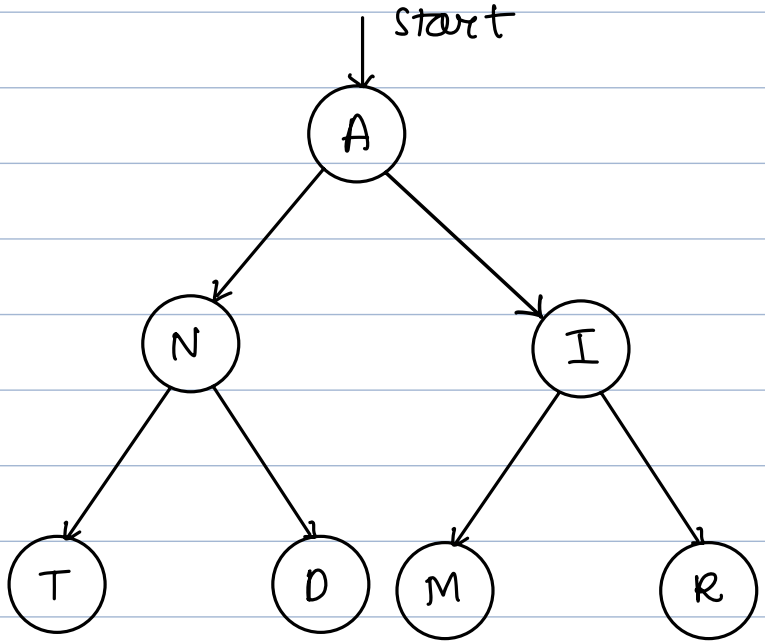
Example

	Output
N = 1	( )
N = 2	(( )) ( ) ( )
N = 3	((( ))) , ( ) (( )) , (( )) ( ) , (( ) ( )) ( ) ( ) ( )

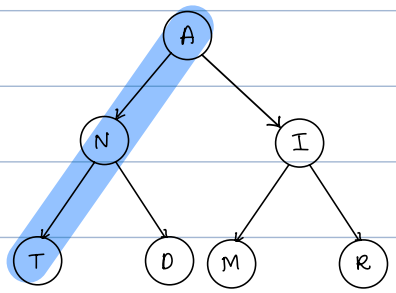
N = 2



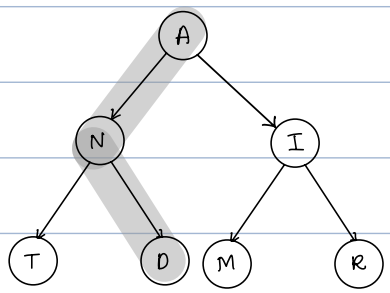
what is Backtracking ?



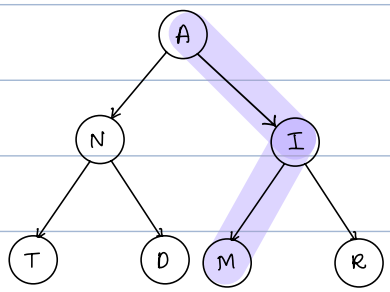
search for AIM



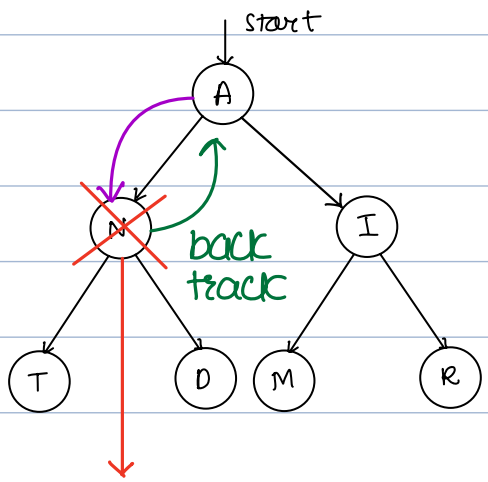
ANT != AIM



AND != AIM



AIM == AIM



Try out all possibilities to find a solution.
If sol" is not possible from a path we exit early { backtrack }

Backtrack is a technique

I can never find AIM, if I doesn't match with N

Pseudocode

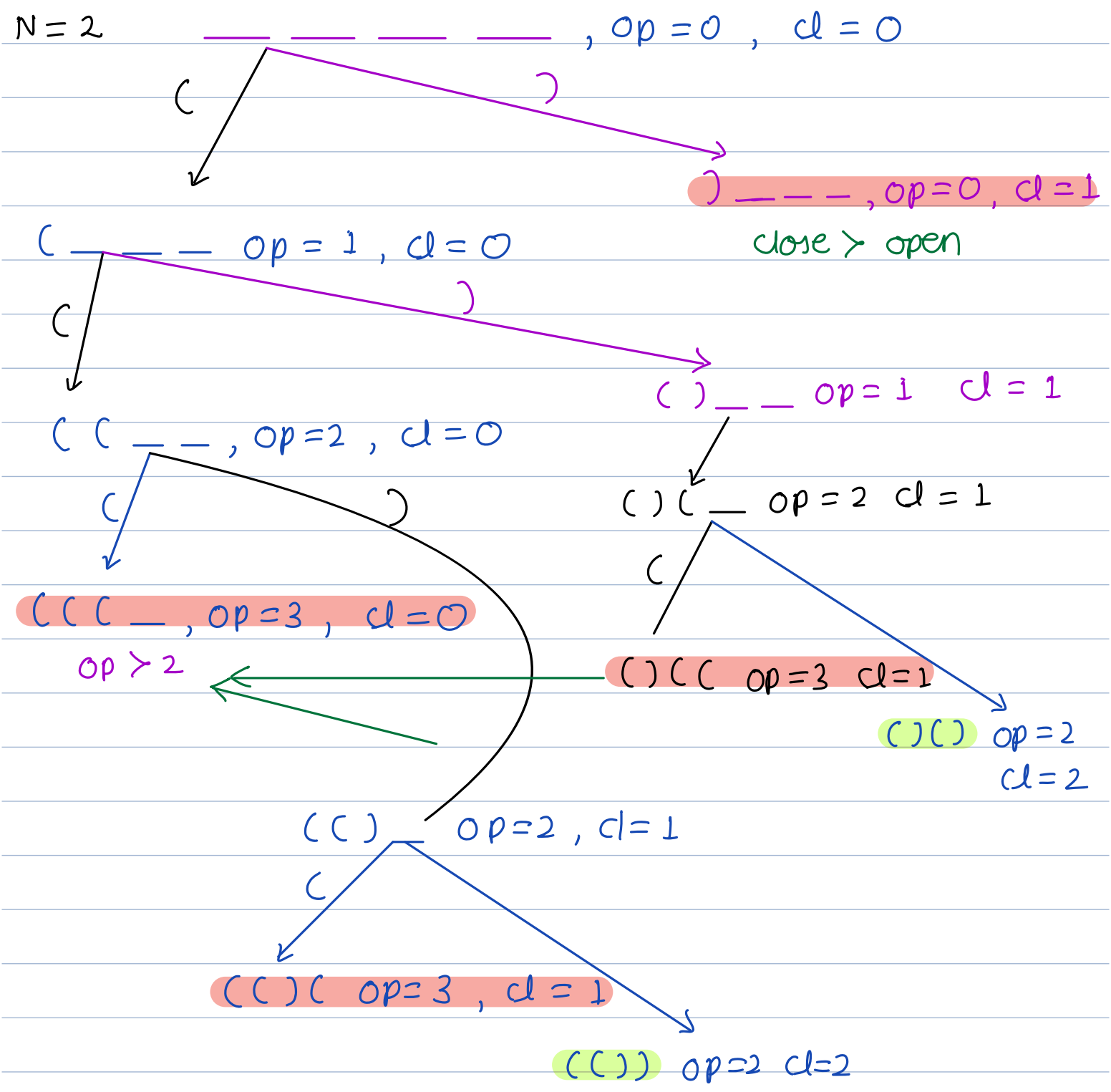
```
void valid (int N, String s, int open, int close) {
    // Base condition
    // opening brackets cannot exceed N
    if (open > N) return
    // closing bracket > opening
    if (close > open) return

    // cond" for valid any                                   OR
    if (open == N && close == N) { // s.len = 2N
        print(s)
        return
    }

    // main logic
    // add opening bracket
    valid(N, s + '(', open + 1, close)
    // add closing bracket
    valid(N, s + ')', open, close + 1)
}
```

Output

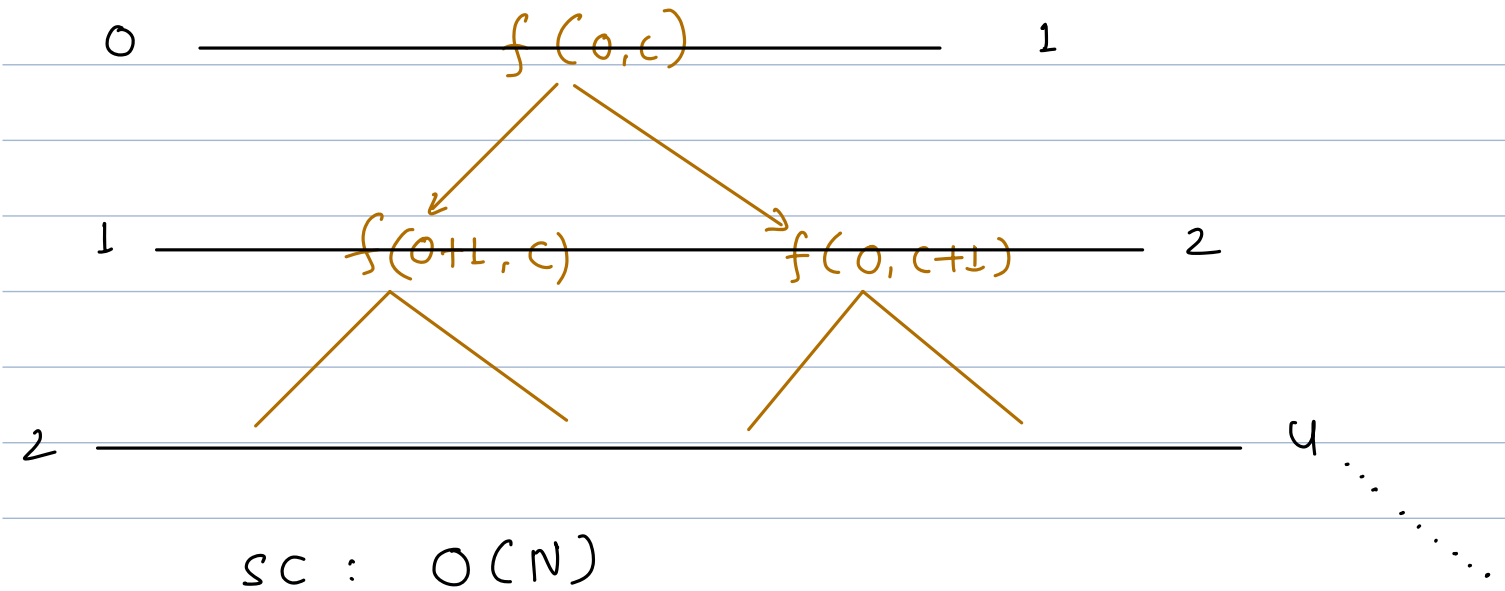
(())

()()

```
void valid (int N, String s, int open, int close) {
    if (open > N) return
    if (close > open) return
    if (open == N && close == N) {
        print(s)
        return
    }
    valid(N, s + '(', open + 1, close)
    valid(N, s + ')', open, close + 1)
}
```



TC for valid Parenthesis. TC : O(2^N)



SC : O(N)