



Chapel: Productive Parallel Programming at Scale

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What You Will Learn

- Concepts of distributed parallel programming models
 - MPI, SPMD
 - CUDA, SIMT
 - PGAS
- What is the Chapel programming language?
 - Learn enough to get started writing programs
- How does the Chapel language fit in to the landscape of parallel programming tools?



About Me



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High Performance Computing (HPC) Programming Models by Example



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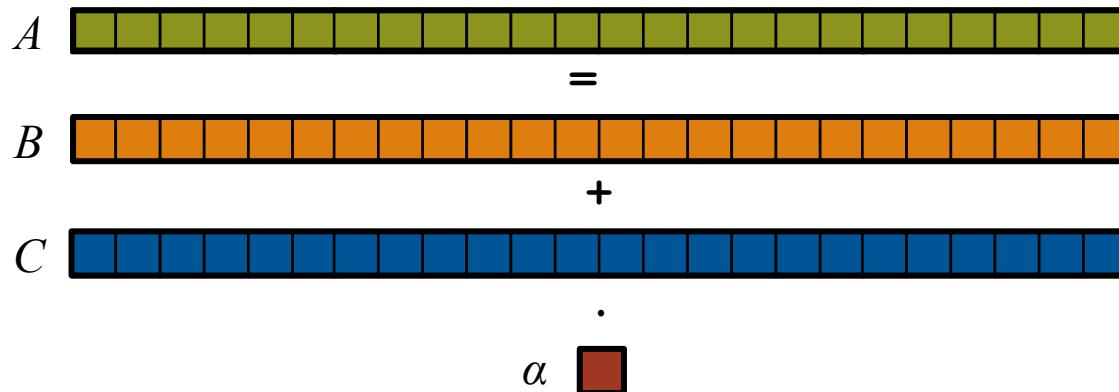
| ANALYZE

STREAM Triad: a trivial parallel computation

Given: m -element vectors A, B, C

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures:

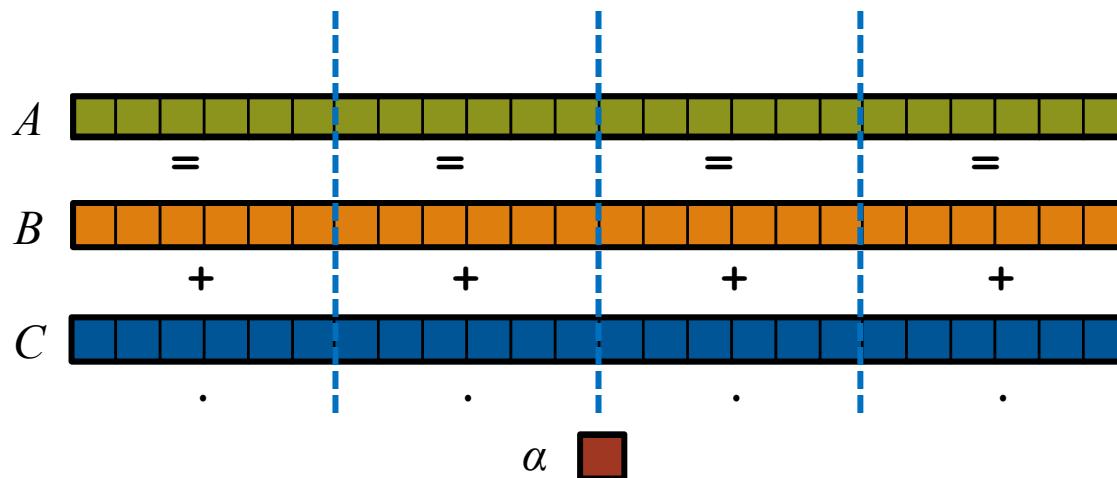


STREAM Triad: a trivial parallel computation

Given: m -element vectors A, B, C

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel:

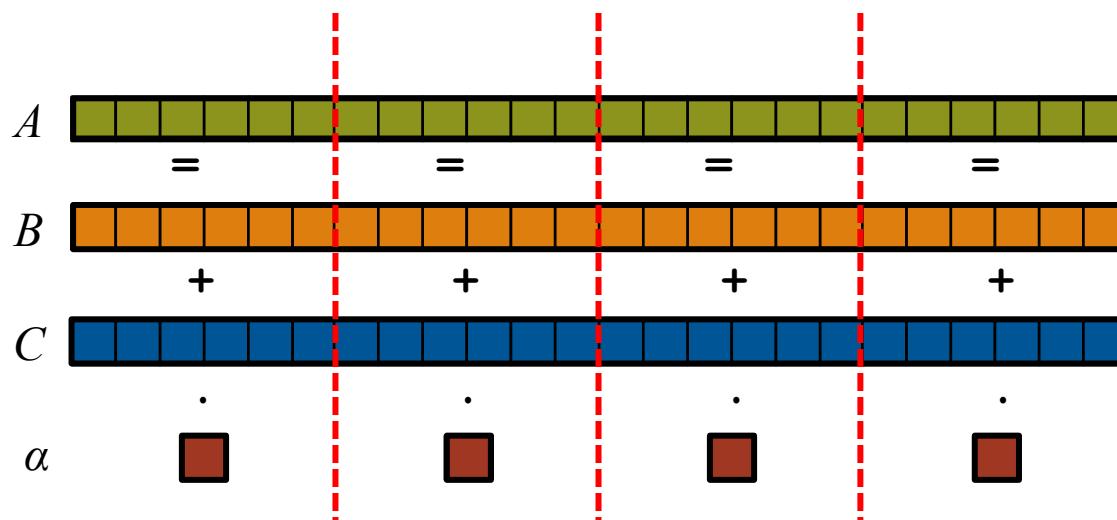


STREAM Triad: a trivial parallel computation

Given: m -element vectors A, B, C

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory):

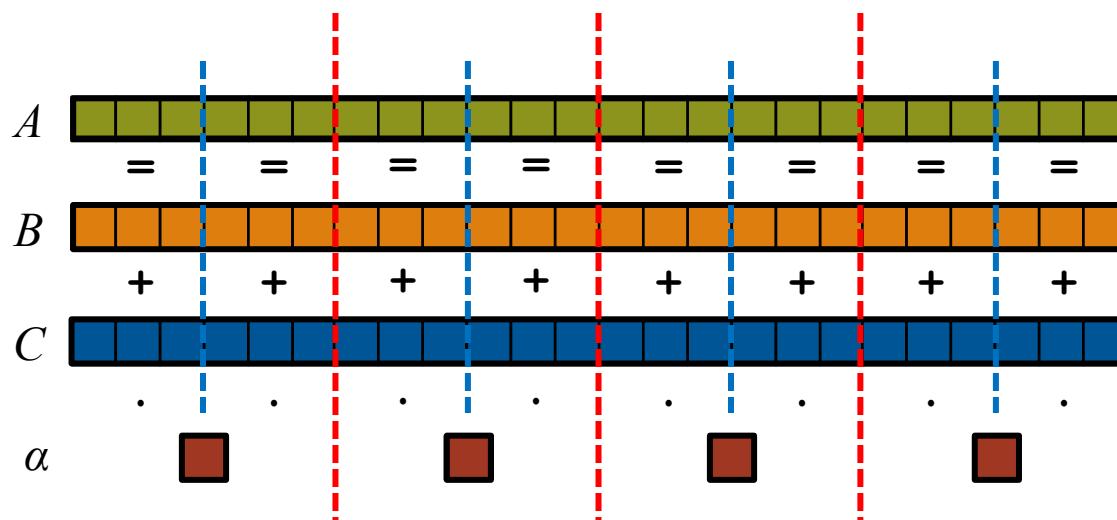


STREAM Triad: a trivial parallel computation

Given: m -element vectors A, B, C

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory multicore):



STREAM Triad: MPI

MPI

```
#include <hpcc.h>

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Parms *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
                0, comm );

    return errCount;
}

int HPCC_Stream(HPCC_Parms *params, int doIO) {
    register int j;
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3,
                                       sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );
```

```
if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to allocate memory
(%d).\n", VectorSize );
        fclose( outFile );
    }
    return 1;
}

for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 0.0;
}

scalar = 3.0;

for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);
```

STREAM Triad: MPI

```
#include <hpcc.h>

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
                0, comm );
}

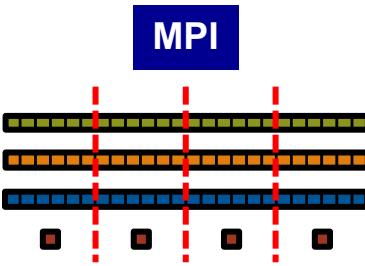
return errCount;
}

int HPCC_Stream(HPCC_Params *params, int doIO) {
    register int j;
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3,
                                       sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );
}

```



"SPMD":
Single Program
Multiple Data

```
if (!a)
    if (c)
        if (b)
            if (a)
                if (doIO) {
                    fprintf( outFile, "Failed to allocate memory
(%d).\n", VectorSize );
                    fclose( outFile );
                }
                return 1;
            }

for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 0.0;
}

scalar = 3.0;

for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);
}

```

STREAM Triad: MPI+OpenMP

```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Parms *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
                0, comm );

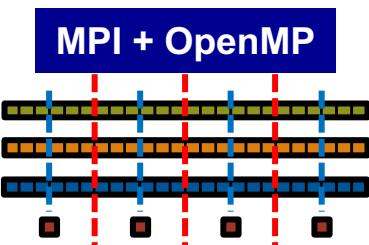
    return errCount;
}

int HPCC_Stream(HPCC_Parms *params, int doIO) {
    register int j;
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3,
                                       sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );
}

```



```
if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to allocate memory
(%d).\n", VectorSize );
        fclose( outFile );
    }
    return 1;
}

#ifndef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 0.0;
}

scalar = 3.0;

#ifndef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);

```

STREAM Triad: MPI+OpenMP



```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Parms *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

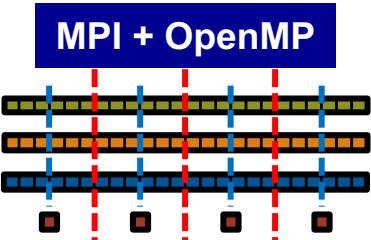
    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
                0, comm );

    return errCount;
}

int HPCC_Stream(HPCC_Parms *params, int doIO) {
    register int j;
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3,
                                       sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );
```



"Hybrid"

multiple models of
parallelism

```
if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to allocate memory
(%d).\n", VectorSize );
        fclose( outFile );
    }
    return 1;
}
```

```
#ifdef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 0.0;
}
```

scalar = 3.0;

```
#ifdef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);
```

shared memory
parallelism



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STREAM Triad: MPI+OpenMP vs. CUDA

```
#define N      2000000
int main() {
    float *d_a, *d_b, *d_c;
    float scalar;

    cudaMalloc((void**)&d_a, sizeof(float)*N);
    cudaMalloc((void**)&d_b, sizeof(float)*N);
    cudaMalloc((void**)&d_c, sizeof(float)*N);

    dim3 dimBlock(128);
    dim3 dimGrid(N/dimBlock.x );
    if( N % dimBlock.x != 0 ) dimGrid

    set_array<<<dimGrid, dimBlock>>>(d_b, .5f, N);
    set_array<<<dimGrid, dimBlock>>>(d_c, .5f, N);

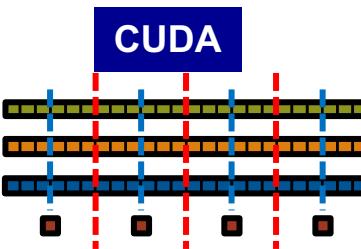
    scalar=3.0f;
    STREAM_Triad<<<dimGrid, dimBlock>>>(d_b, d_c, d_a, scalar, N);
    cudaThreadSynchronize();

    cudaFree(d_a);
    cudaFree(d_b);
    cudaFree(d_c);

}

__global__ void set_array(float *a, float value, int len) {
    int idx = threadIdx.x + blockIdx.x * blockDim.x;
    if (idx < len) a[idx] = value;
}

__global__ void STREAM_Triad( float *a, float *b, float *c,
                           float scalar, int len) {
    int idx = threadIdx.x + blockIdx.x * blockDim.x;
    if (idx < len) c[idx] = a[idx]+scalar*b[idx];
}
```



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STREAM Triad: MPI+OpenMP vs. CUDA

```
#define N      2000000
int main() {
    float *d_a, *d_b, *d_c;
    float scalar;

    cudaMalloc((void**)&d_a, sizeof(float)*N);
    cudaMalloc((void**)&d_b, sizeof(float)*N);
    cudaMalloc((void**)&d_c, sizeof(float)*N);

    dim3 dimBlock(128);
    dim3 dimGrid(N/dimBlock.x );
    if( N % dimBlock.x != 0 ) dimGrid

    set_array<<<dimGrid, dimBlock>>>(d_b, .5f, N);
    set_array<<<dimGrid, dimBlock>>>(d_c, .5f, N);

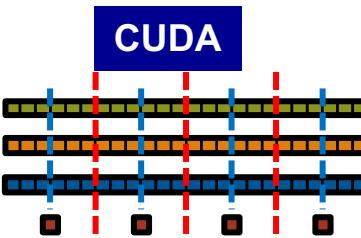
    scalar=3.0f;
    STREAM_Triad<<<dimGrid, dimBlock>>>(d_b, d_c, d_a, scalar, N);
    cudaThreadSynchronize();

    cudaFree(d_a);
    cudaFree(d_b);
    cudaFree(d_c);

}

__global__ void set_array(float *a, float value, int len) {
    int idx = threadIdx.x + blockIdx.x * blockDim.x;
    if (idx < len) a[idx] = value;
}

__global__ void STREAM_Triad( float *a, float *b, float *c,
                           float scalar, int len) {
    int idx = threadIdx.x + blockIdx.x * blockDim.x;
    if (idx < len) c[idx] = a[idx]+scalar*b[idx];
}
```



"SIMT"
Single Instruction
Multiple Threads

kernels

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STREAM Triad: MPI+OpenMP vs. CUDA

MPI + OpenMP

```
#ifdef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params)
{
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;
    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm );

    return errCount;
}

int HPCC_LocalVectorSize( double *a, int len )
{
    double scalar;
    VectorSize = HPCC_LocalVectorSize( params, 3, sizeof(double), 0 );
    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );

    if (!a || !b || !c) {
        if (c) HPCC_free(c);
        if (b) HPCC_free(b);
        if (a) HPCC_free(a);
        if (doIO) {
            fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
            fclose( outFile );
        }
        return 1;
    }

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    for (j=0; j<VectorSize; j++) {
        b[j] = 2.0;
        c[j] = 0.0;
    }

    scalar = 3.0;

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    for (j=0; j<VectorSize; j++)
        a[j] = b[j]+scalar*c[j];

    HPCC_free(c);
    HPCC_free(b);
    HPCC_free(a);
    return 0;
}
```

Why do we need so much code?

CUDA

```
#define N 2000000

int main() {
    float *d_a, *d_b, *d_c;
    float scalar;

    cudaMalloc((void**)&d_a, sizeof(float)*N);
    cudaMalloc((void**)&d_b, sizeof(float)*N);
    cudaMalloc((void**)&d_c, sizeof(float)*N);

    dim3 dimBlock(128);
    dim3 dimGrid(1/(dimBlock.x));
    if( N % dimBlock.x != 0 ) dimGrid.x++;

    set_array<<<dimGrid,dimBlock>>>(d_b, .5f, N);
    set_array<<<dimGrid,dimBlock>>>(d_c, .5f, N);

    scalar=3.0f;
    STREAM_Triad<<<dimGrid,dimBlock>>>(d_b, d_c, d_a, scalar, N);
    cudaThreadSynchronize();

    cudaFree(d_a);
    cudaFree(d_b);
    cudaFree(d_c);

    __global__ void set_array(float *a, float value, int len) {
        int idx = threadIdx.x + blockIdx.x * blockDim.x;
        if (idx < len) a[idx] = value;
    }

    __global__ void STREAM_Triad( float *a, float *b, float *c,
                                float scalar, int len) {
        int idx = threadIdx.x + blockIdx.x * blockDim.x;
        if (idx < len) c[idx] = a[idx]+scalar*b[idx];
    }
}
```

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Why so many programming models?

HPC tends to approach programming models bottom-up:

Given a system and its core capabilities...

...provide features that can access the available performance.

- portability? generality? programmability? ...not strictly required.

Type of HW Parallelism	Programming Model	Unit of Parallelism
Inter-node	MPI	executable
Intra-node/multicore	OpenMP / pthreads	iteration/task
Instruction-level vectors/threads	pragmas, intrinsics	iteration
GPU/accelerator	CUDA / Open[CL MP ACC]	iteration

benefits: lots of control; decent generality; easy to implement
downsides: lots of user-managed detail; brittle to changes

Motivation for Chapel

Q: Can a single language be...

- ...as productive as Python?
- ...as fast as Fortran?
- ...as portable as C?
- ...as scalable as MPI?
- ...as fun as <your favorite language here>?

A: We believe so.



What is Chapel?

Chapel: A productive parallel programming language

- portable
- open-source
- a collaborative effort



Goals:

- Support general parallel programming
 - “any parallel algorithm on any parallel hardware”
- Make parallel programming at scale far more productive

What does “Productivity” mean to you?

Recent Graduates:

“something similar to what I used in school: Python, Matlab, Java, ...”

Seasoned HPC Programmers:

“that sugary stuff that I don’t need because I ~~was born to suffer~~
want full control
to ensure performance”

Computational Scientists:

“something that lets me express my parallel computations
without having to wrestle with architecture-specific details”

Chapel Team:

“something that lets computational scientists express what they want,
without taking away the control that HPC programmers need,
implemented in a language as attractive as recent graduates want.”



Rewinding a few slides...

MPI + OpenMP

```
#ifdef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;
    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm );

    return errCount;
}

int HPCC_LocalVectorSize(HPCC_Params *params, int len)
{
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3, sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );

    if (!a || !b || !c) {
        if (c) HPCC_free(c);
        if (b) HPCC_free(b);
        if (a) HPCC_free(a);
        if (doIO) {
            fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
            fclose( outFile );
        }
        return 1;
    }

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    for (j=0; j<VectorSize; j++) {
        b[j] = 2.0;
        c[j] = 0.0;
    }

    scalar = 3.0;

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    for (j=0; j<VectorSize; j++)
        a[j] = b[j]+scalar*c[j];

    HPCC_free(c);
    HPCC_free(b);
    HPCC_free(a);

    return 0;
}
```

Why do we need so much code?

CUDA

```
#define N 2000000
```

```
int main() {
    float *d_a, *d_b, *d_c;
    float scalar;

    cudaMalloc((void**)&d_a, sizeof(float)*N);
    cudaMalloc((void**)&d_b, sizeof(float)*N);
    cudaMalloc((void**)&d_c, sizeof(float)*N);

    dim3 dimBlock(128);
    dim3 dimGrid(1/(dimBlock.x));
    if( N % dimBlock.x != 0 ) dimGrid.x++;

    set_array<<<dimGrid,dimBlock>>>(d_b, .5f, N);
    set_array<<<dimGrid,dimBlock>>>(d_c, .5f, N);

    scalar=3.0f;
    STREAM_Triad<<<dimGrid,dimBlock>>>(d_b, d_c, d_a, scalar, N);
    cudaThreadSynchronize();

    cudaFree(d_a);
    cudaFree(d_b);
    cudaFree(d_c);
```

```
_global_ void set_array(float *a, float value, int len) {
    int idx = threadIdx.x + blockIdx.x * blockDim.x;
    if (idx < len) a[idx] = value;
}

_global_ void STREAM_Triad( float *a, float *b, float *c,
                           float scalar, int len) {
    int idx = threadIdx.x + blockIdx.x * blockDim.x;
    if (idx < len) c[idx] = a[idx]+scalar*b[idx];
}
```

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STREAM Triad: Chapel

MPI + OpenMP

```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params,
int myRank, commSize;
int rv, errCount;
MPI_Comm comm = MPI_COMM_WORLD;
MPI_Comm_size(comm, &commSize);
MPI_Comm_rank(comm, &myRank);

rv = HPCC_Stream( params, 0 == myRank );
MPI_Reduce( &rv, &errCount, 1, MPI_
return errCount;

int HPCC_Stream(HPCC_Params *params,
register int j;
double scalar;
VectorSize = HPCC_LocalVectorSize();
a = HPCC_XMALLOC( double, VectorSize );
b = HPCC_XMALLOC( double, VectorSize );
c = HPCC_XMALLOC( double, VectorSize );

if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to allocate memory (%d).\n" VectorSize );
        cudaThreadSynchronize();
        fclose( outFile );
    }
}
```

Chapel

```
config const m = 1000,
alpha = 3.0;

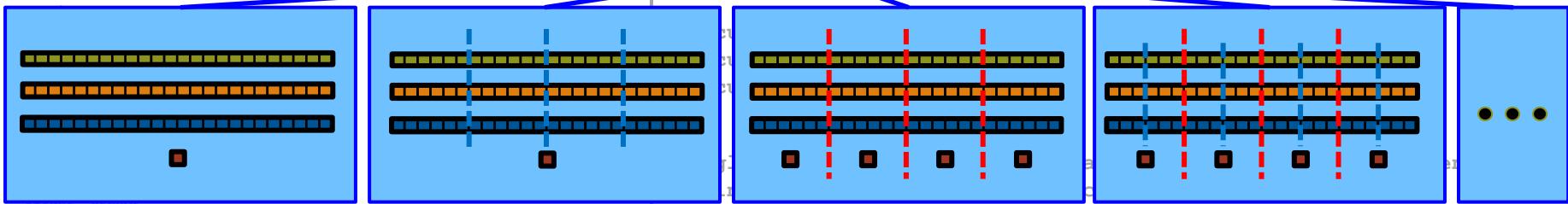
const ProblemSpace = {1..m} dmapped ...;

var A, B, C: [ProblemSpace] real;

B = 2.0;
C = 3.0;

A = B + alpha * C;
```

the special sauce



Philosophy: Good, *top-down* language design can tease system-specific implementation details away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.

Outline

- ✓ Motivation for Chapel
- Survey of Chapel Concepts
- Chapel Project and Characterizations
- Chapel Resources

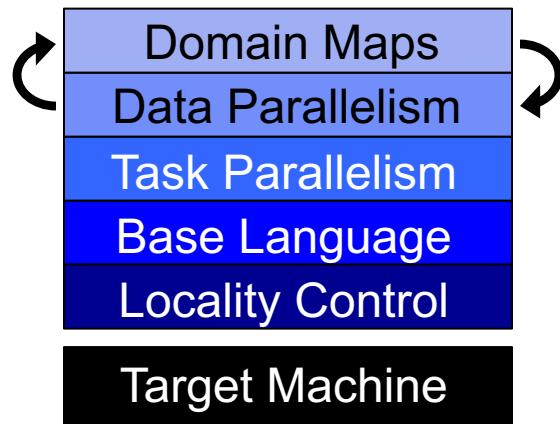


Chapel's Multiresolution Philosophy

Multiresolution Design: Support multiple tiers of features

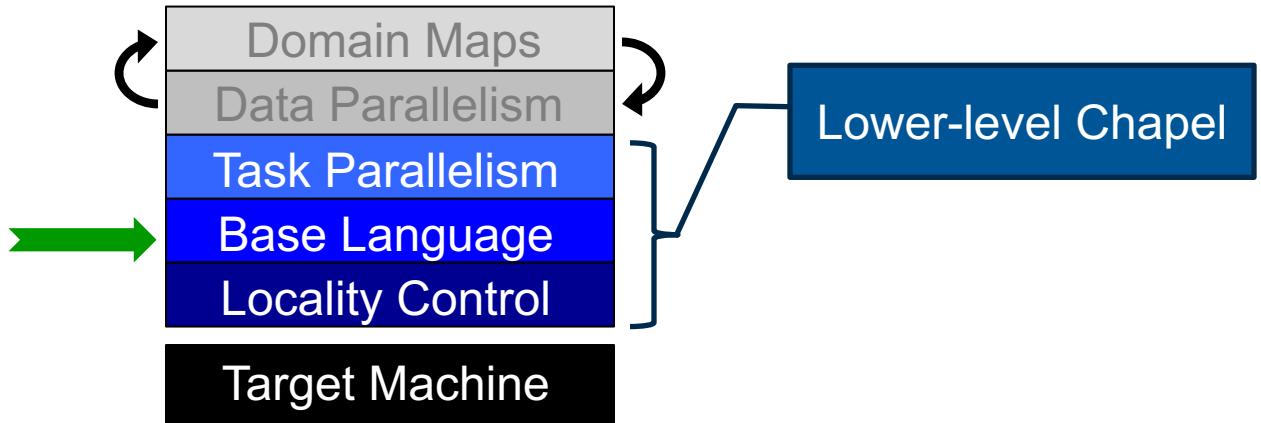
- higher levels for programmability, productivity
- lower levels for greater degrees of control

Chapel language concepts



- build the higher-level concepts in terms of the lower
- permit the user to intermix layers arbitrarily

Base Language



Base Language Features, by example

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <=> next;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```

Base Language Features, by example

CLU-style iterators

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <=gt; next;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```

Base Language Features, by example

```
iter fib(n) {  
    var current = 0,  
        next = 1;  
  
    for i in 1..n {  
        yield current;  
        current += next;  
        current <=> next;  
    }  
}
```

Configuration declarations
(to avoid command-line argument parsing)
. ./a.out -n=1000000

```
config const n = 10;  
  
for f in fib(n) do  
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```

Base Language Features, by example

Static type inference for:

- arguments
- return types
- variables

```
iter fib(n) {  
    var current = 0,  
        next = 1;  
  
    for i in 1..n {  
        yield current;  
        current += next;  
        current <=> next;  
    }  
}
```

```
config const n = 10;  
  
for f in fib(n) do  
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```

Base Language Features, by example

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <=> next;
    }
}
```

config const n = 10;

```
for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...

Zippered iteration

Base Language Features, by example

Range types and operators

```
iter fib(n) {  
    var current = 0  
        next = 1;  
  
    for i in 1..n {  
        yield current;  
        current += next;  
        current <=gt; next;  
    }  
}
```

```
config const n = 10;  
  
for (i,f) in zip(0..#n, fib(n)) do  
writeln("fib #", i, " is ", f);
```

```
fib #0 is 0  
fib #1 is 1  
fib #2 is 1  
fib #3 is 2  
fib #4 is 3  
fib #5 is 5  
fib #6 is 8  
...
```

Base Language Features, by example

```
iter fib(n) {  
    var current = 0,  
        next = 1;  
  
    for i in 1..n {  
        yield current;  
        current += next;  
        current <=> next;  
    }  
}
```

```
config const n = 10;  
  
for (i,f) in zip(0..#n, fib(n)) do  
writeln("fib #", i, " is ", f);
```

fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...



Base Language Features, by example

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <=> next;
    }
}
```

```
config const n = 10;

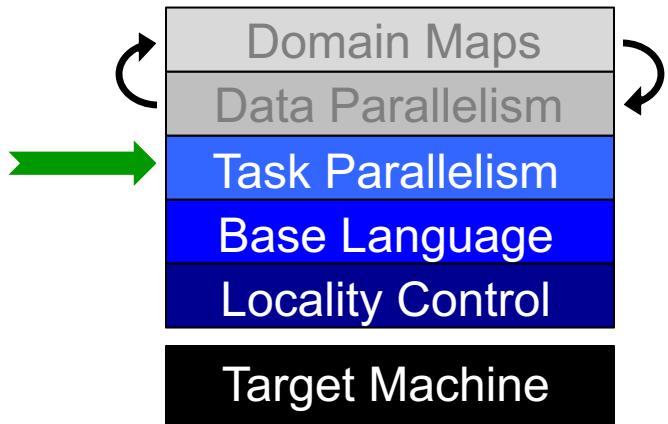
for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```

Other Base Language Features

- **interoperability features**
- **OOP** (value- and reference-based)
- **overloading, where clauses**
- **argument intents, default values, match-by-name**
- **compile-time features for meta-programming**
 - e.g., compile-time functions to compute types, values; reflection
- **modules** (for namespace management)
- **rank-independent programming features**
- ...

Task Parallelism



Task Parallelism: Begin Statements

```
// create a fire-and-forget task for a statement
begin writeln("hello world");
writeln("goodbye");
```

Possible outputs:

hello world
goodbye

goodbye
hello world

Task Parallelism: Coforall Loops

```
// create a task per iteration
coforall t in 0..#numTasks {
    writeln("Hello from task ", t, " of ", numTasks);
} // implicit join of the numTasks tasks here

writeln("All tasks done");
```

Sample output:

```
Hello from task 2 of 4
Hello from task 0 of 4
Hello from task 3 of 4
Hello from task 1 of 4
All tasks done
```



Task Parallelism: Data-Driven Synchronization

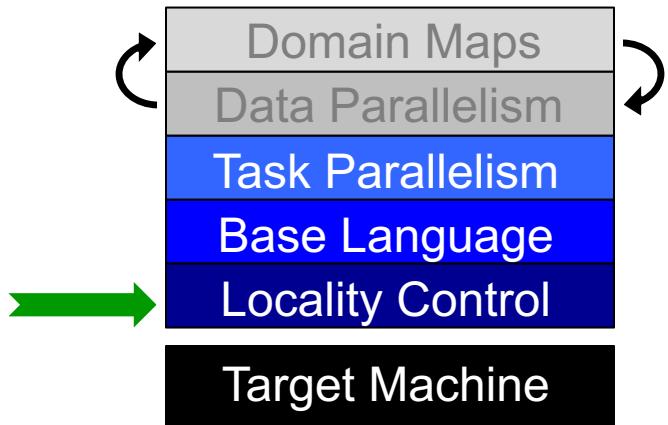
- **atomic variables:** support atomic operations
 - e.g., compare-and-swap; atomic sum, multiply, etc.
 - similar to C/C++

- **sync variables:** store full-empty state along with value
 - by default, reads/writes block until full/empty, leave in opposite state

Other Task Parallel Concepts

- **cobegins**: create tasks using compound statements
- **single variables**: like sync variables, but write-once
- **sync statements**: join unstructured tasks
- **serial statements**: conditionally squash parallelism

Locality Control



The Locale Type

Definition:

- Abstract unit of target architecture
- Supports reasoning about locality
 - defines “here vs. there” / “local vs. remote”
- Capable of running tasks and storing variables
 - i.e., has processors and memory

Typically: A compute node (multicore processor or SMP)



Getting started with locales

- Specify # of locales when running Chapel programs

```
% a.out --numLocales=8
```

```
% a.out -nl 8
```

- Chapel provides built-in locale variables

```
config const numLocales: int = ...;  
const Locales: [0..#numLocales] locale = ...;
```

Locales



- main () starts execution as a task on locale #0



Locale Operations

- Locale methods support queries about the target system:

```
proc locale.physicalMemory(...) { ... }
proc locale.numCores { ... }
proc locale.id { ... }
proc locale.name { ... }
```

- On-clauses support placement of computations:

```
writeln("on locale 0");
on Locales[1] do
    writeln("now on locale 1");
writeln("on locale 0 again");
```

```
on A[i,j] do
    bigComputation(A);
on node.left do
    search(node.left);
```

Parallelism and Locality: Orthogonal in Chapel

- This is a **parallel**, but local program:

```
coforall i in 1..msgs do  
    writeln("Hello from task ", i);
```

- This is a **distributed**, but serial program:

```
writeln("Hello from locale 0!");  
on Locales[1] do writeln("Hello from locale 1!");  
on Locales[2] do writeln("Hello from locale 2!");
```

- This is a **distributed parallel** program:

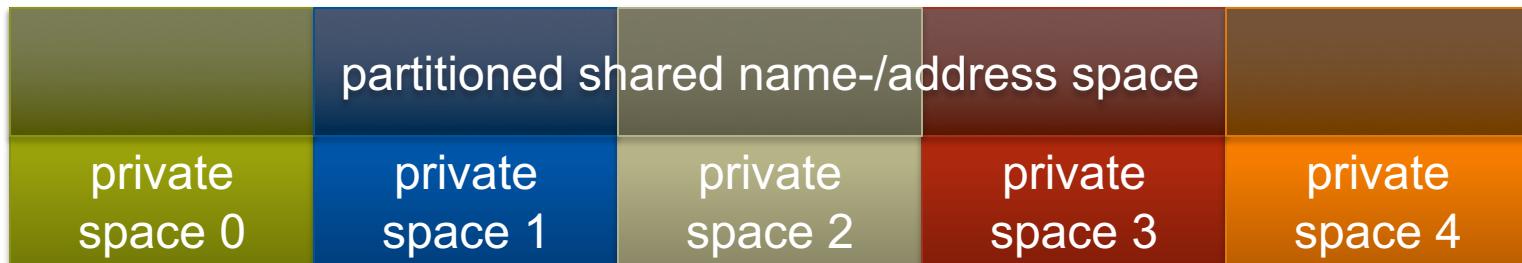
```
coforall i in 1..msgs do  
    on Locales[i%numLocales] do  
        writeln("Hello from task ", i,  
               " running on locale ", here.id);
```

Partitioned Global Address Space (PGAS) Languages



(Or perhaps: partitioned global namespace languages)

- **abstract concept:**
 - support a shared namespace on distributed memory
 - permit parallel tasks to access remote variables by naming them
 - establish a strong sense of ownership
 - every variable has a well-defined location
 - local variables are cheaper to access than remote ones
- **traditional PGAS languages have been SPMD in nature**
 - best-known examples: Fortran Co-Arrays, UPC



COMPUTE

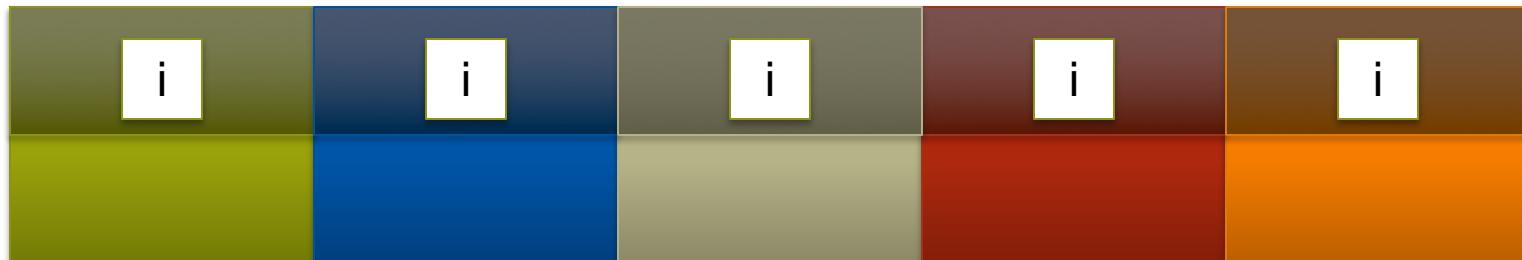
STORE

ANALYZE

SPMD PGAS Languages

(using a pseudo-language, not Chapel)

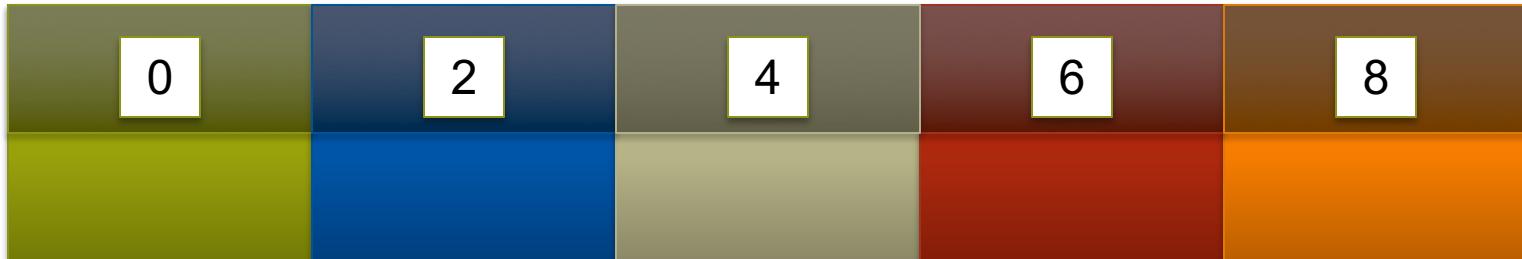
```
shared int i(*) ; // declare a shared variable i
```



SPMD PGAS Languages (using a pseudo-language, not Chapel)

```
shared int i(*) ; // declare a shared variable i  
function main() {  
    i = 2*this_image(); // each image initializes its copy
```

i=



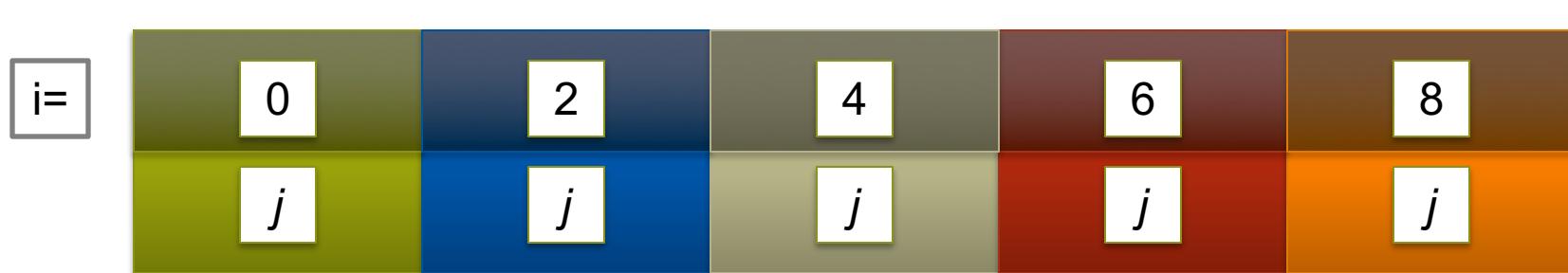
COMPUTE

STORE

ANALYZE

SPMD PGAS Languages

```
shared int i(*) ; // declare a shared variable i  
function main() {  
    i = 2*this_image(); // each image initializes its copy  
  
    private int j; // declare a private variable j
```

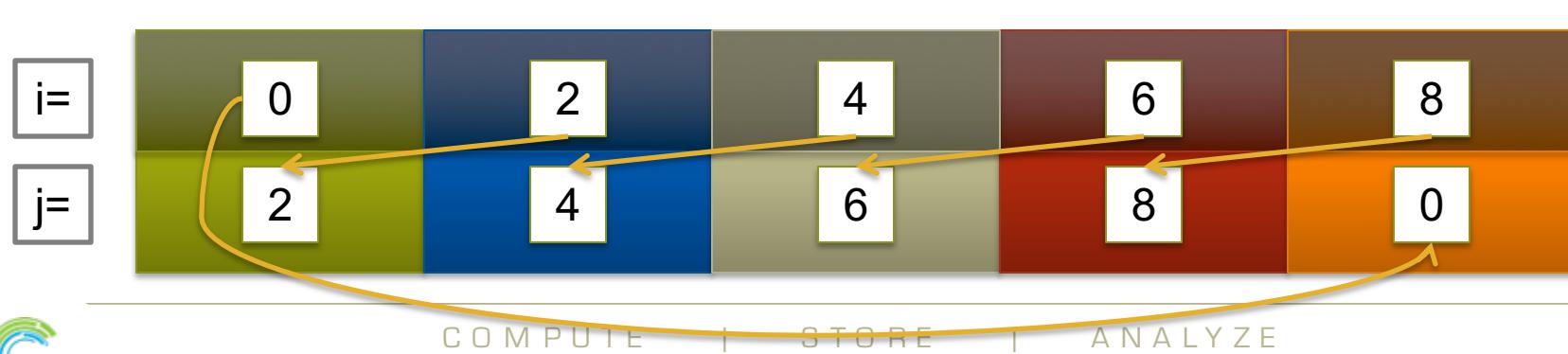


SPMD PGAS Languages (using a pseudo-language, not Chapel)

```

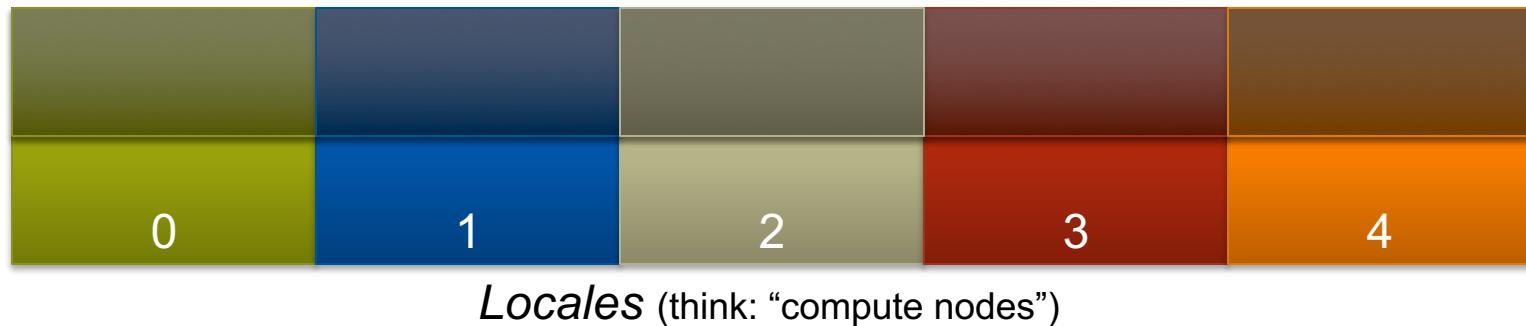
shared int i(*) ; // declare a shared variable i
function main() {
    i = 2*this_image() ; // each image initializes its copy
    barrier();
    private int j; // declare a private variable j
    j = i( (this_image()+1) % num_images() );
    // ^ access our neighbor's copy of i
    // communication implemented by compiler + runtime
    // How did we know our neighbor had an i?
    // Because it's SPMD - we're all running the same
    // program. (Simple, but restrictive)

```



Chapel and PGAS

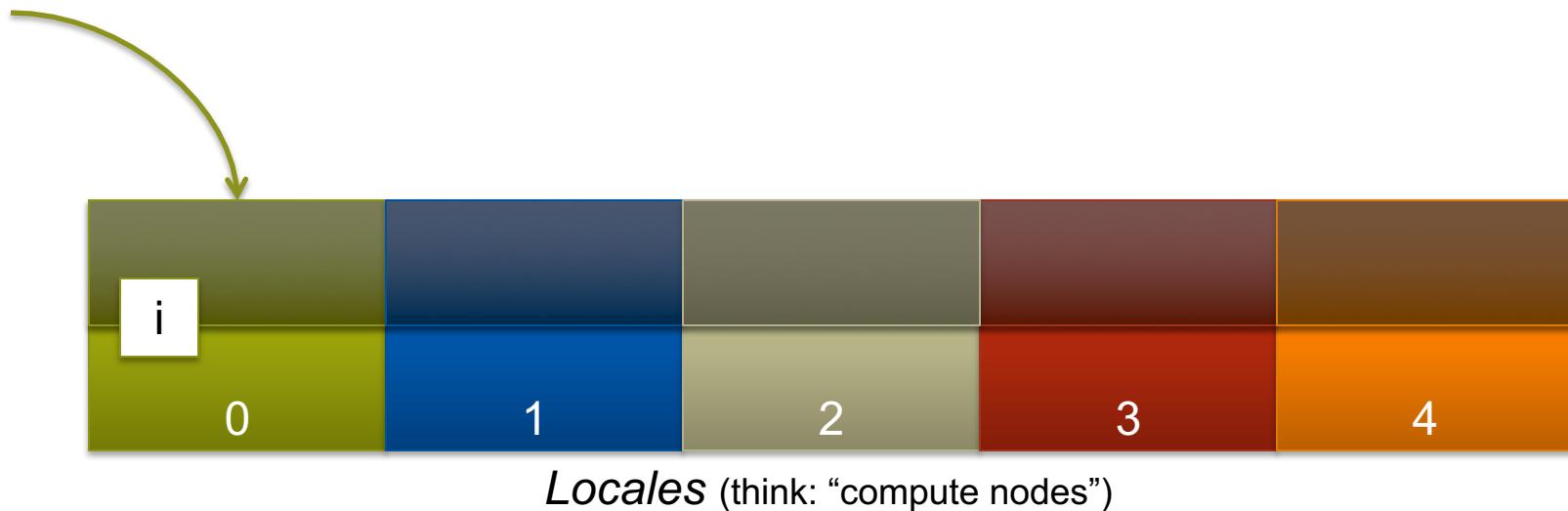
- Chapel is PGAS, but unlike most, it's not inherently SPMD
 - never think about “the other copies of the program”
 - “global name/address space” comes from lexical scoping
 - as in traditional languages, each declaration yields one variable
 - variables are stored on the locale where the task declaring it is executing



COMPUTE | STORE | ANALYZE

Chapel: Scoping and Locality

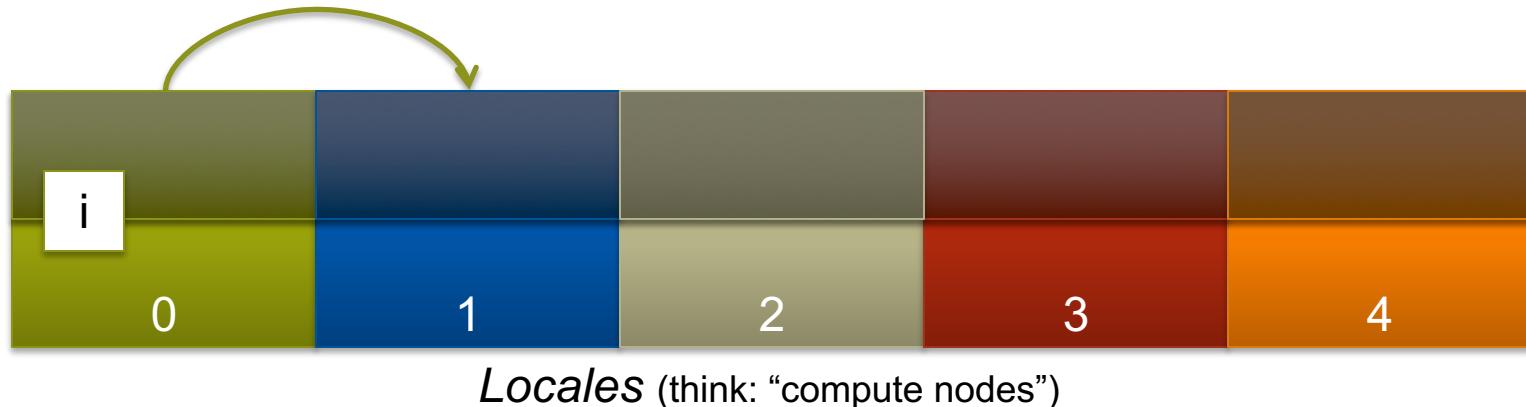
```
var i: int;
```



COMPUTE | STORE | ANALYZE

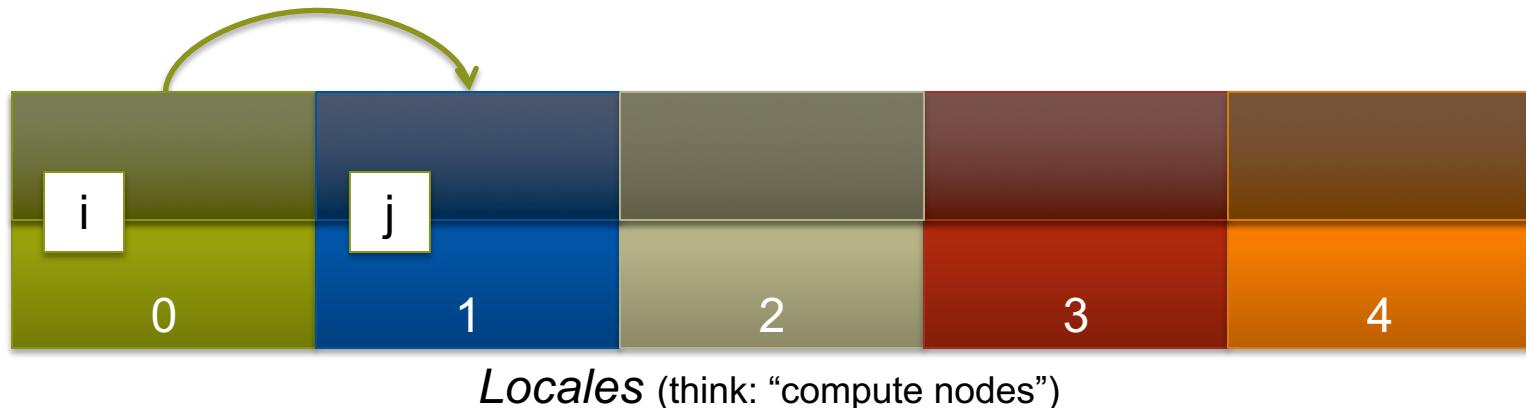
Chapel: Scoping and Locality

```
var i: int;  
on Locales[1] {
```



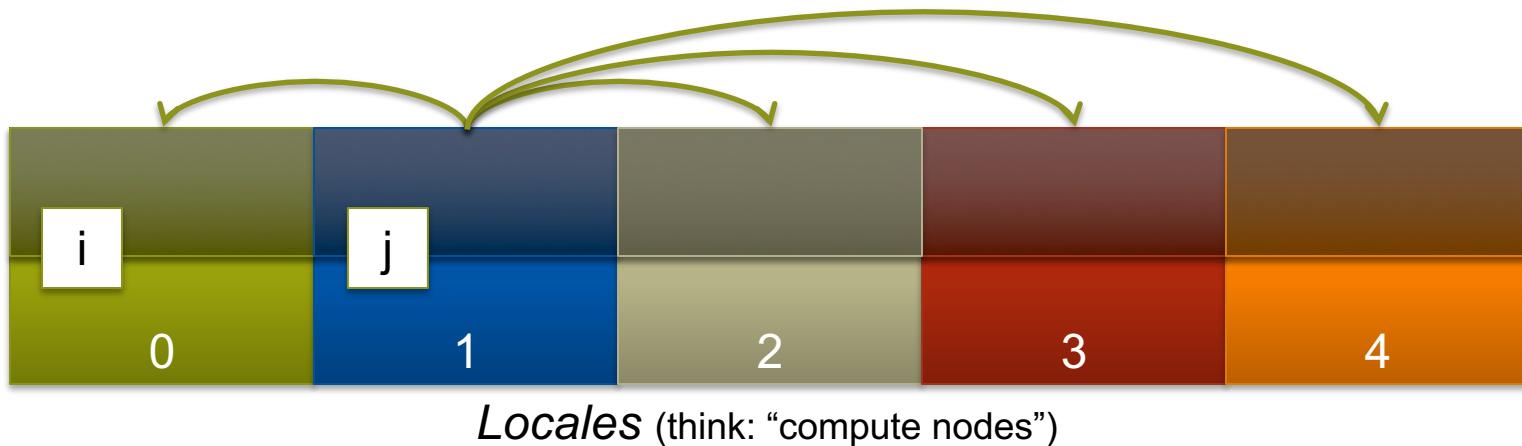
Chapel: Scoping and Locality

```
var i: int;  
on Locales[1] {  
    var j: int;
```



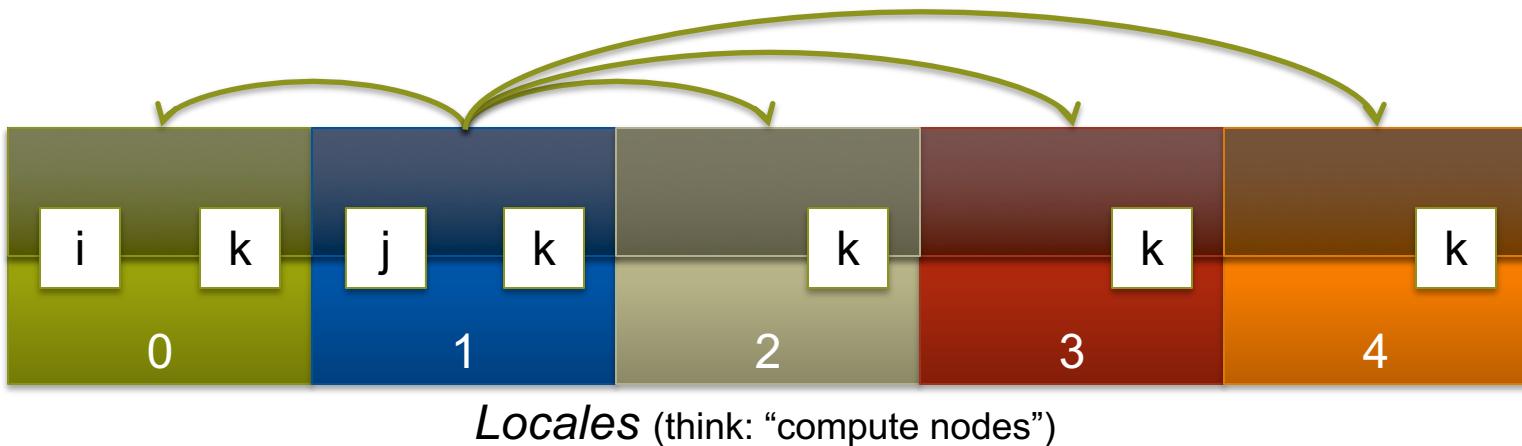
Chapel: Scoping and Locality

```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {
```



Chapel: Scoping and Locality

```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
  
            ...  
        }  
    }  
}
```

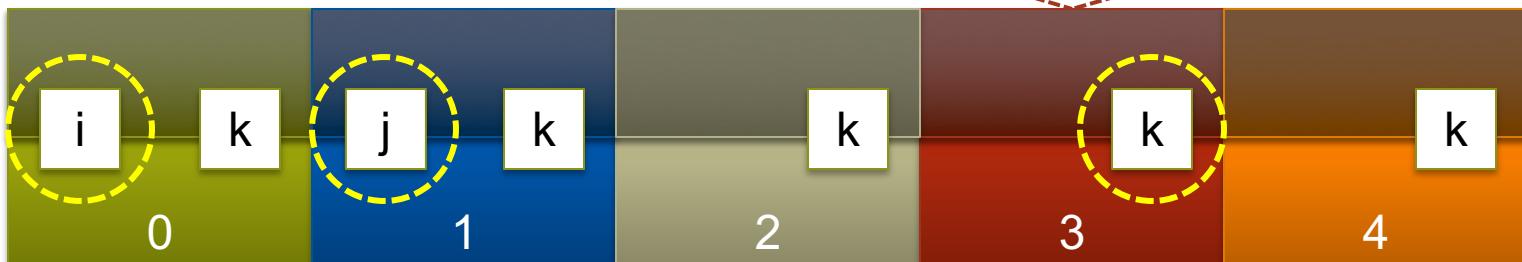


Chapel: Scoping and Locality

```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
            k = 2*i + j;  
        }  
    }  
}
```

OK to access i , j , and k wherever they live

$k = 2*i + j;$



Locales (think: “compute nodes”)

Chapel: Scoping and Locality

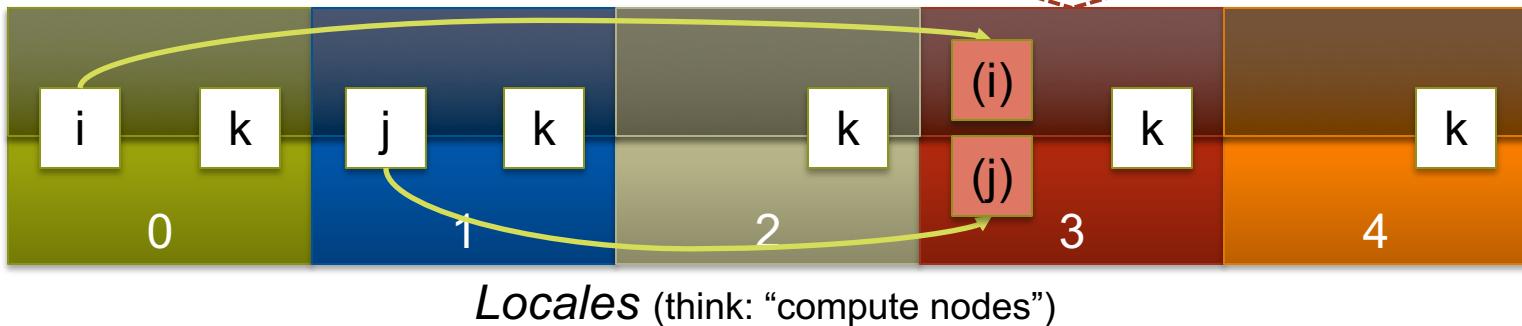
```

var i: int;
on Locales[1] {
    var j: int;
    coforall loc in Locales {
        on loc {
            var k: int;
            k = 2*i + j;
        }
    }
}

```

here, *i* and *j* are remote, so
the compiler + runtime will
transfer their values

$k = 2*i + j;$



Locales (think: “compute nodes”)

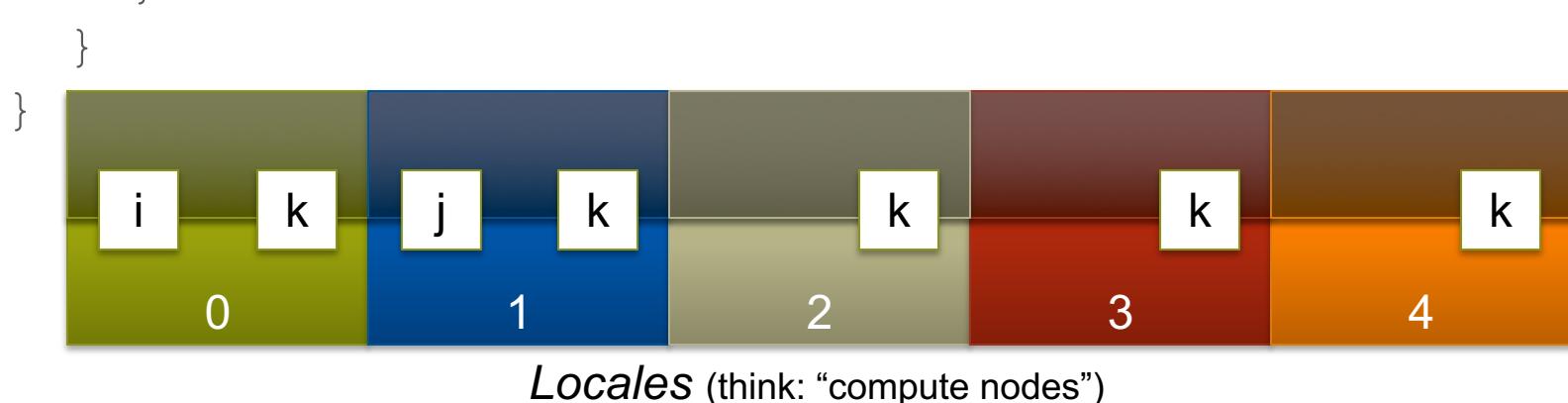
Chapel: Locality queries

```

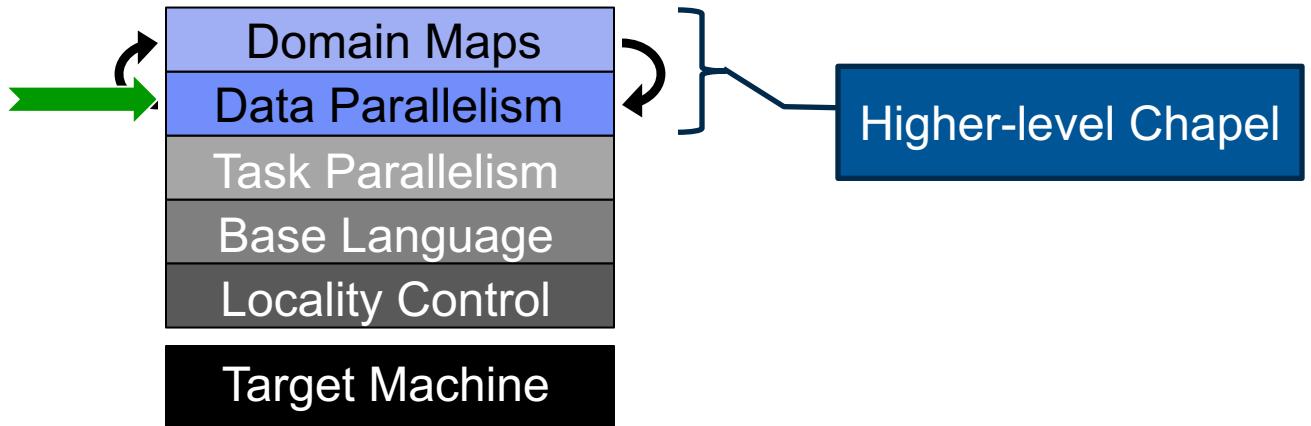
var i: int;
on Locales[1] {
    var j: int;
    coforall loc in Locales {
        on loc {
            var k: int;

                ...here...           // query the locale on which this task is running
                ...j.locale...         // query the locale on which j is stored
        }
    }
}

```



Data Parallelism

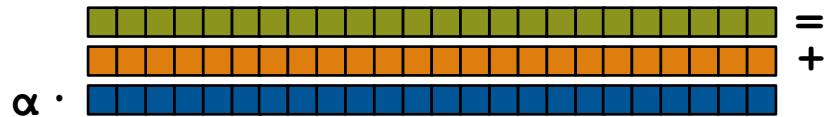


Data Parallelism By Example: STREAM Triad

```
const ProblemSpace = {1..m};
```



```
var A, B, C: [ProblemSpace] real;
```



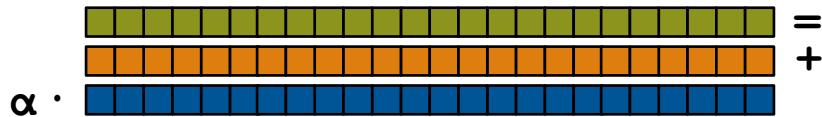
```
forall (a,b,c) in zip(A,B,C) do  
  a = b + alpha*c;
```

Data Parallelism By Example: STREAM Triad

```
const ProblemSpace = {1..m};
```



```
var A, B, C: [ProblemSpace] real;
```

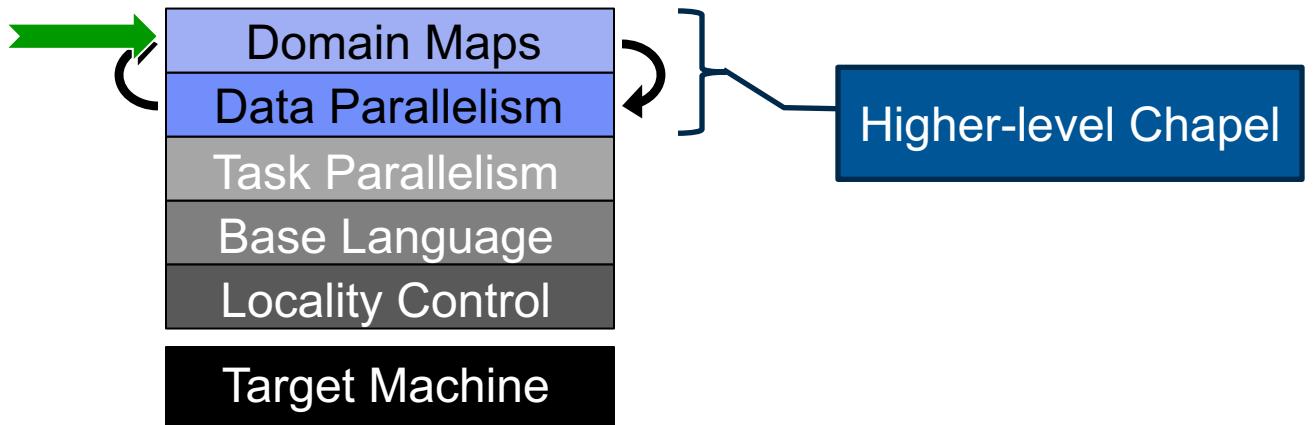


```
A = B + alpha * C; // equivalent to the zippered forall
```

Other Data Parallel Features

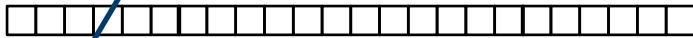
- **Rich Domain/Array Types:**
 - multidimensional
 - strided
 - sparse
 - associative
- **Slicing:** Refer to subarrays using ranges/domains
 - ... `A[2..n-1, lo..#b]` ...
 - ... `A[ElementsOfInterest]` ...
- **Promotion:** Call scalar functions with array arguments
 - ... `pow(A, B)` ... // equivalent to: `forall (a,b) in zip(A,B) do pow(a,b)`
- **Reductions/Scans:** Apply operations across collections
 - ... `+ reduce A` ...
 - ... `myReduceOp reduce A` ...

Domain Maps

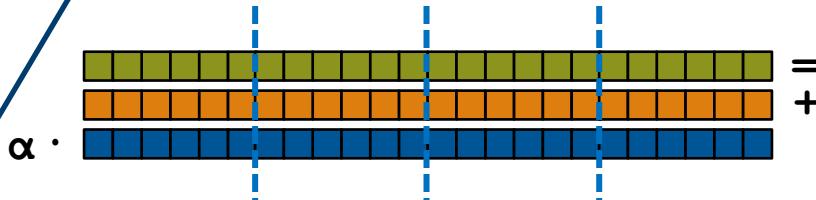


STREAM Triad: Chapel (multicore)

```
const ProblemSpace = {1..m};
```



```
var A, B, C: [ProblemSpace] real;
```

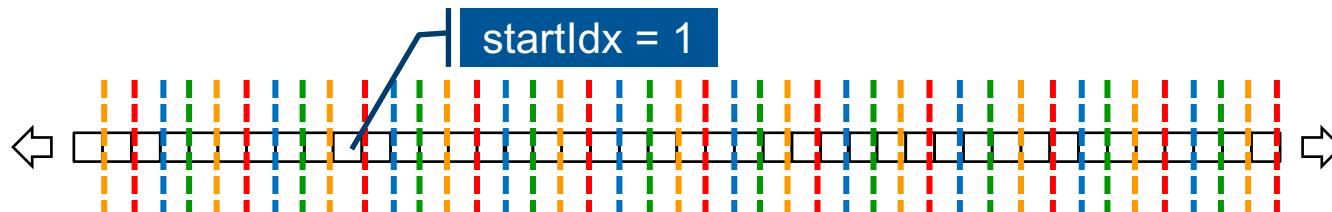


```
A = B + alpha * C;
```

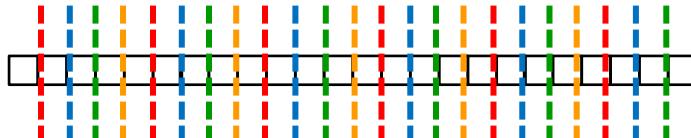
No domain map specified \Rightarrow use default layout

- current locale owns all domain indices and array values
- computation will execute using local processors only

STREAM Triad: Chapel (multilocale, cyclic)



```
const ProblemSpace = {1..m}
    dmapped Cyclic(startIdx=1);
```

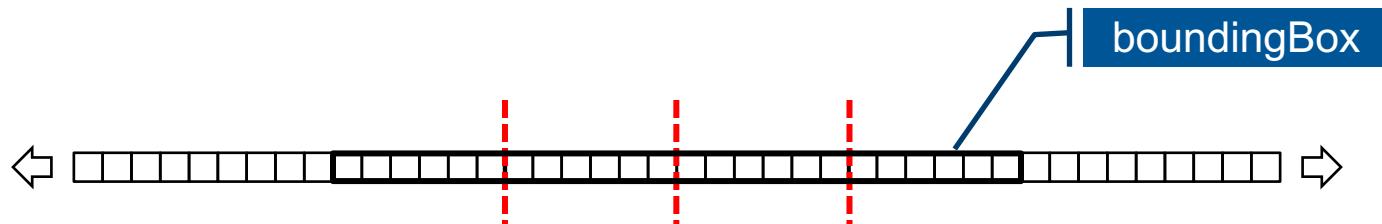


```
var A, B, C: [ProblemSpace] real;
```

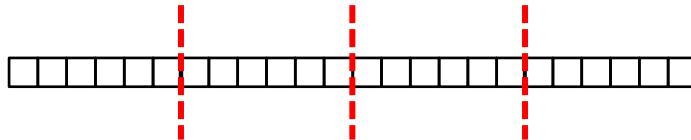


```
A = B + alpha * C;
```

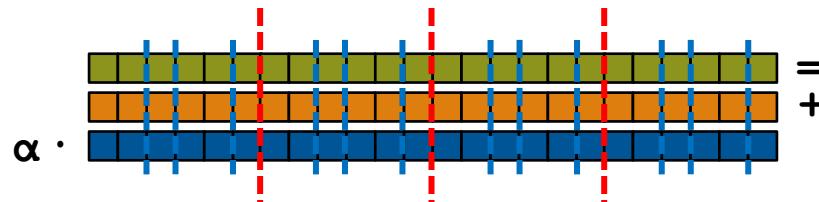
STREAM Triad: Chapel (multilocale, blocked)



```
const ProblemSpace = {1..m}
    dmapped Block (boundingBox={1..m});
```



```
var A, B, C: [ProblemSpace] real;
```



```
A = B + alpha * C;
```

STREAM Triad: Chapel

MPI + OpenMP

```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params,
int myRank, commSize;
int rv, errCount;
MPI_Comm comm = MPI_COMM_WORLD;
MPI_Comm_size(comm, &commSize);
MPI_Comm_rank(comm, &myRank);

rv = HPCC_Stream( params, 0 == myRank );
MPI_Reduce( &rv, &errCount, 1, MPI_
return errCount;

int HPCC_Stream(HPCC_Params *params,
register int j;
double scalar;
VectorSize = HPCC_LocalVectorSize();
a = HPCC_XMALLOC( double, VectorSize );
b = HPCC_XMALLOC( double, VectorSize );
c = HPCC_XMALLOC( double, VectorSize );

if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to allocate memory (%d).\n" VectorSize );
        cudaThreadSynchronize();
        fclose( outFile );
    }
}
```

Chapel

```
config const m = 1000,
alpha = 3.0;

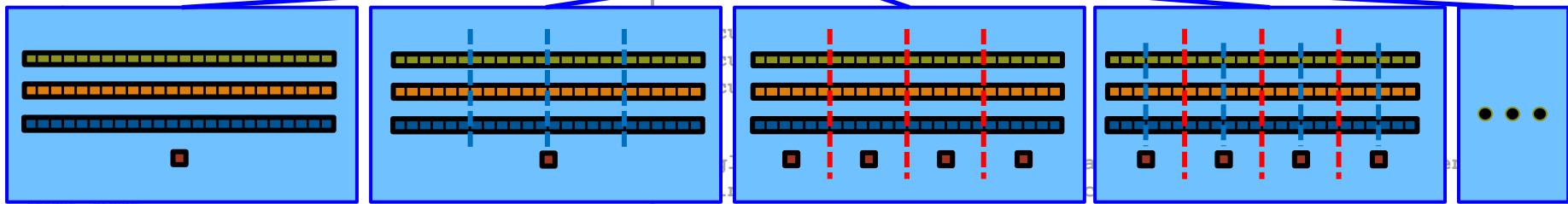
const ProblemSpace = {1..m} dmapped ...;

var A, B, C: [ProblemSpace] real;

B = 2.0;
C = 3.0;

A = B + alpha * C;
```

the special sauce



Philosophy: Good, *top-down* language design can tease system-specific implementation details away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.

Language Summary

Parallel programmers deserve better programming models

Higher-level programming models can help insulate algorithms from parallel implementation details

- yet, without necessarily abdicating control
- Chapel does this via its multiresolution design

We believe Chapel can greatly improve productivity

- ...for current and emerging parallel architectures
- ...for HPC users as well as mainstream uses of parallelism



Outline

- ✓ Motivation for Chapel
- ✓ Survey of Chapel Concepts
- Chapel Project and Characterizations
- Chapel Resources



Computer Language Benchmarks Game

Chapel was recently added to the game:

As of December 2017:

- for performance:

1 top entries: pidigits

4 top-5 entries:

fannkuch-redux, chameneos-redux,
meteor, thread-ring

2 top-10 entries:, mandelbrot, fasta

3 top-20 entries: spectral-norm,
regex-redux

- for code compactness:

2 top entries: n-body, thread-ring

1 top-5 entries: spectral-norm

2 top-10 entry: pidigits, regex-redux

3 top-20 entries: mandelbrot,
chameneos-redux, meteor

The Computer Language
Benchmarks Game

64-bit quad core data set

Will your toy benchmark program be faster if you write it in a different programming language? It depends how you write it!

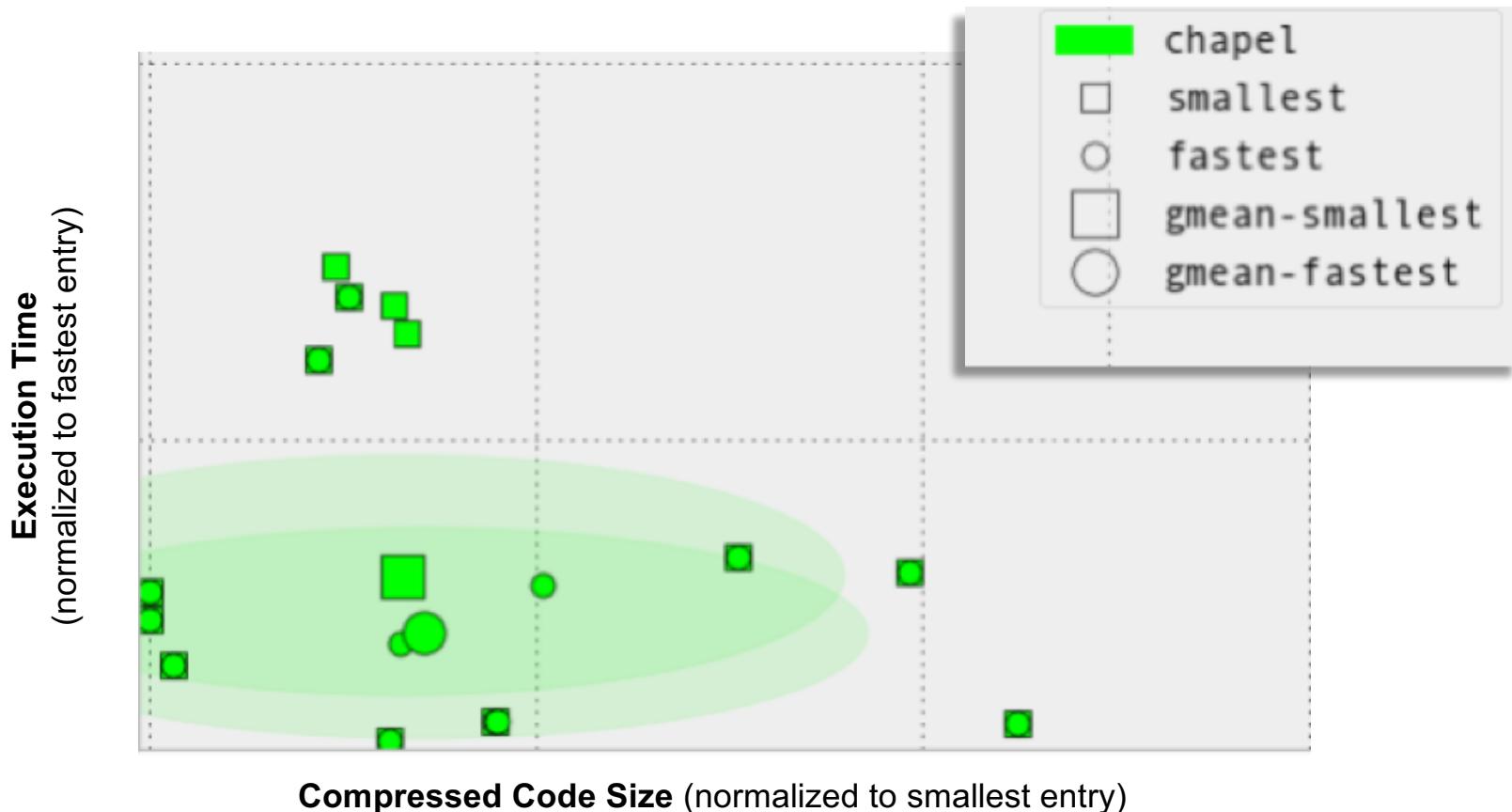
Which programs are fast?

Which are succinct? Which are efficient?

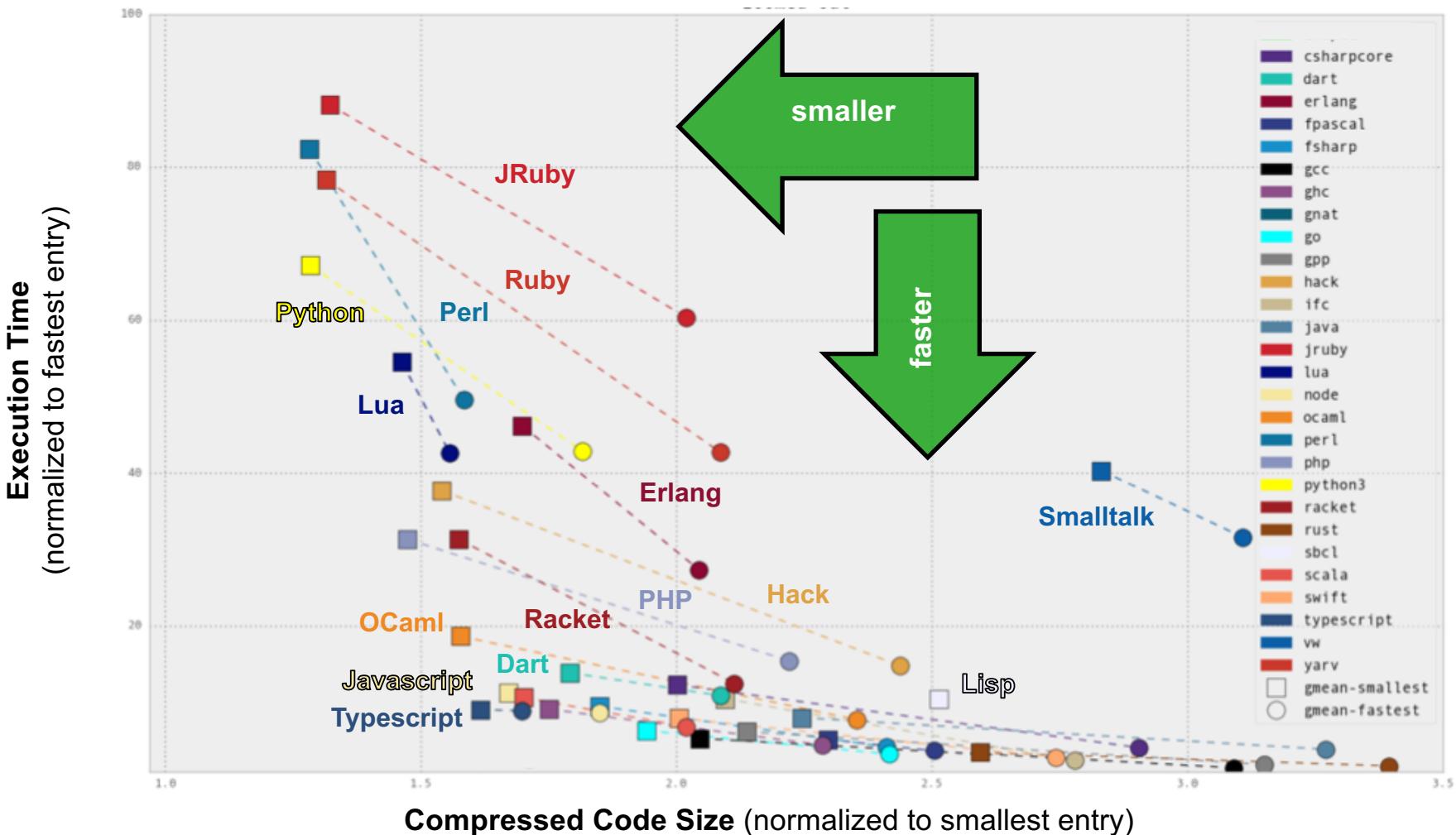
<u>Ada</u>	<u>C</u>	<u>Chapel</u>	<u>C#</u>	<u>C++</u>	<u>Dart</u>
<u>Erlang</u>	<u>F#</u>	<u>Fortran</u>	<u>Go</u>	<u>Hack</u>	
<u>Haskell</u>	<u>Java</u>	<u>JavaScript</u>	<u>Lisp</u>	<u>Lua</u>	
<u>OCaml</u>	<u>Pascal</u>	<u>Perl</u>	<u>PHP</u>	<u>Python</u>	
<u>Racket</u>	<u>Ruby</u>	<u>Rust</u>	<u>Smalltalk</u>	<u>Swift</u>	
			<u>TypeScript</u>		



Scatter plots of CLBG code size x speed



CLBG Cross-Language Summary (Oct 2017 standings)

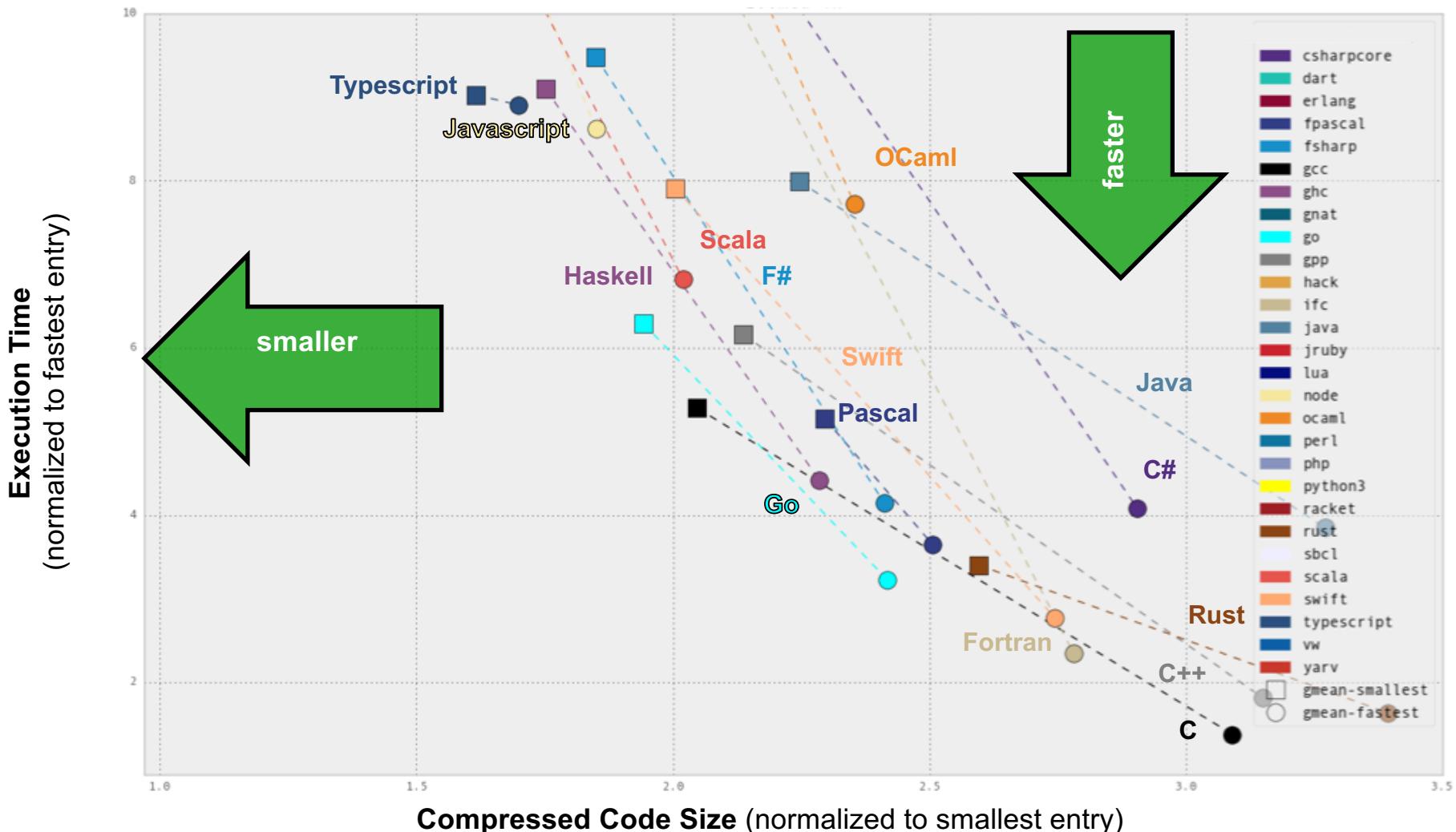


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CLBG Cross-Language Summary (Oct 2017 standings)



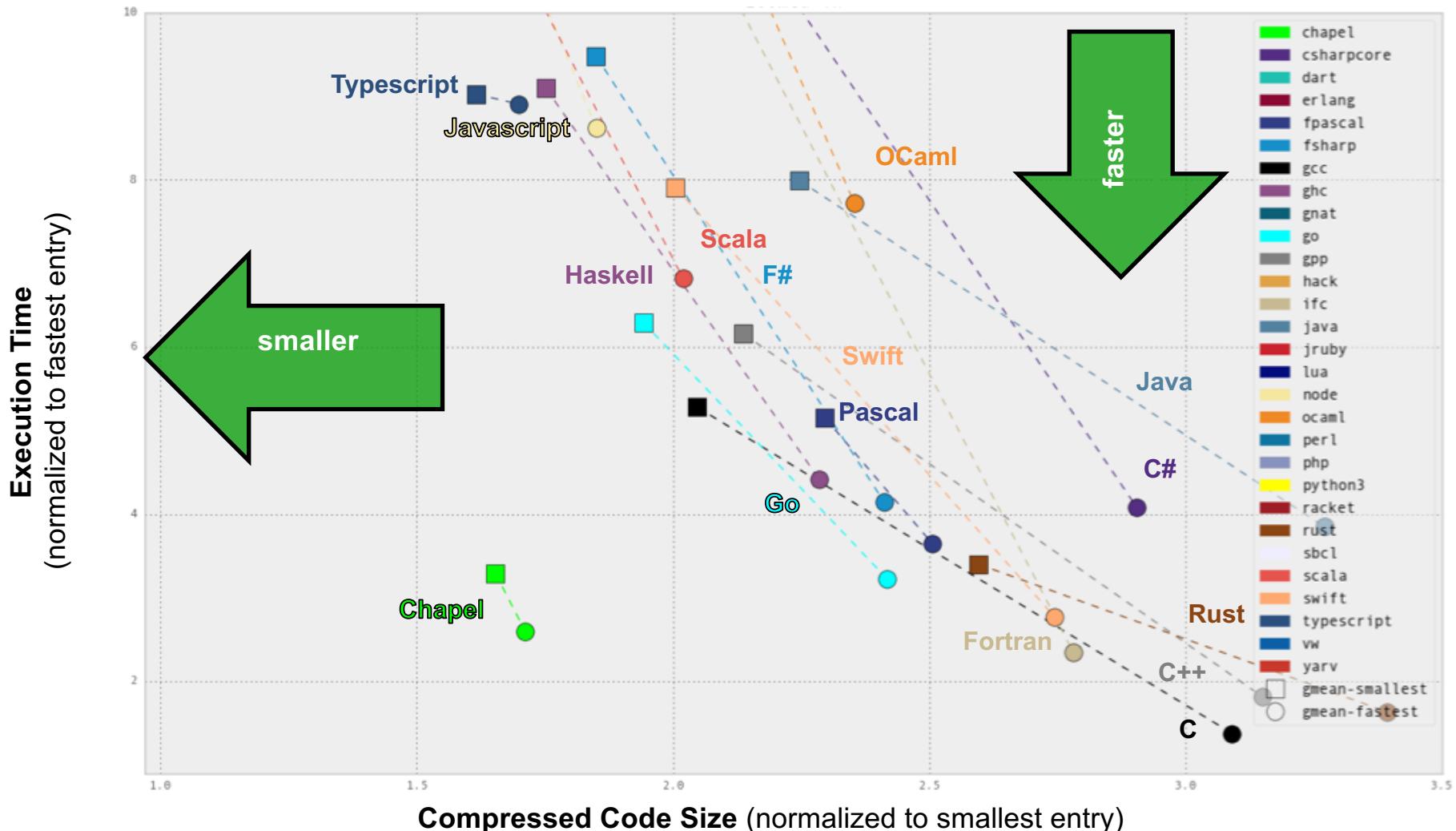
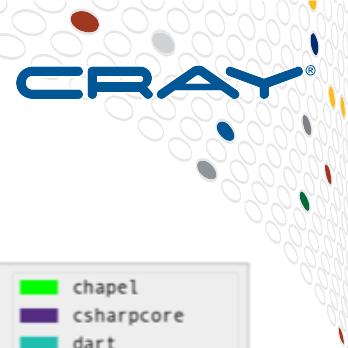
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ANALYZE

CLBG Cross-Language Summary

(Oct 2017 standings)



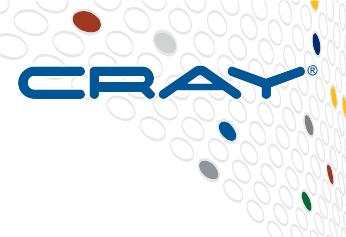
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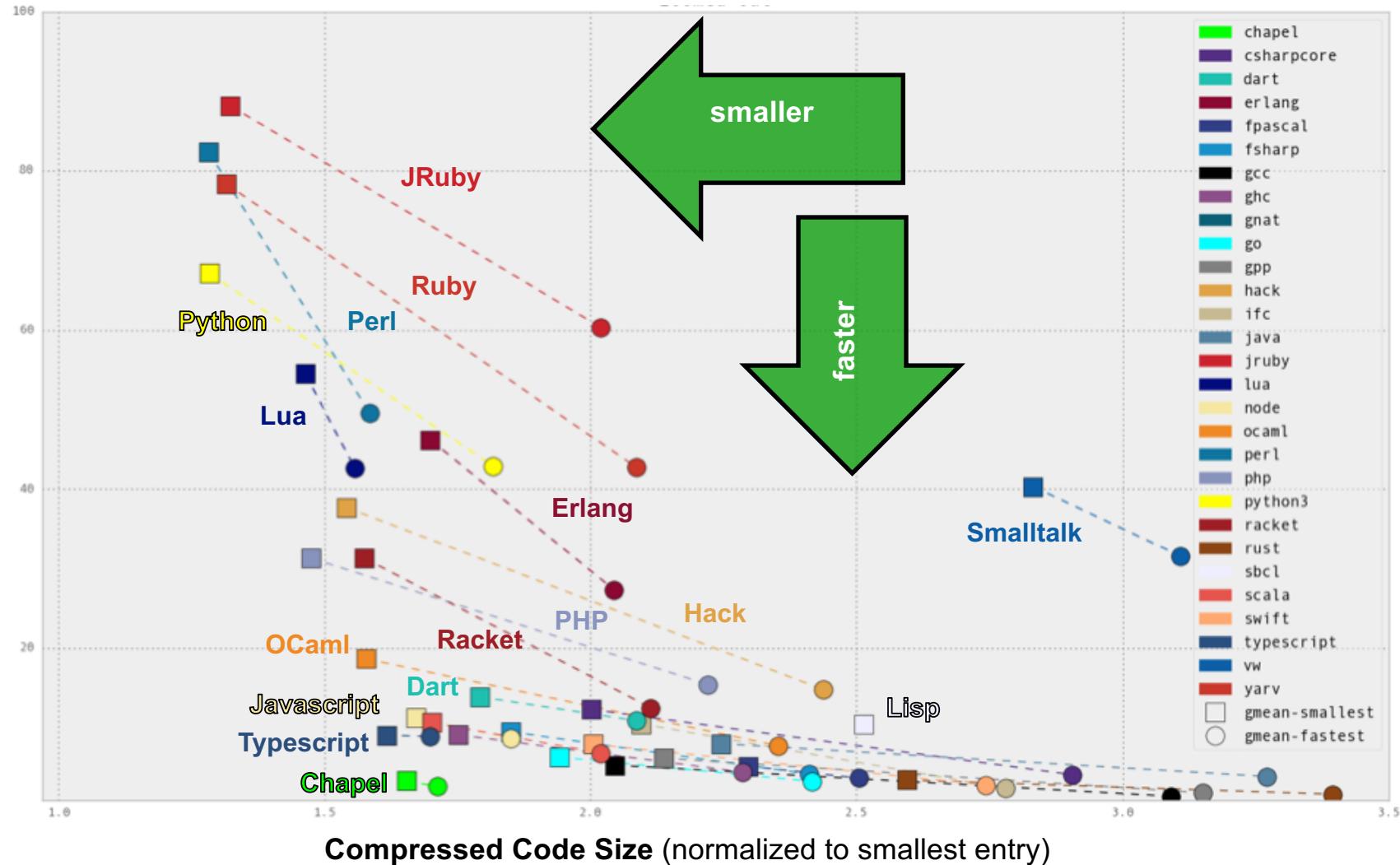
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CLBG Cross-Language Summary

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Execution Time
(normalized to fastest entry)



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Chapel is Portable

- **Chapel is designed to be hardware-independent**
- **The current release requires:**
 - a C/C++ compiler
 - a *NIX environment (Linux, OS X, BSD, Cygwin, ...)
 - POSIX threads
 - RDMA, MPI, or UDP (for distributed memory execution)
- **Chapel can run on...**
 - ...laptops and workstations
 - ...commodity clusters
 - ...the cloud
 - ...HPC systems from Cray and other vendors
 - ...modern processors like Intel Xeon Phi, GPUs*, etc.

* = academic work only; not yet supported in the official release



Chapel is Open-Source

- Chapel's development is hosted at GitHub
 - <https://github.com/chapel-lang>
- Chapel is licensed as Apache v2.0 software
- Instructions for download + install are online
 - see <http://chapel.cray.com/download.html> to get started



The Chapel Team at Cray (Summer 2016)





Chapel is a Collaborative Effort



Lawrence Berkeley
National Laboratory



Yale

(and several others...)

<https://chapel-lang.org/collaborations.html>



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Chapel Status

- **Currently being picked up by early adopters**
 - over two releases, 3000+ downloads per year
 - Users who try it generally like what they see
- **Most current features are functional and working well**
- **Performance is compelling in many cases**



A notable early adopter

Chapel in the (Cosmological) Wild

Nikhil Padmanabhan, Yale University Professor, Physics & Astronomy

Abstract: This talk aims to present my personal experiences using Chapel in my research. My research interests are in observational cosmology; more specifically, I use large surveys of galaxies to constrain the evolution of the

The image shows a YouTube video player interface. At the top left is the YouTube logo. To its right is a search bar with a magnifying glass icon. Below the search bar, there's a navigation menu with 'Videos', 'Playlists', and 'Channels'. On the left side of the main content area is a thumbnail image of a man (Nikhil Padmanabhan) standing in front of a projection screen, gesturing with his hands. The projection screen displays a complex, multi-colored visualization, likely related to cosmological data. In the bottom right corner of the thumbnail, the duration '56:14' is visible. The main title of the video is 'CHIUW 2016 keynote: "Chapel in the (Cosmological) Wild", Nikhil Padmanabhan'. Below the title, it says 'Chapel Parallel Programming Language', '1 month ago • 86 views', and a brief description: 'This is Nikhil Padmanabhan's keynote talk from CHIUW 2016: the 3rd Annual Chapel Implementers and Users workshop. The slides are availabl...'

Outline

- ✓ Motivation for Chapel
- ✓ Survey of Chapel Concepts
- ✓ Chapel Project and Characterizations
- Chapel Resources

Chapel Websites

Project page: <http://chapel.cray.com>

- overview, papers, presentations, language spec, ...

GitHub: <https://github.com/chapel-lang>

- download Chapel; browse source repository; contribute code

Facebook: <https://www.facebook.com/ChapelLanguage>

Twitter: <https://twitter.com/ChapelLanguage>



The screenshot shows the Chapel Programming Language page on Facebook. It features a large green and blue 'C' logo. A sidebar on the left lists 'Chapel highlights' such as syntax constructs for task parallelism, support for local/affinity resources, static type inference, and modules for namespace management. A central post reads: 'Chapel Programming Language is on Facebook. To connect with Chapel Programming Language, sign up for Facebook today.' Below the post is some sample Chapel code:

```

taskParallel.chpl
forall i in 1..n
    const numTasks = config const n = 1000;
    forall tid in
        writer("Hello from task %d of %d running on %s\n", tid, n,
            helloFromTask);
    end
end

dataParallel.chpl
use CyclicList;
config const n = 1000;
var D = (1..n, 1..n) mapped Cyclic(startIdx = (1,1));
var A: [D] real;
    
```

...and much, much more.



The screenshot shows the Chapel Language Twitter profile (@ChapelLanguage). It has 4 tweets, 1 follower, and 19 following. A recent tweet from March 8, 2016, states: 'Chapel is a productive parallel programming language designed for large-scale computing whose development is being led by @cray_inc'. Below the tweet is a link to their GitHub page: [chapel.sourceforge.net/perf/chapcs/?s...](https://github.com/chapel-lang/chapel). To the right is a line graph titled 'Binary Tree Shootout Benchmark (msec)' comparing different memory allocators: 'jemalloc', 'tcmalloc', 'jemalloc', and 'tcmalloc'. The graph shows performance metrics over time.



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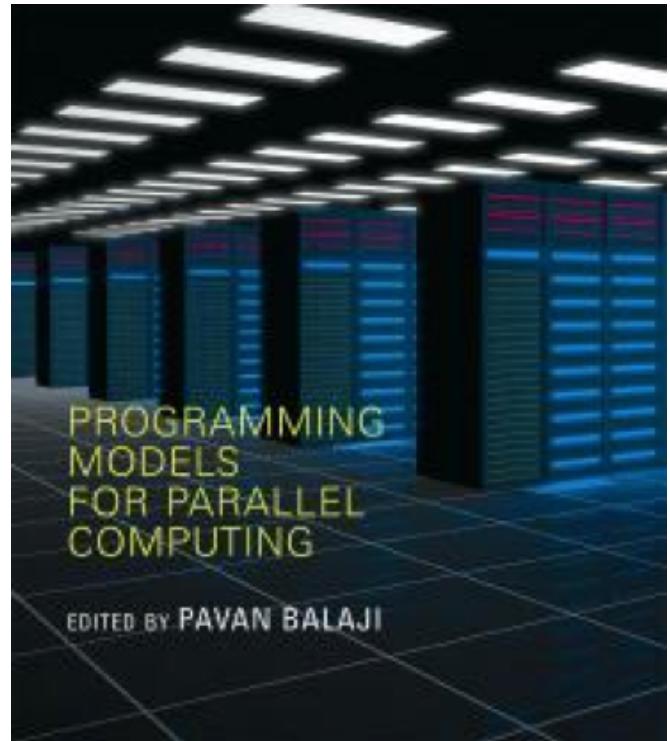
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Suggested Reading

Chapel chapter from *Programming Models for Parallel Computing*

- a detailed overview of Chapel's history, motivating themes, features
- published by MIT Press, November 2015
- edited by Pavan Balaji (Argonne)
- chapter is now also available [online](#)



Other Chapel papers/publications available at <http://chapel.cray.com/papers.html>

Chapel Blog Articles

[Chapel: Productive Parallel Programming](#), Cray Blog, May 2013.

- *a short-and-sweet introduction to Chapel*

[Chapel Springs into a Summer of Code](#), Cray Blog, April 2016.

- *coverage of recent events*

[Six Ways to Say “Hello” in Chapel](#) (parts [1](#), [2](#), [3](#)), Cray Blog, Sep-Oct 2015.

- *a series of articles illustrating the basics of parallelism and locality in Chapel*

[Why Chapel?](#) (parts [1](#), [2](#), [3](#)), Cray Blog, Jun-Oct 2014.

- *a series of articles answering common questions about why we are pursuing Chapel in spite of the inherent challenges*

[\[Ten\] Myths About Scalable Programming Languages](#), [IEEE TCSC Blog](#) (index available on chapel.cray.com “blog articles” page), Apr-Nov 2012.

- *a series of technical opinion pieces designed to argue against standard reasons given for not developing high-level parallel languages*



Chapel Mailing Lists

low-traffic / read-only:

`chapel-announce@lists.sourceforge.net`: announcements about Chapel

community lists:

`chapel-users@lists.sourceforge.net`: user-oriented discussion list

`chapel-developers@lists.sourceforge.net`: developer discussions

`chapel-education@lists.sourceforge.net`: educator discussions

`chapel-bugs@lists.sourceforge.net`: public bug forum

(subscribe at SourceForge: <http://sourceforge.net/p/chapel/mailman/>)

To mail the Cray team:

`chapel_info@cray.com`: contact the team at Cray

`chapel_bugs@cray.com`: for reporting non-public bugs

or use IRC (#chapel on chat.freenode.net) or StackOverflow





Chapel: Productive Parallel Programming at Scale

Questions?



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