

# CS 450 – Numerical Analysis

## Chapter 4: Eigenvalue Problems<sup>†</sup>

Prof. Michael T. Heath

Department of Computer Science  
University of Illinois at Urbana-Champaign  
heath@illinois.edu

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## Eigenvalue Problems

# Eigenvalues and Eigenvectors

- ▶ Standard *eigenvalue problem*: Given  $n \times n$  matrix  $\mathbf{A}$ , find scalar  $\lambda$  and nonzero vector  $\mathbf{x}$  such that

$$\mathbf{A} \mathbf{x} = \lambda \mathbf{x}$$

- ▶  $\lambda$  is *eigenvalue*, and  $\mathbf{x}$  is corresponding *eigenvector*
- ▶  $\lambda$  may be complex even if  $\mathbf{A}$  is real
- ▶ *Spectrum*  $= \lambda(\mathbf{A})$  = set of all eigenvalues of  $\mathbf{A}$
- ▶ *Spectral radius*  $= \rho(\mathbf{A}) = \max\{|\lambda| : \lambda \in \lambda(\mathbf{A})\}$

## Geometric Interpretation

- ▶ Matrix expands or shrinks any vector lying in direction of eigenvector by scalar factor
- ▶ Scalar expansion or contraction factor is given by corresponding eigenvalue  $\lambda$
- ▶ Eigenvalues and eigenvectors decompose complicated behavior of general linear transformation into simpler actions

## Examples: Eigenvalues and Eigenvectors

$$\blacktriangleright \mathbf{A} = \begin{bmatrix} 1 & 0 \\ 0 & 2 \end{bmatrix}: \lambda_1 = 1, \mathbf{x}_1 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \quad \lambda_2 = 2, \mathbf{x}_2 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

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$$\blacktriangleright \mathbf{A} = \begin{bmatrix} 3 & -1 \\ -1 & 3 \end{bmatrix}: \lambda_1 = 2, \mathbf{x}_1 = \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \quad \lambda_2 = 4, \mathbf{x}_2 = \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

$$\blacktriangleright \mathbf{A} = \begin{bmatrix} 1.5 & 0.5 \\ 0.5 & 1.5 \end{bmatrix}: \lambda_1 = 2, \mathbf{x}_1 = \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \quad \lambda_2 = 1, \mathbf{x}_2 = \begin{bmatrix} -1 \\ 1 \end{bmatrix}$$

$$\blacktriangleright \mathbf{A} = \begin{bmatrix} 0 & 1 \\ -1 & 0 \end{bmatrix}: \lambda_1 = i, \mathbf{x}_1 = \begin{bmatrix} 1 \\ i \end{bmatrix}, \quad \lambda_2 = -i, \mathbf{x}_2 = \begin{bmatrix} i \\ 1 \end{bmatrix}$$

where  $i = \sqrt{-1}$

## Characteristic Polynomial and Multiplicity

# Characteristic Polynomial

- ▶ Equation  $\mathbf{Ax} = \lambda\mathbf{x}$  is equivalent to

$$(\mathbf{A} - \lambda\mathbf{I})\mathbf{x} = \mathbf{0}$$

which has nonzero solution  $\mathbf{x}$  if, and only if, its matrix is singular

- ▶ Eigenvalues of  $\mathbf{A}$  are roots  $\lambda_i$  of *characteristic polynomial*

$$\det(\mathbf{A} - \lambda\mathbf{I}) = 0$$

in  $\lambda$  of degree  $n$

- ▶ *Fundamental Theorem of Algebra* implies that  $n \times n$  matrix  $\mathbf{A}$  always has  $n$  eigenvalues, but they may not be real nor distinct
- ▶ Complex eigenvalues of real matrix occur in complex conjugate pairs: if  $\alpha + i\beta$  is eigenvalue of real matrix, then so is  $\alpha - i\beta$ , where  $i = \sqrt{-1}$

## Example: Characteristic Polynomial

- ▶ Characteristic polynomial of previous example matrix is

$$\det \left( \begin{bmatrix} 3 & -1 \\ -1 & 3 \end{bmatrix} - \lambda \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \right) =$$

$$\det \left( \begin{bmatrix} 3 - \lambda & -1 \\ -1 & 3 - \lambda \end{bmatrix} \right) =$$

$$(3 - \lambda)(3 - \lambda) - (-1)(-1) = \lambda^2 - 6\lambda + 8 = 0$$

so eigenvalues are given by

$$\lambda = \frac{6 \pm \sqrt{36 - 32}}{2}, \quad \text{or} \quad \lambda_1 = 2, \quad \lambda_2 = 4$$



## Characteristic Polynomial, continued

- ▶ Computing eigenvalues using characteristic polynomial is *not* recommended because of
  - ▶ work in computing coefficients of characteristic polynomial
  - ▶ sensitivity of coefficients of characteristic polynomial
  - ▶ work in solving for roots of characteristic polynomial
- ▶ Characteristic polynomial is powerful theoretical tool but usually not useful computationally

## Example: Characteristic Polynomial

- ▶ Consider

$$\mathbf{A} = \begin{bmatrix} 1 & \epsilon \\ \epsilon & 1 \end{bmatrix} \quad \text{--- } 0.1$$

where  $\epsilon$  is positive number slightly smaller than  $\sqrt{\epsilon_{\text{mach}}}$

- ▶ Exact eigenvalues of  $\mathbf{A}$  are  $1 + \epsilon$  and  $1 - \epsilon$   
 $1.1 \quad 0.9$
- ▶ Computing characteristic polynomial in floating-point arithmetic, we obtain

$$\det(\mathbf{A} - \lambda \mathbf{I}) = \lambda^2 - 2\lambda + (1 - \epsilon^2) = \lambda^2 - 2\lambda + 1$$

which has 1 as double root

- ▶ Thus, eigenvalues cannot be resolved by this method even though they are distinct in working precision

## Computing Eigenvalues and Eigenvectors

## Problem Transformations

- ▶ *Shift*: If  $\mathbf{A}\mathbf{x} = \lambda\mathbf{x}$  and  $\sigma$  is any scalar, then  $(\mathbf{A} - \sigma\mathbf{I})\mathbf{x} = (\lambda - \sigma)\mathbf{x}$ , so eigenvalues of shifted matrix are shifted eigenvalues of original matrix
- ▶ *Inversion*: If  $\mathbf{A}$  is nonsingular and  $\mathbf{A}\mathbf{x} = \lambda\mathbf{x}$  with  $\mathbf{x} \neq \mathbf{0}$ , then  $\lambda \neq 0$  and  $\mathbf{A}^{-1}\mathbf{x} = (1/\lambda)\mathbf{x}$ , so eigenvalues of inverse are reciprocals of eigenvalues of original matrix
- ▶ *Powers*: If  $\mathbf{A}\mathbf{x} = \lambda\mathbf{x}$ , then  $\mathbf{A}^k\mathbf{x} = \lambda^k\mathbf{x}$ , so eigenvalues of power of matrix are same power of eigenvalues of original matrix
- ▶ *Polynomial*: If  $\mathbf{A}\mathbf{x} = \lambda\mathbf{x}$  and  $p(t)$  is polynomial, then  $p(\mathbf{A})\mathbf{x} = p(\lambda)\mathbf{x}$ , so eigenvalues of polynomial in matrix are values of polynomial evaluated at eigenvalues of original matrix

# Similarity Transformation

- ▶  $B$  is *similar* to  $A$  if there is nonsingular matrix  $T$  such that

$$B = T^{-1} A T$$

- ▶ Then

$$B\mathbf{y} = \lambda\mathbf{y} \Rightarrow T^{-1}AT\mathbf{y} = \lambda\mathbf{y} \Rightarrow A(T\mathbf{y}) = \lambda(T\mathbf{y})$$

so  $A$  and  $B$  have same eigenvalues, and if  $\mathbf{y}$  is eigenvector of  $B$ , then  $\mathbf{x} = T\mathbf{y}$  is eigenvector of  $A$

- ▶ Similarity transformations preserve eigenvalues, and eigenvectors are easily recovered

## Example: Similarity Transformation

- ▶ From eigenvalues and eigenvectors for previous example,

$$\begin{bmatrix} 3 & -1 \\ -1 & 3 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 2 & 0 \\ 0 & 4 \end{bmatrix}$$

and hence

$$\begin{bmatrix} 0.5 & 0.5 \\ 0.5 & -0.5 \end{bmatrix} \begin{bmatrix} 3 & -1 \\ -1 & 3 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} = \begin{bmatrix} 2 & 0 \\ 0 & 4 \end{bmatrix}$$

- ▶ So original matrix is similar to diagonal matrix, and eigenvectors form columns of similarity transformation matrix

## Diagonal Form

- ▶ Eigenvalues of diagonal matrix are diagonal entries, and eigenvectors are columns of identity matrix
- ▶ Diagonal form is desirable in simplifying eigenvalue problems for general matrices by similarity transformations
- ▶ But not all matrices are diagonalizable by similarity transformation
- ▶ Closest one can get, in general, is *Jordan form*, which is nearly diagonal but may have some nonzero entries on first superdiagonal, corresponding to one or more multiple eigenvalues

## Triangular Form

- ▶ Any matrix can be transformed into triangular (*Schur*) form by similarity, and eigenvalues of triangular matrix are diagonal entries
- ▶ Eigenvectors of triangular matrix less obvious, but still straightforward to compute
- ▶ If

$$\mathbf{A} - \lambda \mathbf{I} = \begin{bmatrix} \mathbf{U}_{11} & \mathbf{u} & \mathbf{U}_{13} \\ \mathbf{0} & 0 & \mathbf{v}^T \\ \mathbf{0} & \mathbf{0} & \mathbf{U}_{33} \end{bmatrix}$$

is triangular, then  $\mathbf{U}_{11}\mathbf{y} = \mathbf{u}$  can be solved for  $\mathbf{y}$ , so that

$$\mathbf{x} = \begin{bmatrix} \mathbf{y} \\ -1 \\ \mathbf{0} \end{bmatrix}$$

is corresponding eigenvector



# Relevant Properties of Matrices

- Properties of matrix **A** relevant to eigenvalue problems

Property	Definition
diagonal	$a_{ij} = 0$ for $i \neq j$
<del>tridiagonal</del>	<del><math>a_{ij} = 0</math> for <math> i - j  &gt; 1</math></del>
triangular	$a_{ij} = 0$ for $i > j$ (upper) $a_{ij} = 0$ for $i < j$ (lower)
Hessenberg	$a_{ij} = 0$ for $i > j + 1$ (upper) $a_{ij} = 0$ for $i < j - 1$ (lower)
orthogonal	$\mathbf{A}^T \mathbf{A} = \mathbf{A} \mathbf{A}^T = \mathbf{I}$
unitary	$\mathbf{A}^H \mathbf{A} = \mathbf{A} \mathbf{A}^H = \mathbf{I}$
symmetric	$\mathbf{A} = \mathbf{A}^T$
Hermitian	$\mathbf{A} = \mathbf{A}^H$
normal	$\mathbf{A}^H \mathbf{A} = \mathbf{A} \mathbf{A}^H$



## Eigenspaces and Invariant Subspaces

- ▶ Eigenvectors can be scaled arbitrarily: if  $\mathbf{Ax} = \lambda\mathbf{x}$ , then  $\mathbf{A}(\gamma\mathbf{x}) = \lambda(\gamma\mathbf{x})$  for any scalar  $\gamma$ , so  $\gamma\mathbf{x}$  is also eigenvector corresponding to  $\lambda$
- ▶ Eigenvectors are usually *normalized* by requiring some norm of eigenvector to be 1
- ▶  $\text{Eigenspace} = \mathcal{S}_\lambda = \{\mathbf{x} : \mathbf{Ax} = \lambda\mathbf{x}\}$
- ▶ Subspace  $\mathcal{S}$  of  $\mathbb{R}^n$  (or  $\mathbb{C}^n$ ) is *invariant* if  $\mathbf{AS} \subseteq \mathcal{S}$
- ▶ For eigenvectors  $\mathbf{x}_1 \cdots \mathbf{x}_p$ ,  $\text{span}([\mathbf{x}_1 \cdots \mathbf{x}_p])$  is invariant subspace

## Power Iteration

## Power Iteration

- ▶ Simplest method for computing one eigenvalue-eigenvector pair is *power iteration*, which repeatedly multiplies matrix times initial starting vector
- ▶ Assume  $\mathbf{A}$  has unique eigenvalue of maximum modulus, say  $\lambda_1$ , with corresponding eigenvector  $\mathbf{v}_1$
- ▶ Then, starting from nonzero vector  $\mathbf{x}_0$ , iteration scheme

$$\mathbf{x}_k = \mathbf{A}\mathbf{x}_{k-1}$$

converges to multiple of eigenvector  $\mathbf{v}_1$  corresponding to *dominant* eigenvalue  $\lambda_1$

## Convergence of Power Iteration

- ▶ To see why power iteration converges to dominant eigenvector, express starting vector  $\mathbf{x}_0$  as linear combination

$$\mathbf{x}_0 = \sum_{i=1}^n \alpha_i \mathbf{v}_i$$

where  $\mathbf{v}_i$  are eigenvectors of  $\mathbf{A}$

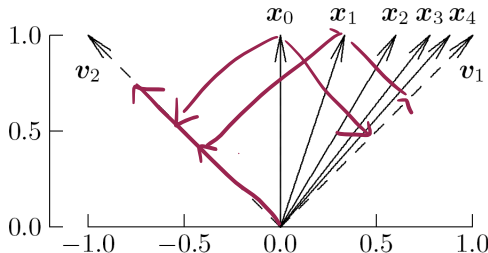
- ▶ Then

$$\begin{aligned} \mathbf{x}_k &= \mathbf{A}\mathbf{x}_{k-1} = \mathbf{A}^2\mathbf{x}_{k-2} = \cdots = \mathbf{A}^k\mathbf{x}_0 = \\ &= \sum_{i=1}^n \lambda_i^k \alpha_i \mathbf{v}_i = \lambda_1^k \left( \alpha_1 \mathbf{v}_1 + \sum_{i=2}^n (\lambda_i/\lambda_1)^k \alpha_i \mathbf{v}_i \right) \end{aligned}$$

- ▶ Since  $|\lambda_i/\lambda_1| < 1$  for  $i > 1$ , successively higher powers go to zero, leaving only component corresponding to  $\mathbf{v}_1$

## Geometric Interpretation

- Behavior of power iteration depicted geometrically



- Initial vector  $\mathbf{x}_0 = \mathbf{v}_1 + \mathbf{v}_2$  contains equal components in eigenvectors  $\mathbf{v}_1$  and  $\mathbf{v}_2$  (dashed arrows)
- Repeated multiplication by  $\mathbf{A}$  causes component in  $\mathbf{v}_1$  (corresponding to larger eigenvalue, 2) to dominate, so sequence of vectors  $\mathbf{x}_k$  converges to  $\mathbf{v}_1$

## Example: Power Iteration

- ▶ Ratio of values of given component of  $\mathbf{x}_k$  from one iteration to next converges to dominant eigenvalue  $\lambda_1$

- ▶ For example, if  $\mathbf{A} = \begin{bmatrix} 1.5 & 0.5 \\ 0.5 & 1.5 \end{bmatrix}$  and  $\mathbf{x}_0 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$ , we obtain

$k$	$\mathbf{x}_k^T$		ratio
0	0.0	1.0	
1	0.5	1.5	1.500
2	1.5	2.5	1.667
3	3.5	4.5	1.800
4	7.5	8.5	1.889
5	15.5	16.5	1.941
6	31.5	32.5	1.970
7	63.5	64.5	1.985
8	127.5	128.5	1.992

- ▶ Ratio is converging to dominant eigenvalue, which is 2

## Normalized Power Iteration

- ▶ Geometric growth of components at each iteration risks eventual overflow (or underflow if  $\lambda_1 < 1$ )
- ▶ Approximate eigenvector should be normalized at each iteration, say, by requiring its largest component to be 1 in modulus, giving iteration scheme

$$\begin{aligned}\mathbf{y}_k &= \mathbf{A}\mathbf{x}_{k-1} \\ \mathbf{x}_k &= \mathbf{y}_k / \|\mathbf{y}_k\|_\infty\end{aligned}$$

- ▶ With normalization,  $\|\mathbf{y}_k\|_\infty \rightarrow |\lambda_1|$ , and  $\mathbf{x}_k \rightarrow \mathbf{v}_1 / \|\mathbf{v}_1\|_\infty$



## Example: Normalized Power Iteration

- ▶ Repeating previous example with normalized scheme,

$k$	$\mathbf{x}_k^T$		$\ \mathbf{y}_k\ _\infty$
0	0.000	1.0	
1	0.333	1.0	1.500
2	0.600	1.0	1.667
3	0.778	1.0	1.800
4	0.882	1.0	1.889
5	0.939	1.0	1.941
6	0.969	1.0	1.970
7	0.984	1.0	1.985
8	0.992	1.0	1.992

⟨ interactive example ⟩

## Power Iteration with Shift

- ▶ Convergence rate of power iteration depends on ratio  $|\lambda_2/\lambda_1|$ , where  $\lambda_2$  is eigenvalue having second largest modulus
- ▶ May be possible to choose shift,  $\mathbf{A} - \sigma \mathbf{I}$ , such that

$$\left| \frac{\lambda_2 - \sigma}{\lambda_1 - \sigma} \right| < \left| \frac{\lambda_2}{\lambda_1} \right|$$

so convergence is accelerated

- ▶ Shift must then be added to result to obtain eigenvalue of original matrix
- ▶ In earlier example, for instance, if we pick shift of  $\sigma = 1$ , (which is equal to other eigenvalue) then ratio becomes zero and method converges in one iteration
- ▶ In general, we would not be able to make such fortuitous choice, but shifts can still be extremely useful in some contexts, as we will see later

## Limitations of Power Iteration

Power iteration can fail for various reasons

- ▶ Starting vector may have *no* component in dominant eigenvector  $\mathbf{v}_1$  (i.e.,  $\alpha_1 = 0$ ) — not problem in practice because rounding error usually introduces such component in any case
- ▶ There may be more than one eigenvalue having same (maximum) modulus, in which case iteration may converge to linear combination of corresponding eigenvectors
- ▶ For real matrix and starting vector, iteration can never converge to complex vector

## Inverse and Rayleigh Quotient Iterations

## Inverse Iteration

- ▶ To compute smallest eigenvalue of matrix rather than largest, can make use of fact that eigenvalues of  $\mathbf{A}^{-1}$  are reciprocals of those of  $\mathbf{A}$ , so smallest eigenvalue of  $\mathbf{A}$  is reciprocal of largest eigenvalue of  $\mathbf{A}^{-1}$
- ▶ This leads to *inverse iteration* scheme

$$\begin{aligned}\mathbf{A}\mathbf{y}_k &= \mathbf{x}_{k-1} \\ \mathbf{x}_k &= \mathbf{y}_k / \|\mathbf{y}_k\|_\infty\end{aligned}$$

which is equivalent to power iteration applied to  $\mathbf{A}^{-1}$

- ▶ Inverse of  $\mathbf{A}$  not computed explicitly, but factorization of  $\mathbf{A}$  used to solve system of linear equations at each iteration
- ▶ Inverse iteration converges to eigenvector corresponding to *smallest* eigenvalue of  $\mathbf{A}$
- ▶ Eigenvalue obtained is dominant eigenvalue of  $\mathbf{A}^{-1}$ , and hence its reciprocal is smallest eigenvalue of  $\mathbf{A}$  in modulus

## Example: Inverse Iteration

- ▶ Applying inverse iteration to previous example to compute smallest eigenvalue yields sequence

$k$	$\mathbf{x}_k^T$		$\ \mathbf{y}_k\ _\infty$
0	0.000	1.0	
1	-0.333	1.0	0.750
2	-0.600	1.0	0.833
3	-0.778	1.0	0.900
4	-0.882	1.0	0.944
5	-0.939	1.0	0.971
6	-0.969	1.0	0.985

which is indeed converging to 1 (which is its own reciprocal in this case)

⟨ interactive example ⟩

## Inverse Iteration with Shift

- ▶ As before, shifting strategy, working with  $\mathbf{A} - \sigma \mathbf{I}$  for some scalar  $\sigma$ , can greatly improve convergence
- ▶ Inverse iteration is particularly useful for computing eigenvector corresponding to approximate eigenvalue, since it converges rapidly when applied to shifted matrix  $\mathbf{A} - \lambda \mathbf{I}$ , where  $\lambda$  is approximate eigenvalue
- ▶ Inverse iteration is also useful for computing eigenvalue closest to given value  $\beta$ , since if  $\beta$  is used as shift, then desired eigenvalue corresponds to smallest eigenvalue of shifted matrix

## Rayleigh Quotient

- ▶ Given approximate eigenvector  $\mathbf{x}$  for real matrix  $\mathbf{A}$ , determining best estimate for corresponding eigenvalue  $\lambda$  can be considered as  $n \times 1$  linear least squares approximation problem

$$\mathbf{x}\lambda \cong \mathbf{A}\mathbf{x}$$

- ▶ From normal equation  $\mathbf{x}^T \mathbf{x} \lambda = \mathbf{x}^T \mathbf{A} \mathbf{x}$ , least squares solution is given by

$$\lambda = \frac{\mathbf{x}^T \mathbf{A} \mathbf{x}}{\mathbf{x}^T \mathbf{x}}$$

- ▶ This quantity, known as *Rayleigh quotient*, has many useful properties



## Example: Rayleigh Quotient

- ▶ Rayleigh quotient can accelerate convergence of iterative methods such as power iteration, since Rayleigh quotient  $\mathbf{x}_k^T \mathbf{A} \mathbf{x}_k / \mathbf{x}_k^T \mathbf{x}_k$  gives better approximation to eigenvalue at iteration  $k$  than does basic method alone
- ▶ For previous example using power iteration, value of Rayleigh quotient at each iteration is shown below

$k$	$\mathbf{x}_k^T$		$\ \mathbf{y}_k\ _\infty$	$\mathbf{x}_k^T \mathbf{A} \mathbf{x}_k / \mathbf{x}_k^T \mathbf{x}_k$
0	0.000	1.0		
1	0.333	1.0	1.500	1.500
2	0.600	1.0	1.667	1.800
3	0.778	1.0	1.800	1.941
4	0.882	1.0	1.889	1.985
5	0.939	1.0	1.941	1.996
6	0.969	1.0	1.970	1.999

## Rayleigh Quotient Iteration

- ▶ Given approximate eigenvector, Rayleigh quotient yields good estimate for corresponding eigenvalue
- ▶ Conversely, inverse iteration converges rapidly to eigenvector if approximate eigenvalue is used as shift, with one iteration often sufficing
- ▶ These two ideas combined in *Rayleigh quotient iteration*

$$\sigma_k = \mathbf{x}_k^T \mathbf{A} \mathbf{x}_k / \mathbf{x}_k^T \mathbf{x}_k$$

$$(\mathbf{A} - \sigma_k \mathbf{I}) \mathbf{y}_{k+1} = \mathbf{x}_k$$

$$\mathbf{x}_{k+1} = \mathbf{y}_{k+1} / \|\mathbf{y}_{k+1}\|_\infty$$

starting from given nonzero vector  $\mathbf{x}_0$

## Example: Rayleigh Quotient Iteration

- ▶ Using same matrix as previous examples and randomly chosen starting vector  $\mathbf{x}_0$ , Rayleigh quotient iteration converges in two iterations

$k$	$\mathbf{x}_k^T$		$\sigma_k$
0	0.807	0.397	1.896
1	0.924	1.000	1.998
2	1.000	1.000	2.000

## Deflation

# Deflation

- ▶ After eigenvalue  $\lambda_1$  and corresponding eigenvector  $\mathbf{x}_1$  have been computed, then additional eigenvalues  $\lambda_2, \dots, \lambda_n$  of  $\mathbf{A}$  can be computed by *deflation*, which effectively removes known eigenvalue
- ▶ Let  $\mathbf{H}$  be any nonsingular matrix such that  $\mathbf{H}\mathbf{x}_1 = \alpha\mathbf{e}_1$ , scalar multiple of first column of identity matrix (Householder transformation is good choice for  $\mathbf{H}$ )
- ▶ Then similarity transformation determined by  $\mathbf{H}$  transforms  $\mathbf{A}$  into form

$$\mathbf{H}\mathbf{A}\mathbf{H}^{-1} = \begin{bmatrix} \lambda_1 & \mathbf{b}^T \\ \mathbf{0} & \mathbf{B} \end{bmatrix} \begin{matrix} \mathbf{x}_1 \\ \vdots \\ \mathbf{x}_n \end{matrix} = \lambda_1 \mathbf{x}_1 + \mathbf{B} \begin{matrix} \vdots \\ \mathbf{x}_n \end{matrix}$$

where  $\mathbf{B}$  is matrix of order  $n - 1$  having eigenvalues  $\lambda_2, \dots, \lambda_n$

$$\begin{pmatrix} \alpha & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 \end{pmatrix}$$

## Deflation, continued

- ▶ Alternative approach lets  $\mathbf{u}_1$  be any vector such that  $\mathbf{u}_1^T \mathbf{x}_1 = \lambda_1$
- ▶ Then  $\mathbf{A} - \mathbf{x}_1 \mathbf{u}_1^T$  has eigenvalues  $0, \lambda_2, \dots, \lambda_n$
- ▶ Possible choices for  $\mathbf{u}_1$  include
  - ▶  $\mathbf{u}_1 = \lambda_1 \mathbf{x}_1$ , if  $\mathbf{A}$  is symmetric and  $\mathbf{x}_1$  is normalized so that  $\|\mathbf{x}_1\|_2 = 1$
  - ▶  $\mathbf{u}_1 = \lambda_1 \mathbf{y}_1$ , where  $\mathbf{y}_1$  is corresponding left eigenvector (i.e.,  $\mathbf{A}^T \mathbf{y}_1 = \lambda_1 \mathbf{y}_1$ ) normalized so that  $\mathbf{y}_1^T \mathbf{x}_1 = 1$
  - ▶  $\mathbf{u}_1 = \mathbf{A}^T \mathbf{e}_k$ , if  $\mathbf{x}_1$  is normalized so that  $\|\mathbf{x}_1\|_\infty = 1$  and  $k$ th component of  $\mathbf{x}_1$  is 1

## QR Iteration

## Simultaneous Iteration

- ▶ Simplest method for computing many eigenvalue-eigenvector pairs is *simultaneous iteration*, which repeatedly multiplies matrix times matrix of initial starting vectors

- ▶ Starting from  $n \times p$  matrix  $\mathbf{X}_0$  of rank  $p$ , iteration scheme is

$$\mathbf{X}_k = \mathbf{A}\mathbf{X}_{k-1}$$

- ▶  $\text{span}(\mathbf{X}_k)$  converges to invariant subspace determined by  $p$  largest eigenvalues of  $\mathbf{A}$ , provided  $|\lambda_p| > |\lambda_{p+1}|$
- ▶ Also called *subspace iteration*



# Orthogonal Iteration

$\begin{pmatrix} 1 & 1 & 1 \end{pmatrix}$

- ▶ As with power iteration, normalization is needed with simultaneous iteration
- ▶ Each column of  $\mathbf{X}_k$  converges to dominant eigenvector, so columns of  $\mathbf{X}_k$  become increasingly ill-conditioned basis for  $\text{span}(\mathbf{X}_k)$
- ▶ Both issues can be addressed by computing QR factorization at each iteration

$\begin{pmatrix} 1 & 1 & 1 \\ 0 & 0 & 0 \end{pmatrix}$

$$\begin{aligned}\hat{\mathbf{Q}}_k \mathbf{R}_k &= \mathbf{X}_{k-1} \\ \mathbf{X}_k &= \mathbf{A} \hat{\mathbf{Q}}_k\end{aligned}$$

where  $\hat{\mathbf{Q}}_k \mathbf{R}_k$  is **reduced** QR factorization of  $\mathbf{X}_{k-1}$

- ▶ This *orthogonal iteration* converges to block triangular form, and leading block is triangular if moduli of consecutive eigenvalues are distinct

## QR Iteration

- ▶ For  $p = n$  and  $\mathbf{X}_0 = \mathbf{I}$ , matrices

$$\mathbf{A}_k = \hat{\mathbf{Q}}_k^H \mathbf{A} \hat{\mathbf{Q}}_k$$

generated by orthogonal iteration converge to triangular or block triangular form, yielding all eigenvalues of  $\mathbf{A}$

- ▶ *QR iteration* computes successive matrices  $\mathbf{A}_k$  without forming above product explicitly
- ▶ Starting with  $\mathbf{A}_0 = \mathbf{A}$ , at iteration  $k$  compute QR factorization

$$\mathbf{Q}_k \mathbf{R}_k = \mathbf{A}_{k-1}$$

and form reverse product

$$\mathbf{A}_k = \mathbf{R}_k \mathbf{Q}_k$$

## QR Iteration, continued

- ▶ Successive matrices  $\mathbf{A}_k$  are unitarily similar to each other

$$\mathbf{A}_k = \mathbf{R}_k \mathbf{Q}_k = \mathbf{Q}_k^H \mathbf{A}_{k-1} \mathbf{Q}_k$$

- ▶ Diagonal entries (or eigenvalues of diagonal blocks) of  $\mathbf{A}_k$  converge to eigenvalues of  $\mathbf{A}$
- ▶ Product of orthogonal matrices  $\mathbf{Q}_k$  converges to matrix of corresponding eigenvectors
- ▶ If  $\mathbf{A}$  is symmetric, then symmetry is preserved by QR iteration, so  $\mathbf{A}_k$  converge to matrix that is both triangular and symmetric, hence diagonal

## Example: QR Iteration

- ▶ Let  $A_0 = \begin{bmatrix} 7 & 2 \\ 2 & 4 \end{bmatrix}$
- ▶ Compute QR factorization

$$A_0 = Q_1 R_1 = \begin{bmatrix} .962 & -.275 \\ .275 & .962 \end{bmatrix} \begin{bmatrix} 7.28 & 3.02 \\ 0 & 3.30 \end{bmatrix}$$

and form reverse product

$$A_1 = R_1 Q_1 = \begin{bmatrix} 7.83 & .906 \\ .906 & 3.17 \end{bmatrix}$$

- ▶ Off-diagonal entries are now smaller, and diagonal entries closer to eigenvalues, 8 and 3
- ▶ Process continues until matrix is within tolerance of being diagonal, and diagonal entries then closely approximate eigenvalues