Homework 3

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1 Problem 1

1.1 (a)

Due to the chain rule, as well as f(x) is a nonnegative combination

$$\partial f(x) = A^{\top} \partial ||Ax - b||_2 + \partial ||x||_2$$

Obviously, when $x \neq 0$, $\partial ||x|| = \frac{x}{||x||_2}$. When x = 0, by definition, we have

$$||y||_2 \ge g^\top y, \forall y \in \mathbb{R}^n$$

which leads to $||g||_2 \le 1$. Hence, $\partial ||x||_2 = \{g|||g||_2 \le 1\}$. Hence,

$$\partial f(x) = \begin{cases} A^{\top}(Ax - b) / \|Ax - b\|_2 + x / \|x\|_2, & \text{if } x \neq 0, Ax - b \neq 0 \\ \{-A^{\top}b / \|b\|_2 + g | \|g\|_2 \leq 1\}, & \text{if } x = 0, b \neq 0 \\ \{A^{\top}g + x / \|x\|_2 | \|g\|_2 \leq 1\}, & \text{if } x \neq 0, Ax - b = 0 \\ \{A^{\top}g_1 + g_2 | \|g_1\|_2, \|g_2\|_2 \leq 1\}, & \text{if } x = 0, b = 0 \end{cases}$$

1.2 (b)

We first compute $\partial \|z\|_{\infty}$. Since $\|z\|_{\infty} = \sup_{\|w\|_1} z \top w$, by the **Danskin-Bertseka's Theorem for subdifferentials**, since $\{w: \|w\|_1 \le 1\}$ is compact, and $w^\top z$ is convex in z for every w, we have

$$\partial ||z||_{\infty} = \operatorname{conv}\{w : ||w||_{1} \le 1, z^{\top}w = ||z||_{\infty}\}$$
$$= \{w : ||w||_{1} \le 1, z^{\top}w = ||z||_{\infty}\}$$

For given x, we minimize $||Ay - x||_{\infty}$ w.r.t. y, which is a LP problem. Then

$$\partial_{x,y} \|Ay - x\|_{\infty} = \partial_{x,y} \| \begin{pmatrix} -I & A \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} \|_{\infty} = \begin{pmatrix} -I \\ A^{\top} \end{pmatrix} = \{ \begin{pmatrix} -I \\ A^{\top} \end{pmatrix} w : \|w\|_{1} \le 1, w^{\top} (Ay - x) = \|Ay - x\|_{\infty} \}$$

Since we set $A^{\top}w = 0$, hence the subgradient is

$$\{w: A^{\top}w = 0, \|w\|_1 \le 1, w^{\top}\hat{x} = \|A\hat{y} - \hat{x}\|_{\infty}\}$$

at point \hat{x} .

Now we need to prove that such w exist. We need first find \hat{y} to minimize $||Ay - x||_{\infty}$. If $||\hat{x} - A\hat{y}||_{\infty} = 0$, we set w = 0. Otherwise, the problem can be reformulated as

$$\min_{\substack{t,y\\ \text{s.t.}}} t$$
s.t. $-t\mathbb{1} \leq Ay - \hat{x} \leq t\mathbb{1}$ (1.1)

whose Lagrangian can be written as

$$L(t, y, a, b) = t + a^{\mathsf{T}} (Ay - \hat{x} - t\mathbb{1}) + b^{\mathsf{T}} (-t\mathbb{1} - Ay + \hat{x})$$
(1.2)

where $a, b \succeq 0$.

The dual problem is

$$\min_{a,b} \quad \hat{x}^{\top}(a-b) \tag{1.3a}$$

s.t.
$$A^{\top}(a-b) = 0$$
, $\mathbb{1}^{\top}(a+b) = 1$ (1.3b)

$$a, b \succeq 0$$
 (1.3c)

Assuming the optimal solution is t^* , y^* , a^* , b^* , then the KKT condition amounts to

$$A^{\top}(a^{\star} - b^{\star}) = 0 \tag{1.4a}$$

$$\mathbb{1}^{\top}(a^{\star} + b^{\star}) = 1 \tag{1.4b}$$

$$(a^{\star})^{\top} (Ay^{\star} - \hat{x} - t^{\star} \mathbb{1}) = 0 \tag{1.4c}$$

$$(b^*)^{\top} (-t^* \mathbb{1} - Ay^* + \hat{x}) = 0$$
 (1.4d)

We claim that $w = a^* - b^*$ satisfy 1.2. According to 1.4a, $A^T w = 0$; and from 1.4b, 1.4c and 1.4d, we have

$$w^{\top} \hat{x} = (a^{\star} - b^{\star})^{\top} (Ay^{\star} - \hat{x}) = (a^{\star} + b^{\star})^{\top} \mathbb{1}t^{\star} = t^{\star}$$

Besides, if $a_i^{\star} \neq 0$, then from 1.4c we have $(Ay^{\star} - \hat{x})_i = t$, combined with 1.4d and t > 0, we have $b_i^{\star} = 0$. Similarly, if $b_i^{\star} \neq 0$, then $a_i^{\star} = 0$. Hence,

$$||w||_1 = ||a^* - b^*||_1 = ||a^*||_1 + ||b^*||_1 = \mathbb{1}^\top (a^* + b^*) = 1$$

Therefore, the existence is proved. Form the proof, we can easily find that by solving the dual 1.3, $\partial f(\hat{x})$ is naturally obtained by setting $w = a^* - b^*$.

2 Problem 2

2.1 (a)

Define

$$\delta(x) = \begin{cases} 0, & \|x\|_{\infty} \le 1\\ +\infty, & \text{otherwise} \end{cases}$$

Hence, $f(x) = ||x||_1 + \delta(x)$. If we optimize elementwisely, i.e.

$$\min_{u_i} |u_i| + \delta(u_i) + \frac{1}{2}|u_i - x_i|^2$$

which means

$$\min_{u_i} |u_i| + \frac{1}{2}|u_i - x_i|^2
\text{s.t.} -1 \le u_i \le 1$$

Taking the derivative w.r.t. u_i , we get

$$u_i - x_i + \partial |u_i| = 0$$

The solution is $u_i = \operatorname{sign}(x_i) \max(|x_i| - 1, 0)$. Considering $|u_i| \le 1$, since the problem is quadratic and have positive leading coefficient, the solution should be modified to $u_i = \operatorname{sign}(x_i) \min(\max(|x_i| - 1, 0), 1)$.

Hence

$$\operatorname{prox}_f(x)_i = \operatorname{sign}(x_i) \min(\max(|x_i| - 1, 0), 1)$$

2.2 (b)

Take the subgradient, we have

$$\partial f(u) + u - x = 0$$

Assume e_i is the vector whose *i*-th item is 1 and the others are 0, according to the relation of taking maximum and subgradient, we have

$$\partial f(u) = \operatorname{conv}\{e_i|u_i = f(u)\}$$

where the equality holds since each x_k is linear (differentiable). Hence, we have

$$u_j = \begin{cases} x_j, & u_j \neq f(u) \\ x_j - c_j & u_j = f(u) \end{cases}$$

where c_j satisfies $\sum_{j:u_j=f(u)} c_j = 1$.

Assuming the solution of $\sum_{k} \max(x_k - t, 0) = 1$ w.r.t. $t \in \mathbb{R}$ is \hat{t} . Clearly, \hat{t} exist and is unique. This equation can be solved by bisecting t with the initial interval $[\min_i x_i - \frac{1}{n}, \max_i x_i].$

We thus have $prox_f(x)_i = min(x_i, \hat{t})$.

2.3 (c)

Since

$$prox_f(x) = \min_{u} \left\{ ||Au - b||_1 + \frac{1}{2} ||u - x||_2^2 \right\}$$

which can be formulated as

$$\min_{u,z} \quad \|z\|_1 + \frac{1}{2} \|u - x\|_2^2 \tag{2.1a}$$

$$s.t. \quad z = Au - b \tag{2.1b}$$

Define the optimal solution of 2.2 by (\hat{z}, \hat{u}) . The existence of such solution follows from the fact that $||x||_1$ is convex and positive. Clearly $\hat{u} = \operatorname{prox}_f(x)$, fixing $z = \hat{z}$, we have that \hat{u} is the optimal solution of

$$\min_{u,z} \quad \frac{1}{2} ||u - x||_2^2$$
s.t.
$$Au = \hat{z} + b$$

Thanks to the strong duality, there exsit w such that

$$\hat{u} \in \arg\min_{u} \left\{ \frac{1}{2} \|u - x\|_{2}^{2} + \langle y, Au - \hat{z} - b \rangle \right\}$$
 (2.3a)

$$A\hat{u} = \hat{z} + b \tag{2.3b}$$

By 2.3a, we have

$$\hat{u} = x - A^{\mathsf{T}} y \tag{2.4}$$

Substituting it into 2.3b, we have

$$A(x - A^{\top}y) = \hat{z} + b$$

Since $AA^{\top} = D$, and D is a diagonal matrix whose diagonal elements are positive,

$$y = D^{-1}(Ax - \hat{z} - b)$$

which combined with 2.4, we have

$$\hat{u} = x + A^{\mathsf{T}} D^{-1} (\hat{z} + b - Ax) \tag{2.5}$$

Substituting it into 2.2, we have

$$\hat{z} = \arg\min_{z} \left\{ \|z\|_{1} + \frac{1}{2} \|A^{\top} D^{-1} (z + b - Ax)\|_{2}^{2} \right\}$$
 (2.6)

Taking the derivative w.r.t. z, assuming $d_i = D_{ii}$, we have

$$\partial ||z_i||_1 + (z+b-Ax)_i/d_i = 0$$

which leads to

$$z_i = \operatorname{sign}(Ax - b)_i \max(|(Ax - b)|_i - d_i, 0)$$

Substituting it into 2.5

$$\operatorname{prox}_f(x) = \hat{u} = x + A^{\top} D^{-1} \left(b - Ax + \operatorname{sign}(Ax - b) \odot \max(|Ax - b| - \operatorname{diag}(D), 0) \right)$$

Problem 3 3

Taking the derivative w.r.t. y, we have

$$y - w + \frac{1}{(1 - \alpha)\sigma n} \sum_{i=1}^{n} \partial (y_i - t)_{+}$$

where

$$\partial(x)_{+} = \begin{cases} 1, & x > 0 \\ [0, 1], & x = 0 \\ o, & x < 0 \end{cases}$$

WLOG, assuming $w_1 \leq w_2 \leq \cdots, w_n$, and $w_0 = -\infty, w_{n+1} = \infty$. For each t, we can readily find u, v, such that $w_u \leq t < w_{u+1}, w_v \leq t + \frac{1}{(1-\alpha)\sigma n} < w_{v+1}$. Then, for $i \leq u$, $y_i = w_i$; for $i \geq v+1$, $y_i = w_i - \frac{1}{(1-\alpha)\sigma n} > t$; for $u+1 \leq i \leq v$, $y_i = w_i$

must be t. Substituting it into the original problem, we have

$$\min_{t} \quad t + \sum_{i=v+1}^{n} \frac{w_i - t}{(1 - \alpha)n} + \sum_{i=u+1}^{v} (w_i - t)^2 + \left(\frac{1}{(1 - \alpha)\sigma n}\right)^2 (n - v)$$

i.e. given u, v

$$\min_{t} (v - u)t^{2} + \left(1 - \frac{n - v}{(1 - \alpha)n} - 2\sum_{i=u+1}^{v} w_{i}\right)t$$

Keeping u, v, it is a quadratic minimization. Denote $w'_i = w_i - \frac{1}{(1-\alpha)\sigma n}$. Then w_i, w'_i divide the real line into at most 2n + 1 interval. Each interval corresponds to a set of (u,v). For each set of (u,v), we can minimize t by performing a quadratic function minimization.

If we precompute and store $\sum_{i=1}^{m} w_i$ for $m = 1, 2, \dots, n$, the algorithm's complexity can be O(n).

Problem 4 4

Assuming $X = S\Lambda S^{T}$, where Λ is diagonal, since

$$\log \det(X) = \log \det(S\Lambda S^{\top}) = \log \det \Lambda$$

 $\log \det(X)$ is determined by the eigenvalues of X. A standard result shows that the optimal solution can be written as $U^* = SDS^T$, where D is diagonal. Denote $d_i =$ $D_{ii}, \lambda_i = \Lambda_{ii}$. We have

$$f(U) + \frac{1}{2} ||U - X||_F^2$$

$$= f(D) + \frac{1}{2} ||\Lambda - D||_F^2$$

$$= \sum_{i=1}^n \left(-\log d_i + \frac{1}{2} |\lambda_i - d_i|^2 \right)$$

For each i, we optimize separately. Note that since $X, U \in S_{++}^n$, $\lambda_i, d_i > 0$, taking the derivative w.r.t. d_i and setting it to 0, we have

$$-\frac{1}{d_i} + d_i - \lambda_i = 0$$

which leads to $d_i = \frac{\lambda_i + \sqrt{\lambda_i^2 + 4}}{2}$. (omit the negative root) If we denote $g(X) = Sg(\Lambda)S^{\top}$ and operations are taken element-wise on Λ . Define $g(x) = \frac{x + \sqrt{x^2 + 4}}{2}, x \in \mathbb{R}$, then

$$\operatorname{prox}_f(X) = g(X)$$