

Linked List 3

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AGENDA:

—————>

Doubly Linked List

—————>

Insert a node

—————>

Delete first occurrence

—————>

LRU cache *

—————>

Shallow Copy vs Deep Copy

—————>

Copy List.

GREAT
JOB!

Current PSP

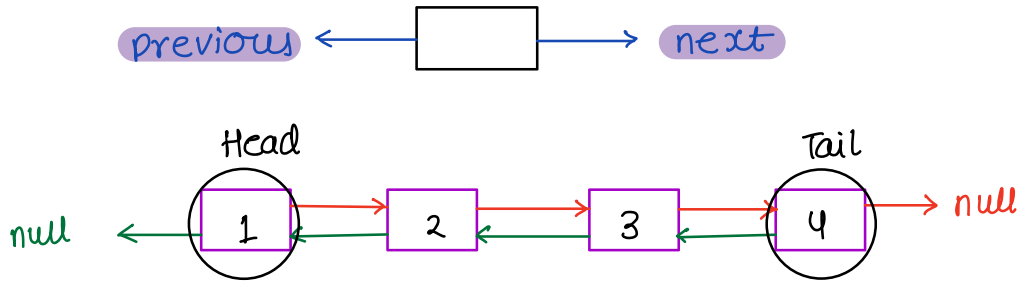


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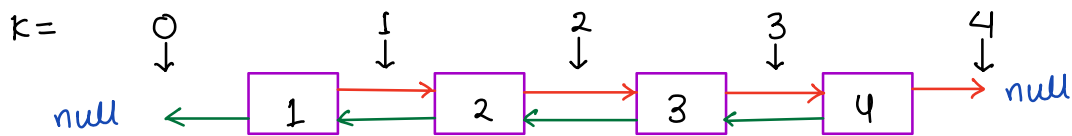


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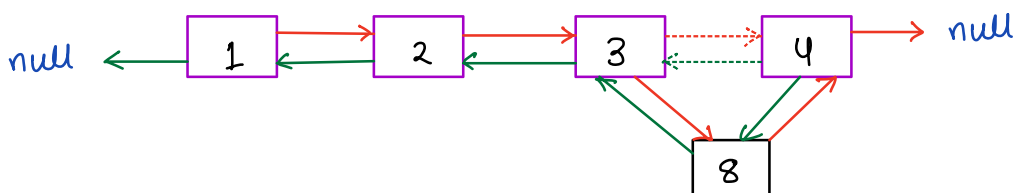
Doubly Linked List



Given a doubly LL. Insert a node with data X at a position K $[0 \leq K \leq N]$



$X = 8$, $K = 3$



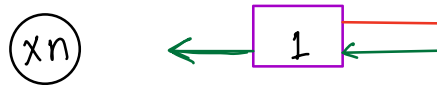
Corner Cases

- ① Head == null
- ② $K == 0$
- ③ insert at tail

Node

```
insertNode ( head , k , x ) {  
  xn = new DLL(x) → new node to be  
                        inverted
```

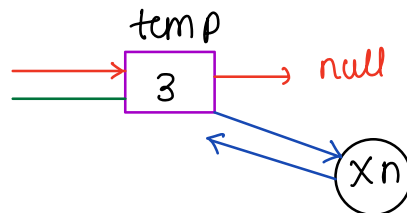
```
  if ( head == null ) return xn
```



```
  if ( k == 0 ) {  
    xn.next = head  
    head.prev = xn  
    return xn  
  }
```

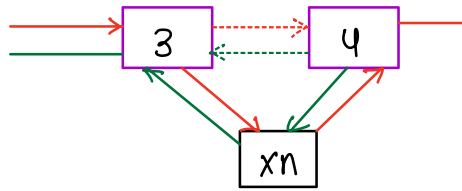
```
  temp = head
```

```
  for i → 1 to k-1 {  
    temp = temp.next  
  }
```



```
  if ( temp.next == null ) { // insert at tail  
    temp.next = xn  
    xn.prev = temp  
    return head  
  }
```

temp. next



```

nxt = temp.next
temp.next = xn
xn.prev = temp
xn.next = nxt
nxt.prev = xn

```

```

return head

```

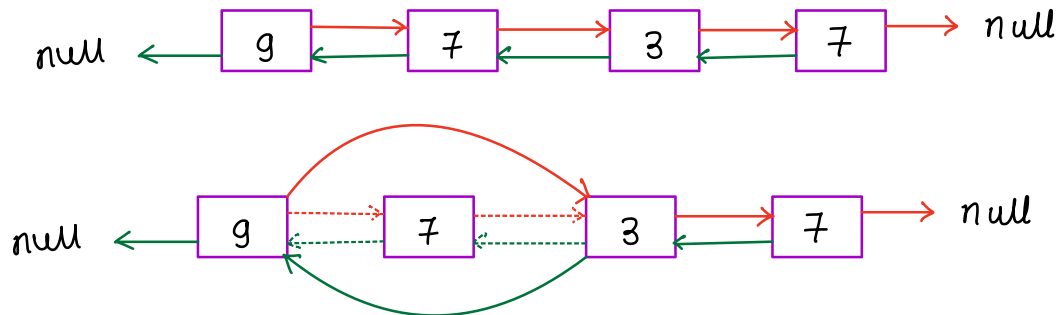
TC: $O(k)$

SC: $O(1)$

Q> Delete the first occurrence of data X in a given doubly linked list.

If not present then, no updates

$X = 7$



Corner Cases

① X is not present

② $\leftarrow \text{null} \quad \boxed{X} \rightarrow \text{null}$

Node delete Node (head , X) {

temp = head

```
while ( temp != null ) {  
    if ( temp.data == X ) {  
        break  
    }  
    temp = temp.next  
}
```

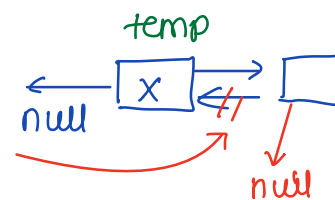
TC: OCN)

```
// X not found  
if ( temp == null ) {  
    return head  
}
```



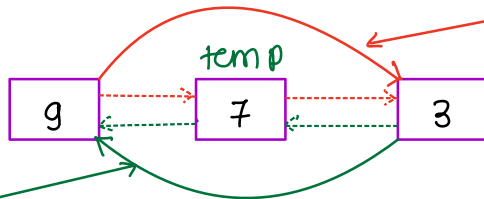
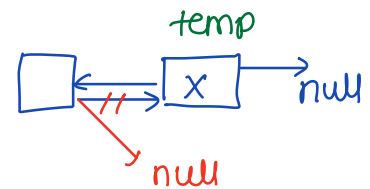
```
// Single node X  
if ( temp.next == null &&  
    temp.prev == null ) {  
    return null  
}
```

```
// X is at head  
if ( temp.prev == null ) {  
    temp.next.prev = null  
    return temp.next  
}
```



// X is at tail

```
if (temp.next == null) {  
    temp.prev.next = null  
    return head  
}
```



```
temp.prev.next = temp.next  
temp.next.prev = temp.prev
```

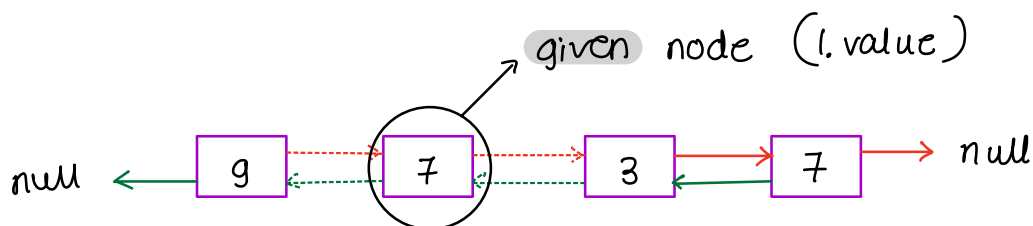
return head

TC: $O(N)$

SC: $O(1)$

Given the node X to be deleted in a DLL

TC to delete node X \rightarrow TC: $O(1)$



{least Recently used} LRU cache ****

Q> Given a running stream of integer and a fixed memory of size M . we have to maintain the most recent M elements.

In case the memory is full, we have to delete least recent element and insert current data into the memory {as most recent}

↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓
10 15 19 20 18 23 20 19 17 17 10

$M = 5$

~~10~~ ~~15~~ ~~19~~ ~~20~~ ~~18~~ 23 20 19 ~~17~~ 17 10

if my memory is full

if x is present

- ① delete x
- ② Insert x as most recent

if x is not present

- ① delete oldest element
- ② Insert x as most recent element.

Break 22:20

What is the best data structure to do the checking part i.e. if x is present or not?

~~HashSet~~ or **HashMap**

To maintain order \longrightarrow ~~Array~~, ~~ArrayList~~, **Linked List**

Tc to delete $O(N)$ $O(N)$ $O(N)$

If node to be deleted is known

Tc: $O(1)$

which is **DLL**

HashMap < Integer, Node >

head = null

tail = null

global

void addX (X) {

if (hm.containsKey(x)) {

node = hm.get(x) // return Node

deleteNode (node)

insertAtLast (node)

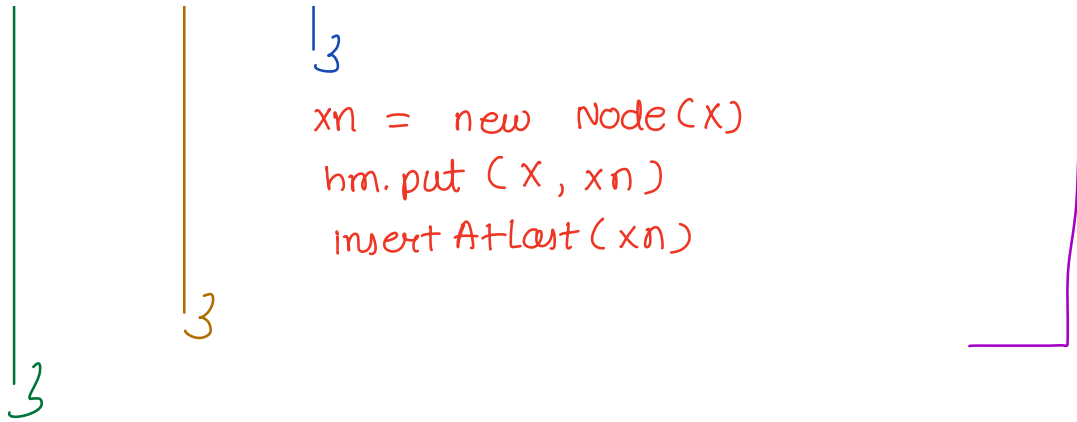
}

else {

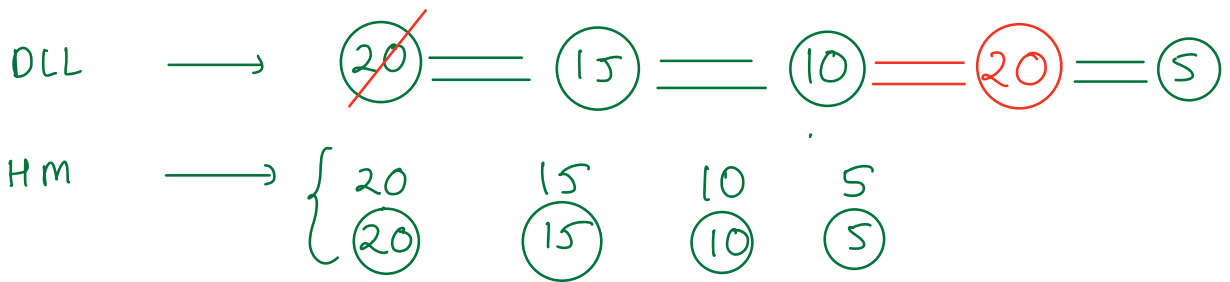
if (hm.size() == M) {

hm.remove (head.data)

deleteNode (head)



$M = 5$



TC per addition of x : $O(1)$

SC : $O(M)$

Deep Copy & Shallow Copy

```
x = new Node(5)
```

```
y = x
```

```
y.val = 10
```

```
print(x.val) → 10
```

shallow

```
x = new Node(5)
```

```
y = new Node(x.val)
```

```
y.val = 10
```

```
print(x.val) → 5
```

Deep Copy

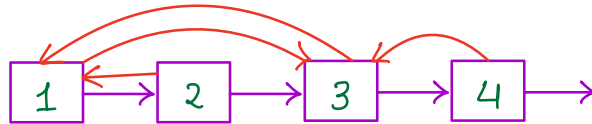
```
class Node {  
    data  
    next  
    random  
}
```

Q> Create a deep copy of given LL with random pointers

Random →

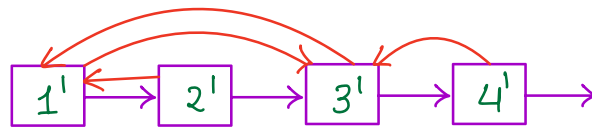
Next →

Head



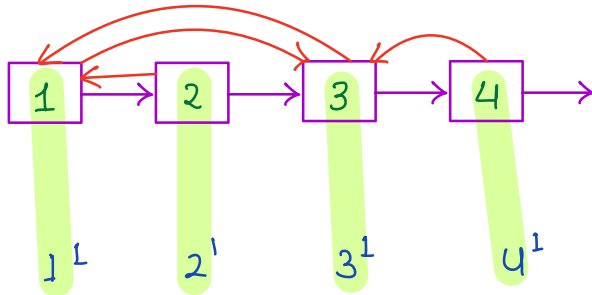
1/p

Head¹



ans

Head



HM { Node : Node }

original copy

Pseudocode

```
Node deepCopy ( Node head) {  
    HashMap<Node, Node> hm = new HashMap<>()  
    hm.put (null, null)  
    HM { Node : Node }  
        ↓       ↓  
    original  copy  
  
    temp = head  
    while (temp != null) {  
        hm.put (temp, new Node(temp.data))  
        temp = temp.next  
    }  
  
    temp = head  
  
    while (temp != null) {  
        // update next pointer of copy  
        hm.get(temp).next = hm.get(temp.next)  
        // update next pointer of copy  
        hm.get(temp).random = hm.get(temp.random)  
        temp = temp.next  
    }  
  
    return hm.get(head)  
}
```

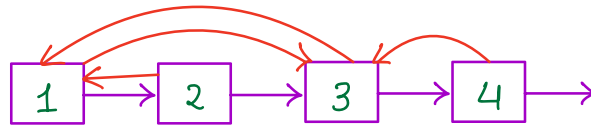
Tc: $O(N)$

Sc: $O(N)$

1

↓
O(1)

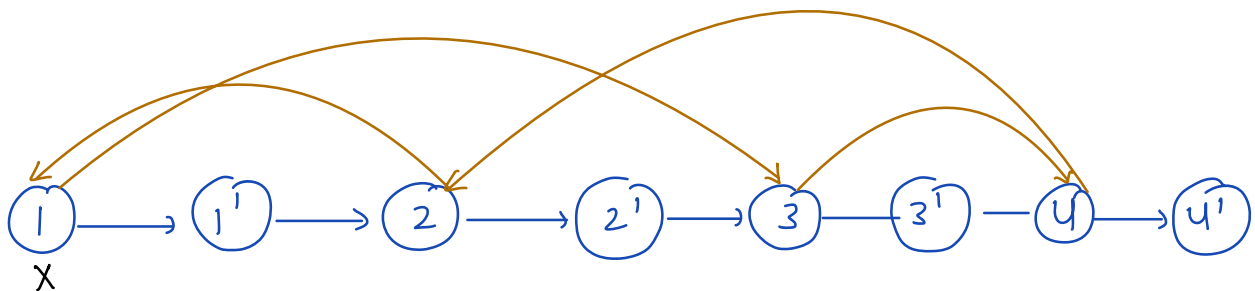
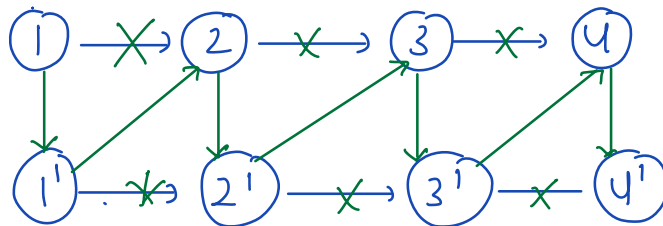
Head



step 1



step 2



How to connect random pointers

```

while (x != null)
{
    y = x.next
    y.random = x.random.next
    x = x.next.next
}
  
```

