

# SI231 Matrix Computations

## Lecture 8: Least Squares Revisited

Ziping Zhao

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School of Information Science and Technology  
ShanghaiTech University, Shanghai, China

# Lecture 8: Least Squares Revisited

- Part I: regularization
- Part II: sparsity
  - $\ell_0$  minimization
  - greedy pursuit,  $\ell_1$  minimization, and variations
  - majorization-minimization for  $\ell_2$ - $\ell_1$  minimization
  - dictionary learning
- Part III: LS with errors in  $\mathbf{A}$ 
  - total LS
  - robust LS, and its equivalence to regularization

# Part I: Regularization

## Sensitivity to Noise

- **Question:** how sensitive is the LS solution when there is noise?
- **Model:**

$$\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} + \boldsymbol{\nu},$$

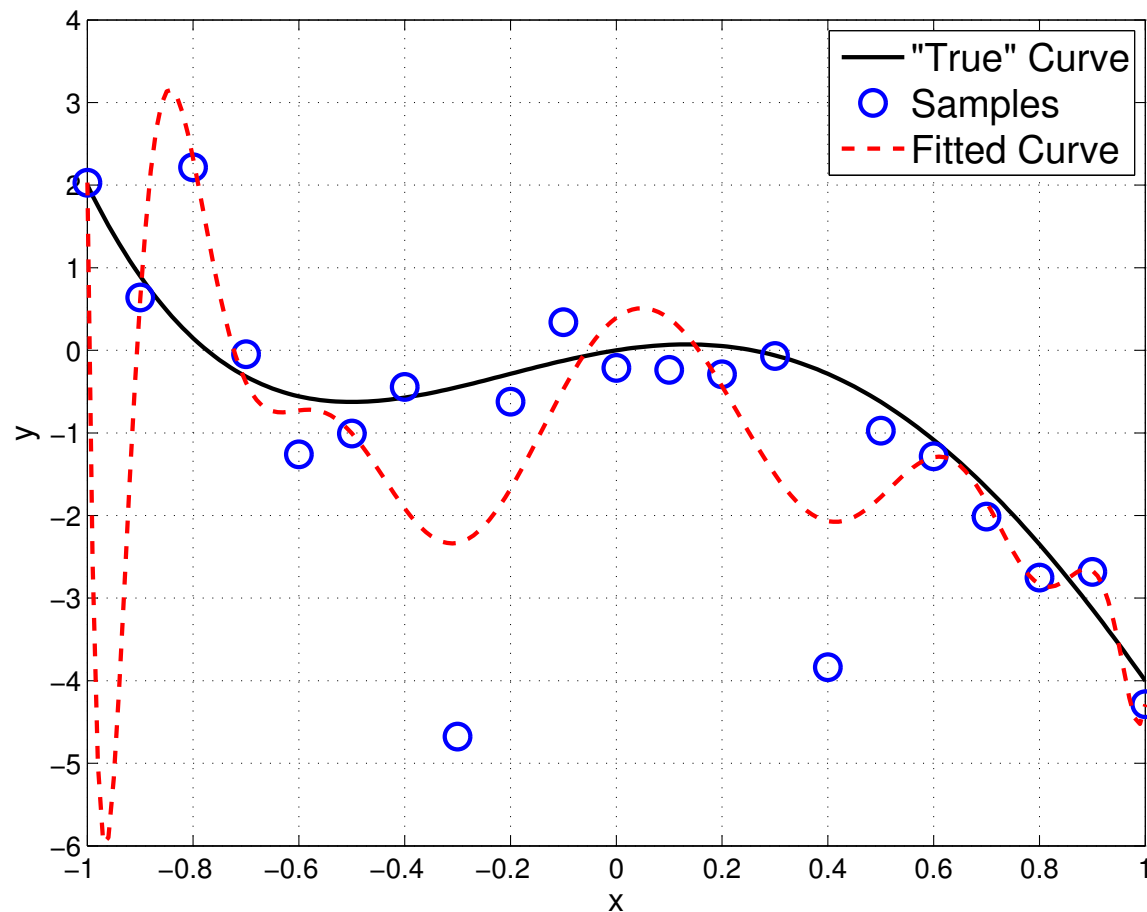
where  $\bar{\mathbf{x}}$  is the true result;  $\mathbf{A} \in \mathbb{R}^{m \times n}$  has full column rank;  $\boldsymbol{\nu}$  is noise, modeled as a random vector with mean zero and covariance  $\gamma^2 \mathbf{I}$ .

- **Mean square error (MSE) analysis:** from  $\mathbf{x}_{\text{LS}} = \mathbf{A}^\dagger \mathbf{y} = \bar{\mathbf{x}} + \mathbf{A}^\dagger \boldsymbol{\nu}$  we get

$$\begin{aligned} \mathbb{E}[\|\mathbf{x}_{\text{LS}} - \bar{\mathbf{x}}\|_2^2] &= \mathbb{E}[\|\mathbf{A}^\dagger \boldsymbol{\nu}\|_2^2] = \mathbb{E}[\text{tr}(\mathbf{A}^\dagger \boldsymbol{\nu} \boldsymbol{\nu}^T (\mathbf{A}^\dagger)^T)] = \text{tr}(\mathbf{A}^\dagger \mathbb{E}[\boldsymbol{\nu} \boldsymbol{\nu}^T] (\mathbf{A}^\dagger)^T) \\ &= \gamma^2 \text{tr}(\mathbf{A}^\dagger (\mathbf{A}^\dagger)^T) = \gamma^2 \text{tr}((\mathbf{A}^T \mathbf{A})^{-1}) \\ &= \gamma^2 \sum_{i=1}^n \frac{1}{\sigma_i^2(\mathbf{A})} \end{aligned}$$

- **Observation:** the MSE becomes very large if some  $\sigma_i(\mathbf{A})$ 's are close to zero.

## Toy Demonstration: Curve Fitting



The same curve fitting example in [Lecture 3: Least Squares](#). The “true” curve is the true  $f(x)$  with model order  $n = 4$ . In practice, the model order may not be known and we may have to guess. The fitted curve above is done by LS with a guessed model order  $n = 16$ .

## $\ell_2$ -Regularized LS

- **Intuition:** replace  $\mathbf{x}_{LS} = \mathbf{A}^\dagger \mathbf{y} = (\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T \mathbf{y}$  by

$$\mathbf{x}_{RLS} = (\mathbf{A}^T \mathbf{A} + \lambda \mathbf{I})^{-1} \mathbf{A}^T \mathbf{y},$$

for some  $\lambda > 0$ , where the term  $\lambda \mathbf{I}$  is added to improve the system conditioning, thereby attempting to reduce noise sensitivity

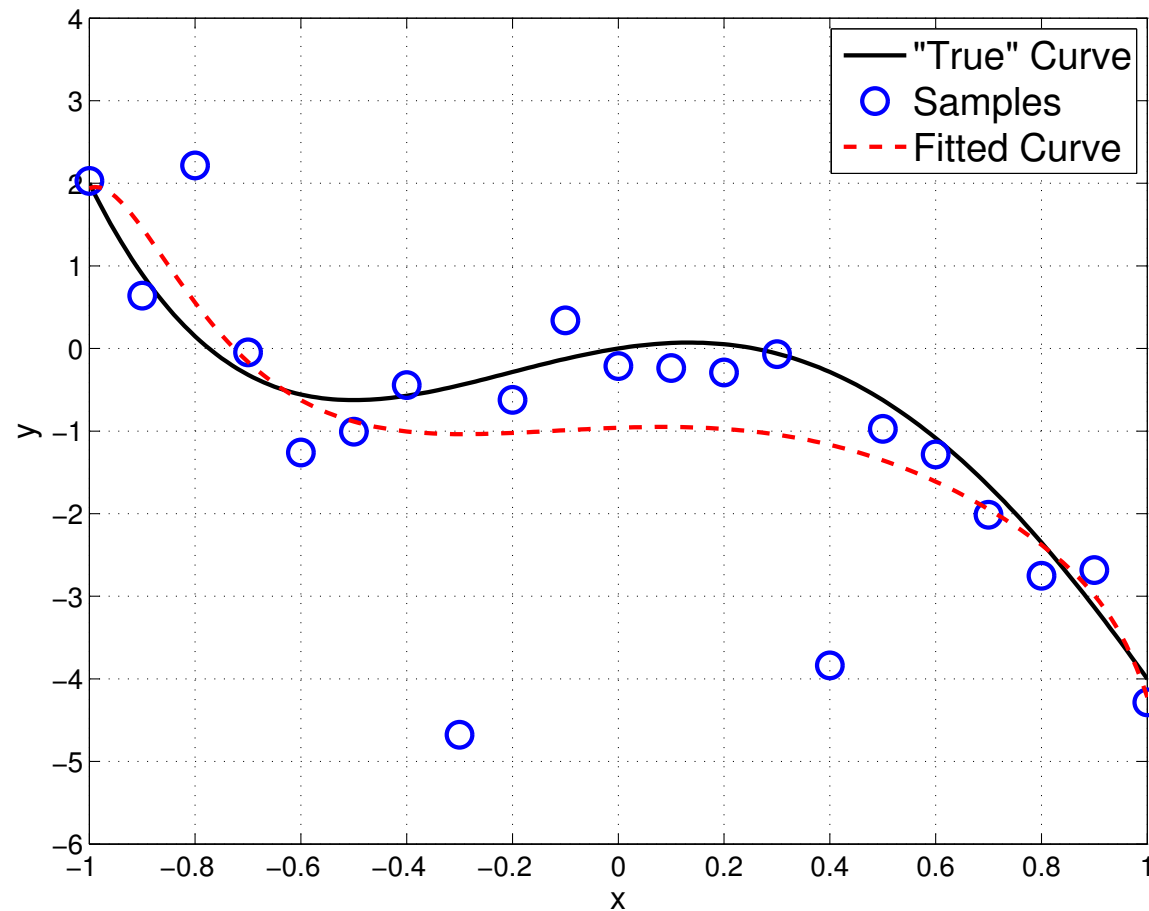
- how may we make sense out of such a modification?
- $\ell_2$ -regularized LS: find an  $\mathbf{x}$  that solves

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 + \lambda \|\mathbf{x}\|_2^2$$

for some pre-determined  $\lambda > 0$ .

- the solution is uniquely given by  $\mathbf{x}_{RLS} = (\mathbf{A}^T \mathbf{A} + \lambda \mathbf{I})^{-1} \mathbf{A}^T \mathbf{y}$
- the formulation says that we try to minimize both  $\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2$  and  $\|\mathbf{x}\|_2^2$ , and  $\lambda$  controls which one should be more emphasized in the minimization

# Toy Demonstration: Curve Fitting



The fitted curve is done by  $\ell_2$ -regularized LS with a guessed model order  $n = 18$  and with  $\lambda = 0.1$ .

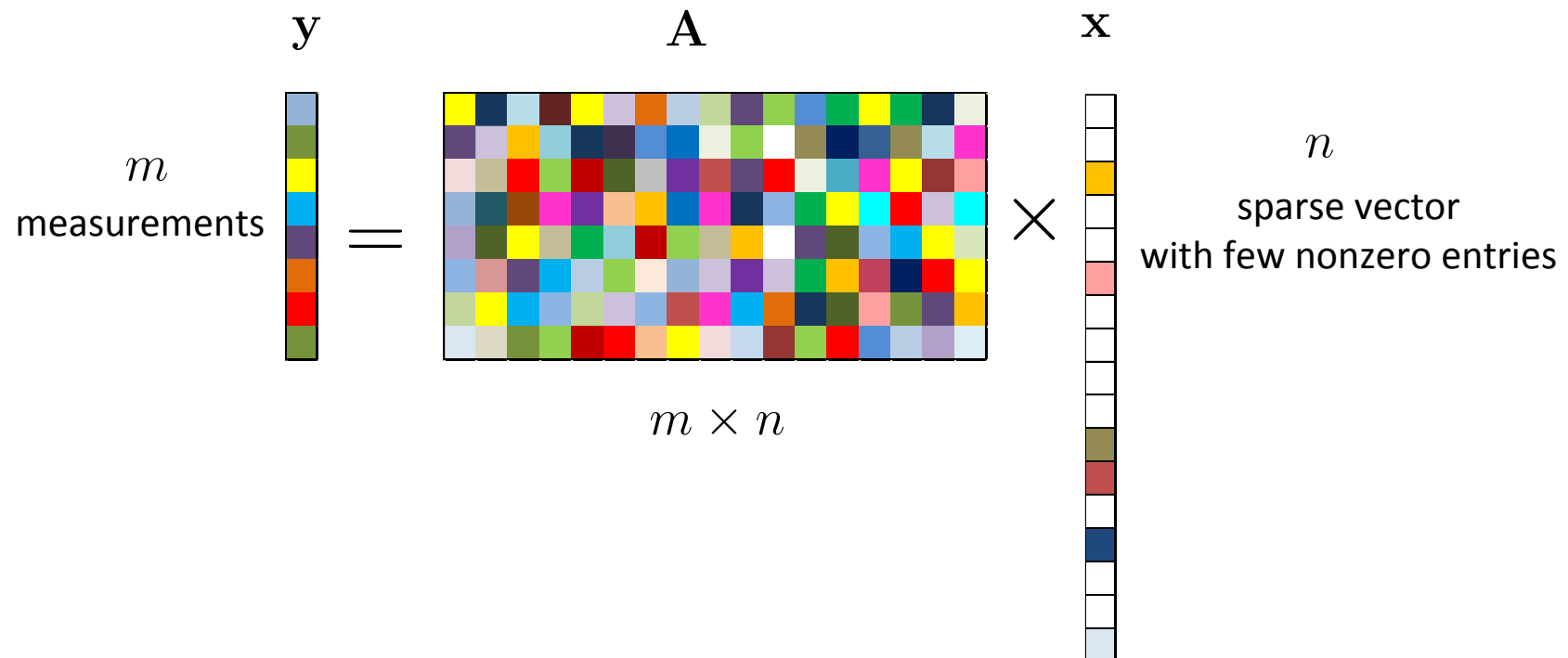
# Part II: Sparsity



# The Sparse Recovery Problem

**Problem:** given  $\mathbf{y} \in \mathbb{R}^m$ ,  $\mathbf{A} \in \mathbb{R}^{m \times n}$ ,  $m < n$ , find a **sparsest**  $\mathbf{x} \in \mathbb{R}^n$  such that

$$\mathbf{y} = \mathbf{A}\mathbf{x}.$$



- by sparsest, we mean that  $\mathbf{x}$  should have as many zero elements as possible.

# A Sparsity Optimization Formulation

- let

$$\|\mathbf{x}\|_0 = \sum_{i=1}^n \mathbb{1}\{x_i \neq 0\}$$

denote the cardinality function

– commonly called the “ $\ell_0$ -norm”, though it is not a norm.

- Minimum  $\ell_0$ -norm formulation:

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^n} \quad & \|\mathbf{x}\|_0 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}. \end{aligned} \tag{*}$$

- **Question:** suppose that  $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$ , where  $\bar{\mathbf{x}}$  is the vector we seek to recover. Can the min.  $\ell_0$ -norm problem recover  $\bar{\mathbf{x}}$  in an exact and unique fashion?
  - an answer lies in the notion of **spark**, which may be seen as a strong definition of rank

# Spark

**Spark:** the spark of  $\mathbf{A}$ , denoted by  $\text{spark}(\mathbf{A})$ , is the **minimal** number of **linearly dependent** columns of  $\mathbf{A}$ , i.e.,

$$\text{spark}(\mathbf{A}) = \min_{\mathbf{x} \neq \mathbf{0}} \|\mathbf{x}\|_0 \quad \text{s.t.} \quad \mathbf{A}\mathbf{x} = \mathbf{0}.$$

- let  $\text{spark}(\mathbf{A}) = k$ . Then,  $k$  is the smallest number such that there exists a linearly dependent  $\{\mathbf{a}_{i_1}, \dots, \mathbf{a}_{i_k}\}$  for some  $\{i_1, \dots, i_k\} \subseteq \{1, \dots, n\}$ <sup>1</sup>.
  - $\{\mathbf{a}_{i_1}, \dots, \mathbf{a}_{i_{k-1}}\}$  is linearly independent for any  $\{i_1, \dots, i_{k-1}\} \subseteq \{1, \dots, n\}$
- **Comparison with rank:** the rank of  $\mathbf{A}$ , denoted by  $\text{rank}(\mathbf{A})$ , is the **maximal** number of **linearly independent** columns of  $\mathbf{A}$ .
- let  $\text{rank}(\mathbf{A}) = r$ . Then,  $k$  is the largest number such that there exists a linearly independent  $\{\mathbf{a}_{i_1}, \dots, \mathbf{a}_{i_r}\}$  for some  $\{i_1, \dots, i_r\} \subseteq \{1, \dots, n\}$ .
  - $\{\mathbf{a}_{i_1}, \dots, \mathbf{a}_{i_{r+1}}\}$  is linearly dependent for any  $\{i_1, \dots, i_{r+1}\} \subseteq \{1, \dots, n\}$
- **Kruskal rank:** this is an alternative definition of rank. The Kruskal rank of  $\mathbf{A}$ , denoted by  $\text{krank}(\mathbf{A})$ , has its definition equivalent to  $\text{krank}(\mathbf{A}) = \text{spark}(\mathbf{A}) - 1$ .

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<sup>1</sup>We leave it implicit that  $i_k \neq i_j$  for any  $k \neq j$ .

## Spark

- if any collection of  $m$  vectors in  $\{\mathbf{a}_1, \dots, \mathbf{a}_n\} \subseteq \mathbb{R}^m$ , with  $n \geq m$ , is linearly independent, then

$$\text{spark}(\mathbf{A}) = m + 1, \quad \text{rank}(\mathbf{A}) = m.$$

- an example is Vandemonde matrices with distinct roots
- some specifically designed bases also have this property
- but there also exist instances in which rank and spark are very different
  - let  $\{\mathbf{v}_1, \dots, \mathbf{v}_r\} \in \mathbb{R}^m$  be linearly independent, and let  $\mathbf{A} = [\mathbf{v}_1, \dots, \mathbf{v}_r, \mathbf{v}_1]$ .
  - we have  $\text{rank}(\mathbf{A}) = r$ , but  $\text{spark}(\mathbf{A}) = 2$
- to conclude, spark may be seen as a stronger definition of rank, and

$$\text{krank}(\mathbf{A}) = \text{spark}(\mathbf{A}) - 1 \leq \text{rank}(\mathbf{A})$$

## Perfect Recovery Guarantee of the Min. $\ell_0$ -Norm Problem

**Theorem 8.1.** Suppose that  $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$ . Then,  $\bar{\mathbf{x}}$  is the unique solution to the minimum  $\ell_0$ -norm problem if

$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}\text{spark}(\mathbf{A}).$$

- **Implication:** any collection of  $2\|\bar{\mathbf{x}}\|_0$  columns of  $\mathbf{A}$  is linearly independent
  - for  $\bar{\mathbf{x}}'$  with  $\|\bar{\mathbf{x}}'\|_0 = \|\bar{\mathbf{x}}\|_0$ ,  $\mathbf{A}\bar{\mathbf{x}}' \neq \mathbf{A}\bar{\mathbf{x}}$
- **Implication:** if  $\bar{\mathbf{x}}$  is sufficiently sparse, then the minimum  $\ell_0$ -norm problem (\*) perfectly recovers  $\bar{\mathbf{x}}$
- **Proof sketch:**
  1. let  $\mathbf{x}^*$  be a solution to the min.  $\ell_0$ -norm problem. Let  $\mathbf{e} = \bar{\mathbf{x}} - \mathbf{x}^*$ .
  2.  $\mathbf{0} = \mathbf{A}\bar{\mathbf{x}} - \mathbf{A}\mathbf{x}^* = \mathbf{A}\mathbf{e}$ ;  $\|\mathbf{e}\|_0 \leq \|\bar{\mathbf{x}}\|_0 + \|\mathbf{x}^*\|_0 \leq 2\|\bar{\mathbf{x}}\|_0$ .
  3.  $\mathbf{A}\mathbf{e} = \mathbf{0}$ ,  $\|\mathbf{e}\|_0 \leq 2\|\bar{\mathbf{x}}\|_0 \implies \text{spark}(\mathbf{A}) \leq \|\mathbf{e}\|_0 \leq 2\|\bar{\mathbf{x}}\|_0$

# Perfect Recovery Guarantee of the Min. $\ell_0$ -Norm Problem

- **coherence:** the coherence of  $\mathbf{A}$  is defined as

$$\mu(\mathbf{A}) = \max_{j \neq k} \frac{|\mathbf{a}_j^T \mathbf{a}_k|}{\|\mathbf{a}_j\|_2 \|\mathbf{a}_k\|_2}.$$

– measures how similar the columns of  $\mathbf{A}$  are in the worst-case sense.

- a weaker version of Theorem 8.1:

**Corollary 8.1.** Suppose that  $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$ . Then,  $\bar{\mathbf{x}}$  is the unique solution to the minimum  $\ell_0$ -norm problem if

$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}(1 + \mu(\mathbf{A})^{-1}).$$

- **Implication:** perfect recovery may depend on how incoherent  $\mathbf{A}$  is.
- proof idea: show that  $\text{spark}(\mathbf{A}) \geq 1 + \mu(\mathbf{A})^{-1}$

# On Solving the Minimum $\ell_0$ -Norm Problem

**Question:** How should we solve the minimum  $\ell_0$ -norm problem

$$\begin{aligned} \min_{\mathbf{x}} \quad & \|\mathbf{x}\|_0 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}, \end{aligned}$$

or can it be efficiently solved?

- $\ell_0$ -norm minimization does not lead to a simple solution as in 2-norm min.
- the minimum  $\ell_0$ -norm problem is **NP-hard** in general
  - what does that mean?
    - \* given any  $\mathbf{y}, \mathbf{A}$ , the problem is unlikely to be exactly solvable in polynomial time (i.e., in a complexity of  $\mathcal{O}(n^p)$  for any  $p > 0$ )

## Brute Force Search for the Minimum $\ell_0$ -Norm Problem

- notation:  $\mathbf{A}_{\mathcal{I}}$  denotes a submatrix of  $\mathbf{A}$  obtained by keeping the columns indicated by  $\mathcal{I}$
- we may solve the  $\ell_0$ -norm minimization problem via brute force search:

**input:**  $\mathbf{A}, \mathbf{y}$   
for all  $\mathcal{I} \subseteq \{1, 2, \dots, n\}$  do  
    if  $\mathbf{y} = \mathbf{A}_{\mathcal{I}}\tilde{\mathbf{x}}$  has a solution for some  $\tilde{\mathbf{x}} \in \mathbb{R}^{|\mathcal{I}|}$   
        record  $(\tilde{\mathbf{x}}, \mathcal{I})$  as one of candidate solutions  
end  
**output:** a candidate solution  $(\tilde{\mathbf{x}}, \mathcal{I})$  whose  $|\mathcal{I}|$  is the smallest

- example: for  $n = 3$ , we test  $\mathcal{I} = \{1\}, \mathcal{I} = \{2\}, \mathcal{I} = \{3\}, \mathcal{I} = \{1, 2\}, \mathcal{I} = \{2, 3\}, \mathcal{I} = \{1, 3\}, \mathcal{I} = \{1, 2, 3\}$
- manageable for very small  $n$ , too expensive even for moderate  $n$
- how about a greedy search that searches less?



# Greedy Pursuit

- consider a greedy search called the **orthogonal matching pursuit (OMP)**

**Algorithm:** OMP

**input:**  $\mathbf{A}, \mathbf{y}$

set  $\mathcal{I} = \emptyset, \hat{\mathbf{x}} = \mathbf{0}$

repeat

$$\mathbf{r} = \mathbf{y} - \mathbf{A}\hat{\mathbf{x}}$$

$$k = \arg \max_{j \in \{1, \dots, n\}} |\mathbf{a}_j^T \mathbf{r}| / \|\mathbf{a}_j\|_2$$

$$\mathcal{I} := \mathcal{I} \cup \{k\}$$

$$\hat{\mathbf{x}} := \arg \min_{\mathbf{x} \in \mathbb{R}^n, x_i=0 \ \forall i \notin \mathcal{I}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2$$

until a stopping rule is satisfied, e.g.,  $\|\mathbf{y} - \mathbf{A}\hat{\mathbf{x}}\|_2$  is sufficiently small

**output:**  $\hat{\mathbf{x}}$

- note: there are many other greedy search strategies

# Perfect Recovery Guarantee of Greedy Pursuit

- again, a key question is the conditions under which OMP admits perfect recovery
- there are many such theoretical conditions, not only for OMP but also for other greedy algorithms
- one such result is as follows:

**Theorem 8.2.** Suppose that  $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$ . Then, OMP recovers  $\bar{\mathbf{x}}$  if

$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}(1 + \mu(\mathbf{A})^{-1}).$$

- proof idea: show that OMP is guaranteed to pick a correct column at every stage.

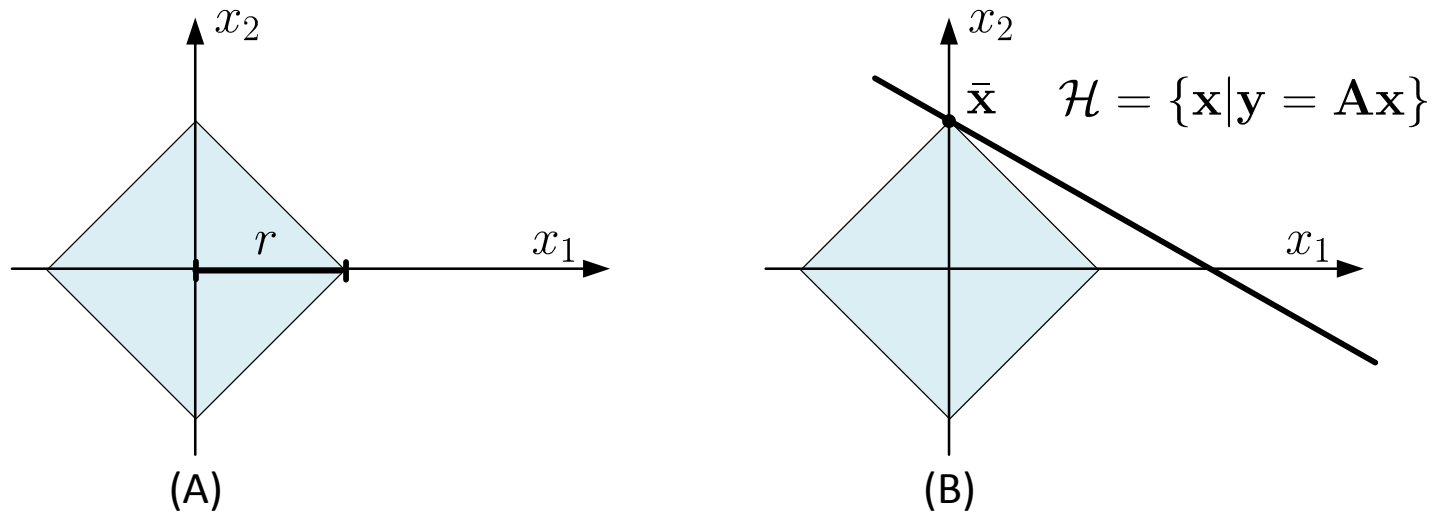
# Convex Relaxation

Another approximation approach is to replace  $\|\mathbf{x}\|_0$  by a convex function:

$$\begin{aligned} \min_{\mathbf{x}} \quad & \|\mathbf{x}\|_1 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}. \end{aligned}$$

- also known as basis pursuit in the literature
- convex, a linear program
- no closed-form solution (while the minimum 2-norm problem has)
- but the success of this minimum 1-norm problem, both in theory and practice, has motivated a large body of work on computationally efficient algorithms for it

## Illustration of 1-Norm Geometry



- Fig. A shows the 1-norm ball of radius  $r$  in  $\mathbb{R}^2$ . Note that the 1-norm ball is “pointy” along the axes.
- Fig. B shows the 1-norm recovery solution. The point  $\bar{\mathbf{x}}$  is a “sparse” vector; the line  $\mathcal{H}$  is the set of all  $\mathbf{x}$  that satisfy  $\mathbf{y} = \mathbf{A}\mathbf{x}$ .

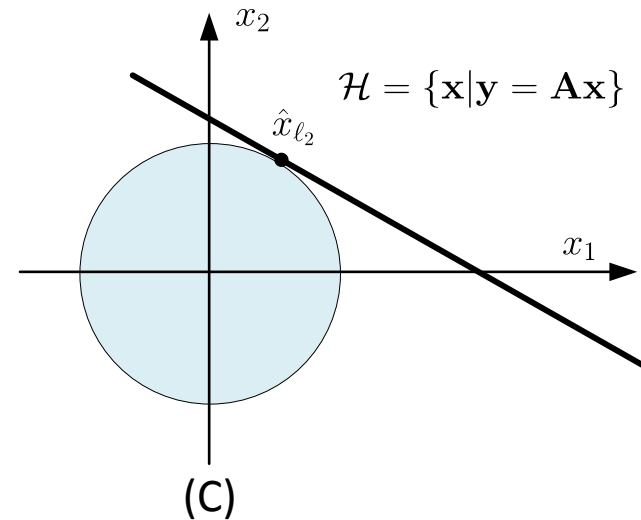
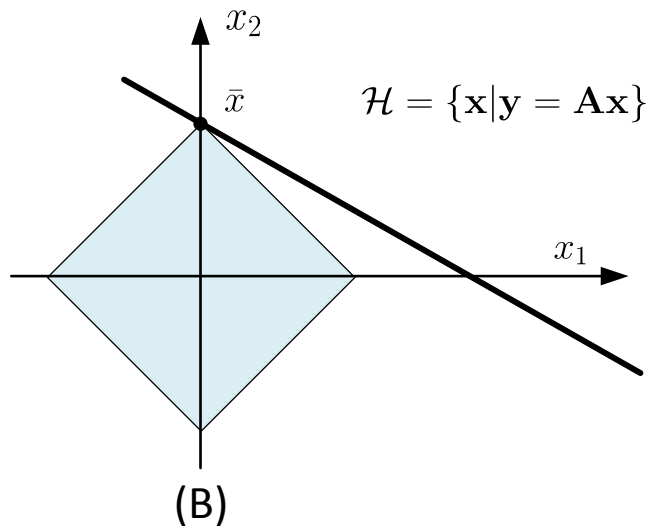
# Convex Relaxation

if replace  $\|\mathbf{x}\|_0$  by the  $\|\mathbf{x}\|_2$ :

$$\begin{aligned} \min_{\mathbf{x}} \quad & \|\mathbf{x}\|_2 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}. \end{aligned}$$

- also known as method of frames
- convex, a quadratic program
- closed-form solution (the minimum energy solution)
- but cannot promote sparsity

## Illustration of 1-Norm Geometry



- The 1-norm recovery problem is to pick out a point in  $\mathcal{H}$  that has the minimum 1-norm. We can see that  $\bar{\mathbf{x}}$  is such a point.
- Fig. C shows the geometry when 2-norm is used. We can see that the solution  $\hat{\mathbf{x}}$  may not be sparse.

# Perfect Recovery Guarantee of the Min. 1-Norm Problem

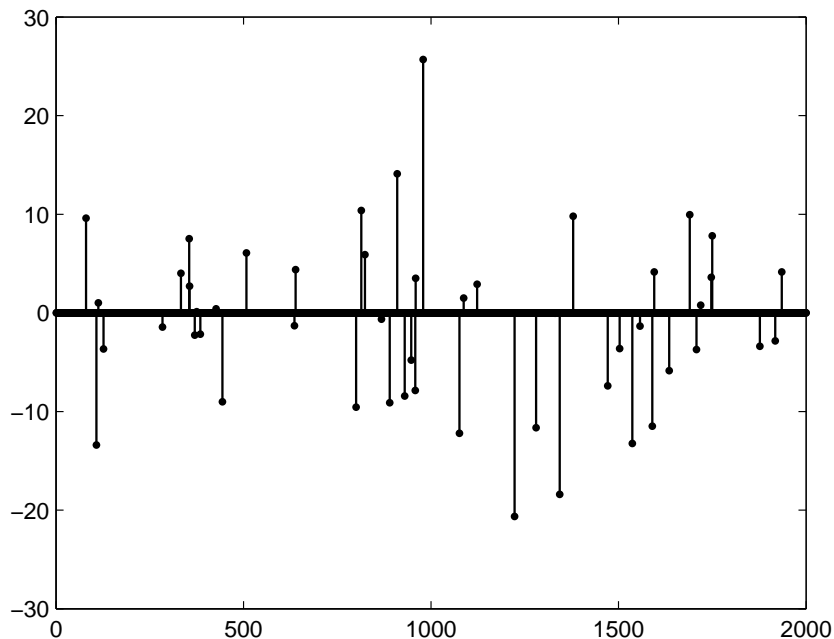
- again, researchers studied conditions under which the minimum 1-norm problem admits perfect recovery
- this has been an exciting topic, with many provable conditions such as the restricted isometry property (RIP), the nullspace property (NSP), ...
  - see the literature for details, and here is one: [\[Yin'13\]](#)
- a simple one is as follows:

**Theorem 8.3.** Suppose that  $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$ . Then,  $\bar{\mathbf{x}}$  is the unique solution to the minimum 1-norm problem if

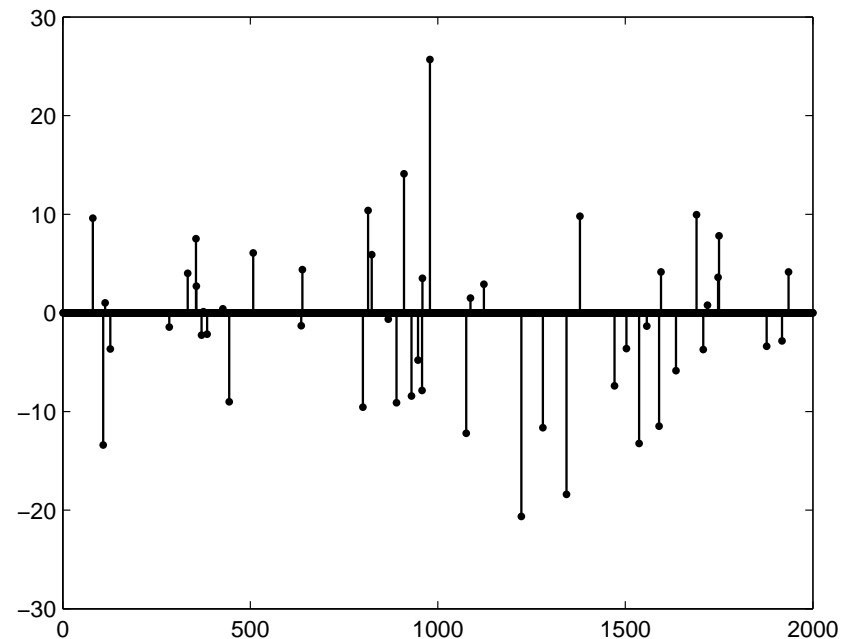
$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}(1 + \mu(\mathbf{A})^{-1}).$$

# Toy Demonstration: Sparse Signal Reconstruction

- Sparse vector  $\mathbf{x} \in \mathbb{R}^n$  with  $n = 2000$  and  $\|\mathbf{x}\|_0 = 50$ .
- $m = 400$  noise-free observations of  $\mathbf{y} = \mathbf{A}\mathbf{x}$ ,  $a_{ij}$  is randomly generated.

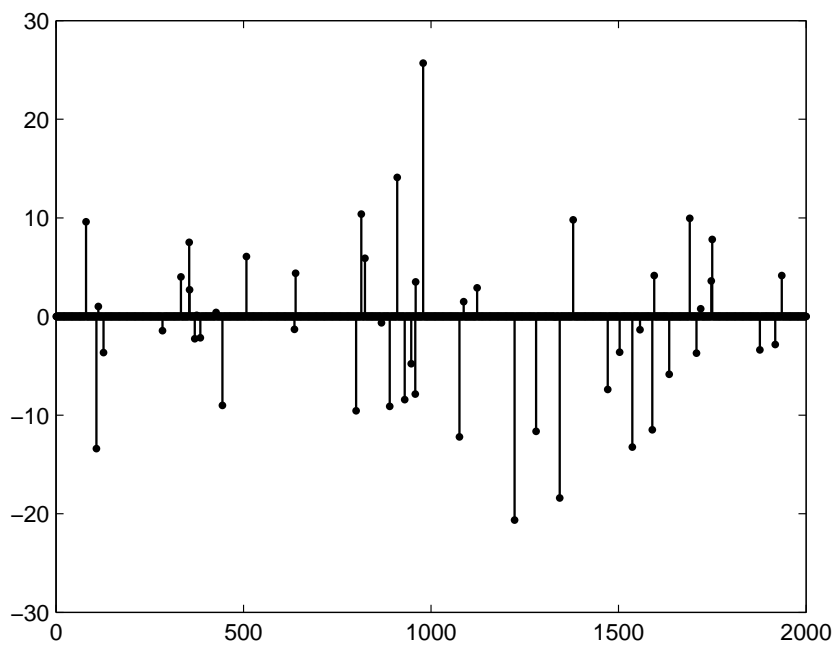


(a) Sparse source signal

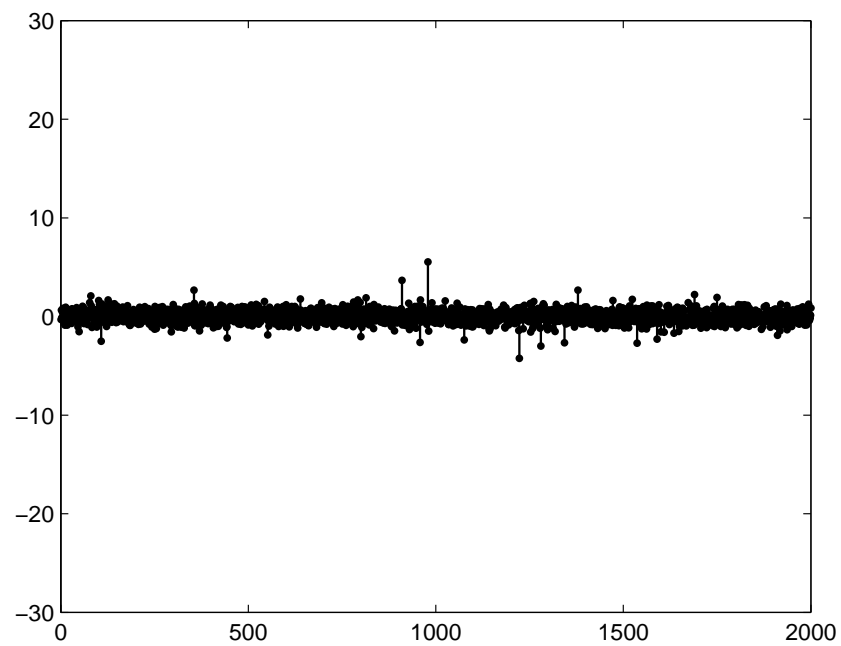


(b) Recovery by 1-norm minimization





(c) Sparse source signal



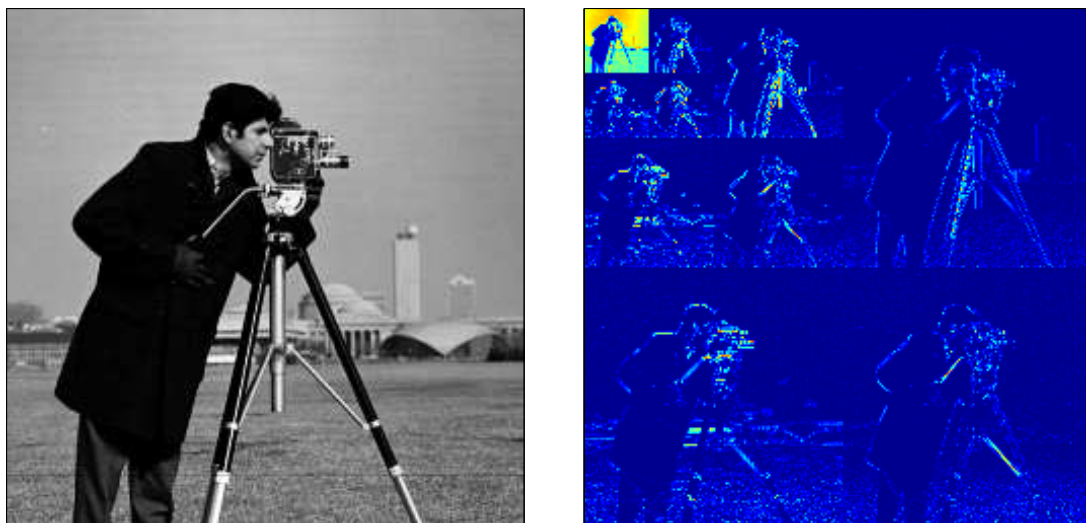
(d) Recovery by 2-norm minimization

## Application: Compressive sensing (CS)

- Consider a signal  $\tilde{\mathbf{x}} \in \mathbb{R}^n$  that has a sparse representation  $\mathbf{x} \in \mathbb{R}^n$  in the domain of the representation matrix  $\Psi \in \mathbb{R}^{n \times n}$  (e.g. DCT or wavelet), i.e.,

$$\tilde{\mathbf{x}} = \Psi \mathbf{x},$$

where  $\mathbf{x}$  is sparse.



Left: the original image  $\tilde{\mathbf{x}}$ . Right: the corresponding coefficient  $\mathbf{x}$  in the wavelet domain, which is sparse. Source: [\[Romberg-Wakin'07\]](#)

- compressive sensing is also called [compressive sampling](#)

## Application: CS

- To acquire  $\mathbf{x}$ , we use a sensing matrix  $\Phi \in \mathbb{R}^{m \times n}$  to observe  $\mathbf{x}$

$$\mathbf{y} = \Phi \tilde{\mathbf{x}} = \Phi \Psi \mathbf{x}.$$

Here, we have  $m \ll n$ , i.e., much fewer observations than the no. of unknowns

- Such a  $\mathbf{y}$  will be good for compression, transmission and storage.
- $\tilde{\mathbf{x}}$  is recovered by recovering  $\mathbf{x}$ :

$$\begin{aligned} \min \quad & \|\mathbf{x}\|_0 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}, \end{aligned}$$

where  $\mathbf{A} = \Phi \Psi$

- how to choose  $\Phi$ ? CS research suggests that i.i.d. random  $\Phi$  (a universal sensing matrix) will work well!

## Application: CS

$$\begin{aligned}
 y_1 &= \langle \text{img}_1, \text{noise}_1 \rangle \\
 y_2 &= \langle \text{img}_2, \text{noise}_2 \rangle \\
 y_3 &= \langle \text{img}_3, \text{noise}_3 \rangle \\
 &\vdots \\
 y_M &= \langle \text{img}_M, \text{noise}_M \rangle
 \end{aligned}$$



original (25k wavelets)

(b) original image



perfect recovery

(c)  $\ell_1$  recovery

(a) measurements ( $y_i = \langle \tilde{\mathbf{x}}, \Phi(i, :) \rangle$ ) via i.i.d. random  $\Phi$

Source: [\[Romberg-Wakin'07\]](#)

## Variations

- when  $\mathbf{y}$  is contaminated by noise, or when  $\mathbf{y} = \mathbf{A}\mathbf{x}$  does not exactly hold, some variants of the previous min. 1-norm formulation may be considered:

- **basis pursuit denoising**: given  $\epsilon > 0$ , solve

$$\min_{\mathbf{x}} \|\mathbf{x}\|_1 \quad \text{s.t.} \quad \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 \leq \epsilon$$

- **$\ell_1$ -regularized LS**: given  $\lambda > 0$ , solve

$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \lambda \|\mathbf{x}\|_1$$

- **Lasso**: given  $\tau > 0$ , solve

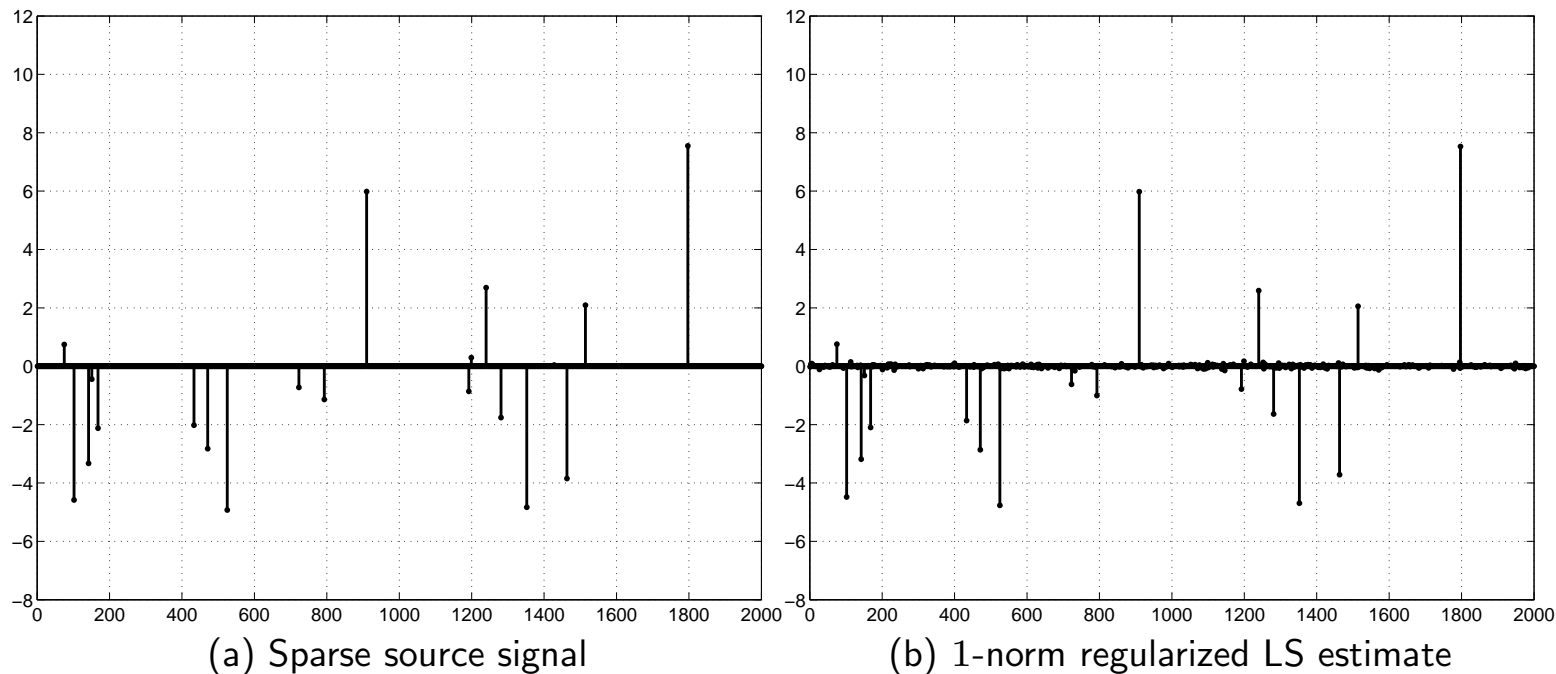
$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 \quad \text{s.t.} \quad \|\mathbf{x}\|_1 \leq \tau$$

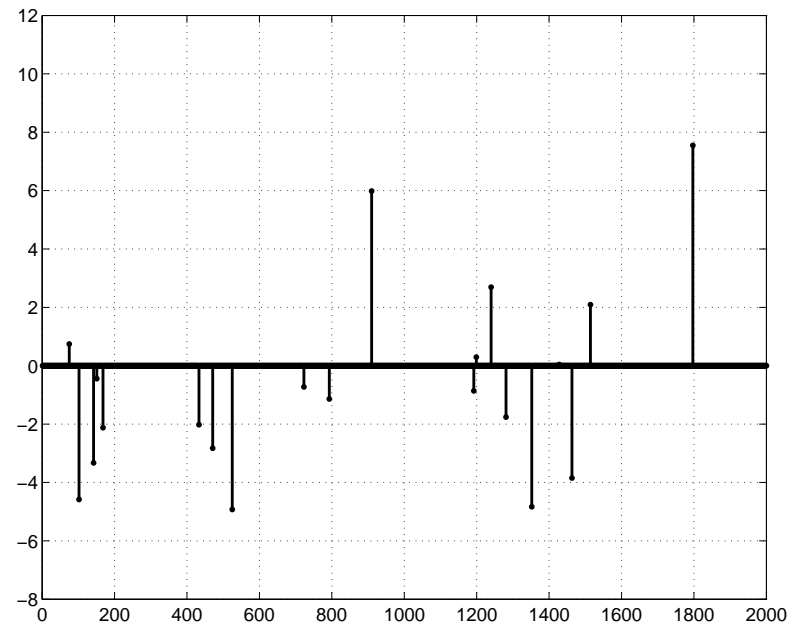
- when outliers exist in  $\mathbf{y}$  (i.e., some elements of  $\mathbf{y}$  are badly corrupted), we also want the residual  $\mathbf{r} = \mathbf{y} - \mathbf{A}\mathbf{x}$  to be sparse; so,

$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_1 + \lambda \|\mathbf{x}\|_1.$$

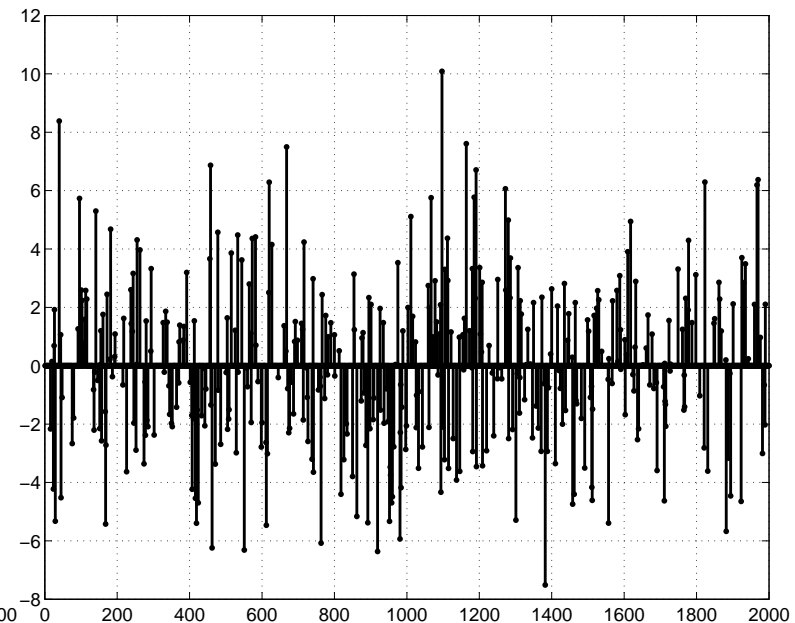
## Toy Demonstration: Noisy Sparse Signal Reconstruction

- Sparse signal  $\mathbf{x} \in \mathbb{R}^n$  with  $n = 2000$  and  $\|\mathbf{x}\|_0 = 20$ .
- $m = 400$  noisy observations of  $\mathbf{y} = \mathbf{A}\mathbf{x} + \boldsymbol{\nu}$ , both  $a_{ij}$  and  $\nu_i$  are randomly generated.
- 1-norm regularized LS  $\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \lambda \|\mathbf{x}\|_1$  is used with  $\lambda = 0.1$ .



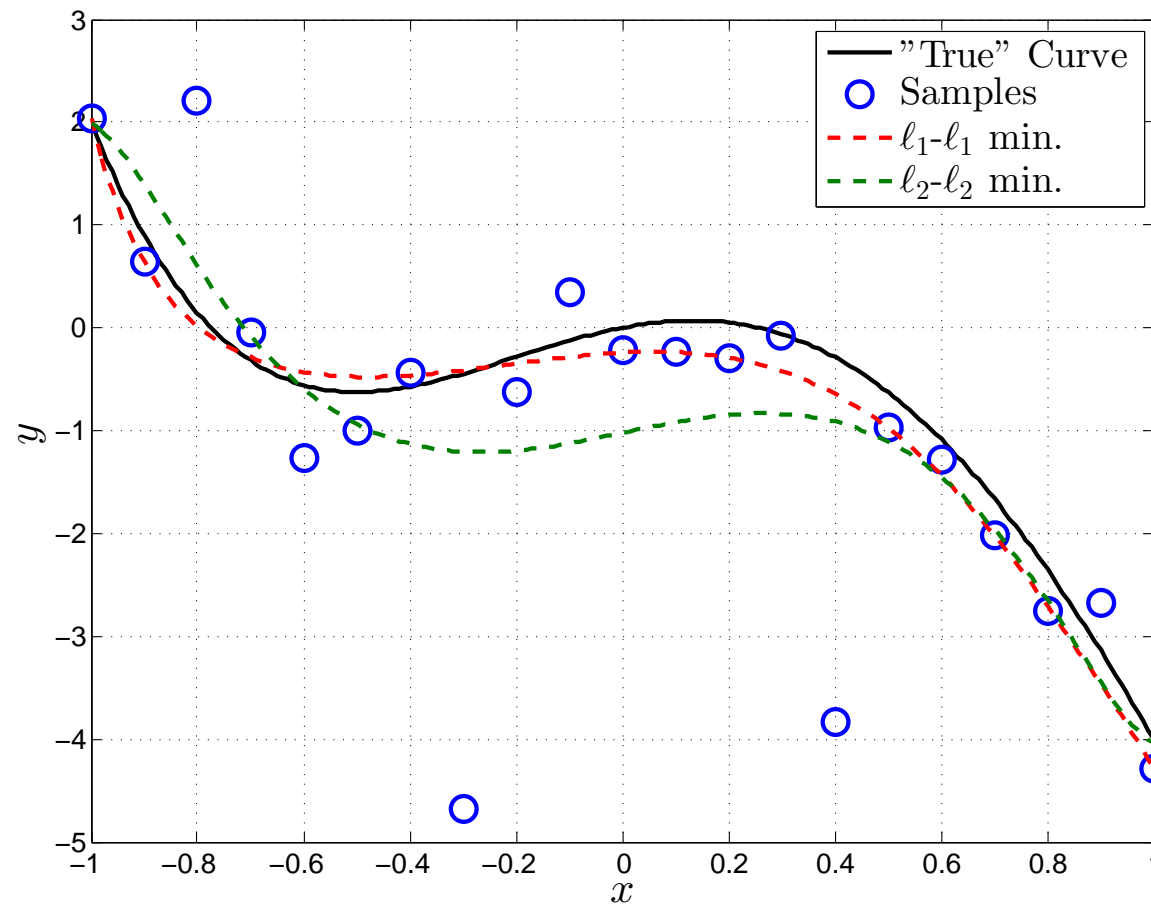


(c) Sparse source signal



(d) LS estimate

## Toy Demonstration: Curve Fitting



The same curve fitting problem in [Lecture 3: Least Squares](#). The guessed model order is  $n = 18$ .

$$\ell_2\text{-}\ell_2 \text{ min.: } \min \|\mathbf{y} - \mathbf{Ax}\|_2^2 + \lambda \|\mathbf{x}\|_2^2$$

$$\ell_1\text{-}\ell_1 \text{ min.: } \min \|\mathbf{y} - \mathbf{Ax}\|_1 + \lambda \|\mathbf{x}\|_1$$



# Total Variation (TV) Denoising

- Scenario:

- estimate  $\mathbf{x} \in \mathbb{R}^n$  from a noisy measurement  $\mathbf{x}_{\text{cor}} = \mathbf{x} + \boldsymbol{\nu}$ .
- $\mathbf{x}$  is known to be piecewise linear, i.e., for most  $i$  we have

$$x_i - x_{i-1} = x_{i+1} - x_i \iff -x_{i+1} + 2x_i - x_{i-1} = 0.$$

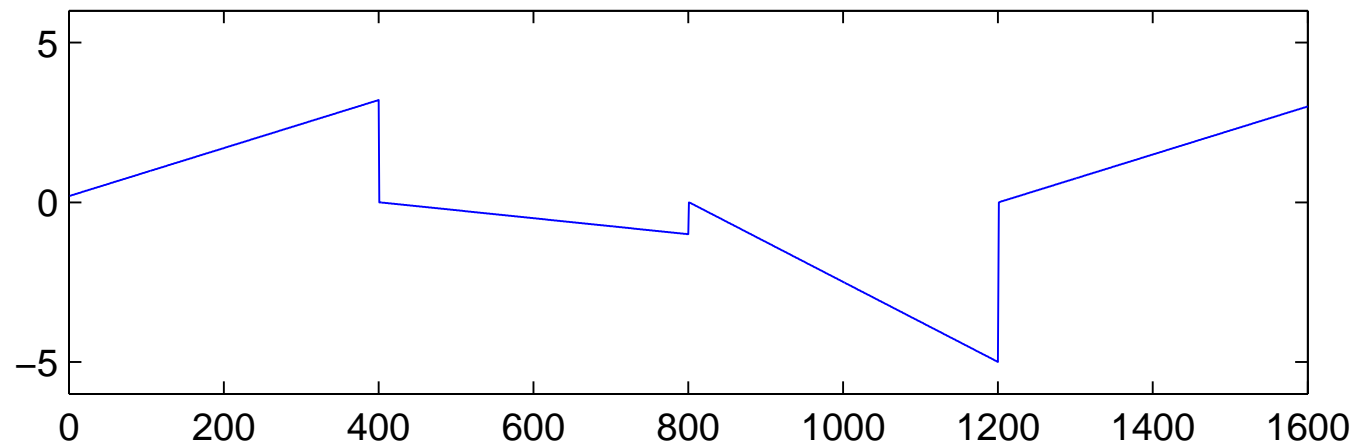
- equivalently,  $\mathbf{D}\mathbf{x}$  is sparse, where

$$\mathbf{D} = \begin{bmatrix} -1 & 2 & 1 & 0 & \dots \\ 0 & -1 & 2 & 1 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \dots & \dots & -1 & 2 & 1 \end{bmatrix}.$$

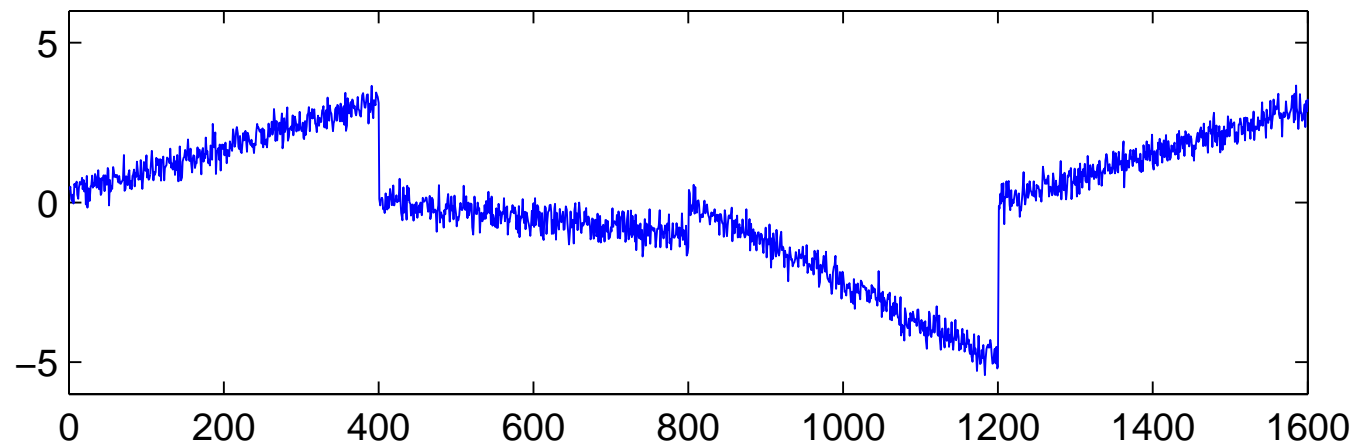
- TV denoising: estimate  $\mathbf{x}$  by solving

$$\min_{\mathbf{x}} \|\mathbf{x}_{\text{cor}} - \mathbf{x}\|_2^2 + \lambda \|\mathbf{D}\mathbf{x}\|_1$$

Source

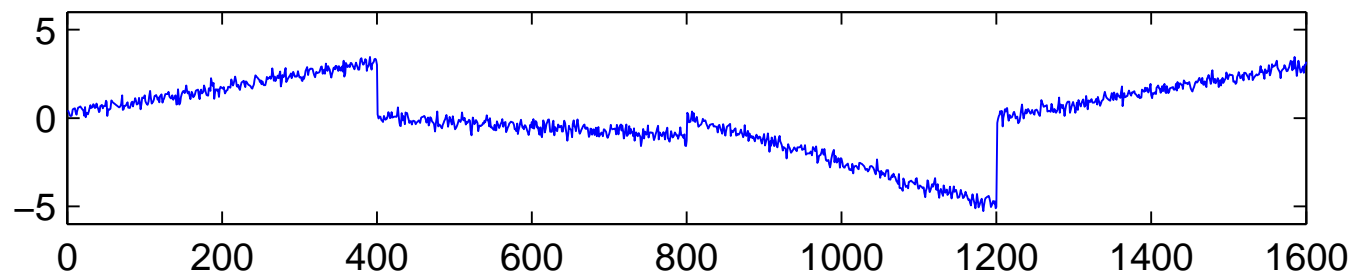


Corrupted by noise

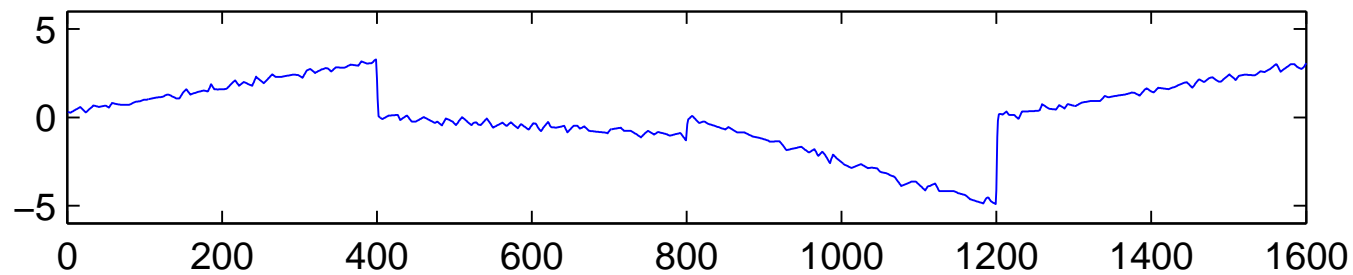


Original  $\mathbf{x}$  and corrupted  $\mathbf{x}_{\text{cor}}$

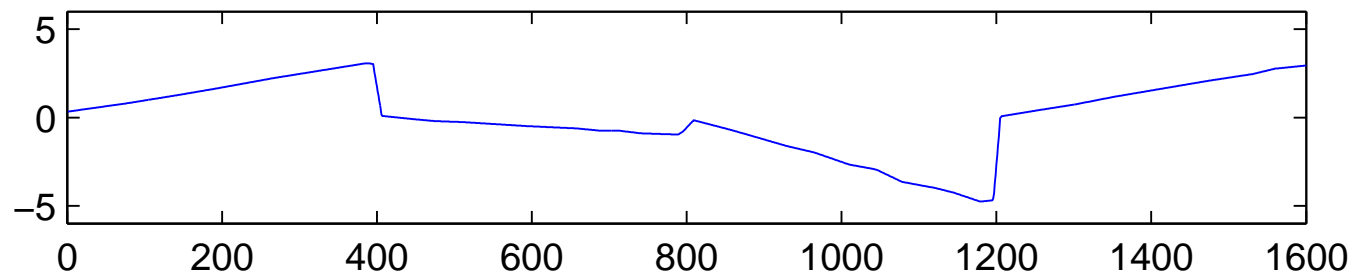
$\hat{x}$  with  $\lambda = 0.1$



$\hat{x}$  with  $\lambda = 1$

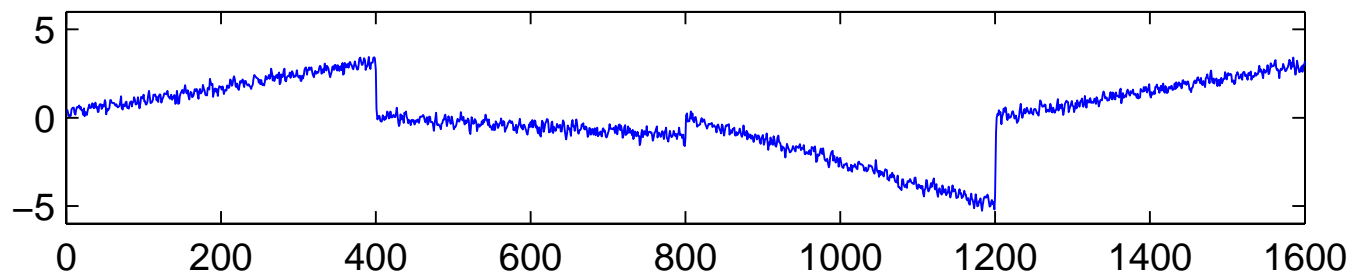


$\hat{x}$  with  $\lambda = 10$

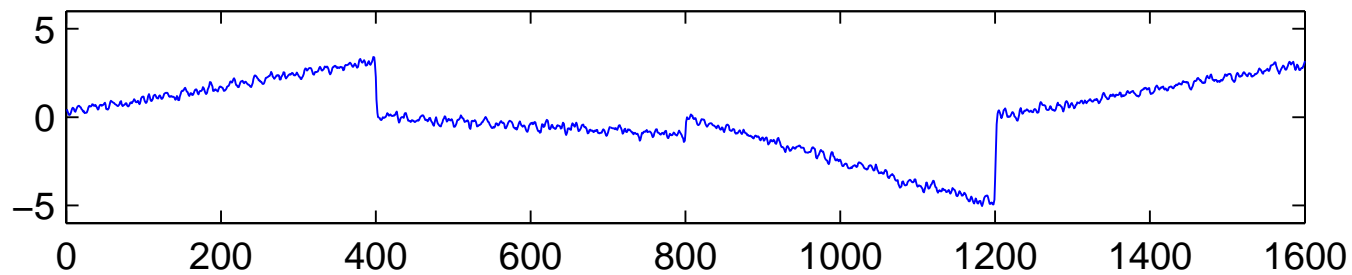


TV denoised signals for various  $\lambda$ 's.

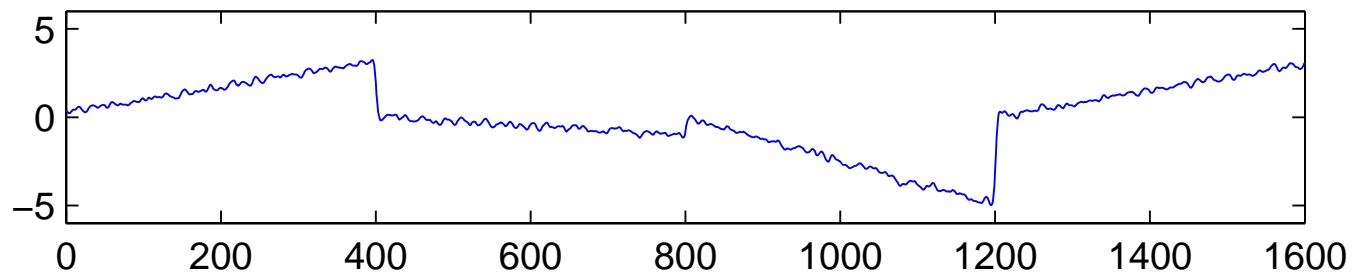
$\hat{x}$  with  $\lambda = 0.1$



$\hat{x}$  with  $\lambda = 1$



$\hat{x}$  with  $\lambda = 10$



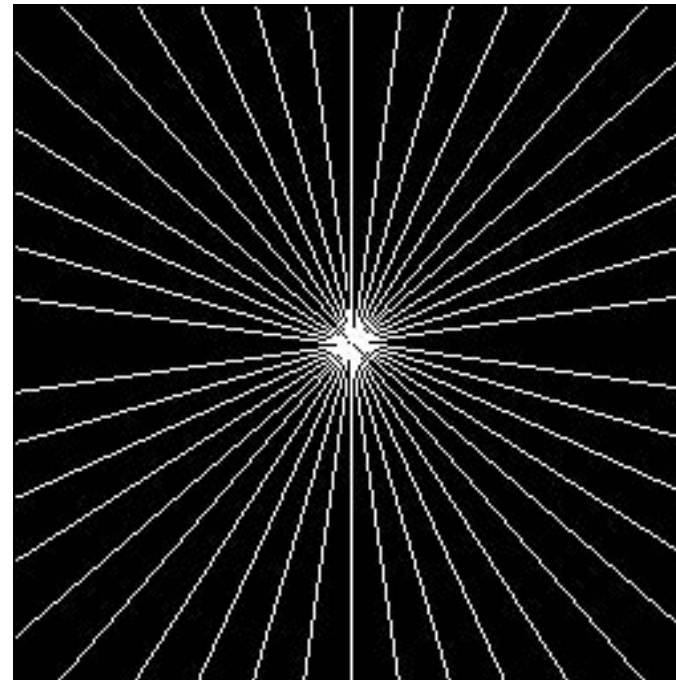
TV denoised signals via  $\ell_2$  regularization and for various  $\lambda$ 's.

# Application: Magnetic Resonance Imaging (MRI)

Problem: MRI image reconstruction.



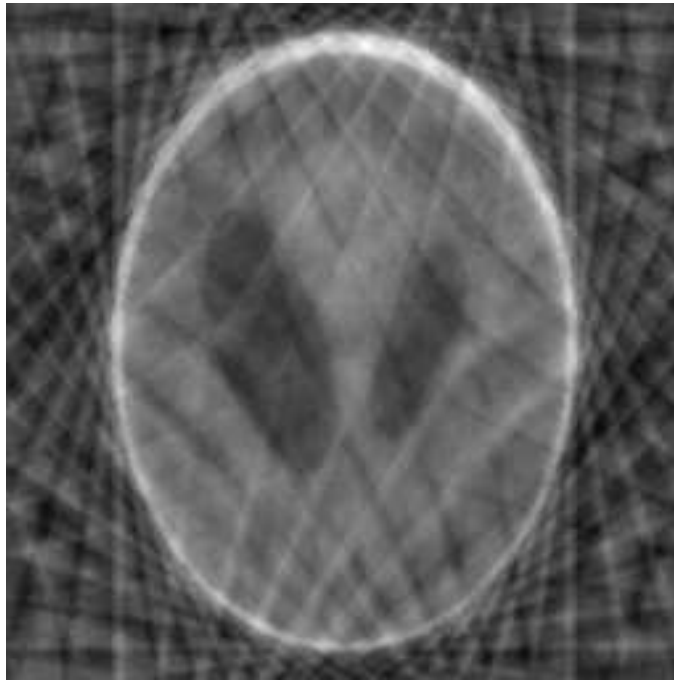
(a)



(b)

Fig. a shows the original test image. Fig. b shows the sampling region in the frequency domain. Fourier coefficients are sampled along 22 approximately radial lines. Source: [\[Candès-Romberg-Tao'06\]](#)

## Application: MRI



(c)



(d)

Fig. c is the recovery by filling the unobserved Fourier coefficients to zero. Fig. d is the recovery by a TV minimization problem. Source: [\[Candès-Romberg-Tao'06\]](#)

# Efficient Computations of the $\ell_2 - \ell_1$ Minimization Solution

- consider the  $\ell_2 - \ell_1$  minimization problem

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{y} - \mathbf{Ax}\|_2^2 + \lambda \|\mathbf{x}\|_1.$$

- as mentioned, the problem is convex and there are many optimization algorithms custom-designed for it
  - some keywords for such algorithms: majorization-minimization (MM), ADMM, fast proximal gradient (or the so-called FISTA), Frank-Wolfe,...
- Aim: get some flavor of one particular algorithm, namely, MM, that is sufficiently “matrix” and is suitable for large-scale problems

## MM for $\ell_2 - \ell_1$ Minimization: LS as an Example

- to see the insight of MM, we start with the plain old LS

$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2.$$

- observe that for a given  $\bar{\mathbf{x}}$ , one has

$$\begin{aligned} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 &= \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}} - \mathbf{A}(\mathbf{x} - \bar{\mathbf{x}})\|_2^2 \\ &= \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + \|\mathbf{A}(\mathbf{x} - \bar{\mathbf{x}})\|_2^2 \\ &\leq \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + c\|\mathbf{x} - \bar{\mathbf{x}}\|_2^2 \end{aligned}$$

for any  $\mathbf{x} \in \mathbb{R}^n$  and for any  $c \geq \sigma_{\max}^2(\mathbf{A})$



## MM for $\ell_2 - \ell_1$ Minimization: LS as an Example

- let  $c \geq \sigma_{\max}^2(\mathbf{A})$ , and let

$$g(\mathbf{x}, \bar{\mathbf{x}}) = \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + c\|\mathbf{x} - \bar{\mathbf{x}}\|_2^2$$

- we have

$$\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 \leq g(\mathbf{x}, \bar{\mathbf{x}}), \quad \text{for any } \mathbf{x}, \bar{\mathbf{x}} \in \mathbb{R}^n$$

$$\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 = g(\mathbf{x}, \mathbf{x}), \quad \text{for any } \mathbf{x} \in \mathbb{R}^n$$

- also,

$$\arg \min_{\mathbf{x} \in \mathbb{R}^n} g(\mathbf{x}, \bar{\mathbf{x}}) = \frac{1}{c} \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + \bar{\mathbf{x}}$$

- **Idea:** given an initial point  $\mathbf{x}^{(0)}$ , do

$$\mathbf{x}^{(k+1)} = \arg \min_{\mathbf{x} \in \mathbb{R}^n} g(\mathbf{x}, \mathbf{x}^{(k)}) = \frac{1}{c} \mathbf{A}^T (\mathbf{y} - \mathbf{A}\mathbf{x}^{(k)}) + \mathbf{x}^{(k)}, \quad k = 1, 2, \dots$$

- note: not very interesting at this moment as the above iteration is the same as gradient descent with step size  $1/c$

## MM for $\ell_2 - \ell_1$ Minimization: General MM Principle

- the example shown above is an instance of MM
- general MM principle:
  - consider a general optimization problem

$$\min_{\mathbf{x} \in \mathcal{C}} f(\mathbf{x})$$

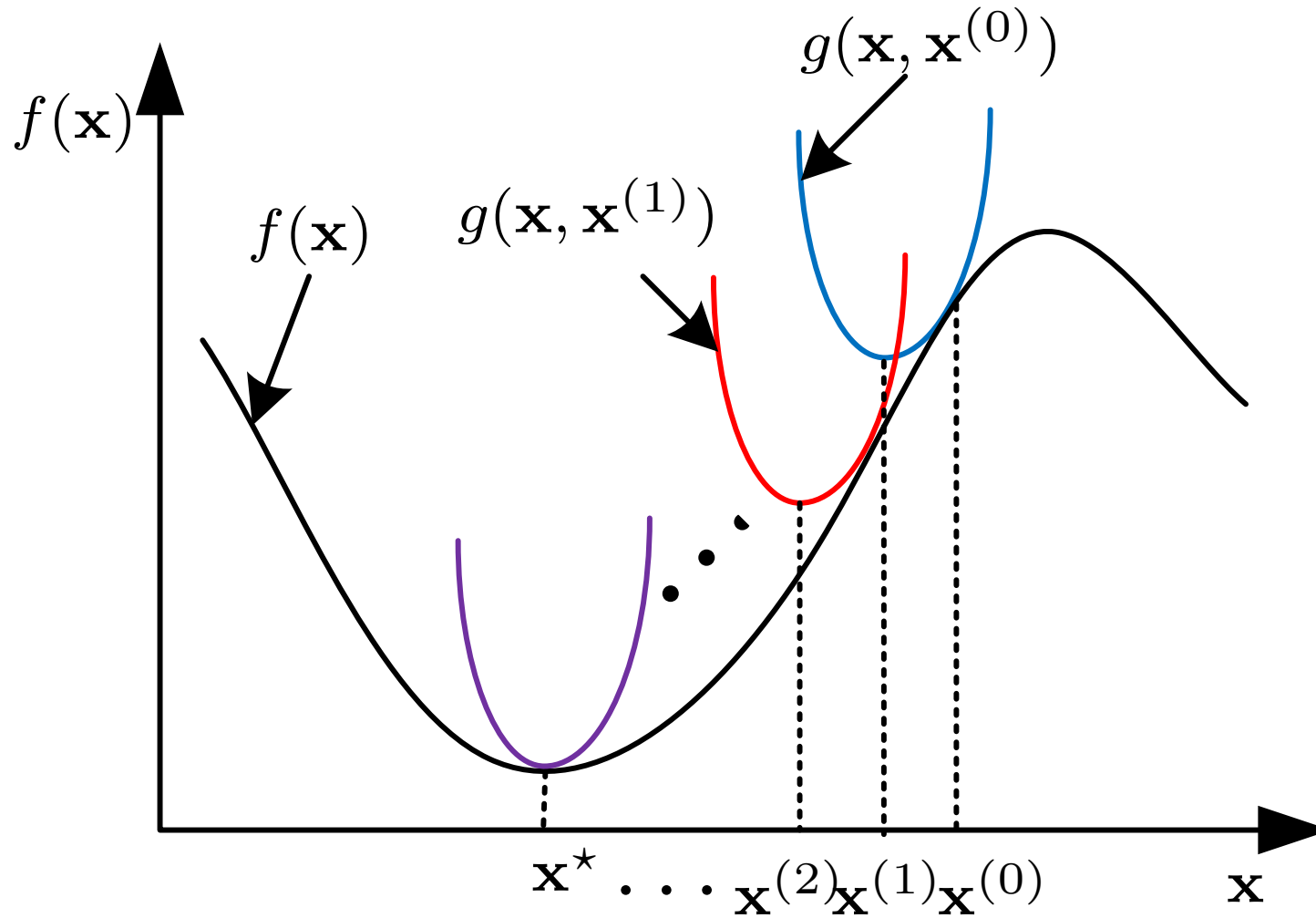
and suppose that  $f$  is hard to minimize directly

- let  $g(\mathbf{x}, \bar{\mathbf{x}})$  be a **surrogate function** that is easy to minimize and satisfies

$$f(\mathbf{x}) \leq g(\mathbf{x}, \bar{\mathbf{x}}) \text{ for all } \mathbf{x}, \bar{\mathbf{x}}, \quad f(\mathbf{x}) = g(\mathbf{x}, \mathbf{x}) \text{ for all } \mathbf{x}$$

- MM algorithm:  $\mathbf{x}^{(k+1)} = \arg \min_{\mathbf{x} \in \mathcal{C}} g(\mathbf{x}, \mathbf{x}^{(k)}), k = 1, 2, \dots$
- as a basic result,  $f(\mathbf{x}^{(0)}) \geq f(\mathbf{x}^{(1)}) \geq f(\mathbf{x}^{(2)}) \dots$
- suppose that  $f$  is convex and  $\mathcal{C}$  is convex. MM is guaranteed to converge to an optimal solution under some mild assumption **[Razaviyayn-Hong-Luo'13]**

## MM for $\ell_2 - \ell_1$ Minimization: General MM Principle



## MM for $\ell_2 - \ell_1$ Minimization

- now consider applying MM to the  $\ell_2 - \ell_1$  minimization problem

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \lambda \|\mathbf{x}\|_1.$$

- let  $c \geq \sigma_{\max}^2(\mathbf{A})$ , and let

$$g(\mathbf{x}, \bar{\mathbf{x}}) = \frac{1}{2} (\|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + c\|\mathbf{x} - \bar{\mathbf{x}}\|_2^2) + \lambda \|\mathbf{x}\|_1$$

– simply plug the same surrogate for  $\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2$  we saw previously

- it can be shown that

$$\mathbf{x}^{(k+1)} = \text{soft} \left( \frac{1}{c} \mathbf{A}^T (\mathbf{y} - \mathbf{A}\mathbf{x}^{(k)}) + \mathbf{x}^{(k)}, \lambda/c \right)$$

where  $\text{soft}$  is called the soft-thresholding operator and is defined as follows: if  $\mathbf{z} = \text{soft}(\mathbf{x}, \delta)$  then  $z_i = \text{sign}(x_i) \max\{|x_i| - \delta, 0\}$

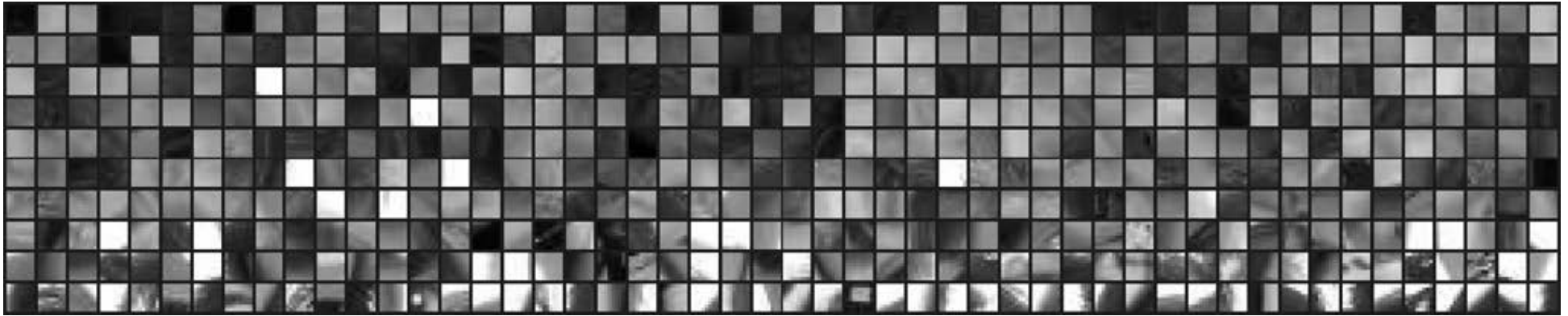
# Dictionary Learning

- previously  $\mathbf{A}$  is assumed to be given
- how about learning a fat  $\mathbf{A}$  from data, as in matrix factorization?
- Dictionary learning (DL): given  $\tau > 0$  and  $\mathbf{Y} \in \mathbb{R}^{m \times n}$ , solve

$$\begin{aligned} \min_{\mathbf{A} \in \mathbb{R}^{m \times k}, \mathbf{B} \in \mathbb{R}^{k \times n}} \quad & \sum_{i=1}^n \|\mathbf{y}_i - \mathbf{A}\mathbf{b}_i\|_2^2 \\ \text{s.t.} \quad & \|\mathbf{b}_i\|_0 \leq \tau, \quad i = 1, \dots, n \end{aligned}$$

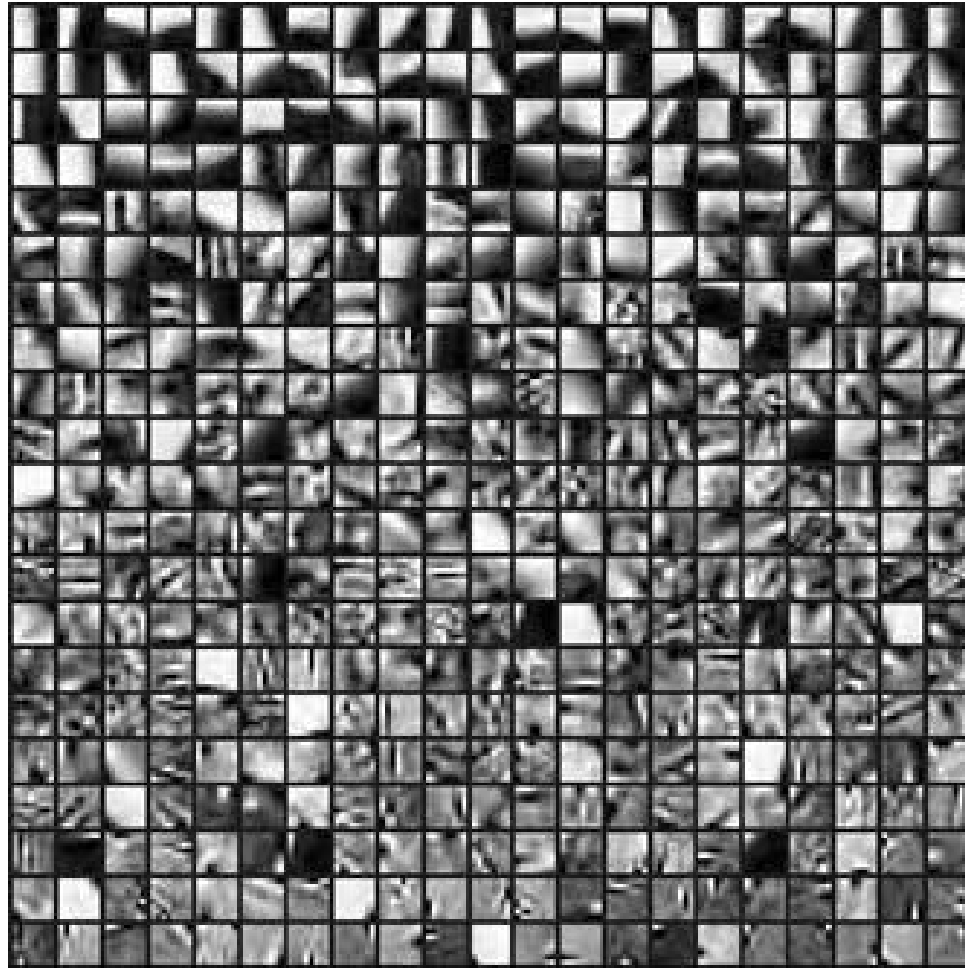
- DL considers  $k \geq m$ , and  $\mathbf{A}$  is called an overcomplete dictionary
- DL is handled by alternating optimization—the same approach in matrix fac.

# Dictionary Learning



A collection of  $n = 500$  random image blocks. Source: [\[Aharon-Elad-Bruckstein'06\]](#).

# Dictionary Learning



The learned dictionary ( $k = 421$ ). Source: [\[Aharon-Elad-Bruckstein'06\]](#).

# Part III: LS with Errors in $A$



## LS with Errors in $\mathbf{A}$

- **Scenario:** errors exist in the system matrix  $\mathbf{A}$
- **Aim:** mitigate the effects of the system matrix errors on the LS solution
- there are many ways to do so, and we look at two
- **Total LS (TLS):**

$$\min_{\mathbf{x} \in \mathbb{R}^n, \mathbf{\Delta}_A \in \mathbb{R}^{m \times n}} \|\mathbf{y} - (\mathbf{A} + \mathbf{\Delta}_A)\mathbf{x}\|_2^2 + \|\mathbf{\Delta}_A\|_F^2$$

- minimally perturb the system matrix by  $\mathbf{\Delta}_A$  for best fitting in the Euclidean sense

- **Robust LS:**

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\mathbf{\Delta}_A \in \mathcal{U}} \|\mathbf{y} - (\mathbf{A} + \mathbf{\Delta}_A)\mathbf{x}\|_2^2$$

for some pre-determined uncertainty set  $\mathcal{U} \subset \mathbb{R}^{m \times n}$

- robustify the LS via a worst-case means

## Total LS

$$\min_{\mathbf{x} \in \mathbb{R}^n, \Delta_{\mathbf{A}} \in \mathbb{R}^{m \times n}} \|\mathbf{y} - (\mathbf{A} + \Delta_{\mathbf{A}})\mathbf{x}\|_2^2 + \|\Delta_{\mathbf{A}}\|_F^2$$

- does not seem to have a closed-form solution at first sight
- turns out to have a closed-form solution under some mild assumptions
- assume  $\mathbf{A}$  to be of full column rank with  $m \geq n + 1$
- let  $\mathbf{C} = [\mathbf{A} \ \mathbf{y}]$ , and let  $\mathbf{v}_{n+1}$  be the  $(n + 1)$ th right singular vector of  $\mathbf{C}$ . If

$$\text{rank}(\mathbf{C}) = n + 1, \quad v_{n+1,n+1} \neq 0,$$

then

$$\mathbf{x}_{\text{TLS}} = -\frac{1}{v_{n+1,n+1}} \begin{bmatrix} v_{1,n+1} \\ \vdots \\ v_{n,n+1} \end{bmatrix}$$

is a TLS solution

- see [\[Golub-Van Loan'13\]](#) for further discussion on issues like  $v_{n+1,n+1} \neq 0$

## Proof Sketch of the TLS Solution

- idea: turn the TLS problem to a low-rank matrix approximation problem
- by a change of variables

$$\mathbf{C} = [ \mathbf{A} \ \mathbf{y} ] \in \mathbb{R}^{m \times (n+1)}, \quad \mathbf{D} = [ \Delta_{\mathbf{A}} \ (\mathbf{A} + \Delta_{\mathbf{A}})\mathbf{x} ] \in \mathbb{R}^{m \times (n+1)},$$

the TLS problem can be formulated as

$$\min_{\mathbf{x}, \mathbf{D}} \|\mathbf{C} - \mathbf{D}\|_F^2 \quad \text{s.t.} \quad \mathbf{D} \begin{bmatrix} \mathbf{x} \\ -1 \end{bmatrix} = \mathbf{0} \quad (\dagger)$$

- the constraint in  $(\dagger)$ , together with  $m \geq n + 1$ , implies  $\text{rank}(\mathbf{D}) \leq n$
- or, we can equivalently rewrite  $(\dagger)$  as

$$\min_{\mathbf{x}, \mathbf{D}} \|\mathbf{C} - \mathbf{D}\|_F^2 \quad \text{s.t.} \quad \text{rank}(\mathbf{D}) \leq n, \quad \mathbf{D} \begin{bmatrix} \mathbf{x} \\ -1 \end{bmatrix} = \mathbf{0}$$

## Proof Sketch of the TLS Solution

- consider a *relaxation* of  $(\dagger)$ :

$$\min_{\mathbf{D}} \|\mathbf{C} - \mathbf{D}\|_F^2 \quad \text{s.t. } \text{rank}(\mathbf{D}) \leq n, \quad (\ddagger)$$

where we drop the constraint  $\mathbf{D} \begin{bmatrix} \mathbf{x} \\ -1 \end{bmatrix} = \mathbf{0}$

- let  $\mathbf{D}^*$  be a solution to  $(\ddagger)$ . If there exists an  $\mathbf{x}^*$  such that  $\mathbf{D}^* \begin{bmatrix} \mathbf{x}^* \\ -1 \end{bmatrix} = \mathbf{0}$ ,  $(\mathbf{D}^*, \mathbf{x}^*)$  is also a solution to  $(\dagger)$  and  $\mathbf{x}^*$  is a TLS solution
- let  $\mathbf{C} = \sum_{i=1}^{n+1} \sigma_i \mathbf{u}_i \mathbf{v}_i^T$  be the SVD
- by the Eckart-Young-Mirsky theorem, a solution to  $(\ddagger)$  is  $\mathbf{D}^* = \sum_{i=1}^n \sigma_i \mathbf{u}_i \mathbf{v}_i^T$ .
- as a basic fact of SVD, we have  $\mathbf{D}^* \mathbf{v}_{n+1} = \mathbf{0}$ .
- thus, if  $v_{n+1,n+1} \neq 0$ , we have the desired TLS solution

## Robust LS

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\Delta_{\mathbf{A}} \in \mathcal{U}} \|\mathbf{y} - (\mathbf{A} + \Delta_{\mathbf{A}})\mathbf{x}\|_2$$

- consider the case of  $\mathcal{U} = \{\Delta_{\mathbf{A}} \in \mathbb{R}^{m \times n} \mid \|\Delta_{\mathbf{A}}\|_2 \leq \lambda\}$  for some  $\lambda > 0$
- the robust LS problem can be shown to be equivalent to

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2 + \lambda \|\mathbf{x}\|_2$$

- Observations and Implications:
  - the equivalent form of the robust LS is very similar to (but not exactly the same as) the previous  $\ell_2$ -regularized LS
  - robustification is equivalent to regularization
- it can be shown that the same equivalence holds if we replace the uncertainty set by  $\mathcal{U} = \{\Delta_{\mathbf{A}} \in \mathbb{R}^{m \times n} \mid \|\Delta_{\mathbf{A}}\|_F \leq \lambda\}$

## Proof Sketch of the Robust LS Equivalence Result

- by the definition of induced norms, we have

$$\|\Delta_{\mathbf{A}}\|_2 \leq \lambda \iff \|\Delta_{\mathbf{A}}\mathbf{x}\|_2 \leq \lambda\|\mathbf{x}\|_2 \quad \text{for all } \mathbf{x} \in \mathbb{R}^n$$

- then, for any  $\mathbf{x} \in \mathbb{R}^n$  and for any  $\Delta_{\mathbf{A}} \in \mathcal{U}$ ,

$$\begin{aligned} \|\mathbf{y} - (\mathbf{A} + \Delta_{\mathbf{A}})\mathbf{x}\|_2 &\leq \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2 + \|\Delta_{\mathbf{A}}\mathbf{x}\|_2 \\ &\leq \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2 + \lambda\|\mathbf{x}\|_2, \end{aligned} \tag{*}$$

and note that the 1st equality above holds if  $\mathbf{y} - \mathbf{A}\mathbf{x} = -\alpha\Delta_{\mathbf{A}}\mathbf{x}$  for some  $\alpha \geq 0$ , and the 2nd equality above holds if  $\mathbf{x}$  is the 1st right singular vector of  $\Delta_{\mathbf{A}}$

- consider the case of  $\mathbf{x} \neq \mathbf{0}$ ,  $\mathbf{y} - \mathbf{A}\mathbf{x} \neq \mathbf{0}$ . It can be verified that

$$\Delta_{\mathbf{A}} = -\frac{\lambda}{\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2\|\mathbf{x}\|_2}(\mathbf{y} - \mathbf{A}\mathbf{x})\mathbf{x}^T$$

attains the equalities in (\*) and lies in  $\mathcal{U}$

- the other cases of  $\mathbf{x}$  are handled in a similar fashion

## More Robust LS Equivalences

- denote  $\mathcal{U}_{q,p} = \{\Delta_{\mathbf{A}} \in \mathbb{R}^{m \times n} \mid \|\Delta_{\mathbf{A}} \mathbf{x}\|_p \leq \lambda \|\mathbf{x}\|_q \ \forall \mathbf{x}\}$ , where  $p, q \geq 1$ . We have

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\Delta_{\mathbf{A}} \in \mathcal{U}_{q,p}} \|\mathbf{y} - (\mathbf{A} + \Delta_{\mathbf{A}}) \mathbf{x}\|_p = \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A} \mathbf{x}\|_p + \lambda \|\mathbf{x}\|_q$$

- proof: almost the same as the previous case
- some interesting special cases:

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\Delta_{\mathbf{A}} \in \mathcal{U}_{2,1}} \|\mathbf{y} - (\mathbf{A} + \Delta_{\mathbf{A}}) \mathbf{x}\|_2 = \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A} \mathbf{x}\|_2 + \lambda \|\mathbf{x}\|_1$$

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\substack{\Delta_{\mathbf{A}} \in \mathbb{R}^{m \times n} \\ \|\delta_i\|_1 \leq \lambda \ \forall i}} \|\mathbf{y} - (\mathbf{A} + \Delta_{\mathbf{A}}) \mathbf{x}\|_1 = \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A} \mathbf{x}\|_1 + \lambda \|\mathbf{x}\|_1$$

note  $\mathcal{U}_{1,1} = \{\Delta_{\mathbf{A}} \in \mathbb{R}^{m \times n} \mid \|\delta_i\|_1 \leq \lambda \ \forall i\}$

- Implication:**  $\ell_1$  regularization may also be seen as an act of robustification
- suggested reading: [\[Bertsimas-Copenhaver'17\]](#), including extension to PCA

## More on LS

cf. Chapter 6 in **[Golub-Van Loan'13]**



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