

Braeden Pope
Algorithm Design
Sort Design

```
1 PROMPT for file name
2 READ file at file name into list

3 i_pivot = len(list) - 1

4 WHILE i_pivot >= 0
5     i_largest = 0
6     FOR i_check in range(len(list))
7
8         IF list[i_check] > i_largest
9             i_largest = list[i_check]
10
11         IF i_largest > list[i_pivot]
12             temp = list[i_pivot]
13             list[i_pivot] = i_largest
14             list[i_check] = temp
15
16     i_pivot -= 1
17
18 PUT list on the screen
```

The algorithmic efficiency for this algorithm is $O(\log n)$, as after each iteration, the part of the list that gets checked (the unsorted side) becomes progressively smaller, and therefore quicker to check. This lines up with the $O(\log n)$ algorithmic efficiency.

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|-----------------------------|
| Sort Design Trace Table |
| Test Case: [52, 15, 39, 26] |

| Line | i_pivot | i_largest | i_check | list[i_check] | list[i_pivot] | temp |
|------|---------|-----------|---------|---------------|---------------|------|
| 1 | / | / | / | / | / | / |
| 2 | / | / | / | / | / | / |
| 3 | 3 | / | / | / | / | / |
| 4 | 3 | / | / | / | / | / |
| 5 | 3 | 0 | / | / | / | / |
| 6 | 3 | 0 | 0 | / | / | / |
| 7 | 3 | 0 | 0 | / | / | / |
| 8 | 3 | 0 | 0 | 52 | / | / |
| 9 | 3 | 52 | 0 | 52 | / | / |
| 10 | 3 | 52 | 0 | 52 | / | / |
| 11 | 3 | 52 | 0 | 52 | 26 | / |
| 12 | 3 | 52 | 0 | 52 | 26 | 26 |
| 13 | 3 | 52 | 0 | 52 | 26 | 26 |
| 14 | 3 | 52 | 0 | 52 | 52 | 26 |
| 15 | 3 | 52 | 0 | 26 | / | / |
| 16 | 2 | 52 | 0 | / | / | / |
| 17 | / | / | / | / | / | / |
| 18 | / | / | / | / | / | / |