



Note: Attempt all the questions. Use blue or black pen only.

1. Give a dynamic programming solution to find Longest Common Subsequence. [Note: only write down the base case and formula/equation to update the values in the dynamic programming table]. [0.5]

$$c[i, j] = \begin{cases} 0 & \text{if } i = 0 \text{ or } j = 0, \\ c[i - 1, j - 1] + 1 & \text{if } i, j > 0 \text{ and } x_i = y_j, \\ \max \{c[i, j - 1], c[i - 1, j]\} & \text{if } i, j > 0 \text{ and } x_i \neq y_j. \end{cases}$$

2. Given a procedure, RANDOMLY-PERMUTE, that produces a uniform random permutation, a permutation as likely as any other permutation of any input array A. What is the probability that the algorithm generates a permutation such that for $1 \leq i < n$, $A[i] < A[i+1]$? [0.5]

$$\frac{1}{n!}$$

3. Given an array of size n write down an efficient randomized algorithm to find the $(n/2)^{\text{th}}$ smallest element in the array. Write down the expected time complexity of the algorithm. [1]

Quick_Select($S, n/2$)

Expected time complexity: $O(n)$

4. Let $A[1..n]$ be an array of n distinct numbers. Given a number x , use indicator random variables to compute the expected number of elements in an array that are greater than x . [1]

$$E[Z] = \sum_{i=1}^n E[z_{ix}] = \sum_{i=1}^n pr[z_{ix}] = \sum_{i=1}^n \frac{1}{2} = \frac{n}{2} \text{ where } z_{ix} \text{ is the event when } x < z_i$$

5. What is the best-case complexity of the Partition Algorithm in Quick Sort? Justify your answer. [1]
 $O(n)$. The algorithm scans all the elements of the array to find the correct location of the pivot hence the complexity is $O(n)$
6. The worst-case complexity of the Quick Select is $O(n^2)$. Justify your answer [no partial marking] (T/F) [0.5]
In the case of the sorted array, the bad partition will result in a recurrence relation of the form $T(n) = T(n-1) + O(n)$; hence the worst-case complexity is $O(n^2)$.
7. The average-case complexity of the Quick Sort is $O(n \lg n)$. Justify your answer [no partial marking] (T/F) [0.5]
An average-case would be a combination of good and bad partitions; hence, the average case complexity is $O(\log n)$.