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CUS 1188

Problem Set 6

$$1) T(n) = 2T\left(\frac{n}{2}\right) + 1000n$$

Every time the recurrence is called for a value  $k$ , the recurrence will be

$$T(k) = 2T\left(\frac{k}{2}\right) + 1000k$$

By plugging in the value  $n/2$  into the recurrence, the whole equation becomes

$$T(n) = 2\left[2T\left(\frac{n}{2^2}\right) + \frac{1000n}{2^1}\right] + 1000n$$

By expanding

$$2^2 T\left(\frac{n}{2^2}\right) + 2^1 \left(\frac{1000n}{2^1}\right) + 1000n$$

If the recurrence is unfolded again with  $n/2$  it becomes

$$2^2 \left[2T\left(\frac{n}{2^3}\right) + \frac{1000n}{2^2}\right] + 2^1 \left(\frac{1000n}{2^1}\right) + 1000n$$

Expanding again yields

$$2^3 T\left(\frac{n}{2^3}\right) + 2^2 \left(\frac{1000n}{2^2}\right) + 2^1 \left(\frac{1000n}{2^1}\right) + 1000n$$

The pattern becomes clear and can be written as

$$2^k T\left(\frac{n}{2^k}\right) + k(1000n)$$

To set  $T(n)$  to 1,  $n$  must be 1 or in other words,  $2^k$  must be  $n$ ;

$2^k = n$ , solving for  $n$  gives  $k$  to be  $\log_2 n$

$2^{\log_2 n}$  cancels with  $\left(\frac{n}{2^{\log_2 n}}\right)$  to just 1

$$1 + \log_2 n(1000n)$$

$$T(n) \in O(n \log n)$$

$$2) T(n) = 7T\left(\frac{n}{2}\right) + 18n^2$$

$$T(k) = 7T\left(\frac{k}{2}\right) + 18k^2$$

$$7\left[7T\left(\frac{n}{2^2}\right) + \frac{18n^2}{2^1}\right] + 18n^2$$

Expanding gives

$$7^2 T\left(\frac{n}{2^2}\right) + \frac{7^1}{2^1} (18n^2) + 18n^2$$

By unfolding again we get

$$7^2 \left[7T\left(\frac{n}{2^3}\right) + \frac{18n^2}{2^2}\right] + \frac{7^1}{2^1} (18n^2) + 18n^2$$

Expanding again gives

$$7^3 T\left(\frac{n}{2^3}\right) + \frac{7^2}{2^2} (18n^2) + \frac{7^1}{2^1} (18n^2) + 18n^2$$

The pattern can be expressed like so

$$7^k T\left(\frac{n}{2^k}\right) + \sum_{i=0}^{k-1} (18n^2) \left(\frac{7}{2}\right)^i$$

The summation can be written according to this formula

$$\sum_{k=0}^{n-1} ar^k = a \left( \frac{1-r^n}{1-r} \right)$$

as

$$(18n^2) \left( \frac{1 - \left(\frac{7}{2}\right)^k}{1 - \left(\frac{7}{2}\right)} \right)$$

Similar to earlier,  $2^k$  must be n to cancel  $\frac{n}{2^k}$

$2^k = n$ , solving for n gives k to be  $\log_2 n$

$$7^{\log_2 n} + (18n^2) \left( \frac{1 - \left(\frac{7}{2}\right)^{\log_2 n}}{1 - \left(\frac{7}{2}\right)} \right)$$

Because of the log rule  $a^{\log b} = b^{\log a}$

$$n^{\log 7} + (18n^2)^{\left(\frac{1-n^{\log(\frac{7}{2})}}{1-(\frac{7}{2})}\right)}$$

$\log 7$  is roughly 3, so  $\log(\frac{7}{2})$  is roughly 2; substituting these and expanding roughly equals

$$n^3 + \left(\frac{18n^2 - 18n^4}{1-(\frac{7}{2})}\right) \in O(n^4)$$

3)

```
public int recursive(int n) {  
    int sum=0;  
    for (int i=1; i<= n; i++){  
        sum= sum +1;  
    }  
    if (n>1) {  
        return recursive(n/2);  
    }  
    else {  
        return 1  
    }  
}
```

The recursive function is only called once per recurrence as well as n iterations therefore

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$

Unfolding the recurrence gives

$$T(n) = 2T\left(\frac{n}{2^2}\right) + 2^1 \frac{n}{2^1} + n$$

Unfolding again gives

$$2T\left(\frac{n}{2^3}\right) + 2^2 \frac{n}{2^2} + 2^1 \frac{n}{2^1} + n$$

The pattern is simple

$$2T\left(\frac{n}{2^k}\right) + kn$$

Substituting k for logn gives

$$n \log n \in O(n \log n)$$

