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Problem Set 5

1. $2T\left(\frac{n}{2}\right) + 1000n \quad \forall n \geq 2$

Base case:

$$n = 2$$

Assume $O(n \log(n))$; prove $2T\left(\frac{n}{2}\right) + 1000\left(\frac{n}{2}\right) \leq c\left(\left(\frac{n}{2}\right)\log\left(\frac{n}{2}\right)\right)$

$$2 + 2002 \leq c(2)$$

$$c = 1001$$

Substitution:

$$2T\left(\frac{n}{2}\right) + 1000n \leq c(n \log(n))$$

Values smaller than n can be assumed $\leq c(n \log(n))$, so it's substituted in for $\left(\frac{n}{2}\right)$

$$2c\left(\frac{n}{2}\right)\log\left(\frac{n}{2}\right) + 1000n$$

$\log\left(\frac{n}{2}\right)$ due to the division rule can be expanded like so

$$2c\left(\frac{n}{2}\right)\log(n) - 1 + 1000n$$

The multiplication and division of 2 can be cancelled and the cn term can be distributed

$$cn \log n - (cn + 1000n)$$

Because the $cn + 1000n$ can be assumed to be a positive value, the whole term must be

$$\leq c(n \log(n)) \text{ therefore } 2T\left(\frac{n}{2}\right) + 1000n \leq c(n \log(n)) \text{ for any } c \text{ and the recursion } \in O(n \log(n))$$

2. $7T\left(\frac{n}{2}\right) + 18n^2 \quad \forall n \geq 2$

Base case:

$$n = 2$$

Assume $O(n^3)$; prove $7\left(\frac{n}{2}\right) + 18(2)^2 \leq c(2)^3$

$$79 \leq c8$$

$$c = 10$$

Substitution:

$$7\left(\frac{n}{2}\right) + 18(n)^2 \leq c(n)^3$$

Values smaller than n can be assumed $\leq cn^3$, so it's substituted in for $\left(\frac{n}{2}\right)$

$$7\frac{cn^3}{2} + 18n^2 \leq cn^3$$

The conjugate of $7\frac{n^3}{2}$ is added and subtracted to remove the denominator from the n^3 term

$$7cn^3 - 7\frac{cn^3}{2} + 18n^2 \leq cn^3$$

The 7 coefficient is cancelled through division

$$cn^3 - \left(\frac{cn^3}{2} - \frac{18}{7}(n^2)\right)$$

The $\frac{cn^3}{2} - \frac{18}{7}(n^2)$ term can be assumed to be positive, therefore $cn^3 - \left(\frac{cn^3}{2} - \frac{18}{7}(n^2)\right) \leq cn^3$

3.

```
public int recursive(int n) {  
    int sum=0;  
    for (int i=1; i<= n; i++){  
        sum= sum +1;  
    }  
    if (n>1) {  
        return recursive(n/2);  
    }  
    else {  
        return 1  
    }  
}
```

The initial for loop iterates from 1 to n , or n times and completes $\frac{n}{2}$ recursions; therefore

$$T(n) \approx T\left(\frac{n}{2}\right) + n.$$

Similar to the above problem, if we assume $O(n \log(n))$, we can substitute n for $n \log(n)$ for all values $< n$

$$c\left(\frac{n}{2}\right)\log\left(\frac{n}{2}\right) + n$$

Like above, expand the log and then add and subtract the conjugate

$$cn \log n - \left(\frac{cn \log n}{2} + cn - n \right)$$

Again the term being subtracted is positive and therefore

$$cn \log n - \left(\frac{cn \log n}{2} + cn - n \right) \leq cn \log n$$