

**Formal Language And Automata Theory (CS21004) - Class Test 1**

*IIT Kharagpur, CSE Dept., Spring'23*

Answer all questions. In case of reasonable doubt, make practical assumptions.  
Marks will be deducted for sketchy proofs and claims without proper reasoning.

Marks = 20

Time = 1 hour

1. Construct a DFA for

(a)  $L = \{ab^5wb^2 \mid w \in \{a, b\}^*\}$ . [2.5]

(b)  $L = \{a^n \mid n \geq 0, n \neq 4\}$ . [2.5]

2. Let  $L$  be a regular language on some alphabet  $\Sigma$  and let  $\Sigma_1 \subset \Sigma$  be some smaller alphabet. Prove that  $L_1 = L \cap \Sigma_1^*$  is also regular. [5]

3. For each of the following languages, construct a regular expression that generates it:

(a) the set of binary strings that have both 00 and 11 as substrings; [2.5]

(b) the set of strings over the alphabet  $\{x, y, z\}$  in which each  $y$  is immediately followed by  $x$ ; [2.5]

4. For languages  $A$  and  $B$ , let the perfect shuffle of  $A$  and  $B$  be the language

$$\{w \mid w = a_1b_1 \cdots a_kb_k, \text{ where } a_1 \cdots a_k \in A \text{ and } b_1 \cdots b_k \in B, \text{ each } a_i, b_i \in \Sigma\}$$

Show that the class of regular languages is closed under perfect shuffle. [5]

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