

## CSC 320 - Tutorial

1. Deterministic Finite Automaton
2. Non-deterministic Finite Automaton

### Deterministic Finite Automaton (DFA)

- Is expressed as a **5-tuple**  $(Q, \Sigma, \delta, q_0, F)$ 
  - $Q$ : finite set of states
  - $\Sigma$ : alphabet finite set
  - $\delta$ : transition function  $(Q \times \Sigma) \rightarrow Q$
  - $q_0$ : start state  $q_0 \in Q$
  - $F$ : sets of accept/final states  $F \subseteq Q$
- The language  $L$  of a deterministic finite automata  $M$ ,  $L(M)$  is exactly the set of all strings that  $M$  accepts  $\therefore M$  recognizes  $L(M)$
- A given language  $L$  is **regular** iff it is recognized by some deterministic finite automaton

### Non-deterministic Finite Automata (NFA)

- Is expressed as a **5-tuple**  $(Q, \Sigma, \delta, q_0, F)$ 
  - a.  $Q$ : finite set of states
  - b.  $\Sigma$ : alphabet
  - c.  $\delta$ : transition function  $Q \times (\Sigma \cup \{\epsilon\}) \rightarrow P(Q)$
  - d.  $q_0$ : start state  $q_0 \in Q$
  - e.  $F$ : sets of accept/final states  $F \subseteq Q$
- The language  $L$  of a non-deterministic finite automata  $N$ ,  $L(N)$  is exactly the set of all strings that  $N$  accepts
- For every DFA  $M$  there exists an equivalent NFA  $N$  (ie.  $L(M) = L(N)$ )
- For every NFA  $N$  there exists an equivalent DFA  $M$

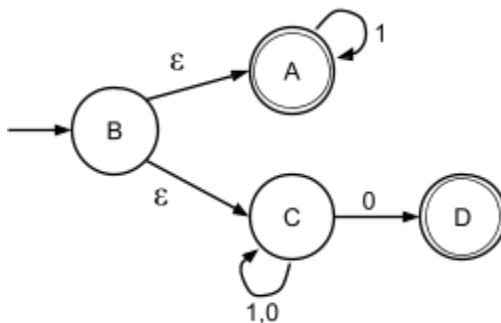
## Questions

1. Give the formal specification of a DFA for the following languages:

a.  $L_1 = \{0\}^*$  over  $\Sigma = \{0\}$

b.  $L_2 = \{w \in \{a, b\}^* \mid w \text{ is a string NOT in } L((ab^+)^*)\}$

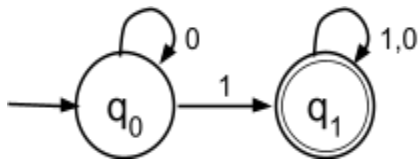
2. Consider the state diagram below:



a. Is this state machine a DFA or an NFA? How can you tell?

- b. Is the string 0011 accepted by this state machine? What about 1100?
- c. What is the language of this machine?
3. Given an example of how regular languages  $L_1$  and  $L_2$  are closed under intersection using the DFAs  $M_1$  and  $M_2$  below (proof by construction) where  $L_1 = L(M_1)$  and  $L_2 = L(M_2)$

$M_1$



$M_2$



