

③ physical address = virtual address + base

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③  $0 \leq \text{virtual address} < \text{bounds}$ .

## CALCULATION BASED THINGS

Segment	Base	Size	Grows Positive?	Protection
code	32K	2K	1	Read-execute
Heap	34K	2K	1	Read-Write
Stack	28K	2K	0	Read-Write

Segment	Bits	
code	00	(13) $\text{PTEAddr} = \text{PTBR} + (\text{VPN} * \text{size of (PTE)})$
Heap	01	(14) $\text{header\_t} * \text{hptr} = (\text{void} * \text{ptr}) - \text{sizeof}(\text{header\_t})$
Stack	10	(15) $\text{AMAT} = (\text{P}_{\text{hit}} * \text{T}_{\text{m}}) + (\text{P}_{\text{miss}} * \text{T}_{\text{d}})$ <div style="display: flex; justify-content: space-around; font-size: small;"> <span>time to access memory</span> <span>time to access disk</span> </div>

④ Virtual Address space is  $n$  KB

$$\log_2(n) = x \text{ bits}$$

so you'll need  $x$  bits of address. (first two bits for segment)

⑤ Convert the virtual address into binary the left most <sup>few</sup> bits will tell you the segment.

⑥ To find number of bits for a segment, if there are  $n$  segments so,

$$\log_2(n) = x \text{ bits}$$

⑦  $\text{Total No of Pages} = \frac{\text{Address Space}}{\text{Page Size}}$  if 4K/2K = 2

⑧  $\text{offset} = \log_2(\text{Total Pages})$

⑨  $\text{VPN \# of Bits} = \frac{\text{VA}}{\text{bits}} = \text{offset} \log_2(\text{Total Pages})$

⑩  $\text{Page Table Size} = \text{No of Pages} \times \text{Size of each PTE}$

⑪  $\text{Internal Fragmentation} = (\text{No of Pages} \times \text{Page Size}) - \text{used memory size}$

⑫  $\text{Physical Memory} = \text{PFN} + \text{offset}$