Backtracking

Content

- Quizzes
- Subset
- Permutation

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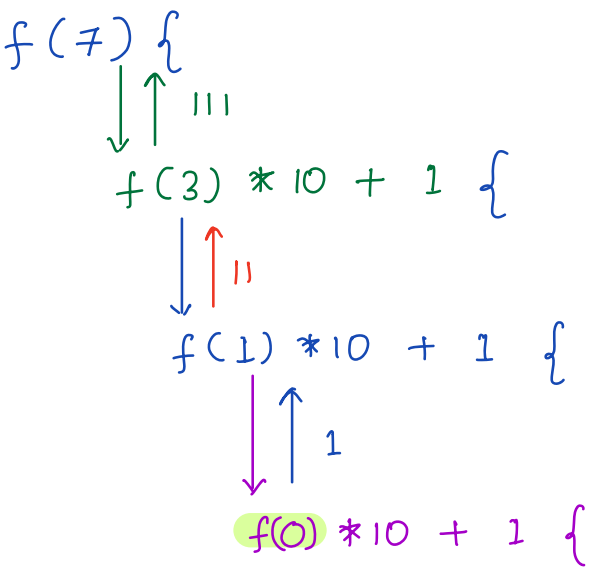
last time → Monday

64 % → 66.4 % → 75 %.

NOTE: Attempt all re-attempts for all contest.

```
int magicfun(int N) {
    if (N == 0) return 0
    else return magicfun(N/2) * 10 + N % 2
}
```

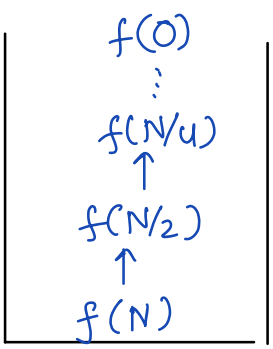
Output of above code for N = 7



TC : O(log(N))

SC : O(log(N))

NOTE : space complexity for a recursive code is never constant due to stack space.



```
void fun (char s[], int x) {
    print(s)
    char temp
    if (x < s.length/2) {
        temp = s[x]
        s[x] = s[s.length - x - 1]
        s[s.length - x - 1] = temp
        fun(s, x+1)
    }
}
```

swap (x, n-x-1)

Output for fun("SCROLL", 0) n = 6

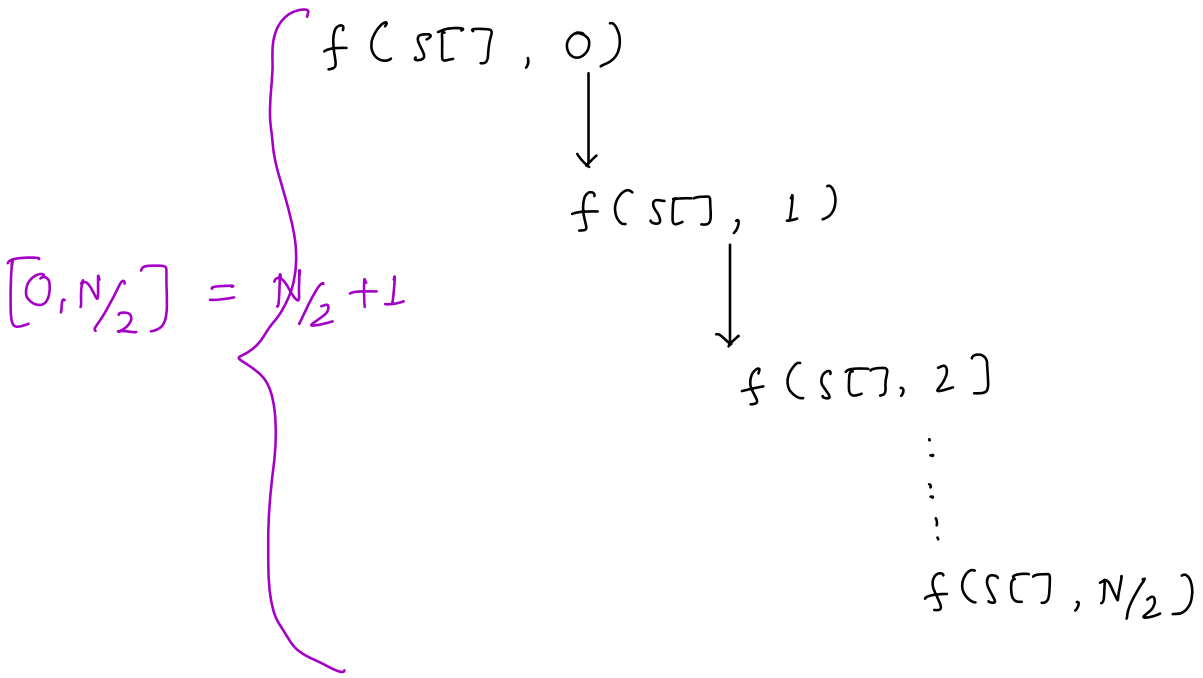
```
f(SCROLL, 0) {
    print(SCROLL)
    swap(0, 5)
    f(LCROLS, 1) {
        print(LCROLS)
        swap(1, 4)
        f(LLROCS, 2) {
            print(LLROCS)
            swap(2, 3)
            f(LLORCS, 3) {
                print(LLORCS)
```

output
SCROLL
LCROLS
LLROCS
LLORCS

TC : $O(\frac{N}{2} * N)$ $O(N^2)$

SC : O(N)

TC = # function calls * TC per call



TC = # fn call * TC per call

$= \left(\frac{N}{2} + 1\right) * N$

$= \frac{N^2}{2} + N \qquad O(N^2)$

SC = Max stack size at any instant

$= O\left(\frac{N}{2} + 1\right)$

$= O(N)$

Subset

* subset is not a subarray.

A = [1 2 3]

subsets = [ ]
[1]
[2]
[3]
[1 2]
[1 3]
[2 3]
[1 2 3]

⟶ 1 2 3 3 2 1 2 1 3 are same subset.

⟶ Every subarray is a subset ⟶ true

Every subset is a subarray ⟶ false

Given an array of distinct integer. Print all subsets using recursion.

I/P A = [1 2 3]

subsets = [ ]
[1]
[2]
[3]
[1 2]
[1 3]
[2 3]
[1 2 3]

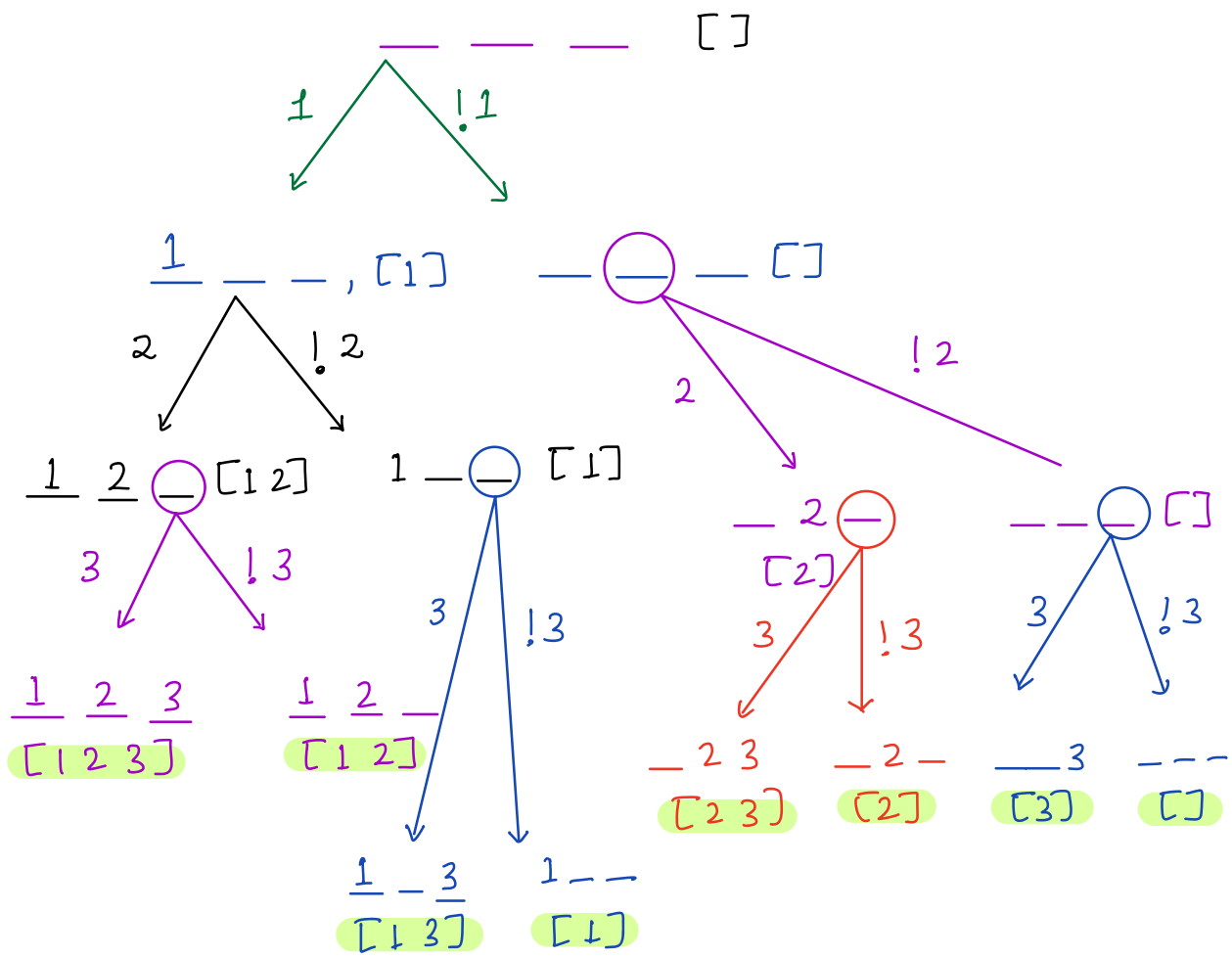
I/P A = [5 4]

subsets = [ ]
[4]
[5]
[4 5]

n	= 1	subsets	= 2
n	= 2	subsets	= 4
n	= 3	subsets	= 8
n	= n	subsets	$= 2^n$

Technique ***** ⟶ TAKE / DONT TAKE
select / dont select.
pick / dont pick

A = [1 2 3]



Pseudocode

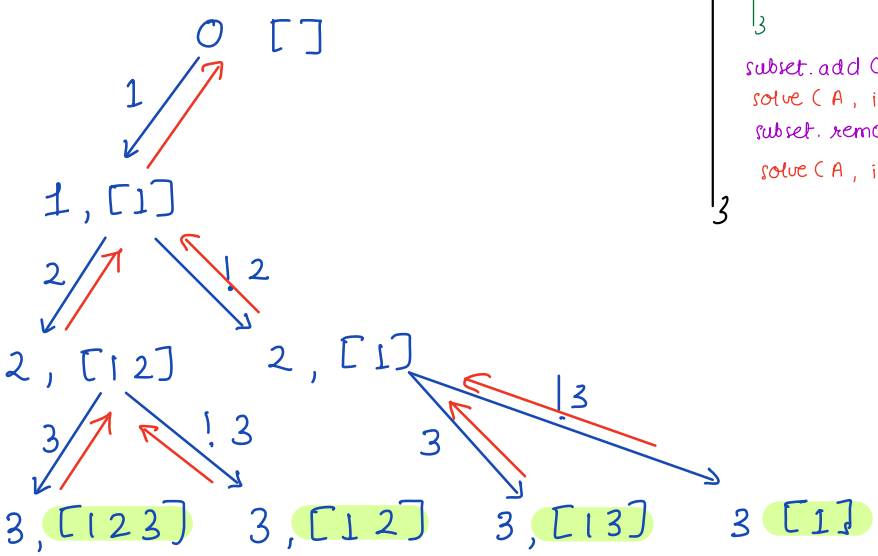
```
List<List<Integer>> subsets   // init. global

void solve(A[], index, subset) {
    // Base condition
    if (index == N) {
        subsets.add(subset)    subset.copy()
        return
    }
```

```
        // take A[index]
        subset.add(A[index])
        solve(A, index+1, subset)
        subset.remove() // remove last element

        // dont take A[index]
        solve(A, index+1, subset)
}
```

solve([1 2 3], 0, [])

subset → [] [] [] [] [] [] [] []

A = [1 2 3] (0 1 2)




```
void solve(A[], index, subset) {
    if (index == N) {
        subset.add(subset)   subset.copy()
        return
    }
    subset.add(A[index])
    solve(A, index+1, subset)
    subset.remove() // remove last element
    solve(A, index+1, subset)
}
```

TC : $O(N * 2^N)$

SC : $O(N)$

Break : 22 : 50

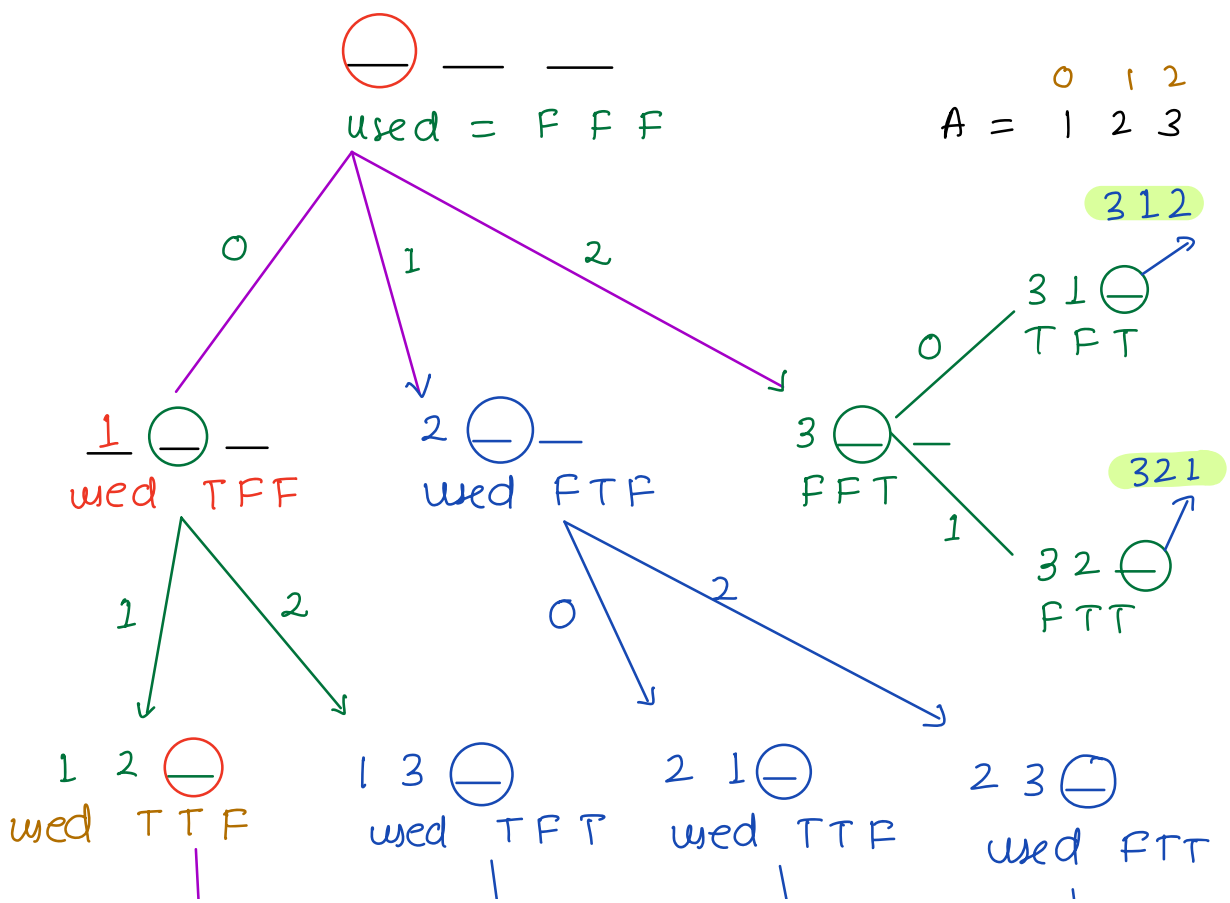
Permutation Ways to arrange an array.

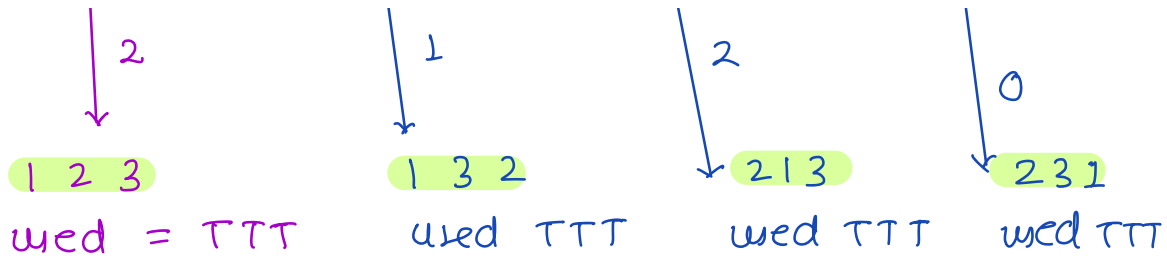
A = [1 2 3]

Permutations of A.

1	2	3
1	3	2
2	1	3
2	3	1
3	1	2
3	2	1

Given an array A[]. Print all the permutation of A[], {distinct integer}.





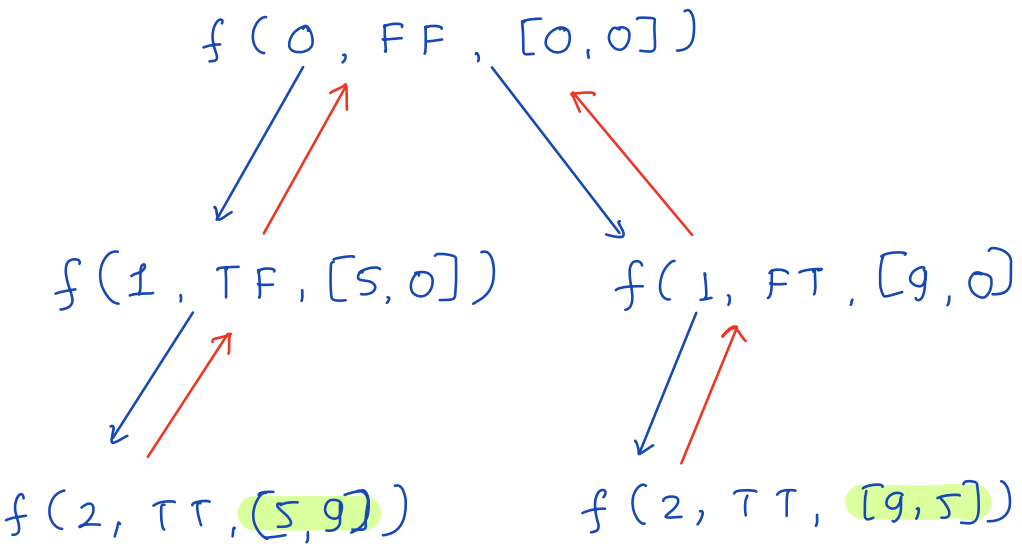
Pseudocode

```
void permutation ( A[] , pos , used[] , perm[] )
    if ( pos == N ) {
        print ( perm )
        return
    }

    // identify free values.
    for i ⟶ 0 to N-1 {
        if ( used [i] == false ) {
            used [i] = true          } do
            perm [pos] = A[i]
            permutation ( A , pos+1 , used , perm )
            used [i] = false         } undo
            perm [pos] = 0
        }
    }
}
```

pos → 0 ; used[] → [F F F] ; perm[] → [0 0 0]

0 1

A = [5 9] used = F F

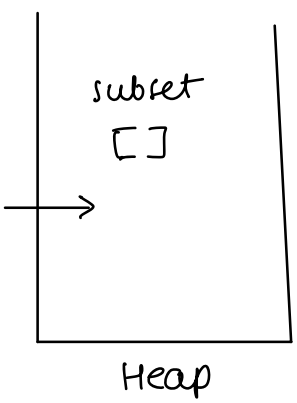


output
[5, 9]
[9, 5]

TC : O(N . N!)

SC : O(N + N + N) = O(N)

stack used perm



Doubt Session

Using dynamic array.

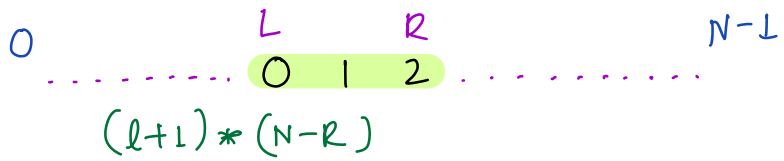
```
void permutation ( A[] , used[], , perm )        <- perm: List
    if ( perm.size() == N ) {
        print ( perm )
        return
    }

    // identify free values.
    for i ⟶ 0 to N-1 {
        if ( used[i] == false ) {
            used[i] = true                  }
            perm.add(A[i])                  } do
            permutation ( A , used , perm )
            used[i] = false                 }
            perm.remove()                   } undo
        }
    }
}
```

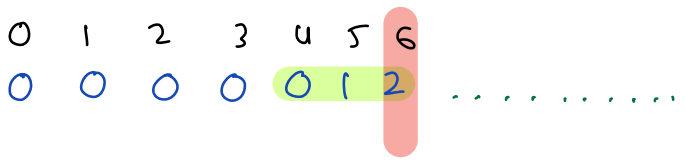
remove last index.

Nice Subarray Count

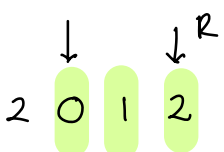
A = 0 1 2

How many subarrays are there with atleast one 0, 1, 2 ?





2 0 1 2 ⟶ 3



ans += L+1

1 + 2 = 3

f ⟶ R 0 to N-1

while subarray is nice

left++

// subarray is not nice

```
        count += L
    }
```