

Minimizing Schedule Changes

Brandon Tang, Isabelle Telford, Grace Ekstrom, Brendan Willis

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Introduction

Our project had 2 parts:

1. Build a course schedule for the UMass Amherst Mathematics and Statistics Department that maximizes the faculty members' course preference and time availability
2. Take changes that need to be made to the schedule, and minimize how many other changes are made as a result

Goals

- ▶ We want to find what course schedule maximizes faculty preferences of course and time, while leaving only as many lower-level and large lecture courses as necessary unassigned to an instructor
- ▶ From the course schedule we want to minimize changes made to the original schedule, while satisfying issues coming up after the first schedule has been made
- ▶ Examples of this are a faculty member no longer being able to teach at the given time or the given course, or going suddenly on leave

Schedule Builder

The first portion of our project was constructing a linear program that creates a schedule maximizing the instructor preferences.

We will use this schedule as the starting point for changes that need to be made later.

Basic Information:

- ▶ 70 instructors (18 VAPs, 14 lecturers, 38 tenure-track professors)
- ▶ 133 total sections of 61 total courses to be taught
- ▶ 14 time slots (MWF 8-8:50 considered 1 time slot)

Schedule Builder

Decisions

We want to decide who is teaching what class and when, and which constraints to fully satisfy

Variables

- Assignment Variables: these are all binary variables that equal 0 if person i isn't teaching class j (at section level, not course) at time h , and 1 if they are.

x_{ijh} = VAP $i \in \{1, \dots, n_1\}$ teaching class $j \in \{1, \dots, m\}$ at time $h \in \{1, \dots, l\}$

y_{ijh} = lecturer $i \in \{1, \dots, n_2\}$ teaching class j at time h

z_{ijh} = tenure-track professor $i \in \{1, \dots, n_3\}$ teaching class j at time h

- Softness Variables: these are binary and non-negative integer variables added in to allow some constraints to not be fully satisfied

Schedule Builder

Goal

We want to find what schedule maximizes instructor time and course preferences

Objective Function

Let a_{ij} = preference weight for person i teaching course j ,
and b_{ih} = preference weight for person i teaching at time h .
Then the objective is to:

$$\begin{aligned} \max \quad & \sum_{i=1}^{n_1} \sum_{j=1}^m \sum_{h=1}^l f(a_{ij}, b_{ih}) x_{ijh} + \sum_{i=1}^{n_2} \sum_{j=1}^m \sum_{h=1}^l f(a_{ij}, b_{ih}) y_{ijh} \\ & + \sum_{i=1}^{n_3} \sum_{j=1}^m \sum_{h=1}^l f(a_{ij}, b_{ih}) z_{ijh} - \text{softness variables} \end{aligned}$$

Where $f(a_{ij}, b_{ih})$ is some function of the preferences. We tested
 $f(a, b) = ab$ and $f(a, b) = a + b$

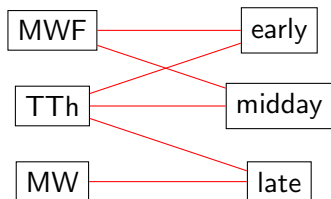
Schedule Builder

Determining a

- ▶ A weight of 0 is given to any course an instructor explicitly said they did not want to teach
- ▶ A weight of 0.25 is given to any course an instructor did not mention in the spreadsheet
- ▶ In general, a weight of 1 is given to a person and course if it was their top choice when ranking upper/lower/lecture courses.
- ▶ Similarly, a weight of 0.75 is given to second choices, and 0.5 to third choices
- ▶ However, if a person did not volunteer to teach large lectures, their preferences for those were reduced to 0.3, 0.2, and 0.1 respectively.

Schedule Builder

Determining b



1. Rank MWF/TTh/MW 1-3 (1 best)
2. Rank early/midday/late 1-3 (1 best)
3. Multiply the values as shown in the image
4. Assign the following preferences by the product:
 - ▶ $1 = 1$
 - ▶ $2 = 0.8$
 - ▶ $3 = 0.6$
 - ▶ $4 = 0.4$
 - ▶ $6 = 0.2$
 - ▶ $9 = 0$

Schedule Builder

Constraints

1. Each course should be assigned exactly one time slot and one instructor, with the exception of some lower-level courses or large lectures which may be filled in by new instructors later.

$$\sum_{i=1}^{n_1} \sum_{h=1}^l x_{ijh} + \sum_{i=1}^{n_2} \sum_{h=1}^l y_{ijh} + \sum_{i=1}^{n_3} \sum_{h=1}^l z_{ijh} = 1 - p_j \quad \forall j \quad (1)$$

$$p_j = 0 \quad \forall j \in \{\text{upper-level or grad class}\} \quad (2)$$

Schedule Builder Constraints

2. Each person teaches at most one class during any time-slot.

$$\sum_{j=1}^m x_{ijh} \leq 1 \quad \forall i \in \{1, \dots, n_1\}, h \quad (3)$$

$$\sum_{j=1}^m y_{ijh} \leq 1 \quad \forall i \in \{1, \dots, n_2\}, h \quad (4)$$

$$\sum_{j=1}^m z_{ijh} \leq 1 \quad \forall i \in \{1, \dots, n_3\}, h \quad (5)$$

Schedule Builder Constraints

3. Pre-assigned course-instructor c, i pairings are upheld.

Let $C = \{j : j \text{ is a section of course } c\}$

$$\sum_{j \in C} \sum_{h=1}^I x_{ijh} = r_{ic} \quad (6)$$

if the instructor is a VAP,

$$\sum_{j \in C} \sum_{h=1}^I y_{ijh} = r_{ic} \quad (7)$$

if the instructor is a lecturer, and

$$\sum_{j \in C} \sum_{h=1}^I z_{ijh} = r_{ic} \quad (8)$$

if the instructor is a tenure-track professor, where r_{ic} = the number of sections of course c instructor i is teaching.

Schedule Builder Constraints

4. Instructors that cannot teach certain classes should not be assigned to teach those classes.

For fixed i^* and j^* being the instructor-course pair that cannot be assigned:

$$\sum_{h=1}^I x_{i^*j^*h} = 0 \quad (9)$$

if the instructor is a VAP,

$$\sum_{h=1}^I y_{i^*j^*h} = 0 \quad (10)$$

if the instructor is a lecturer, and

$$\sum_{h=1}^I z_{i^*j^*h} = 0 \quad (11)$$

if the instructor is a tenure-track professor.

Schedule Builder Constraints

5. Instructors that cannot teach at certain times should not be assigned to teach during those times.

For fixed i^* and h^* being the instructor-time pair that cannot be assigned:

$$\sum_{j=1}^m x_{i^*jh^*} = 0 \quad (12)$$

if the instructor is a VAP,

$$\sum_{j=1}^m y_{i^*jh^*} = 0 \quad (13)$$

if the instructor is a lecturer, and

$$\sum_{j=1}^m z_{i^*jh^*} = 0 \quad (14)$$

if the instructor is a tenure-track professor.

Schedule Builder Constraints

6. Lecturers cannot teach upper-level classes.

$$\sum_{h=1}^I x_{ijh} = 0 \quad \forall i \in \{1, \dots, n_1\}, j \in \{\text{upper-level}\} \quad (15)$$

7. Each instructor teaches the correct number of courses. As this number cannot be generalized for everyone, we will assign each person a parameter L_{xi} , L_{yi} , or L_{zi} accordingly that represents the load or number of classes that person is able to teach. (SOFT, may teach 1 less course)

$$\sum_{j=1}^m \sum_{h=1}^I x_{ijh} = L_{xi} \quad \forall i \in \{0, \dots, n_1\} \quad (16)$$

$$\sum_{j=1}^m \sum_{h=1}^I y_{ijh} = L_{yi} \quad \forall i \in \{0, \dots, n_2\} \quad (17)$$

$$\sum_{j=1}^m \sum_{h=1}^I z_{ijh} = L_{zi} \quad \forall i \in \{0, \dots, n_3\} \quad (18)$$

Schedule Builder Constraints

8. Instructors should not be assigned to teach both MWF/MW and TTh.

$$\sum_{j=1}^m x_{ijh'} + \sum_{j=1}^m x_{ijh''} \leq 1 + q_{xi} \quad \forall i, h', h'' \quad (19)$$

$$\sum_{j=1}^m y_{ijh'} + \sum_{j=1}^m y_{ijh''} \leq 1 + q_{yi} \quad \forall i, h', h'' \quad (20)$$

$$\sum_{j=1}^m z_{ijh'} + \sum_{j=1}^m z_{ijh''} \leq 1 \quad \forall i, h', h'' \quad (21)$$

Where $h' = \{1, \dots, 8\}$, $h'' = \{9, \dots, 14\}$, and $q_{xi}, q_{yi} = 1$ if a VAP or lecturer respectively would like to teach all 5 days, 0 otherwise.

Schedule Builder Constraints

9. Instructors should not be assigned to teach both early MWF and late MW. (SOFT for lecturers)

$$\sum_{j=1}^m x_{ij1} + \sum_{j=1}^m x_{ijw} \leq 1 \quad \forall i \quad (22)$$

$$\sum_{j=1}^m y_{ij1} + \sum_{j=1}^m y_{ijw} \leq 1 \quad \forall i \quad (23)$$

$$\sum_{j=1}^m z_{ij1} + \sum_{j=1}^m z_{ijw} \leq 1 \quad \forall i \quad (24)$$

Where w = the last MW time slot.

Schedule Builder Constraints

10. Instructors should not be assigned to teach both early and late TTh. (SOFT for lecturers)

$$\sum_{j=1}^m x_{ij(w+1)} + \sum_{j=1}^m x_{ijl} \leq 1 \quad \forall i \quad (25)$$

$$\sum_{j=1}^m y_{ij(w+1)} + \sum_{j=1}^m y_{ijl} \leq 1 \quad \forall i \quad (26)$$

$$\sum_{j=1}^m z_{ij(w+1)} + \sum_{j=1}^m z_{ijl} \leq 1 \quad \forall i \quad (27)$$

Where w = the last MW time slot, l = the last TTh time slot.

Schedule Builder Constraints

11. For each time slot, there should be an appropriate number of lower- and upper-level courses being taught so we can assign classes to rooms. (SOFT with large penalty)

$$\left\lfloor \frac{\alpha}{\beta} \right\rfloor \leq \sum_{i=1}^{n_1} \sum_{j=1}^{\alpha} x_{ijh} + \sum_{i=1}^{n_2} \sum_{j=1}^{\alpha} y_{ijh} + \sum_{i=1}^{n_3} \sum_{j=1}^{\alpha} z_{ijh} \leq \left\lceil \frac{\alpha}{\beta} \right\rceil \quad (28)$$
$$\forall h \in \{\text{not first or last of the day}\}$$

$$\left\lfloor \frac{m - \alpha}{\beta} \right\rfloor \leq \sum_{i=1}^{n_1} \sum_{j=\alpha+1}^m x_{ijh} + \sum_{i=1}^{n_2} \sum_{j=\alpha+1}^m y_{ijh} + \sum_{i=1}^{n_3} \sum_{j=\alpha+1}^m z_{ijh} \leq \left\lceil \frac{m - \alpha}{\beta} \right\rceil \quad \forall h \in \{\text{not first or last of the day}\} \quad (29)$$

Where α = the number of lower-level and large lecture sections being taught, and β = the number of time slots that are not the first or last of the day.

Schedule Builder Constraints

11. For less-popular time slots, we don't need to lower-bound the number of classes.

$$\sum_{i=1}^{n_1} \sum_{j=1}^{\alpha} x_{ijh} + \sum_{i=1}^{n_2} \sum_{j=1}^{\alpha} y_{ijh} + \sum_{i=1}^{n_3} \sum_{j=1}^{\alpha} z_{ijh} \leq \left\lceil \frac{\alpha}{\beta} \right\rceil$$

$\forall h \in \{\text{first or last of the day}\}$
(30)

$$\sum_{i=1}^{n_1} \sum_{j=\alpha+1}^m x_{ijh} + \sum_{i=1}^{n_2} \sum_{j=\alpha+1}^m y_{ijh} + \sum_{i=1}^{n_3} \sum_{j=\alpha+1}^m z_{ijh} \leq \left\lceil \frac{m - \alpha}{\beta} \right\rceil$$

$\forall h \in \{\text{first or last of the day}\}$
(31)

Schedule Builder Constraints

12. We do not want sections of core courses (lower-level, large lecture) to overlap too much. The number of sections of a given course that can overlap will depend on how many sections of that class there are, so we assign a parameter: s_c = max number of sections of core course c that can be taught during a time-slot.

Let $C = \{j : j \text{ is a section of course } c\}$

$$\sum_{j \in C} \sum_{i=1}^{n_1} x_{ijh} + \sum_{j \in C} \sum_{i=1}^{n_2} y_{ijh} + \sum_{j \in C} \sum_{i=1}^{n_3} z_{ijh} \leq s_j \quad \forall h, c \in \{\text{core courses}\}$$

(32)

Schedule Builder Constraints

13. We want courses that are commonly taken together to have little overlap so that students are able to fit all of those courses in their schedule. For some courses that have many sections, there can be some overlap between that class and another, but for upper-level courses that have only one or two sections, we want these sections to be at different times. Take sets D_1, D_2, \dots to be sets of j that belong to lower-level courses often taken together.

$$\sum_{j \in D_k} \left(\sum_{i=1}^{n_1} x_{ijh} + \sum_{i=1}^{n_2} y_{ijh} + \sum_{i=1}^{n_3} z_{ijh} \right) \leq 2 \quad \forall h, k \quad (33)$$

Take sets F_1, F_2, \dots to be sets of j that belong to upper level courses often taken together.

$$\sum_{j \in F_k} \left(\sum_{i=1}^{n_1} x_{ijh} + \sum_{i=1}^{n_2} y_{ijh} + \sum_{i=1}^{n_3} z_{ijh} \right) \leq 1 \quad \forall h, k \quad (34)$$

Schedule Builder Constraints

14. If someone is up for promotion but has not taught a variety of courses yet, we want them to teach at least one new course during this semester.

Let

$\alpha_i = \{c : \text{person } i \text{ has taught course } c \text{ before}\}$

$C_{\alpha_i} = \{j : j \text{ is a section of a course in } \alpha_i\}$

$$\sum_{j \in C_{\alpha_i}} \sum_{h=1}^l x_{ijh} + 1 \leq \sum_{j=1}^m \sum_{h=1}^l x_{ijh} \quad (35)$$

$\forall i \in \{\text{VAPs to be promoted and needs new courses}\}$

$$\sum_{j \in C_{\alpha_i}} \sum_{h=1}^l y_{ijh} + 1 \leq \sum_{j=1}^m \sum_{h=1}^l y_{ijh} \quad (36)$$

$\forall i \in \{\text{lecturers to be promoted and needs new courses}\}$

Schedule Builder Constraints

15. To ensure that courses for which some sections are left open still have a qualified chair, we want at least one of the sections to be taught by a tenure-track professor or permanent lecturer. (SOFT, and forced to not need to hold if a VAP is pre-assigned to a course)

$$\sum_{i=1}^{n_3} \sum_{j \in C} \sum_{h=1}^l z_{ijh} + \sum_{i=1}^{n_2} \sum_{j \in C} \sum_{h=1}^l y_{ijh} \geq 1 \quad (37)$$

$$\forall c \in \{\text{lower-level or big lecture}\}$$

Schedule Builder Results

[illegible]

Figure 1: Optimal Schedule 1 ($a * b$)

Schedule Builder Results

[illegible]

Figure 2: Optimal Schedule 2 ($a + b$)

Schedule Builder Results

A Closer Look at Some Instructors

Professor 8

- ▶ 2 classes
- ▶ Upper-level preferences: 37, (41,34), (28,39)
- ▶ Does not volunteer for large lectures
- ▶ Lower-level preferences: 18, 13
- ▶ No classes they prefer not to teach
- ▶ Time preferences: Early (not 8am MWF), midday, late
- ▶ Day preferences: TTh, MW, MWF
- ▶ Has reasons to only teach 2 days a week (didn't implement)
- ▶ No specific time they can't teach
- ▶ **Schedule 1:** teaching 37 (TTh 8:30-9:45) and 18 (TTh 10-11:15)
- ▶ **Schedule 2:** teaching 37 (TTh 8:30-9:45) and 38 (TTh 10-11:15)

Schedule Builder Results

A Closer Look at Some Instructors

Lecturer 11

- ▶ 3 classes
- ▶ Does not want to teach upper-level courses
- ▶ Large lecture preferences: (4,8), (9,1,2), (11,6,7) (volunteers)
- ▶ Lower-level preferences: 10, 40 (not lower-level), 14
- ▶ Time preferences: Early, midday, late
- ▶ Day preferences: TTh, MWF, MW
- ▶ No specific time they can't teach
- ▶ **Schedule 1:** teaching 14 (MWF 9:05-9:55), 14 (MWF 10:10-11), 16 (MWF 1:25-2:15)
- ▶ **Schedule 2:** teaching 16 (MWF 10:10-11), 14 (MWF 12:20-1:10), 14 (MWF 1:25-2:15)

Schedule Builder Results

A Closer Look at Some Instructors

VAP 2

- ▶ 2 classes
- ▶ Upper-level preferences: (28,39,41,36,23,37), (30,27), (34,26)
- ▶ Does not volunteer for large lectures
- ▶ Lower-level preferences: (18,13), (17,16), 14
- ▶ No classes they prefer not to teach
- ▶ Time preferences: Early, midday, late
- ▶ Day preferences: TTh, MW, MWF
- ▶ Can't teach Tuesdays 4-5pm
- ▶ **Schedule 1:** teaching 27 (MW 4-5:15), 27 (TTh 10-11:25)
- ▶ **Schedule 2:** teaching 27 (TTh 8:30-9:45), 27 (TTh 10-11:25)

Change Minimizer

Decisions

We want to decide who is teaching what class and when in the new schedule, and what has changes

Variables

- ▶ Assignment Variables: We will have the same x_{ijh} , y_{ijh} , z_{ijh} , and p_j variables in this program as in the schedule builder.
- ▶ Change Variables: We have binary variables (1 = yes, changed, 0 = no, stayed the same):
 - \hat{x}_{ijh} = whether the status of VAP i teaching class j at time h has changed
 - \hat{y}_{ijh} = whether the status of lecturer i teaching class j at time h has changed
 - \hat{z}_{ijh} = whether the status of tenure-track professor i teaching class j at time h has changed
- ▶ Softness Variables: more variables to allow some constraints to not be fully satisfied if necessary

Change Minimizer

Goal

We want to find the schedule the minimizes changes made to an original schedule.

Objective Function

$$\begin{aligned} \min \quad & \sum_{i=1}^n \sum_{j=1}^m \sum_{h=1}^l \hat{x}_{ijh} + \sum_{i=1}^n \sum_{j=1}^m \sum_{h=1}^l \hat{y}_{ijh} \\ & + \sum_{i=1}^n \sum_{j=1}^m \sum_{h=1}^l \hat{z}_{ijh} + \text{softness variables} \end{aligned} \tag{38}$$

Change Minimizer

Constraints:

1. We will take an already-created schedule, and assign the following binary (1 = yes, 0 = no) parameters from that:
 \bar{x}_{ijh} = lecturer i teaches class j at time h in original schedule
 \bar{y}_{ijh} = VAP i teaches class j at time h in original schedule
 \bar{z}_{ijh} = tenure-track professor i teaches class j at time h in original schedule
Then we want the $\hat{x}, \hat{y}, \hat{z}$ variables to take on the right values, so we have:

$$\hat{x}_{ijh} = x_{ijh} + \bar{x}_{ijh} - 2\bar{x}_{ijh}x_{ijh} \quad \forall i, j, h \quad (39)$$

$$\hat{y}_{ijh} = y_{ijh} + \bar{y}_{ijh} - 2\bar{y}_{ijh}y_{ijh} \quad \forall i, j, h \quad (40)$$

$$\hat{z}_{ijh} = z_{ijh} + \bar{z}_{ijh} - 2\bar{z}_{ijh}z_{ijh} \quad \forall i, j, h \quad (41)$$

Change Minimizer Constraints

2. All constraints from the schedule builder will apply for this program as well.
3. For instructor i^* who has an issue with teaching class j^* , we add a constraint as in (4) of the schedule builder.
4. For instructor i^* who has an issue with teaching at time h^* , we add a constraint as in (5) of the schedule builder.

Change Minimizer Constraints

5. For an instructor i^* who is suddenly on leave, we add a constraint

$$\sum_{j=1}^m \sum_{h=1}^l x_{i^*jh} = 0, \quad (42)$$

$$\sum_{j=1}^m \sum_{h=1}^l y_{i^*jh} = 0, \text{ or} \quad (43)$$

$$\sum_{j=1}^m \sum_{h=1}^l z_{i^*jh} = 0 \quad (44)$$

to match the level of instructor they are.

Change Minimizer Constraints

6. If pre-registration, all courses c should keep its original instructor i^* and have the time changed. (SOFT) Where $C = \{j : j \text{ is a section of course type } c\}$,

$$\sum_{j \in C} \sum_{h=1}^I x_{i^*jh} = 1 \quad \forall c \quad (45)$$

$$\sum_{j \in C} \sum_{h=1}^I y_{i^*jh} = 1 \quad \forall c \quad (46)$$

$$\sum_{j \in C} \sum_{h=1}^I z_{i^*jh} = 1 \quad \forall c \quad (47)$$

7. If post-registration, a course c should keep its time-slot h^* and have the instructor changed. (SOFT)

$$\sum_{j \in C} \sum_{i=1}^{n_1} x_{ijh^*} + \sum_{j \in C} \sum_{i=1}^{n_2} y_{ijh^*} + \sum_{j \in C} \sum_{i=1}^{n_3} z_{ijh^*} = 1 \quad \forall c \quad (48)$$

Change Minimizer Results

Changes Made to Schedule 1 ($a * b$) :

1. VAP 5 teaching course 26
2. VAP 14 teaching at time 7
3. Lecturer 11 teaching at time 3
4. Professor 14 teaching at time 11
5. Professor 2 on leave

Change Minimizer Results

[illegible]

Figure 3: New Schedule 1 - Pre Registration

Change Minimizer Results

[illegible]

Figure 4: New Schedule 1 - Post Registration

Change Minimizer Results

Changes Made to Schedule 2 ($a + b$) :

1. VAP 5 teaching course 24
2. VAP 14 teaching at time 7
3. Lecturer 11 teaching at time 3
4. Professor 13 teaching at time 10
5. Professor 2 on leave

Change Minimizer Results

[illegible]

Figure 5: New Schedule 2 - Pre Registration

Change Minimizer Results

[illegible]

Figure 6: New Schedule 2 - Post Registration

Comparing Change Minimizers

	$a * b$	$a + b$
Pre-Reg.	5 exact 16 same unassigned 49 same instructor	5 exact 17 same unassigned 46 same instructor
Post-Reg.	33 exact 15 same unassigned 108 at same time	39 exact 16 same unassigned 108 at same time

- ▶ Much better post-registration
- ▶ 133 sections total, so over 90% of assigned classes stay at same time post-registration
- ▶ $a + b$ leads to slightly better results, but more analysis needed to see if it is significant

Change Minimizer Results

A Closer Look at Some Instructors

Lecturer 11

- ▶ **Forced Change:** No teaching MWF 10:10-11
- ▶ **Schedule 1:** teaching 14 (MWF 9:05-9:55), 14 (MWF 10:10-11), 16 (MWF 1:25-2:15)
- ▶ **Schedule 1 PRE:** teaching 14 (MWF 8-8:50), 14 (12:20-1:10)
- ▶ **Schedule 1 POST:** teaching 14 (MWF 9:05-9:55), 3 (MWF 11:15-12:05), 14 (MWF 1:25-2:15)
- ▶ **Schedule 2:** teaching 16 (MWF 10:10-11), 14 (MWF 12:20-1:10), 14 (MWF 1:25-2:15)
- ▶ **Schedule 2 PRE:** teaching 14 (TTh 8:30-9:45), 14 (TTh 2:30-3:45)
- ▶ **Schedule 2 POST:** teaching 17 (MWF 8-8:50), 14 (MWF 1:25-2:15), 14 (TTh 8:30-9:45)

Change Minimizer Results

A Closer Look at Some Instructors

Professor 8

- ▶ No forced change, happy with original schedule
- ▶ **Schedule 1:** teaching 37 (TTh 8:30-9:45) and 18 (TTh 10-11:15)
- ▶ **Schedule 1 PRE:** teaching 27 (MWF 8-8:50)
- ▶ **Schedule 1 POST:** teaching 14 (MWF 10:10-11), 27 (TTh 8:30-9:45)
- ▶ **Schedule 2:** teaching 37 (TTh 8:30-9:45) and 38 (TTh 10-11:15)
- ▶ **Schedule 2 PRE:** teaching 44 (MWF 8-8:50), 31 (MWF 10:10-11)
- ▶ **Schedule 2 POST:** teaching 27 (TTh 8:30-9:45), 26 (TTh 1:00-2:15)
- ▶ None of new classes were highly ranked

Conclusions

- ▶ Schedule builder seems to create a reasonable schedule
- ▶ Even "minimizing" changes still creates a lot of changes
- ▶ Easier to keep courses at the same time than keep professors teaching the same courses
- ▶ Might want to make some schedule builder constraints soft(er) in change minimizer if the number of changes is more important
- ▶ Generating a schedule can be automated in a way that seems to be effective, but hard to know if the lack of "human touch" makes a difference without seeing the hand-made schedule
- ▶ At this stage, the change minimizer might be better used as a tool to see a possible way of accommodating changes
 - ▶ Our code might create a domino effect of changes to satisfy constraints when you could just switch two people and technically break a constraint, but save a lot of changes from being made

Future Ideas

- ▶ Implement constraints for same vs. different classes, back to back classes, forcing 2 days/week if needed
- ▶ Compare our output with actual schedule for next semester
- ▶ Make preference data less arbitrary by changing how the department gathers preferences: ask them to explicitly weight time slots and courses 0-1
- ▶ Experiment with other ways of formulating change minimizer objective function to get better results
 - ▶ More softness of schedule builder constraints
 - ▶ Potentially include preference data
- ▶ Look at individual-level results of more instructors to identify any potential biases or issues

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