

Arrays 1

Content

- Find maximum subarray sum.
- Perform multiple queries.
- Rainwater trapping.

- Rules →
- ① Private chat for answers
 - ② Use question
 - ③ Be patient.

Maximum Subarray Sum

continuous part of Arr

Given $\text{int}[] A$. Find max subarray sum of all subarrays

$A = \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 \\ -2 & 3 & 4 & -1 & 5 & -10 & 7 \end{matrix}$ ans = 11

$A = \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 \\ -3 & 4 & 6 & 8 & -10 & 2 & 7 \end{matrix}$ ans = 18

Quiz

$A = \begin{matrix} 4 & 5 & 2 & 1 & 6 \end{matrix}$ ans = 18

Observation — $\forall_i A_i > 0$ entire array sum = ans.

$A = \begin{matrix} -4 & -3 & -6 & -9 & -2 \end{matrix}$

Observation — $\forall_i A_i < 0$ pick the largest no.

Bruteforce

calculate sum of all subarrays and check max

ans = $-\infty$ // INT_MIN

```
for i → 0 to N-1 { // start index of sub
  for j → i to N-1 { // end index of sub
    sum = 0
    for k → i to j {
```

```

    sum += A[k]
    ans = max(ans, sum)
}
print(ans)

```

TC: $O(N^3)$
SC: $O(1)$

can be optimised via
Prefix sum.

TC: $O(N^2)$
SC: $O(N)$

Carry forward \longrightarrow calculate & use immediately.

```

ans = -∞
for i  $\longrightarrow$  0 to N-1 {
    sum = 0 // carry forward
    for j  $\longrightarrow$  i to N-1 {
        sum += A[j]
        ans = max(ans, sum)
    }
}
print(ans)

```

TC: $O(N^2)$
SC: $O(1)$

Example:
 i: -1 10 -2
 j: -1 10 -2
 sum = 0 -2
 ans = ~~0~~ -1
 10

Kadane's

Case 1 > all no. are +ve

$A = 1 \ 2 \ 3$

+ve +ve +ve

Case 2 > all no. are -ve

$A = -1 \ -2 \ -3$

pick the largest

Case 3 > -ve -ve -ve +ve +ve +ve -ve -ve -ve

Case 4 > Generic case

→ take as in increased overall sum
 $\cancel{7} \ -1 \ 1 \ 2 \ 3$

→ don't take
 $\cancel{7} \ -10 \ 1 \ 2 \ 3$

↓

	0	1	2	3	4	5	6	7	8
	-2	3	4	-1	5	-10	7	-9	12

sum	0	-2 ⁰	3	7	6	11	1	8	-1 ¹⁰	12
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ans	-∞	-2	3	7	7	11	11	11		12
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Pseudocode

TC : $O(N)$

SC : $O(1)$

ans = $-\infty$

sum = 0

for $i \rightarrow 0$ to $N-1$ {

sum += $A[i]$

ans = max(ans, sum)

if (sum < 0) {

sum = 0

}

return ans

$A =$

	0	1	2	3
	-3	-5	-2	-6
sum = 0	-3 $\rightarrow 0$	-5 $\rightarrow 0$	-2 $\rightarrow 0$	-6 $\rightarrow 0$
ans = $-\infty$	-3	-3	-2	-2

		0	1	2	3	4	5	6	7	8
$A =$		6	-3	9	-15	3	12	0	-24	3
sum	0	6	3	12	-3 $\rightarrow 0$	3	15	15	-9 $\rightarrow 0$	3
ans	$-\infty$	6	6	12	12	12	15	15	15	15

-2 3 4 -1 5 -10 7

ans = $-\infty$

sum = 0, start, L, R

final ans

for i \rightarrow 0 to N-1 {

sum += A[i]

if (sum > ans) {

ans = sum

L = start

R = i

}

if (sum < 0) {

sum = 0

start = i+1

}

}

// Print the array b/w L to R

Perform multiple queries from i to last index

Given an $\text{int}[]$ A where every element is 0

Return the final array after performing multiple queries

Query (i, x) \longrightarrow Add x to all no. from i to $N-1$

	0	1	2	3	4	5	6
$A[]$	0	0	0	0	0	0	0
$i \quad x$							
Query (1, 3)		+3	+3	+3	+3	+3	+3
Query (4, 2)					+2	+2	+2
Query (3, 1)				+1	+1	+1	+1
	0	3	3	4	6	6	6

Bruteforce \longrightarrow

For every query

Add x to all no. from i to $N-1$

TC: $O(Q * N)$

SC: $O(1)$

	0	1	2	3	4	5	6
$A[]$	0	0	0	0	0	0	0
$i \quad x$							
Query (1, 3)		+3					
Query (4, 2)					+2		
Query (3, 1)				+1			
	0	3	0	1	2	0	0

Prefix sum

0 3 3 4 6 6 6

Pseudocode

Queries $[][[]]$

$\begin{matrix} i & x \\ \{0, 1\} \end{matrix}$

for $q \rightarrow 0$ to $Q-1$ {

$i = \text{Queries}[q][0]$

$x = \text{Queries}[q][1]$

$A[i] += x$

TC: $O(Q+N)$

SC: $O(1)$

for ($i \rightarrow 1$ to $N-1$) {

$A[i] = A[i] + A[i-1]$

print(A)

Perform multiple queries from i to j index

Given an $\text{int}[]$ A where every element is 0

Return the final array after performing multiple queries

Query (i, j, x) \longrightarrow Add x to all no. from i to j

	0	1	2	3	4	5	6
$A[]$	0	0	0	0	0	0	0
i							
j							
x							
Query (1, 3, 2)		+2	+2	+2			
Query (2, 5, 3)			+3	+3	+3	+3	
Query (5, 6, -1)						-1	-1
	0	2	5	5	3	2	-1

	0	1	2	3	4	5	6	7
	0	0	0	0	0	0	0	0
i								
j								
x								
1 4 3		3	3	3	3			
0 5 -1	-1	-1	-1	-1	-1			
2 2 4			4					
4 6 3					3	3	3	
	-1	2	6	2	5	3	3	0

Bruteforce

For each query \longrightarrow

Add x from index i to j

TC : $O(QN)$

SC : $O(1)$

	0	1	2	3	4	5
	0	0	0	0	0	0

Query (1, 3, 2)



		+2	+2	+2	+2	+2
					-2	-2

	0	2	2	2	0	0
--	---	---	---	---	---	---

Query (1, 2)

Query (4, -2)

	0	1	2	3	4	5
	0	0	0	0	0	0

+2

-2

	0	2	0	0	-2	0
--	---	---	---	---	----	---

ps



0	2	2	2	0	0
---	---	---	---	---	---

Pseudocode

Queries [][][] { i, j, x }

```

for q -> 0 to Q-1 {
    i = Queries[q][0]
    j = Queries[q][1]
    x = Queries[q][2]
    A[i] += x
    if (j+1 < N) A[j+1] -= x
}

```

TC: O(Q+N)

SC: O(1)

Break

```

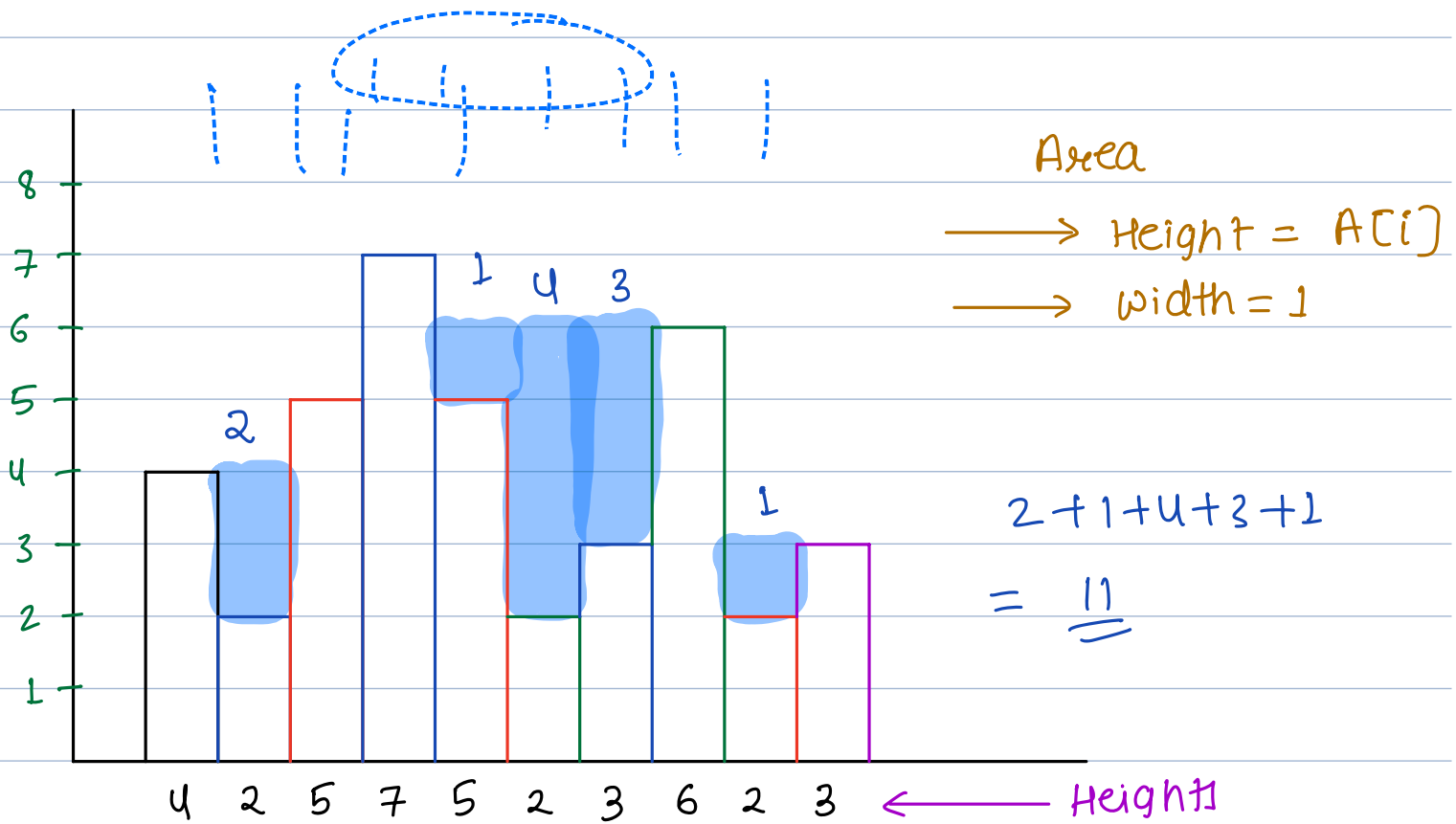
for ( i -> 1 to N-1 ) {
    A[i] = A[i] + A[i-1]
}

```

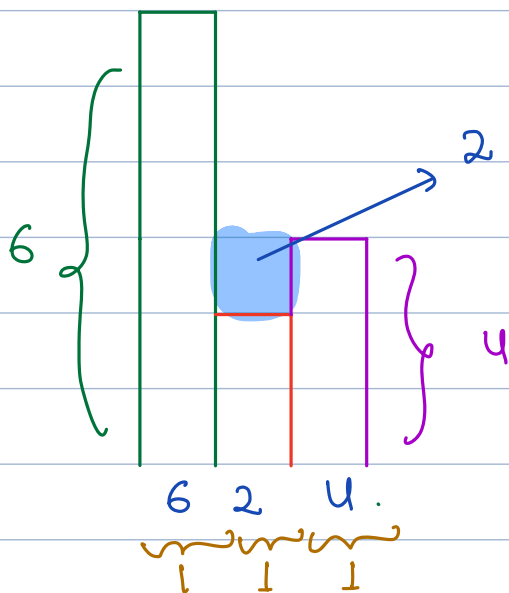
22:38

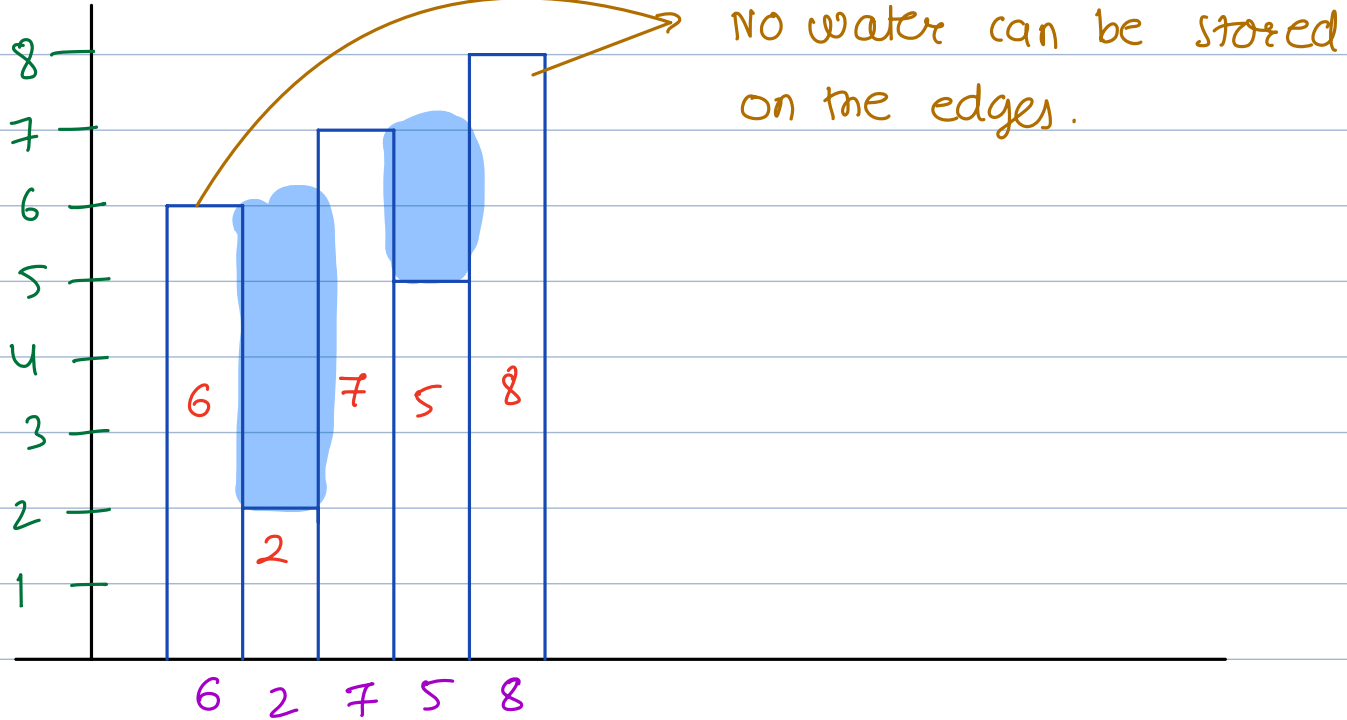
print (A)

Q> Rain water Trapping

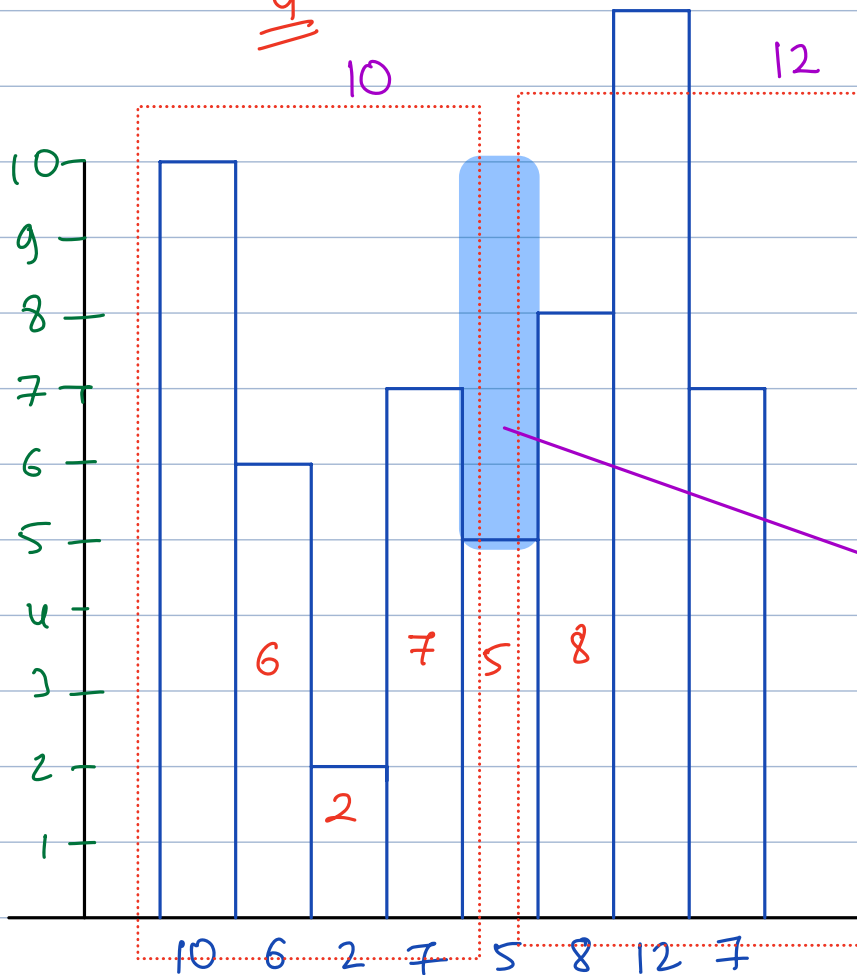


Given $A[N]$ each value represents building height
 calculate the total water accumulated coz of rain





$$\begin{aligned}
 & \underbrace{6 \quad 8} \\
 & 6 - 2 \\
 & \underline{\underline{4}}
 \end{aligned}$$



Till height 10
I can fill the water

So above building 5
water stored

$$\longrightarrow 10 - 5 = \underline{\underline{5}}$$

Bruteforce

water = 0

TC $O(N^2)$

SC $O(1)$

for $i \rightarrow 0$ to $N-1$ {

leftMax = $-\infty$ // 0 will work too.

for $j \rightarrow i$ to 0 {

leftMax = max(leftMax, A[j])

}

rightMax = $-\infty$ // 0 will work too.

for $j \rightarrow i$ to $N-1$ {

rightMax = max(rightMax, A[j])

}

waterHeight = min(leftMax, rightMax)

area = waterHeight - A[i]

water += area

}

print(water)

↓

6 2 4

leftMax

6 6 6

rightMax

6 4 4

waterHeight

6 4 4

area

6-6 2 4-4

water

0 2 2

	4	2	5	7	5	2	3	6	2	3	
left max	4	4	5	7	7	7	7	7	7	7	→ prefix Max
right max	7	7	7	7	6	6	6	6	3	3	→ suffix Max
water height	4	4	5	7	6	6	6	6	3	3	
area	0	2	0	0	1	4	3	0	1	0	
water	0	2	2	2	3	7	10	10	11	11	

prefix $R \rightarrow L$ = suffix

Pseudocode

```

prefixMax = new int[N]
leftMax = -∞
for i → 0 to N-1 {
    leftMax = max(leftMax, A[i])
    prefixMax[i] = leftMax
}

```

```

suffixMax = new int[N]
rightMax = -∞
for i → N-1 to 0 {
    rightMax = max(rightMax, A[i])
    suffixMax[i] = rightMax
}

```

water = 0

for $i \longrightarrow 0$ to $N-1$ {

leftMax = prefixMax[i]

rightMax = suffixMax[i]

waterHeight = min(leftMax, rightMax)

area = waterHeight - A[i]

water += area

}