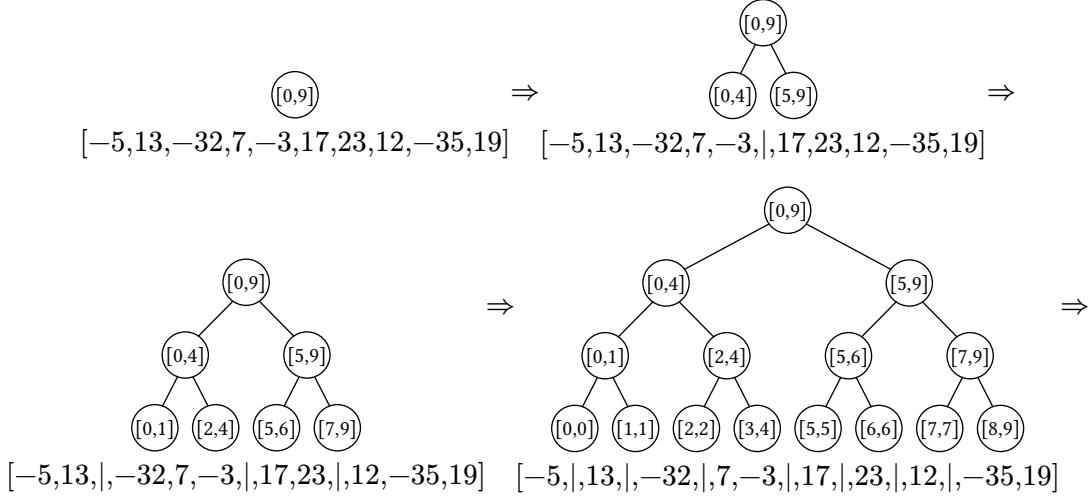


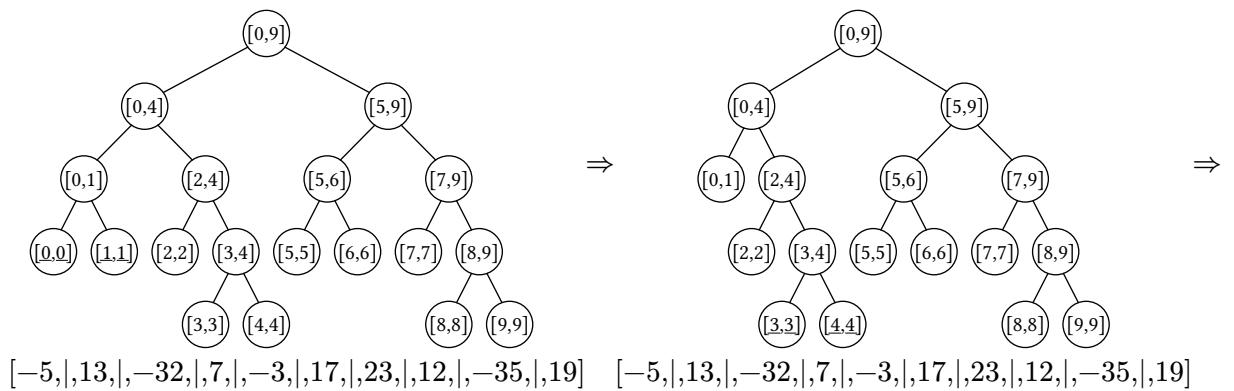
Aufgabe 1

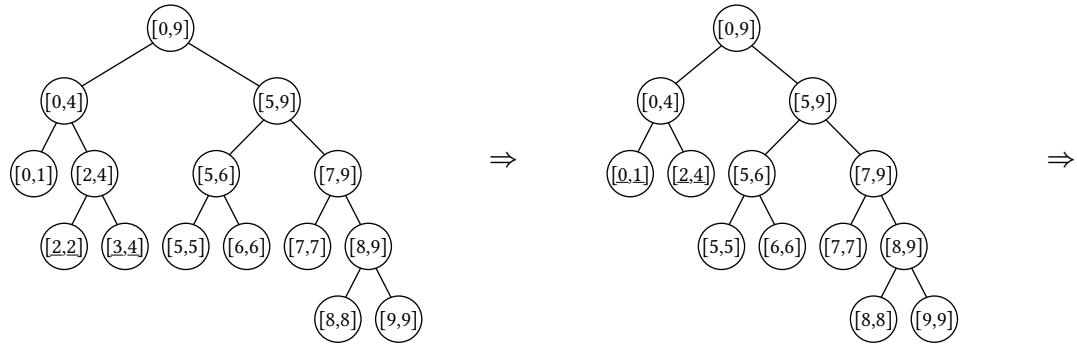
Sortiere $[-5, 13, -32, 7, -3, 17, 23, 12, -35, 19]$ mit ...

- MergeSort
 1. Divide

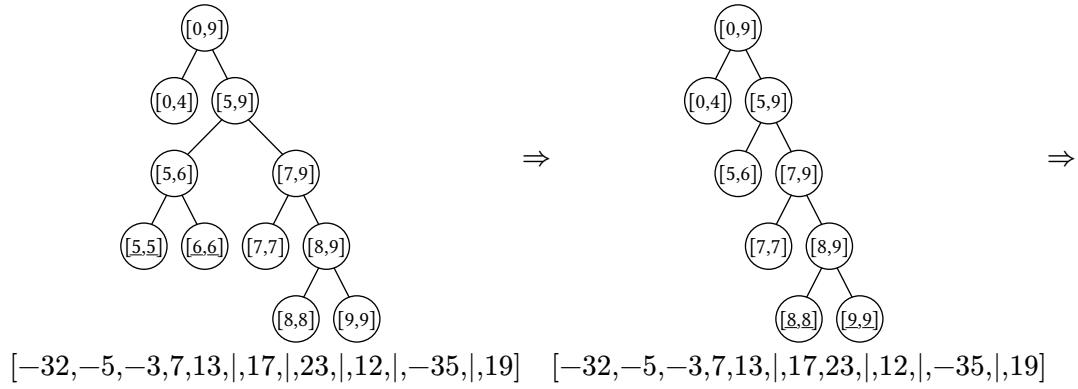


2. Conquer / Merge

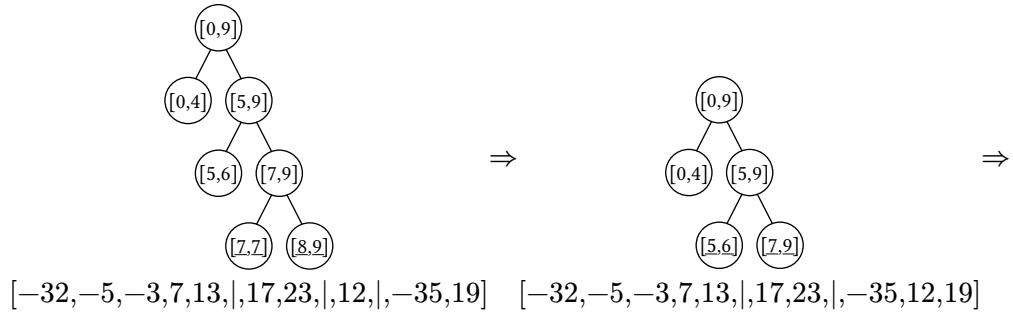




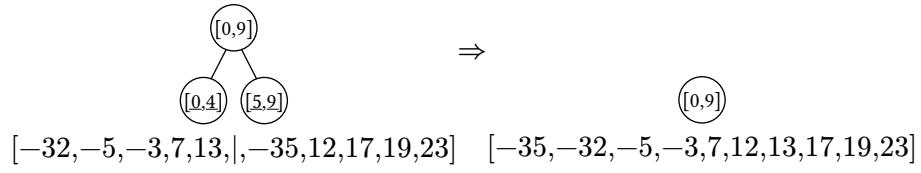
$[-5,13,|,7,|,-32,-3,|,17,|,23,|,12,|,-35,|,19]$ $[-5,13,|,-32,-3,7,|,17,|,23,|,12,|,-35,|,19]$



$[-32,-5,-3,7,13,|,17,|,23,|,12,|,-35,|,19]$ $[-32,-5,-3,7,13,|,17,23,|,12,|,-35,|,19]$



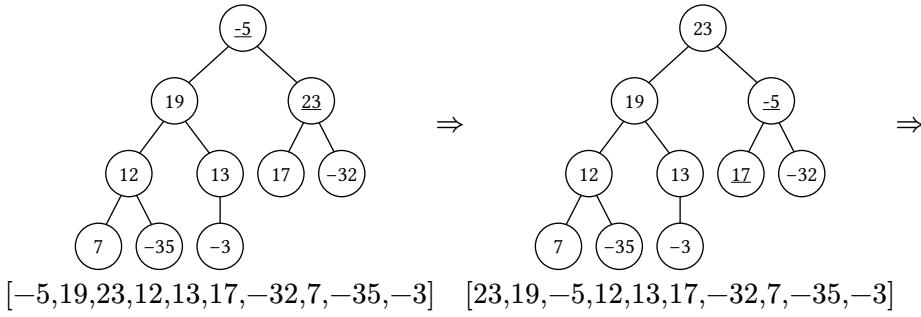
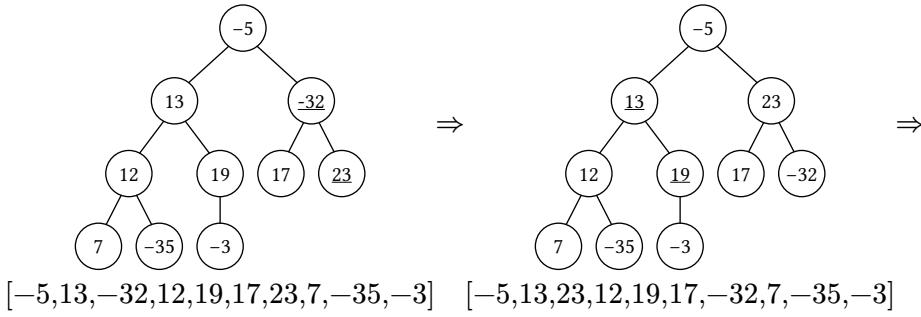
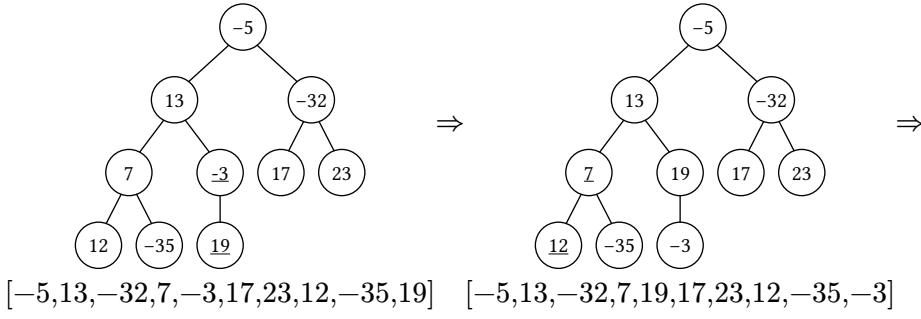
$[-32,-5,-3,7,13,|,17,23,|,12,|,-35,19]$ $[-32,-5,-3,7,13,|,17,23,|,-35,12,19]$



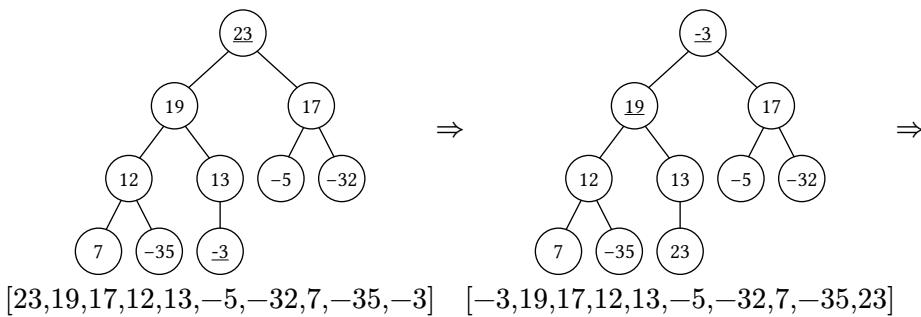
$[-32,-5,-3,7,13,|,-35,12,17,19,23]$ $[-35,-32,-5,-3,7,12,13,17,19,23]$

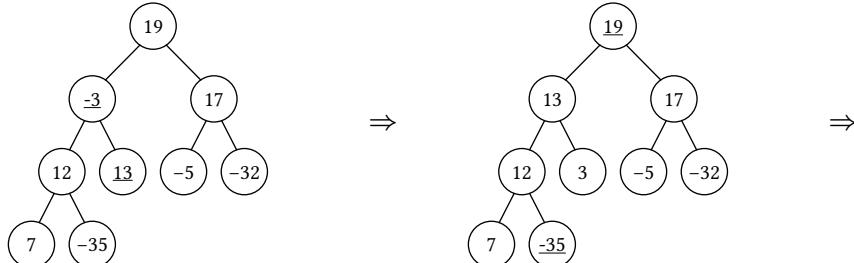
- HeapSort

1. Heapify

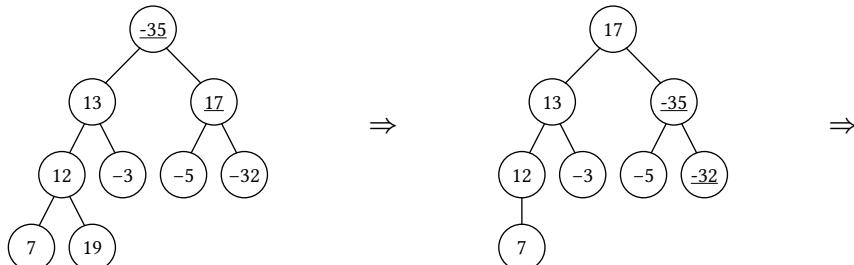


2. Sort

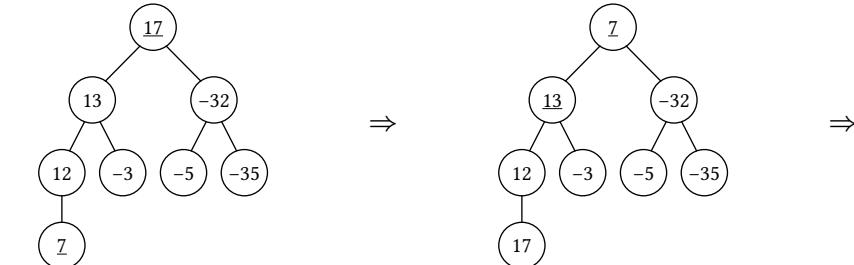




$[19, -3, 17, 12, 13, -5, -32, 7, -35, 23]$ $[19, 13, 17, 12, -3, -5, -32, 7, -35, 23]$



$[-35, 13, 17, 12, -3, -5, -32, 7, 19, 23]$ $[17, 13, -35, 12, -3, -5, -32, 7, 19, 23]$



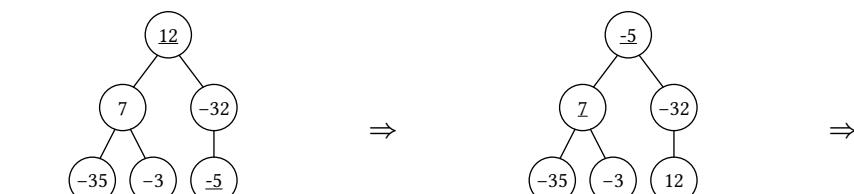
$[17, 13, -32, 12, -3, -5, -35, 7, 19, 23]$ $[7, 13, -32, 12, -3, -5, -35, 17, 19, 23]$



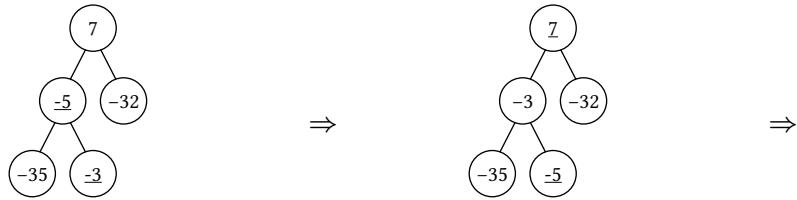
$[13, 7, -32, 12, -3, -5, -35, 17, 19, 23]$ $[13, 12, -32, 7, -3, -5, -35, 17, 19, 23]$



$[-35, 12, -32, 7, -3, -5, 13, 17, 19, 23]$ $[12, -35, -32, 7, -3, -5, 13, 17, 19, 23]$



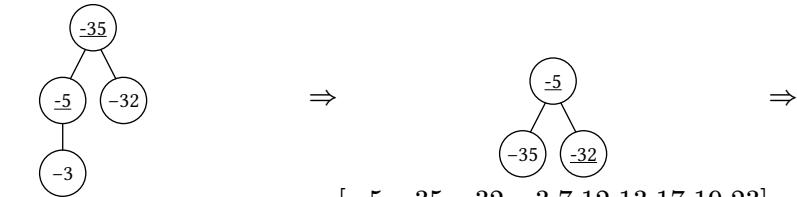
$[12, 7, -32, -35, -3, -5, 13, 17, 19, 23]$ $[-5, 7, -32, -35, -3, 12, 13, 17, 19, 23]$



$$[7, -5, -32, -35, -3, 12, 13, 17, 19, 23] \quad [7, -3, -32, -35, -5, 12, 13, 17, 19, 23]$$



$$[-5, -3, -32, -35, 7, 12, 13, 17, 19, 23] \quad [-3, -5, -32, -35, 7, 12, 13, 17, 19, 23]$$



$$[-35, -5, -32, -3, 7, 12, 13, 17, 19, 23] \quad [-5, -35, -32, -3, 7, 12, 13, 17, 19, 23]$$

,



$$[-32, -35, -5, -3, 7, 12, 13, 17, 19, 23] \quad [-35, -32, -5, -3, 7, 12, 13, 17, 19, 23]$$

$$\Rightarrow \text{[-35, -32, -5, -3, 7, 12, 13, 17, 19, 23]}$$

$$\Rightarrow [-35, -32, -5, -3, 7, 12, 13, 17, 19, 23]$$

Aufgabe 2

Zeige folgende Aussagen:

1. Ein Heap mit n Elementen hat die Höhe $\lfloor \log n \rfloor$

Sei $h \in \mathbb{N}$ die Höhe des Heap

Sei $n_i \in \mathbb{N}_0$ die Anzahl auf der i -ten Ebene

Es gilt:

$$\forall i \in \mathbb{N} : 1 \leq n_i \leq 2^i$$

Da ein Heap linksvoll ist muss auch gelten:

$$\begin{aligned} n = \sum_{i=0}^h n_i &= \sum_{i=0}^{h-1} 2^i + n_h \stackrel{\text{endl. geom. Reihe}}{=} \frac{2^{h-1+1} - 1}{2 - 1} + n_h = \frac{2^h - 1}{2 - 1} + n_h = 2^h - 1 + n_h \\ \Rightarrow n = 2^h + n_h - 1 &\Rightarrow 2^h - 1 < n \leq 2^{h+1} - 1 \Rightarrow \#\text{Höhe } h = 2^h \wedge \#\text{Höhe } h = 2^{h+1} - 1 \\ \Rightarrow \log(2^h - 1) &< \log n \leq \log(2^{h+1} - 1) \\ \Rightarrow \log(2^h - 1) &< \log n \leq \log(2^{h+1} - 1) < \log(2^{h+1}) \Rightarrow \log(2^h - 1) < \log n < \log(2^{h+1}) \\ \Rightarrow \log(2^h) &\leq \log n < \log(2^{h+1}) \Rightarrow h \leq \log n < h + 1 \Rightarrow h = \lfloor \log n \rfloor \end{aligned}$$

2. Ein Heap mit n Elementen hat höchstens $\lceil \frac{n}{2^{h+1}} \rceil$ viele Knoten der Höhe h

Sei H ein linksvoller Heap mit n Elementen. Dann gilt: $\max(\#\text{Knoten Höhe } h) = \lceil \frac{n}{2^{h+1}} \rceil$

Sei $\text{index}_{\text{Vorgänger}}(i) = \lfloor \frac{i-1}{2} \rfloor$ und $A(i) = i + 1$ die Anzahl der Elemente bis zu einem Index
IA: $h = 0$

$$\begin{aligned} \#\text{Blätter} &= n - A(\text{index}_{\text{Vorgänger}}(n-1)) = n - \left(\left\lfloor \frac{n-2}{2} \right\rfloor + 1 \right) \\ &= \begin{cases} n - (\frac{n-2}{2} + 1), & n \bmod 2 = 0 \\ n - (\frac{n-1-2}{2} + 1), & n \bmod 2 = 1 \end{cases} \\ &= \begin{cases} \frac{n}{2}, & n \bmod 2 = 0 \\ \frac{n}{2} + \frac{1}{2}, & n \bmod 2 = 1 \end{cases} = \lceil \frac{n}{2} \rceil \end{aligned}$$

IV: angenommen die Behauptung gilt für ein festes aber beliebiges $h-1$

IS: $h-1 \rightarrow h$

Aus dem Heap H wird durch streichen aller Blätter der Heap H' . Beobachtung: alle Knoten in H'

die Höhe $h-1$ haben in H die Höhe h (eigentlich noch durch Induktion Beweisen)

nach IV:

$$\begin{aligned} \#\text{Blätter}^{H'} &= \left\lceil \frac{n - \lceil \frac{n}{2} \rceil}{2^{(h-1)+1}} \right\rceil = \left\lceil \frac{\lfloor \frac{n}{2} \rfloor}{2^h} \right\rceil \leq \left\lceil \frac{\frac{n}{2}}{2^h} \right\rceil = \left\lceil \frac{n}{2^{h+1}} \right\rceil \Rightarrow \#\text{Knoten mit Höhe } h \leq \lceil \frac{n}{2^{h+1}} \rceil \\ &\Rightarrow \text{Die Aussage gilt für alle } h \end{aligned}$$

3. Für alle x mit $|x| < 1$: $\sum_{k=0}^{\infty} kx^k = \frac{x}{(1-x)^2}$

$$\forall x \in \mathbb{R} \wedge |x| < 1 : \sum_{k=0}^{\infty} kx^k = \frac{x}{(1-x)^2}$$

geometrische Reihe: $\sum_{k=0}^{\infty} x^k = \frac{1}{1-x}$

$$\frac{\partial}{\partial x} \sum_{k=0}^{\infty} x^k = \frac{\partial}{\partial x} \frac{1}{1-x} \Leftrightarrow \sum_{k=0}^{\infty} kx^{k-1} = \frac{1}{(1-x)^2}$$

$$\Rightarrow \forall 0 < |x| < 1 : x \sum_{k=0}^{\infty} kx^{k-1} = x \frac{1}{(1-x)^2} \Leftrightarrow \sum_{k=0}^{\infty} kx^k = \frac{x}{(1-x)^2}$$

Aufgabe 3

$o_{i,j}$ gilt: $o_{i,j} = \sum_{k=1}^n (m_{i,k} \cdot n_{k,i})$

1.

- Variante 1:

$$O_{11} = M_{11} \cdot N_{11} + M_{12} \cdot N_{12} = \text{MN1} + \text{MN2} =$$

$$\begin{pmatrix} m_{1,1} & m_{1,2} & \dots & m_{1,\frac{n}{2}} \\ m_{2,1} & m_{2,2} & \dots & m_{2,\frac{n}{2}} \\ \vdots & \vdots & \ddots & \vdots \\ m_{\frac{n}{2},1} & \dots & \dots & m_{\frac{n}{2},\frac{n}{2}} \end{pmatrix} \cdot \begin{pmatrix} n_{1,1} & m_{1,2} & \dots & n_{1,\frac{n}{2}} \\ n_{2,1} & m_{2,2} & \dots & n_{2,\frac{n}{2}} \\ \vdots & \vdots & \ddots & \vdots \\ n_{\frac{n}{2},1} & \dots & \dots & n_{\frac{n}{2},\frac{n}{2}} \end{pmatrix} + \begin{pmatrix} m_{1,\frac{n}{2}} & m_{1,\frac{n}{2}+1} & \dots & m_{1,n} \\ m_{2,\frac{n}{2}} & m_{2,\frac{n}{2}+1} & \dots & m_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ m_{\frac{n}{2},\frac{n}{2}} & \dots & \dots & m_{\frac{n}{2},n} \end{pmatrix} \cdot \begin{pmatrix} m_{\frac{n}{2}+1,1} & m_{\frac{n}{2}+1,2} & \dots & m_{\frac{n}{2}+1,\frac{n}{2}} \\ m_{\frac{n}{2}+2,1} & m_{\frac{n}{2}+2,2} & \dots & m_{\frac{n}{2}+2,\frac{n}{2}} \\ \vdots & \vdots & \ddots & \vdots \\ m_{n,1} & \dots & \dots & m_{n,\frac{n}{2}} \end{pmatrix}$$

$$i, j \in \left[1, \frac{n}{2}\right] : \text{MN1}_{i,j} = \sum_{k=1}^{\frac{n}{2}} m_{i,k} \cdot n_{k,i} \wedge \text{MN2}_{i,j} = \sum_{k=\frac{n}{2}+1}^n m_{i,k} \cdot n_{k,i}$$

$$\Rightarrow \text{MN1}_{i,j} + \text{MN2}_{i,j} = \sum_{k=1}^{\frac{n}{2}} m_{i,k} \cdot n_{k,i} + \sum_{k=\frac{n}{2}+1}^n m_{i,k} \cdot n_{k,i} = \sum_{k=1}^n m_{i,k} \cdot n_{k,i} = o_{i,j}$$

gilt für $O_{1,2}, O_{2,1}, O_{2,2}$ analog

- Variante 2:

$$H_1 := (M_{1,1} + M_{2,2}) \cdot (N_{1,1} + N_{2,2})$$

$$H_2 := (M_{2,1} + M_{2,2}) \cdot N_{1,1}$$

$$H_3 := M_{1,1} \cdot (N_{1,2} - N_{2,2})$$

$$H_4 := M_{2,2} \cdot (N_{2,1} - N_{1,1})$$

$$H_5 := (M_{1,1} + M_{1,2}) \cdot N_{2,2}$$

$$H_6 := (M_{2,1} - M_{1,1}) \cdot (N_{1,1} + N_{1,2})$$

$$H_7 := (M_{1,2} - M_{2,2}) \cdot (N_{2,1} + N_{2,2})$$

$$O_{1,1} := H_1 + H_4 - H_5 + H_7$$

$$O_{1,2} := H_3 + H_5$$

$$O_{2,1} := H_2 + H_4$$

$$O_{2,2} := H_1 - H_2 + H_3 + H_6$$

Umformungen:

$$\begin{aligned}
O_{1,1} &= H_1 + H_4 - H_5 + H_7 \\
&= (M_{1,1} + M_{2,2}) \cdot (N_{1,1} + N_{2,2}) + M_{2,2} \cdot (N_{2,1} - N_{1,1}) - (M_{1,1} + M_{1,2}) \cdot N_{2,2} + \\
&\quad (M_{1,2} - M_{2,2}) \cdot (N_{2,1} + N_{2,2}) \\
&= M_{1,1} \cdot N_{1,1} + M_{1,1} \cdot N_{2,2} + M_{2,2} \cdot N_{1,1} + M_{2,2} \cdot N_{2,2} + M_{2,2} \cdot N_{2,1} - M_{2,2} \cdot N_{1,1} - \\
&\quad M_{1,1} \cdot N_{2,2} - M_{1,2} \cdot N_{2,2} + M_{1,2} \cdot N_{2,1} + M_{1,2} \cdot N_{2,2} - M_{2,2} \cdot N_{2,1} - M_{2,2} \cdot N_{2,2} \\
&= (M_{1,1} \cdot N_{1,1} + M_{1,2} \cdot N_{2,1}) + (M_{1,1} \cdot N_{2,2} - M_{1,1} \cdot N_{2,2}) + (M_{2,2} \cdot N_{1,1} - M_{2,2} \cdot N_{1,1}) \\
&\quad + (M_{2,2} \cdot N_{2,2} - M_{2,2} \cdot N_{2,2}) + (M_{2,2} \cdot N_{2,1} - M_{2,2} \cdot N_{2,1}) + (-M_{1,2} \cdot N_{2,2} + M_{1,2} \cdot N_{2,2}) \\
&= M_{1,1} \cdot N_{1,1} + M_{1,2} \cdot N_{2,1}
\end{aligned}$$

$$\begin{aligned}
O_{1,2} &= H_3 + H_5 \\
&= M_{1,1} \cdot (N_{1,2} - N_{2,2}) + (M_{1,1} + M_{1,2}) \cdot N_{2,2} \\
&= M_{1,1} \cdot N_{1,2} - M_{1,1} \cdot N_{2,2} + M_{1,1} \cdot N_{2,2} + M_{1,2} \cdot N_{2,2} \\
&= (M_{1,1} \cdot N_{1,2} + M_{1,2} \cdot N_{2,2}) - M_{1,1} \cdot N_{2,2} + M_{1,1} \cdot N_{2,2} \\
&= M_{1,1} \cdot N_{1,2} + M_{1,2} \cdot N_{2,2}
\end{aligned}$$

$$\begin{aligned}
O_{2,1} &= H_2 + H_4 \\
&= (M_{2,1} + M_{2,2}) \cdot N_{1,1} + M_{2,2} \cdot (N_{2,1} - N_{1,1}) \\
&= M_{2,1} \cdot N_{1,1} + M_{2,2} \cdot N_{1,1} + M_{2,2} \cdot N_{2,1} - M_{2,2} \cdot N_{1,1} \\
&= M_{2,1} \cdot N_{1,1} + M_{2,2} \cdot N_{2,1}
\end{aligned}$$

$$\begin{aligned}
O_{2,2} &= H_1 - H_2 + H_3 + H_6 \\
&= (M_{1,1} + M_{2,2}) \cdot (N_{1,1} + N_{2,2}) - (M_{2,1} + M_{2,2}) \cdot N_{1,1} + M_{1,1} \cdot (N_{1,2} - N_{2,2}) + \\
&\quad (M_{2,1} - M_{1,1}) \cdot (N_{1,1} + N_{1,2}) \\
&= (M_{1,1} \cdot N_{1,1} + M_{1,1} \cdot N_{2,2} + M_{2,2} \cdot N_{1,1} + M_{2,2} \cdot N_{2,2}) - (M_{2,1} \cdot N_{1,1} + M_{2,2} \cdot N_{1,1}) + \\
&\quad (M_{1,1} \cdot N_{1,2} - M_{1,1} \cdot N_{2,2}) + (M_{2,1} \cdot N_{1,1} + M_{2,1} \cdot N_{1,2} - M_{1,1} \cdot N_{1,1} - M_{1,1} \cdot N_{1,2}) \\
&= M_{1,1} \cdot N_{1,1} + M_{1,1} \cdot N_{2,2} + M_{2,2} \cdot N_{1,1} + M_{2,2} \cdot N_{2,2} - M_{2,1} \cdot N_{1,1} - M_{2,2} \cdot N_{1,1} + \\
&\quad M_{1,1} \cdot N_{1,2} - M_{1,1} \cdot N_{2,2} + M_{2,1} \cdot N_{1,1} + M_{2,1} \cdot N_{1,2} - M_{1,1} \cdot N_{1,1} - M_{1,1} \cdot N_{1,2} \\
&= (M_{2,1} \cdot N_{1,2} + M_{2,2} \cdot N_{2,2}) + (M_{1,1} \cdot N_{1,1} - M_{1,1} \cdot N_{1,1}) + (M_{2,2} \cdot N_{1,1} - M_{2,2} \cdot N_{1,1}) + \\
&\quad (-M_{2,1} \cdot N_{1,1} + M_{2,1} \cdot N_{1,1}) + (M_{1,1} \cdot N_{1,2} - M_{1,1} \cdot N_{1,2}) - M_{1,1} \cdot N_{2,2} + M_{1,1} \cdot N_{2,2} \\
&= M_{2,1} \cdot N_{1,2} + M_{2,2} \cdot N_{2,2}
\end{aligned}$$

2.

für n gelte: $n, k \in \mathbb{N}_0 : n = 2^k$

- Variante 1:

Matrix Additionen werden durch die Konstante $4n^2$ abgeschätzt

$$\Rightarrow T(n) = 8T\left(\frac{n}{2}\right) + 4n^2$$

mit Master Methode: $4n^2 \in O(n^{\log_2 8}) \Rightarrow T(n) \in \Theta(n^3)$

- Variante 2:

Matrix Additionen werden durch die Konstante $18n^2$ abgeschätzt

$$\Rightarrow 7T\left(\frac{n}{2}\right) + 18n^2$$

mit Master Methode: $18n^2 \in O(n^{\log_2 7}) \Rightarrow T(n) \in \Theta(n^{\log_2 7})$

Aufgabe 4

```
main.cpp:  
#include "matrix.hpp"  
#include <chrono>  
#include <iostream>  
#include <climits>  
  
int main() {  
    for (int i = 2; i < INT_MAX; i = i * 2) {  
        std::cout << i << std::endl;  
        Matrix M(i,i);  
        Matrix N(i,i);  
  
        M.randomFill();  
        N.randomFill();  
  
        auto tempRecStart = std::chrono::system_clock::now();  
        Matrix res1 = M.mult(N);  
        auto durationRec = std::chrono::system_clock::now() - tempRecStart;  
  
        auto tempNormStart = std::chrono::system_clock::now();  
        Matrix res2 = M*N;  
        auto durationNorm = std::chrono::system_clock::now() - tempNormStart;  
  
        if (durationRec <= durationNorm) {  
            std::cout << "break even: " << i << std::endl;  
            break;  
        }  
    }  
}  
  
matrix.hpp:  
#pragma once  
class Matrix {  
private:  
    unsigned int m, n;  
    int** val;  
  
public:  
    Matrix();  
    Matrix(unsigned int m, unsigned n);  
    Matrix(const Matrix& other);  
    ~Matrix();  
    void print() const;  
    void input();
```

```

void randomFill();
Matrix mult(const Matrix&) const;
void operator=(const Matrix& B);
Matrix operator+(const Matrix& B) const;
Matrix operator-(const Matrix& B) const;
Matrix operator*(const Matrix& B) const;
const int* operator[](int) const;
void setCell(unsigned int, unsigned int, int);

friend void recursiveMult(const Matrix& M, const Matrix& N, Matrix& O);
};

matrix.cpp:
#include "matrix.hpp"
#include <iostream>
#include <string>
#include <random>
#include <climits>
#include <cmath>
#include <exception>

Matrix::Matrix(): m(0), n(0), val(nullptr) {}
Matrix::Matrix(unsigned int m, unsigned int n): m(m), n(n) {
    this->val = new int*[m];
    for (int i = 0; i < m; i++) {
        this->val[i] = new int[n];
        for (int j = 0; j < n; j++)
            this->val[i][j] = 0;
    }
}

Matrix::Matrix(const Matrix& other): m(other.m), n(other.n) {
    this->val = new int*[m];
    for (int i = 0; i < m; i++) {
        this->val[i] = new int[n];
        for (int j = 0; j < n; j++)
            this->val[i][j] = other.val[i][j];
    }
}

Matrix::~Matrix() {
    if (this->val != nullptr) {
        for (int i = 0; i < m; i++)
            delete[] this->val[i];
        delete[] this->val;
    }
}

```

```

}

void Matrix::print() const {
    std::cout << "-----" << std::endl;
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < m; j++)
            std::cout << this->val[i][j] << "\t";
        std::cout << std::endl;
    }
    std::cout << "-----" << std::endl;
}

void Matrix::operator= (const Matrix& other) {
    if (val != nullptr)
        delete this->val;
    this->m = other.m;
    this->n = other.n;
    this->val = new int*[m];
    for (int i = 0; i < m; i++) {
        this->val[i] = new int[n];
        for (int j = 0; j < n; j++)
            this->val[i][j] = other.val[i][j];
    }
}

const int* Matrix::operator[] (int index) const {
    if (index < 0 || index >= m)
        throw std::runtime_error("index out of bounds");
    return this->val[index];
}

void Matrix::input() {
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < m; j++) {
            int tmp = 0;

            std::cout << "please give a value for input: (" << i << "/" << j << ")";
            ;
            while (true) {
                std::cin >> tmp;
                if (std::cin.fail()) {
                    std::cin.clear();
                    std::cin.ignore(10000000000, '\n');
                    std::cout << "please give a valid value for input: (" << i <<
"/" << j << ")": ";
                }
            }
        }
    }
}

```

```

        } else
            break;
    }

    this->val[i][j] = tmp;
}
}

void Matrix::randomFill() {
    static std::random_device rd;
    static std::mt19937 gen (rd());
    static std::uniform_int_distribution<int> dist(-100, 100);

    for (int i = 0; i < m; i++) {
        for (int j = 0; j < n; j++) {
            this->val[i][j] = dist(gen);
        }
    }
}

Matrix Matrix::operator+ (const Matrix& other) const {
    if (this->n != other.n || this->m != other.m)
        throw std::runtime_error("cannot add two Matrices that don't have the same
dimension");
    unsigned int iterations = 0;

    Matrix res(this->m, this->n);
    for (int i = 0; i < this->m; i++) {
        for (int j = 0; j < this->n; j++) {
            iterations++;
            res.val[i][j] = this->val[i][j] + other.val[i][j];
        }
    }
    return res;
}

Matrix Matrix::operator- (const Matrix& other) const {
    if (this->n != other.n || this->m != other.m)
        throw std::runtime_error("cannot add two Matrices that don't have the same
dimension");
    unsigned int iterations = 0;

    Matrix res(this->m, this->n);
    for (int i = 0; i < this->m; i++) {
        for (int j = 0; j < this->n; j++) {

```

```

        iterations++;
        res.val[i][j] = this->val[i][j] - other.val[i][j];
    }
}
return res;
}

Matrix Matrix::operator* (const Matrix& other) const {
    if (this->n != other.m || this->m != other.n)
        throw std::runtime_error("cannot multiply two Matrices where rows and column
not match");
    Matrix res(other.m, this->n);
    int iterations = 0;

    for (int row = 0; row < res.n; row++) {
        for (int column = 0; column < res.m; column++) {
            int cij = 0;
            for (int index = 0; index < res.m; index++) {
                iterations++;
                cij += this->val[row][index] * other.val[index][column];
            }
            res.val[row][column] = cij;
        }
    }

    return res;
}

void Matrix::setCell(unsigned int y, unsigned int x, int value) {
    if (y >= m || x >= n)
        throw std::runtime_error("index out bounds error");
    this->val[y][x] = value;
}

void recursiveMult(const Matrix& M, const Matrix& N, Matrix& O) {
    if (M.n == 2) {
        O.setCell(0, 0, M[0][0] * N[0][0] + M[0][1] * N[1][0]);
        O.setCell(0, 1, M[0][0] * N[0][1] + M[0][1] * N[1][1]);
        O.setCell(1, 0, M[1][0] * N[0][0] + M[1][1] * N[1][0]);
        O.setCell(1, 1, M[1][0] * N[0][1] + M[1][1] * N[1][1]);
        return;
    }

    int nHalf = M.n/2;
    Matrix M11(nHalf, nHalf);

```

```

Matrix M12(nHalf, nHalf);
Matrix M21(nHalf, nHalf);
Matrix M22(nHalf, nHalf);
Matrix N11(nHalf, nHalf);
Matrix N12(nHalf, nHalf);
Matrix N21(nHalf, nHalf);
Matrix N22(nHalf, nHalf);

for (int i = 0; i < nHalf; i++) {
    for (int j = 0; j < nHalf; j++) {
        M11.setCell(i, j, M[i][j]);
        M12.setCell(i, j, M[i][nHalf + j]);
        M21.setCell(i, j, M[nHalf + i][j]);
        M22.setCell(i, j, M[nHalf + i][nHalf + j]);
        N11.setCell(i, j, N[i][j]);
        N12.setCell(i, j, N[i][nHalf + j]);
        N21.setCell(i, j, N[nHalf + i][j]);
        N22.setCell(i, j, N[nHalf + i][nHalf + j]);
    }
}

Matrix H1(nHalf, nHalf);
Matrix H2(nHalf, nHalf);
Matrix H3(nHalf, nHalf);
Matrix H4(nHalf, nHalf);
Matrix H5(nHalf, nHalf);
Matrix H6(nHalf, nHalf);
Matrix H7(nHalf, nHalf);

recursiveMult(M11 + M22, N11 + N22, H1);
recursiveMult(M21 + M22, N11, H2);
recursiveMult(M11, N12 - N22, H3);
recursiveMult(M22, N21 - N11, H4);
recursiveMult(M11 + M12, N22, H5);
recursiveMult(M21 - M11, N11 + N12, H6);
recursiveMult(M12 - M22, N21 + N22, H7);

for (int i = 0; i < nHalf; i++) {
    for (int j = 0; j < nHalf; j++) {
        O.setCell(i, j, H1[i][j] + H4[i][j] - H5[i][j] + H7[i][j]);
        O.setCell(i, nHalf + j, H3[i][j] + H5[i][j]);
        O.setCell(nHalf + i, j, H2[i][j] + H4[i][j]);
        O.setCell(nHalf + i, nHalf + j, H1[i][j] - H2[i][j] + H3[i][j] + H6[i][j]);
    }
}

```

```
    }
}

}

Matrix Matrix::mult(const Matrix& other) const {
    if (n == 0 || m == 0 || other.n == 0 || other.m == 0) {
        throw std::runtime_error("cannot multiply empty Matrices");
    }else if (n != other.n || m != other.m || n != m) {
        throw std::runtime_error("cannot multiply non square shaped matrices
recursively");
    } else if (std::fmod(log2(n), 1) != 0) {
        throw std::runtime_error("cannot multiply matrices that are not in
R^(2^ix2^i) recursively");
    }

    Matrix O(n,n);
    recursiveMult(*this, other, 0);

    return O;
}
```