

CMP 334 (2/27/19)

Quiz 2 (truth tables, formulas, circuits)

Signed binary arithmetic (review)

Ripple-carry adders

ALU building block circuits:

Inverters, decoders, multiplexers

Condition flags and comparisons

Comparisons and conditional branches

TOY assembly language

HW 7 ($W \leftarrow X + Y + Z$)

Relative conditional branch op: **bc**

Quiz 2 #1 solution

a	b	c	f
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Quiz 2 #1 solution

a	b	c	f	
0	0	0	0	\overline{abc}
0	0	1	0	$\overline{ab}c$
0	1	0	0	$\overline{a}b\overline{c}$
0	1	1	1	$\overline{a}bc$
1	0	0	1	$a\overline{b}\overline{c}$
1	0	1	0	$a\overline{b}c$
1	1	0	0	$ab\overline{c}$
1	1	1	1	abc

Quiz 2 #1 solution

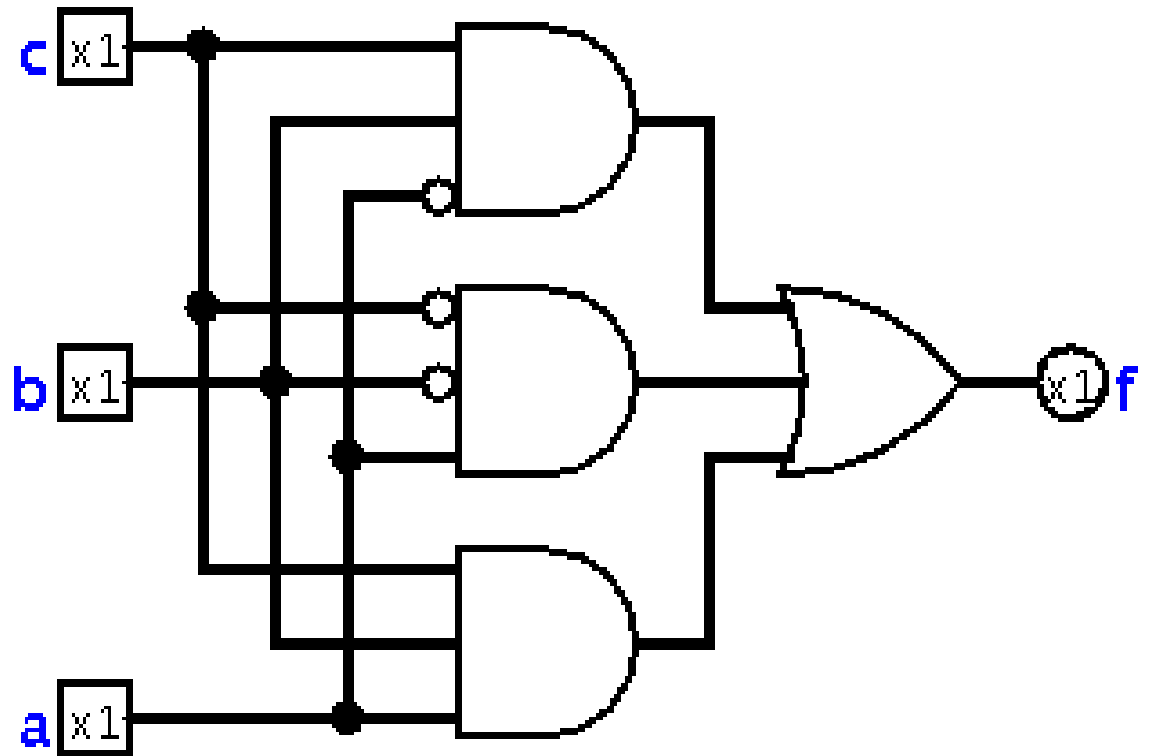
a	b	c	f	
0	0	0	0	$\overline{a}bc$
0	0	1	0	$\overline{a}b\overline{c}$
0	1	0	0	$\overline{a}b\overline{c}$
0	1	1	1	$\overline{a}bc$
1	0	0	1	$a\overline{b}\overline{c}$
1	0	1	0	$a\overline{b}c$
1	1	0	0	$ab\overline{c}$
1	1	1	1	abc

$$f = \overline{a}bc + a\overline{b}\overline{c} + abc$$

Quiz 2 #1 solution

a	b	c	f	
0	0	0	0	$\overline{a}\overline{b}\overline{c}$
0	0	1	0	$\overline{a}\overline{b}c$
0	1	0	0	$\overline{a}b\overline{c}$
0	1	1	1	$\overline{a}bc$
1	0	0	1	$a\overline{b}\overline{c}$
1	0	1	0	$a\overline{b}c$
1	1	0	0	$ab\overline{c}$
1	1	1	1	abc

$$f = \overline{a}bc + a\overline{b}c + abc$$



Quiz 2 #2 solution

n_3	n_2	n_1	n_0
0	0	0	0
0	0	0	1
0	0	1	0
0	0	1	1
0	1	0	0
0	1	0	1
0	1	1	0
0	1	1	1
1	0	0	0
1	0	0	1
1	0	1	0
1	0	1	1
1	1	0	0
1	1	0	1
1	1	1	0
1	1	1	1

Quiz 2 #2 solution

n_3	n_2	n_1	n_0	n
0	0	0	0	0
0	0	0	1	1
0	0	1	0	2
0	0	1	1	3
0	1	0	0	4
0	1	0	1	5
0	1	1	0	6
0	1	1	1	7
1	0	0	0	8
1	0	0	1	9
1	0	1	0	A
1	0	1	1	B
1	1	0	0	C
1	1	0	1	D
1	1	1	0	E
1	1	1	1	F

Quiz 2 #2 solution

n_3	n_2	n_1	n_0	n	$n//5$
0	0	0	0	0	0
0	0	0	1	1	0
0	0	1	0	2	0
0	0	1	1	3	0
0	1	0	0	4	0
0	1	0	1	5	1
0	1	1	0	6	1
0	1	1	1	7	1
1	0	0	0	8	1
1	0	0	1	9	1
1	0	1	0	A	2
1	0	1	1	B	2
1	1	0	0	C	2
1	1	0	1	D	2
1	1	1	0	E	2
1	1	1	1	F	3

Quiz 2 #2 solution

n_3	n_2	n_1	n_0	n	$n//5$	$n\%5$
0	0	0	0	0	0	0
0	0	0	1	1	0	1
0	0	1	0	2	0	2
0	0	1	1	3	0	3
0	1	0	0	4	0	4
0	1	0	1	5	1	0
0	1	1	0	6	1	1
0	1	1	1	7	1	2
1	0	0	0	8	1	3
1	0	0	1	9	1	4
1	0	1	0	A	2	0
1	0	1	1	B	2	1
1	1	0	0	C	2	2
1	1	0	1	D	2	3
1	1	1	0	E	2	4
1	1	1	1	F	3	0

Quiz 2 #2 solution

n_3	n_2	n_1	n_0	n	$n//5$	$n\%5$	$n\%5=4$
0	0	0	0	0	0	0	0
0	0	0	1	1	0	1	0
0	0	1	0	2	0	2	0
0	0	1	1	3	0	3	0
0	1	0	0	4	0	4	1
0	1	0	1	5	1	0	0
0	1	1	0	6	1	1	0
0	1	1	1	7	1	2	0
1	0	0	0	8	1	3	0
1	0	0	1	9	1	4	1
1	0	1	0	A	2	0	0
1	0	1	1	B	2	1	0
1	1	0	0	C	2	2	0
1	1	0	1	D	2	3	0
1	1	1	0	E	2	4	1
1	1	1	1	F	3	0	0

Quiz 2 #2 solution

n_3	n_2	n_1	n_0	n	$n//5$	$n\%5$	$n\%5=4$	
0	0	0	0	0	0	0	0	$\bar{n}_3\bar{n}_2\bar{n}_1\bar{n}_0$
0	0	0	1	1	0	1	0	$\bar{n}_3\bar{n}_2\bar{n}_1n_0$
0	0	1	0	2	0	2	0	$\bar{n}_3\bar{n}_2n_1\bar{n}_0$
0	0	1	1	3	0	3	0	$\bar{n}_3\bar{n}_2n_1n_0$
0	1	0	0	4	0	4	1	$\bar{n}_3n_2\bar{n}_1\bar{n}_0$
0	1	0	1	5	1	0	0	$\bar{n}_3n_2\bar{n}_1n_0$
0	1	1	0	6	1	1	0	$\bar{n}_3n_2n_1\bar{n}_0$
0	1	1	1	7	1	2	0	$\bar{n}_3n_2n_1n_0$
1	0	0	0	8	1	3	0	$n_3\bar{n}_2\bar{n}_1\bar{n}_0$
1	0	0	1	9	1	4	1	$n_3\bar{n}_2\bar{n}_1n_0$
1	0	1	0	A	2	0	0	$n_3\bar{n}_2n_1\bar{n}_0$
1	0	1	1	B	2	1	0	$n_3\bar{n}_2n_1n_0$
1	1	0	0	C	2	2	0	$n_3n_2\bar{n}_1\bar{n}_0$
1	1	0	1	D	2	3	0	$n_3n_2\bar{n}_1n_0$
1	1	1	0	E	2	4	1	$n_3n_2n_1\bar{n}_0$
1	1	1	1	F	3	0	0	$n_3n_2n_1n_0$

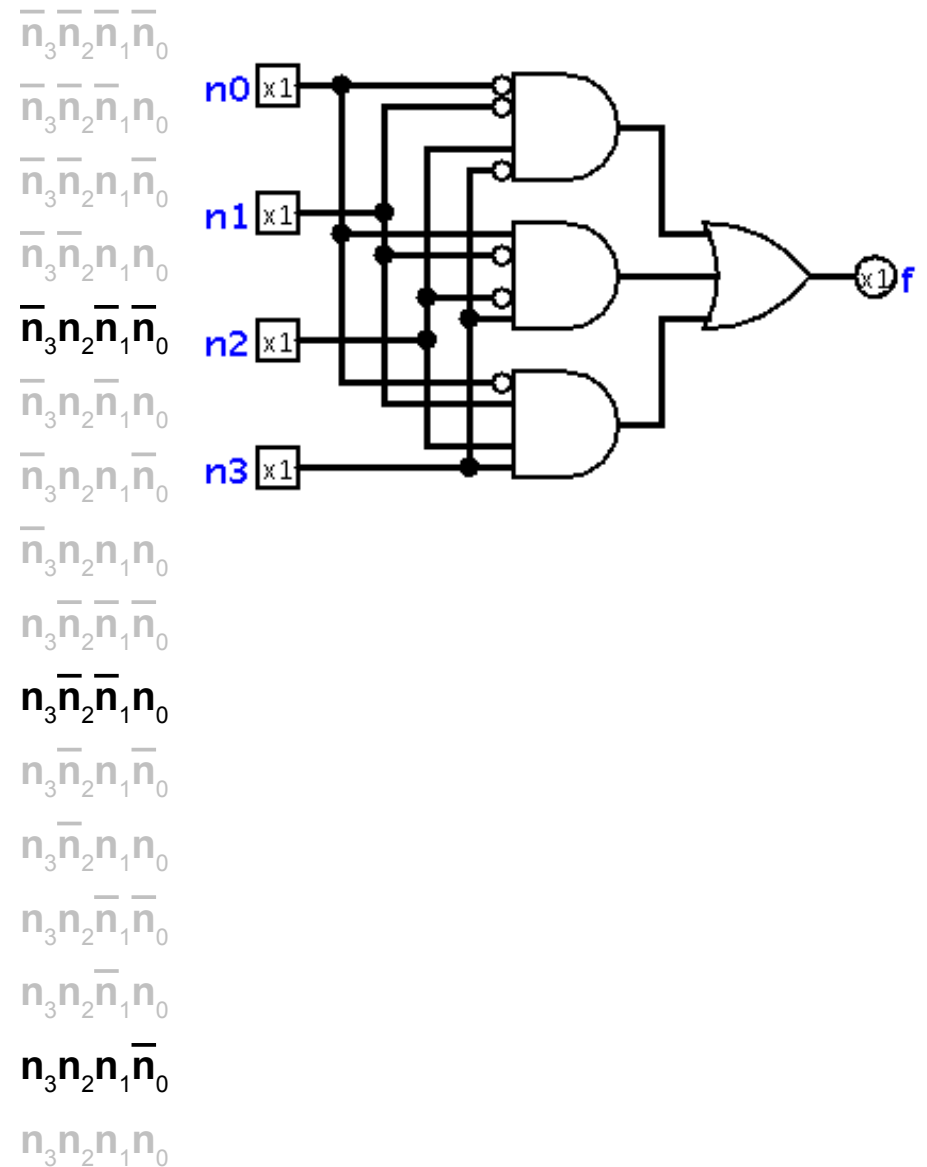
Quiz 2 #2 solution

n_3	n_2	n_1	n_0	n	n//5	n%5	n%5=4	$f = \bar{n}_3 n_2 \bar{n}_1 \bar{n}_0 + n_3 \bar{n}_2 \bar{n}_1 n_0 + n_3 n_2 n_1 \bar{n}_0$
0	0	0	0	0	0	0	0	$\bar{n}_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$
0	0	0	1	1	0	1	0	$\bar{n}_3 \bar{n}_2 \bar{n}_1 n_0$
0	0	1	0	2	0	2	0	$\bar{n}_3 \bar{n}_2 n_1 \bar{n}_0$
0	0	1	1	3	0	3	0	$\bar{n}_3 \bar{n}_2 n_1 n_0$
0	1	0	0	4	0	4	1	$\bar{n}_3 n_2 \bar{n}_1 \bar{n}_0$
0	1	0	1	5	1	0	0	$\bar{n}_3 n_2 \bar{n}_1 n_0$
0	1	1	0	6	1	1	0	$\bar{n}_3 n_2 n_1 \bar{n}_0$
0	1	1	1	7	1	2	0	$\bar{n}_3 n_2 n_1 n_0$
1	0	0	0	8	1	3	0	$n_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$
1	0	0	1	9	1	4	1	$n_3 \bar{n}_2 \bar{n}_1 n_0$
1	0	1	0	A	2	0	0	$n_3 \bar{n}_2 n_1 \bar{n}_0$
1	0	1	1	B	2	1	0	$n_3 \bar{n}_2 n_1 n_0$
1	1	0	0	C	2	2	0	$n_3 n_2 \bar{n}_1 \bar{n}_0$
1	1	0	1	D	2	3	0	$n_3 n_2 \bar{n}_1 n_0$
1	1	1	0	E	2	4	1	$n_3 n_2 n_1 \bar{n}_0$
1	1	1	1	F	3	0	0	$n_3 n_2 n_1 n_0$

Quiz 2 #2 solution

n_3	n_2	n_1	n_0	n	n//5	n%5	n%5=4
0	0	0	0	0	0	0	0
0	0	0	1	1	0	1	0
0	0	1	0	2	0	2	0
0	0	1	1	3	0	3	0
0	1	0	0	4	0	4	1
0	1	0	1	5	1	0	0
0	1	1	0	6	1	1	0
0	1	1	1	7	1	2	0
1	0	0	0	8	1	3	0
1	0	0	1	9	1	4	1
1	0	1	0	A	2	0	0
1	0	1	1	B	2	1	0
1	1	0	0	C	2	2	0
1	1	0	1	D	2	3	0
1	1	1	0	E	2	4	1
1	1	1	1	F	3	0	0

$$f = \bar{n}_3 n_2 \bar{n}_1 \bar{n}_0 + n_3 \bar{n}_2 \bar{n}_1 n_0 + n_3 n_2 n_1 \bar{n}_0$$



Quiz 2 take aways

n_3	n_2	n_1	n_0
0	0	0	0
0	0	0	1
0	0	1	0
0	0	1	1
0	1	0	0
0	1	0	1
0	1	1	0
0	1	1	1
1	0	0	0
1	0	0	1
1	0	1	0
1	0	1	1
1	1	0	0
1	1	0	1
1	1	1	0
1	1	1	1

Quiz 2 take aways

n_3	n_2	n_1	n_0	$n_3 n_2 n_1 n_0$
0	0	0	0	$\bar{n}_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$
0	0	0	1	$\bar{n}_3 \bar{n}_2 \bar{n}_1 n_0$
0	0	1	0	$\bar{n}_3 \bar{n}_2 n_1 \bar{n}_0$
0	0	1	1	$\bar{n}_3 \bar{n}_2 n_1 n_0$
0	1	0	0	$\bar{n}_3 n_2 \bar{n}_1 \bar{n}_0$
0	1	0	1	$\bar{n}_3 n_2 \bar{n}_1 n_0$
0	1	1	0	$\bar{n}_3 n_2 n_1 \bar{n}_0$
0	1	1	1	$\bar{n}_3 n_2 n_1 n_0$
1	0	0	0	$n_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$
1	0	0	1	$n_3 \bar{n}_2 \bar{n}_1 n_0$
1	0	1	0	$n_3 \bar{n}_2 n_1 \bar{n}_0$
1	0	1	1	$n_3 \bar{n}_2 n_1 n_0$
1	1	0	0	$n_3 n_2 \bar{n}_1 \bar{n}_0$
1	1	0	1	$n_3 n_2 \bar{n}_1 n_0$
1	1	1	0	$n_3 n_2 n_1 \bar{n}_0$
1	1	1	1	$n_3 n_2 n_1 n_0$

Quiz 2 take aways

n_3	n_2	n_1	n_0	$n_3 n_2 n_1 n_0$	n
0	0	0	0	$\bar{n}_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	0
0	0	0	1	$\bar{n}_3 \bar{n}_2 \bar{n}_1 n_0$	1
0	0	1	0	$\bar{n}_3 \bar{n}_2 n_1 \bar{n}_0$	2
0	0	1	1	$\bar{n}_3 \bar{n}_2 n_1 n_0$	3
0	1	0	0	$\bar{n}_3 n_2 \bar{n}_1 \bar{n}_0$	4
0	1	0	1	$\bar{n}_3 n_2 \bar{n}_1 n_0$	5
0	1	1	0	$\bar{n}_3 n_2 n_1 \bar{n}_0$	6
0	1	1	1	$\bar{n}_3 n_2 n_1 n_0$	7
1	0	0	0	$n_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	8
1	0	0	1	$n_3 \bar{n}_2 \bar{n}_1 n_0$	9
1	0	1	0	$n_3 \bar{n}_2 n_1 \bar{n}_0$	A
1	0	1	1	$n_3 \bar{n}_2 n_1 n_0$	B
1	1	0	0	$n_3 n_2 \bar{n}_1 \bar{n}_0$	C
1	1	0	1	$n_3 n_2 \bar{n}_1 n_0$	D
1	1	1	0	$n_3 n_2 n_1 \bar{n}_0$	E
1	1	1	1	$n_3 n_2 n_1 n_0$	F

Quiz 2 take aways

n_3	n_2	n_1	n_0	$n_3 n_2 n_1 n_0$	n
0	0	0	0	$\bar{n}_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	0
0	0	0	1	$\bar{n}_3 \bar{n}_2 \bar{n}_1 n_0$	1
0	0	1	0	$\bar{n}_3 \bar{n}_2 n_1 \bar{n}_0$	2
0	0	1	1	$\bar{n}_3 \bar{n}_2 n_1 n_0$	3
0	1	0	0	$\bar{n}_3 n_2 \bar{n}_1 \bar{n}_0$	4
0	1	0	1	$\bar{n}_3 n_2 \bar{n}_1 n_0$	5
0	1	1	0	$\bar{n}_3 n_2 n_1 \bar{n}_0$	6
0	1	1	1	$\bar{n}_3 n_2 n_1 n_0$	7
1	0	0	0	$n_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	8
1	0	0	1	$n_3 \bar{n}_2 \bar{n}_1 n_0$	9
1	0	1	0	$n_3 \bar{n}_2 n_1 \bar{n}_0$	A
1	0	1	1	$n_3 \bar{n}_2 n_1 n_0$	B
1	1	0	0	$n_3 n_2 \bar{n}_1 \bar{n}_0$	C
1	1	0	1	$n_3 n_2 \bar{n}_1 n_0$	D
1	1	1	0	$n_3 n_2 n_1 \bar{n}_0$	E
1	1	1	1	$n_3 n_2 n_1 n_0$	F

$$n = n_3 \cdot 2^3 + n_2 \cdot 2^2 + n_1 \cdot 2^1 + n_0 \cdot 2^0$$

Quiz 2 take aways

n_3	n_2	n_1	n_0	$n_3 n_2 n_1 n_0$	n	s
0	0	0	0	$\bar{n}_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	0	0
0	0	0	1	$\bar{n}_3 \bar{n}_2 \bar{n}_1 n_0$	1	1
0	0	1	0	$\bar{n}_3 \bar{n}_2 n_1 \bar{n}_0$	2	2
0	0	1	1	$\bar{n}_3 \bar{n}_2 n_1 n_0$	3	3
0	1	0	0	$\bar{n}_3 n_2 \bar{n}_1 \bar{n}_0$	4	4
0	1	0	1	$\bar{n}_3 n_2 \bar{n}_1 n_0$	5	5
0	1	1	0	$\bar{n}_3 n_2 n_1 \bar{n}_0$	6	6
0	1	1	1	$\bar{n}_3 n_2 n_1 n_0$	7	7
1	0	0	0	$n_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	8	-8
1	0	0	1	$n_3 \bar{n}_2 \bar{n}_1 n_0$	9	-7
1	0	1	0	$n_3 \bar{n}_2 n_1 \bar{n}_0$	A	-6
1	0	1	1	$n_3 \bar{n}_2 n_1 n_0$	B	-5
1	1	0	0	$n_3 n_2 \bar{n}_1 \bar{n}_0$	C	-4
1	1	0	1	$n_3 n_2 \bar{n}_1 n_0$	D	-3
1	1	1	0	$n_3 n_2 n_1 \bar{n}_0$	E	-2
1	1	1	1	$n_3 n_2 n_1 n_0$	F	-1

$$n = n_3 \cdot 2^3 + n_2 \cdot 2^2 + n_1 \cdot 2^1 + n_0 \cdot 2^0$$

Quiz 2 take aways

n_3	n_2	n_1	n_0	$n_3 n_2 n_1 n_0$	n	s
0	0	0	0	$\bar{n}_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	0	0
0	0	0	1	$\bar{n}_3 \bar{n}_2 \bar{n}_1 n_0$	1	1
0	0	1	0	$\bar{n}_3 \bar{n}_2 n_1 \bar{n}_0$	2	2
0	0	1	1	$\bar{n}_3 \bar{n}_2 n_1 n_0$	3	3
0	1	0	0	$\bar{n}_3 n_2 \bar{n}_1 \bar{n}_0$	4	4
0	1	0	1	$\bar{n}_3 n_2 \bar{n}_1 n_0$	5	5
0	1	1	0	$\bar{n}_3 n_2 n_1 \bar{n}_0$	6	6
0	1	1	1	$\bar{n}_3 n_2 n_1 n_0$	7	7
1	0	0	0	$n_3 \bar{n}_2 \bar{n}_1 \bar{n}_0$	8	-8
1	0	0	1	$n_3 \bar{n}_2 \bar{n}_1 n_0$	9	-7
1	0	1	0	$n_3 \bar{n}_2 n_1 \bar{n}_0$	A	-6
1	0	1	1	$n_3 \bar{n}_2 n_1 n_0$	B	-5
1	1	0	0	$n_3 n_2 \bar{n}_1 \bar{n}_0$	C	-4
1	1	0	1	$n_3 n_2 \bar{n}_1 n_0$	D	-3
1	1	1	0	$n_3 n_2 n_1 \bar{n}_0$	E	-2
1	1	1	1	$n_3 n_2 n_1 n_0$	F	-1

$$n = n_3 \cdot 2^3 + n_2 \cdot 2^2 + n_1 \cdot 2^1 + n_0 \cdot 2^0$$

$$s = -n_3 \cdot 2^3 + n_2 \cdot 2^2 + n_1 \cdot 2^1 + n_0 \cdot 2^0$$

HW 8: Two's Complement Circuit

Use the four step *Combinational Circuit Design Process* presented in class to design circuits that take a 3-bit unsigned binary integer **X** as input and produces as output a 3-bit unsigned binary integer **Y** that is the two's complement of **X**.

$$\mathbf{Y} = \mathbf{2}^3 - \mathbf{X} = \overline{\mathbf{X}} + \mathbf{1}$$

Do not minimize the circuits for **Y** as a part of this assignments.

Extra credit: minimize the circuits for **Y**.

N-Bit Integers (N = 8)

Unsigned Integers

...	0100000000
...	0011111111
...	0010000000
...	0001111111
...	0000000101
...	0000000100
...	0000000011
...	0000000010
...	0000000001
...	0000000000
...	1111111111
...	1111111110
...	1111111101
...	1111111100
...	1111111011
...	1110000000
...	1101111111
...	1100000000
...	1011111111

$$= 2^N$$

$$= 2^N - 1$$

$$= 2^{N-1}$$

$$= 2^{N-1} - 1$$

$$= 4 = 2^2$$

$$= 3 = 2^2 - 1$$

$$= 2 = 2^1$$

$$= 1 = 2^0 = 2^1 - 1$$

$$= 0 = 2^0 - 1$$

$$= -1 = -2^0$$

$$= -2 = -2^1$$

$$= -3 = -2^1 - 1$$

$$= -4 = -2^2$$

$$= -5 = -2^2 - 1$$

$$= -2^{N-1}$$

$$= -2^{N-1} - 1$$

$$= -2^N$$

$$= -2^N - 1$$

Signed Integers

Additive Inverse

.001000	+8
.000111	+7
.000110	+6
.000101	+5
.000100	+4
.000011	+3
.000010	+2
.000001	+1
.000000	+0
.111111	-1
.111110	-2
.111101	-3
.111100	-4
.111011	-5
.111010	-6
.111001	-7
.111000	-8

THEOREM:

If X and Y are signed integers then

$$X - Y = X + \bar{Y} + 1$$

Corollary:

$$-Y = \bar{Y} + 1$$

$$- .0101 = \overline{.0101} + 1 = .1010 + 1 = .1011$$

$$- .0001 = \overline{.0001} + 1 = .1110 + 1 = .1111$$

$$- .0000 = \overline{.0000} + 1 = .1111 + 1 = .0000$$

$$- .1111 = \overline{.1111} + 1 = .0000 + 1 = .0001$$

$$- .1011 = \overline{.1011} + 1 = .0100 + 1 = .0101$$

n-bit Binary Numbers

Unsigned: $b_{n-1} b_{n-2} \dots b_1 b_0$ ($b_i = 0$ or $b_i = 1$)

value: $b_{n-1}2^{n-1} + b_{n-2}2^{n-2} + \dots + b_12^1 + b_02^0$

range: $[0 \dots 2^n - 1]$

n-bit sum: $A \oplus B = A + B - c \cdot 2^n$

n-bit diff: $A \ominus B \equiv A \oplus (2^n - B) = A \oplus \overline{B} + 1 = A - B + 2^n - c \cdot 2^n$

Signed: $b_{n-1} b_{n-2} \dots b_1 b_0$ ($b_i = 0$ or $b_i = 1$)

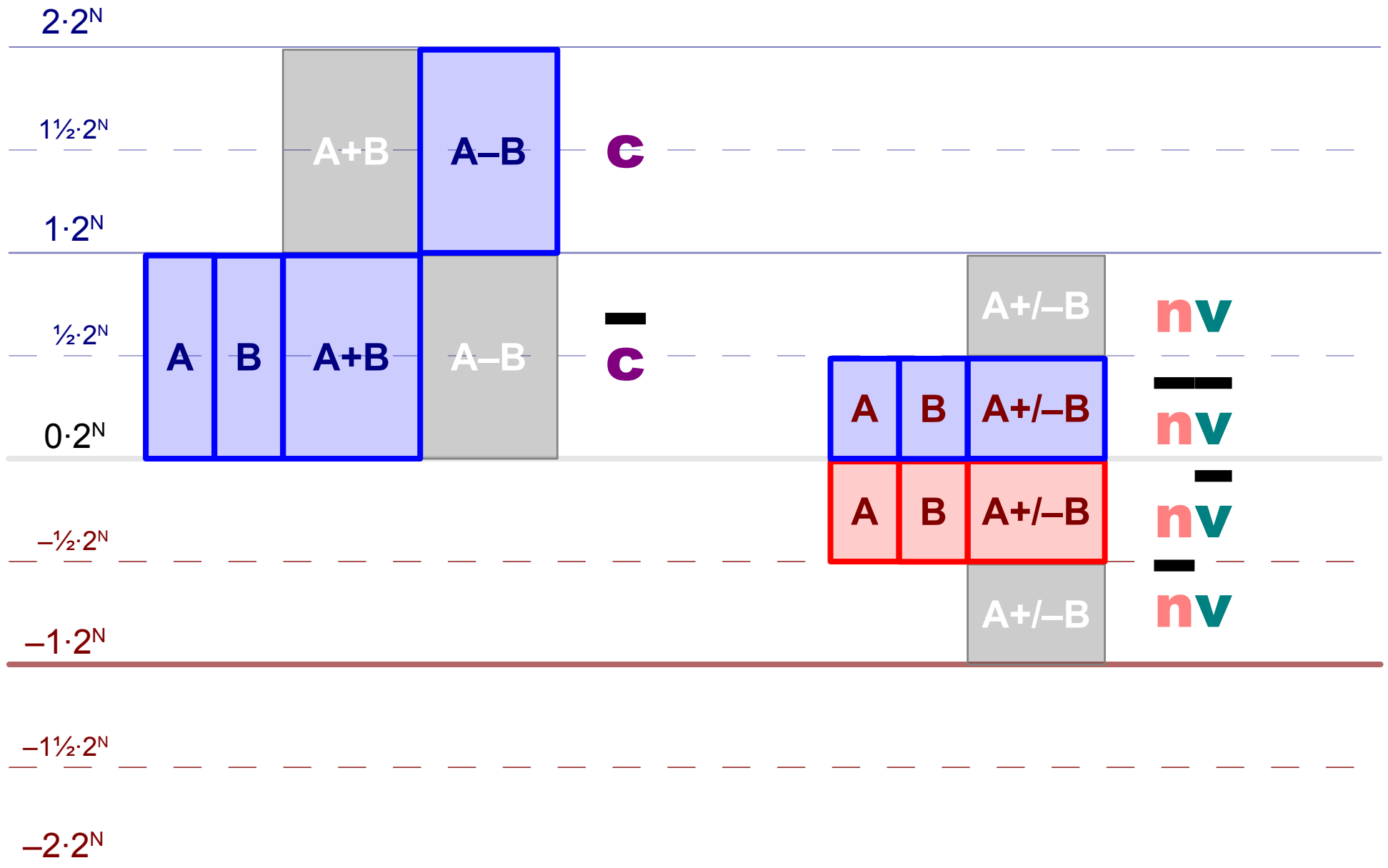
Value: $-b_{n-1}2^{n-1} + b_{n-2}2^{n-2} + \dots + b_12^1 + b_02^0$

Range: $[-2^{n-1} \dots 2^{n-1} - 1]$

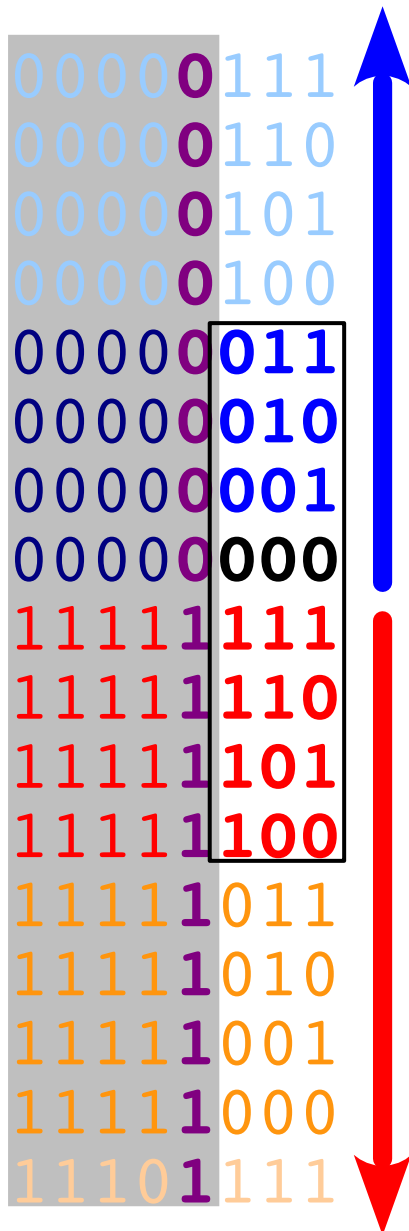
n-bit sum: $A \oplus B = A + B$ iff $v=0$

n-bit diff: $A \ominus B = A - B$ iff $v=0$

Condition Flags



3-Bit Signed Binary Integers



0000	0111
0000	0110
0000	0101
0000	0100
0000	0011
0000	0010
0000	0001
0000	0000
1111	1111
1111	1110
1111	1101
1111	1100
1111	1011
1111	1010
1111	1001
1111	1000
1110	1111

8 supported values: **-4** .. **3**

Sign: high bit: **0** positive, **1** negative

Not closed under addition

incorrect results: **-8** .. **-5**, **4** .. **6**

Not closed under subtraction

incorrect results: **-7** .. **-5**, **4** .. **7**

Correct: $X \oplus Y = Z$, $X \oplus Y = Z$

$X \oplus Y$, $X \oplus Y$ always correct. Why?

Incorrect: $X \oplus Y = Z$, $X \oplus Y = Z$

$X \ominus Y = Z$, $X \ominus Y = Z$

Honesty Criteria

The n-bit result **r** of a binary operation on n-bit values **a** and **b** is **honest** (**deceptive**) if it is *the same as* (*different from*) the whole number result of the same operation on the same values.

(n-bit) unsigned addition is *honest* iff (c = 0)

Carry flag is not set

(n-bit) unsigned subtraction is *honest* iff (c = 1)

Carry flag is set

(n-bit) signed addition is *honest* iff (v = 0)

a and **b** have different signs or **a**, **b**, and **r** have same sign

(n-bit) signed subtraction is *honest* iff (v = 0)

a and $\overline{\overline{\overline{\overline{\mathbf{b}}}}}$ have different signs or **a**, $\overline{\overline{\overline{\overline{\mathbf{b}}}}}$, and **r** have same sign

HW 9: Signed Binary Arithmetic

For each of the $\langle X, Y \rangle$ pairs in the table below:

- Convert X and $Y \rightarrow$ binary
- Compute $X+Y$ (the 8-bit **sum**)
- Compute \bar{Y} (the 2's complement of Y)
- Compute $X-Y \equiv X+\bar{Y}$ (the 8-bit **difference**)
- Indicate the signs of X , Y , $X+Y$, \bar{Y} , and $X-Y$
- Convert $X+Y$, \bar{Y} , and $X-Y \rightarrow$ hexadecimal
- Indicate condition flag (**z**, **n**, **c**, **v**) values for $X+Y$, $X-Y$
- Is $X+Y$ honest? is $X-Y$ honest?

Where $\langle X, Y \rangle =$

- | | |
|---------------------------------|---------------------------------|
| 1) $\langle 0x4F, 0x6D \rangle$ | 2) $\langle 0xB3, 0x17 \rangle$ |
| 3) $\langle 0xA3, 0x95 \rangle$ | 4) $\langle 0x6E, 0x3A \rangle$ |

Signed Arithmetic: X1 i)

X	Y	X + Y	~Y	X - Y
---	---	-------	----	-------

0x8C

0x6F

10001100

01101111

10001100
 01101111

 00001100
 01111011

01101111
 10010000
 00000001

 10010001

10001100
 10010001

 100011101

0xFB

0x91

0x1D

zncv

(s) honest
(u) honest

zncv

(s) deceptive
(u) honest

Signed Arithmetic: X2 i)

X	Y	X + Y	~Y	X - Y
0x54	0xF3			
01010100	11110011	<div> <div>01010100</div> <div>11110011</div> <hr/> <div>11110000</div> <div>101000111</div> </div>	<div> <div>11110011</div> <hr/> <div>00001100</div> <div>00000001</div> <hr/> <div>00001101</div> </div>	<div> <div>01010100</div> <div>00001101</div> <hr/> <div>00110001</div> </div>
		0x47	0x0D	0x61
		<div> <div>zncv</div> </div>		<div> <div>zncv</div> </div>
		(s) honest (u) deceptive		(s) honest (u) deceptive

Signed Arithmetic: X3 i)

X	Y	X + Y	~Y	X - Y
---	---	-------	----	-------

0x9E

0xCC

10011110

11001100

1	0	0	1	1	1	0	0
1	0	0	1	1	1	0	0
<hr/>							
1	0	1	1	0	1	0	1

1	1	0	0	1	1	0	0
0	0	1	1	0	0	1	1
<hr/>							
0	0	0	0	0	0	1	1
<hr/>							
0	0	1	1	0	1	0	0

1	0	0	1	1	1	0	0
0	0	1	1	0	1	0	0
<hr/>							
0	1	1	0	1	0	0	1

0x6A

0x34

0xD2

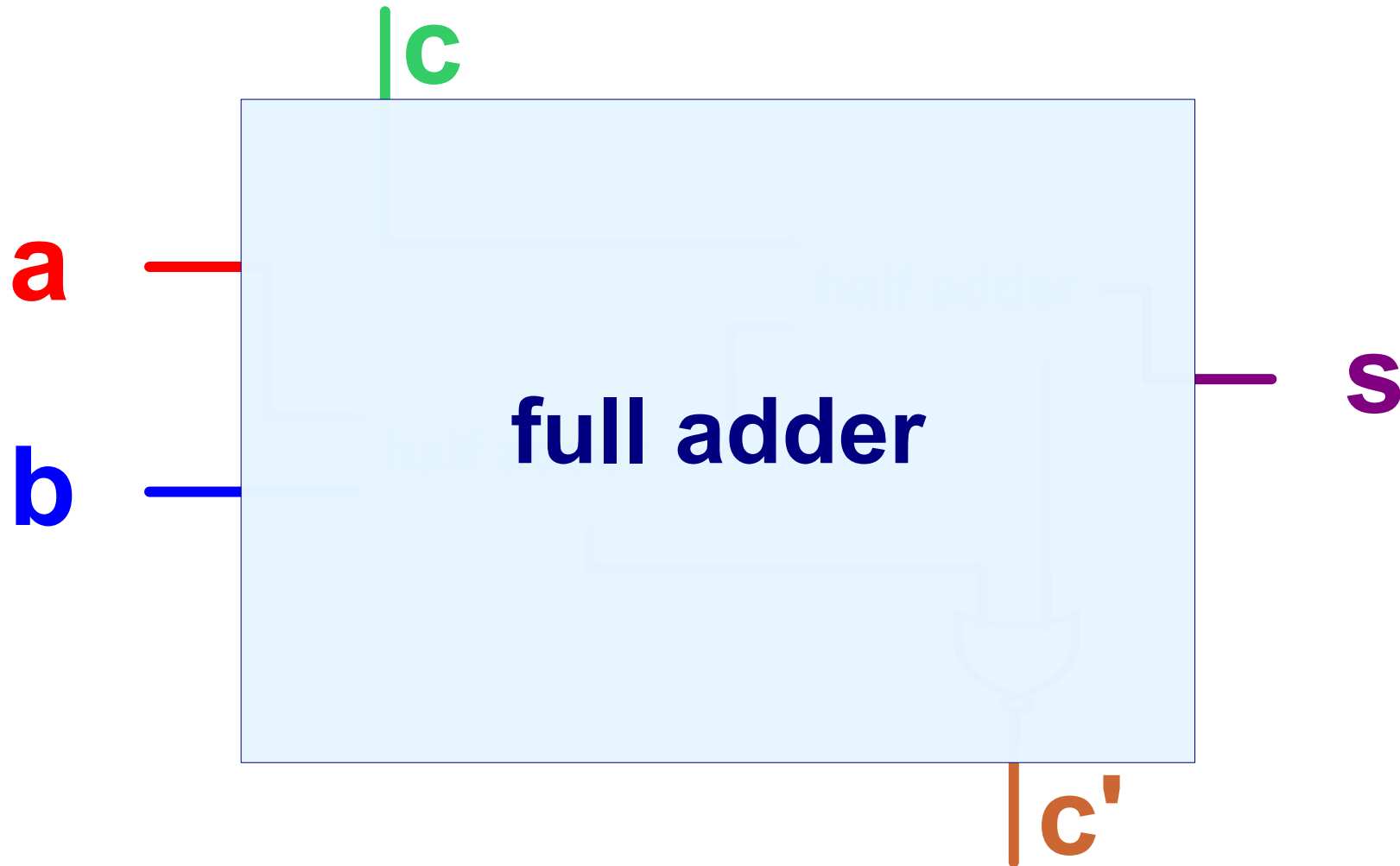
zncv

(s) honest
(u) deceptive

zncv

(s) honest
(u) deceptive

Full Adder Black Box



Sum 3 1-bit inputs to give a 2-bit output

Full Adder Truth Table

The diagram illustrates the eight possible input combinations for a full adder. Each case is shown as a vertical addition problem. The top number is the first operand, the middle number is the second operand, and the bottom number is the sum. The carry-in is indicated by a green '1' or '0' above the second operand, and the carry-out is indicated by a black '1' or '0' to the left of the sum. Arrows show the carry flow from the first sum to the second sum.

$\begin{array}{r} \text{+} \quad \textcolor{red}{1} \\ \textcolor{blue}{1} \\ \hline 1 \quad \textcolor{green}{0} \\ \textbf{0} \end{array}$	$\begin{array}{r} \text{+} \quad \textcolor{red}{1} \\ \textcolor{blue}{0} \\ \hline 0 \quad \textcolor{green}{0} \\ \textbf{1} \end{array}$	$\begin{array}{r} \text{+} \quad \textcolor{red}{0} \\ \textcolor{blue}{1} \\ \hline 0 \quad \textcolor{green}{0} \\ \textbf{1} \end{array}$	$\begin{array}{r} \text{+} \quad \textcolor{red}{0} \\ \textcolor{blue}{0} \\ \hline 0 \quad \textcolor{green}{0} \\ \textbf{0} \end{array}$
$\begin{array}{r} \text{+} \quad \textcolor{red}{1} \\ \textcolor{blue}{1} \\ \hline 1 \quad \textcolor{green}{1} \\ \textbf{1} \end{array}$	$\begin{array}{r} \text{+} \quad \textcolor{red}{1} \\ \textcolor{blue}{0} \\ \hline 1 \quad \textcolor{green}{1} \\ \textbf{0} \end{array}$	$\begin{array}{r} \text{+} \quad \textcolor{red}{0} \\ \textcolor{blue}{1} \\ \hline 1 \quad \textcolor{green}{1} \\ \textbf{0} \end{array}$	$\begin{array}{r} \text{+} \quad \textcolor{red}{0} \\ \textcolor{blue}{0} \\ \hline 0 \quad \textcolor{green}{1} \\ \textbf{1} \end{array}$

carry out

carry in

Full Adder Truth Table

#	a	b	c	c'	s
0	0	0	0	0	0
1	0	0	1	0	1
2	0	1	0	0	1
3	0	1	1	1	0
4	1	0	0	0	1
5	1	0	1	1	0
6	1	1	0	1	0
7	1	1	1	1	1

Full Adder Boolean Formulas

#	a	b	c	c'	c'	s	s
0	0	0	0	0		0	
1	0	0	1	0		1	$\overline{a}bc$
2	0	1	0	0		1	$a\overline{b}c$
3	0	1	1	1	$\overline{a}bc$	0	
4	1	0	0	0		1	$a\overline{b}\overline{c}$
5	1	0	1	1	$a\overline{b}c$	0	
6	1	1	0	1	$ab\overline{c}$	0	
7	1	1	1	1	abc	1	abc

$$s = \overline{a}bc + \overline{a}b\overline{c} + a\overline{b}c + abc$$

$$c' = \overline{a}bc + a\overline{b}c + ab\overline{c} + abc$$

Full Adder Boolean Formulas

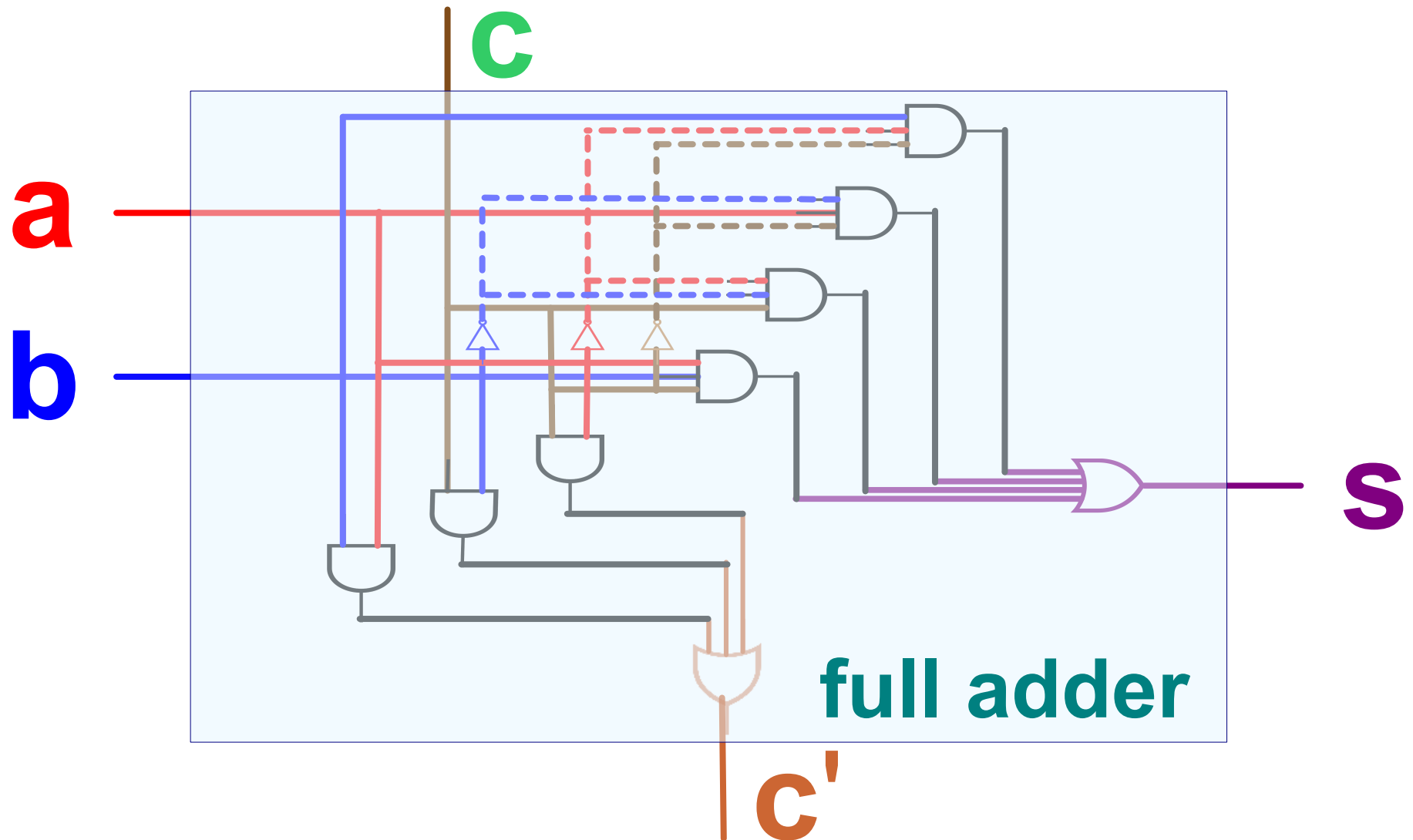
#	a	b	c	c'	s
0	0	0	0	0	0
1	0	0	1	0	1
2	0	1	0	0	1
3	0	1	1	1	0
4	1	0	0	0	1
5	1	0	1	1	0
6	1	1	0	1	0
7	1	1	1	1	1

$$s = \overline{a}bc + a\overline{b}c + a\overline{b}\overline{c} + abc$$

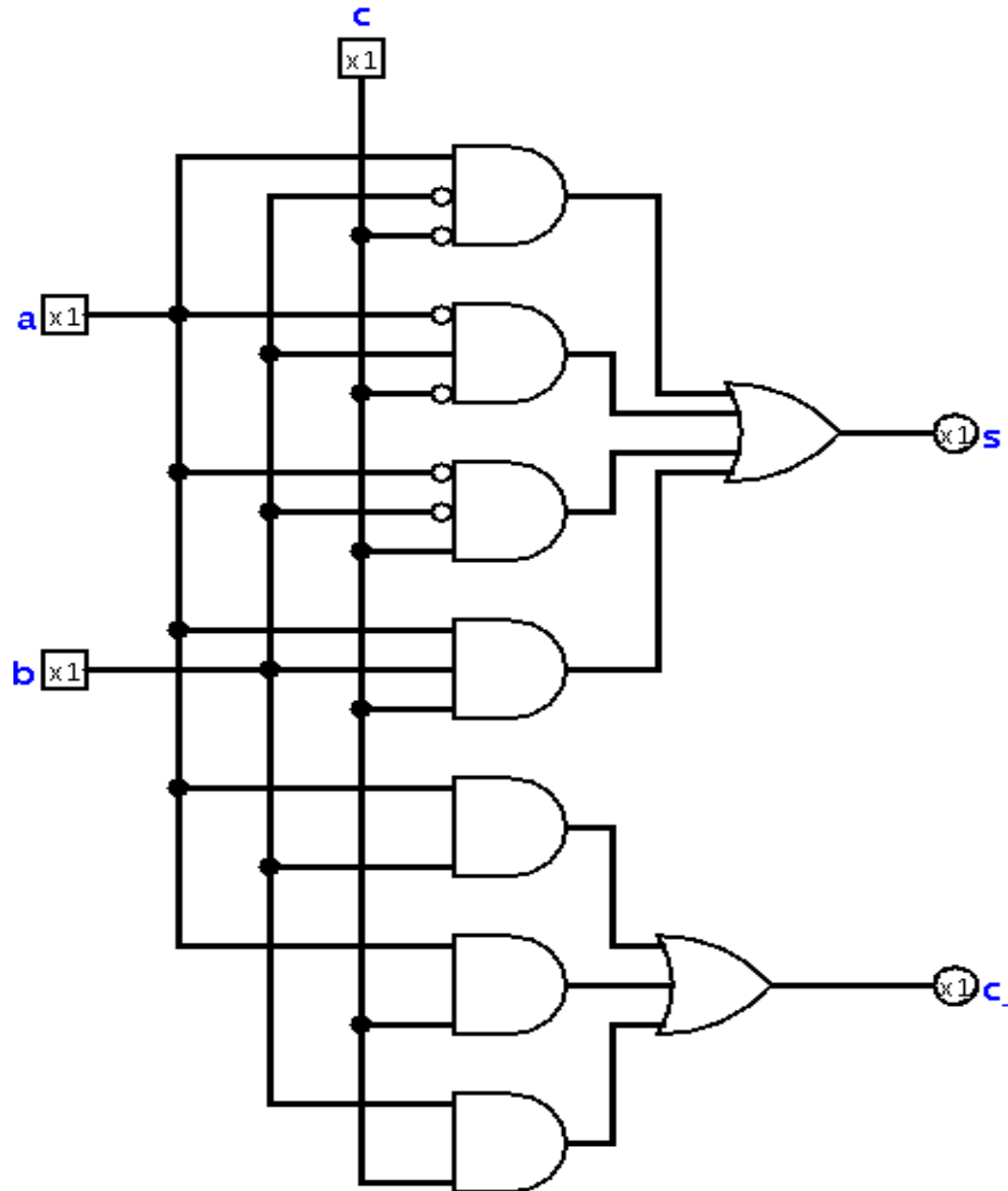
$$c' = bc + ac + ab$$

$$s = abc + \overline{a}bc + \overline{a}\overline{b}c + a\overline{b}c$$

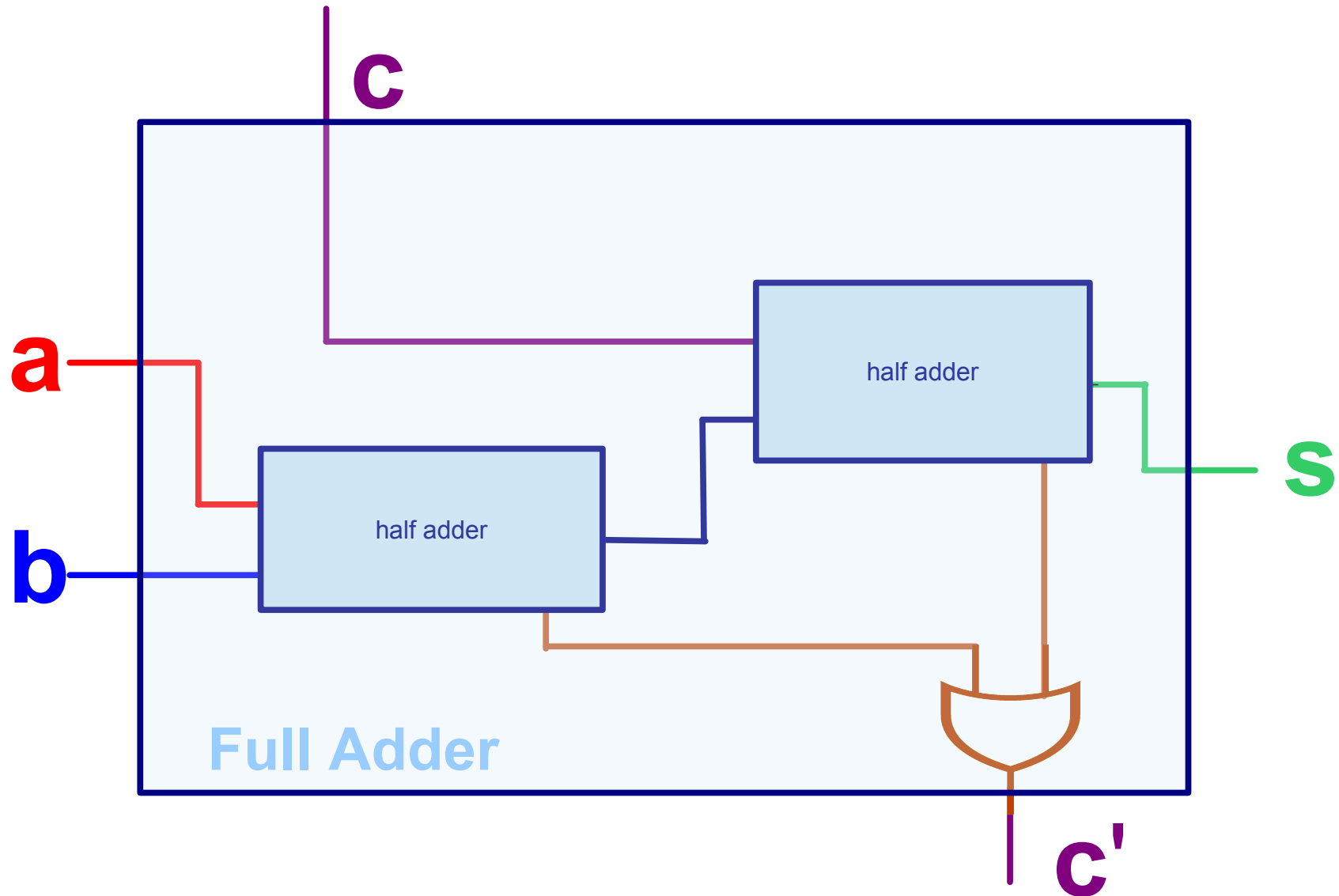
$$c' = ab + ac + bc$$



Full Adder Circuit



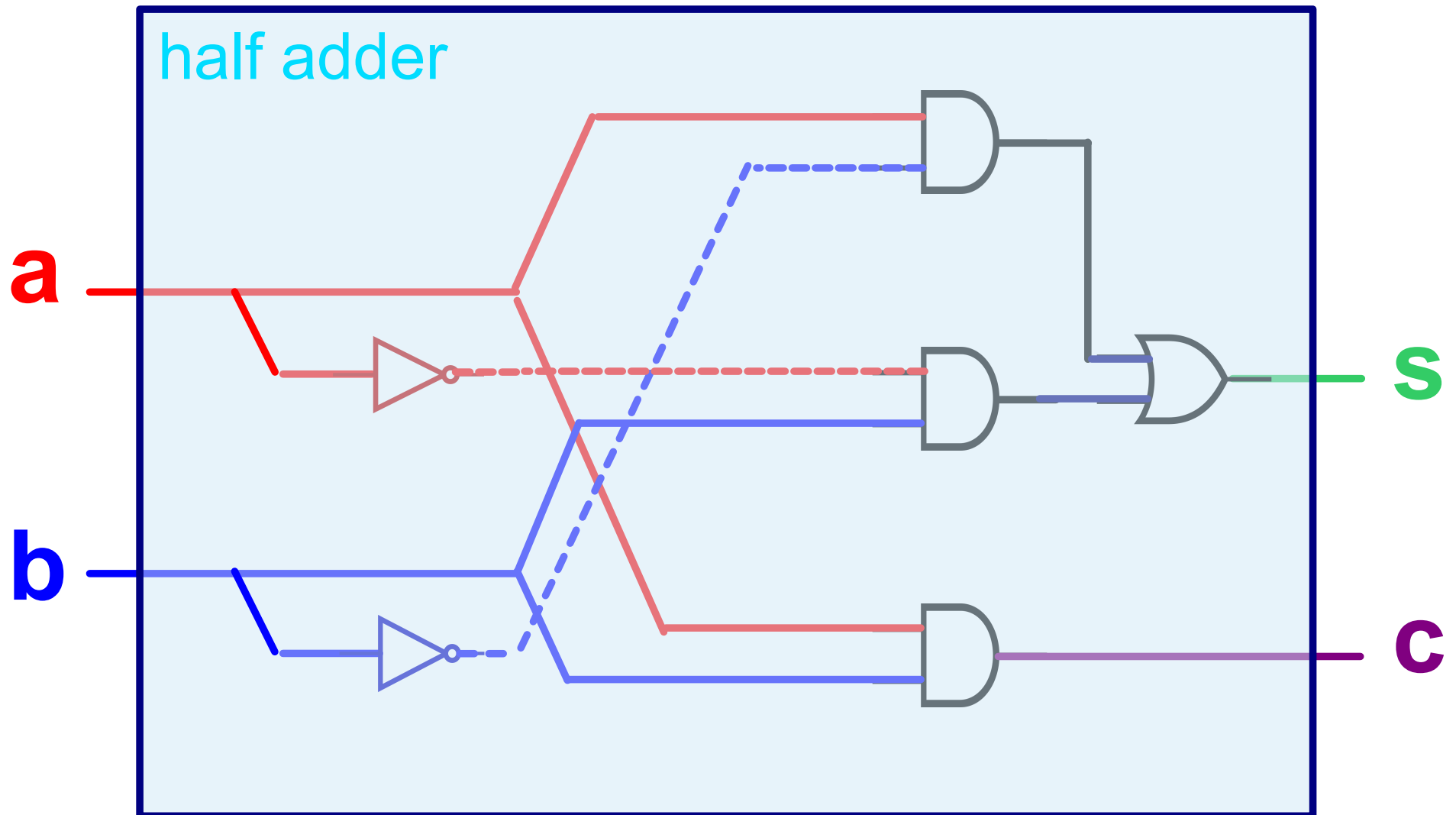
A Full Adder built of Half Adders



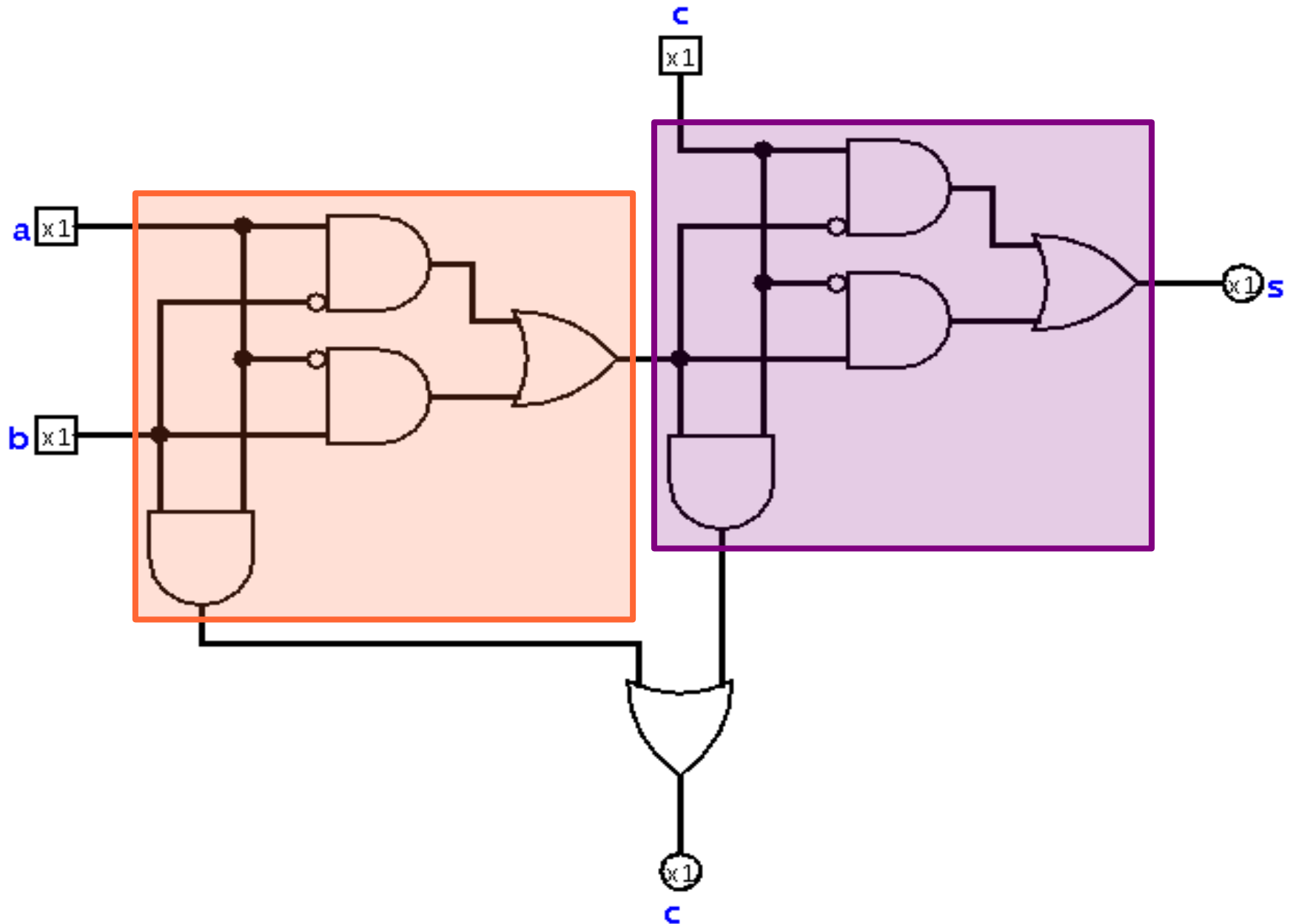
Half Adder Circuit



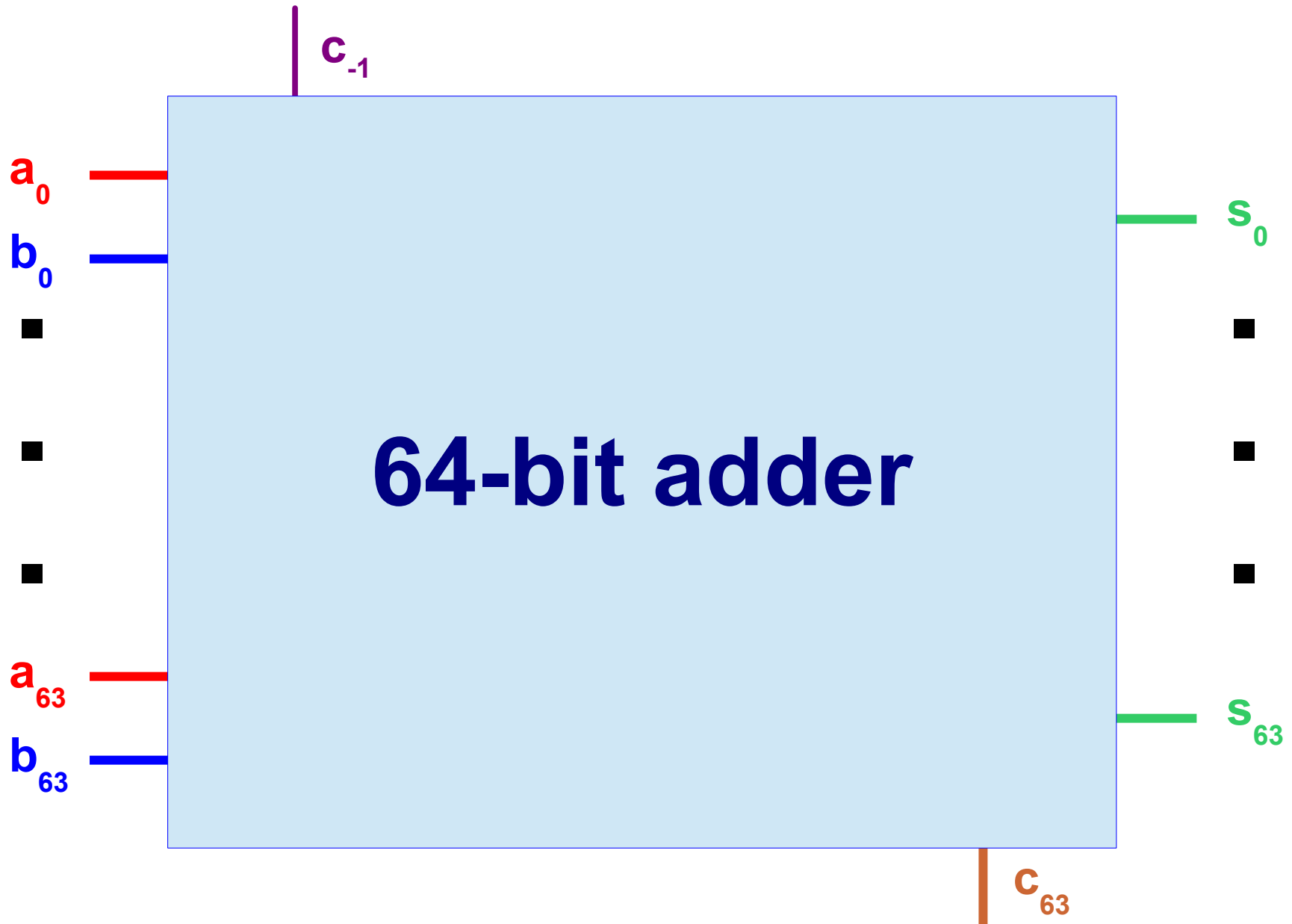
Half Adder Circuit



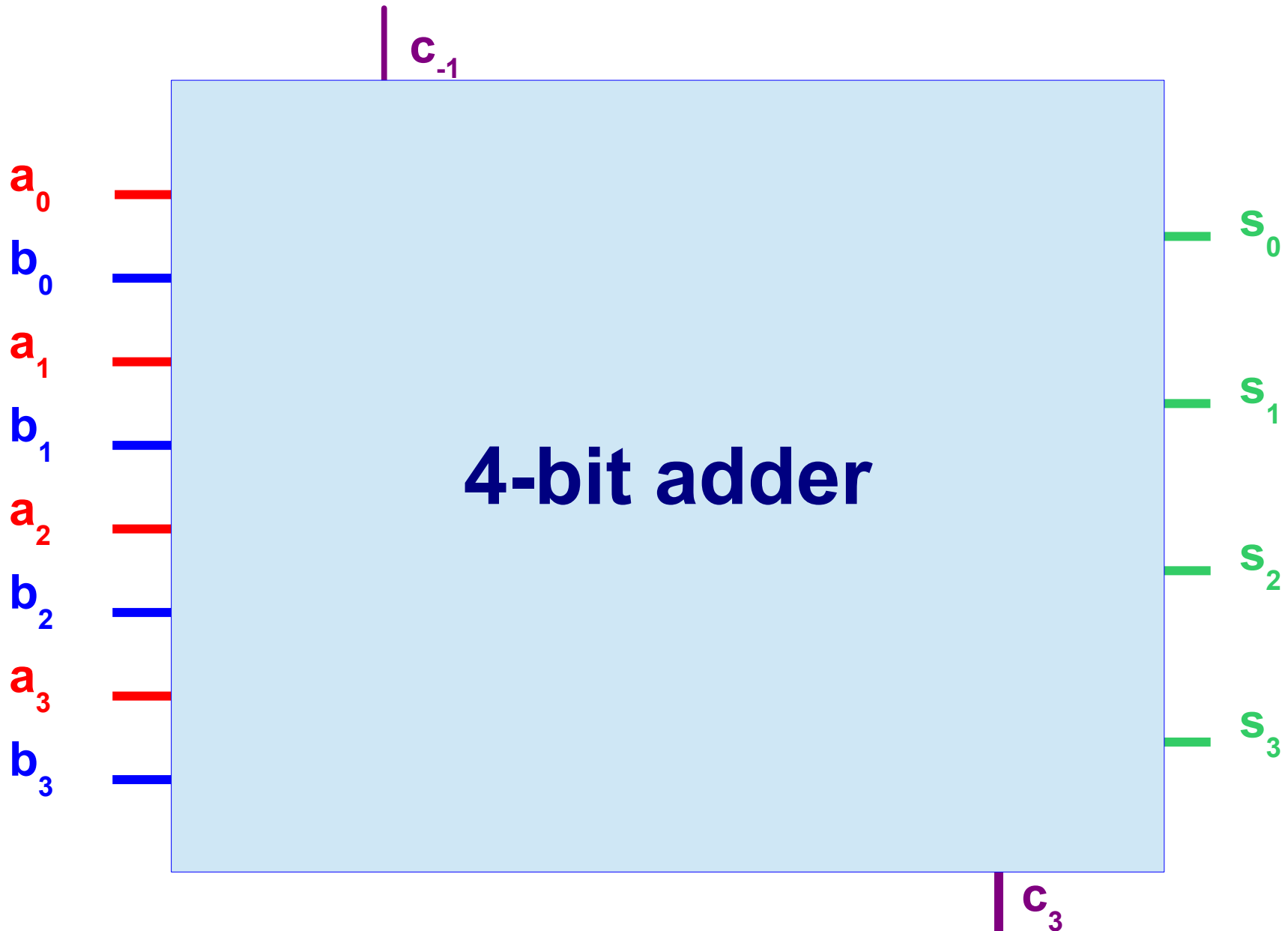
Full Adder built of Half Adders



64-bit Word Adder



4-bit Word Adder



4-Bit Ripple-Carry Addition

	1	1	1	0
+	0	1	1	1
				0

4-Bit Ripple-Carry Addition

	1	1	1	0
+	0	1	1	1
			0	0
				1

4-Bit Ripple-Carry Addition

	1	1	1	0
+	0	1	1	1
		1	0	0
			0	1

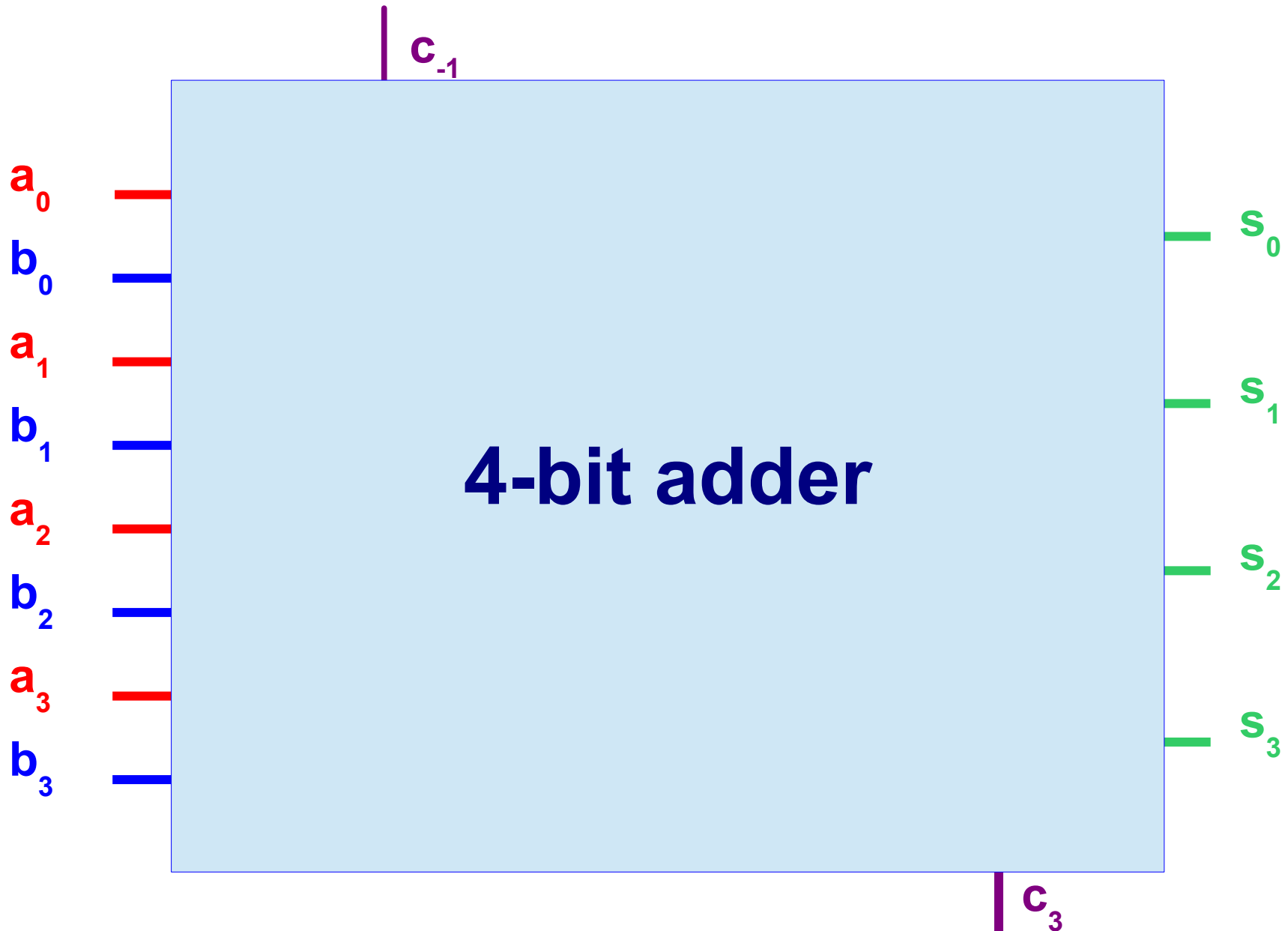
4-Bit Ripple-Carry Addition

	1	1	1	0
+	0	1	1	1
	1	1	0	0
		1	0	1

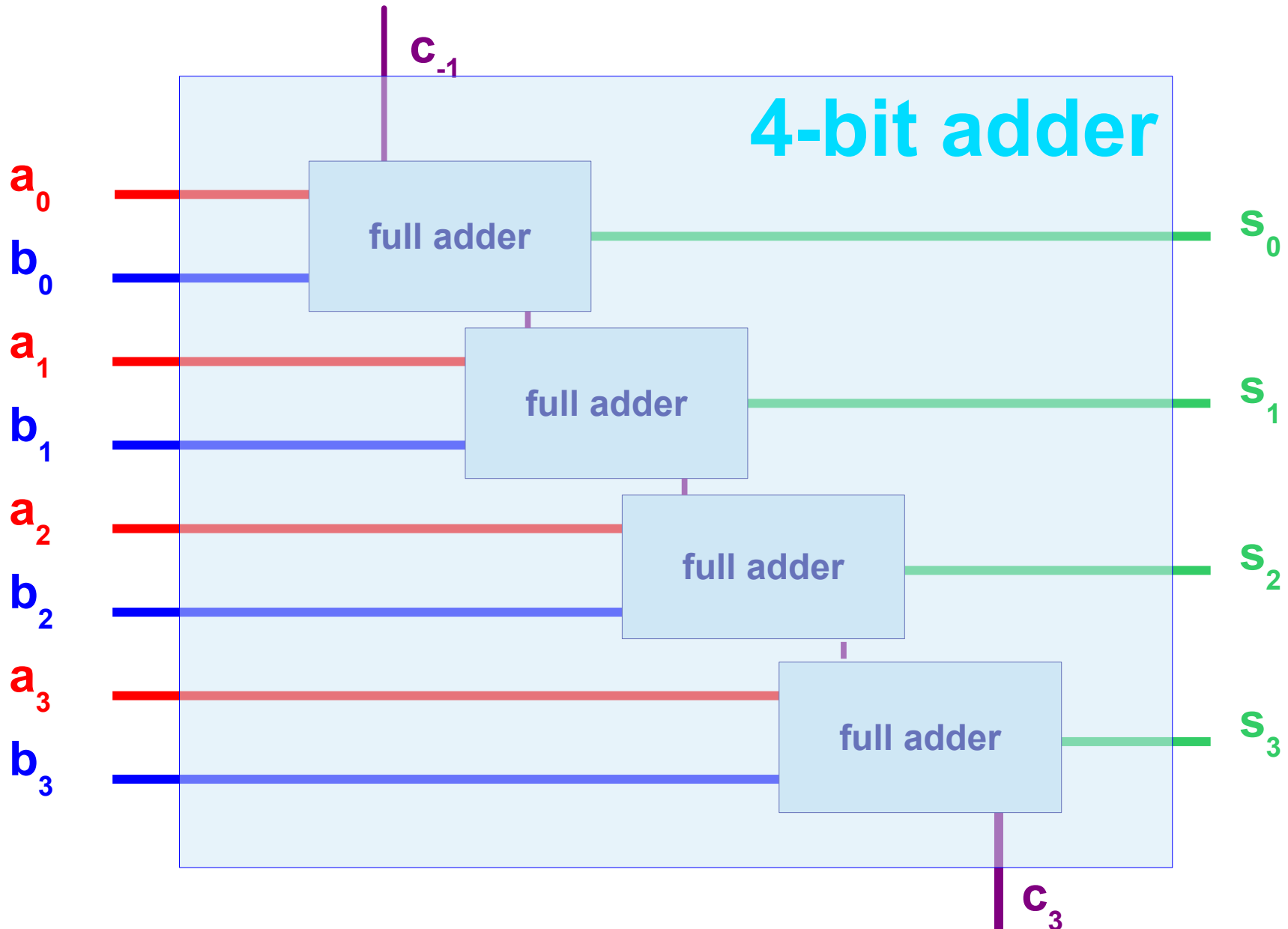
4-Bit Ripple-Carry Addition

		1	1	1	0
		0	1	1	1
+					
	1	1	1	0	0
		0	1	0	1

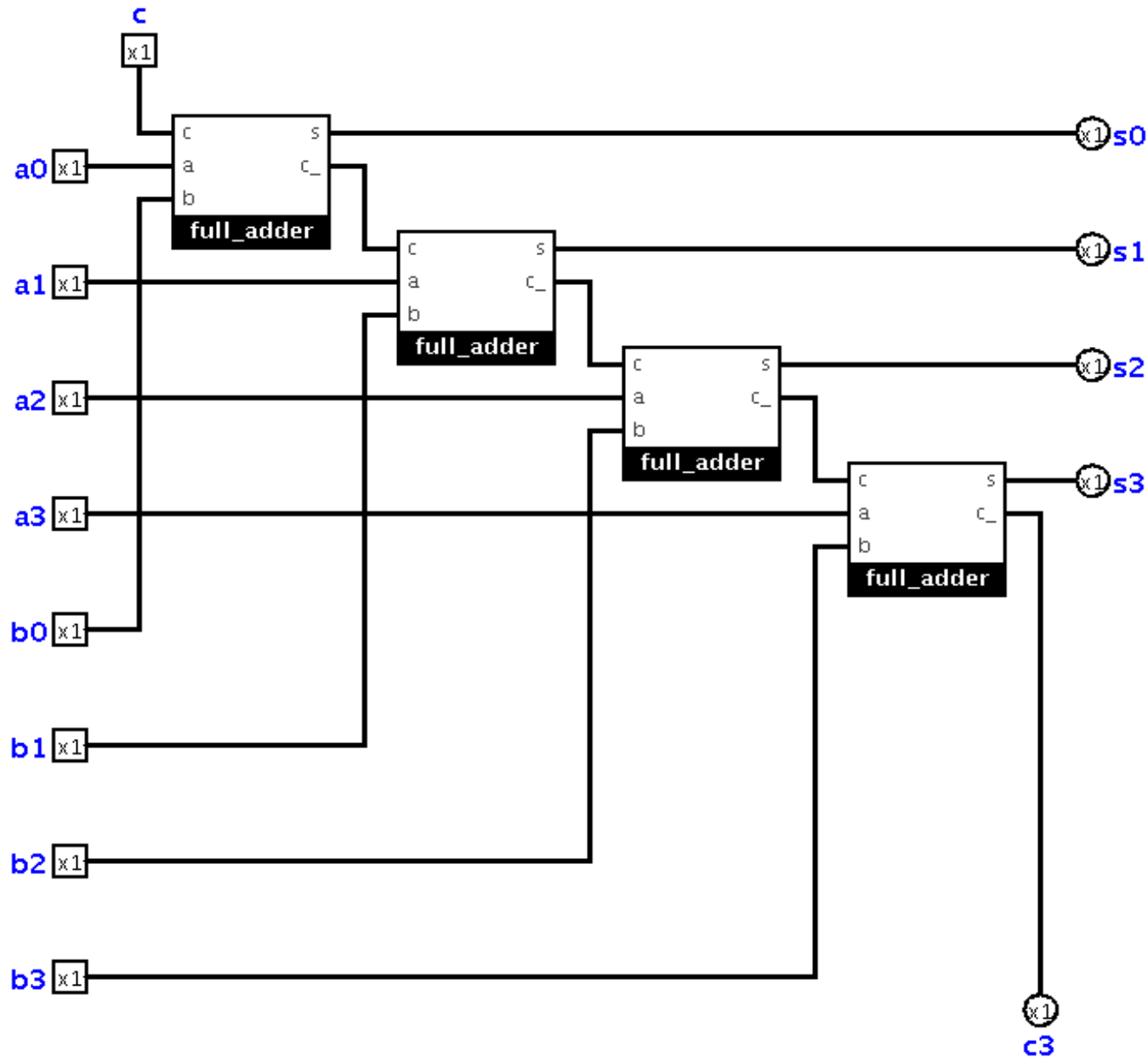
4-bit Ripple-Carry Adder



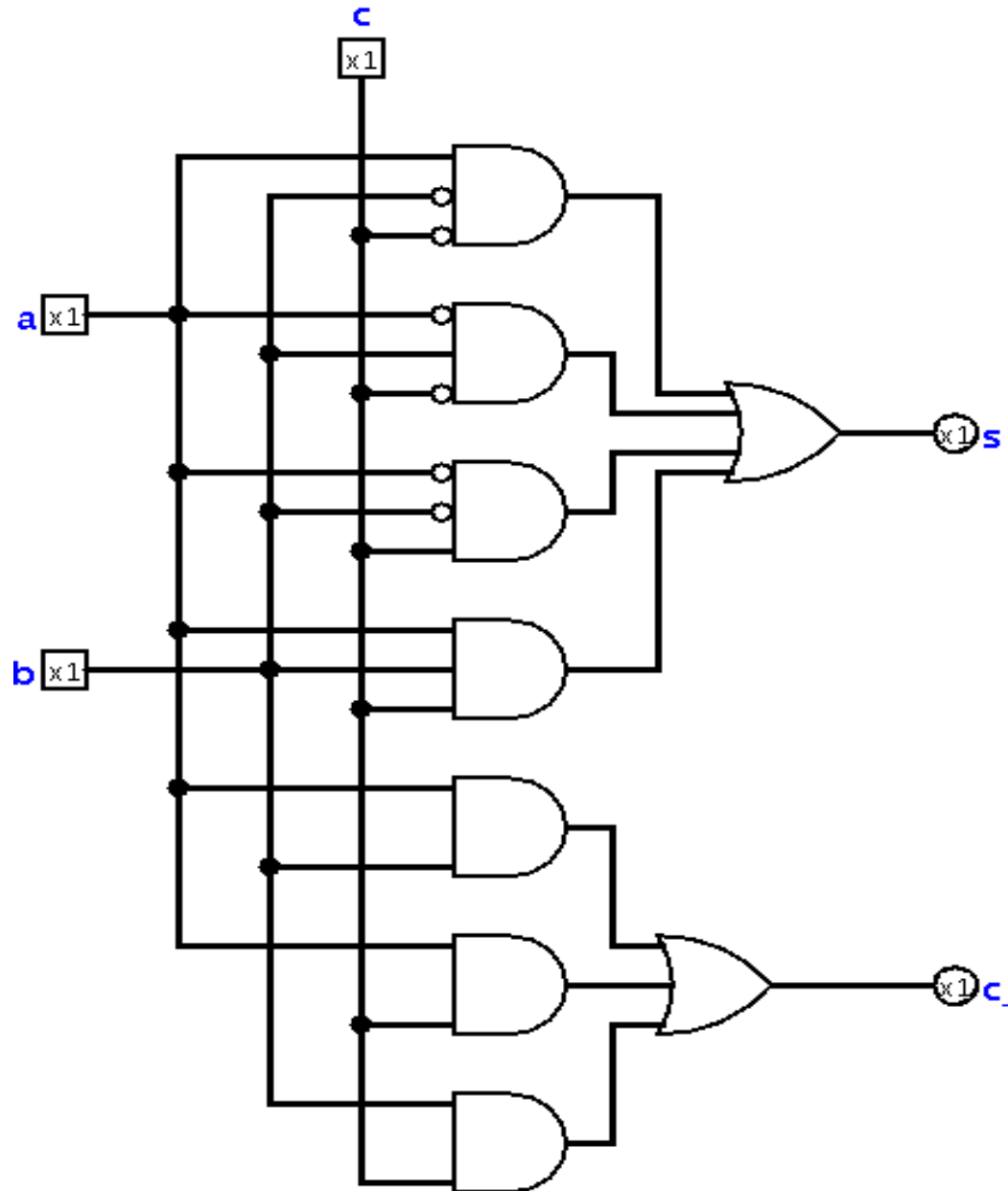
4-bit Ripple-Carry Adder



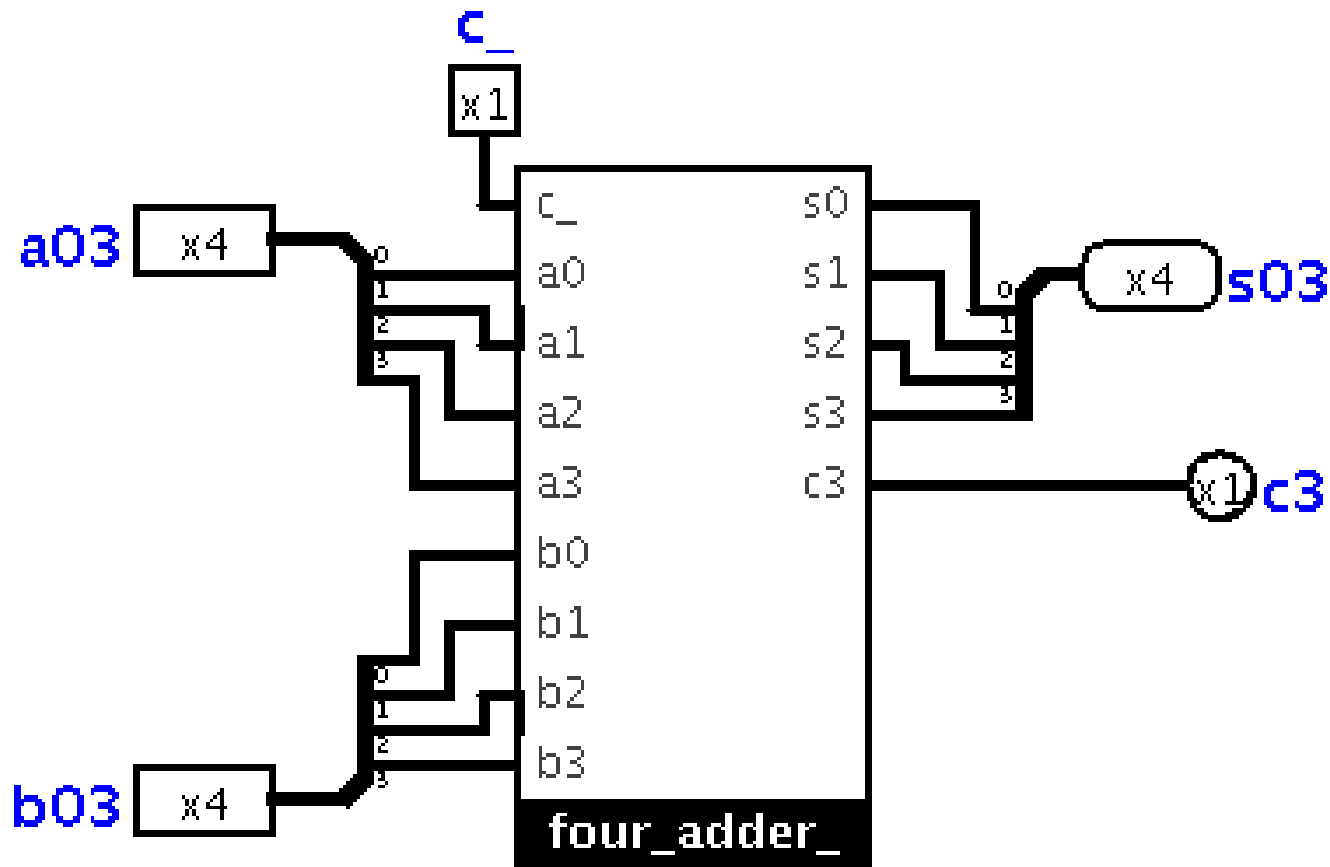
4-bit Ripple-Carry Adder



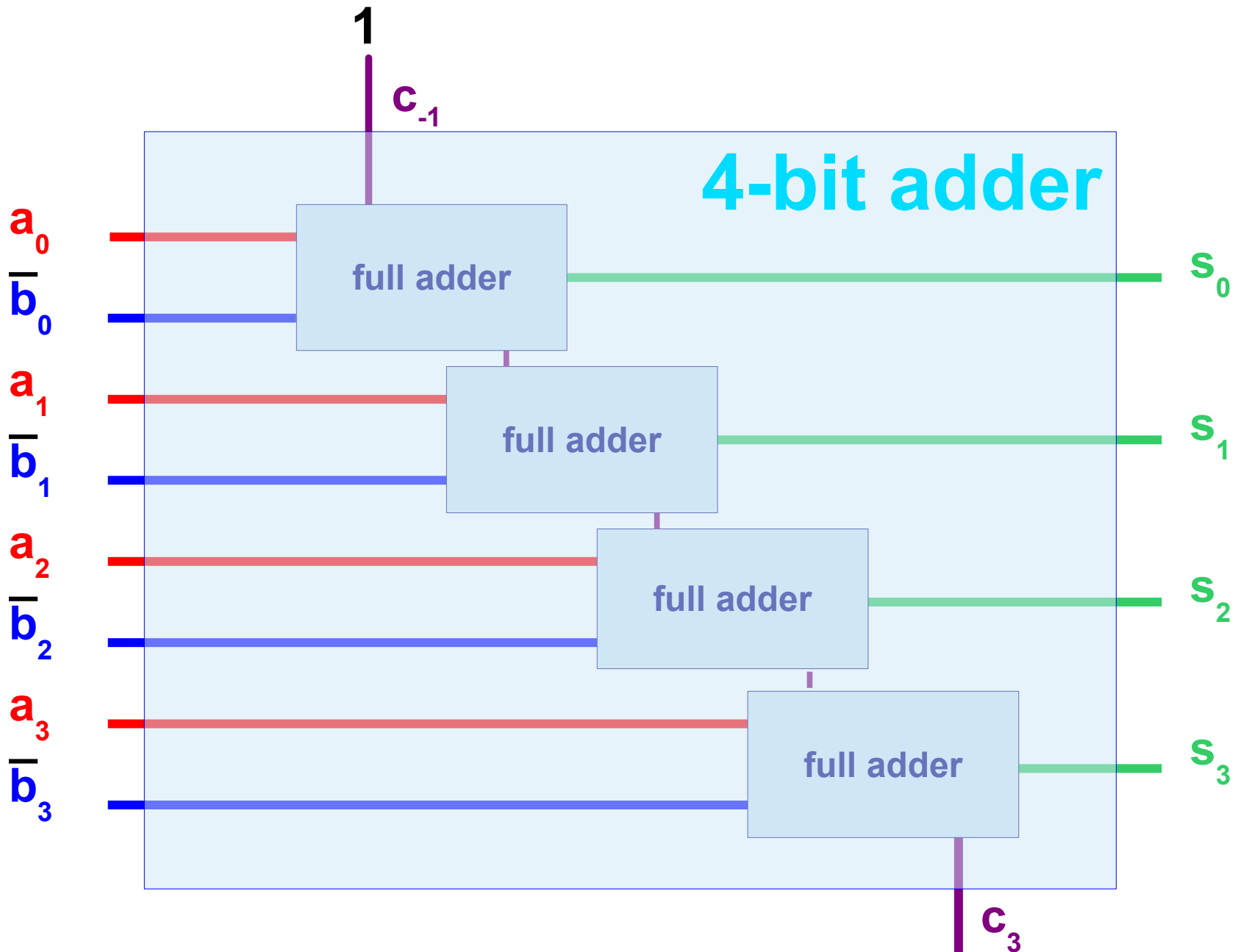
Full Adder Circuit



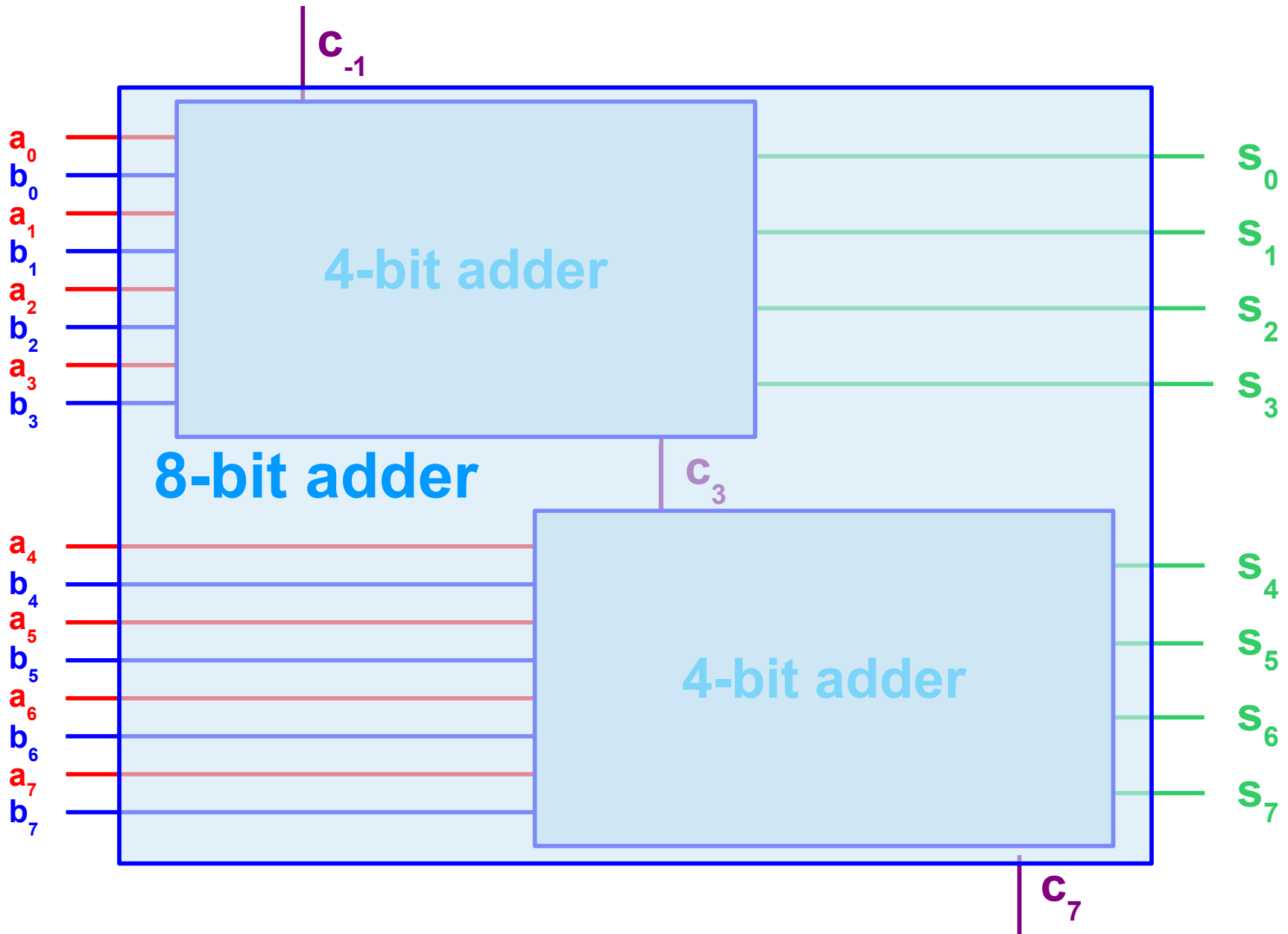
4-bit Ripple-Carry Adder



4-bit Subtraction



8-bit Ripple-Carry Adder



16-bit Ripple-Carry Adder

