

max array size $\equiv 10^5 \equiv 100100$.

(i) `vector<int> g[100100];` \rightarrow basic adj. list for graphs.
`int vis[100100];`

(ii) To reverse a vector: \rightarrow this faster than using `cmp` bool of `sort` for descending order.
`reverse(g.begin(), g.end())`

(iii) Min/Max heap.

\downarrow done by priority-queue `q`;

****** for this whenever we are pushing into the queue,
we push `(-x)`

****** and whenever we want the min element. `int min = -q.front()`.

(iv) using `ii = pair<int, int>;`
using `lli = long long;`

(v) `vector<int> a(n, val);`

↓ initializing and storing
val in each element.

————— X —————

*** Whenever there is a new number & a difference or something of that sort, we sort the array and use lower_bound & upper_bound.

1) `int ind = lower_bound(a.begin(), a.end(), v) - a.begin();`

↓
This gives the "first" index from left → right where the $a[ind] \geq v$. greater than/

∴ ind -- is surely less than v. equal to.

2) `ind = upper_bound(a.begin(), a.end(), v) - a.begin();`

↓
This gives "first" index strictly greater than v. \Rightarrow $\begin{matrix} 1 & 8 & 8 & 8 & 16 \\ \text{1} & \text{2} & \text{3} & \text{4} & \end{matrix}$

$upper(8) = 4$

$lower(8) = 1$

ind -- may be equal to v or less than $\leq v$

Both are exactly the same when no's
arent repeated in the array &
query doesnt match any no.

(V) Now if u have any seq. of Objects($\text{int}, \text{pair}(\text{int}, \text{int})$) and u have to search if they exist in the container, storing in a vector would take $O(n)$ but if u use map or set, we can use the

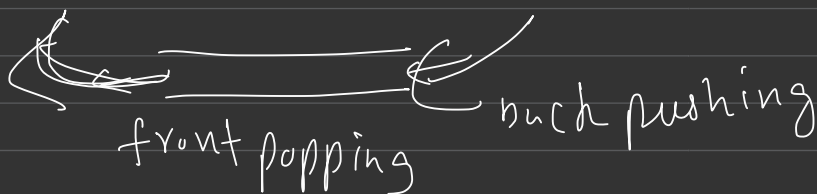
`set<pair<int, int>> s;`

`s.find(pair<int, int>)` fn which takes $\log n$ time to get.

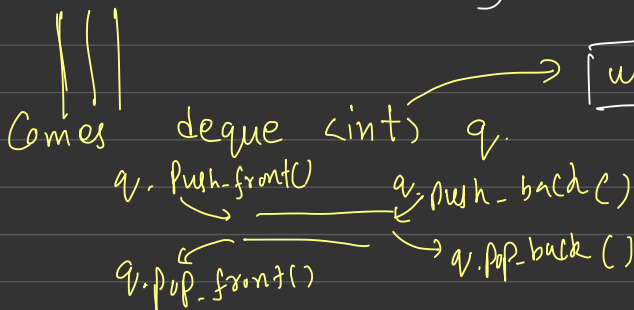
if the query doesn't have the element, it returns `s.end()` so we can check with an if condⁿ whether an element exists.!! in map,

`mp.find(k)` where k is the key!
not whole element.

iv) queue \rightarrow only has `q.push()`



only 1 dirⁿ. We cant push to the front of the queue.



used in 0-1 bfs.

for 0 edged vertex, we push in the front, normal edges in the back.

Adj. List for weighted graphs:

vector <pair<int, int>> g[100100];

g[a].push_back({b, c});

For Kruskal's algo: (edge list).

vector <pair<int, pair<int, int>>> q;

q.push_back({ $\overset{\downarrow}{w}, \overset{\downarrow}{a} \xrightarrow{w} \overset{\downarrow}{b}$, {a, b}});

sort(q.begin(), q.end());

2D graph settings: