

1. A) $f(n) = \Omega(g(n))$ because as n approaches infinity, $f(n)$ is growing faster than $g(n)$.
B) $f(n) = O(g(n))$ because as n approaches infinity, $g(n)$ is growing faster than $f(n)$.
C) $f(n) = \Omega(g(n))$ because as n approaches infinity, $f(n)$ is growing faster than $g(n)$.
D) $f(n) = O(g(n))$ because as n approaches infinity, $g(n)$ is growing faster than $f(n)$.
E) $f(n) = \Theta(g(n))$ because as n approaches infinity, $g(n)$ is growing equally.
F) $f(n) = \Omega(g(n))$ because as n approaches infinity, $f(n)$ is growing faster than $g(n)$.

2. A) If $f_1(n) = \Theta(g(n))$ and $f_2(n) = \Theta(g(n))$ then $f_1(n) + f_2(n) = \Theta(g(n))$

$f(n) = \Theta(g(n))$ if there exist positive constants c_1, c_2 and n_0 , such that $0 \leq c_1g(n) \leq f(n) \leq c_2g(n)$ for all $n \geq n_0$.

$f_1(n) = \Theta(g(n))$ so there must exist c_1 and c_2 such that $c_1g(n) \leq f_1(n) \leq c_2g(n)$

$f_2(n) = \Theta(g(n))$ so there must exist c_3 and c_4 such that $c_3g(n) \leq f_2(n) \leq c_4g(n)$

$f_1(n) + f_2(n) = c_1g(n) + c_3g(n) \leq f_1(n) + f_2(n) \leq c_2g(n) + c_4g(n)$

We can simplify the equation above into the following:

$(c_1 + c_3)g(n) \leq f_1(n) + f_2(n) \leq (c_2 + c_4)g(n)$

$c_5g(n) \leq f_1(n) + f_2(n) \leq c_6g(n)$ which satisfies the definition of theta.

Therefore, $f_1(n) + f_2(n) = \Theta(g(n))$ hence TRUE.

- B) If $f_1(n) = O(g(n))$ and $f_2(n) = O(g(n))$ then $f_1(n) = O(f_2(n))$

This statement is clearly false.

A counter example would be:

Suppose $f_1(n) = n$ and $f_2(n) = \log n$ and $g(n) = n^2$.

$f_1(n) \leq g(n)$ and $f_2(n) \leq g(n)$ so $f_1(n) = O(g(n))$ and $f_2(n) = O(g(n))$ are true.

If $f_1(n)$ is equal to $O(f_2(n))$ then $f_1(n) \leq cf_2(n)$ must be true.

However, $n \leq c \log n$ will never be true for any value of $c > 0$.

Therefore, $f_1(n) = O(f_2(n))$ is FALSE.

3. mergeSort.cpp and insertSort.cpp

4. Modified mergeSort2.cpp and insertSort2.cpp

A) Modified insertSort2 source code and mergeSort2 source code.

```
#include <iostream>
#include <string>
#include <ctime>
#include <cstdlib>

using namespace std;

// modified insertion sort method
void insertionSort(int list[], int size) {
    for (int x = 1; x < size; x++) {
        int current = x;
        int index = x - 1;

        // compare current value with our analyzed section
        while (index >= 0 && list[current] < list[index]) {
            int temp = list[index];
            list[index] = list[current];
            list[current] = temp;

            // update index
            current = index;
            index--;
        }
    }
}

int main() {
    for (int i = 0; i < 10; i++) {
        srand(time(0));
        int size;
        int* list;
        cout << "Enter a number between 10000 and 1000000: " << endl;
        cin >> size;

        // allocate memory for our dynamic array
        list = new int[size];

        // get elements for dynamic array
        for (int x = 0; x < size; x++) {
            int temp = rand() % 10001;
            list[x] = temp;
        }

        // Sort values and output the time
        clock_t start;
        start = clock();
```

```

        insertionSort(list, size);
        start = clock() - start;
        cout << endl << "It took " << double(start) / CLOCKS_PER_SEC << " seconds to sort" <<
endl << endl;

        // deallocate
        delete[] list;
    }
    return 0;
}

```

```

#include <iostream>
#include <string>
#include <ctime>

```

```

using namespace std;

```

```

// helper function

```

```

void merge(int list[], int start, int mid, int end) {
    // temp list to hold values in sorted order
    int * tempList = new int[end];
    int index = 0;
    int x = start;
    int y = mid;

    // iterate through sub arrays
    // will compare the current index and pass the appropriate value to the temp list
    while (x < mid && y < end) {
        if (list[x] <= list[y]) {
            tempList[index] = list[x];

            // update index
            index++;
            x++;
        }
        else if (list[y] <= list[x]) {
            tempList[index] = list[y];
            index++;
            y++;
        }
    }
}

```

```

// loops to make sure the subarrays were fully traversed
// pass the remaining values
while (x < mid) {
    tempList[index] = list[x];
    index++;
    x++;
}

```

```

    }

    while (y < end) {
        tempList[index] = list[y];
        index++;
        y++;
    }

    // copy values from temp dynamic array back to input lis
    for (x = start, y = 0; y < index; x++, y++) {
        list[x] = tempList[y];
    }

    // deallocate
    delete[] tempList;
}

// merge sort function
void mergeSort(int list[], int start, int end) {
    if (start + 1 == end) {
        return;
    }

    int mid = (start + end) / 2;
    // recursive call for the first subarray
    mergeSort(list, start, mid);
    // recursive call for the second subarray
    mergeSort(list, mid, end);
    // call helper function to merge in proper order.
    merge(list, start, mid, end);
}

int main() {
    for (int i = 0; i < 10; i++) {
        srand(time(0));
        int size;
        int* list;
        cout << "Enter a number between 10000 and 1000000: " << endl;
        cin >> size;
        // allocate memory for dynamic array
        list = new int[size];
        // get elements for dynamic array
        for (int x = 0; x < size; x++) {
            int temp = rand() % 10001;
            list[x] = temp;
        }
        // sort values and output time cost
        clock_t start;
        start = clock();
    }
}

```

```

mergeSort(list, 0, size);
start = clock() - start;
cout << endl << "It took " << double(start) / CLOCKS_PER_SEC << " seconds to sort" <<
endl << endl;

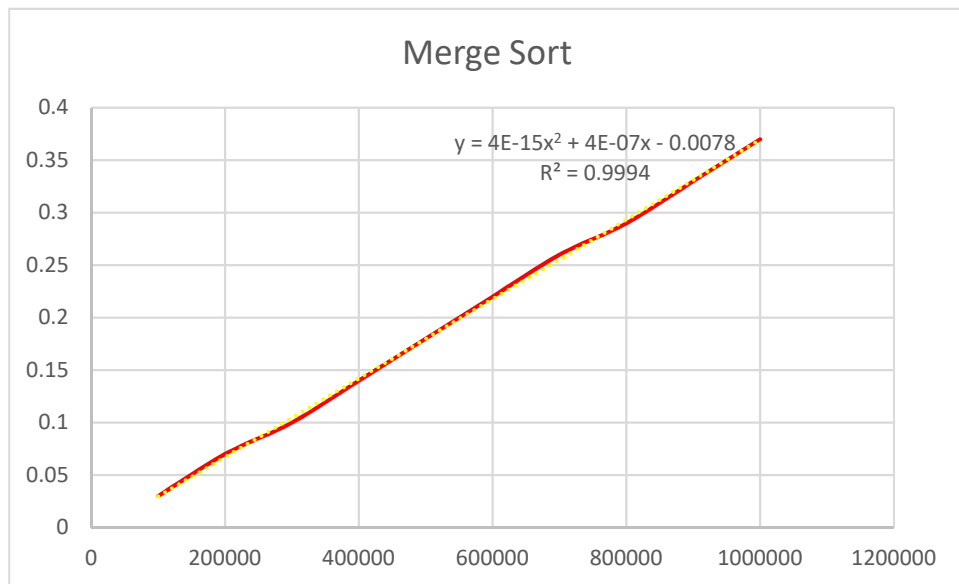
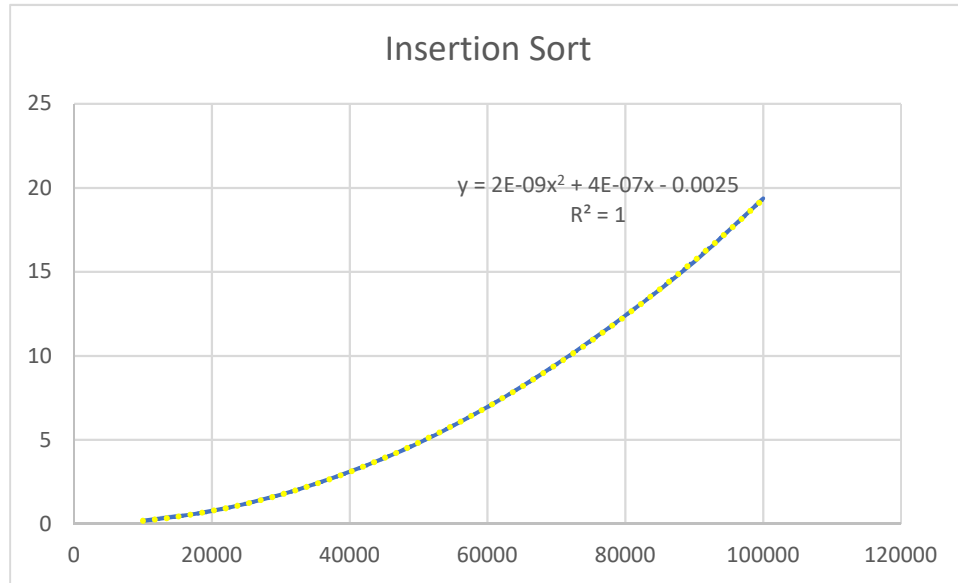
// deallocate
delete[] list;
}
return 0;
}

```

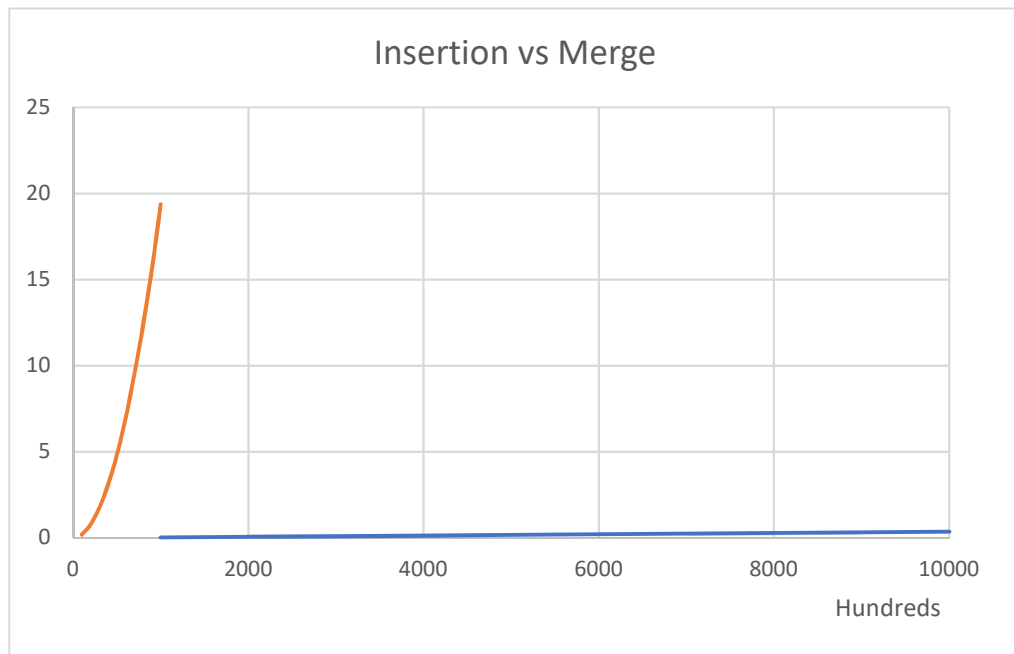
B)

insert sort			merge sort	
n	time		n	time
10000	0.19		100000	0.03
20000	0.78		200000	0.07
30000	1.75		300000	0.1
40000	3.11		400000	0.14
50000	4.82		500000	0.18
60000	6.97		600000	0.22
70000	9.51		700000	0.26
80000	12.4		800000	0.29
90000	15.61		900000	0.33
100000	19.37		1000000	0.37

c)



D)



E) Theoretical runtime of insertion sort is $O(n^2)$ and merge sort is $O(n \log n)$. The experimental running times are on par with the theoretical run times. Merge sort runs faster and takes less time than the insertion sort.