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**UNIVERSITÄT  
BERN**

# **3D Metric Fields**

## **A Novel Approach to a New Idea**

### **Bachelor Thesis**

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# Abstract

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# Chapter 1

## Introduction

In many applications of computer graphics or engineering, the geometry of objects has to be described in some way. Finite elements encode various information such as physical properties to be used in e.g. stress simulations. While it is possible for simple objects such as spheres or planes to describe them with analytical formulas, this is unfeasible for arbitrary objects of any shape. To describe any geometric object, it is discretized and encoded as a mesh. Volumetric meshes are used to encode the surface and the interior of an object. The most common types of volumetric meshes are tetrahedral and hexahedral meshes. Hexahedral meshes, that is a decomposition into cube-like elements, are usually preferred due to their superior numerical accuracy and faster compute time. However, automatically generating high-quality meshes remains a hard challenge in computer graphics. A promising research direction is through the use of frame fields [11].

A *frame field* can be seen as a generalization of vector fields. They prescribe a *frame* to each point on an object, that is, three linearly-independent vectors. Each frame locally represents the orientation and deformation of a cube. Consequently, frame fields must be so called boundary-aligned: At the boundary of an object, one vector of the frame must be aligned with the surface normal as hexahedral elements (*short: hex element*) should align at the boundary. Then, the objective is to find the “best” frame field. Usually, “best” is measured in the smoothness of the frame field. Smoothness is measured with the Dirichlet energy  $\int_{\Omega} \|\nabla \phi\|^2 dx$ , where  $\phi$  is a frame, and we measure how much these frames twist and rotate within the object.

To allow for non-uniform and anisotropic hex elements, we introduce *metric fields*. Optimizing the frame field in this new metric (instead of the usual euclidean metric) then results in elements that are uniform and isotropic in this new metric. To this end, the expression for the Dirichlet energy must be changed. We will show how to calculate a special rotation matrix  $R$ , which describes how frames are aligned in the new metric.

In this thesis, we rederive how to pull back the metric with the rotation coefficient  $R$  such that we can measure the Dirichlet energy in the usual cartesian coordinates as described in *Metric-Driven 3D Frame Field Generation*[4]. The necessary mathematical background is introduced in Section 2. In Section 3, by focusing on integrability of the frame field, we derive the connection  $\omega$  under our new metric field and how to evaluate it. Section 4 discusses how the metric field is piecewise-linearly discretized as a tetrahedral mesh and how the rotation coefficient  $R$  to align frames in the metric is calculated within a specific tetrahedron. Because in general,  $R$  is needed between arbitrary points in the metric field, a robust tetrahedron finder algorithm is presented in Section 5. We briefly explain how the actual frame field optimization works in Section ???. The whole frame field generation is tested with some experiments in Section ??. We conclude in Section 6.





## Chapter 2

# Mathematical Background

We will make heavy use of differential geometry in the following chapters. To get us all on the same page, I introduce the basic concepts what we will use, but I refrain from giving any proofs. I will give definitions only as far as we need it. These definitions will by no means be exhaustive. The following is an incomplete summary of what we need presented in *Introduction to smooth manifolds*[?].

**Manifold** A manifold  $\mathcal{M}$  is a space that locally looks like Euclidean space. More exactly, a  $n$ -manifold is a topological space, where each point on the manifold has an open neighborhood that is locally homeomorphic to an open subset of Euclidean space  $\mathbb{R}^n$ . A manifold can be equipped with additional structure. For example, we can work on *smooth manifolds*. In simple terms, a manifold is *smooth* if it is similar enough to  $\mathbb{R}^n$  that we can do Calculus like differentiation or integration on it. For this, each point on the manifold must be locally *diffeomorphic* to an open subset of  $\mathbb{R}^n$  space.

**Tangent space, Tangent bundle** There are many equivalent definitions for the tangent space. One definition is for each point  $p$  in the manifold  $\mathcal{M}$ , the tangent space  $T_p\mathcal{M}$  consists of  $\gamma'(0)$  for all differentiable paths  $\gamma : (-\varepsilon, \varepsilon) \rightarrow \mathcal{M}$  with  $p = \gamma(0)$ . The tangent space is a vector space which has the same dimension as its manifold, which is 3 in our case. These tangent spaces can be “glued” together to form the *tangent bundle*  $T\mathcal{M} = \sqcup_{p \in \mathcal{M}} T_p\mathcal{M}$ , which itself is a manifold of dimension  $2n$ . An element of  $T\mathcal{M}$  can be written as  $(p, v)$  with  $p \in \mathcal{M}$  and  $v \in T_p\mathcal{M}$ . This admits a natural projection  $\pi : T\mathcal{M} \rightarrow \mathcal{M}$ , which sends each vector  $v \in T_p\mathcal{M}$  to the point  $p$  where it is tangent:  $\pi(p, v) = p$ . A *section*  $\sigma : \mathcal{M} \rightarrow T\mathcal{M}$  is a continuous map, with  $\pi \circ \sigma = \text{Id}_{\mathcal{M}}$ . Sections of  $T\mathcal{M}$  are tangent vector fields on  $\mathcal{M}$ .

**Cotangent space, Cotangent bundle** The dual space  $V^*$  of a vector space  $V$  consists of all linear maps  $\omega : V \rightarrow \mathbb{R}$ . We call these functionals *covectors* on  $V$ .  $V^*$  is itself a vector space, with the same dimension as  $V$  and operations like addition and scalar multiplication can be performed on its elements. Any element in a vector space can be expressed as a finite linear combination of its basis. This basis is called the *dual basis*. Thus, we call the dual space of the vector space  $T_p\mathcal{M}$  its *cotangent space*, denoted by  $T_p^*\mathcal{M}$ . As before, the disjoint union of  $T_p^*\mathcal{M}$  forms the *cotangent bundle*:  $T^*\mathcal{M} = \sqcup_{p \in \mathcal{M}} T_p^*\mathcal{M}$ . Defined analogously from above, sections  $\sigma$  on  $T^*\mathcal{M}$  define *covector fields* or *1-forms*.

**Tensors** Before we can introduce differential forms in the next paragraph, we need to go a little bit into *tensors*. In simple words, tensors are real-valued, multilinear functions. A map  $F : V_1 \times \dots \times V_k \rightarrow W$  is multilinear, if  $F$  is linear in each component. For example, the dot product in  $\mathbb{R}^n$  is a tensor. It takes two vectors and is linear in each component - bilinear. Another example is the *Tensor Product of Covectors*: Let  $V$  be a vector space and take two covectors  $\omega, \eta \in V^*$ . Define the new function  $\omega \otimes \eta : V \times V \rightarrow \mathbb{R}$  by  $\omega \otimes \eta(v_1, v_2) = \omega(v_1)\eta(v_2)$ . It is multilinear, because  $\omega$  and  $\eta$  are linear. We look at a special class of tensors, the *alternating tensors*. A tensor is alternating, if it changes sign whenever two arguments are interchanged, i.e.  $\omega(v_1, v_2) = -\omega(v_2, v_1)$ . A covariant tensor field over a manifold defines a covariant

tensor at each point on the manifold, covariant because the tensor is over the cotangent space  $T_p^*\mathcal{M}$ . An alternating tensor field is called a *differential form*.

**Differential Forms, Exterior Derivative** Recall that a section from  $T^*\mathcal{M}$  is called a differential 1-form, or just 1-form. Define the *wedge product* (or *exterior product*) between two 1-forms:

$$(\omega \wedge \eta)_p = \omega_p \wedge \eta_p$$

Notice the similarity to the *Tensor Product of Covectors*: We get a new map, (a 2-form):

$$\omega \wedge \eta : T\mathcal{M} \times T\mathcal{M} \rightarrow \mathbb{R}$$

The wedge product is antisymmetric, therefore  $\omega \wedge \eta = -\eta \wedge \omega$  for 1-forms  $\omega$  and  $\eta$ . We denote by  $\Omega^k(\mathcal{M})$  the space of differential  $k$ -forms on  $\mathcal{M}$ . There is a natural differential operator  $d$  on differential forms we call *exterior derivative*. It maps  $k$ -forms to  $(k+1)$ -forms, i.e.  $d : \Omega^k(\mathcal{M}) \rightarrow \Omega^{k+1}(\mathcal{M})$  and has the the following properties:

- $df$  is the ordinary differential of a smooth function  $f$ . Smooth functions are 0-forms.
- $d(d\alpha) = 0$
- $d(\alpha \wedge \beta) = (d\alpha \wedge \beta) + (-1)^k(\alpha \wedge d\beta)$  for a  $k$ -form  $\alpha$ . (Leibnitz Rule)

In section 3 we will need what  $d\omega$  is for some smooth 1-form  $\omega$ , the calculation will be done there. For now, just note that any arbitrary smooth 1-form can be written as  $\omega = Fdx + Gdy + Hdz$  for some appropriate smooth functions  $F, G, H$ .

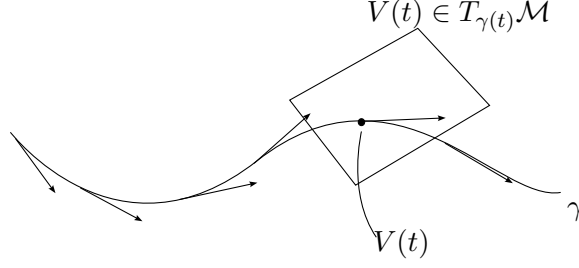
**Riemannian metric,  $g$ -orthonormality** Inner products are examples of symmetric tensors. They allow us to define lengths and angles between vectors. We can apply this idea to manifolds. A Riemannian metric  $g$  is a symmetric positive-definite tensor field at each point. If  $\mathcal{M}$  is a manifold, the pair  $(\mathcal{M}, g)$  is called a *Riemannian manifold*. Let  $g$  be the Riemannian metric on  $\mathcal{M}$  and  $p \in \mathcal{M}$ , then  $g_p$  is an inner product on  $T_p\mathcal{M}$ . We write  $\langle \cdot, \cdot \rangle_g$  to denote this inner product. Any Riemannian metric can be written as positive-definite symmetric matrix, which allows for this simple form:  $\langle v, w \rangle_g = v^\top g w$ .

Such a new metric allows for the definition of  *$g$ -orthonormality*: A basis  $[e_1, e_2, e_3]$  of  $T_p\mathcal{M}$  is  *$g$ -orthonormal* if  $\langle e_i, e_j \rangle_g = \delta_{ij}$ .

**Connection, Covariant derivative, Parallel transport** A connection defines how two different tangent spaces are connected to each other, such that tangent vector fields can be differentiated. There is an infinite amount of connections on a manifold. An *affine connection*  $\nabla$  is a bilinear map that takes two tangent vector fields  $X, Y$  and maps it to a new tangent vector field  $\nabla_X Y$  on  $\mathcal{M}$  such that

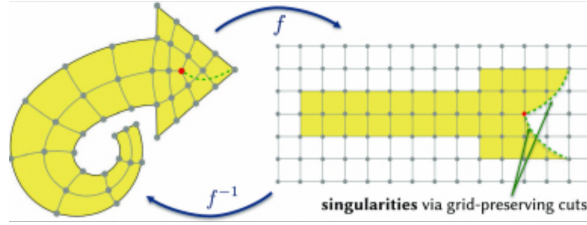
- $\nabla_{fX} Y = f \nabla_X Y$ , where  $f \in C^\infty(\mathcal{M}, \mathbb{R})$
- $\nabla_X(fY) = f \nabla_X Y + (Xf)Y$  for  $f \in C^\infty(\mathcal{M}, \mathbb{R})$ , it satisfies the Leibnitz rule in the second variable

We call  $\nabla_X Y$  the *covariant derivative of  $Y$  in the direction of  $X$* . A connection  $\nabla$  defines the parallel transport of a vector along a curve. Given a curve  $\gamma : [0, 1] \rightarrow \mathcal{M}$  and a vector  $v_0 \in T_{\gamma(0)}\mathcal{M}$ , there exists a unique parallel vector field  $V : [0, 1] \rightarrow T\mathcal{M}$  along  $\gamma$  such that  $V(0) = v_0$  [6]. Recall: a vector field  $V$  along  $\gamma$  means  $\pi(V(t)) = \gamma(t)$ . The uniqueness and existence is a consequence of the vector field  $V$  being the solution of  $\nabla_{\dot{\gamma}(t)} V(t) = 0$  defining a linear ordinary differential equation with initial condition  $V(0) = v_0$ . This vector field  $V(t)$  is called the *parallel transport* of  $v_0$  along  $\gamma$ . It is “parallel” in the sense that the transported vector does not change within the tangent space. See figure 2.1 for an illustration in 2D.



**Figure 2.1.** A parallel vector field  $V(t)$  along a curve  $\gamma$ . Each  $V(t) \in T_{\gamma(t)}\mathcal{M}$  and  $\nabla_{\dot{\gamma}(t)}V(t) = 0$ , so each vector that is drawn is parallel to each other.

**Frame field, vector field, integrability** Recall from the introduction that we are trying to find a hexahedral mesh for some geometric object. In 2D, integer-grid maps have proven to be reliable in generating the analogue of hexahedral meshes, quad meshes [2]. The idea is to use the same idea but in three dimensions. The objective is to find a parametrization  $\phi : \mathcal{M} \rightarrow \mathbb{R}^3$ , which embeds the 3-dimensional object onto an 3-dimensional voxel grid, where the inverse  $\phi^{-1}$  maps the deformed grid back onto the object to recover a shape-aligned hexahedral mesh, see figure 2.2. This is called an integer-grid map



**Figure 2.2.** Idea of quadrangulation with integer-grid maps in 2D

and by minimizing distortion and satisfying boundary alignment, a hexahedral mesh could be extracted. However, this is a hard mixed-integer and non-convex optimization problem for which no current optimization technique is available that would result in an acceptable solution. By taking the Jacobian of the parametrization,  $\nabla\phi$ , we get a mapping  $\nabla\phi : \mathcal{M} \rightarrow \mathbb{R}^{3 \times 3}$  of *frames*, a *frame field*. The idea of frame fields is then to search for an approximation of  $\nabla\phi$ . If  $F : \mathcal{M} \rightarrow \mathbb{R}^{3 \times 3}$  is the approximation of  $\nabla\phi$ , we can solve for  $\phi : \mathcal{M} \rightarrow \mathbb{R}^3$  with

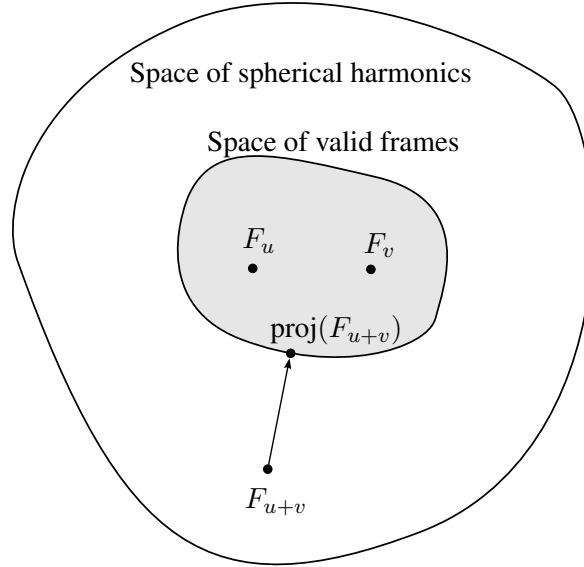
$$\min_{\phi} \int_{\mathcal{M}} \|\nabla\phi F - \text{Id}\|^2$$

where  $\|\cdot\|$  is the Frobenius norm. If this is sufficiently small, then  $F^{-1} \approx \nabla\phi$  and the extracted hexahedral mesh closely follow the integer-grid isolines. We treat a frame  $F$  as a set of 3 linearly independent vectors  $\{F_1, F_2, F_3\}$  which we can collect into a matrix  $F = (F_1, F_2, F_3) \in \mathbb{R}^{3 \times 3}$ . A frame locally represents the edges of a deformed cube. We can think of a frame field as the composition of three vector fields. However, frame fields can contain singularities that the underlying vector fields do not contain [8]. Notice that many of these frames are equivalent. There are  $2^3$  choices of the sign, and  $3!$  possible permutations, which gives  $2^3 \cdot 3! = 48$  equivalent frames. Since we want non-degenerate frames and the same orientation through the grid, the constraint  $\det(F) > 0$  is imposed, which leaves 24 equivalent frames. Equivalence of frames is then defined as

$$F_u \sim F_v \iff \exists R \in \mathcal{O} : F_u = RF_v$$

where  $\mathcal{O}$  is the *chiral cubical symmetry group*[8]. These symmetries make the optimization for frame fields more complicated. To handle these symmetries, we use the spherical harmonics based representation [5]. The idea is to use the polynomial  $x^4 + y^4 + z^4$  to express the frames as rotations. Under the restriction of the polynomial to the sphere, that is  $\mathbb{S}^2 \rightarrow \mathbb{R}$ , the octahedral frames are invariant under the chiral cubical symmetry group. However, the spherical harmonics manifold is 9-dimensional so not

all vectors in the spherical harmonics represent valid frames, since rotations only exhibit 3 degrees of freedom. While optimizing the frame field, some kind of “average” of two frames will be done. This averaging makes sense in the spherical harmonics space, but the result may be outside the valid space of frames. A projection from the spherical harmonics space back to the valid space of frames will be needed, see fig. 2.3 [12].

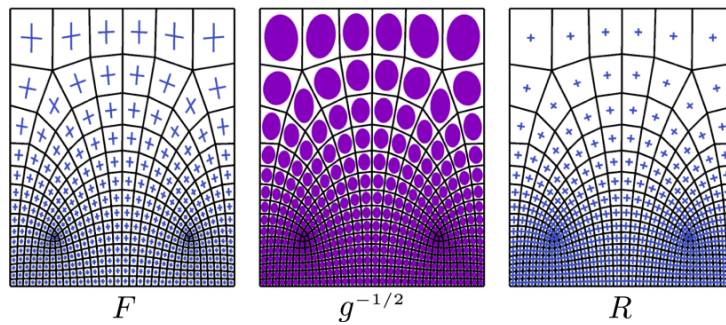


**Figure 2.3.** Combination of two valid frames leads to a resulting frame outside the valid frames. A projection to the nearest valid frame is done.

When a frame  $F$  has orthonormal columns, the resulting hex elements look like unit cubes. By introducing a metric  $g$ , we can control the size and shear of the cubes by relaxing the orthonormality constraint to  $g$ -orthonormality when

$$F^{-1}gF = \text{Id}$$

holds. We can decompose the frame field  $F$  into a rotational part  $R : \mathcal{M} \rightarrow \text{SO}(3)$  and a symmetric metric part  $g^{-1/2}$  (akin to the polar decomposition of linear transformations)[9], that is  $F = g^{-1/2}R$ .



**Figure 2.4.** Factorization of a  $g$ -orthogonal frame field into a symmetric metric part  $g^{-1/2}$  and a rotational part  $R$  (Figure from [4]).

## Chapter 3

### Connection One-Form $\omega$

A vector field  $U$  is integrable, if and only if  $\nabla \times U = 0$ , which means the vector field has vanishing curl everywhere. We can express this more naturally with the language of differential forms: The curl can be written as the exterior derivative  $d$  of a one-form  $\alpha$ . A one-form (more generally, a differential form) is closed, if  $d\alpha = 0$ . Therefore, the local integrability can be expressed as the closedness of a one-form. We want  $F^{-1}$  (TODO: why  $F^{-1}$ ) to be integrable. To achieve local integrability for, it suffices to make  $R$  locally integrable. We can think of a rotation field  $R$  as the composition of 3 vector fields

$$R = \begin{bmatrix} | & | & | \\ R_1 & R_2 & R_3 \\ | & | & | \end{bmatrix}$$

where  $R_i : \mathbb{R}^3 \rightarrow \mathbb{R}^3$  is a vector field. We can therefore construct a vector-valued one-form, given  $p = (x, y, z)^\top$  in Euclidean coordinates

$$\alpha \triangleq F^{-1}dp = R^\top g^{1/2}dp$$

where  $dp = (dx, dy, dz)^\top$  is the common orthonormal one-form basis of  $\Omega^1(\mathcal{M})$ .

$$R \text{ locally integrable} \iff \mathbf{0} = d\alpha$$

Some reformulations yield:

$$\begin{aligned} \mathbf{0} = d\alpha &= d(R^\top g^{1/2}dp) \stackrel{(1)}{=} dR^\top \wedge (g^{1/2}dp) + R^\top d(g^{1/2}dp) \\ &= R^\top (\omega \wedge (g^{1/2}dp) + d(g^{1/2}dp)) \end{aligned}$$

where for (1), the Leibnitz Rule for the exterior derivative is applied, and we define

$$\omega = RdR^\top \in \mathfrak{so}(3).$$

**Remark.**  $\omega$  is an antisymmetric matrix-valued one-form (every element is a one-form).

*Proof.* We differentiate the orthogonality condition of the rotation matrix  $R$  (taking the derivative w.r.t. every element):

$$\begin{aligned} Id &= RR^\top \\ d(Id) &= d(RR^\top) \\ \mathbf{0} &= dRR^\top + RdR^\top \\ \mathbf{0} &= (RdR^\top)^\top + RdR^\top \\ -(RdR^\top)^\top &= RdR^\top \end{aligned}$$

□

Elements of the Lie algebra  $\mathfrak{so}(3)$  can be thought as infinitesimal rotations and  $\omega$  can be used as a connection one-form. The Lie algebra  $\mathfrak{so}(3)$  has manifold structure, but for our purposes, it suffices to know that elements are antisymmetric matrices. To make  $R$  locally integrable, we find  $\omega$  such that the one-form  $\alpha$  is closed ( $d\alpha = 0$ , curl-free), then try to match the  $\omega$  with  $R$ , which can be expressed as

$$\min_{R \in SO(3)} \|RdR^T - \omega\|^2.$$

**Connection evaluation** To find  $\omega$ , we use the above equation

$$\mathbf{0} = R^\top (\omega \wedge (g^{1/2} dp) + d(g^{1/2} dp)) \iff \mathbf{0} = \omega \wedge (g^{1/2} dp) + d(g^{1/2} dp)$$

and reformulate into a linear system. We represent the antisymmetric matrix-valued one-form  $\omega$

$$\omega = \begin{bmatrix} 0 & \omega_{12} & -\omega_{31} \\ -\omega_{12} & 0 & \omega_{23} \\ \omega_{31} & -\omega_{23} & 0 \end{bmatrix}$$

by  $\begin{bmatrix} \omega_{23} & \omega_{31} & \omega_{12} \end{bmatrix} = \begin{bmatrix} dx & dy & dz \end{bmatrix} W$ . We write  $W = [W_1, W_2, W_3]$ ,  $W_i \in \mathbb{R}^3$ . That is,  $W$  is the matrix with the coefficients for the one-forms, e.g.  $W_1 = [(\omega^{23})_1, (\omega^{23})_2, (\omega^{23})_3]^\top$ . Recall, a one-form can be expressed as

$$\omega_{ij} = (\omega^{ij})_1 dx + (\omega^{ij})_2 dy + (\omega^{ij})_3 dz$$

So e.g.

$$\omega_{23} = \begin{bmatrix} dx & dy & dz \end{bmatrix} W_1 = (\omega^{23})_1 dx + (\omega^{23})_2 dy + (\omega^{23})_3 dz$$

We also write  $A = g^{1/2} = [A^1, A^2, A^3]$ . We use the equation  $\mathbf{0} = \omega \wedge (g^{1/2} dp) + d(g^{1/2} dp)$ . Let us start with the first part:

$$\begin{aligned} \omega \wedge g^{1/2} dp &= \begin{bmatrix} 0 & \omega_{12} & -\omega_{31} \\ -\omega_{12} & 0 & \omega_{23} \\ \omega_{31} & -\omega_{23} & 0 \end{bmatrix} \wedge \begin{bmatrix} A_1^1 & A_1^2 & A_1^3 \\ A_2^1 & A_2^2 & A_2^3 \\ A_3^1 & A_3^2 & A_3^3 \end{bmatrix} \begin{bmatrix} dx \\ dy \\ dz \end{bmatrix} \\ &= \begin{bmatrix} +\omega_{12} \wedge (A_2^1 dx + A_2^2 dy + A_2^3 dz) - \omega_{31} \wedge (A_3^1 dx + A_3^2 dy + A_3^3 dz) \\ -\omega_{12} \wedge (A_1^1 dx + A_1^2 dy + A_1^3 dz) + \omega_{23} \wedge (A_3^1 dx + A_3^2 dy + A_3^3 dz) \\ +\omega_{31} \wedge (A_1^1 dx + A_1^2 dy + A_1^3 dz) - \omega_{23} \wedge (A_2^1 dx + A_2^2 dy + A_2^3 dz) \end{bmatrix} \end{aligned}$$

It will get really messy if we calculate each component here, so let us calculate one component separately first:

$$\begin{aligned} \omega_{ij} \wedge (A_k^1 dx + A_k^2 dy + A_k^3 dz) &= ((\omega^{ij})_1 dx + (\omega^{ij})_2 dy + (\omega^{ij})_3 dz) \wedge (A_k^1 dx + A_k^2 dy + A_k^3 dz) \\ &= (\omega^{ij})_1 A_k^2 dx \wedge dy + (\omega^{ij})_1 A_k^3 dx \wedge dz \\ &\quad + (\omega^{ij})_2 A_k^1 dy \wedge dx + (\omega^{ij})_2 A_k^3 dy \wedge dz \\ &\quad + (\omega^{ij})_3 A_k^1 dz \wedge dx + (\omega^{ij})_3 A_k^2 dz \wedge dy \\ &= ((\omega^{ij})_1 A_k^2 - (\omega^{ij})_2 A_k^1) dx \wedge dy \\ &\quad + ((\omega^{ij})_2 A_k^3 - (\omega^{ij})_3 A_k^2) dy \wedge dz \\ &\quad + ((\omega^{ij})_3 A_k^1 - (\omega^{ij})_1 A_k^3) dz \wedge dx \end{aligned}$$

where we use the fact that  $dx \wedge dx = 0$  and  $dx \wedge dy = -dy \wedge dx$ . We can clean up the above expression using the cross product:

$$\begin{bmatrix} ((\omega^{ij})_2 A_k^3 - (\omega^{ij})_3 A_k^2) \\ ((\omega^{ij})_3 A_k^1 - (\omega^{ij})_1 A_k^3) \\ ((\omega^{ij})_1 A_k^2 - (\omega^{ij})_2 A_k^1) \end{bmatrix}^\top \begin{bmatrix} dy \wedge dz \\ dz \wedge dx \\ dx \wedge dy \end{bmatrix} = [W_i \times A^k]^\top \begin{bmatrix} dy \wedge dz \\ dz \wedge dx \\ dx \wedge dy \end{bmatrix}$$

We use the fact that  $[A_1, A_2, A_3] = [A^1, A^2, A^3]$  because  $A$  is symmetric, and  $W_1$  corresponds to  $\omega_{23}$ ,  $W_2$  to  $\omega_{31}$  and  $W_3$  to  $\omega_{12}$ . Second part:

$$d(g^{1/2}dp) = d(Adp) = d\left(\begin{bmatrix} A_1^1 & A_1^2 & A_1^3 \\ A_2^1 & A_2^2 & A_2^3 \\ A_3^1 & A_3^2 & A_3^3 \end{bmatrix} \begin{bmatrix} dx \\ dy \\ dz \end{bmatrix}\right)$$

Again, we can do this separately for each row:

$$\begin{aligned} & d(A_k^1 dx + A_k^2 dy + A_k^3 dz) \\ &= dA_k^1 \wedge dx + dA_k^2 \wedge dy + dA_k^3 \wedge dz \\ &= \frac{\partial A_k^1}{\partial x} dx \wedge dx + \frac{\partial A_k^1}{\partial y} dy \wedge dx + \frac{\partial A_k^1}{\partial z} dz \wedge dx \\ &+ \frac{\partial A_k^2}{\partial x} dx \wedge dy + \frac{\partial A_k^2}{\partial y} dy \wedge dy + \frac{\partial A_k^2}{\partial z} dz \wedge dy \\ &+ \frac{\partial A_k^3}{\partial x} dx \wedge dz + \frac{\partial A_k^3}{\partial y} dy \wedge dz + \frac{\partial A_k^3}{\partial z} dz \wedge dz \\ &= \left(\frac{\partial A_k^3}{\partial y} - \frac{\partial A_k^2}{\partial z}\right) dy \wedge dz + \left(\frac{\partial A_k^1}{\partial z} - \frac{\partial A_k^3}{\partial x}\right) dz \wedge dx + \left(\frac{\partial A_k^2}{\partial x} - \frac{\partial A_k^1}{\partial y}\right) dx \wedge dy \\ &= (\nabla \times A_k)^\top \begin{bmatrix} dy \wedge dz \\ dz \wedge dx \\ dx \wedge dy \end{bmatrix} \end{aligned}$$

Finally, we can put everything together:

$$\begin{aligned} \mathbf{0} &= \begin{bmatrix} (W_3 \times A^2 - W_2 \times A^3 + \nabla \times A^1)^\top \\ (W_1 \times A^3 - W_3 \times A^1 + \nabla \times A^2)^\top \\ (W_2 \times A^1 - W_1 \times A^2 + \nabla \times A^3)^\top \end{bmatrix} \begin{bmatrix} dy \wedge dz \\ dz \wedge dx \\ dx \wedge dy \end{bmatrix} \\ &\iff \begin{bmatrix} (W_2 \times A^3 - W_3 \times A^2)^\top \\ (W_3 \times A^1 - W_1 \times A^3)^\top \\ (W_1 \times A^2 - W_2 \times A^1)^\top \end{bmatrix} \begin{bmatrix} dy \wedge dz \\ dz \wedge dx \\ dx \wedge dy \end{bmatrix} = \begin{bmatrix} (\nabla \times A^1)^\top \\ (\nabla \times A^2)^\top \\ (\nabla \times A^3)^\top \end{bmatrix} \begin{bmatrix} dy \wedge dz \\ dz \wedge dx \\ dx \wedge dy \end{bmatrix} \end{aligned}$$

Take the curl to the other side and switch order on the left-hand side to cancel the  $-1$ . As we are only interested in the 9 components of  $W$ , we omit the two-form basis and transform into a 9x9 linear system for  $W$ . We define  $A_\times$  and  $\text{vec}(\cdot)$  as

$$A_\times = \begin{bmatrix} 0 & -A_\times^3 & A_\times^2 \\ A_\times^3 & 0 & -A_\times^1 \\ -A_\times^2 & A_\times^1 & 0 \end{bmatrix}, \text{vec}(W) = \begin{bmatrix} W_1 \\ W_2 \\ W_3 \end{bmatrix}$$

with  $A_\times^i$  defined as

$$A_\times^i = [A_1^i \quad A_2^i \quad A_3^i]_\times = \begin{bmatrix} 0 & -A_3^i & A_2^i \\ A_3^i & 0 & -A_1^i \\ -A_2^i & A_1^i & 0 \end{bmatrix}$$

and  $\text{vec}(\cdot)$  turns a 3x3-matrix into a 9x1-vector by stacking the columns. With these two definitions, we can transform the above equality into a linear system

$$A_\times \text{vec}(W) = \text{vec}(\nabla \times A)$$

where  $\nabla \times A$  is just the curl applied to each column. This transformation can be checked by laboriously plugging in the definitions and comparing the coefficients. With tedious calculations, one can show that  $\det(A_\times) = -2 \det(A)^3 = -2 \det(g)^{3/2} < 0$ , which means this is a linear system that is solvable and can be used to calculate  $W$  at a point.





## Chapter 4

# Calculation of rotation coefficient

We are interested in the rotation. Let  $p, q$  be points on the manifold and  $\ell$  a path connecting them. Let  $R : \mathcal{M} \rightarrow SO(3)$  be a rotation field, i.e.  $R(p), R(q)$  are orthogonal frames. Under the initial condition that  $R(p) = \text{Id}$ , the equation

$$R(q) = \exp(R_{pq}(\omega))R(p)$$

is a differential equation that corresponds to the parallel transport under the connection  $\omega$  of the frame along  $\ell$ . To recover this rotation  $R_{pq}$ , we integrate  $\omega$  along  $\ell$ .

**Discretization** We parametrize the path by  $\ell(0) = a, \ell(1) = b$ . We resort to numerical integration for  $R_{ab}$  and cut the path into  $n$  small segments, i.e.

$$R_{ab} = R_n R_{n-1} \cdots R_1$$

where  $R_i = \exp(-\omega(\dot{\ell}(i\gamma))\gamma) = \exp((\int W^\top dp)_\times)$ , and  $\gamma$  is the length of a segment and  $\dot{\ell}(s) = \frac{\partial \ell}{\partial s}(s)$ . Calculating the exponential map of an antisymmetric matrix (which  $\omega \in \mathfrak{so}(3)$  is) can be done with Rodrigues' formula:

$$\exp(u_\times) = \text{Id} + \sin(\theta)\hat{u}_\times + (1 - \cos(\theta))\hat{u}_\times^2$$

where  $\theta = \|u\|_2$  is the rotation angle and  $\hat{u} = u/\theta$  is the rotation axis. We use the trapezoidal rule to evaluate the a short interval of the integral of  $W$ , which is given by

$$R_{ab} = \exp\left(\left(\frac{1}{2}(W_a + W_b)^\top(b - a)\right)_\times\right)$$

where  $W_a, W_b$  is solved for by the 9x9 linear system given by  $A_x$  and  $\nabla \times A$ .

**Piecewise linear discretization** We discretize our metric field with a tetrahedral mesh  $\mathcal{T}$ . At each vertex, we attach a metric and linearly interpolate with barycentric coordinates within a tet.

Let  $A_i \in \mathbb{R}^{3 \times 3}, i \in \{1, 2, 3, 4\}$  be the square root metrics at the vertices  $v_i \in \mathbb{R}^3$  of a tet, such that  $A_i^2 = g(v_i)$ . We represent a tet given by its four vertices by a 3x4 matrix, i.e.

$$\begin{pmatrix} | & | & | & | \\ v_1 & v_2 & v_3 & v_4 \\ | & | & | & | \end{pmatrix} \in \mathbb{R}^{3 \times 4}.$$

Any point  $p$  within the tet can then be represented as

$$p = \alpha v_1 + \beta v_2 + \gamma v_3 + \delta v_4$$

with  $\alpha, \beta, \gamma, \delta \geq 0$  and  $\alpha + \beta + \gamma + \delta = 1$ . This is a linear transformation between two coordinate systems, which we can write in matrix form as

$$\underbrace{\begin{pmatrix} | & | & | & | \\ v_1 & v_2 & v_3 & v_4 \\ | & | & | & | \\ 1 & 1 & 1 & 1 \end{pmatrix}}_{T^{-1}} \underbrace{\begin{pmatrix} \alpha \\ \beta \\ \gamma \\ \delta \end{pmatrix}}_{\lambda} = \underbrace{\begin{pmatrix} x \\ y \\ z \\ 1 \end{pmatrix}}_p \iff T^{-1}\lambda = p \iff \lambda = Tp$$

$T$  always exists because  $v_1, \dots, v_4$  are linearly independent, else it would not be a tetrahedron. By denoting  $T = \{t_{ij}\}_{i,j \in \{1, \dots, 4\}}$ , we can write our barycentric functions as

$$\begin{aligned} \alpha(x, y, z) &= t_{11}x + t_{12}y + t_{13}z + t_{14} \\ \beta(x, y, z) &= t_{21}x + t_{22}y + t_{23}z + t_{24} \\ \gamma(x, y, z) &= t_{31}x + t_{32}y + t_{33}z + t_{34} \\ \delta(x, y, z) &= t_{41}x + t_{42}y + t_{43}z + t_{44} \end{aligned}$$

The convex combination

$$A(x, y, z) = \alpha A_1 + \beta A_2 + \gamma A_3 + \delta A_4$$

is the metric prescribed in the tet. To find  $\nabla \times A$ , let  $A = (A^1, A^2, A^3)$ . We will need the derivatives for the curl, so let

$$(A_j^i)_x \triangleq \frac{\partial A_j^i}{\partial x}$$

be the derivative with respect to  $x$  of entry  $i, j$ . E.g.  $(A_j^i)_x$  is given by

$$(A_j^i)_x = \alpha_x (A_1)_j^i + \beta_x (A_2)_j^i + \gamma_x (A_3)_j^i + \delta_x (A_4)_j^i = t_{11}(A_1)_j^i + t_{21}(A_2)_j^i + t_{31}(A_3)_j^i + t_{41}(A_4)_j^i$$

If we write  $T = (T^1, T^2, T^3, T^4)$  and collect  $(A_k)_j^i$  into a vector

$$\bar{A}_j^i = \begin{pmatrix} (A_1)_j^i \\ (A_2)_j^i \\ (A_3)_j^i \\ (A_4)_j^i \end{pmatrix}$$

this can be shortened to  $\bar{A}_j^{i\top} T^1 = (A_j^i)_x$ . Analogously, we get

$$(A_j^i)_y = \bar{A}_j^{i\top} T^2 \text{ and } (A_j^i)_z = \bar{A}_j^{i\top} T^3$$

The curl is then given by

$$\nabla \times A^i = \begin{pmatrix} (A_3^i)_y - (A_2^i)_z \\ (A_1^i)_z - (A_3^i)_x \\ (A_2^i)_x - (A_1^i)_y \end{pmatrix} = \begin{pmatrix} \bar{A}_3^{i\top} T^2 - \bar{A}_2^{i\top} T^3 \\ \bar{A}_1^{i\top} T^3 - \bar{A}_3^{i\top} T^1 \\ \bar{A}_2^{i\top} T^1 - \bar{A}_1^{i\top} T^2 \end{pmatrix}$$

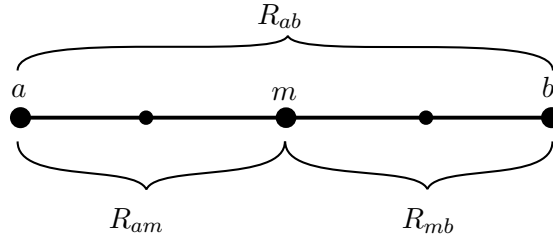
and  $\nabla \times A = \nabla \times (A^1, A^2, A^3)$ . Notice that the curl is constant within a tetrahedron.

**Recursive Subdivision** Whenever the rotation coefficient  $R_{pq}$  is needed between some points  $p$  and  $q$ , it is unclear ahead of time how many sampling points on the line  $\ell$  are needed such that  $R_{pq}$  accurately describes how the frame rotates along  $\ell$ . We apply a recursive subdivision scheme to recursively sample more points on  $\ell$  only where is needed until sampling more points leads to no noticable improvement

anymore (see figure 4.1). We begin by calculating  $R_{ab}$  with just the endpoints. This coefficient is then compared to the result if the midpoint was sampled as well, so i.e. if

$$\frac{\|R_{ab} - R_{am} \cdot R_{mb}\|_2^2}{\ell^2} < \varepsilon$$

then no measurable improvement happened. We divide by the length of the segment  $\vec{ab}$ , because the segments get smaller and we want the tolerance  $\varepsilon$  to remain the same. This approach has the advantage of only sampling points where there is improvement. See section ?? for results.



**Figure 4.1.** Points are recursively sampled as midpoints between the line  $\vec{ab}$ . If sampling more points within some line segment leads to noticeable improvement, more points are sampled.

The following code does what is described above.

---

**Algorithm 1** Recursive Subdivision

---

```

1: recursiveDivide ( $a, b$ )
2:    $R_1 = \text{Rotation}(a, b)$ 
3:    $\ell = \text{length}(a - b)$ 
4:    $\text{midpoint} = \frac{a+b}{2}$ 
5:   if  $\frac{\|R_1 - \text{Rotation}(a, \text{midpoint}) \cdot \text{Rotation}(\text{midpoint}, b)\|_2^2}{\ell^2} < \varepsilon$ 
6:     return  $R_1$ 
7:   else
8:     return  $\text{recursiveDivide}(a, \text{midpoint}) \cdot \text{recursiveDivide}(\text{midpoint}, b)$ 

```

---



## Chapter 5

# Algorithm for $R$ between two arbitrary points in a mesh

To measure the Dirichlet energy, we need to calculate the rotation coefficient between two arbitrary points  $q$  and  $p$  that do not necessarily lie within the same tetrahedron. Since the metric and curl is different in each tet, we need to be able to efficiently determine all tets that get intersected by the straight line from  $q$  to  $p$ , and use the correct metric for each corresponding line segment. The calculation for the coefficient then works in the following way:

---

**Algorithm 2** Rotation coefficient  $R$  between  $q$  and  $p$ 

---

```
1: Input  $(q, p)$ 
2:   LINESEGMENTS  $\leftarrow tetFinder(q, p)$  //returns all tets intersected by the line  $\vec{pq}$  with the line segments within them
3:    $R \leftarrow Id$ 
4:   for each SEGMENT in LINESEGMENTS
5:      $R \leftarrow R \cdot calcCoeff(SEGMENT)$ 
6: return  $R$ 
```

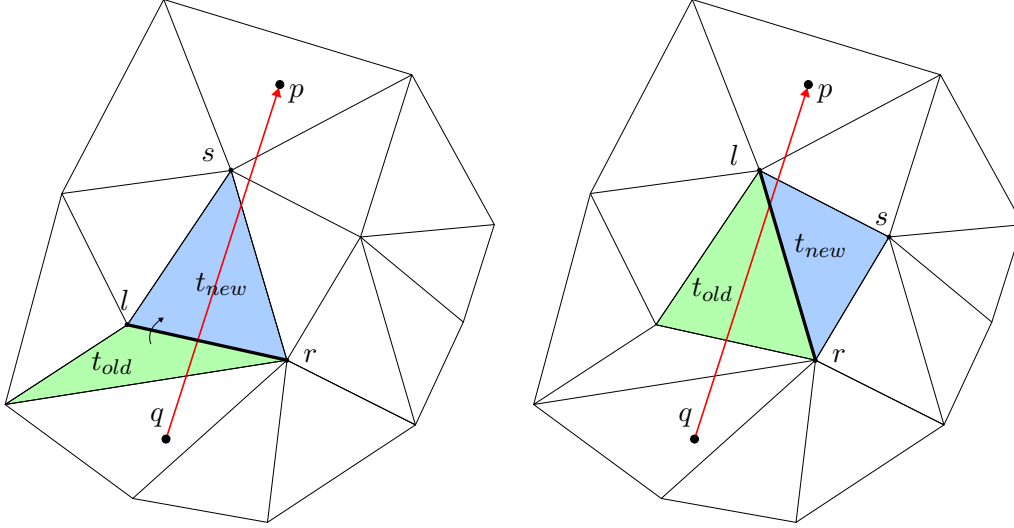
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The missing component here is how to efficiently find all tetrahedra that get intersected. One possibility would be to use ray-triangle intersection and test against the whole mesh, but this is not practical, as we have local information that we can exploit.

We use the idea of the straight walk from *Walking in a Triangulation*[3], which relies only on so called *orientation tests* to determine which triangles we traverse.

**Framework** Let  $\mathcal{T}$  be a triangulation of a domain  $\Omega$  that is convex. The straight walk traverses all triangles that get intersected by the line segment from  $q$  to  $p$ . The algorithm first makes an initialization step to get into a valid state, then the straight walk can start. To get a feeling how the algorithm works, let us go through an example in 2D. If the algorithm was in a valid state before, the line from  $q$  to  $p$  intersects with some edge  $\vec{lr}$ . Two triangles share this edge. We test on which side point  $p$  lies of this edge to decide whether the walk continues. If the walk continues, we jump through the edge to hop from the old triangle to a new one. This triangle is defined by three vertices  $(l, r, s)$ . We decide if the new candidate point  $s$  lies on the left side or right side of the line from  $q$  to  $p$ . If  $s$  lies on the left, point  $l$  is moved, else point  $r$  is moved. A new edge intersected with the ray  $\vec{qp}$  is found and the walk repeats. This process is illustrated in Figure 5.1

Notice how the ray  $\vec{qp}$  always intersects the edge  $\vec{lr}$  at each update step. We can use this observation to add each edge at each update step to a list. When the algorithm terminates, we can just iterate over this list, find the intersection point of the ray  $\vec{qp}$  with the edge and calculate the rotation coefficient for this segment. The straight walk in 3d works similarly. The initialization step consists of finding a starting tet  $t$  where  $q$  is contained. Then, we find the face of the tet  $t$  that gets intersected by the ray  $\vec{qp}$ . Again, at each



**Figure 5.1.** Straight walk step

step, we know that the ray goes out of our current tet  $t$  through some face  $e$  defined by vertices  $uvw$ . We decide if the walk continues by checking on which side  $p$  lies relative to  $e$ . If the walk should continue, we hop through  $e$  to a new tet  $t_{new}$ . With two orientation tests we decide which of the vertices  $u, v, w$  gets moved to the new candidate point  $s$ . This defines a new face  $e_{new}$  where our ray intersects and the walk repeats. Degenerate cases such as when the ray  $\vec{qp}$  goes exactly through a vertex or when the ray lies within a face, the algorithm may get into an invalid state, where the algorithm may then traverse through cells that do not get intersected by the ray. We need to detect and escape those degenerate cases through some additional checks. These are not described in the paper[3]. In most cases, a simple reinitialization is enough when the invalid state is detected.

**A note on orientation tests** To determine on which side some point  $s$  lies relative to two other points  $q, p$  (that represent a line) in 2D, the geometric *orientation* predicate is used. It corresponds to evaluating the sign of a determinant. Analogously in 3D, the orientation predicate tests whether a fourth point lies above or below a plane defined by three other points. How “above” the plane is defined depends on the ordering within the determinant. In this case, it is above if point  $a$  sees the triangle  $bcd$  when turning counterclockwise (see figure 5.2).

$$orientation(\alpha, \beta, \gamma) = sign \left( \begin{vmatrix} \beta_x - \alpha_x & \gamma_x - \alpha_x \\ \beta_y - \alpha_y & \gamma_y - \alpha_y \end{vmatrix} \right)$$

$$orientation(\alpha, \beta, \gamma, \delta) = sign \left( \begin{vmatrix} \beta_x - \alpha_x & \gamma_x - \alpha_x & \delta_x - \alpha_x \\ \beta_y - \alpha_y & \gamma_y - \alpha_y & \delta_y - \alpha_y \\ \beta_z - \alpha_z & \gamma_z - \alpha_z & \delta_z - \alpha_z \end{vmatrix} \right)$$

It is important that the sign of the determinant is evaluated exactly. If geometric predicates are implemented with floating point arithmetic, the answers may be inconsistent and wrong. Because of the finite mantissa in floating point representation [1], many numbers cannot be represented exactly. The machine needs to round a number  $x$  to its nearest number that is exactly representable. Many times, roundoff errors occur, a famous example of this is  $0.1 + 0.2 = 0.30000000000000004$ . We call the distance between two exactly representable floating point numbers *machine epsilon*  $\epsilon$ . See figure 5.3 for how this machine epsilon can cause trouble (Example from <http://groups.csail.mit.edu/graphics/classes/6.838/S98/meetings/m12/pred/m12.htm>). Logical decision based on floating point arithmetic is correct most of the times, but not always. To fix this issue, we make use of exact predicates that are implemented with arbitrary precision [13]. Inputs to the orientation test subroutines are assumed to be exact, and we can be sure that the result has the correct sign.

---

**Algorithm 3** tetFinder

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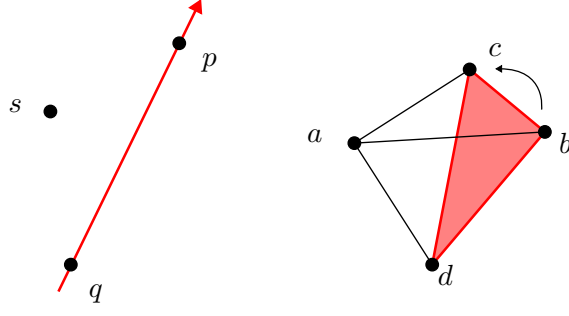
```
1: Input  $(q, p)$ 
2:    $t = \text{LocateTet}(q)$ 
3:    $\text{initialization}(t)$ 

4:   //  $qp$  intersects triangle  $uvw$ 
5:   //  $wvqp, vuqp, uwqp$  are positively oriented
6:    $\text{LINESEGMENTS} = []$ 
7:    $\text{PREV} = q$ 
8:    $\text{CURR} = \text{intersection}(qp, uvw)$ 

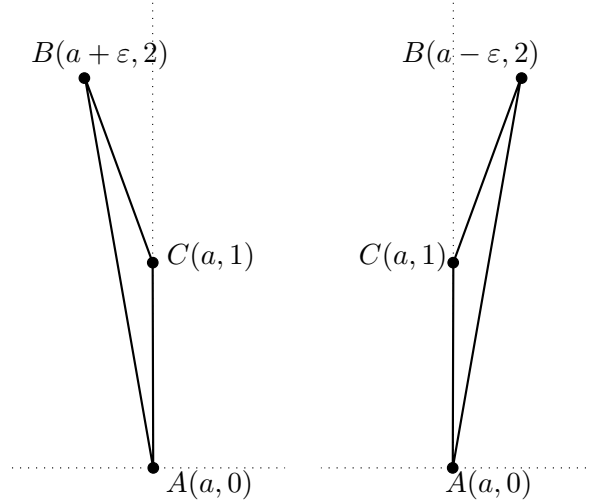
9:    $\text{LINESEGMENTS.add}([t, \text{PREV}, \text{CURR}])$ 

10:  while  $\text{orientation}(u, w, v, p) > 0$  {
11:     $t = \text{neighbor}(t \text{ through } uvw)$ 
12:     $s = \text{vertex of } t, s \neq u, s \neq v, s \neq w$ 
13:     $\text{PREV} = \text{CURR}$ 
14:    if  $\text{orientation}(u, s, q, p) > 0$  //  $qp$  does not intersect triangle  $usw$ 
15:      if  $\text{orientation}(v, s, q, p) > 0$  //  $qp$  intersects triangle  $vsw$ 
16:         $u = s$ 
17:      else //  $qp$  intersects triangle  $usv$ 
18:         $w = s$ 
19:    else //  $qp$  does not intersect  $usv$ 
20:      if  $\text{orientation}(w, s, q, p) > 0$  //  $qp$  intersects triangle  $usw$ 
21:         $v = s$ 
22:      else //  $qp$  intersects triangle  $vsw$ 
23:         $u = s$ 
24:     $\text{CURR} = \text{intersection}(qp, uvw)$ 
25:     $\text{LINESEGMENTS.add}([t, \text{PREV}, \text{CURR}])$ 
26:  } //  $t$  contains  $p$ 
27:   $\text{LINESEGMENTS.add}([t, \text{PREV}, p])$ 
```

---



**Figure 5.2.** Orientation predicate: Point  $a$  lies above the plane  $bcd$  because it sees the points in counter-clockwise order,  $\text{orientation}(a, b, c, d) > 0$



**Figure 5.3.** The orientation predicate implemented with floating point arithmetic cannot distinguish these two cases. The orientation test calculates  $(a + \varepsilon) - a$ . If  $a = 0$ , the orientation test is strictly positive. If  $a = 1$ , then the result is 0, because during the calculation the machine had to round.

The final algorithm to find all intersected tetrahedra between two points is as follows.  
The algorithm uses several subroutines which are described here:

- `neighbor(t through uvw)` returns the tetrahedron sharing face  $uvw$  with tetrahedron  $t$
- `vertex of t,  $s \neq u, s \neq v, s \neq w$`  returns the remaining vertex of a tetrahedron whose other three vertices are known
- `initialization()`



## **Chapter 6**

# **Conclusion**

The conclusion looks back at the entire work, gives a critical look, summarizes, and discusses extensions and future work.



## **Appendix A**

### **Extra material**

Extra material may be placed in an appendix that appears after the conclusion.



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*Erklärung gemäss Art. 30 RSL Phil.-nat. 18*

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