

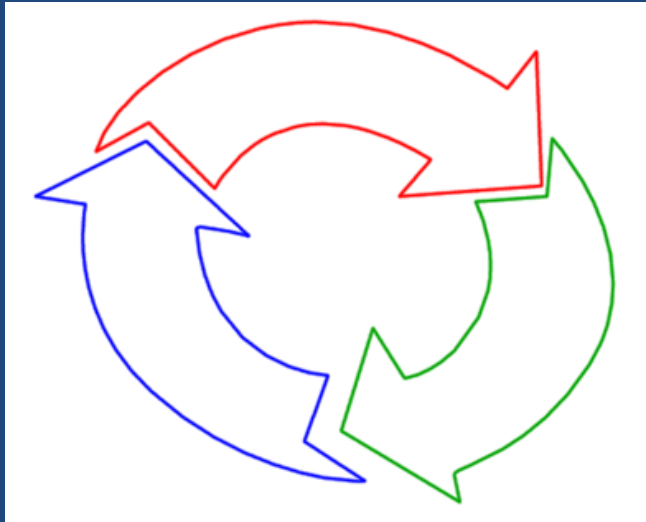
DS

Circular Linked Lists
Circular Arrays
Stacks and Queues

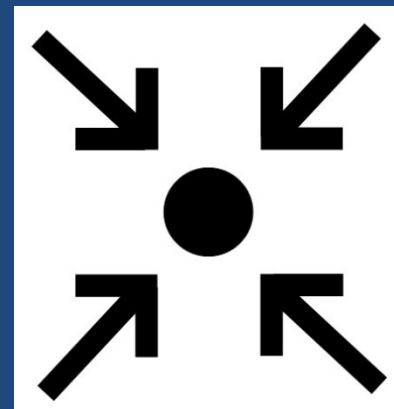
Circular Linked Lists

Circular Linked Lists

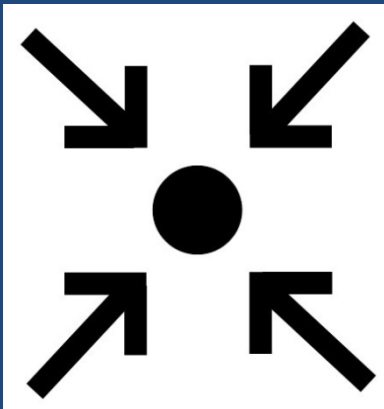
The list
never
ends



The final *next* pointer
points to
the first item in the list



If the list is doubly-linked, the first item's *prev* pointer points to the last item

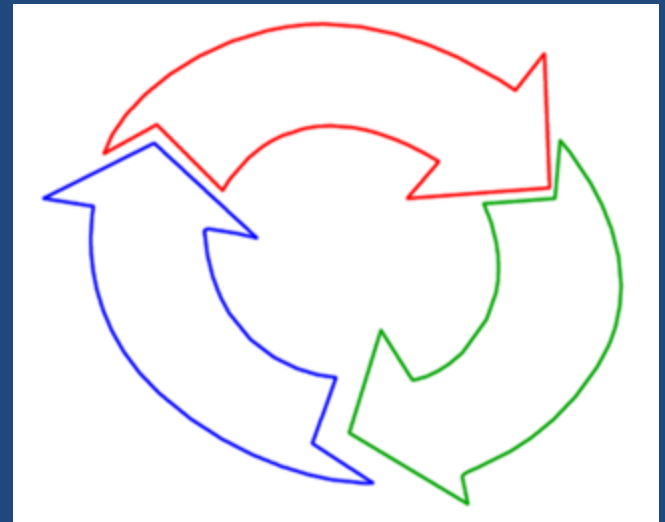


There is no *tail* pointer



Traversal

Traversal loops back
to the start



Compare
the next pointer
of the current node
to
the head pointer



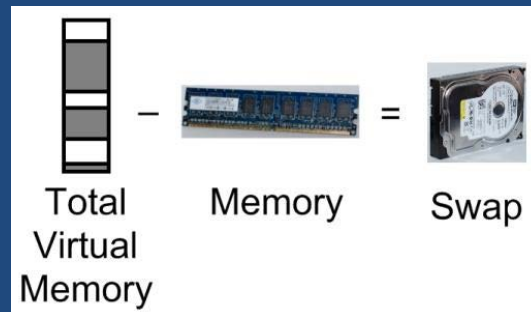
You can't just compare the
current node to the head
pointer or you'll never start!



Example of Usefulness of a Circular List

Determine which memory page is
least recently used, so that it can
be swapped out

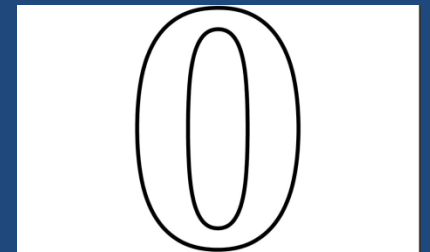
(You took about this in OSF)





Uses *reference number*

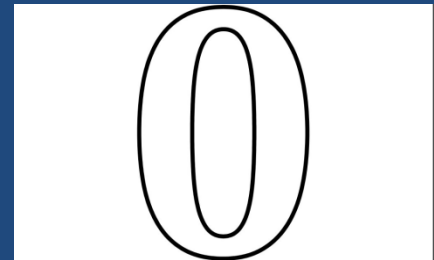
Set to 1 when the page is accessed
and 0 otherwise



(Internal to the memory manager)

Commentary: Mechanism of Swapping

When swapping out a page,
the memory manager
searches
for a page with
a *reference number of 0*



While it's searching, it resets the
reference number of the pages
it found from 1 to 0



Commentary: Why Circular?

Will always find
a page to swap out
because it'll either find
one with a value of 0
or
it will find where it started
(which was reset from 1 to 0)

Commentary: Optimization of the Search

It'll then reset the head of the list to point to the location we've just swapped out



The actual location
of the start of the list
isn't important
since we're just using it
to try to find a swappable page



Search from here next time
because
we didn't necessarily
complete the circuit this time



Circular Arrays

Circular Arrays

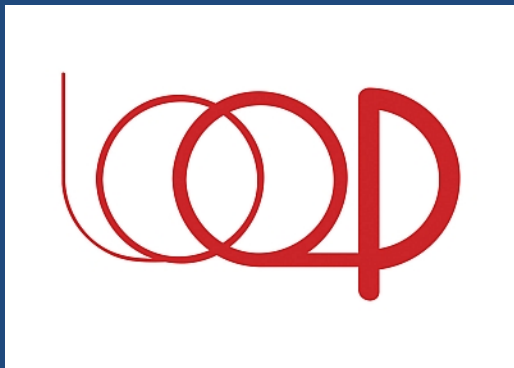
Arrays in which
the “first” element in the array
is not necessarily
the physically-first one
and the “last” element
is not necessarily
the physically-last one



Have
a head and tail **index**
(not pointer)
that indicate where
the “first”
and “last” elements
actually are



When you reach
the physical end of the array



in traversal,
loop around
to the physical start
of the array

Circular Array Terminology

The “physical-first” element of a circular array called *fred* would be *fred[0]*

The “logical-first” element of a circular array is the first element that would be processed

Similarly:

- The “physical-last” element of a circular array of size 10 called *fred* would be *fred[9]*
- The “logical-last” element of a circular array is the last element that would be processed

More Circular Array Terminology

The “head” of a circular array is the index of the “logical-first” element

The “tail” of a circular array is the index of one **past** the “logical-last” element

You could do these with pointers if you want, as well

Circular Array: Full vs. Empty (one definition)

A circular array is full if the tail is one position before the head

This must account for wrapping around the physical end of the array

A circular array is empty if the head and tail are equal

This implies that we cannot use up all of the elements in the array

This is one method of
implementation

You also could keep a separate variable indicating whether the list is full, empty, or neither

Examples of Full and Empty in this Circular Array

Let's say we have an array of size 5 (indices 0 to 4)

“X” will indicate a used element

Index	0	1	2	3	4
Value					

Example #1

Index	0	1	2	3	4
Value	X	X	X		

Head = 0

Tail = 3

This is not empty and not full

Example #2

Index	0	1	2	3	4
Value	X	X	X	X	

head = 0

tail = 4

This is full

Example #3

Index	0	1	2	3	4
Value	X		X	X	X

head = 2

tail = 1

This is full

Example #4

Index	0	1	2	3	4
Value					

head = 3 (or any valid index)

tail = 3

This is empty

Example #5

Index	0	1	2	3	4
Value		X	X	X	X

head = 1

tail = 0

This is full

Stacks and Queues

Reading

A Book on C, Sections 10.5 to
10.7

Why?

Queues in real life:

- e.g. lineups while shopping or filling the car with gas
- e.g. sending an application in



Stacks in real life are less obvious

- e.g. an “IN box” that stuff piles up in
- e.g. a narrow drawer that you throw stuff into



Sorting

These are sorted lists in which the order of the list is time-based



The sorting is implicit,
in that
it happens through
the act of adding to the list



There is the potential
of delayed processing
of the items



	GATE #	STATUS
ON	A23	DELAYED
ES	C72	DELAYED
	B34	DELAYED
	A14	DELAYED
	C89	DELAYED
	G12	DELAYED
	C5	DELAYED
	D13	DELAYED
0	A4	DELAYED
	B22	DELAYED
	A33	DELAYED

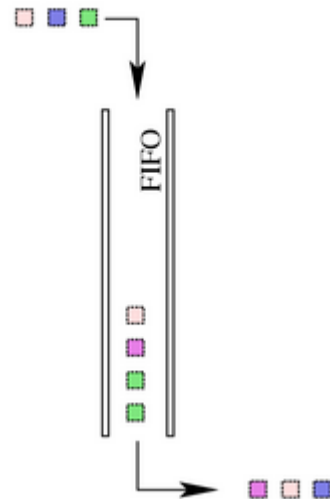
So there may be more than
one item in the list

FIFO vs. LIFO

In a queue, the **first** item into the queue is the first item processed from the queue

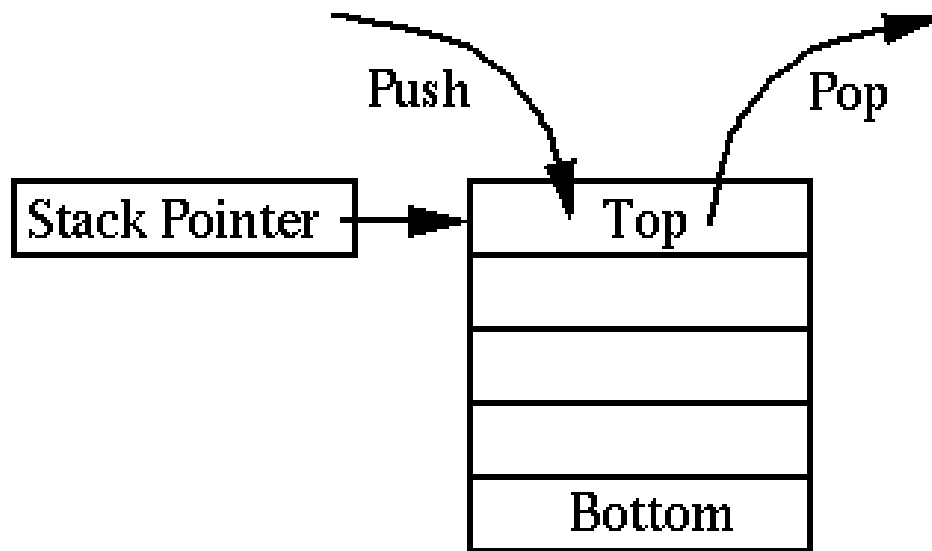
- This is called “FIFO” for “First In, First Out”

First-in First-out (FIFO)

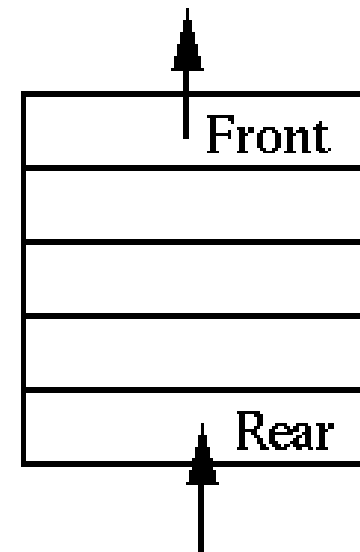


In a stack, the **last** item into the stack is the first item processed from the stack

- This is called “LIFO” for “Last In, First Out”



A LIFO Stack



A FIFO queue

Stack Terminology

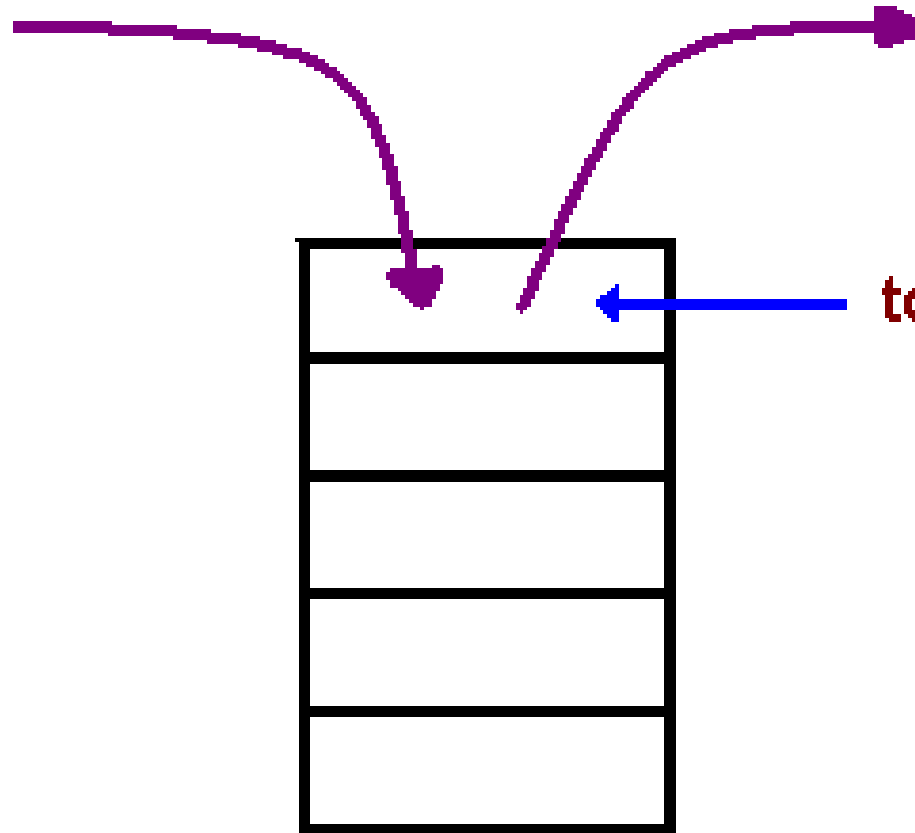
Putting an item on the top of the stack is called “pushing” it

Taking an item off the top of the stack is called “popping” it

push

pop

top



Viewing an item on the stack is called “peeking” at it



Queue Terminology

If you put an item at the end of a queue, you “enqueue” it

If you take an item from the start of a queue, you “dequeue” it

Looking at an item in the queue
without processing it is called
“peeking” at it
(same as stack)

Implementation

So far, stacks and queues have been referred to as “lists” but **not** “linked lists”

You **can** use a linked list to implement them but you don’t **have** to

Other alternatives if imposing a
maximum size is OK:

arrays

circular arrays (queues only)

Actions in a Queue

Queue items are always

added to the tail

and

processed from the head

Processing a queue element
always involves
deleting it

Typically, there is no need for:

- deletion of items except at the head
- adding of items except at the tail
 - explicit sorting

Actions in a Stack

Actions are restricted in a stack,
similar to in a queue

Difference: Items are always
added to the tail and
processed from the **tail**

So, you only have to support
deletion from the **tail**

Stack Overflow

If using an array to implement a stack, you have to prevent stack overflow

Don't add past the maximum size of the array

Summary

Circular Linked Lists

Circular Arrays

Stacks

Queues

all have their uses