

Spanner: Google's Globally-Distributed Database

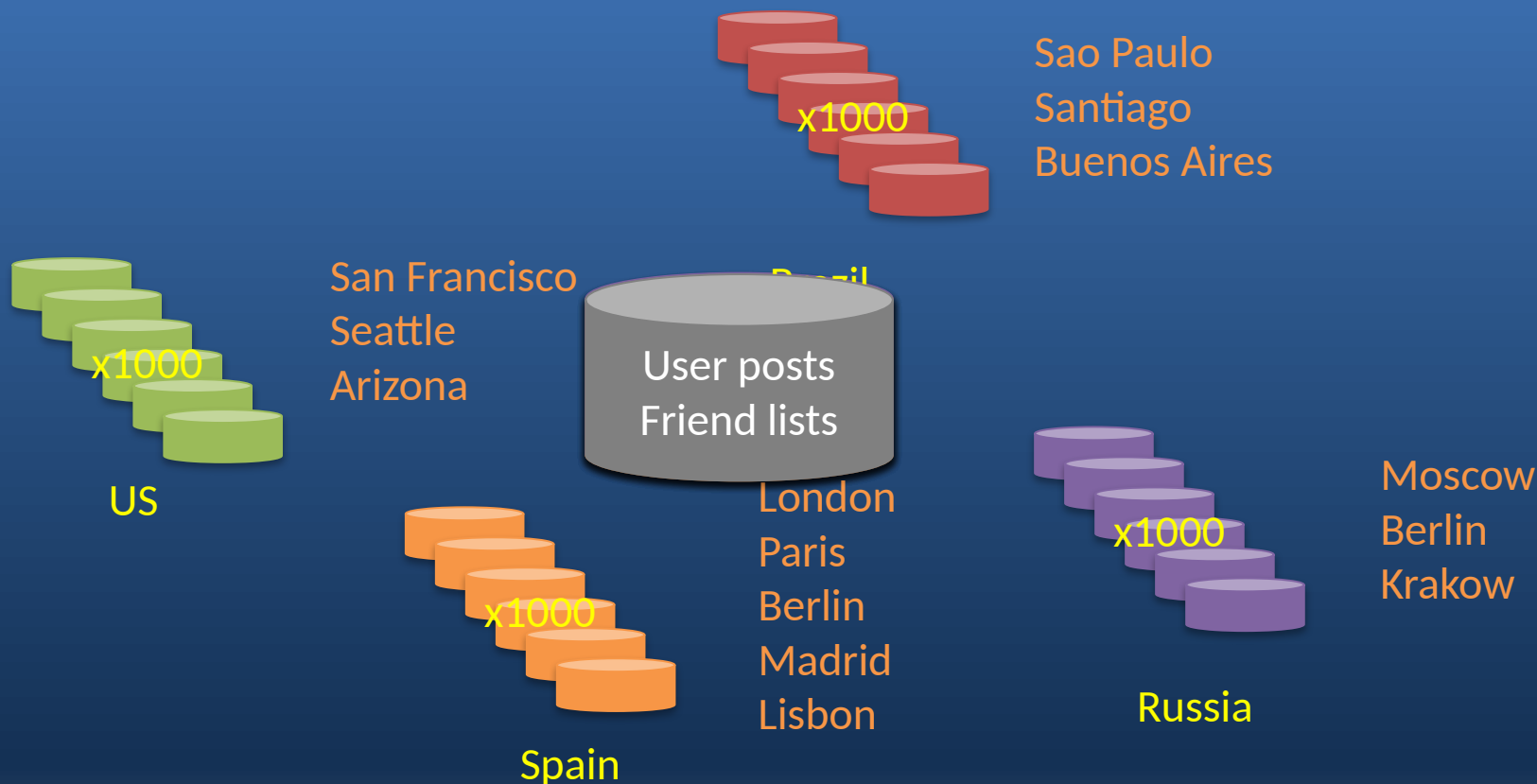
Wilson Hsieh
representing a host of authors
OSDI 2012



What is Spanner?

- Distributed multiversion database
 - General-purpose transactions (ACID)
 - SQL query language
 - Schematized tables
 - Semi-relational data model
- Running in production
 - Storage for Google's ad data
 - Replaced a sharded MySQL database

Example: Social Network



Overview

- Feature: Lock-free distributed read transactions
- Property: External consistency of distributed transactions
 - First system at global scale
- Implementation: Integration of concurrency control, replication, and 2PC
 - Correctness and performance
- Enabling technology: TrueTime
 - Interval-based global time

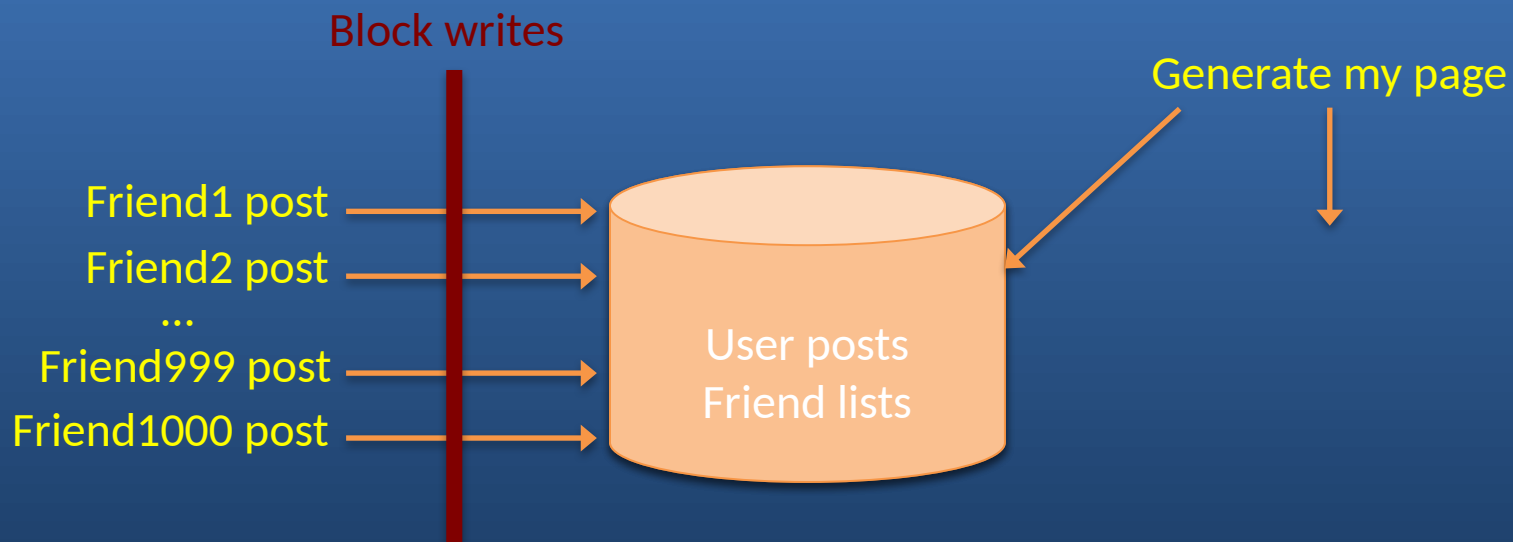
Read Transactions

- Generate a page of friends' recent posts
 - Consistent view of friend list and their posts

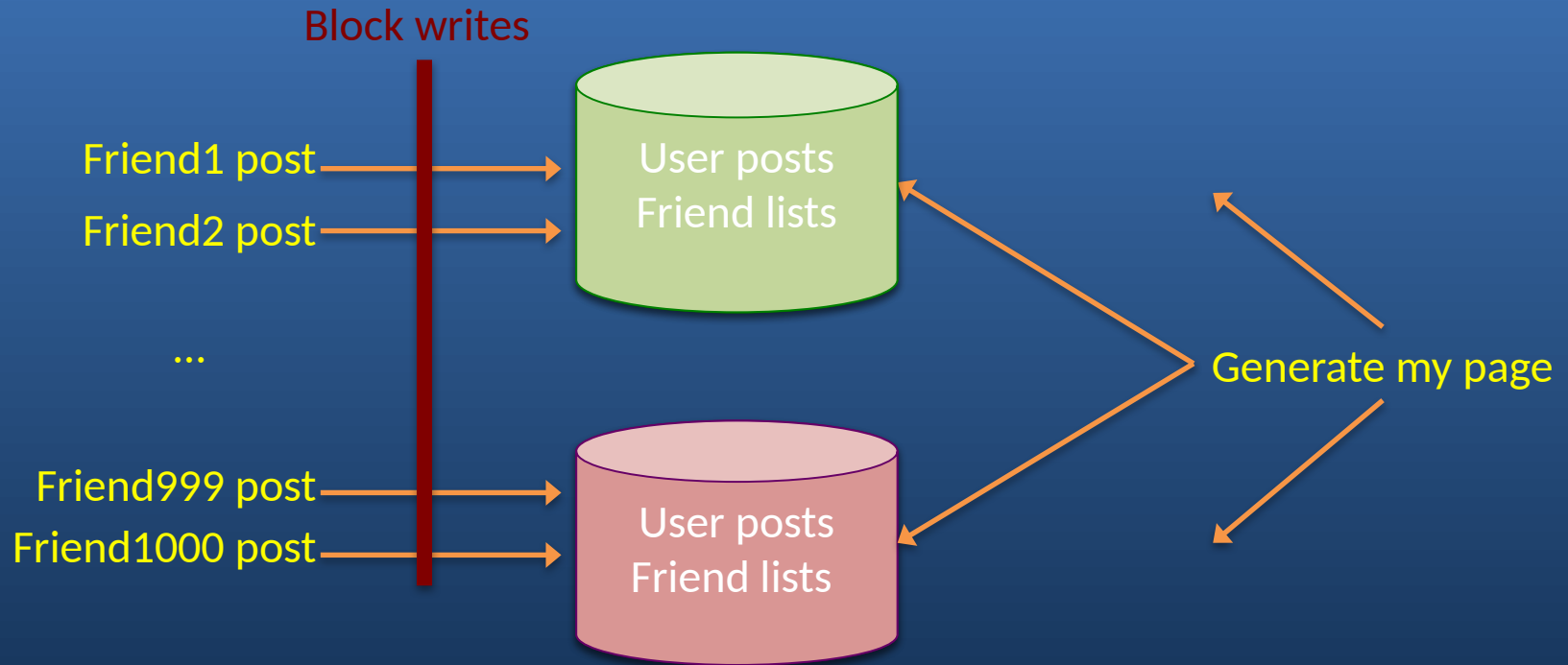
Why consistency matters

1. Remove untrustworthy person X as friend
2. Post P: “My government is repressive...”

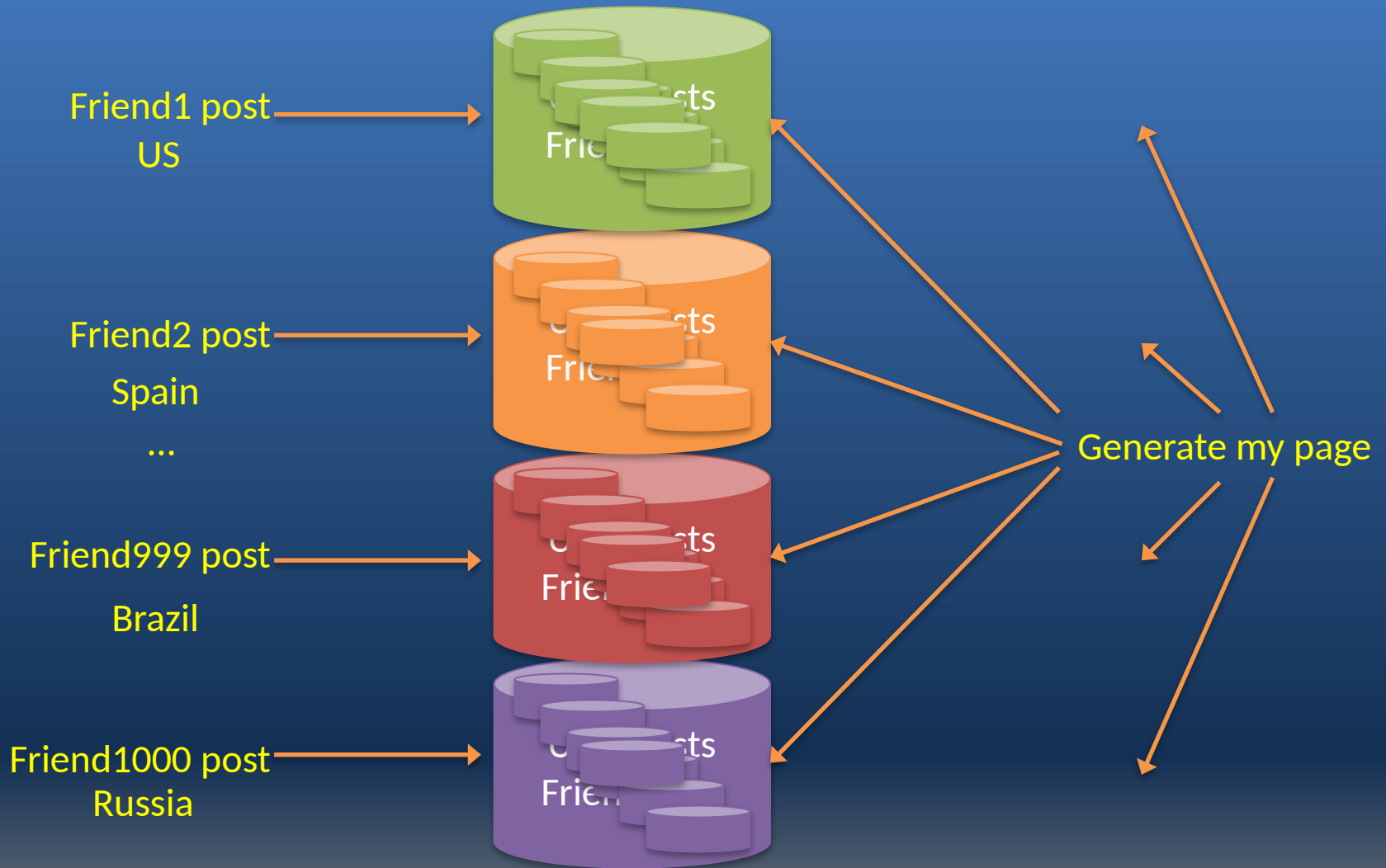
Single Machine



Multiple Machines



Multiple Datacenters



Version Management

- Transactions that write use strict 2PL
 - Each transaction T is assigned a timestamp s
 - Data written by T is timestamped with s

Time	<8	8	15
My friends	[X]	[]	
My posts			[P]
X's friends	[me]	[]	

Synchronizing Snapshots

Global wall-clock time

==

External Consistency:

Commit order respects global wall-time order

==

Timestamp order respects global wall-time order
given

timestamp order == commit order

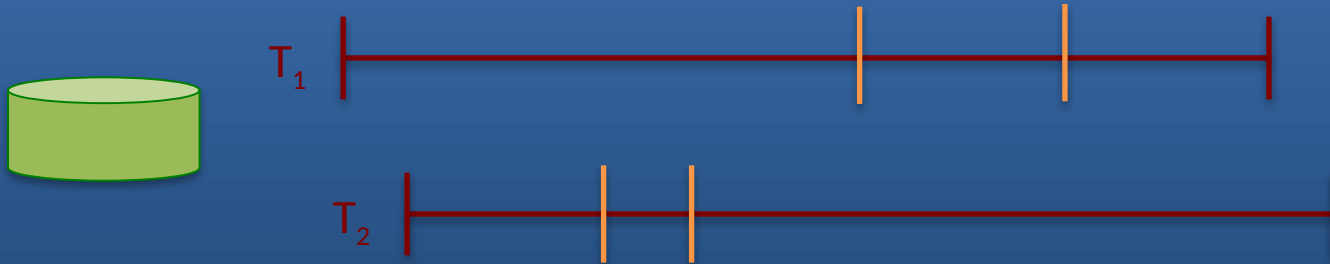
Timestamps, Global Clock

- Strict two-phase locking for write transactions
- Assign timestamp while locks are held



Timestamp Invariants

- Timestamp order == commit order

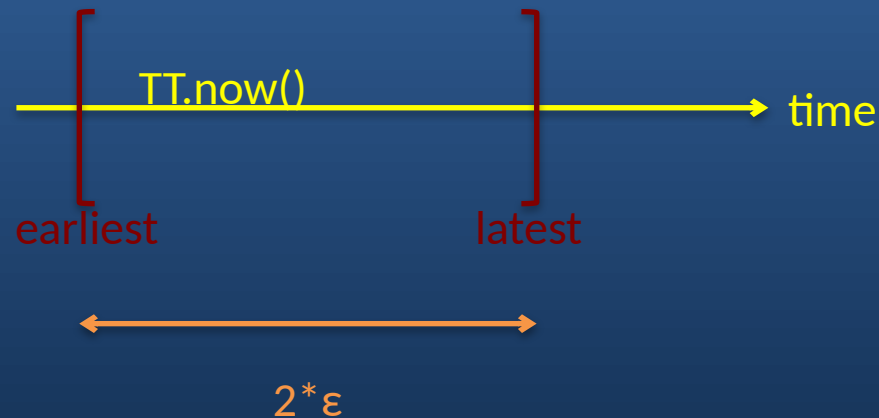


- Timestamp order respects global wall-time order

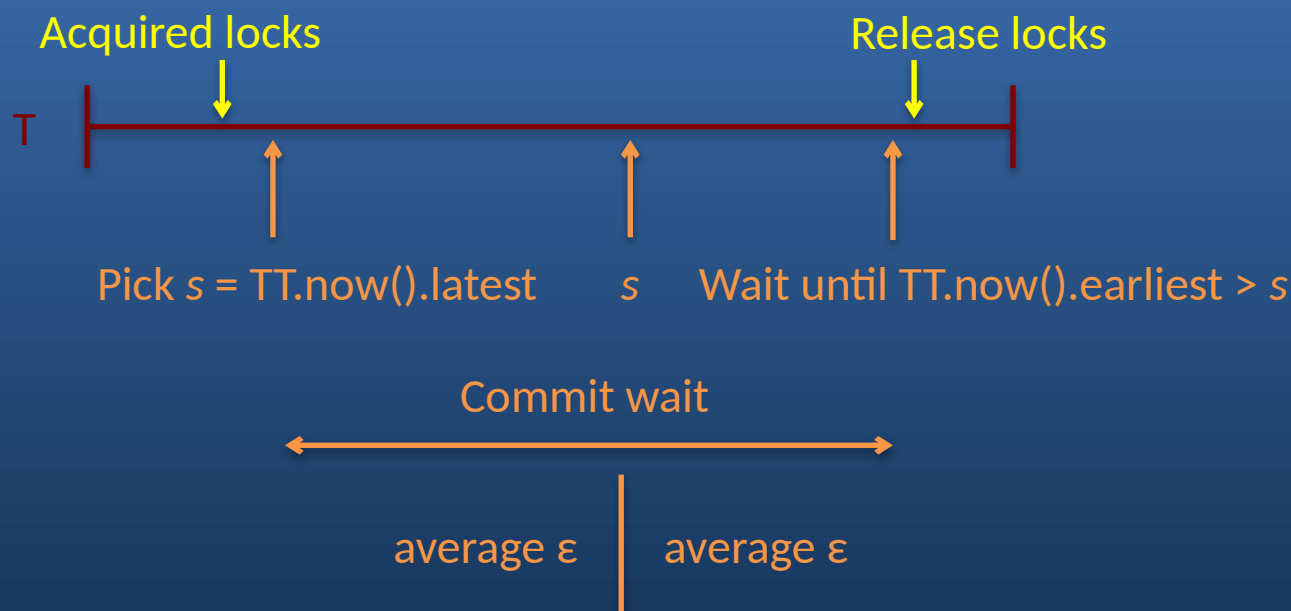
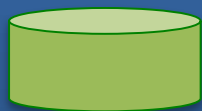


TrueTime

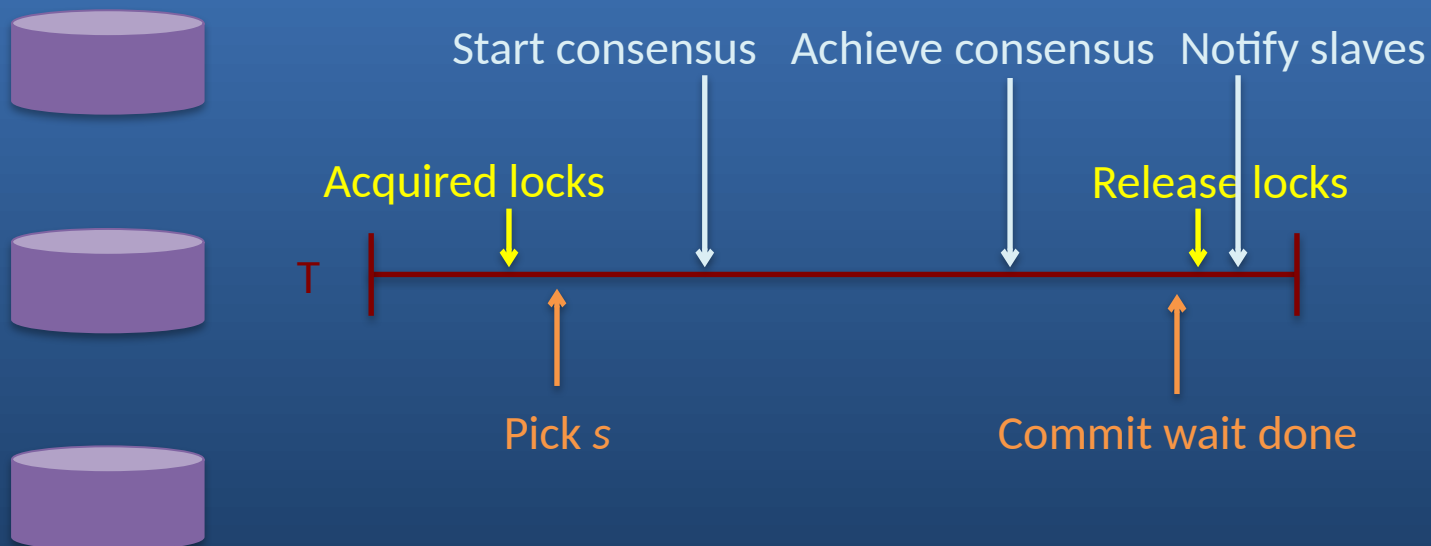
- “Global wall-clock time” with bounded uncertainty



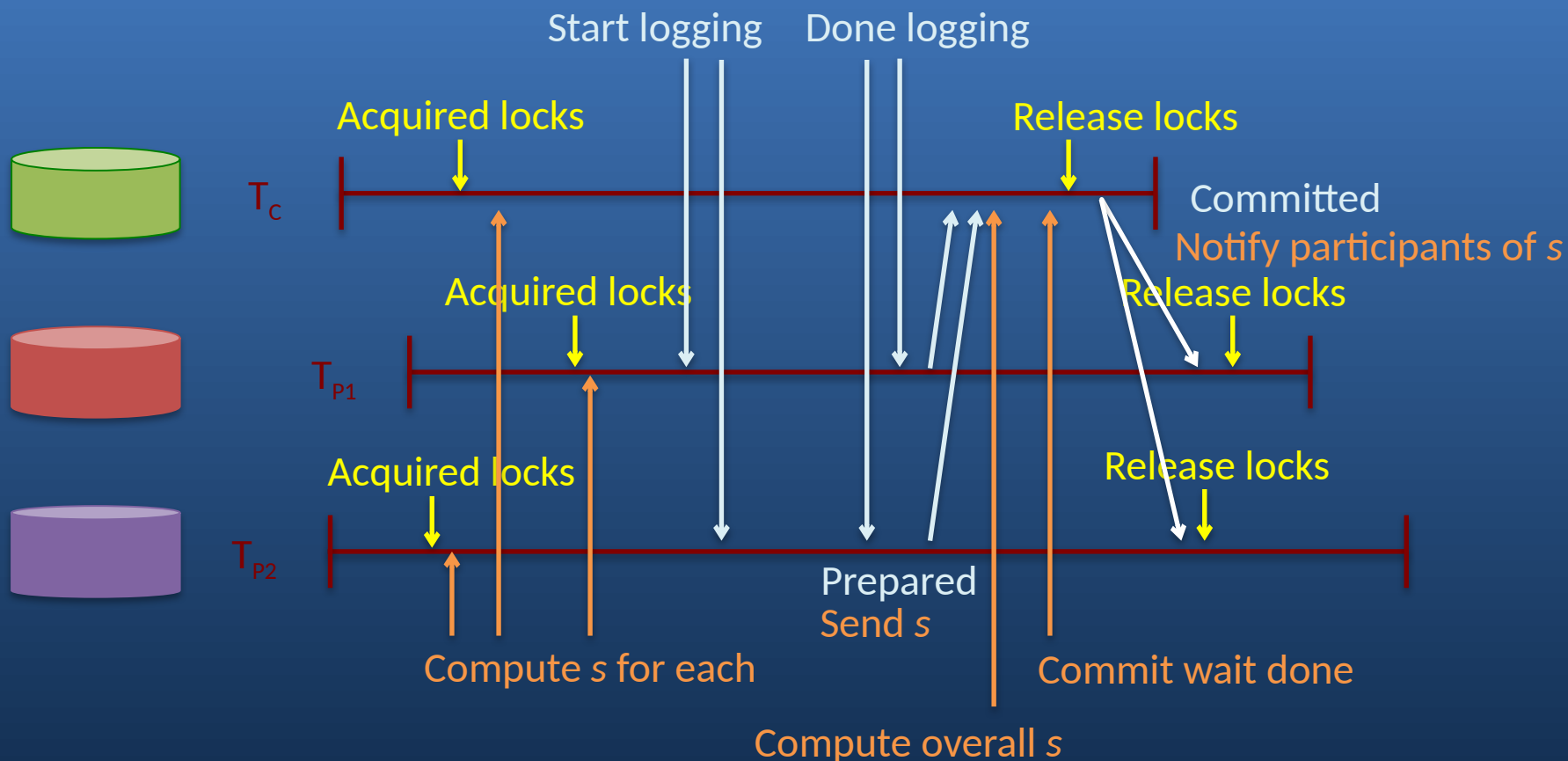
Timestamps and TrueTime



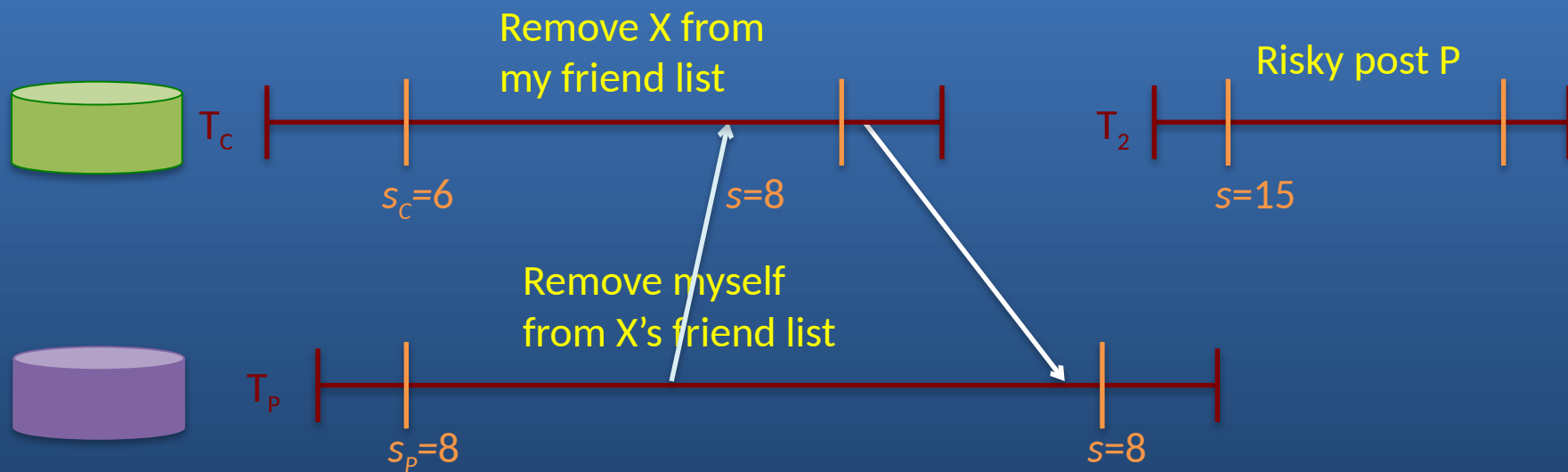
Commit Wait and Replication






Commit Wait and 2-Phase Commit



Example



	Time	<8	8	15
 My friends		[X]	[]	
 My posts				[P]
 X's friends		[me]	[]	

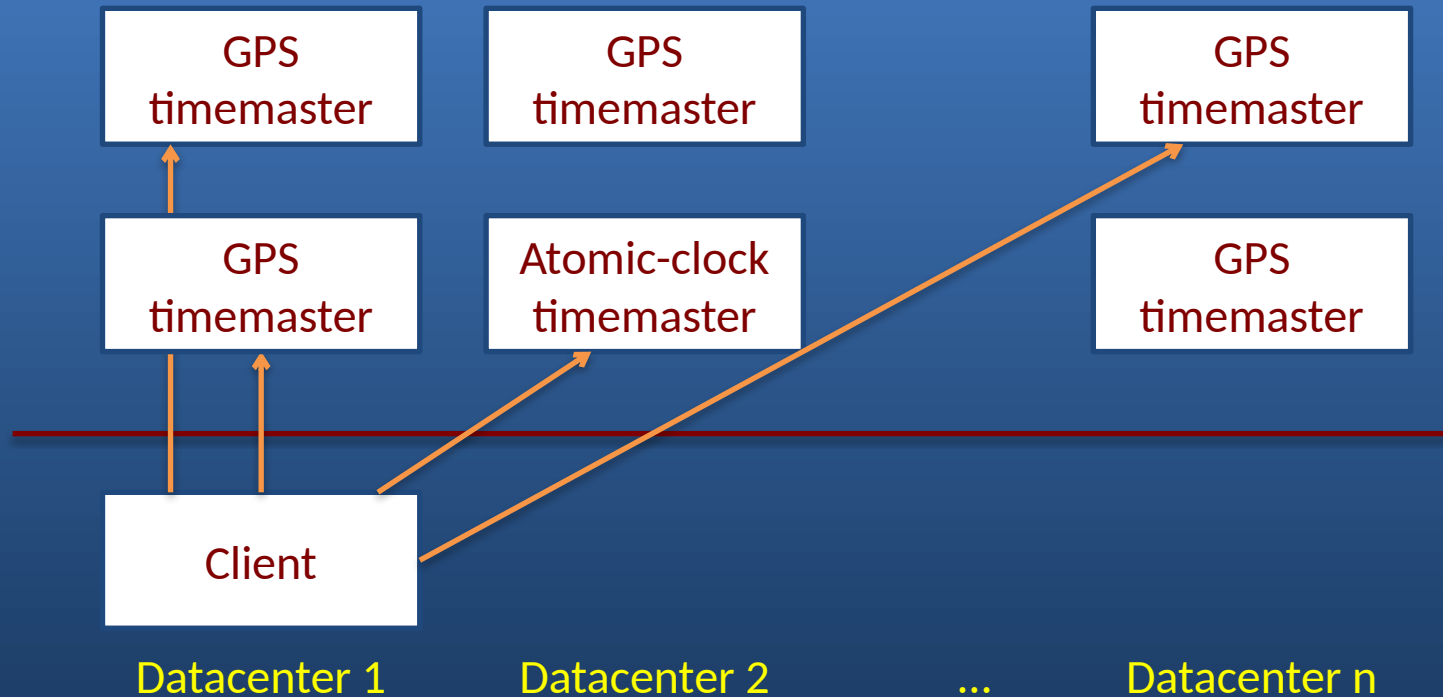
What Have We Covered?

- Lock-free read transactions across datacenters
- External consistency
- Timestamp assignment
- TrueTime
 - Uncertainty in time can be waited out

What Haven't We Covered?

- How to read at the present time
- Atomic schema changes
 - Mostly non-blocking
 - Commit in the future
- Non-blocking reads in the past
 - At any sufficiently up-to-date replica

TrueTime Architecture

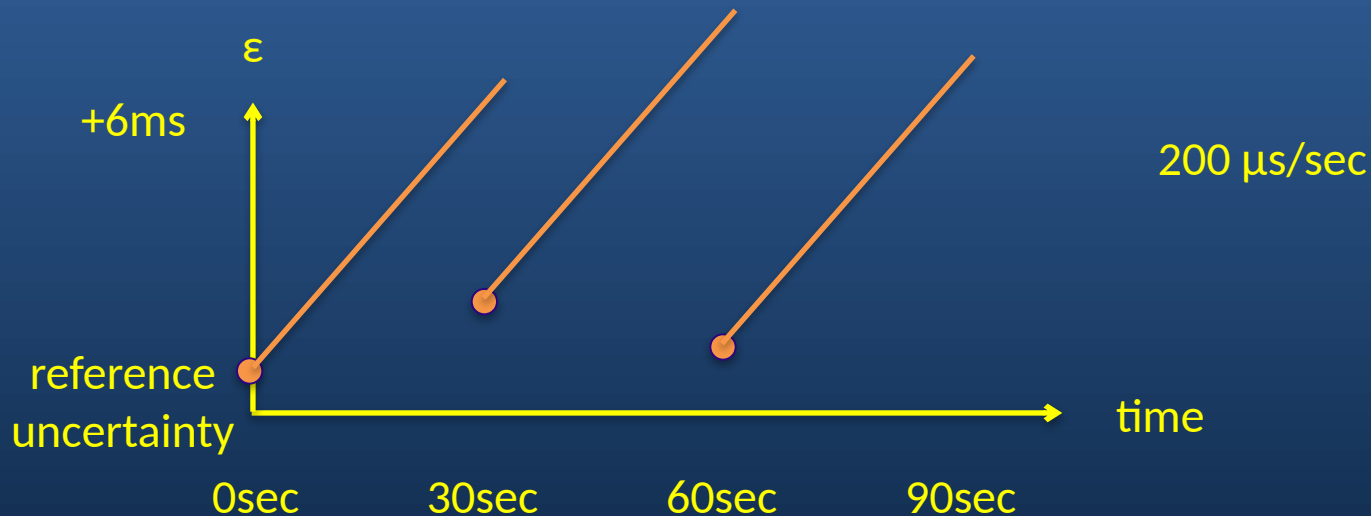


Compute reference [earliest, latest] = now $\pm \epsilon$

TrueTime implementation

$\text{now} = \text{reference now} + \text{local-clock offset}$

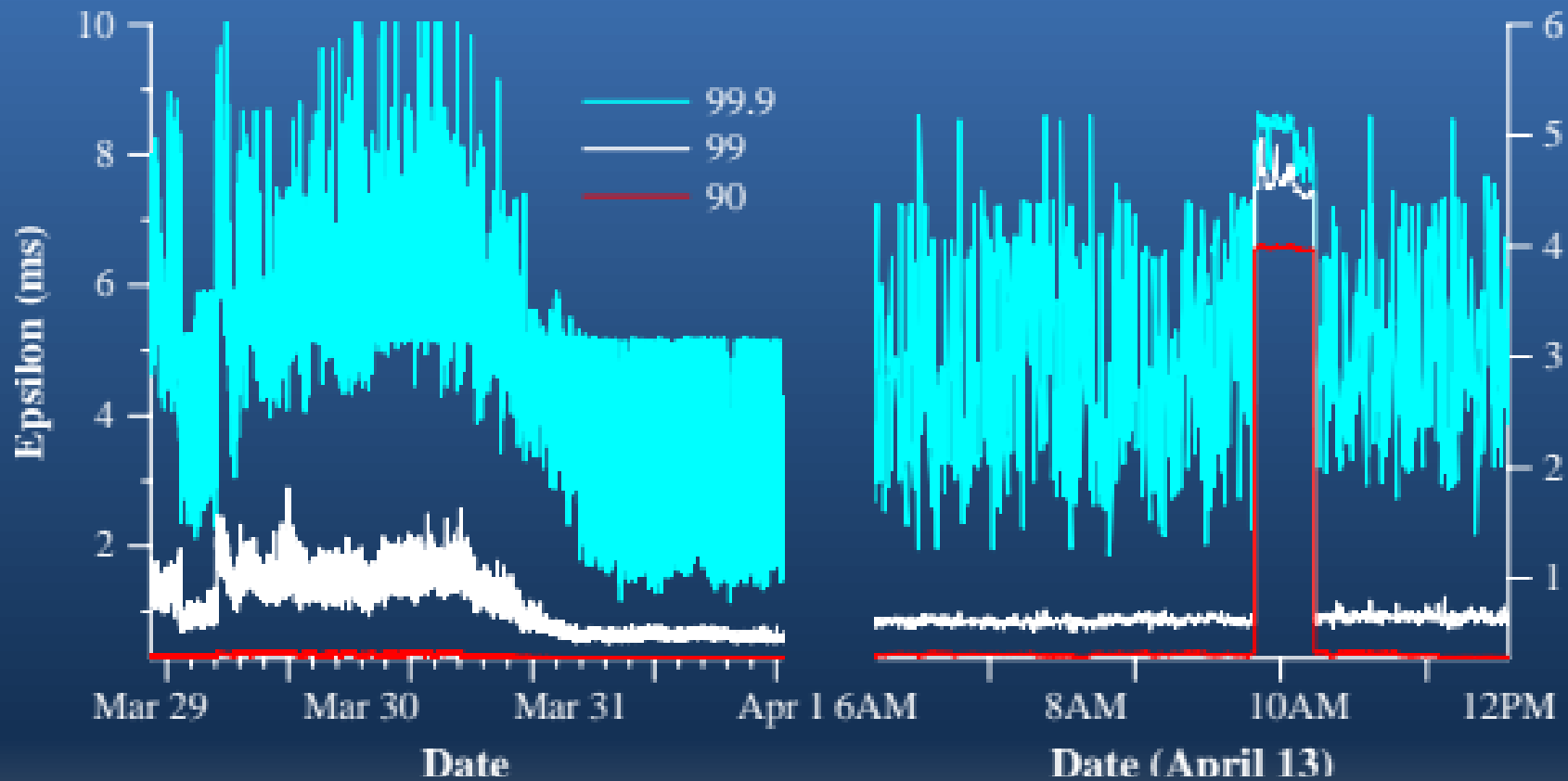
$\varepsilon = \text{reference } \varepsilon + \text{worst-case local-clock drift}$



What If a Clock Goes Rogue?

- Timestamp assignment would violate external consistency
- Empirically unlikely based on 1 year of data
 - Bad CPUs 6 times more likely than bad clocks

Network-Induced Uncertainty



What's in the Literature

- External consistency/linearizability
- Distributed databases
- Concurrency control
- Replication
- Time (NTP, Marzullo)

Future Work

- Improving TrueTime
 - Lower $\varepsilon < 1$ ms
- Building out database features
 - Finish implementing basic features
 - Efficiently support rich query patterns

Conclusions

- Reify clock uncertainty in time APIs
 - Known unknowns are better than unknown unknowns
 - Rethink algorithms to make use of uncertainty
- Stronger semantics are achievable
 - Greater scale \neq weaker semantics

Thanks

- To the Spanner team and customers
 - To our shepherd and reviewers
 - To lots of Googlers for feedback
 - To you for listening!
-
- Questions?