Quantum Circuit Implementing Grover's Algorithm

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Introduction

This is a tutorial on the use of Quantum Computing Mathematica add-on to implement Grover's Algorithm using Quantum Gates.

Grover's algorithm is a quantum algorithm for searching an unsorted database with N entries in $O(N^{1/2})$ time and using $O(\log N)$ storage space (Grover L.K.: A fast quantum mechanical algorithm for database search, Proceedings, 28th Annual ACM Symposium on the Theory of Computing, (May 1996) p. 212).

Grover's algorithm can be generalized to "Quantum Counting" (Brassard G., P. Hoyer and Al. Tapp, arXiv:quant-ph/9805082) and "Quantum Amplitude Amplification and Estimation" (Brassard G., P. Hoyer, M. Mosca and A. Tapp, arXiv:quant-ph/0005055).

A Mathematica "Demonstration" implementing Grover's Algorithm can be found in this link:

Alexander Prokopenya, "Quantum Circuit Implementing Grover's Search Algorithm" from The Wolfram Demonstrations Project

http://demonstrations.wolfram.com/QuantumCircuitImplementingGroversSearchAlgorithm/

Load the Package

First load the Quantum' Computing' package. Write:

Needs["Quantum'Computing'"];

then press at the same time the keys SHFT-ENTER to evaluate. Mathematica will load the package:

Needs["Quantum`Computing`"]

```
Quantum`Computing` Version 2.2.0. (July 2010)

A Mathematica package for Quantum Computing
  in Dirac bra-ket notation and plotting of quantum circuits
by José Luis Gómez-Muñoz

Execute SetComputingAliases[] in order to use
  the keyboard to enter quantum objects in Dirac's notation

SetComputingAliases[] must be executed again in each new notebook that is created
```

In order to use the keyboard to enter quantum objects write:

SetComputingAliases[];

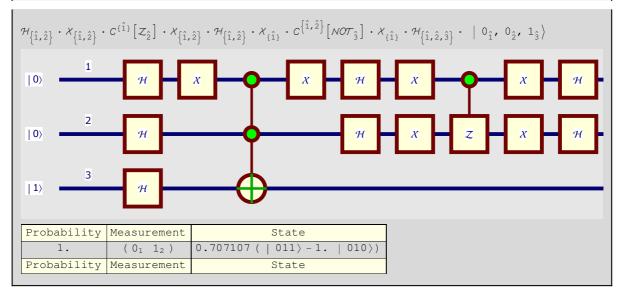
then press at the same time the keys shell-enter to evaluate. The semicolon prevents *Mathematica* from printing the help message. Remember that SetComputingAliases[] must be evaluated again in each new notebook:

SetComputingAliases[];

Circuit implementation of Gover's Algorithm for two qubits

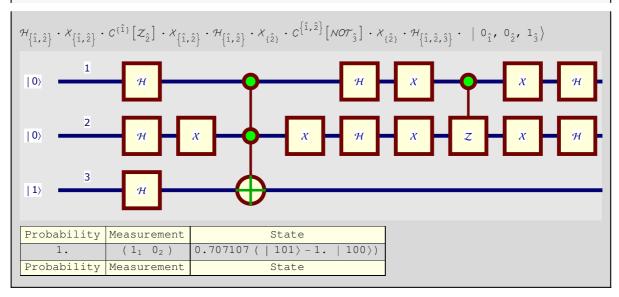
Following commands show and calculate the quantum circuit implementation of Grover's Algorithm for selecting the ket that represents the number k=1 as a two-qbit binary number (0_1 1_2). Please continue reading below for a detailed explanation.

```
n = 2;
k = 1;
it = 1;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
\texttt{circuit} = \left(\mathcal{H}_{n} \cdot \mathcal{X}_{n} \cdot C^{n-1}\left[\mathcal{Z}_{\hat{n}}\right] \cdot \mathcal{X}_{n} \cdot \mathcal{H}_{n} \cdot \mathcal{X}_{\texttt{mylist}} \cdot C^{n}\left[\mathcal{NOT}_{\hat{n+1}}\right] \cdot \mathcal{X}_{\texttt{mylist}}\right)^{\texttt{it}} \cdot
       \mathcal{H}_{n+1} \, \cdot \, \left( \, \, \left| \, \, \mathbf{0}_{\hat{1}} \right\rangle \right)^{\otimes n} \, \cdot \, \, \left| \, \, \mathbf{1}_{\hat{n+1}} \right\rangle;
Column[{
     circuit,
     QuantumPlot[
       circuit, ImageSize \rightarrow {600, Automatic}],
     TraditionalForm[
       N[QuantumEvaluate[QubitMeasurement[circuit, n, FactorKet → False]]]]
       }]
```



Following commands show and calculate the quantum circuit implementation of Grover's Algorithm for selecting the ket that represents the number k=2 as a binary number (1_1 0_2). See the difference with the previous circuit. The working of this circuit is explained below in this document.

```
n = 2;
k = 2;
it = 1;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
\texttt{circuit} = \left(\mathcal{H}_n \, \cdot \, \mathcal{X}_n \, \cdot \, C^{n-1} \left[ \, \mathcal{Z}_{\hat{n}} \, \right] \, \cdot \, \mathcal{X}_n \, \cdot \, \mathcal{H}_n \, \cdot \, \mathcal{X}_{\texttt{mylist}} \, \cdot \, C^n \left[ \, \mathcal{NOT}_{\hat{n+1}}^{\, \hat{}} \, \right] \, \cdot \, \mathcal{X}_{\texttt{mylist}} \right)^{\texttt{it}} \, \cdot \,
       \mathcal{H}_{n+1} \cdot \left( \left| 0_{\hat{1}} \right\rangle \right)^{\otimes n} \cdot \left| 1_{\hat{n+1}} \right\rangle;
Column[{
     circuit,
     QuantumPlot[
       circuit, ImageSize → {600, Automatic}],
     TraditionalForm[
       N[QuantumEvaluate[QubitMeasurement[circuit, n, FactorKet → False]]]]
```



Step-by-step explanation of the circuit implementing Grover's Algorithm for two qubits

Initial State

An equally weighted superposition of all basis-states is generated using three Hadamard gates.

Notice that at this stage all the kets having a positive sign have the ancillary qubit set to $0_{\hat{3}}$, and all the kets having a negative sign have the ancillary qubit set to 13

The state of the system is stored in $|g1\rangle$

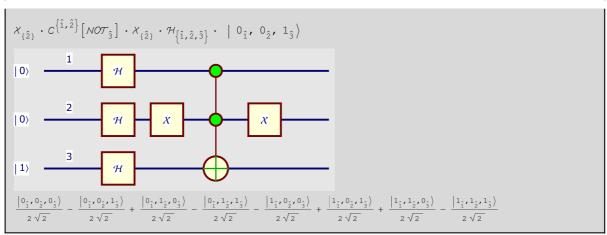
```
n = 2;
circuitpart = \mathcal{H}_{n+1} \cdot \left( \left| 0_{\hat{1}} \right\rangle \right)^{\otimes^n} \cdot \left| 1_{\hat{n+1}} \right\rangle;
Column[{
    circuitpart,
    QuantumPlot[circuitpart],
     | g1 > = QuantumEvaluate[circuitpart] } ]
```

Changing the phase of those kets that meet the search criteria

The search criteria in this example is to represent the number 2 in binary digits: $(1_{\hat{1}}, 0_{\hat{2}})$, those kets that meet the criteria get their phase changed (the signs plus and minus are exchanged). The conditional negation of the ancillary qubit gives this change of sign.

The state of the system is stored in $|g2\rangle$

```
\begin{split} &n=2;\\ &k=2;\\ &mylist=Flatten[Position[IntegerDigits[k,\,2,\,n],\,0]];\\ &circuitpart=\mathcal{X}_{mylist}\cdot\mathcal{C}^n\left[\mathcal{NOT}_{\hat{n+1}}\right]\cdot\mathcal{X}_{mylist}\cdot\mathcal{H}_{n+1}\cdot\left(\;\big|\;0_{\hat{1}}\rangle\right)^{\otimes^n}\cdot\;\big|\;1_{\hat{n+1}}\rangle;\\ &column[\{&circuitpart,\\ &QuantumPlot[circuitpart],\\ &\big|\;g2\rangle=QuantumEvaluate[circuitpart]\}] \end{split}
```



It is very important to understand the **difference** between the states just after the Hadamard gates ($|g1\rangle$), and the state after the conditional change of sign (phase) for those kets that meet the criteria ($|g2\rangle$):

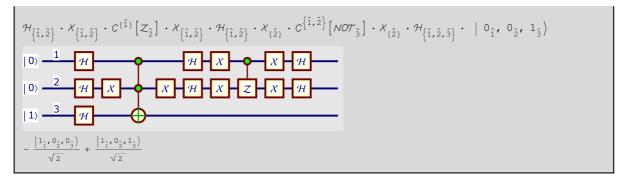
```
TraditionalForm[
 Column[{"The sign was changed (+\leftrightarrow -) only on those kets that meet the
       search criteria", |g1\rangle, |g2\rangle}, Dividers \rightarrow All]]
 The sign was changed (+\leftrightarrow -) only on those kets that meet the search criteria
```

```
+ |110>
                                          \frac{|011\rangle}{2\sqrt{2}} + \frac{|100\rangle}{2\sqrt{2}} - \frac{|101\rangle}{2\sqrt{2}}
                            |010>
              |001>
2\sqrt{2}
             2\sqrt{2}
                                       2\sqrt{2}
                                                                                                    2\sqrt{2}
                            2\sqrt{2}
                       + |010>
                                                                  + |101>
                                                                                 + |110>
          |001>
                                       |011>
                                                     _ |100>
                                                                                                 |111>
|000>
2\sqrt{2}
             2\sqrt{2}
                            2\sqrt{2}
                                          2\sqrt{2}
                                                        2\sqrt{2}
                                                                       2\sqrt{2}
                                                                                     2\sqrt{2}
                                                                                                    2\sqrt{2}
```

Increasing the amplitude of those kets that had their phase changed

The rest of the circuit increases the amplitude of those kets that had their phase changed. In the case of two qubits, the other kets get their amplitude reduced to zero.

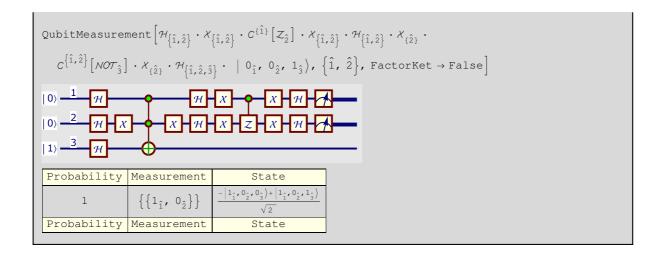
```
n = 2;
k = 2;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
circuitpart =
     \mathcal{H}_{n} \cdot \mathcal{X}_{n} \cdot C^{n-1} \left[ \mathcal{Z}_{\hat{n}} \right] \cdot \mathcal{X}_{n} \cdot \mathcal{H}_{n} \cdot \mathcal{X}_{mylist} \cdot C^{n} \left[ NOT_{\hat{n},1} \right] \cdot \mathcal{X}_{mylist} \cdot \mathcal{H}_{n+1} \cdot \left( \begin{array}{c} 0_{\hat{1}} \end{array} \right) \right) \otimes^{n} \cdot \left[ 1_{\hat{n},1} \right);
Column[{
     circuitpart,
     QuantumPlot[circuitpart],
     QuantumEvaluate[circuitpart]}]
```



Measurement gives with those kets that had their amplitude increased

In the case of **two qubits**, measurment will give those values that satisfy the search criteria with a 100% of probability

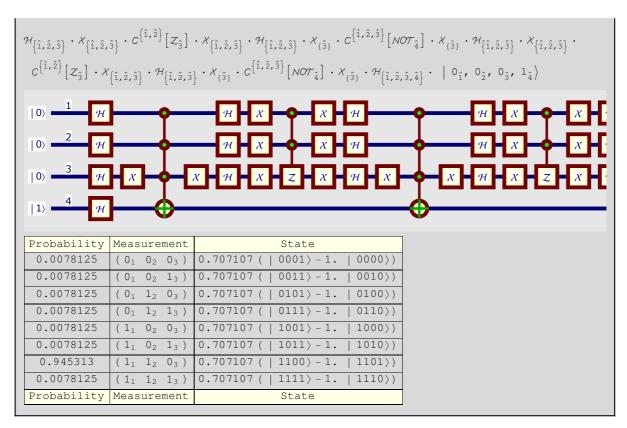
```
n = 2;
  k = 2;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
\texttt{circuitpart} = \texttt{QubitMeasurement} \left[ \mathcal{H}_n \, \cdot \, \mathcal{X}_n \, \cdot \, \mathcal{C}^{n\text{--}1} \left[ \, \mathcal{Z}_{\hat{n}} \, \right] \, \cdot \, \mathcal{X}_n \, \cdot \, \mathcal{H}_n \, \cdot \, \mathcal{X}_{\texttt{mylist}} \, \cdot \, \, \mathcal{X}_{\texttt{mylist}} \, \cdot 
                                                                  C^{n}[NOT_{n+1}] \cdot X_{mylist} \cdot \mathcal{H}_{n+1} \cdot (\mid 0_{\hat{1}}))^{\otimes n} \cdot \mid 1_{n+1} \rangle, n, FactorKet \rightarrow False;
Column[{
                                 circuitpart,
                                   QuantumPlot[circuitpart],
                                   QuantumEvaluate[circuitpart]}]
```



Circuit implementation of Gover's Algorithm for three qubits

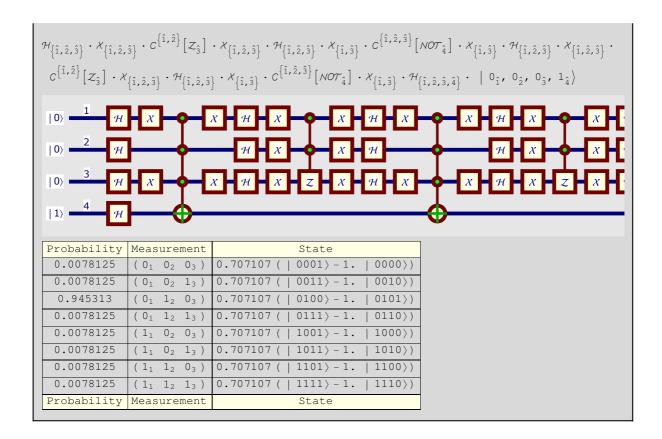
Following commands show and calculate the quantum circuit implementation of Grover's Algorithm for selecting the ket that represents the number k=6 as a three-qubit binary number (1_1 1_2 0_3). Please continue reading below for a detailed explanation. The use of the standard Mathematica command Expand[] makes the calculation faster because it expands the power of the operators:

```
n = 3;
k = 6;
 it = 2;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
\texttt{circuit} = \texttt{Expand} \left[ \left( \mathcal{H}_{n} \cdot \mathcal{X}_{n} \cdot \mathcal{C}^{n-1} \left[ \mathcal{Z}_{\hat{n}} \right] \cdot \mathcal{X}_{n} \cdot \mathcal{H}_{n} \cdot \mathcal{X}_{\texttt{mylist}} \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n+1}} \right] \cdot \mathcal{X}_{\texttt{mylist}} \right)^{\texttt{it}} \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{X}_{\texttt{mylist}} \right] \cdot \mathcal{X}_{\texttt{mylist}} \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{N}_{\texttt{mylist}} \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{N}_{\texttt{mylist}} \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{N}_{\texttt{mylist}} \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{C}^{n} \left[ \mathcal{NOT}_{\hat{n}} \right] \cdot \mathcal{C}^{n} \left[ 
                                                           \mathcal{H}_{n+1} \cdot \left( \left| 0_{\hat{1}} \right\rangle \right)^{\otimes n} \cdot \left| 1_{\hat{n+1}} \right\rangle \right];
Column[{
                              circuit,
                              QuantumPlot[
                                             circuit, ImageSize → {600, Automatic}],
                              TraditionalForm[
                                             {\tt N[QuantumEvaluate[QubitMeasurement[circuit, n, FactorKet $\rightarrow$ False]]]]}
                                             }]
```



Following commands show and calculate the quantum circuit implementation of Grover's Algorithm for selecting the ket that represents the number k=2 as a three-qubit binary number (0_1 1_2 0_3). Compare it with the previous circuit. This circuit is explained below in this document. The use of the standard Mathematica command Expand[] makes the calculation faster because it expands the power of the operators:

```
n = 3;
 k = 2;
 it = 2;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
\texttt{circuit} = \texttt{Expand} \left[ \left( \mathcal{H}_{n} \cdot \mathcal{X}_{n} \cdot \mathit{C}^{n-1} \left[ \mathcal{Z}_{\hat{n}} \right] \cdot \mathcal{X}_{n} \cdot \mathcal{H}_{n} \cdot \mathcal{X}_{\texttt{mylist}} \cdot \mathit{C}^{n} \left[ \mathit{NOT}_{\hat{n+1}} \right] \cdot \mathcal{X}_{\texttt{mylist}} \right)^{\texttt{it}} \cdot \mathcal{X}_{\texttt{mylist}} \right] \cdot \mathcal{X}_{\texttt{mylist}} \cdot \mathcal{X}_{\texttt{myli
                                                          \mathcal{H}_{n+1} \cdot \left( \left| 0_{\hat{1}} \right\rangle \right)^{\otimes n} \cdot \left| 1_{\hat{n+1}} \right\rangle \right];
 Column[{
                              circuit,
                              QuantumPlot[
                                            circuit, ImageSize → {600, Automatic}],
                              TraditionalForm[
                                            N[QuantumEvaluate[QubitMeasurement[circuit, n, FactorKet → False]]]]
                                            }]
```



Step-by-step explanation of the circuit implementing Grover's Algorithm for three qubits

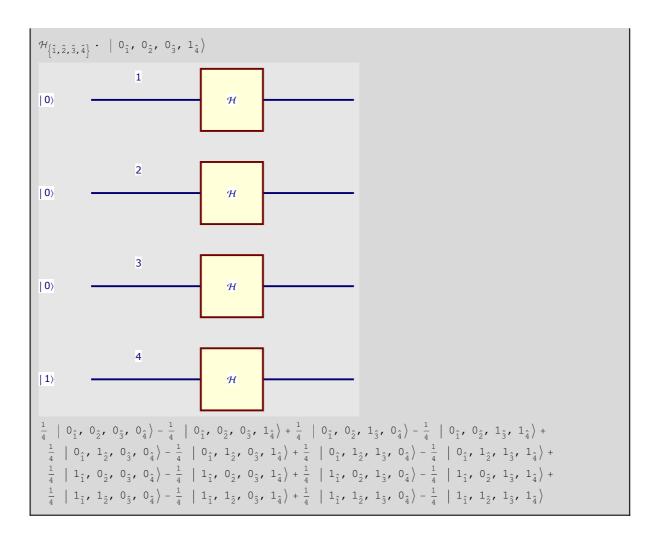
Initial State

An equally weighted superposition of all basis-states is generated using four Hadamard gates.

Notice that at this stage all the kets having a positive sign have the ancillary qubit set to $0_{\hat{a}}$, and all the kets having a negative sign have the ancillary qubit set to $\mathbf{1}_{\hat{\mathbf{A}}}$

The state of the system is stored in $| h1 \rangle$

```
n = 3;
circuitpart = \mathcal{H}_{n+1} \cdot ( \mid 0_{\hat{1}} \rangle)^{\otimes^n} \cdot \mid 1_{\hat{n+1}} \rangle;
Column[{
   circuitpart,
   QuantumPlot[circuitpart],
     | h1 > = QuantumEvaluate[circuitpart] }]
```

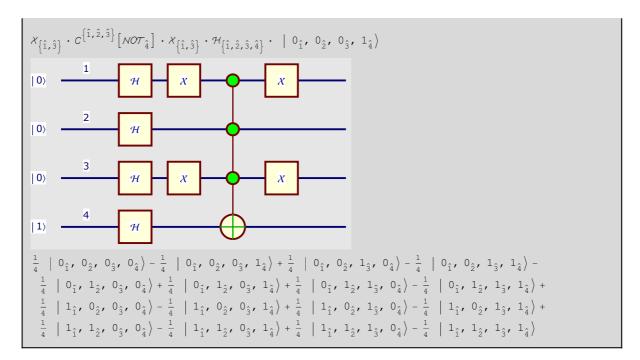


Changing the phase of those kets that meet the search criteria

The search criteria in this example is to represent the number 2 in three binary digits: $(0_1 \quad 1_2 \quad 0_3)$, those kets that meet the criteria get their phase changed (the signs plus and minus are exchanged). The conditional negation of the ancillary qubit gives this change of sign.

The state of the system is stored in $| h2 \rangle$

```
n = 3;
k = 2;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
\texttt{circuitpart} = \textit{X}_{\texttt{mylist}} \, \cdot \, \textit{C}^{n} \left[ \textit{NOT}_{\hat{n+1}} \right] \, \cdot \, \textit{X}_{\texttt{mylist}} \, \cdot \, \mathcal{H}_{n+1} \, \cdot \, \left( \, \, \big| \, \, \mathbf{0}_{\hat{1}} \right) \right) \, e^{n} \, \cdot \, \, \, \left| \, \, \mathbf{1}_{\hat{n+1}} \right\rangle;
Column[{
     circuitpart,
     QuantumPlot[circuitpart],
       | h2 > = QuantumEvaluate[circuitpart]}]
```



It is very important to understand the **difference** between the states just after the Hadamard gates (| h1 >), and the state after the conditional change of sign (phase) for those kets that meet the criteria (| h2)):

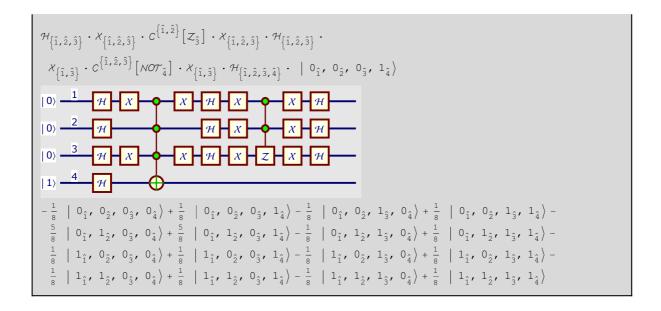
```
TraditionalForm[
                 Column[{"The sign was changed (+\leftrightarrow -) only on those kets that meet the
                                                                                 search criteria", | h1 \rangle, | h2 \rangle}, Dividers \rightarrow All]]
          The sign was changed (+\leftrightarrow -) only on those kets that meet the search criteria
            \frac{1}{4} \mid 0000 \rangle - \frac{1}{4} \mid 0001 \rangle + \frac{1}{4} \mid 0010 \rangle - \frac{1}{4} \mid 0011 \rangle + \frac{1}{4} \mid 0100 \rangle - \frac{1}{4} \mid 0101 \rangle + \frac{1}{4} \mid 0110 \rangle - \frac{1}{4} \mid 0111 \rangle + \frac{1}{4} \mid 0110 \rangle - \frac{1}{4} \mid 0111 \rangle + \frac{1}{4} \mid 0
```

$\frac{1}{4} \mid 1000\rangle - \frac{1}{4} \mid 1001\rangle + \frac{1}{4} \mid 1010\rangle - \frac{1}{4} \mid 1011\rangle + \frac{1}{4} \mid 1100\rangle - \frac{1}{4} \mid 1101\rangle + \frac{1}{4} \mid 1100\rangle - \frac{1}{4} \mid 1111\rangle + \frac{1}{4} \mid 1100\rangle - \frac{1}{4} \mid 1110\rangle - \frac{1}{4} \mid 1111\rangle + \frac{1}{4} \mid 0000\rangle - \frac{1}{4} \mid 0001\rangle + \frac{1}{4} \mid 0010\rangle - \frac{1}{4} \mid 0011\rangle - \frac{1}{4} \mid 0100\rangle + \frac{1}{4} \mid 0101\rangle + \frac{1}{4} \mid 0110\rangle - \frac{1}{4} \mid 0111\rangle + \frac{1}{4} \mid 0101\rangle + \frac{1}{4} \mid 0110\rangle - \frac{1}{4} \mid 0111\rangle + \frac{1}{4} \mid 0110\rangle - \frac{1}{4} \mid 0111\rangle + \frac{1}{4} \mid 0110\rangle - \frac{1}{4} \mid 0111\rangle + \frac{1}{4} \mid 0111\rangle + \frac{1}{4} \mid 0110\rangle - \frac{1}{4} \mid 0111\rangle + \frac{1}{4} \mid$ $\frac{1}{4} \mid 1000 \rangle - \frac{1}{4} \mid 1001 \rangle + \frac{1}{4} \mid 1010 \rangle - \frac{1}{4} \mid 1011 \rangle + \frac{1}{4} \mid 1100 \rangle - \frac{1}{4} \mid 1101 \rangle + \frac{1}{4} \mid 1110 \rangle - \frac{1}{4} \mid 1111 \rangle$

Increasing the amplitude of those kets that had their phase changed

The rest of the circuit increases the amplitude of those kets that had their phase changed. The state of the system is stored in | h3 >

```
n = 3;
k = 2;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
     \mathcal{H}_{n} \cdot \mathcal{X}_{n} \cdot \mathcal{C}^{n-1}\left[\mathcal{Z}_{\hat{n}}\right] \cdot \mathcal{X}_{n} \cdot \mathcal{H}_{n} \cdot \mathcal{X}_{mylist} \cdot \mathcal{C}^{n}\left[\mathcal{NOT}_{\hat{n+1}}\right] \cdot \mathcal{X}_{mylist} \cdot \mathcal{H}_{n+1} \cdot \left(\begin{array}{c|c} \mathbf{0}_{\hat{1}} \end{array}\right)\right) \otimes^{n} \cdot \left[\begin{array}{c|c} \mathbf{1}_{\hat{n+1}} \end{array}\right];
Column[{
      circuitpart,
     QuantumPlot[circuitpart],
        | h3 > = QuantumEvaluate[circuitpart] }]
```



Measurement gives with those kets that had their amplitude increased

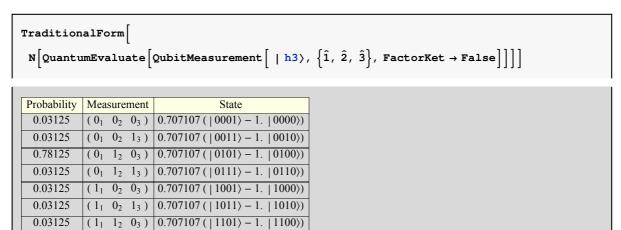
0.03125

Probability

 $(1_1 \ 1_2 \ 1_3)$

Measurement

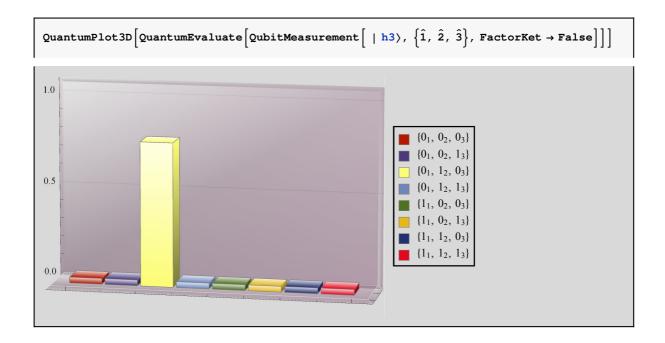
The result of a measurement on the state | h3 has a high probability of giving the state that meets the search criteria:



The result of a measurement on the state | h3 has a high probability of giving the state that meets the search criteria:

 $0.707107 (| 11111 \rangle - 1. | 1110 \rangle)$

State



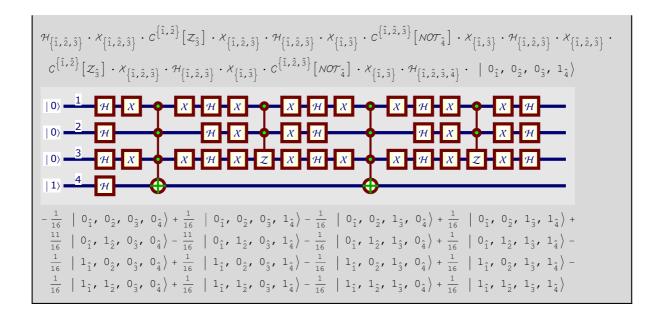
A second iteration gives the optimal increase of probability for three qubits

A second iteration further increments the amplitude of the kets that meets the criteria.

The optimal number of iterations is given by IntegerPart $\left[\pi / \left(4 \operatorname{ArcSin}\left(\frac{1}{\sqrt{2^n}}\right)\right)\right]$. For three qubits, this optimal is two; further iterations would decrease (instead of increase) the desired amplitude.

The state of the system is stored in $|h4\rangle$

```
n = 3;
k = 2;
mylist = Flatten[Position[IntegerDigits[k, 2, n], 0]];
circuitpart = Expand
  Column[{
 circuitpart,
 QuantumPlot[circuitpart, ImageSize → 500],
```



Measurement gives with those kets that had their amplitude increased

The result of a measurement on the state $|h4\rangle$ has the highest probability of giving the state that meets the search criteria:

$$\begin{split} &\text{TraditionalForm} \Big[\\ &\text{N} \Big[\text{QuantumEvaluate} \Big[\text{QubitMeasurement} \Big[\mid \text{h4} \rangle, \; \Big\{ \hat{1}, \; \hat{2}, \; \hat{3} \Big\}, \; \text{FactorKet} \to \text{False} \Big] \Big] \Big] \Big] \end{aligned}$$

Probability	Measurement	State
0.0078125	$(0_1 \ 0_2 \ 0_3)$	0.707107 (0001 > - 1. 0000 >)
0.0078125	$(0_1 \ 0_2 \ 1_3)$	0.707107 (0011 > - 1. 0010 >)
0.945313	$(0_1 \ 1_2 \ 0_3)$	0.707107 (0100> - 1. 0101>)
		0.707107 (0111) – 1. 0110))
0.0078125	$(1_1 \ 0_2 \ 0_3)$	0.707107 (1001 \rangle - 1. 1000 \rangle)
		0.707107 (1011) – 1. 1010))
		0.707107 (1101 > - 1. 1100 >)
	$(1_1 \ 1_2 \ 1_3)$	0.707107 (1111) – 1. 1110))
Probability	Measurement	State

A Mathematica "Demonstration" implementing Grover's Algorithm can be found in this link:

Alexander Prokopenya, "Quantum Circuit Implementing Grover's Search Algorithm" from The Wolfram Demonstrations

http://demonstrations.wolfram.com/QuantumCircuitImplementingGroversSearchAlgorithm/

by José Luis Gómez-Muñoz http://homepage.cem.itesm.mx/lgomez/quantum/ jose.luis.gomez@itesm.mx