



NPTEL ONLINE CERTIFICATION COURSES

Blockchain and its applications **Prof. Sandip Chakraborty**

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Lecture 28: Paxos

CONCEPTS COVERED

Paxos - CFT Consensus





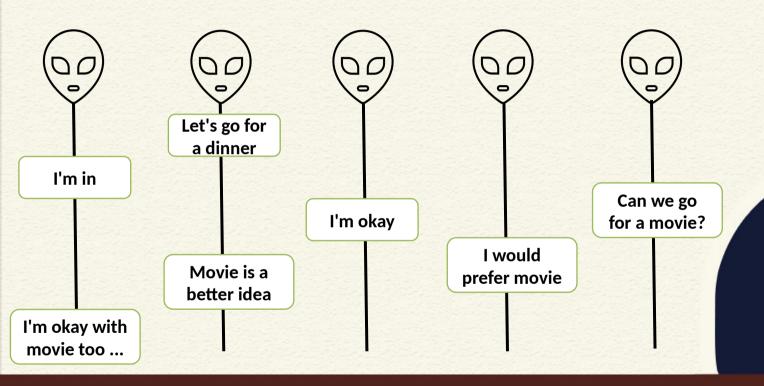
KEYWORDS

- Paxos
- CFT





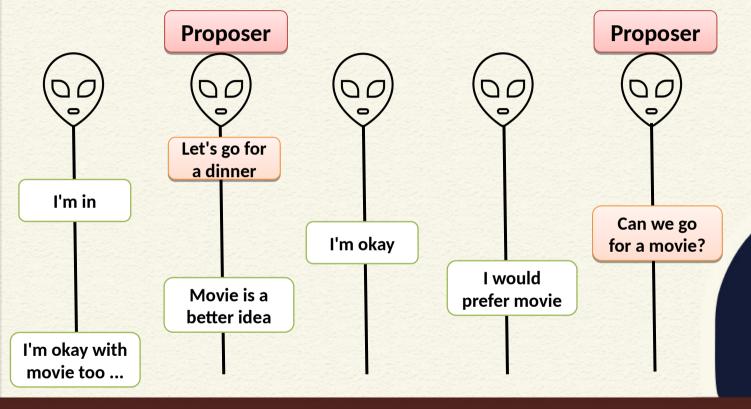
Paxos - The Roles of Individuals







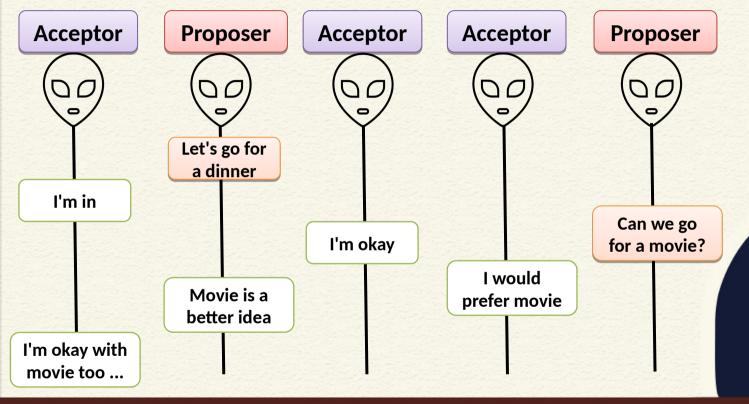
Paxos - The Roles of Individuals





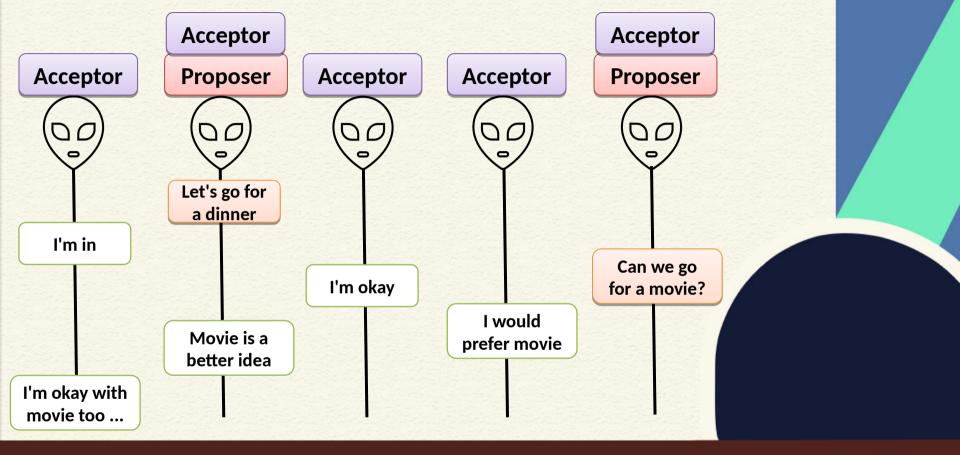


Paxos - The Roles of Individuals



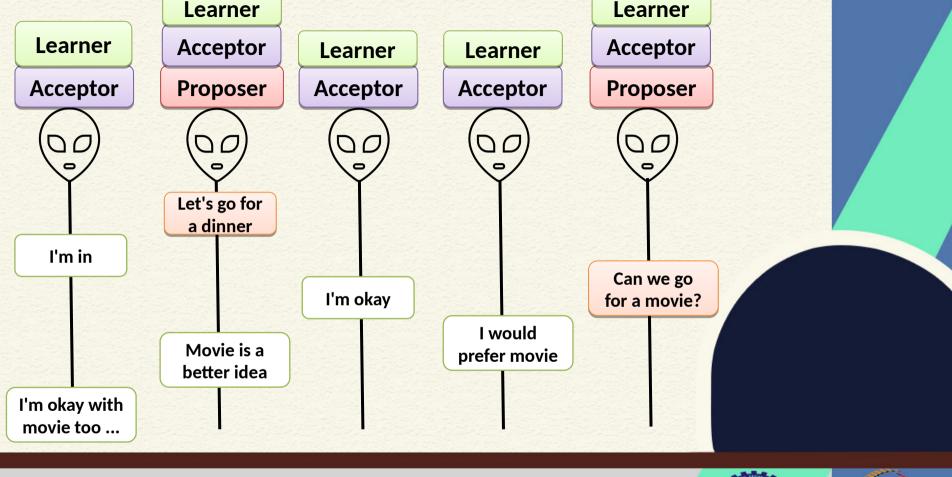






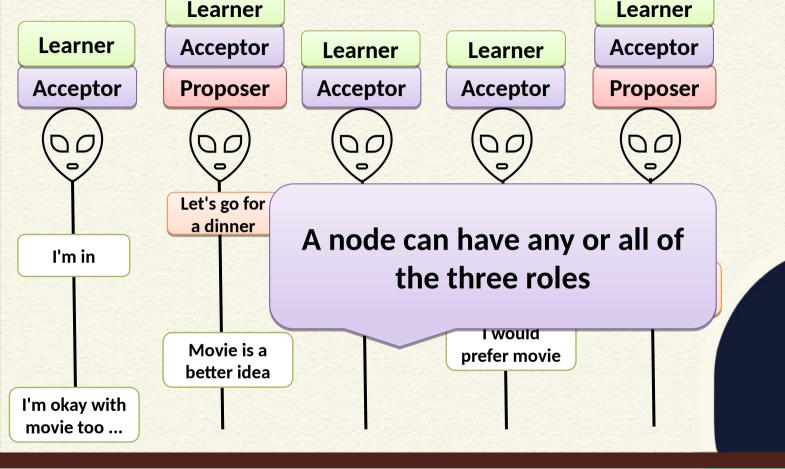






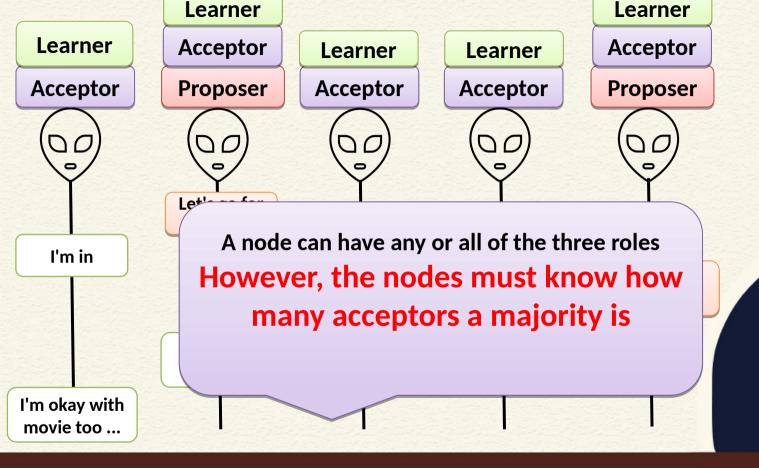






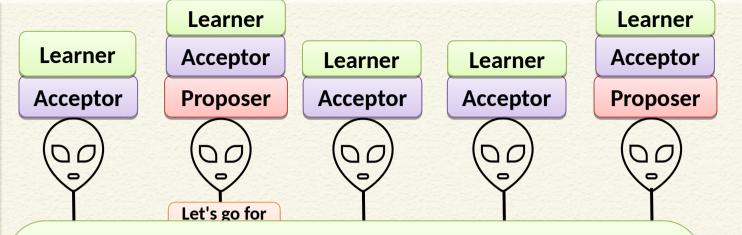












Two majorities will always overlap in atleast one nodes

5 acceptors, majority = 3,

2 proposers:

To accept based on majority voting, at least one acceptor need to choose between one of the two proposals





Paxos Basics

- Paxos is based on state-machine replication
 - Proposers and Acceptors maintain a state of the running epochs
 - Uses a variable IDp where p is an epoch number maintains the state
- A Paxos run aims at reaching a single consensus
 - Once a consensus is reached, Paxos cannot progress to another consensus
 - To reach multiple consensus, you need to run Paxos in rounds (Multi-Paxos)





Paxos Algorithm Proposer Acceptor

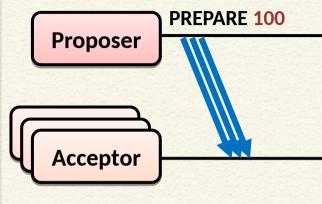
Proposer



- Proposer wants to propose its choice (values):
 - Sends PREPARE IDp to a majority (or all) of the acceptors



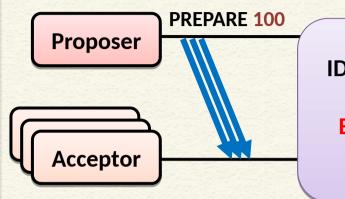




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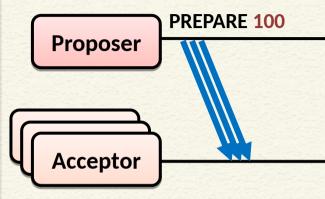
IDp must be unique across proposers for each PREPARE message

Ex. Use Hash(timestamp+Proposer ID) to generate p

- Proposer wants to propose its choice (values):
 - Sends PREPARE IDp to a majority (or all) of the acceptors



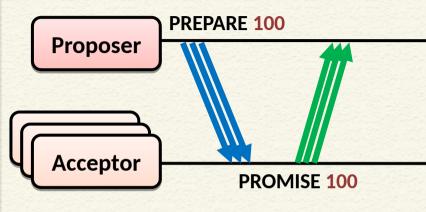




- Acceptor received a PREPARE message with IDp:
 - Did it promised to ignore requests with this IDp?
 - YES: Ignore
 - NO: Will promise to ignore any request lower than IDp
 - (?) Reply with PROMISE IDp



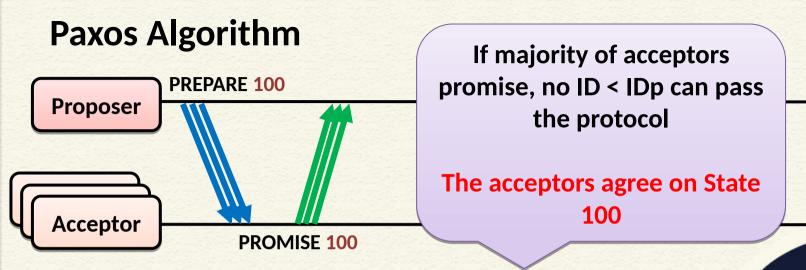




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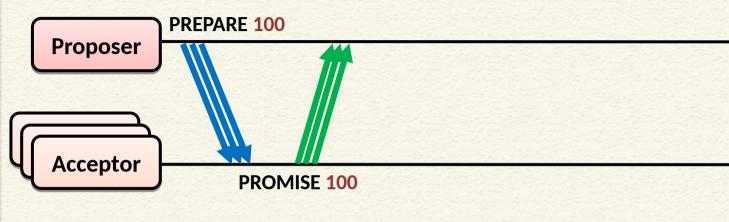




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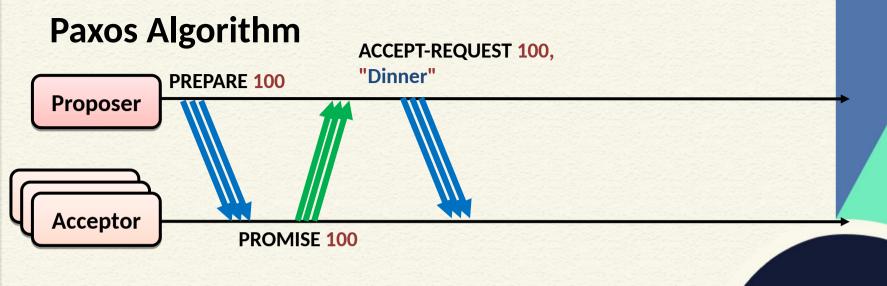




- Proposer gets majority of PROMISE messages for a specific IDp:
 - Sends ACCEPT-REQUEST IDp, <u>VALUE</u> to a majority (or all) of <u>Acceptors</u>
 - (?) It picks any value of its choice



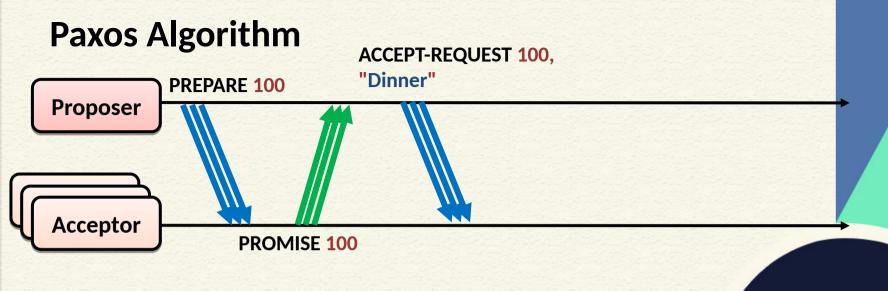




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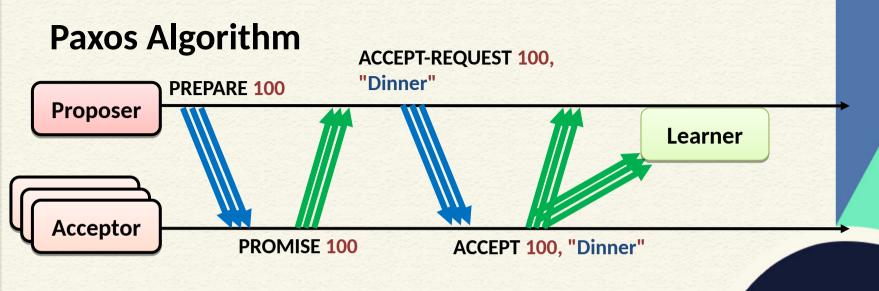




- Acceptor receives an ACCEPT-REQUEST IDp, VALUE :
 - Did it promised to ignore request with this IDp?
 - YES: Ignore
 - NO: Reply with ACCEPT IDp, <u>VALUE</u>; Also send it to all learners



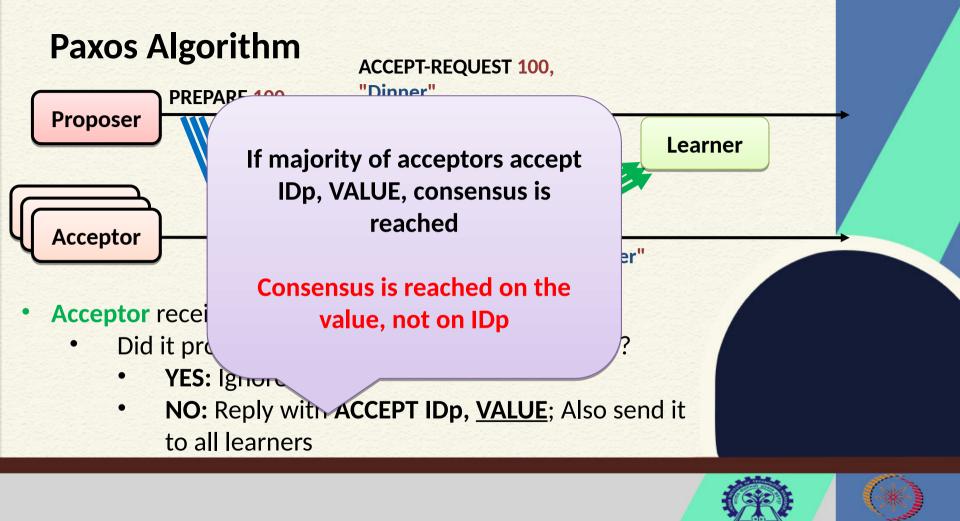


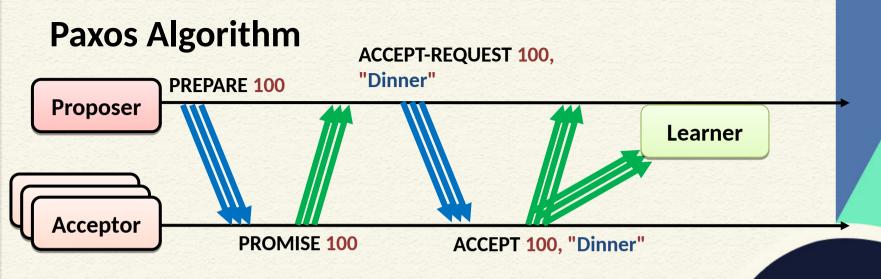


- Acceptor receives an ACCEPT-REQUEST IDp, VALUE :
 - Did it promised to ignore request with this IDp?
 - YES: Ignore
 - NO: Reply with ACCEPT IDp, <u>VALUE</u>; Also send it to all learners









- Proposer or Learner gets ACCEPT message with IDp, VALUE:
 - If a proposer/learner gets majority of accept for a specific IDp, they know that consensus is reached for the value (not IDp).

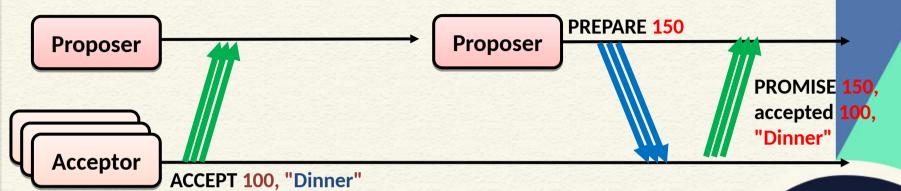




Paxos - Multiple Proposers PREPARE 50 Proposer **Proposer** Acceptor ACCEPT 100, "Dinner"



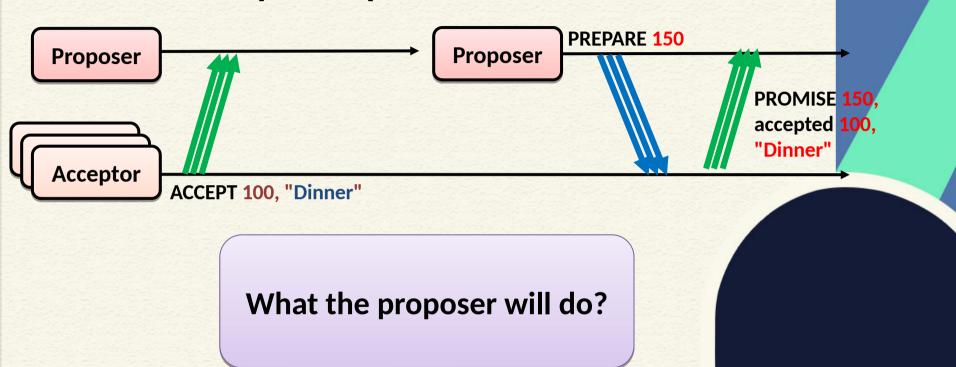




- Acceptor received a PREPARE message with IDp:
 - Did it promised to ignore requests with this IDp?
 - YES: Ignore
 - NO: Will promise to ignore any request lower than IDp
 - Has it ever accepted anything? (Assume accepted ID = IDa)
 - YES: Reply with PROMISE IDp accepted IDa, VALUE
 - NO: Reply with PROMISE IDp

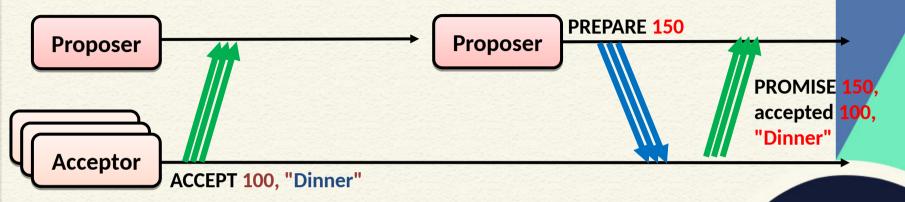








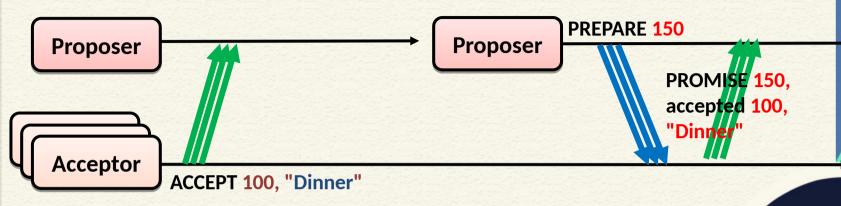




- Proposer gets majority of PROMISE messages for a specific IDp:
 - It sends ACCEPT-REQUEST IDp, <u>VALUE</u> to a majority (or all) of <u>Acceptors</u>
 - (?) It picks any value it wants



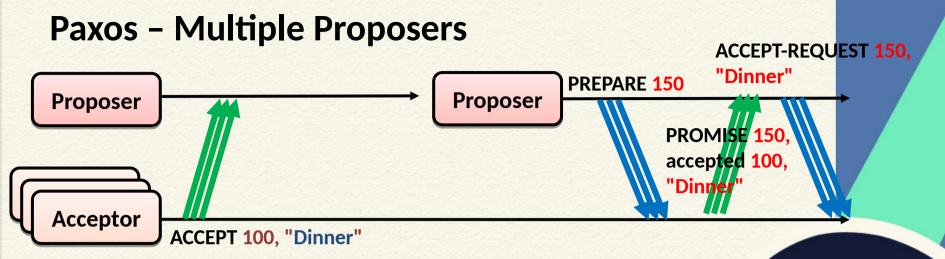




- Proposer gets majority of PROMISE messages for a specific IDp:
 - It sends **ACCEPT-REQUEST IDp**, <u>VALUE</u> to a majority (or all) of **Acceptors**
 - Has it got any already accepted value from promises?
 - YES: Picks the value with the highest IDa
 - NO: Picks the value of its choice







- Proposer gets majority of PROMISE messages for a specific IDp:
 - It sends ACCEPT-REQUEST IDp, <u>VALUE</u> to a majority (or all) of <u>Acceptors</u>
 - Has it got any already accepted value from promises?
 - **YES:** Picks the value with the highest IDa
 - NO: Picks the value of its choice





Conclusion

- Paxos works in two rounds
 - Agreement on the state (ID)
 - Agreement on the value
- Safety and liveness of Paxos?









