Analiticcl

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Introduction

What is analitical?

- Analitical is a string-matching / fuzzy-matching system
- Intended for text normalisation like:
 - Diachronic variation
 - post-OCR/HTR variation
 - spelling correction (especially non-word errors)
- Lexicon-based; fuzzy lookups against a lexicon

Introduction: Context

- Developed in the Golden Agents projects
- ▶ Builds upon prior research (Reynaert 2010; Reynaert 2004)

Introduction: Implementation

- Built with performance and scalability in mind
 - Multi-threaded (parallellisation)
 - Written in Rust, compiles to native code
 - Low-level command-line tool and programming library
 - for Rust and for Python, via a binding
 - Unit/integration tests, CI, benchmarks
- ► Feature-rich
 - Highly parametrised
 - Flexible usage
- Source: https://github.com/proycon/analiticcl
- License: GNU GPLv3

Installation

Download, compile and install:

- \$ cargo install analiticcl
- and/or for the Python binding:
- \$ pip install analiticcl

Fuzzy string matching

Core function: Fuzzy string matching

- Given a string, find the forms in the lexicon that are closest (query mode)
 - error correction
- Given a text, find corrections for arbitrary substrings in the text (search mode)
 - error detection and correction
- Given a lexicon entry, find close variants in the text (learn mode)

Example: Query mode

```
$ analiticcl query --interactive --lexicon examples/eng.aspell.lexicon
                  --alphabet examples/simple.alphabet.tsv
Initializing model...
Loading lexicons...
Building model...
Computing anagram values for all items in the lexicon...
 - Found 119773 instances
Adding all instances to the index...
 - Found 108802 anagrams
Creating sorted secondary index...
Sorting secondary index...
 . . .
Querying the model...
(accepting standard input; enter input to match, one per line)
seperate
               separate 0.734375 \
seperate
               operate 0.6875 \
               desperate 0.6875 \
               temperate 0.6875
                                               serrate 0.65625
```

Variant matching: Naive approach

A naive approach to variant matching:

- ► Given *m* input words
- ► Compute edit distance (levenshtein) between each input word to all words in the lexicon (n)
- ► High computational cost! O(mn), and the levenshtein algorithm itself has already a O(I) (I=length) time complexity.
 - Does not scale

Variant matching: anagram hashing (1)

Anagram hashing (Reynaert 2010; Reynaert 2004) aims to drastically reduce the variant search space.

- ▶ Provides a fast *heuristic* for edit distance
- Analitical reimplements and improves upon the idea implemented in earlier tool TICCL

Variant matching: anagram hashing (2)

An Anagram Value (AV)..

- is computed for each 'word' in the input and in the lexicon
- uniquely represents the combination of characters in the word (unordered)
 - ightharpoonup AV(east) = AV(eats)
- has compositional properties:
 - $ightharpoonup AV(eat) \cdot AV(s) = AV(eats)$
 - $ightharpoonup \frac{AV(eats)}{AV(s)} = AV(eat)$
- each anagram value can be unambiguously decomposed to all its constituents
- no collisions between anagrams guaranteed (in this reimplementation)
- anagrams themselves deliberately collide
- serves as the key in a hash map (stores the lexicon)

Variant matching: hash function

Computation of the Anagram Value is simple composition of **prime** factors:

- Input: alphabet
- ► Each 'letter' in the alphabet is assigned a successive **prime number**, this is the Anagram Value of the 'letter'.
 - Example:

$$AV(a) = 2$$
, $AV(b) = 3$, $AV(c) = 7$, $AV(d) = 11$, $AV(e) = 17$

- ► The use of prime number guarantees no collisions between anagrams
- ▶ Novel with respect to Reynaert's approach.
- ▶ Simple hashing function (I=length, c_i =character at index i):

$$\prod_{i=0}^{l} AV(c_i)$$

- ► Caveat: May result in very large integers!
 - Exceeds 64-bit register
 - ► Requires an efficient big integer implementation

Variant matching: Search (1)

Loading stage: Compute Anagram Index and secondary index

- Compute Anagram Value for each entry in the lexicon and store in a hash map (the anagram index)
- ▶ Mapping the anagram value to all instances of the anagram:

$$AV(a, e, s, t) \mapsto [east, seat, eats]$$

Compute a secondary index mapping to sorted anagram values:

$$(n,|s|)\mapsto L$$

- where s is a string, |s| its length in characters, and n it's length in words/tokens
- where L is a sorted list of anagram values
- example: $(1,4) \mapsto [AV(a,e,s,t),...]$

Variant matching: Search (2)

Search stage: Given a 'word' to correct:

- we compute the anagram value for the input
- we look up this anagram value in the anagram index (if it exists) and gather the variant candidates associated with the anagram value
- we compute all deletions within a certain distance (e.g. by removing any 2 characters).
 - Example with 1 character:

$$del(AV(a,e,s,t)) = [AV(a,e,s),AV(e,s,t),AV(a,s,t),AV(a,e,t)]$$

► This is an arithmetic operation on the anagram values (division)

Variant matching: Search (3)

- For all of the anagram values resulting from these deletions we look which anagram values in our index match or contain the value under consideration. We again gather the candidates that result from all matches.
 - ▶ Match or contain: AV_a contains AV_b when

$$AV_a \mod AV_b = 0$$

- ▶ To facilitate this lookup, we make use of the *secondary index*
- Uses a binary search to find the anagrams that we should check our anagram value against (i.e. to check whether it is a subset of the anagram)
- Prevents needing to exhaustively try all anagram values in our index.

Variant matching: Search (4)

We have collected all possibly relevant variant instances: a considerably smaller set than the entire set we'd get if we didn't have the anagram heuristic! Now the set is reduced we apply more conventional measures:

- ► We compute several similarity metrics between the input and the possible variants:
 - Damerau-Levenshtein
 - Longest common substring
 - Longest common prefix/suffix
 - Casing difference
- ▶ A score is computed that is a weighted linear combination of the above components
 - the actual weights are configurable.
 - an exact match always has score 1.0.
- ► A cut-off value prunes the list of candidates that score too low
- Optionally, if a confusable list was provided, we adjust the score for known confusables

Feature: Confusable lists

- ▶ A list of *confusable patterns* with a weight
- Allows favouring or penalizing certain edits
- Example: OCR pattern: -[f]+[s]
- Example: historical dutch pattern: -[uy]+[ui] (huys -> huis)
- Allows context matching
- ► Taken into account as part of the similarity score function

Input and output

Analitical takes simple TSV files (tab separated values) as input:

- Lexicon
 - List of preferably validated words/multi-word expressions
 - May contain frequency information
- ▶ Variant list: explicitly relates variants to preferred forms.
 - Each variant carries a score expression how likely the variant maps to the preferred word
 - May also contain frequency information
 - Error list; a form of a variant lists where the variants are considered errors
 - Example: separate seperate 1.0 seperete 1.0
 - ▶ This is also the output form in *learn* mode
- Language model: for context-sensitive error detection/correction
- Multiple lexicon/variants lists supported
- Output is TSV or JSON

Background lexicon

- Analitical depends greatly on the quality of your input (lexicons)
- A good background corpus is required (out of vocabulary problem)
- ..otherwise analitical will eagerly mismatch to words it does know

Error Detection (1)

- ▶ In Query mode, input is a word/phrase you want to correct as a whole
- In Search mode, input is running text: analitical detects which parts of the input (words or higher order n-grams) need to be corrected.
- An additional and complex challenge!
- ▶ N-grams: consider splits and merges:
 - ▶ thehouse \rightarrow the house , teahouse ?
 - ightharpoonup tea house ightharpoonup teahouse ?
- ► Context is often a determining factor

Error Detection (2)

Given an input sentence:

- Extract all segments of the input, i.e. all n-grams up until a certain order
- 2. Do variant lookup for each (like query mode)
- Express all segments, their variants, their scores as transitions in a Finite State Transducer (FST)
 - Scores are expressed as costs
- 4. Extract the best path (lowest cost) with a beam search

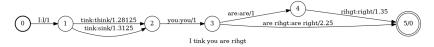


Figure 1: FST

Error Detection (3)

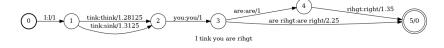


Figure 2: FST

- Scores are re-expressed as a cost (to be minimised)
- Base cost (integer) covers the number of input tokens spanned
 - establish a common ground for comparison between n-grams
 - n-grams compete
- ▶ Variant cost (fraction): inverse of the variant score: (0.0 best, approaching 1.0 as scores get worse)

$$cost = 1 - score$$

- ▶ **Joint variant score**: Sum of all costs on a complete path.
- Extract the 'cheapest' path(s)

Error Detection (4): Context

- 1. Extract the best *n* solutions from the FST (e.g. n = 250)
- 2. Compute the perplexity for each; using Language Model
- Compute a weighted combined score of the perplexity and the joint variant score
 - Not trivial, strikes a balance between LM and variant model
 - Compute normalised joint variant score:

$$variantscore_i = ln(\frac{cost_{best}}{cost_i})$$

Compute normalised LM score:

$$Imscore_i = \ln(\frac{PP_{best}}{PP_i})$$

► Weighted geometric mean:

$$score_i = \lambda_1 variantscore_i + \lambda_2 lmscore_i$$

4. Select the best scoring solution (minimize score)