Lab Report 1

April 30th, 2017

Bradley Evans and Dharti Tarapara

CS153 Operating Systems

Part 1: Adding System Calls

Hello World

Implementing a hello world function and executing it in the xv6 shell.

The Hello World function is implemented by changing the following files.

```
    defs.h

            A function prototype is added here.
            void hello(void);

    syscall.h

            Define a system call here.
            #define SYS_hello 22

    user.h

            Define a function prototype.
             int hello(void);

    proc.c.

            Define the actual hello function.
            void hello(void) { cprintf("hellol\n"); }

    sysproc.c.

            Define the system call here. This will simply call our hello() program in proc.c.
            int sys_hello(void) { hello(); return 0; }

    hello.c
```

This is added to the root directory.

```
#include "types.h"
#include "stat.h"
#include "user.h"

int main(int argc, char * argv[]) {
  hello();
  exit(-1);

  return 0;
}
```

Implemented in this way, the user can now execute a hello from the command line. This will trigger a system call, which will eventually call hello() from proc.c, displaying the hello world message.

Editing wait() and exit()

The $\overline{\texttt{exit()}}$ and $\overline{\texttt{wait()}}$ functions were modified to take integers.

Each instance of exit() in user programs was changed to exit(e) to reflect the change in type of exit from void to int. exit (and, in the same way, wait) now return a status.

To make use of this new parameter, exit(int) was changed to the following:

```
exit(int status)
 struct proc *p;
 int fd;
 // cprintf("exit status %d", proc->status);
 if(proc == initproc)
   panic("init exiting");
 // Close all open files.
 for(fd = 0; fd < NOFILE; fd++){
   if(proc->ofile[fd]){
     fileclose(proc->ofile[fd]);
     proc->ofile[fd] = 0;
 begin_op();
 end_op();
 proc->cwd = 0;
  acquire(&ptable.lock);
 // Parent might be sleeping in wait().
 wakeup1(proc->parent);
 // Pass abandoned children to init.
 for(p = ptable.proc; p < &ptable.proc[NPROC]; p++){</pre>
   if(p->parent == proc){
     p->parent = initproc;
     if(p->state == ZOMBIE)
       wakeup1(initproc);
```

```
proc->status = status; // MOD - 4/18
// Jump into the scheduler, never to return.
proc->state = ZOMBIE;
sched();
panic("zombie exit");
}
```

Also, wait() was changed to the following:

```
int
wait(int * status)
 struct proc *p;
 int havekids, pid;
 acquire(&ptable.lock);
   \ensuremath{//} Scan through table looking for exited children.
   havekids = 0:
   for(p = ptable.proc; p < &ptable.proc[NPROC]; p++){</pre>
     if(p->parent != proc)
       continue;
     havekids = 1;
     if(p->state == ZOMBIE){
       // Found one.
       // MOD - 4/29
       if (p->status != 0) {
          *status = p->status;
       } else *status = 0;
        pid = p->pid;
        kfree(p->kstack);
        freevm(p->pgdir);
        p->pid = 0;
        p->parent = 0;
        p->name[0] = 0;
        p->killed = 0;
        p->state = UNUSED;
        release(&ptable.lock);
        return pid;
   // No point waiting if we don't have any children.
   if(!havekids || proc->killed){
     release(&ptable.lock);
     return -1:
    // Wait for children to exit. (See wakeup1 call in proc_exit.)
    sleep(proc, &ptable.lock); //DOC: wait-sleep
```

Part 2: Schedulers

Implementing Priority Scheduling

Prior to this lab, xv6 used a round robin style scheduler. We modified it to use a priority scheduler, changing the scheduler code to the following.

```
scheduler(void)
 struct proc *p:
 int priority = 0; // hold priority value
   // Enable interrupts on this processor.
   sti();
    /* *** BEGIN MOD: 4/30 PRISCHED
    for(priority = 0; priority < maxpriority; priority++) {
     acquire(&ptable.lock);
     // Loop over process table looking for process to run.
      for(p = ptable.proc; p < &ptable.proc[NPROC]; p++){</pre>
       if(p->state != RUNNABLE)
         continue;
        ^{\prime} // PRISCHED: Now check and see if that process matches our current
        // priority level.
        if (p->priority == priority) \{
         // If it does, get it running.
         proc = p;
          switchuvm(p);
         p->state = RUNNING;
         swtch(&cpu->scheduler, proc->context);
          switchkvm();
        \ensuremath{//} Process is done running for now.
        // It should have changed its p->state before coming back.
       proc = 0;
      release(&ptable.lock);
   }*/
    // Loop over process table looking for process to run.
```

```
acquire(&ptable.lock);
  for(p = ptable.proc; p < &ptable.proc[NPROC]; p++) {</pre>
    if(p->state != RUNNABLE) {
     continue;
   release(&ptable.lock);
   priority = getprocpriority();
    acquire(&ptable.lock);
   if (priority < p->priority) {
     p->priority = priority;
   // Switch to chosen process. It is the process's job
   // to release ptable.lock and then reacquire it
    // before jumping back to us.
   proc = p;
    switchuvm(p);
   p->state = RUNNING;
    swtch(&cpu->scheduler, p->context);
   switchkvm();
   // Process is done running for now.
   // It should have changed its p->state before coming back.
   proc = 0;
 release(&ptable.lock):
}
```

In addition, a priority variable was added to to struct proc to give variables a means of storing a priority code.

The processes need a way to change priority. A mutator is required. In proc.c this is defined as setpriority:

```
int
setpriority(int num)
{
   if (num < 0) {
      num = 0;
   } else if (num > 63) {
      num = 63;
   }s priority

   proc->priority = num;
   return 0;
}
```

This is tied to a $\frac{1}{2}$ system call, added into the system in a similar way to the earlier $\frac{1}{2}$ call.

Also, a getprocpriority function will return the current seniormost priority level.

```
int
getprocpriority(void)
  struct proc *p;
  int priority = 65;
  // go through ptable
  acquire(&ptable.lock);
  for (p = ptable.proc; p < &ptable.proc[NPROC]; p++) {</pre>
   // look for runnable process
if(p->state != RUNNABLE) {
          continue;
   if(priority == 0) {
      priority = p->priority;
    } else {
      if (p->priority < priority) {
        priority = p->priority;
  release(&ptable.lock);
  return priority;
```

Priority Donation and Inheritence

A problem with priority schedulers is deadlock (more often called "unbounded priority inversion"), where lower-priority processes do not yield resources to a higher-priority process. A way to mitigate this is via *priority donation and inheritence*. If a higher priority task (H) attempts to access resource in use by a lower priority task (L), the lower priority task "inherits" a higher priority so that it can complete execution of some critical, uninterruptable portion of its programming. This prevents L from being pre-empted by a medium priority task M, which might then block H from executing.

To make this happen, you could use the setpriority mutator whenever a task enters the queue below a lower priority task. Setpriority would manually change the lower priority's priority value to be equal to that of the high-priority task. When the high-priority task is allowed to execute, it can use setpriority to return the low-priority task to its original priority level.